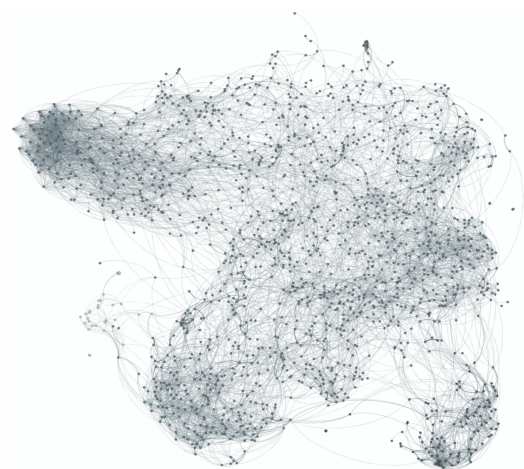


Lecture 2. Fundamental Concepts and Terminologies

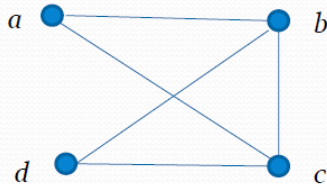


Graph of Harry Potter Fanfiction

Definition 2.1 A graph $G = (V, E)$ consists of a nonempty set V of **vertices** (or **nodes**) and a set E of **edges**.

Example:

This is a graph
with four
vertices and five
edges.



where $V = \{a, b, c, d\}$ and $E = \{(a, b), (a, c), (b, d), (b, c), (d, c)\}$.

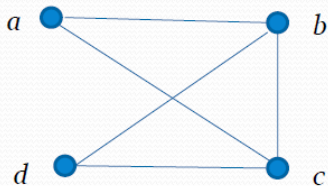
Remark: Since the graph is undirected, $(a, b) = (b, a)$, $(a, c) = (c, a)$, etc.

Neighbors: vertices u and v are **neighbors** (or called **adjacent**) if an edge (u, v) connects them.

Example 2.2. b and c are neighbors.

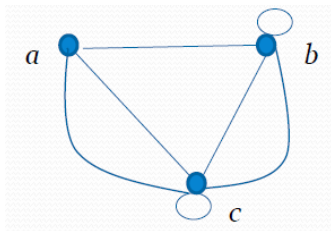
Example:

This is a graph with four vertices and five edges.



Example 1.3. All neighbors of the node $a = \{b, c\}$.

An edge that connects a vertex to itself is called a **loop**.

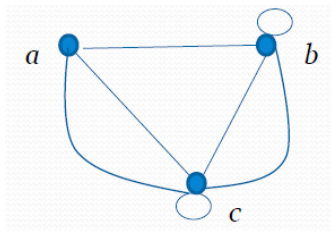


The **degree** of a vertex is the number of edges connected to the vertex ¹, denoted by $\deg(u)$. For example, in the figure above

- $\deg(a) =$
- $\deg(b) =$
- $\deg(c) =$

¹Equivalently, the number of neighbors of the vertex.

An edge that connects a vertex to itself is called a **loop**.



The **degree** of a vertex is the number of edges connected to the vertex ¹, denoted by $\deg(u)$. For example, in the figure above

- $\deg(a) =$
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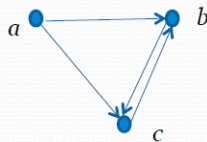
Notice that a loop at a vertex contributes two to the degree of that vertex.

¹Equivalently, the number of neighbors of the vertex.

Definition 2.4 An **directed graph** $G = (V, E)$ consists of a nonempty set V of vertices and a set E of directed edges.

Example:

This is a directed graph with three vertices and four edges.

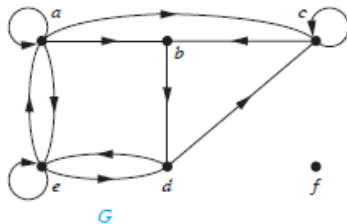


In a directed graph,

- **Indegree** = the number of edges coming in to the vertex v ($\deg^-(v)$)
- **Outdegree** = the number of edges going out of the vertex v ($\deg^+(v)$)

Remark: $(a, b) \neq (b, a)$, $(b, c) \neq (c, b)$, etc.

Example 2.5. Find the in-degree and out-degree of each vertex in the graph G with directed edges shown in



A **path** in a graph G is a sequence of vertices v_i and edges e_i such that the edge e_i connects vertices v_{i-1} and v_i :

$$(v_0, e_1, v_1, e_2, v_2, \dots, v_{n-1}, e_n, v_n)$$

Equivalently, assuming that two consecutive vertices are connected by an edge, we may write it without e_i 's

$$(v_0, v_1, v_2, \dots, v_{n-1}, v_n)$$

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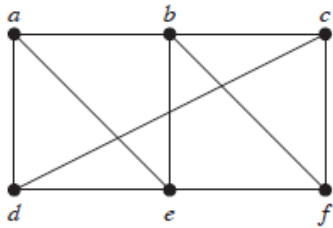
$$(v_0, e_1, v_1, e_2, v_2, \dots, v_{n-1}, e_n, v_n)$$

Equivalently, assuming that two consecutive vertices are connected by an edge, we may write it without e_i 's

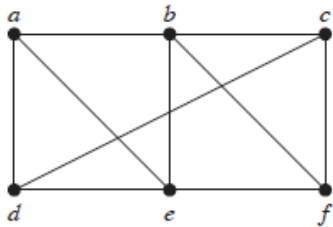
$$(v_0, v_1, v_2, \dots, v_{n-1}, v_n)$$

- The **length of a path** is the number of edges in the path.
- A **circuit** (or **cycle**) is a path that starts and ends on the same vertex ($v_0 = v_n$).
- A path or circuit is **simple** if it does not contain the same edge more than once.

Example 2.6. In the graph

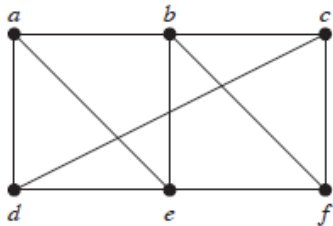


Example 2.6. In the graph



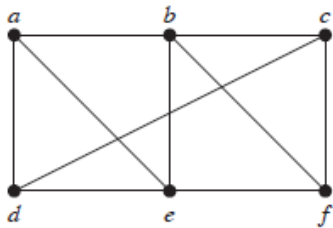
- a, d, c, f, e is a simple path of length 4.

Example 2.6. In the graph



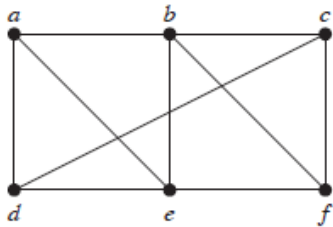
- a, d, c, f, e is a simple path of length 4.
- d, e, c, a is not a path because e is not connected to c .

Example 2.6. In the graph



- a, d, c, f, e is a simple path of length 4.
- d, e, c, a is not a path because e is not connected to c .
- b, c, f, e, b is a circuit of length 4.

Example 2.6. In the graph



- a, d, c, f, e is a simple path of length 4.
- d, e, c, a is not a path because e is not connected to c .
- b, c, f, e, b is a circuit of length 4.
- a, b, e, d, a, b is a path of length 5, but not a simple path.

Example 2.7

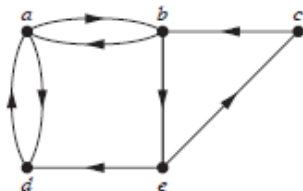
2. Does each of these lists of vertices form a path in the following graph? Which paths are simple? Which are circuits? What are the lengths of those that are paths?

a) a, b, e, c, b

b) a, d, a, d, a

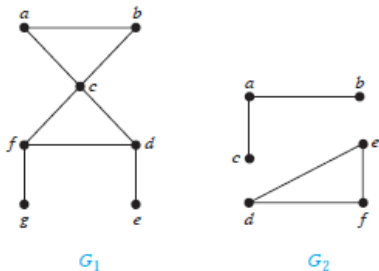
c) a, d, b, e, a

d) a, b, e, c, b, d, a



- An **undirected** graph is **connected** if there is a path connecting any two vertices.
- A **directed** graph is **connected** if the underlying undirected graph is.

- An **undirected** graph is **connected** if there is a path connecting any two vertices.
- A **directed** graph is **connected** if the underlying undirected graph is.
- **Example 2.8.** G_1 is connected because there is a path between any pair of its vertices, as can be easily seen. However G_2 is not connected because there is no path between vertices a and f , for example.



- **Connected component:** A subset of vertices $V_i \subset V$ that is connected. For example, in the above graph, $V_1 = \{a, b, c\}$ and $V_2 = \{d, e, f\}$ are two connected components of G_2 .

Problem 2.9. Is the graph below connected or not?

