## Don Bosco Institute of Technology, Kulra(W) Department of Computer Engineering CSL601: System Programming and Compiler Construction Lab2022-23

03
For any Assembly input for a hypothetical machine, implement pass 2 by using the intermediate code generated as input to pass 2 program which was generated as output in pass 1.
Ronak Surve
64
<ol> <li>Students will be able to learn various file handling operations</li> <li>Students will be able to produce the target code using the intermediate representation generated in Pass 1.</li> </ol>
3) Students will be able to implement the working of Pass 2 of 2 pass Assembler
The theory behind a two-pass assembler is to separate the tasks of assembling a source code written in assembly language into two distinct stages, each with a specific objective. The two passes are:  Pass 1: This pass scans the source code to gather information about the labels used in the code, their addresses, and the size of the program. This information is recorded in a symbol table, which is used in the second pass to resolve symbolic references and replace them with absolute addresses. The main objective of the first pass is to gather information and prepare the necessary data for the second pass.
Pass 2: This pass uses the information gathered in the first pass, such as the symbol table, to generate the final object code. During the second pass, the symbolic references are replaced with the corresponding absolute addresses, and the final object code is generated. The main objective of the second pass is to generate the final object code and perform error checking.  The theory behind the two-pass approach is to optimize the accuracy and efficiency of the assembly process by dividing it into two separate stages. The first pass can detect errors and report them to the user before the final object code is generated, allowing for easier correction and debugging. The second pass can generate optimized and efficient object code using the information gathered in the first pass.  Overall, the two-pass assembler theory is based on the idea of dividing the assembly process into two separate stages,

```
efficiency of the assembly process.
Program Code:
                        # Initialisation
                        MOT = {
                          "MOVEM": 1,
                          "MOVER": 2,
                          "ADD": 3,
                          "SUB": 4,
                          "MUL": 5,
                          "DIV": 6,
                          "BC": 7,
                          "COMP": 8,
                          "READ": 9,
                          "PRINT": 10
                        POT = {
                          "START": 1,
                          "END": 2,
                          "EQU": 3,
                          "ORIGIN": 4,
                          "LTORG": 5
                        DL = {
                          "DS": 1,
                          "DC": 2
                       REG = {
                          "AREG": 1,
                          "BREG": 2,
                          "CREG": 3,
                          "DREG": 4
                        SYMTAB = []
                        LITTAB = []
                        POOLTAB = []
                        pool_1st = 1
                        ASM = []
                       # Initialisation complete
                        # Take input program
                        ASM.append(input("Start typing the assembly program \n"))
                        while ASM[-1].upper() != "END":
                          ASM.append(input())
                        print("Input taken")
                        # Pass 1
                        print("Starting Pass 1\n")
                        loc = 0
                        ic = []
                        for line in ASM:
                          cmd = line.split(" ")
                          if cmd[0].upper() in MOT:
```

```
ic.append([f"{loc}", f"(IS, {MOT[cmd[0].upper()]})"])
    loc += 1
    if cmd[1].upper().replace(",", "") in REG: # e.g. MOVER AREG, B
       ic[-1].append(f"{REG[cmd[1].upper().replace(',', ")]}")
    if cmd[2] is None: # e.g. READ A
      continue
    elif cmd[2][0] == "=": # e.g. MOVEM CREG, ='5'
       LITTAB.append([f"{len(LITTAB) + 1}", f"{cmd[2]}", ""])
       ic[-1].append(f"(L, {len(LITTAB)})")
    else: # e.g. MOVER AREG, B
      SYMTAB.append([f"{len(SYMTAB) + 1}", f"{cmd[2]}", ""])
      ic[-1].append(f"(S, {len(SYMTAB)})")
  elif cmd[0].upper() in POT:
    if POT[cmd[0].upper()] == 1: # START
       loc = int(cmd[1])
      ic.append(["", "(AD, 1)", "", f"(C, {loc})"])
    elif POT[cmd[0].upper()] == 2: # END
      count = 0
       for record in LITTAB:
         if record[-1] == "": # Find all records with empty address fields
           record.pop(-1)
           record.append(f"{loc}")
           ic.append([f"{loc}", "(AD, 2)", "", f"{record[1][-2]}"])#
'record[1][-2]' is value of literal
           count += 1
           loc += 1
       POOLTAB.append([f"{pool_1st}", f"{count}"])
       pool_1st += count
       break
    elif POT[cmd[0].upper()] == 4: # ORIGIN
      ic.append([f"{loc}", "No Intermediate code for ORIGIN Assembler
Directive"])
      loc = int(cmd[1]) # Go to new address
    elif POT[cmd[0].upper()] == 5: # LTORG
       count = 0
       for record in LITTAB:
         if record[-1] == "": # Find all records with empty address fields
           record.pop(-1)
           record.append(f"{loc}")
           ic.append([f"{loc}", "(AD, 5)", "", f"{record[1][-2]}"])#
'record[1][-2]' is value of literal
           count += 1
           loc += 1
       POOLTAB.append([f"{pool_1st}", f"{count}"])
       pool 1st += count
 else:
    isPresent = False
    for record in SYMTAB:
      if record[1] == cmd[0].upper():
         record.pop(-1)
         record.append(f"{loc}")
         isPresent = True
         break
```

```
if not isPresent:
       SYMTAB.append([f"{len(SYMTAB) + 1}", f"{cmd[0]}",
f"{loc}"])
    if cmd[1].upper() in DL:
       ic.append([f"{loc}", f"(DL, {DL[cmd[1].upper()]})", "",
f"{cmd[2]}"])
    elif cmd[1].upper() in MOT:
       ic.append([f"{loc}", f"(IS, {MOT[cmd[1].upper()]})"])
       if cmd[2].upper().replace(",", "") in REG: # e.g. MOVER AREG, B
         ic[-1].append(f"{REG[cmd[2].upper().replace(',', ")]}")
       if cmd[3] is None: # e.g. READ A
         continue
       elif cmd[3][0] == "=": # e.g. MOVEM CREG, ='5'
         LITTAB.append([f"{len(LITTAB) + 1}", f"{cmd[3]}", ""])
         ic[-1].append(f"(L, {len(LITTAB)})")
       else: # e.g. MOVER AREG, B
         SYMTAB.append([f"{len(SYMTAB) + 1}", f"{cmd[3]}", ""])
         ic[-1].append(f"(S, {len(SYMTAB)})")
    elif POT[cmd[1].upper()] == 3: # EOU
       ic.append([loc, "No Intermediate code for EQU Assembler
Directive"])
       SYMTAB[-1].pop(-1)
       for record in SYMTAB:
         if record[1] == cmd[2].upper(): # Change 1st symbol's address to
2nd symbol's address
            SYMTAB[-1].append(record[2])
    loc += 1
for line in ic:
  print(line)
# Pass 2
print("\nStarting Pass 2\n")
mc = []
for record in ic:
  if record[0] == "": # Skip the START statement
     continue
  if record[1][0] != '(': # Skip the EQU & ORIGIN statements
  statement = record[1].replace('(',").replace(')',").split(' ') # Split statement
into statement type and its value
  operand = record[3].replace('(',").replace(')',").split(' ') # Split operand 2
into literal/ symbol type and its value
  if len(operand) > 1: # if operand is a value, skip
    record.pop(3)
    if operand[0] == 'S,':
       record.append(SYMTAB[int(operand[1]) - 1][-1]) # Find address
location
     elif operand[0] == 'L,':
       record.append(LITTAB[int(operand[1]) - 1][-1]) # Find address
  mc.append([f"{record[0]}+", statement[1] if statement[0] == 'IS,' else ",
record[2], record[3]])
```

```
for line in mc:
                                                                    print(line)
                                                               print("\nSymbol Table\n")
                                                               for line in SYMTAB:
                                                                    print(line)
                                                               print("\nLiteral Table\n")
                                                               for line in LITTAB:
                                                                    print(line)
                                                               print("\nPool Table\n")
                                                               for line in POOLTAB:
                                                                    print(line)
Input to the
                                                                  Input taken
                                                                  Starting Pass 1
Program:
                                                               ['', '(AD, 1)', '', '(C, 100)']
['100', '(DL, 2)', '', '10']
['101', '(IS, 2)', '1', '(S, 2)']
['102', '(IS, 1)', '2', '(L, 1)']
['103', '(IS, 3)', '1', '(L, 2)']
['104', '(IS, 4)', '2', '(L, 3)']
['105', '(DL, 2)', '', '20']
['106', 'No Intermediate code for ORIGIN Assembler Directive']
['300', '(AD, 5)', '', '1']
['301', '(AD, 5)', '', '2']
['302', '(AD, 5)', '', '1']
['303', '(IS, 2)', '1', '(S, 3)']
['304', '(IS, 2)', '3', '(S, 4)']
['305', '(IS, 3)', '2', '(L, 4)']
['306', '(DL, 1)', '', '5']
['307', '(DL, 2)', '', '10']
['308', '(AD, 2)', '', '1']
```

```
Output of the
                              Starting Pass 2
program:
                              ['100+', '', '', '10']
                              ['101+', '2', '1', '105']
                              ['102+', '1', '2', '300']
                              ['103+', '3', '1', '301']
                               ['104+', '4', '2', '302']
                              ['105+', '', '', '20']
                              ['300+', '', '', '1']
                              ['301+', '', '', '2']
                              ['302+', '', '', '1']
                              ['303+', '2', '1', '306']
                              ['304+', '2', '3', '307']
                              ['305+', '3', '2', '308']
                              ['306+', '', '', '5']
['307+', '', '', '10']
                              ['308+', '', '', '1']
                              Symbol Table
                               ['1', 'A', '100']
                               ['2', 'B', '105']
                              ['3', 'NUM', '306']
                              ['4', 'LOOP', '307']
                              Literal Table
                              ['1', "='1'", '300']
                              ['2', "='2'", '301']
                              ['3', "='1'", '302']
                              ['4', "='1'", '308']
                              Pool Table
                              ['1', '3']
                              ['4', '1']
                       The outcome of a pass 2 assembler is the final object code, which is a
Outcome of the
                       machine-readable representation of the
Experiment:
                       source code written in assembly language. During the second pass, the
                       assembler uses the information gathered
                       in the first pass, such as the symbol table and the listing of the source code,
                       to generate the object code. The
                       final object code is typically stored in a file that can be loaded into memory
                       and executed by the computer. The
                       second pass also performs error checking and generates error messages for
                       any problems detected in the source code,
                       such as invalid operands or undefined symbols. The final object code
                       generated by the pass 2 assembler can also be
                       optimized for size or performance by applying various code optimization
                       https://www.geeksforgeeks.org/single-pass-two-pass-and-multi-pass-
References:
                       compilers/
```

Course in-charge-Mayura Gavhane