



Chapter 6

Recursion as a Problem-Solving Technique

Backtracking

- Backtracking
 - A strategy for guessing at a solution and backing up when an impasse is reached
- Recursion and backtracking can be combined to solve problems

The Eight Queens Problem

- Problem
 - Place eight queens on the chessboard so that no queen can attack any other queen
- Strategy: guess at a solution
 - There are 4,426,165,368 ways to arrange 8 queens on a chessboard of 64 squares

The Eight Queens Problem

- An observation that eliminates many arrangements from consideration
 - No queen can reside in a row or a column that contains another queen
 - Now: only 40,320 arrangements of queens to be checked for attacks along diagonals

The Eight Queens Problem

- Providing organization for the guessing strategy
 - Place queens one column at a time
 - If you reach an impasse, backtrack to the previous column

The Eight Queens Problem

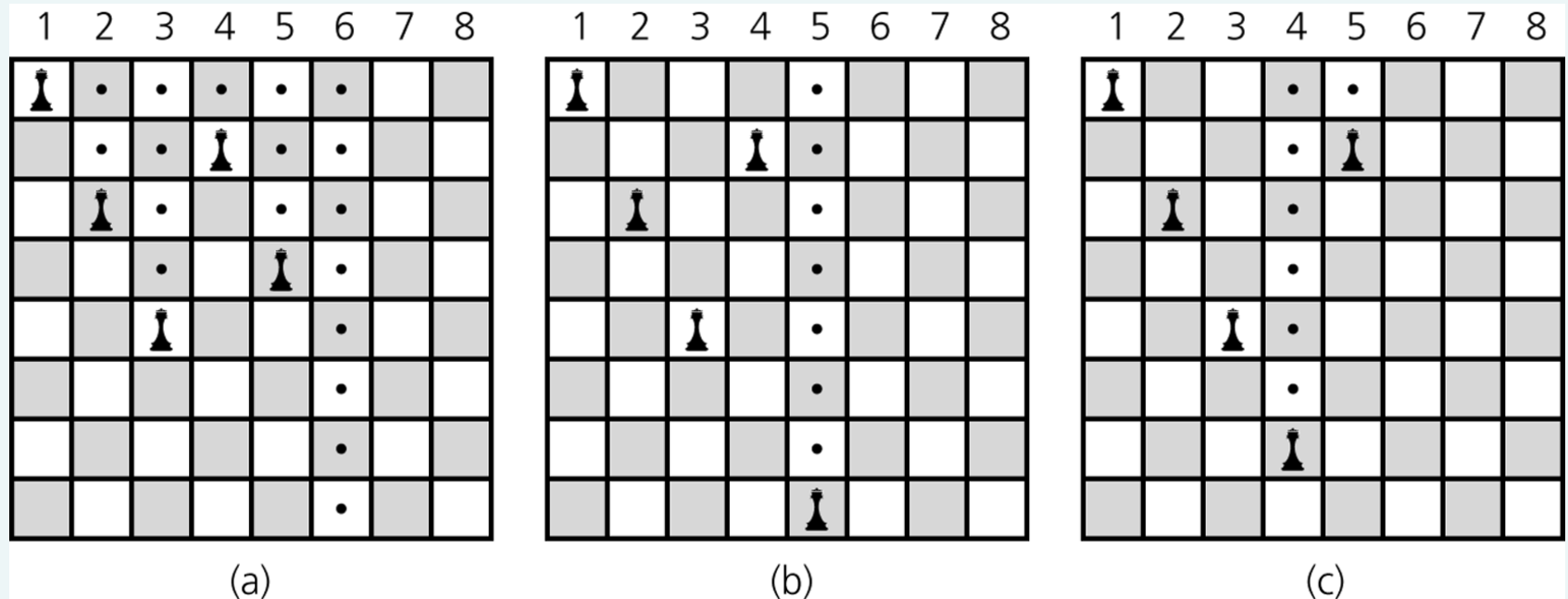


Figure 6-1

- a) Five queens that cannot attack each other, but that can attack all of column 6; b) backtracking to column 5 to try another square for the queen; c) backtracking to column 4 to try another square for the queen and then considering column 5 again

The Eight Queens Problem

- A recursive algorithm that places a queen in a column
 - Base case
 - If there are no more columns to consider
 - You are finished
 - Recursive step
 - If you successfully place a queen in the current column
 - Consider the next column
 - If you cannot place a queen in the current column
 - You need to backtrack

The Eight Queens Problem

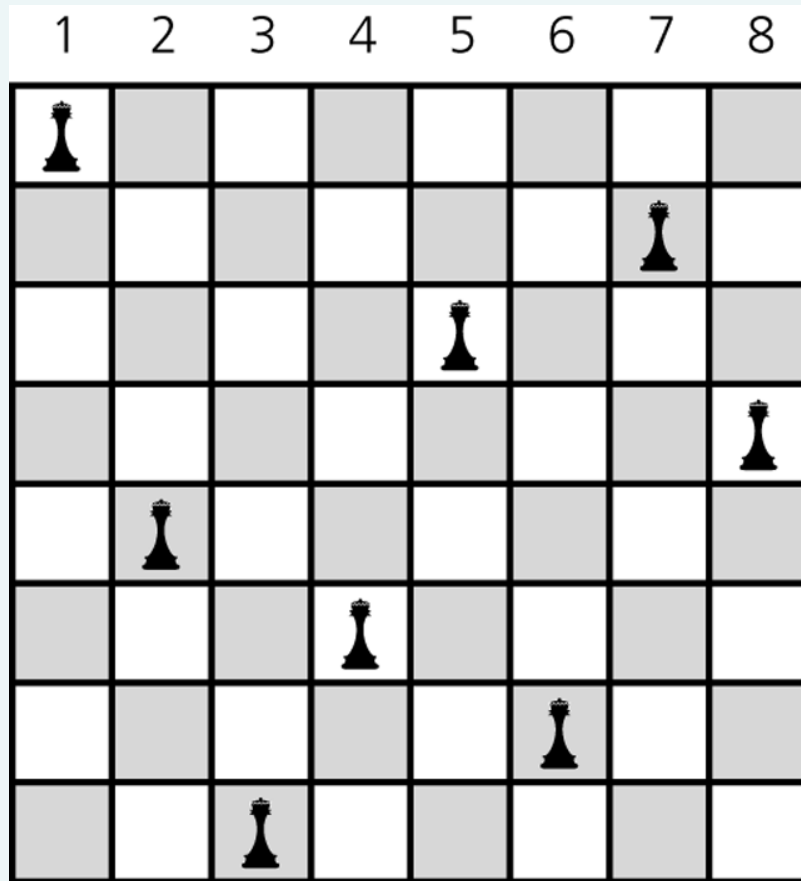


Figure 6-2

A solution to the Eight Queens problem

Defining Languages

- A language
 - A set of strings of symbols
 - Examples: English, Java
 - If a Java program is one long string of characters, the language JavaPrograms is defined as
$$\text{JavaPrograms} = \{\text{strings } w : w \text{ is a syntactically correct Java program}\}$$

Defining Languages

- A language does not have to be a programming or a communication language
 - Example
 - The set of algebraic expressions
$$\text{AlgebraicExpressions} = \{w : w \text{ is an algebraic expression}\}$$

Defining Languages

- Grammar
 - States the rules for forming the strings in a language
- Benefit of recursive grammars
 - Ease of writing a recognition algorithm for the language
 - A recognition algorithm determines whether a given string is in the language

The Basics of Grammars

- Symbols used in grammars
 - $x \mid y$ means x or y
 - xy means x followed by y
(In $x \bullet y$, the symbol \bullet means concatenate, or append)
 - $\langle \text{word} \rangle$ means any instance of word that the definition defines

The Basics of Grammars

- Java identifiers
 - A Java identifier begins with a letter and is followed by zero or more letters and digits

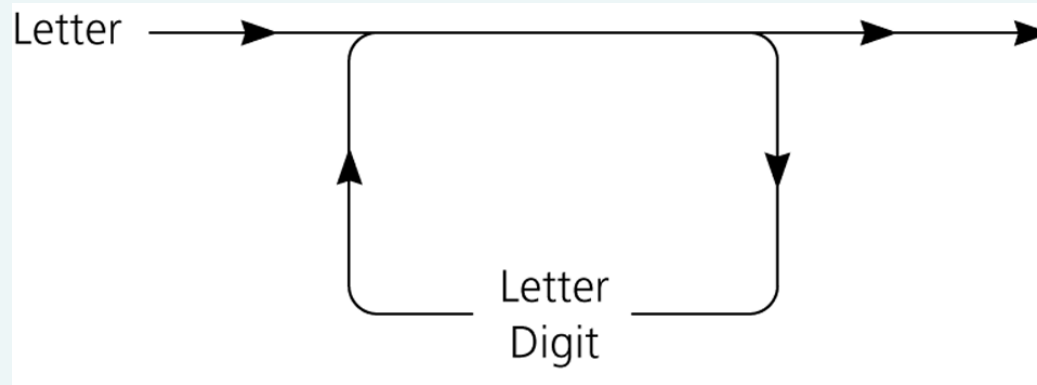


Figure 6-3

A syntax diagram for Java identifiers

The Basics of Grammars

- Java identifiers

- Language

$\text{JavaIds} = \{w : w \text{ is a legal Java identifier}\}$

- Grammar

$\langle \text{identifier} \rangle = \langle \text{letter} \rangle \mid \langle \text{identifier} \rangle \langle \text{letter} \rangle \mid \langle \text{identifier} \rangle \langle \text{digit} \rangle$

$\langle \text{letter} \rangle = a \mid b \mid \dots \mid z \mid A \mid B \mid \dots \mid Z \mid _ \mid \$$

$\langle \text{digit} \rangle = 0 \mid 1 \mid \dots \mid 9$

The Basics of Grammars

- Recognition algorithm

```
isId(w)
  if (w is of length 1) {
    if (w is a letter) {
      return true
    }
    else {
      return false
    }
  }
  else if (the last character of w is a letter or a digit) {
    return isId(w minus its last character)
  }
  else {
    return false
  }
```

Two Simple Languages: Palindromes

- A string that reads the same from left to right as it does from right to left
- Examples: radar, deed
- Language

Palindromes = $\{w : w \text{ reads the same left to right as right to left}\}$

Palindromes

- Grammar

$$\begin{aligned} \langle \text{pal} \rangle = \text{empty string} \mid \langle \text{ch} \rangle \mid a \langle \text{pal} \rangle a \mid b \langle \text{pal} \rangle b \mid \dots \\ \mid Z \langle \text{pal} \rangle Z \end{aligned}$$
$$\langle \text{ch} \rangle = a \mid b \mid \dots \mid z \mid A \mid B \mid \dots \mid Z$$

Palindromes

- Recognition algorithm

`isPal(w)`

```
    if (w is the empty string or w is of length 1) {  
        return true  
    }  
    else if (w's first and last characters are the  
             same letter ) {  
        return isPal(w minus its first and last  
                     characters)  
    }  
    else {  
        return false  
    }
```

Strings of the form A^nB^n

- A^nB^n
 - The string that consists of n consecutive A 's followed by n consecutive B 's
- Language
 - $L = \{w : w \text{ is of the form } A^nB^n \text{ for some } n \geq 0\}$
- Grammar
 - $\langle \text{legal-word} \rangle = \text{empty string} \mid A \langle \text{legal-word} \rangle B$

Strings of the form A^nB^n

- Recognition algorithm

```
isAnBn(w)
    if (the length of w is zero) {
        return true
    }
    else if (w begins with the character A and ends
              with the character B) {
        return isAnBn(w minus its first and last
                       characters)
    }
    else {
        return false
    }
```

Algebraic Expressions

- Three languages for algebraic expressions
 - Infix expressions
 - An operator appears between its operands
 - Example: $a + b$
 - Prefix expressions
 - An operator appears before its operands
 - Example: $+ a b$
 - Postfix expressions
 - An operator appears after its operands
 - Example: $a b +$

Algebraic Expressions

- To convert a fully parenthesized infix expression to a prefix form
 - Move each operator to the position marked by its corresponding open parenthesis
 - Remove the parentheses
 - Example
 - Infix expression: $((a + b) * c$
 - Prefix expression: $* + a b c$

Algebraic Expressions

- To convert a fully parenthesized infix expression to a postfix form
 - Move each operator to the position marked by its corresponding closing parenthesis
 - Remove the parentheses
 - Example
 - Infix form: $((a + b) * c)$
 - Postfix form: $a b + c *$

Algebraic Expressions

- Prefix and postfix expressions
 - Never need
 - Precedence rules
 - Association rules
 - Parentheses
 - Have
 - Simple grammar expressions
 - Straightforward recognition and evaluation algorithms

Prefix Expressions

- Grammar

$\langle \text{prefix} \rangle = \langle \text{identifier} \rangle \mid \langle \text{operator} \rangle \langle \text{prefix} \rangle \langle \text{prefix} \rangle$

$\langle \text{operator} \rangle = + \mid - \mid * \mid /$

$\langle \text{identifier} \rangle = a \mid b \mid \dots \mid z$

- A recognition algorithm

`isPre()`

`size = length of expression strExp`

`lastChar = endPre(0, size - 1)`

`if (lastChar >= 0 and lastChar == size-1 {`

`return true`

`}`

`else {`

`return false`

`}`

Prefix Expressions

- An algorithm that evaluates a prefix expression

```
evaluatePrefix(strExp)
```

```
    ch = first character of expression strExp
```

```
    Delete first character from strExp
```

```
    if (ch is an identifier) {
```

```
        return value of the identifier
```

```
    }
```

```
    else if (ch is an operator named op) {
```

```
        operand1 = evaluatePrefix(strExp)
```

```
        operand2 = evaluatePrefix(strExp)
```

```
        return operand1 op operand2
```

```
    }
```

Postfix Expressions

- Grammar

$\langle \text{postfix} \rangle = \langle \text{identifier} \rangle \mid \langle \text{postfix} \rangle \langle \text{postfix} \rangle \langle \text{operator} \rangle$

$\langle \text{operator} \rangle = + \mid - \mid * \mid /$

$\langle \text{identifier} \rangle = a \mid b \mid \dots \mid z$

- At high-level, an algorithm that converts a prefix expression to postfix form

```
if (exp is a single letter) {  
    return exp  
}  
else {  
    return postfix(prefix1) + postfix(prefix2) +  
        operator  
}
```

Postfix Expressions

- A recursive algorithm that converts a prefix expression to postfix form

```
convert(pre)
    ch = first character of pre
    Delete first character of pre
    if (ch is a lowercase letter) {
        return ch as a string
    }
    else {
        postfix1 = convert(pre)
        postfix2 = convert(pre)
        return postfix1 + postfix2 + ch
    }
```

Fully Parenthesized Expressions

- To avoid ambiguity, infix notation normally requires
 - Precedence rules
 - Rules for association
 - Parentheses
- Fully parenthesized expressions do not require
 - Precedence rules
 - Rules for association

Fully Parenthesized Expressions

- Fully parenthesized expressions
 - A simple grammar
$$\langle \text{infix} \rangle = \langle \text{identifier} \rangle \mid (\langle \text{infix} \rangle \langle \text{operator} \rangle \langle \text{infix} \rangle)$$
$$\langle \text{operator} \rangle = + \mid - \mid * \mid /$$
$$\langle \text{identifier} \rangle = a \mid b \mid \dots \mid z$$
 - Inconvenient for programmers

The Relationship Between Recursion and Mathematical Induction

- A strong relationship exists between recursion and mathematical induction
- Induction can be used to
 - Prove properties about recursive algorithms
 - Prove that a recursive algorithm performs a certain amount of work

The Correctness of the Recursive Factorial Method

- Pseudocode for a recursive method that computes the factorial of a nonnegative integer n

```
fact(n)
  if (n is 0) {
    return 1
  }
  else {
    return n * fact(n - 1)
  }
```


The Correctness of the Recursive Factorial Method

- Induction on n can prove that the method `fact` returns the values

$$\text{fact}(0) = 0! = 1$$

$$\text{fact}(n) = n! = n * (n - 1) * (n - 2) * \dots * 1 \quad \text{if } n > 0$$

The Cost of Towers of Hanoi

- Solution to the Towers of Hanoi problem

```
solveTowers(count, source, destination, spare)
  if (count is 1) {
    Move a disk directly from source to destination
  }
  else {
    solveTowers(count-1, source, spare, destination)
    solveTowers(1, source, destination, spare)
    solveTowers(count-1, spare, destination, source)
  }
```

The Cost of Towers of Hanoi

- Question
 - If you begin with N disks, how many moves does `solveTowers` make to solve the problem?
- Let
 - `moves (N)` be the number of moves made starting with N disks
- When $N = 1$
 - `moves (1) = 1`

The Cost of Towers of Hanoi

- When $N > 1$

$$\text{moves}(N) = \text{moves}(N - 1) + \text{moves}(1) + \text{moves}(N - 1)$$

- Recurrence relation for the number of moves that `solveTowers` requires for N disks

$$\text{moves}(1) = 1$$

$$\text{moves}(N) = 2 * \text{moves}(N - 1) + 1 \quad \text{if } N > 1$$

The Cost of Towers of Hanoi

- A closed-form formula for the number of moves that `solveTowers` requires for N disks

$$\text{moves}(N) = 2^N - 1, \text{ for all } N \geq 1$$

- Induction on N can provide the proof that $\text{moves}(N) = 2^N - 1$

Summary

- Backtracking is a solution strategy that involves both recursion and a sequence of guesses that ultimately lead to a solution
- A grammar is a device for defining a language
 - A language is a set of strings of symbols
 - A recognition algorithm for a language can often be based directly on the grammar of the language
 - Grammars are frequently recursive

Summary

- Different languages of algebraic expressions have their relative advantages and disadvantages
 - Prefix expressions
 - Postfix expressions
 - Infix expressions
- A close relationship exists between mathematical induction and recursion
 - Induction can be used to prove properties about a recursive algorithm