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Lecture 4

Systems of Linear Equations

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Systems of Linear Equations

⊙ Motivation: from Mechanics

The vertical displacement of the beam is represented by a function $y(x)$, where $0 \leq x \leq L$ along the beam of length L .

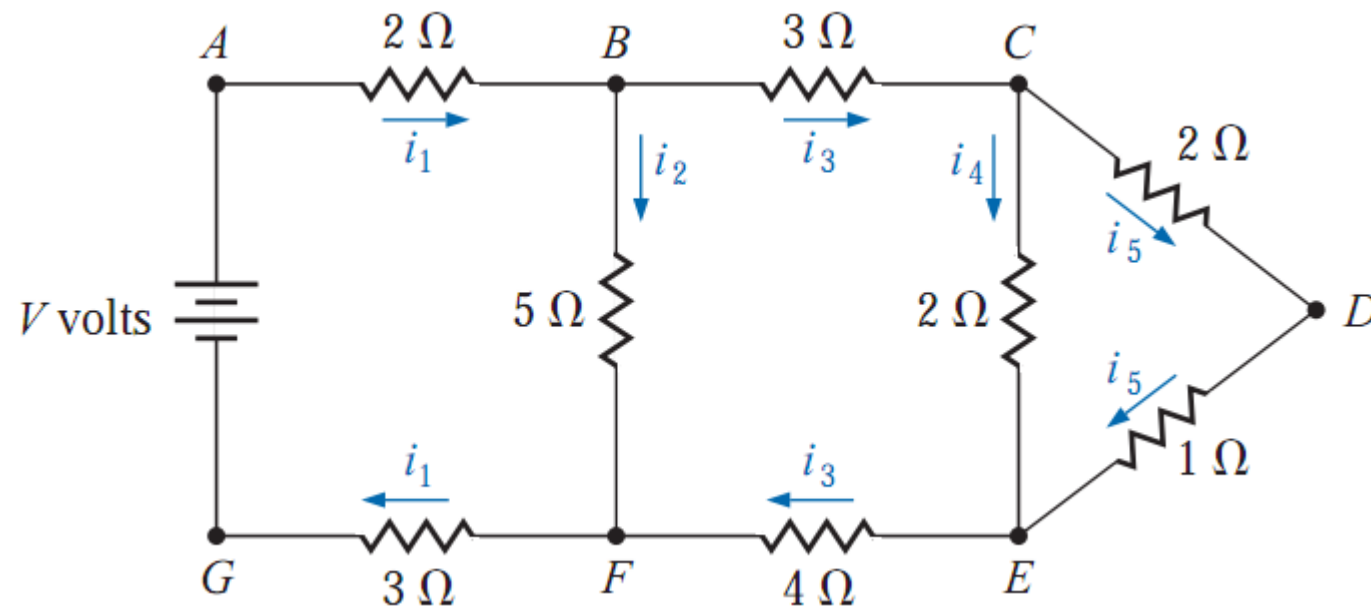
$$EI y'''' = f(x) \quad y_{i-2} - 4y_{i-1} + 6y_i - 4y_{i+1} + y_{i+2} = \frac{h^4}{EI} f(x_i)$$

$$y(0) = y'(0) = y''(L) = y'''(L) = 0$$

$$\begin{bmatrix} 16 & -9 & \frac{8}{3} & -\frac{1}{4} & & & & \\ -4 & 6 & -4 & 1 & & & & \\ & 1 & -4 & 6 & -4 & 1 & & \\ & & 1 & -4 & 6 & -4 & 1 & \\ & & & \ddots & \ddots & \ddots & \ddots & \ddots \\ & & & & 1 & -4 & 6 & -4 & 1 \\ & & & & & 1 & -4 & 6 & -4 & 1 \\ & & & & & & \frac{16}{17} & -\frac{60}{17} & \frac{72}{17} & -\frac{28}{17} \\ & & & & & & -\frac{12}{17} & \frac{96}{17} & -\frac{156}{17} & \frac{72}{17} \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ \vdots \\ \vdots \\ y_{n-1} \\ y_n \end{bmatrix} = \frac{h^4}{EI} \begin{bmatrix} f(x_1) \\ f(x_2) \\ \vdots \\ \vdots \\ \vdots \\ f(x_{n-1}) \\ f(x_n) \end{bmatrix}$$

Systems of Linear Equations

Motivation: from Electrical Circuits



$$5i_1 + 5i_2 = V,$$

$$i_3 - i_4 - i_5 = 0,$$

$$2i_4 - 3i_5 = 0,$$

Kirchhoff's laws of electrical circuits:

$$i_1 - i_2 - i_3 = 0,$$

$$5i_2 - 7i_3 - 2i_4 = 0.$$

Systems of Linear Equations

References

[1] Timothy Sauer, Numerical analysis (2nd ed.), Pearson Education, 2012. **Chapter 2**

[2] Cleve Moler, Numerical Computing with MATLAB, Society for Industrial and Applied Mathematics, 2004. **Chapter 2**

[2] Richard L. Burden, J. Douglas Faires, Numerical analysis (9th ed.), Brooks/Cole, 2011. **Chapter 6**

Systems of Linear Equations

□ Direct Methods

- Gaussian Elimination
- The LU Factorization
- Sources of Error
- The $PA = LU$ Factorization

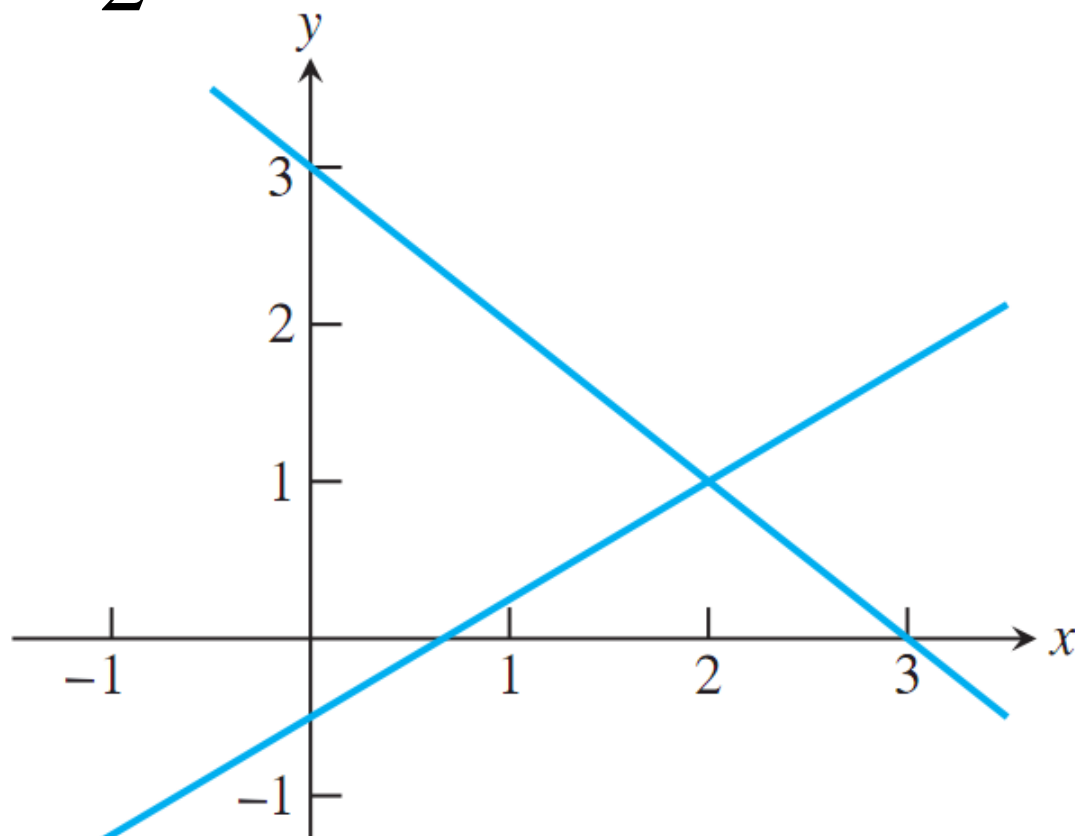
□ Iterative Methods

- Jacobi Method
- Gauss–Seidel Method

Direct Methods

Gaussian Elimination: Example 1

$$\begin{cases} x + y = 3 \\ 3x - 4y = 2 \end{cases}$$



Direct Methods

⊙ Gaussian Elimination: Example 1

$$\begin{cases} x + y = 3 \\ 3x - 4y = 2 \end{cases}$$



Subtracting 3 · [x + y = 3] from the second equation

$$\begin{cases} x + y = 3 \\ -7y = -7 \end{cases} \quad \longrightarrow \quad -7y = -7 \longrightarrow y = 1$$



$$x + y = 3 \longrightarrow x + (1) = 3 \longrightarrow x = 2$$

Direct Methods

Gaussian Elimination: Example 1

$$\begin{cases} x + y = 3 \\ 3x - 4y = 2 \end{cases}$$

The same elimination work can be done in the absence of variables by writing the system in tableau form:

$$\left[\begin{array}{cc|c} 1 & 1 & 3 \\ 3 & -4 & 2 \end{array} \right] \xrightarrow{\substack{\text{subtract } 3 \times \text{row } 1 \\ \text{from row } 2}} \left[\begin{array}{cc|c} 1 & 1 & 3 \\ 0 & -7 & -7 \end{array} \right]$$

Direct Methods

⊙ Gaussian Elimination: Example 2

$$\begin{array}{rcl}
 x + 2y - z & = & 3 \\
 2x + y - 2z & = & 3 \\
 -3x + y + z & = & -6
 \end{array}
 \quad \longleftrightarrow \quad
 \left[\begin{array}{ccc|c}
 1 & 2 & -1 & 3 \\
 2 & 1 & -2 & 3 \\
 -3 & 1 & 1 & -6
 \end{array} \right]$$

To eliminate column 1:

$$\begin{array}{lcl}
 \text{subtract } 2 \times \text{row 1} & & \\
 \longrightarrow \quad \text{from row 2} & \longrightarrow & \left[\begin{array}{ccc|c}
 1 & 2 & -1 & 3 \\
 0 & -3 & 0 & -3 \\
 -3 & 1 & 1 & -6
 \end{array} \right] \\
 \\
 \text{subtract } -3 \times \text{row 1} & & \\
 \longrightarrow \quad \text{from row 3} & \longrightarrow & \left[\begin{array}{ccc|c}
 1 & 2 & -1 & 3 \\
 0 & -3 & 0 & -3 \\
 0 & 7 & -2 & 3
 \end{array} \right]
 \end{array}$$

Direct Methods

⊙ Gaussian Elimination: Example 2

\rightarrow subtract $-3 \times$ row 1
 from row 3 \rightarrow

$$\left[\begin{array}{ccc|c} 1 & 2 & -1 & 3 \\ 0 & -3 & 0 & -3 \\ 0 & 7 & -2 & 3 \end{array} \right]$$

To eliminate column 2:

\rightarrow subtract $-\frac{7}{3} \times$ row 2
 from row 3 \rightarrow

$$\left[\begin{array}{ccc|c} 1 & 2 & -1 & 3 \\ 0 & -3 & 0 & -3 \\ 0 & 0 & -2 & -4 \end{array} \right]$$

$$x + 2y - z = 3$$

$$x = 3 - 2y + z$$



$$-3y = -3$$



$$-3y = -3$$

$$-2z = -4$$

$$-2z = -4$$

Direct Methods

Gaussian Elimination: Operation Counts

For any positive integer n ,

$$1 + 2 + 3 + 4 + \cdots + n = n(n + 1)/2$$

$$1^2 + 2^2 + 3^2 + 4^2 + \cdots + n^2 = n(n + 1)(2n + 1)/6$$

Direct Methods

⊙ Gaussian Elimination: Operation Counts

The general form of the tableau for n equations in n unknowns is

$$\left[\begin{array}{cccc|c} a_{11} & a_{12} & \dots & a_{1n} & b_1 \\ a_{21} & a_{22} & \dots & a_{2n} & b_2 \\ \vdots & \vdots & \dots & \vdots & \vdots \\ a_{n1} & a_{n2} & \dots & a_{nn} & b_n \end{array} \right]$$

To carry out the elimination step:

```
for j = 1 : n-1
    for i = j+1 : n
        eliminate entry a(i,j)
    end
end
```

Direct Methods

⊙ Gaussian Elimination: Operation Counts

$$\begin{array}{cccc|c} a_{11} & a_{12} & \dots & a_{1n} & b_1 \\ a_{21} & a_{22} & \dots & a_{2n} & b_2 \end{array}$$



to eliminate the a_{21} entry

$$\begin{array}{cccc|c} a_{11} & a_{12} & \dots & a_{1n} & b_1 \\ 0 & a_{22} - \frac{a_{21}}{a_{11}}a_{12} & \dots & a_{2n} - \frac{a_{21}}{a_{11}}a_{1n} & b_2 - \frac{a_{21}}{a_{11}}b_1 \end{array}$$

this requires **one division** (to find the multiplier a_{21}/a_{11}),
plus **n multiplications** and **n additions**

Eliminating each entry a_{i1} in the first column uses
 $2n + 1$ operations.

Direct Methods

⊙ Gaussian Elimination: Operation Counts

The row operation used to eliminate entry a_{ij} is

$$\begin{array}{ccccccc|c}
 0 & 0 & a_{jj} & a_{j,j+1} & \dots & a_{jn} & & b_j \\
 \vdots & \vdots & \vdots & \vdots & \dots & \vdots & & \vdots \\
 0 & 0 & 0 & a_{i,j+1} - \frac{a_{ij}}{a_{jj}}a_{j,j+1} & \dots & a_{in} - \frac{a_{ij}}{a_{jj}}a_{jn} & & b_i - \frac{a_{ij}}{a_{jj}}b_j
 \end{array}$$

This requires **one division, $n - j + 1$ multiplications,**
and $n - j + 1$ addition/subtractions.

Direct Methods

⊙ Gaussian Elimination: Operation Counts

The elimination of each a_{ij} requires the following number of operations, including divisions, multiplication, and addition/subtractions:

$$\begin{bmatrix} 0 \\ 2n+1 & 0 \\ 2n+1 & 2(n-1)+1 & 0 \\ 2n+1 & 2(n-1)+1 & 2(n-2)+1 & 0 \\ \vdots & \vdots & \vdots & \ddots & \ddots \\ \vdots & \vdots & \vdots & & \\ 2n+1 & 2(n-1)+1 & 2(n-2)+1 & \cdots & 2(3)+1 & 0 \\ 2n+1 & 2(n-1)+1 & 2(n-2)+1 & \cdots & 2(3)+1 & 2(2)+1 & 0 \end{bmatrix}$$

Direct Methods

Gaussian Elimination: Operation Counts

Starting on the right, we total up the operations as:

$$\begin{aligned}\sum_{j=1}^{n-1} \sum_{i=1}^j 2(j+1) + 1 &= \sum_{j=1}^{n-1} 2j(j+1) + j \\&= 2 \sum_{j=1}^{n-1} j^2 + 3 \sum_{j=1}^{n-1} j = 2 \frac{(n-1)n(2n-1)}{6} + 3 \frac{(n-1)n}{2} \\&= (n-1)n \left[\frac{2n-1}{3} + \frac{3}{2} \right] = \frac{n(n-1)(4n+7)}{6} \\&= \frac{2}{3}n^3 + \frac{1}{2}n^2 - \frac{7}{6}n\end{aligned}$$

Direct Methods

⊙ Gaussian Elimination: Operation Counts

After the elimination is completed, the tableau is:

$$\left[\begin{array}{cccc|c} a_{11} & a_{12} & \cdots & a_{1n} & b_1 \\ 0 & a_{22} & \cdots & a_{2n} & b_2 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & \cdots & a_{nn} & b_n \end{array} \right]$$



$$a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n = b_1$$

$$a_{22}x_2 + \cdots + a_{2n}x_n = b_2$$

$$\vdots$$

$$a_{nn}x_n = b_n$$

where the a_{ij} refer to the revised, not original, entries.

Direct Methods

⊙ Gaussian Elimination: Operation Counts

To carry out the back-substitution step:

$$x_1 = \frac{b_1 - a_{12}x_2 - \cdots - a_{1n}x_n}{a_{11}}$$

$$x_2 = \frac{b_2 - a_{23}x_3 - \cdots - a_{2n}x_n}{a_{22}}$$

$$\vdots$$

$$x_n = \frac{b_n}{a_{nn}}.$$

Direct Methods

⊙ Gaussian Elimination: Operation Counts

We start at the bottom and work our way up to the top equation.

Counting operations yields:

$$\begin{aligned} 1 + 3 + 5 + \cdots + (2n - 1) &= \sum_{i=1}^n 2i - 1 \\ &= 2 \sum_{i=1}^n i - \sum_{i=1}^n 1 = 2 \frac{n(n+1)}{2} - n = n^2 \end{aligned}$$

Direct Methods

• Gaussian Elimination: Operation Counts

For a system of n equations in n variables

- Operation count for the elimination step of Gaussian elimination:

$$\frac{2}{3}n^3 + \frac{1}{2}n^2 - \frac{7}{6}n$$

- Operation count for the back-substitution step of Gaussian elimination:

$$n^2$$

Direct Methods

Gaussian Elimination: Operation Counts

Estimate the time required to carry out back substitution on a system of 500 equations in 500 unknowns, on a computer where elimination takes 1 second.

$$\frac{(500)^2}{2(500)^3/3} = \frac{3}{2(500)} = 0.003 \text{ sec}$$

Smaller powers of n in operation counts can often be safely neglected.

Direct Methods

⊙ LU Factorization: Example 1 revisited

$$\begin{cases} x + y = 3 \\ 3x - 4y = 2 \end{cases}$$



in matrix form

$$\begin{bmatrix} 1 & 1 \\ 3 & -4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 2 \end{bmatrix}$$

Direct Methods

⊙ LU Factorization: Example 1 revisited

$$\begin{bmatrix} 1 & 1 \\ 3 & -4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 2 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 1 \\ 3 & -4 \end{bmatrix} \rightarrow \text{subtract } 3 \times \text{row 1} \\ \text{from row 2} \rightarrow \begin{bmatrix} 1 & 1 \\ 0 & -7 \end{bmatrix} = U$$

To store the multiplier 3 used in the elimination step:

$$L = \begin{bmatrix} 1 & 0 \\ 3 & 1 \end{bmatrix}$$

Direct Methods

⊙ LU Factorization: Example 1 revisited

$$\begin{bmatrix} 1 & 1 \\ 3 & -4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 2 \end{bmatrix}$$

$$L = \begin{bmatrix} 1 & 0 \\ 3 & 1 \end{bmatrix} \quad U = \begin{bmatrix} 1 & 1 \\ 0 & -7 \end{bmatrix}$$

$$LU = \begin{bmatrix} 1 & 0 \\ 3 & 1 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 0 & -7 \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 3 & -4 \end{bmatrix} = A$$

Direct Methods

⊙ Gaussian Elimination: Example 2 revisited

$$x + 2y - z = 3$$

$$2x + y - 2z = 3$$

$$-3x + y + z = -6$$



$$\left[\begin{array}{ccc|c} 1 & 2 & -1 & 3 \\ 2 & 1 & -2 & 3 \\ -3 & 1 & 1 & -6 \end{array} \right]$$

$$\left[\begin{array}{ccc} 1 & 2 & -1 \\ 2 & 1 & -2 \\ -3 & 1 & 1 \end{array} \right]$$

→ subtract $2 \times$ row 1
from row 2 →

→ subtract $-3 \times$ row 1
from row 3 →

→ subtract $-\frac{7}{3} \times$ row 2
from row 3 →

$$\left[\begin{array}{ccc} 1 & 2 & -1 \\ 0 & -3 & 0 \\ -3 & 1 & 1 \end{array} \right]$$

$$\left[\begin{array}{ccc} 1 & 2 & -1 \\ 0 & -3 & 0 \\ 0 & 7 & -2 \end{array} \right]$$

$$\left[\begin{array}{ccc} 1 & 2 & -1 \\ 0 & -3 & 0 \\ 0 & 0 & -2 \end{array} \right] = U$$

Direct Methods

⊙ Gaussian Elimination: Example 2 revisited

$$\begin{bmatrix} 1 & 2 & -1 \\ 2 & 1 & -2 \\ -3 & 1 & 1 \end{bmatrix} \rightarrow \begin{bmatrix} 1 & 2 & -1 \\ 0 & -3 & 0 \\ 0 & 0 & -2 \end{bmatrix} = U$$

$$L = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -3 & -\frac{7}{3} & 1 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -3 & -\frac{7}{3} & 1 \end{bmatrix} \begin{bmatrix} 1 & 2 & -1 \\ 0 & -3 & 0 \\ 0 & 0 & -2 \end{bmatrix} = \begin{bmatrix} 1 & 2 & -1 \\ 2 & 1 & -2 \\ -3 & 1 & 1 \end{bmatrix}$$

Direct Methods

LU Factorization

An $m \times n$ matrix L is **lower triangular** if its entries satisfy $l_{ij} = 0$ for $i < j$.

An $m \times n$ matrix U is **upper triangular** if its entries satisfy $u_{ij} = 0$ for $i > j$.

Direct Methods

⊙ LU Factorization

Fact #1: Let $L_{ij}(-c)$ denote the lower triangular matrix whose only nonzero entries are 1's on the main diagonal and $-c$ in the (i, j) position. Then $A \rightarrow L_{ij}(-c)A$ represents the row operation “subtracting c times row j from row i .”

$$\begin{aligned} A = \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix} &\longrightarrow \begin{bmatrix} 1 & 0 & 0 \\ -c & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} & a_{22} & a_{23} \\ a_{31} & a_{32} & a_{33} \end{bmatrix} \\ &= \begin{bmatrix} a_{11} & a_{12} & a_{13} \\ a_{21} - ca_{11} & a_{22} - ca_{12} & a_{23} - ca_{13} \\ a_{31} & a_{32} & a_{33} \end{bmatrix} \end{aligned}$$

Direct Methods

⊙ LU Factorization

Fact #2: $L_{ij}(-c)^{-1} = L_{ij}(c)$

$$\begin{bmatrix} 1 & 0 & 0 \\ -c & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}^{-1} = \begin{bmatrix} 1 & 0 & 0 \\ c & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

Direct Methods

⊙ LU Factorization: Example 1 revisited

Using Facts #1 & 2:

The elimination step can be represented by

$$L_{21}(-3)A = \begin{bmatrix} 1 & 0 \\ -3 & 1 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 3 & -4 \end{bmatrix} = \begin{bmatrix} 1 & 1 \\ 0 & -7 \end{bmatrix}$$

To multiply both sides on the left by $L_{21}(-3)^{-1}$

$$A = \begin{bmatrix} 1 & 1 \\ 3 & -4 \end{bmatrix} = \begin{bmatrix} 1 & 0 \\ 3 & 1 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 0 & -7 \end{bmatrix}$$

the LU factorization of A

Direct Methods

⊙ LU Factorization

Fact #3: The following matrix product equation holds

$$\begin{bmatrix} 1 & & & \\ c_1 & 1 & & \\ & & 1 & \\ & & & 1 \end{bmatrix} \begin{bmatrix} 1 & & & \\ & 1 & & \\ c_2 & & 1 & \\ & & & 1 \end{bmatrix} \begin{bmatrix} 1 & & & \\ & 1 & & \\ & & 1 & \\ c_3 & & & 1 \end{bmatrix} \\ = \begin{bmatrix} 1 & & & \\ c_1 & 1 & & \\ c_2 & c_3 & 1 & \end{bmatrix}$$

Direct Methods

⦿ LU Factorization: Example 2 revisited

$$\begin{bmatrix} 1 & & \\ & 1 & \\ & \frac{7}{3} & 1 \end{bmatrix} \begin{bmatrix} 1 & & \\ & 1 & \\ & 3 & 1 \end{bmatrix} \begin{bmatrix} 1 & & \\ -2 & 1 & \\ & & 1 \end{bmatrix} \begin{bmatrix} 1 & 2 & -1 \\ 2 & 1 & -2 \\ -3 & 1 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 2 & -1 \\ 0 & -3 & 0 \\ 0 & 0 & -2 \end{bmatrix} = U$$

$$\begin{aligned} A &= \begin{bmatrix} 1 & & \\ 2 & 1 & \\ & & 1 \end{bmatrix} \begin{bmatrix} 1 & & \\ & 1 & \\ -3 & & 1 \end{bmatrix} \begin{bmatrix} 1 & & \\ & 1 & \\ & -\frac{7}{3} & 1 \end{bmatrix} \begin{bmatrix} 1 & 2 & -1 \\ 0 & -3 & 0 \\ 0 & 0 & -2 \end{bmatrix} \\ &= \begin{bmatrix} 1 & & \\ 2 & 1 & \\ -3 & -\frac{7}{3} & 1 \end{bmatrix} \begin{bmatrix} 1 & 2 & -1 \\ 0 & -3 & 0 \\ 0 & 0 & -2 \end{bmatrix} = LU \end{aligned}$$

Direct Methods

⊙ LU Factorization: Application

Once L and U are known, the problem $Ax = b$ can be written as $LUx = b$.

Define a new “auxiliary” vector $c = Ux$. Then back substitution is a two-step procedure:

- (a) Solve $Lc = b$ for c .
- (b) Solve $Ux = c$ for x .

Direct Methods

⊙ LU Factorization: Example 1 revisited

$$\begin{bmatrix} 1 & 1 \\ 3 & -4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 2 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 1 \\ 3 & -4 \end{bmatrix} = LU = \begin{bmatrix} 1 & 0 \\ 3 & 1 \end{bmatrix} \begin{bmatrix} 1 & 1 \\ 0 & -7 \end{bmatrix}$$

$$\begin{bmatrix} 1 & 0 \\ 3 & 1 \end{bmatrix} \begin{bmatrix} c_1 \\ c_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 2 \end{bmatrix} \quad \Rightarrow c_1 = 3, c_2 = -7$$

$$\begin{bmatrix} 1 & 1 \\ 0 & -7 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 3 \\ -7 \end{bmatrix} \Rightarrow x_2 = 1, x_1 = 2$$

Direct Methods

• LU Factorization: Complexity

Suppose we need to solve a number of different problems with the same A and different b .

$$Ax = b_1$$

- **Classical Gaussian elimination:**

$$Ax = b_2$$

Approximately $2kn^3/3$ operations

\vdots

- **The LU approach:**

$$Ax = b_k$$

Approximately $2n^3/3 + 2kn^2$ operations

Direct Methods

⊙ Sources of Error: Error Magnification

The infinity norm, or maximum norm, of the vector $x = (x_1, \dots, x_n)$ is

$$||x||_{\infty} = \max |x_i|, i = 1, \dots, n$$

Let x_a be an approximate solution of the linear system $Ax = b$. The residual is the vector

$$r = b - Ax_a$$

The **backward error**: $||b - Ax_a||_{\infty}$

The **forward error**: $||x - x_a||_{\infty}$

Direct Methods

⊙ Sources of Error: Example 3

Find the forward and backward errors for the approximate solution $[-1, 3.0001]$ of the system

$$x_1 + x_2 = 2$$

$$1.0001x_1 + x_2 = 2.0001$$

The backward error is the infinity norm of the vector

$$\begin{aligned} b - Ax_a &= \begin{bmatrix} 2 \\ 2.0001 \end{bmatrix} - \begin{bmatrix} 1 & 1 \\ 1.0001 & 1 \end{bmatrix} \begin{bmatrix} -1 \\ 3.0001 \end{bmatrix} \\ &= \begin{bmatrix} 2 \\ 2.0001 \end{bmatrix} - \begin{bmatrix} 2.0001 \\ 2 \end{bmatrix} = \begin{bmatrix} -0.0001 \\ 0.0001 \end{bmatrix} \end{aligned}$$

which is **0.0001**

Direct Methods

⊙ Sources of Error: Example 3

Find the forward and backward errors for the approximate solution $[-1, 3.0001]$ of the system

$$x_1 + x_2 = 2$$

$$1.0001x_1 + x_2 = 2.0001$$

The forward error is the infinity norm of the difference

$$x - x_a = \begin{bmatrix} 1 \\ 1 \end{bmatrix} - \begin{bmatrix} -1 \\ 3.0001 \end{bmatrix} = \begin{bmatrix} 2 \\ -2.0001 \end{bmatrix}$$

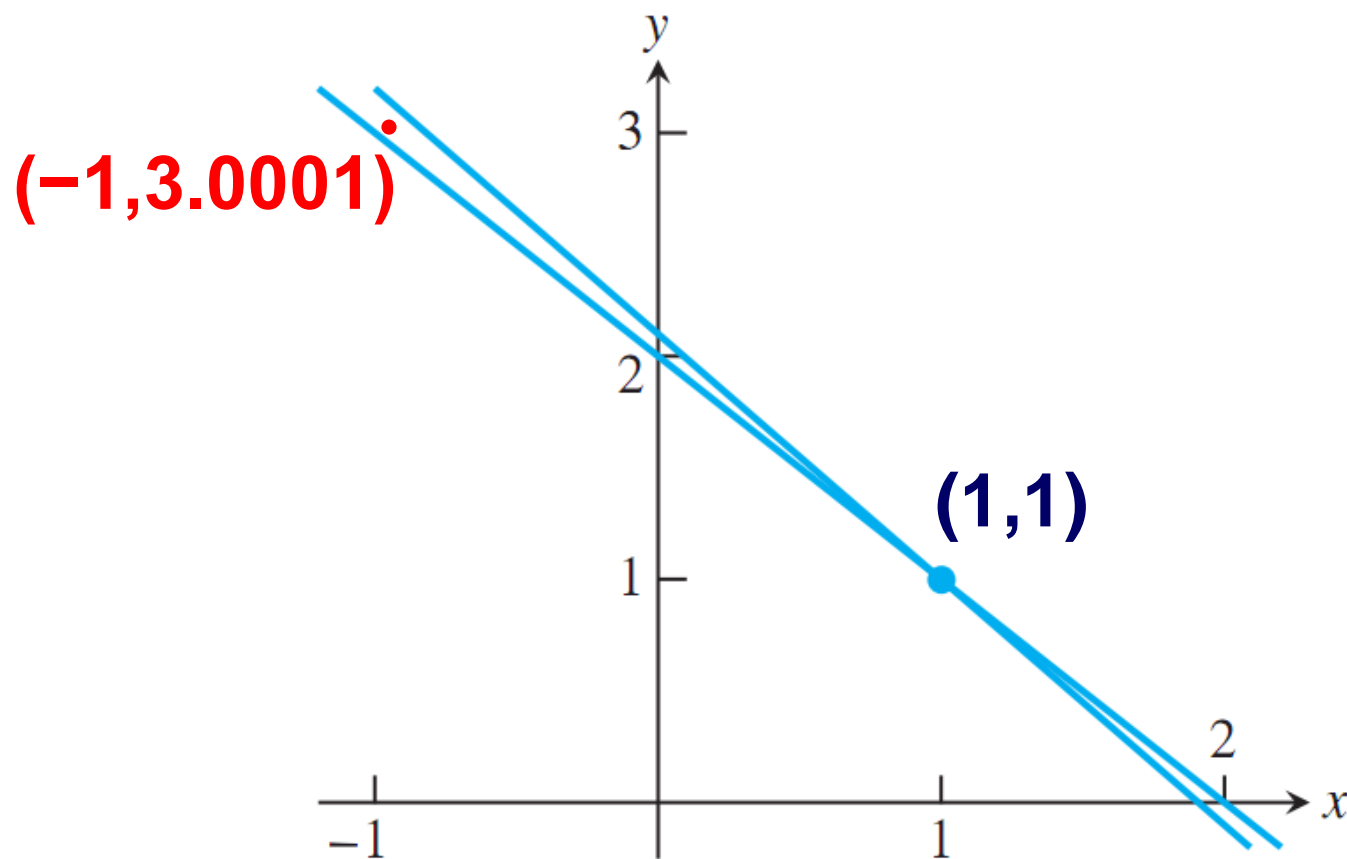
which is 2.0001

Direct Methods

Sources of Error: Example 3

$$x_1 + x_2 = 2$$

$$1.0001x_1 + x_2 = 2.0001$$



Direct Methods

⊙ Sources of Error: Error Magnification

Denote the residual by $r = b - Ax_a$.

The **relative backward error** of system $Ax = b$ is defined to be

$$\frac{||r||_{\infty}}{||b||_{\infty}}$$

and the **relative forward error** is

$$\frac{||x - x_a||_{\infty}}{||x||_{\infty}}$$

The **error magnification factor** for $Ax = b$ is

$$\frac{\text{relative forward error}}{\text{relative backward error}} = \frac{\frac{||x - x_a||_{\infty}}{||x||_{\infty}}}{\frac{||r||_{\infty}}{||b||_{\infty}}}$$

Direct Methods

⊙ Sources of Error: Example 3 revisited

$$x_1 + x_2 = 2$$

$$1.0001x_1 + x_2 = 2.0001$$

The relative forward error is

$$\frac{2.0001}{1} = 2.0001$$

The relative backward error is

$$\frac{0.0001}{2.0001} \approx 0.00005$$

The error magnification factor is

$$2.0001 / (0.0001 / 2.0001) = 40004.0001$$

Direct Methods

⊙ Sources of Error: Error Magnification

The **condition number of a square matrix A** , $\text{cond}(A)$, is the maximum possible error magnification factor for solving $Ax = b$, over all right-hand sides b .

The condition number of the $n \times n$ matrix A is

$$\text{cond}(A) = \|A\| \cdot \|A^{-1}\|$$

Direct Methods

⊙ Sources of Error: Matrix Norm

Analogous to the norm of a vector, the matrix norm of an $n \times n$ matrix A is defined as

$$\|A\|_{\infty} = \text{maximum absolute row sum}$$

that is, total the absolute values of each row, and assign the maximum of these n numbers to be the norm of A .

Direct Methods

⊙ Sources of Error: Example 3 revisited

$$x_1 + x_2 = 2$$

$$1.0001x_1 + x_2 = 2.0001$$

$$A = \begin{bmatrix} 1 & 1 \\ 1.0001 & 1 \end{bmatrix} \quad \|A\| = 2.0001$$

$$A^{-1} = \begin{bmatrix} -10000 & 10000 \\ 10001 & -10000 \end{bmatrix} \quad \|A^{-1}\| = 20001$$

$$\text{cond}(A) = (2.0001)(20001) = 40004.0001$$

This is exactly the error magnification

Direct Methods

⊙ Sources of Error: Example 4

$$10^{-20}x_1 + x_2 = 1$$

$$x_1 + 2x_2 = 4$$

Exact solution

$$\left[\begin{array}{cc|c} 10^{-20} & 1 & 1 \\ 1 & 2 & 4 \end{array} \right] \xrightarrow{\text{subtract } 10^{20} \times \text{row 1} \\ \text{from row 2}} \left[\begin{array}{cc|c} 10^{-20} & 1 & 1 \\ 0 & 2 - 10^{20} & 4 - 10^{20} \end{array} \right]$$

$$(2 - 10^{20})x_2 = 4 - 10^{20} \longrightarrow x_2 = \frac{4 - 10^{20}}{2 - 10^{20}}$$

$$10^{-20}x_1 + \frac{4 - 10^{20}}{2 - 10^{20}} = 1$$

$$x_1 = 10^{20} \left(1 - \frac{4 - 10^{20}}{2 - 10^{20}} \right) = \frac{-2 \times 10^{20}}{2 - 10^{20}}.$$

Direct Methods

⊙ Sources of Error: Example 4

The computer version of Gaussian elimination

$$\left[\begin{array}{cc|c} 10^{-20} & 1 & 1 \\ 1 & 2 & 4 \end{array} \right] \xrightarrow[\text{from row 2}]{\text{subtract } 10^{20} \times \text{row 1}} \left[\begin{array}{cc|c} 10^{-20} & 1 & 1 \\ 0 & 2 - 10^{20} & 4 - 10^{20} \end{array} \right]$$

$$-10^{20}x_2 = -10^{20} \longrightarrow x_2 = 1$$

The machine arithmetic version of the top equation becomes

$$10^{-20}x_1 + 1 = 1 \longrightarrow x_1 = 0$$

Direct Methods

⊙ Sources of Error: Example 4

Repeat the computer version of Gaussian elimination, after changing the order of the two equations

$$\left[\begin{array}{cc|c} 1 & 2 & 4 \\ 10^{-20} & 1 & 1 \end{array} \right] \rightarrow \begin{array}{l} \text{subtract } 10^{-20} \times \text{row 1} \\ \text{from row 2} \end{array}$$

$$\rightarrow \left[\begin{array}{cc|c} 1 & 2 & 4 \\ 0 & 1 - 2 \times 10^{-20} & 1 - 4 \times 10^{-20} \end{array} \right]$$

$$\rightarrow \begin{array}{l} x_1 + 2x_2 = 4 \\ x_2 = 1 \end{array}$$

yield the computed solution $x_1 = 2$ and $x_2 = 1$.

Direct Methods

⊙ The $PA=LU$ Factorization

The key to this improvement is an efficient protocol for **exchanging rows of the coefficient matrix**, called partial pivoting.

The partial pivoting protocol consists of comparing numbers before carrying out each elimination step.

Direct Methods

The PA= LU Factorization

At the start of Gaussian elimination, partial pivoting asks that we select the p th row, where

$$|a_{p1}| \geq |a_{i1}|$$

for all $1 \leq i \leq n$, and **exchange rows 1 and p** .

The multiplier used to eliminate a_{i1} will be

$$m_{i1} = \frac{a_{i1}}{a_{11}} \quad |m_{i1}| \leq 1$$

The same check is applied to every choice of pivot during the algorithm.

Direct Methods

⊙ LU Factorization: Example 1 revisited

$$\begin{bmatrix} 1 & 1 \\ 3 & -4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 2 \end{bmatrix}$$

Since $|a_{21}| > |a_{11}|$, we exchange rows 1 and 2

$$\left[\begin{array}{cc|c} 3 & -4 & 2 \\ 1 & 1 & 3 \end{array} \right] \xrightarrow{\text{subtract } \frac{1}{3} \times \text{row 1} \text{ from row 2}} \left[\begin{array}{cc|c} 3 & -4 & 2 \\ 0 & \frac{7}{3} & \frac{7}{3} \end{array} \right]$$

When we solved this system the first time, the multiplier was 3

Direct Methods

⊙ The PA= LU Factorization

A permutation matrix is an $n \times n$ matrix consisting of all zeros, except for a single 1 in every row and column.

$$\begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}, \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$
$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}, \begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix}$$
$$\begin{bmatrix} 0 & 0 & 1 \\ 0 & 1 & 0 \\ 1 & 0 & 0 \end{bmatrix}, \begin{bmatrix} 0 & 0 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \end{bmatrix}, \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$$

Direct Methods

⊙ The $PA=LU$ Factorization

Let P be the $n \times n$ permutation matrix formed by a particular set of row exchanges applied to the identity matrix. Then, for any $n \times n$ matrix A , PA is the matrix obtained by applying exactly the same set of row exchanges to A .

$$\begin{bmatrix} 1 & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{bmatrix} \begin{bmatrix} a & b & c \\ d & e & f \\ g & h & i \end{bmatrix} = \begin{bmatrix} a & b & c \\ g & h & i \\ d & e & f \end{bmatrix}$$

Direct Methods

The $PA = LU$ Factorization

The $PA = LU$ factorization is simply the LU factorization of a row-exchanged version of A .

Under partial pivoting, the rows that need exchanging are not known at the outset, so we must be careful about fitting the row exchange information into the factorization.

Direct Methods

⊙ The PA=LU Factorization: Example 5

$$A = \begin{bmatrix} 2 & 1 & 5 \\ 4 & 4 & -4 \\ 1 & 3 & 1 \end{bmatrix}$$

$$P = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 2 & 1 & 5 \\ 4 & 4 & -4 \\ 1 & 3 & 1 \end{bmatrix} \longrightarrow \text{exchange rows 1 and 2} \longrightarrow \begin{bmatrix} 4 & 4 & -4 \\ 2 & 1 & 5 \\ 1 & 3 & 1 \end{bmatrix}$$

Direct Methods

⊙ The PA=LU Factorization: Example 5

$$\begin{aligned}
 & \begin{bmatrix} 2 & 1 & 5 \\ 4 & 4 & -4 \\ 1 & 3 & 1 \end{bmatrix} \xrightarrow{\text{exchange rows 1 and 2}} \begin{bmatrix} 4 & 4 & -4 \\ 2 & 1 & 5 \\ 1 & 3 & 1 \end{bmatrix} \\
 & \quad P = \begin{bmatrix} 0 & 1 & 0 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} \\
 & \quad \xrightarrow{\text{subtract } \frac{1}{2} \times \text{row 1 from row 2}} \begin{bmatrix} 4 & 4 & -4 \\ \textcircled{\frac{1}{2}} & -1 & 7 \\ 1 & 3 & 1 \end{bmatrix} \\
 & \quad \xrightarrow{\text{subtract } \frac{1}{4} \times \text{row 1 from row 3}} \begin{bmatrix} 4 & 4 & -4 \\ \textcircled{\frac{1}{2}} & -1 & 7 \\ \textcircled{\frac{1}{4}} & 2 & 2 \end{bmatrix}
 \end{aligned}$$

Direct Methods

⊙ The PA=LU Factorization: Example 5

subtract $\frac{1}{4} \times \text{row 1}$
 → from row 3 →

$$\begin{bmatrix} 4 & 4 & -4 \\ \left(\frac{1}{2}\right) & -1 & 7 \\ \left(\frac{1}{4}\right) & 2 & 2 \end{bmatrix}$$

$$P = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix}$$

→ exchange rows 2 and 3 →

$$\begin{bmatrix} 4 & 4 & -4 \\ \left(\frac{1}{4}\right) & 2 & 2 \\ \left(\frac{1}{2}\right) & -1 & 7 \end{bmatrix}$$

subtract $-\frac{1}{2} \times \text{row 2}$
 → from row 3 →

$$\begin{bmatrix} 4 & 4 & -4 \\ \left(\frac{1}{4}\right) & 2 & 2 \\ \left(\frac{1}{2}\right) & \left(-\frac{1}{2}\right) & 8 \end{bmatrix}$$

Direct Methods

⊙ The PA=LU Factorization: Example 5

$$A = \begin{bmatrix} 2 & 1 & 5 \\ 4 & 4 & -4 \\ 1 & 3 & 1 \end{bmatrix}$$

$$\begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} 2 & 1 & 5 \\ 4 & 4 & -4 \\ 1 & 3 & 1 \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 \\ \frac{1}{4} & 1 & 0 \\ \frac{1}{2} & -\frac{1}{2} & 1 \end{bmatrix} \begin{bmatrix} 4 & 4 & -4 \\ 0 & 2 & 2 \\ 0 & 0 & 8 \end{bmatrix}$$

$P \qquad A \qquad L \qquad U$

Direct Methods

⊙ The PA= LU Factorization: Application

Once P , L and U are known, the problem $Ax = b$ can be written as $LUx = Pb$.

Define a new “auxiliary” vector $c = Ux$. Then back substitution is a two-step procedure:

1. $Lc = Pb$ for c .
2. $Ux = c$ for x .

Direct Methods

⦿ The PA= LU Factorization: Example 5

$$A = \begin{bmatrix} 2 & 1 & 5 \\ 4 & 4 & -4 \\ 1 & 3 & 1 \end{bmatrix}$$

```
>> A = [2,1,5;4,4,-4;1,3,1];
```

```
>> [L,U,P] = lu(A)
```

Systems of Linear Equations

□ Direct Methods

- Gaussian Elimination
- The LU Factorization
- Sources of Error
- The $PA=LU$ Factorization

□ Iterative Methods

- Jacobi Method
- Gauss–Seidel Method

Iterative Methods

Jacobi Method: Basic Idea

The Jacobi Method is a form of fixed-point iteration for a system of equations

**Solve the i th equation for the i th unknown.
Then, iterate as in Fixed-Point Iteration,
starting with an initial guess.**

Iterative Methods

Jacobi Method: Example 6

$$\begin{cases} 3u + v = 5 \\ u + 2v = 5 \end{cases}$$



Begin by solving the first equation for u and the second equation for v .

$$u = \frac{5 - v}{3}$$

$$v = \frac{5 - u}{2}$$

$$\begin{bmatrix} u_0 \\ v_0 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

Iterative Methods

❶ Jacobi Method: Example 6

$$\begin{bmatrix} u_0 \\ v_0 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} u_1 \\ v_1 \end{bmatrix} = \begin{bmatrix} \frac{5-v_0}{3} \\ \frac{5-u_0}{2} \end{bmatrix} = \begin{bmatrix} \frac{5-0}{3} \\ \frac{5-0}{2} \end{bmatrix} = \begin{bmatrix} \frac{5}{3} \\ \frac{5}{2} \end{bmatrix}$$


$$\begin{bmatrix} u_2 \\ v_2 \end{bmatrix} = \begin{bmatrix} \frac{5-v_1}{3} \\ \frac{5-u_1}{2} \end{bmatrix} = \begin{bmatrix} \frac{5-5/2}{3} \\ \frac{5-5/3}{2} \end{bmatrix} = \begin{bmatrix} \frac{5}{6} \\ \frac{5}{3} \end{bmatrix}$$

$$\begin{bmatrix} u_3 \\ v_3 \end{bmatrix} = \begin{bmatrix} \frac{5-5/3}{3} \\ \frac{5-5/6}{2} \end{bmatrix} = \begin{bmatrix} \frac{10}{9} \\ \frac{25}{12} \end{bmatrix}.$$

Iterative Methods

⊙ Jacobi Method: Example 7

$$\begin{cases} u + 2v = 5 \\ 3u + v = 5 \end{cases} \quad \leftarrow \quad \begin{cases} 3u + v = 5 \\ u + 2v = 5 \end{cases}$$

 **Begin by solving the first equation for u and the second equation for v .**

$$\begin{aligned} u &= 5 - 2v \\ v &= 5 - 3u \end{aligned} \quad \begin{bmatrix} u_0 \\ v_0 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

Iterative Methods

Jacobi Method: Example 7

$$\begin{bmatrix} u_0 \\ v_0 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} u_1 \\ v_1 \end{bmatrix} = \begin{bmatrix} 5 - 2v_0 \\ 5 - 3u_0 \end{bmatrix} = \begin{bmatrix} 5 \\ 5 \end{bmatrix}$$

$$\begin{bmatrix} u_2 \\ v_2 \end{bmatrix} = \begin{bmatrix} 5 - 2v_1 \\ 5 - 3u_1 \end{bmatrix} = \begin{bmatrix} -5 \\ -10 \end{bmatrix}$$

$$\begin{bmatrix} u_3 \\ v_3 \end{bmatrix} = \begin{bmatrix} 5 - 2(-10) \\ 5 - 3(-5) \end{bmatrix} = \begin{bmatrix} 25 \\ 20 \end{bmatrix}$$

the iteration diverges

Iterative Methods

Jacobi Method

The $n \times n$ matrix $A = (a_{ij})$ is **strictly diagonally dominant** if, for each $1 \leq i \leq n$,

$$|a_{ii}| > \sum_{j \neq i} |a_{ij}|$$

Each main diagonal entry dominates its row in the sense that it is greater in magnitude than the sum of magnitudes of the remainder of the entries in its row.

Iterative Methods

Jacobi Method: Theory

THEOREM 2.10

If the $n \times n$ matrix A is strictly diagonally dominant, then

- (1) A is a nonsingular matrix
- (2) for every vector b and every starting guess, the Jacobi Method applied to $Ax = b$ converges to the (unique) solution.

Cf. P.107 of Ref. [1]

Iterative Methods

⊙ **Jacobi Method: Algorithm**

Let D denote the main diagonal of A , L denote the lower triangle of A (entries below the main diagonal), and U denote the upper triangle (entries above the main diagonal). Then $A = L + D + U$

$$Ax = b$$

$$(D + L + U)x = b$$

$$Dx = b - (L + U)x$$

$$x = D^{-1}(b - (L + U)x)$$

$$x_0 = \text{initial vector}$$

$$x_{k+1} = D^{-1}(b - (L + U)x_k) \quad \mathbf{for} \quad k = 0, 1, 2, \dots$$

Iterative Methods

⊙ Jacobi Method: Example 6 revisited

$$\begin{cases} 3u + v = 5 \\ u + 2v = 5 \end{cases} \quad \begin{bmatrix} 3 & 1 \\ 1 & 2 \end{bmatrix} \begin{bmatrix} u \\ v \end{bmatrix} = \begin{bmatrix} 5 \\ 5 \end{bmatrix}$$

$$\begin{aligned} \begin{bmatrix} u_{k+1} \\ v_{k+1} \end{bmatrix} &= D^{-1}(b - (L + U)x_k) \\ &= \begin{bmatrix} 1/3 & 0 \\ 0 & 1/2 \end{bmatrix} \left(\begin{bmatrix} 5 \\ 5 \end{bmatrix} - \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \begin{bmatrix} u_k \\ v_k \end{bmatrix} \right) \\ &= \begin{bmatrix} (5 - v_k)/3 \\ (5 - u_k)/2 \end{bmatrix} \end{aligned}$$

Iterative Methods

Gauss–Seidel Method: Basic Idea

The only difference between Gauss–Seidel and Jacobi is that in the former, **the most recently updated values of the unknowns are used at each step**, even if the updating occurs in the current step.

Iterative Methods

⊙ Gauss–Seidel Method: Example 6 revisited

$$\begin{cases} 3u + v = 5 \\ u + 2v = 5 \end{cases} \quad \text{Jacobi Method:} \quad \begin{bmatrix} u_{k+1} \\ v_{k+1} \end{bmatrix} = \begin{bmatrix} (5 - v_k)/3 \\ (5 - u_k)/2 \end{bmatrix}$$

$$\begin{bmatrix} u_0 \\ v_0 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$$

$$\begin{bmatrix} u_1 \\ v_1 \end{bmatrix} = \begin{bmatrix} \frac{5-v_0}{3} \\ \frac{5-u_1}{2} \end{bmatrix} = \begin{bmatrix} \frac{5-0}{3} \\ \frac{5-5/3}{2} \end{bmatrix} = \begin{bmatrix} \frac{5}{3} \\ \frac{5}{3} \end{bmatrix}$$

$$\begin{bmatrix} u_2 \\ v_2 \end{bmatrix} = \begin{bmatrix} \frac{5-v_1}{3} \\ \frac{5-u_2}{2} \end{bmatrix} = \begin{bmatrix} \frac{5-5/3}{3} \\ \frac{5-10/9}{2} \end{bmatrix} = \begin{bmatrix} \frac{10}{9} \\ \frac{35}{18} \end{bmatrix}$$

$$\begin{bmatrix} u_3 \\ v_3 \end{bmatrix} = \begin{bmatrix} \frac{5-v_2}{3} \\ \frac{5-u_3}{2} \end{bmatrix} = \begin{bmatrix} \frac{5-35/18}{3} \\ \frac{5-55/54}{2} \end{bmatrix} = \begin{bmatrix} \frac{55}{54} \\ \frac{215}{108} \end{bmatrix}$$

Iterative Methods

⊙ Gauss–Seidel Method: Algorithm

Let D denote the main diagonal of A , L denote the lower triangle of A (entries below the main diagonal), and U denote the upper triangle (entries above the main diagonal). Then $A = L + D + U$

$$Ax = b$$

$$(D + L + U)x = b$$

$$(L + D)x_{k+1} = -Ux_k + b$$

$$x_0 = \text{initial vector}$$

$$x_{k+1} = D^{-1}(b - Ux_k - Lx_{k+1}) \quad \text{for } k = 0, 1, 2, \dots$$

Iterative Methods

⊙ Gauss–Seidel Method: Example 8

$$\begin{bmatrix} 3 & 1 & -1 \\ 2 & 4 & 1 \\ -1 & 2 & 5 \end{bmatrix} \begin{bmatrix} u \\ v \\ w \end{bmatrix} = \begin{bmatrix} 4 \\ 1 \\ 1 \end{bmatrix} \quad \text{Solution: } [2, -1, 1]$$

$$u_{k+1} = \frac{4 - v_k + w_k}{3}$$

$$v_{k+1} = \frac{1 - 2u_{k+1} - w_k}{4}$$

$$w_{k+1} = \frac{1 + u_{k+1} - 2v_{k+1}}{5}$$

Starting with

$$[u_0, v_0, w_0] = [0, 0, 0]$$

$$\begin{bmatrix} u_1 \\ v_1 \\ w_1 \end{bmatrix} = \begin{bmatrix} \frac{4-0-0}{3} = \frac{4}{3} \\ \frac{1-8/3-0}{4} = -\frac{5}{12} \\ \frac{1+4/3+5/6}{5} = \frac{19}{30} \end{bmatrix}$$

$$\begin{bmatrix} u_2 \\ v_2 \\ w_2 \end{bmatrix} \approx \begin{bmatrix} 1.6833 \\ -0.7500 \\ 0.8367 \end{bmatrix}$$

Summary

☐ Direct Methods

- ✓ Gaussian Elimination
- ✓ The LU Factorization
- ✓ Sources of Error
- ✓ The $PA=LU$ Factorization

☐ Iterative Methods

- ✓ Jacobi Method
- ✓ Gauss–Seidel Method

Thank You !