

## Lecture 12. Logical Equivalence (Section 1.3.1-1.3.4)

## The converse, contrapositive, and inverse

A few statements related to  $p \rightarrow q$ :

- The **converse** of  $p \rightarrow q$  is  $q \rightarrow p$ .
- The **contrapositive** of  $p \rightarrow q$  is  $\neg q \rightarrow \neg p$ .
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- Converse: If graph G doesn't contain an odd length cycle, it can be bipartite.
- Contrapositive: If graph G contain an odd length cycle, it is not bipartite.

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One of the three is equivalent to the original conditional statement, WHICH ONE?  
How do you know?

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$p$	$q$	$p \rightarrow q$	$q \rightarrow p$	$\neg p$	$\neg q$	$\neg q \rightarrow \neg p$	$\neg p \rightarrow \neg q$

- The contrapositive is equivalent to the original conditional statement:

$$p \rightarrow q \Leftrightarrow \neg q \rightarrow \neg p$$

- The converse and inverse are equivalent to each other.

$$q \rightarrow p \Leftrightarrow \neg p \rightarrow \neg q$$

**Example 12.2.** Consider the following theorem from secondary school geometry.

*If a quadrilateral has a pair of parallel sides, then it has a pair of supplementary angles.*<sup>1</sup>

This theorem is of the form  $p \rightarrow q$ . Determine  $p$  and  $q$  and write the theorem in its contrapositive form,  $\neg q \rightarrow \neg p$ , which is logically equivalent to  $p \rightarrow q$

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**Practice 12.3.** Determine if the following compound statements are logically equivalent.

- (a)  $p \rightarrow q$  and  $\neg p \vee q$
- (b)  $\neg(p \rightarrow q)$  and  $\neg p \rightarrow \neg q$

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- The **print** statement is carried out when the *proposition* following the word **if** is *True*.
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- On closer inspection, this big expression is built from two simpler propositions.  
A:  $x > 0$   
B:  $y > 100$
- Then we can rewrite the 'if' condition as  $A \vee (\neg A \wedge B)$ .
- A truth table reveals that this complicated expression is logically equivalent to (what?).



- (Continue from the previous example) ...  $A \vee (\neg A \wedge B)$ . is logically equivalent to (what?).

A	B	$\neg A$	$\neg A \wedge B$	$A \vee (\neg A \wedge B)$	(What?)
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- This means that we can simplify the code snippet without changing the program's behavior:

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### Remarks:

- Rewriting a logical expression involving many variables in the simplest form is both difficult and important.
- Simplifying expressions in software can increase the speed of your program.

**Practice 12.5.** Consider the proposition “If AI takes over the World or outsmarts people, AI will get the ability to feel emotion.” Write “E” for each proposition that is logically equivalent to the given proposition, “C” for each proposition that is logically equivalent to the converse of the given proposition and “N” if neither.

- (a) Unless AI gains the ability to feel an emotion, AI will not dominate the world and cannot surpass humans.
- (b) AI does not take over the World and outsmart people, or AI will get the ability to feel emotion.
- (c) AI takes over the World or outsmarts people, and AI will not get the ability to feel an emotion.
- (d) AI will not get the ability to feel an emotion if AI does not take over the World and outsmart people.
- (e) AI does not take over the World and outsmart people, and AI will not get the ability to feel an emotion.