

CRYPTOGRAPHY #03.3

RSA

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1. RSA – KEM

RSA Key Encapsulation Mechanism assumes usage of RSA for transferring keys used in symmetric encryption. A scheme assumes that RSA keys were generated and both sides established the same hashing function H .

RSA – KEM, Sender:

- 1) Find random value $RAND$, such that: $1 < RAND < n$ (NOTE!!! Using OAEP encoding may assume additional restriction to value $RAND$).
- 2) $BRAND = \text{convertToBytes}(RAND)$
- 3) $BKEY = H(BRAND)$
- 4) $BC = \text{RSA_OAEP_ENCODING}(BRAND)$
- 5) return BC

RSA – KEM, Receiver get BC and than compute:

- 1) $KEY = \text{RSA_OAEP_DECODING}(BC)$
- 2) $BKEY = H(KEY)$

Now both sides have the same key $BKEY$ which can be used for further symmetric communication. Note, that the size of $BKEY$ is strictly related to the number of bytes returned by hashing function H . It is not a problem. The most popular symmetric ciphering scheme, AES, assumes usage of 16B, 24B or 32B key. Thus usage, for example, SHA-256 (which is returning 32B key) in the RSA-KEM scheme, and sometimes truncating the digest (when necessary) is solving this problem.

2. EMSA

Encoding Method for Signature with Appendix it is a scheme that is used to create RSA – PSS (RSA Provable Secure Signature). It is appropriately randomized to reach the highest possible security of a signature. It is also described in RFC 3448.

EMSA encoding parameters:

hLen – length (in bytes) of a hash

sLen – length of the salt. Should be set to *hLen*.

M – bytes of a message to be signed (**input parameter**)

mgf1 – mask generator function.

emLen – assumed length of a signature. Cannot exceed modulus *n*.

```
1) if emLen < hLen + sLen + 2, return error.
2) mHash = Hash(M). mHash length is equal to hLen.
3) Generate a random octet string salt of length sLen.
4) M' = salt || mHash || (0x)00 00 00 00 00 00 00 00. Length of M' is
   equal to 8 + hLen + sLen.
5) H = Hash(M'). Length of H is hLen.
6) Generate an octet string PS consisting of emLen - sLen - hLen - 2
   zero octets.
7) DB = salt || 0x01 || PS. Length of DB is equal to emLen - hLen -
   1.
8) dbMask = mgf1(H, emLen - hLen - 1).
9) maskedDB = DB \xor dbMask.
10) EM = 0xbc || H || maskedDB
11) return EM.
```

EMSA verification parameters:

hLen – length (in bytes) of a hash

sLen – length of the salt. Should be set to *hLen*.

EM – bytes of a signature (**input parameter**)

mgf1 – mask generator function.

emLen – length of *EM*

M – original bytes of message (**input parameter**)

```
1) if emLen < hLen + sLen + 2, return signature_invalid.
2) mHash = Hash(M). mHash length is equal to hLen.
3) if the most significant byte of EM is not equal to 0xbc, return
   signature_invalid.
4) Let maskedDB be the rightmost emLen - hLen - 1 octets of EM, and
   let H be the next hLen octets (SEE POINT 10 IN CODING ALGORITHM).
5) dbMask = mgf1(H, emLen - hLen - 1).
6) DB = maskedDB \xor dbMask.
7) If the emLen - hLen - sLen - 2 rightmost octets of DB are not
   zero, return signature_invalid
8) if the octet at position hLen (when indexing from 1) is not equal
   to 0x01, return signature_invalid
9) Let salt be the last sLen octets of DB.
10) M' = salt || mHash || (0x)00 00 00 00 00 00 00 00.
11) H' = Hash(M'). Length of H' is equal to hLen.
12) if H' = H, return signature_valid.
```

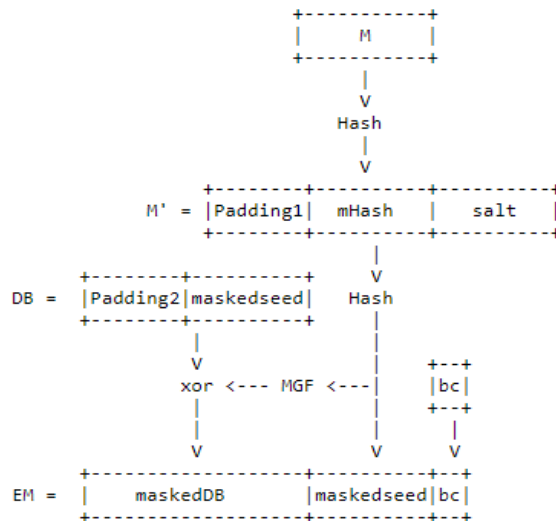


Fig. 1 EMSA coding presented in RFC-3448

3. RSA - PSS

RSA Provable Secure Signature is combining the RSA algorithm and EMSA mentioned in chapter 6. to provide a secure digital signature. The scheme assumes that RSA keys were generated and EMSA parameters established.

RSA – PSS creation for a message m :

- 1) $BEM = \text{EMSA_Encode}(m)$
- 2) $EM = \text{convertToNumber}(BEM)$
- 3) $EM = EM^d \bmod n$
- 4) $SIG = \text{convertToBytes}(EM)$
- 5) return SIG

RSA – PSS verification for a message m and signature SIG :

- 1) $EM = \text{convertToNumber}(SIG)$
- 2) $EM = EM^e \bmod n$
- 3) $BEM = \text{convertToBytes}(EM)$
- 4) return $\text{EMSA_Verify}(BEM, m)$

Exercise 1:

Write two functions that will simulate the RSA – KEM scheme accordingly to chapter 1. The first one will be generating potential 256 bits (32B) symmetric key, ciphering it with RSA – OAEP usage and return as a byte array. The second one will be deciphering, retrieving, and returning this key. Use SHA-256. Function structure:

- 1) `byte[] generateRSAKEM()`
- 2) `byte[] receiveRSAKEM(byte[] cryptogram)`

Verify that both functions work properly, and key generated in *generateRSAKEM* is the same key as in *receiveRSAKEM*.

Verification in *main* function:

```
byte[] cipher = generateRSAKEM();  
byte[] key = receiveRSAKEM (cipher);
```

Exercise 2:

Write two functions: one should generate EMSA-PSS signature and return it in byte form, the second one should verify it, and return true if signature is valid (false otherwise). Functions structure:

- 1) `byte[] createEMSAPSS(byte[] message)`
- 2) `boolean verifyEMSAPSS(byte[] EM, byte[] message)`

Parameters for *createEMSAPSS*:

- 1) *H* – is a SHA-256 hashing function. You can find it in *java.security.MessageDigest* (OR *hashlib* in *python*).
- 2) Accordingly to the previous point, *hLen* = 32.
- 3) *sLen* = 32.
- 4) *emLen* = 255.

Parameters for *verifyEMSAPSS*:

- 1) *H* – is a SHA-256 hashing function.
- 2) Accordingly to the previous point, *hLen* = 32.
- 3) *sLen* = 32.

Exercise 3:

Combine RSA and EMSA-PSS in the same way as described in chapter 3. You should write two functions:

- 1) `byte[] createRSAPSS(byte[] msg)`
- 2) `boolean verifyRSAPSS(byte[] msg, byte[] signature)`

The first one should create RSA-PSS signature and return it in byte form. The second one should verify signature and return true if it is valid (false otherwise).