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Assignment 4

Exercise 1:

1. $\Theta(n^2)$
2. $\Theta(\log(n))$
3. $\Theta(n \log(n))$
4. $\Theta(n^2)$
5. $\Theta(n^4)$
6. $\Theta(\log(n))$
7. $\Theta(n \log(n))$
8. $\Theta(n \log(n))$
9. $\Theta(n^2)$

Exercise 2:

1. -2
2. $2^3 + 5 * 15 -$

Exercise 3:

1. $T_A(n) = T(n/2) + 1$
Using Master Theorem, $n^{\log_2 1}$ vs $1 = n^0$ vs 1
 $\therefore T_A(n) = \Theta(\log n)$
2. $T_B(n) = 2T(n/2) + 1$
Using Master Theorem, $n^{\log_2 2}$ vs. $1 = n$ vs. 1

$n/1$ is a polynomial
 $\therefore T_B(n) = \Theta(n)$

3. Algorithm A is the faster of the two.

Exercise 4: The runtime of the code is $\Theta(\log^2(n))$.

Programming Task:

Description of Algorithm: Our algorithm calculates the maximum area under a histogram given an array of (x, y) coordinates. We implemented a stack to keep track of the indices of our rectangles and arrays to store the area, height, and length of each rectangle. To compute the height, we determined a common function based on the index 'i' and the index of the y-value for the ith rectangle. We did a similar approach for determining the areas and lengths of each rectangle. Lastly, we calculated the area of the top bar by using a common function to compute the area from the rightmost x-value and leftmost x-value (beginning x-value of the ith rectangle).

Code:

```
public static int getMaxArea(int hist[], int total, int n)
{
    // Stack to keep track of indices, 'total' represents
    // the number of elements, and 'n' is the number of
    // rectangles
    Stack<Integer> s = new Stack<>();

    // Arrays to store areas, heights, and lengths of each
    // rectangle
    int[] areas = new int[n];
    int[] heights = new int[n];
    int[] lengths = new int[n];

    int max_area = 0;
    int tp; // To store the index of the top of the stack
    int area_with_top; // To store the area with top bar as
    // the smallest bar

    // Calculating the area, height, and length of each
    // rectangle then storing them in the respective arrays
```

```

for (int i = 0; i < n; i++) {
    areas[i] = ((hist[4 + (4 * i)] - hist[2 + (4 * i)])
        * (hist[3 + (4 * i)]));
    heights[i] = (hist[3 + (4 * i)]);
    lengths[i] = (hist[4 + (4 * i)] - hist[2 + (4 * i)]);
}

int i = 0;
while (i < n) {
    // If the rectangle is taller or equal to the top
    // bar, then push its index onto the stack
    if (s.empty() || heights[s.peek()] <= heights[i])
        s.push(i++);
    /*
     * If the rectangle is shorter than the top bar,
     * then calculate the area of the rectangle with
     * the top as the smallest rectangle
     * 'i' is the right index for the top and the
     * element before the top is the left index
     */
    else {
        tp = s.peek();
        s.pop();

        // Calculate the area with the rectangle at
        // index 'tp' as the smallest rectangle
        // If the stack is empty, multiply height by
        // the x-value of the rightmost vertex
        // Otherwise, subtract the x-values of the
        // rightmost and leftmost vertices
        // (rightmost vertex - leftmost vertex)

        area_with_top = heights[tp]
            * (s.empty() ? hist[2 + (4 * i)] : hist
                [2 + (4 * i)] - hist[4 + (4 * s.peek
                    ())]);

        if (max_area < area_with_top)
            max_area = area_with_top;
    }
}

```

```

    }
}

// Pop the remaining rectangles from the stack
// and calculate the area with every popped rectangle
// as the smallest rectangle
while (s.empty() == false) {
    tp = s.peak();
    s.pop();

    // Same idea and calculation from previous loop
    area_with_top = heights[tp]
        * (s.empty() ? (hist[2 + (4 * i)] - hist
            [0]) : hist[2 + (4 * i)] - hist[4 + (4 *
            s.peak())]);
    if (max_area < area_with_top)
        max_area = area_with_top;
}
return max_area;
}

```

Results Table:

Data Set #	Largest Area
4.3	16
4.4	1,475,958

Table 1: Programming Task