CP468 Assignment 2

Sudoku CSP

Group 8

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1 Problem Statement

This report demonstrates the application of AC-3 algorithms to enforce arcconsistency on an arbitrary 9 x 9 Sudoku puzzle. This 9 x 9 Sudoku puzzle will be structured as a Constraint Satisfaction Problem (CSP). The AC-3 algorithm implementation keeps track of and reports the length of the queue at every step taken by the AC-3 algorithm. After the code execution, it will report whether the equivalent arc-consistent CSP is found. In the instance that the problem is solved, the solution is reported. In the event that the puzzle is not solved, an additional algorithm (backtracking) will be used to solve the puzzle entirely and the solution is reported.

2 CSP Representation:

Basic Game Knowledge: Sudoku is played by assigning numbers ranging from 1-9 on a 9 x 9 grid. This grid is divided into nine 3x3 subgrids/units. The numbers must be assigned in a way that each number only appears once in every row, column, and one time in each of the 9 subgrids.

CSP Elements:

- Variables (X): Each one of the 81 squares, denoted as an x,y pair that corresponds to the coordinate of the square on the 9 x 9 Sudoku board, the values range from x: 0-8, and y: 0-8.
- **Domains** (**D**): Each variable has a domain which represents its possible values.
- Constraints (C): The rules discussed above are represented as constraints in an array, each variable is stored by its x,y coordinate pair and each number must be unique within its row, column and sub-grid. No two cells in the same row can have the same number.

Constraints (In Depth): As discussed in class, a 9x9 Sudoku has 27 AllDiff constraints, 1 for each column, 1 for each row and 1 for each box. For Example:

Row: AllDiff(A1,A2,A3,A4,A5,A6,A7,A8,A9) Column: AllDiff(A1,B1,C1,D1,E1,F1,G1,H1,I1) Box: AllDiff(A1,A2,A3,B1,B2,B3,C1,C2,C3)

These constraints can be decomposed into individual binary constraints for each cell. These will include constraints for every other cell in the row, column and box with respect to each cell. As a general example the cell A1 would have Binary constraints:

```
A1!= A3
A1 != A4
A1!= A5
A1 != A6
A1 != A7
A1!= A8
A1!= A9
Column:
          A1 != B1
A1 != C1
A1 != D1
A1 != E1
A1 != F1
A1 != G1
A1 != H1
A1!= I1
```

Row: A1 != A2

```
Box: A1 != A2 (Repeated for example)
A1 != A3 (Repeated for example)
A1 != B1 (Repeated for example)
A1 != B2
A1 != B3
A1 != C1(Repeated for example)
A1 != C2
A1 != C3
```

This process can be repeated for every cell to create the binary constraints. As noted in the example, some of the constraints are repeated since they are already included in the previous row or column constraints.

Based on the example shown in class we can reduce the domain of cells based on already filled in/confirmed values to ensure arc consistency. Since there is only 1 solution to a Sudoku any value that is either pre filled or confirmed based on a domain of size 1 is absolute and therefore can be used to reduce the domain of related cells to ensure arc consistency. For example, cell E4:

	1	2	3	4	5	6	7	8	9
$\overline{\mathbf{A}}$	-	-	3	-	2	-	6	-	-
$\overline{\mathbf{B}}$	9	-	-	3	-	5	-	-	1
$\overline{\mathbf{C}}$	-	-	1	8	-	6	4	-	- 1
$\overline{\mathbf{D}}$	-	-	8	1	-	2	9	-	- 1
$\overline{\mathbf{E}}$	7	-	-	-	-	-	-	-	8
$\overline{\mathbf{F}}$	-	-	6	7	-	8	2	-	-
$\overline{\mathbf{G}}$	-	-	2	6	-	9	5	-	- 1
H	8	-	-	2	-	3	-	-	9
I	-	-	5	-	1	-	3	-	-

Relevant filled-in cells: Row:

E1 = 7

E9 = 8

Column:

B4 = 3

C4 = 8

D4 = 1

F4 = 7

G4 = 6

H4 = 2

Since these values are filled-in to cells within either the Row, Column or Box of E4 they can be removed from the domain of E4.

Initial Domain:

DE4 = (1,2,3,4,5,6,7,8,9)

Rows:

Within Row E the values 7 and 8 are already filled and can therefore be removed from the domain of E4. Now the domain of E4 is:

$$DE4 = (1,2,3,4,5,6,9)$$

Column:

The same concept can be applied to the column 4, which already has 3,8,1,7,6,2 filled in, 7 and 8 have already been removed from the domain but the other values can now be removed as well: Now the domain of E4 is:

$$DE4 = (4,5,9)$$

Box:

The same concept can also be applied to the box E4 is in. In this case the box contains 1,2,7,8. However, these values have already been removed from the domain of E4 in previous steps, meaning the domain remains:

3 Implementation

```
from collections import deque
   # Helper function to read Sudoku from a file
   def read_puzzle(file_path):
       with open(file_path, 'r') as f:
5
           return [[int(num) for num in line.split()] for line in f]
   # Helper function to print Sudoku grid
   def print_board(board):
9
       for row in board:
10
            print(" ".join(str(num) if num != 0 else '.' for num in row
11
               ))
       print()
12
   # Get all peers for a given cell (row, col)
14
15
   def get_peers(row, col):
       peers = set()
16
17
       for i in range(9):
           peers.add((row, i)) # Same row
18
           peers.add((i, col)) # Same column
19
       # Same 3x3 box
20
       box_row, box_col = 3 * (row // 3), 3 * (col // 3)
21
       for i in range(box_row, box_row + 3):
           for j in range(box_col, box_col + 3):
                peers.add((i, j))
       peers.discard((row, col)) # Remove the cell itself
25
       return peers
26
27
   # Initialize domains for all cells
28
   def initialize_domains(board):
       domains = {}
30
31
       for row in range (9):
           for col in range(9):
32
                if board[row][col] == 0:
                    domains[(row, col)] = set(range(1, 10))
34
                else:
35
                    domains[(row, col)] = {board[row][col]}
36
       return domains
37
38
39
   # AC-3 algorithm implementation
   def ac3(board, domains):
40
       queue = deque([(cell, peer) for cell in domains for peer in
41
           get_peers(*cell)])
       step = 1 # Step counter
42
43
       while queue:
44
           print(f"Step {step}: Queue length: {len(queue)}") # Print
45
               the step and queue length
            cell, peer = queue.popleft()
```

```
if revise(domains, cell, peer):
47
                if not domains[cell]: # Domain is empty -> failure
                    return False
49
                if len(domains[cell]) == 1:
50
                    board[cell[0]][cell[1]] = next(iter(domains[cell]))
51
                          # Update board
                for neighbor in get_peers(*cell) - {peer}:
                    queue.append((neighbor, cell))
53
            step += 1 # Increment step counter
54
55
       return True
56
   # Revise function to enforce arc-consistency
57
   def revise(domains, cell, peer):
58
59
       revised = False
       if len(domains[peer]) == 1:
60
            peer_value = next(iter(domains[peer]))
61
62
            if peer_value in domains[cell]:
                domains[cell].remove(peer_value)
63
64
                revised = True
       return revised
65
   # Function to find an empty cell
67
   def find_empty(board):
68
69
       for i in range(9):
            for j in range(9):
70
                if board[i][j] == 0:
71
                    return i, j
72
       return None
73
74
   # Backtracking algorithm to solve Sudoku if AC-3 doesn't fully
75
   def backtrack(board. domains):
76
        empty = find_empty(board)
77
       if not empty:
78
                         # No empty cells left, puzzle solved
79
           return True
80
       row, col = empty
       for value in sorted(domains[(row, col)]):
81
            if is_consistent(board, row, col, value):
                board[row][col] = value
83
                # Create a copy of domains for backtracking
84
                new_domains = {key: domains[key].copy() for key in
85
                    domains}
                new_domains[(row, col)] = {value}
86
87
                if backtrack(board, new_domains):
88
89
                    return True
                board[row][col] = 0 # Undo assignment
90
       return False
91
92
   # Function to check if a value can be placed in a cell without
       conflicts
   def is_consistent(board, row, col, val):
94
       for peer_row, peer_col in get_peers(row, col):
95
            if board[peer_row][peer_col] == val:
96
97
                return False
       return True
98
99
```

```
# Solve function
100
    def solve_sudoku(file_path):
        board = read_puzzle(file_path)
102
        print("Original Sudoku:")
103
        print_board(board)
104
106
        domains = initialize_domains(board)
        if ac3(board, domains):
            print("Sudoku after AC-3:")
108
109
            print_board(board)
            if find_empty(board):
                 print("AC-3 did not completely solve the puzzle.
111
                     Applying backtracking...")
                 if not backtrack(board, domains):
                     print("No solution exists.")
113
            else:
114
                 print("Solved using AC-3!")
116
117
            print("AC-3 failed to achieve arc-consistency.")
118
        print("Final Solution:")
119
        print_board(board)
120
121
122
    # Example usage
    solve_sudoku('test_file.txt')
```

4 Input and Outputs

Scenario 1: Can be solved by AC-3 NOTE: The top 10 and bottom 10 steps are shown due to the sheer quantity of steps needed. A better example will be provided during live demo of code.

```
Original Sudoku:
2
   5 3 . . 7 . . . .
   6 . . 1 9 5 . . .
3
     98...6
   8 . . . 6 . . . 3
   4 . . 8 . 3 . . 1
     . . . 2 . . . 6
   . 6 . . . . 2 8 .
     . . 4 1 9 . . 5
     . . . 8 . . 7 9
10
11
   Step 1: Queue length: 1620
12
   Step 2: Queue length: 1619
13
   Step 3: Queue length: 1618
14
   Step 4: Queue length: 1617
   Step 5: Queue length: 1616
   Step 6: Queue length: 1615
   Step 7: Queue length: 1614
18
   Step 8: Queue length: 1613
   Step 9: Queue length: 1612
   Step 10: Queue length: 1611
22
   Step 9363: Queue length: 10
```

```
Step 9364: Queue length: 9
24
   Step 9365: Queue length: 8
   Step 9366: Queue length: 7
   Step 9367: Queue length: 6
27
   Step 9368: Queue length: 5
   Step 9369: Queue length: 4
29
   Step 9370: Queue length: 3
   Step 9371: Queue length: 2
31
   Step 9372: Queue length: 1
32
33
   Sudoku after AC-3:
34
   5 3 4 6 7 8 9 1 2
35
   6 7 2 1 9 5 3 4 8
36
   1 9 8 3 4 2 5 6 7
   8 5 9 7 6 1 4 2 3
38
   4 2 6 8 5 3 7 9 1
39
40
   7 1 3 9 2 4 8 5 6
   9 6 1 5 3 7 2 8 4
41
42
   2 8 7 4 1 9 6 3 5
   3 4 5 2 8 6 1 7 9
43
   Solved using AC-3!
45
   Final Solution:
46
47
   5 3 4 6 7 8 9 1 2
   6 7 2 1 9 5 3 4 8
48
   1 9 8 3 4 2 5 6 7
49
   8 5 9 7 6 1 4 2 3
50
   4 2 6 8 5 3 7 9 1
51
   7 1 3 9 2 4 8 5 6
52
   9 6 1 5 3 7 2 8 4
53
   2 8 7 4 1 9 6 3 5
   3 4 5 2 8 6 1 7 9
```

Scenario 2: Can't be solved NOTE: The top 10 and bottom 10 steps are shown due to the sheer quantity of steps needed. A better example will be provided during live demo of code.

```
Original Sudoku:
   . . . . . 6 . . .
   2 . . . 3 . . . 1
   . . 9 . . . 8 . .
     . 5 . . .
               3 .
   . . 4 9 . 2 5 . .
   . . 3 . . . 7 . .
   . . 7 . . . 4 . .
   8 . . . 5 . . . 9
9
   . . . 7 . .
11
   Step 1: Queue length: 1620
12
   Step 2: Queue length: 1619
   Step 3: Queue length: 1618
14
   Step 4: Queue length: 1617
   Step 5: Queue length: 1635
   Step 6: Queue length: 1653
   Step 7: Queue length: 1671
   Step 8: Queue length: 1670
```

```
Step 9: Queue length: 1669
   Step 10: Queue length: 1668
22
   Step 6798: Queue length: 10
23
   Step 6799: Queue length: 9
24
   Step 6800: Queue length: 8
25
   Step 6801: Queue length: 7
   Step 6802: Queue length: 6
   Step 6803: Queue length: 5
   Step 6804: Queue length: 4
   Step 6805: Queue length: 3
30
   Step 6806: Queue length: 2
   Step 6807: Queue length: 1
32
   Sudoku after AC-3:
34
   . . . . . 6 . . .
35
36
   2 . . . 3 . . . 1
   . . 9 . . . 8 . .
37
   . . 5 . . . 3 . .
   . . 4 9 . 2 5 . .
39
   . . 3 . . . 7 . .
   . . 7 . . . 4 . .
41
   8 . . . 5 . . . 9
42
       . 7 . .
43
44
   AC-3 did not completely solve the puzzle. Applying backtracking...
   No solution exists.
46
   Final Solution:
47
   . . . . . 6 . .
48
   2 . . . 3 . . . 1
49
   . . 9 . . . 8 . .
   . . 5 . . . 3 . .
51
   . . 4 9 . 2 5 . .
   . . 3 . . . 7 . .
53
   . . 7 . . . 4 . .
54
55
          . 5 . . . 9
     . . 7 . . . . .
```

Scenario 3: Can be solved by backtracking NOTE: The top 10 and bottom 10 steps are shown due to the sheer quantity of steps needed. A better example will be provided during live demo of code.

```
Original Sudoku:
   1 . 5 . 7 . . . .
   . . 4 . . 3 9 . .
   . . . . . . . 6 8
   . 9 . . . 7 . 2 .
   . . . . . 2 8 . 7
   . . 3 . 4 . . . .
     . 8 . . . . 3 .
   9 4 . 5 . . . . .
9
10
   . . . . . 4 2 . .
11
   Step 1: Queue length: 1620
   Step 2: Queue length: 1619
  Step 3: Queue length: 1618
```

```
| Step 4: Queue length: 1617
15
   Step 5: Queue length: 1616
  Step 6: Queue length: 1615
   Step 7: Queue length: 1614
18
   Step 8: Queue length: 1613
   Step 9: Queue length: 1612
20
   Step 10: Queue length: 1611
   Step 7482: Queue length: 10
   Step 7483: Queue length: 9
   Step 7484: Queue length: 8
25
   Step 7485: Queue length: 7
   Step 7486: Queue length: 6
   Step 7487: Queue length: 5
   Step 7488: Queue length: 4
   Step 7489: Queue length: 3
30
   Step 7490: Queue length: 2
   Step 7491: Queue length: 1
   Sudoku after AC-3:
34
35
   1 . 5 . 7 . 3 4 2
   . . 4 . . 3 9 . .
36
   . . . . . . . 6 8
37
   . 9 . . . 7 . 2 .
   . . . . . 2 8 . 7
39
   . . 3 . 4 . . . .
   . . 8 . . . . 3 .
41
   94.5....
42
   . . . . . 4 2 . .
43
44
   AC-3 did not completely solve the puzzle. Applying backtracking...
  Final Solution:
  1 8 5 9 7 6 3 4 2
   7 6 4 2 8 3 9 5 1
48
   3 2 9 4 1 5 7 6 8
49
   8 9 1 6 5 7 4 2 3
  4 5 6 3 9 2 8 1 7
51
  2 7 3 8 4 1 6 9 5
   6 1 8 7 2 9 5 3 4
53
54
   9 4 2 5 3 8 1 7 6
   5 3 7 1 6 4 2 8 9
```