

MAT3007: Optimization - Assignment 2

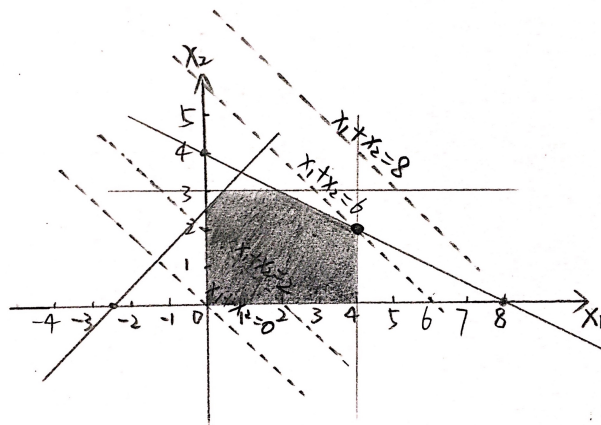
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Problem 1

1. False. If the objective function is not related to solution, then any $x \geq 0$ are optimal solutions. For example, minimize 0, subject to $x \geq 0$, the optimal solution set is $[0, \infty)$, which is unbounded.
2. False. The objective function of LP in its standard form is minimize $c^T x$, if $c = 0$, which means c is a zero vector, then there can be $[0, n]$ $m \leq n$ variables that are positive.
3. True. If there are two optimal solutions, say x and x' , then any convex combination of them are optimal solutions. Which means that if there is more than one optimal solution, then there are uncountably many optimal solutions.

Problem 2



Let x_1 be the horizontal axis, x_2 be the vertical axis.

The optimal solution is $x_1 = 4, x_2 = 2$, which gives an optimal result: 6.

Constrains $x_1 \leq 4$ and $x_1 + 2x_2 \leq 8$ are active at optimal solution.

The vertices of the feasible region are $(0, 0), (0, 2.5), (0.5, 3), (2, 3), (4, 2), (4, 0)$.

Problem 3

1. The optimization problem can be written as the following standard form:

$$\begin{aligned} \text{minimize} \quad & -x_1 - 4x_2 - x_3 \\ \text{s.t.} \quad & 2x_1 + 2x_2 + x_3 + s_1 = 4 \\ & x_1 - x_3 - s_2 = 1 \\ & x_1, x_2, x_3, s_1, s_2 \geq 0 \end{aligned}$$

2. The constrain can also be written as $Ax = b$, where $A = \begin{bmatrix} 2 & 2 & 1 & 1 & 0 \\ 1 & 0 & -1 & 0 & -1 \end{bmatrix}$

$$, x = \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ s_1 \\ s_2 \end{bmatrix}, b = \begin{bmatrix} 4 \\ 1 \end{bmatrix}. \text{ Because of the theorem that if there is an optimal}$$

solution, there is an optimal solution that is a basic feasible solution, I only need to prove that there exist a basic feasible solution with no more than 2 positive variables.

Notice that there are two constrains. So we can choose any two independent columns of A to form A_B . Then $x_B = A_B^{-1}b$. The entries in x_B are the only non-zero entries in the solution x, where the other three entries in x are all zero. Therefore, x can have no more than 2 positive variables. So there must exist an optimal solution with no more than 2 positive variables.

$$\begin{aligned} \text{3. Basic solutions are } & \begin{bmatrix} 1 \\ 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} \frac{5}{3} \\ 0 \\ \frac{2}{3} \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 0 \\ 0 \\ 2 \\ 0 \end{bmatrix}, \begin{bmatrix} 2 \\ 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}, \begin{bmatrix} 0 \\ 2.5 \\ -1 \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 2 \\ 0 \\ 0 \\ -1 \end{bmatrix}, \begin{bmatrix} 0 \\ 0 \\ -1 \\ 5 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 0 \\ 4 \\ 0 \\ -5 \end{bmatrix}, \\ & \begin{bmatrix} 0 \\ 0 \\ 0 \\ 4 \\ -1 \end{bmatrix}. \text{ Basic feasible solutions are } \begin{bmatrix} 1 \\ 1 \\ 0 \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} \frac{5}{3} \\ 0 \\ \frac{2}{3} \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 0 \\ 0 \\ 2 \\ 0 \end{bmatrix}, \begin{bmatrix} 2 \\ 0 \\ 0 \\ 0 \\ 1 \end{bmatrix}. \end{aligned}$$

4. The optimal solution is $x_1 = 1, x_2 = 1, x_3 = 0$, which gives an optimal value -5.

Problem 4

1. I use \mathbf{x} to denote the number of shares of each security to purchase. A_i is the i th column vector of A . The optimization problem can be formed as follows:

$$\begin{aligned} \text{maximize}_{\mathbf{x}} \quad & \min_i -\pi^T \mathbf{x} + A_i^T \mathbf{x}, \quad i = 1, 2, 3, 4, 5 \\ \text{s.t.} \quad & 0 \leq x \leq q \end{aligned}$$

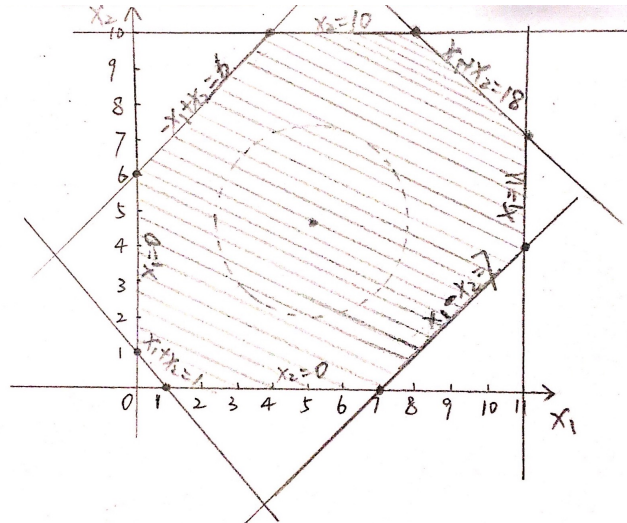
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1  %% Problem 4
2
3  A = [1 1 1 0 0;
4       0 0 0 1 1;
5       1 0 1 0 1;
6       1 1 1 1 0;
7       0 1 0 1 1];
8  p = [0.75; 0.35; 0.40; 0.75; 0.65];
9  q = [10; 5; 10; 10; 5];
10 n = 5;
11
12 cvx_begin quiet
13 variables x(n) t;
14 maximize (t - p'*x)
15 subject to
16     0 <= x <= q;
17     t * ones(n,1) <= A'*x
18 cvx_end
19
20 x
21 t
22 optval == cvx_optval

```

2. The optimal solution is $\mathbf{x} = \begin{bmatrix} 0 \\ 0 \\ 5 \\ 5 \\ 5 \end{bmatrix}$. The optimal value is \$1. This means that buying 5 Security 3, 5 Security 4, 5 Security 5 will get an optimal payoff of \$1.

Problem 5



Let $\mathbf{y} = (y_1, y_2)$ be the center of the fountain, r be the radius. The optimization problem can be formed as:

$$\begin{aligned}
 & \text{maximize}_{y_1, y_2} \quad r \\
 & \text{subject to} \quad y_1 + y_2 - 1 \geq \sqrt{2}r \\
 & \quad \quad \quad 7 - y_1 + y_2 \geq \sqrt{2}r \\
 & \quad \quad \quad -y_1 + 11 \geq r \\
 & \quad \quad \quad 18 - y_1 - y_2 \geq \sqrt{2}r \\
 & \quad \quad \quad 6 + y_1 - y_2 \geq \sqrt{2}r \\
 & \quad \quad \quad 10 - y_2 \geq r \\
 & \quad \quad \quad y_2 \geq r \\
 & \quad \quad \quad y_1 \geq r \\
 & \quad \quad \quad r \geq 0
 \end{aligned}$$

The optimal solution is $\mathbf{y} = (5.3753, 4.8753)$. The optimal value is $r = 4.5962$. Attached is the MatLab code:

```

24 %% Problem 5
25
26 - a = [1 1 -1;
27       -1 1 7;
28       -1 0 11;
29       -1 -1 18;
30       1 -1 6;
31       0 -1 10;
32       0 1 0;
33       1 0 0];
34 - b = [sqrt(2); sqrt(2); 1; sqrt(2); sqrt(2); 1; 1; 1]
35
36 - cvx_begin quiet
37 - variables y(2) r;
38 - maximize r
39 - subject to
40 -     a * [y(1);y(2);1] >= b * r
41 -     r >= 0
42 - cvx_end
43
44 - y
45 - optval = cvx_optval

```