
The University of Queensland
School of Information Technology and Electrical Engineering

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COMS3000/7003 – Tutorial 9, Answers

Q1) Encrypt the following text with a Caesar cipher with a key 'E', i.e. with a shift of 4, which means the letter A in the plaintext will be mapped to the letter E in the ciphertext.

HELLOWORLD

Answer:

Use the following mapping:

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D

Plaintext: H E L L O W O R L D

Ciphertext: L I P P S A S V P H

Q2) Decrypt the following ciphertext, which was encrypted using a Vigenère cipher with the key ART. (Hint: Use the Vigenère table provided below.)

YFN GFM IKK IXA T

A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A
C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B
D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C
E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D
F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E
G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F
H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G
I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H
J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I
K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J
L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K
M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L
N	O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M
O	P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N
P	Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O
Q	R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P
R	S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q
S	T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R
T	U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S
U	V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T
V	W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U
W	X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V
X	Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W
Y	Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X
Z	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y

You need to perform the reverse operation of the encryption. For example, to decrypt the second letter (F), which was encrypted using the key R, you have to look up the letter F in the row R of the Vigenère table and see which letter it maps to in the top row.

YFN GFM IKK IXA T
ART ART ART ART A

YOU GOT ITR IGH T

Q3) The following ciphertext has been encrypted with a Vigenère cipher. Use the Kasiski test to determine the likely key length. (Some repeated trigrams have been highlighted).

KCCPKBGUFDPHQTYAVINRRTMVGRKDNBVFDETDGILTXXRGUD (45 characters per line)
DKOTFMBPVGEGLTGCKQORACQCWDNAWCRXIZAKFTLEWRPTYC
QKYVXCHKFTPONCQQRHJV AJUWETMCMSPKQDYHJV DAHCTRL
SVSKCGCZQQDZXGSFRLSWCWSJTBHAFS IASPRJAHKJRJUMV
GKMITZHFPDISPZLVLGWTFPLKKEBDPGCEBSHCTJRWXBAFS
PEZQNRWXCVCYGAONWDDKACKAWBBIKFTIOVKCGGHJVLNHI
FFSQESVYCLACNVRWBBIREPBWVFEXOSCDYGZWPFDTKFQIY
CWHJVLNHIQIBTKHJVNPIST

The trigram HJV occurs 5 times, at positions 108, 126, 264, 318, and 330.

The distances between each pair of consecutive occurrences are 18, 138, 54, and 12. (You may want to do the same for other repeating trigrams.) The greatest common divisor of these numbers is 6, so that is very likely the keyword length.

The factors of these distances are:

12: 1, 2, 3, 4, 6, 12
18: 1, 2, 3, 6, 18
54: 1, 2, 3, 4, 6, 9, 18, 27, 54
138: 1, 2, 3, 6, 23, 46, 69, 138

The greatest common divisor of these numbers is 6, so that is very likely the keyword length.

Q4) Encrypt the following plaintext word using a simple transposition cipher with 3 columns (no permutation of column order).

Plaintext: TRANSPOSITION

TRA
NSP
OSI
TIO
N

Answer:

Ciphertext: TNOTNRSSIAPIO

Q5) Alice has the following secret message that she wants to send to Bob:

$M1 = 10011101$

Previously, Alice and Bob created a shared random Key $K = 01011000$, which they have not used yet.

a) Encrypt the message $M1$ with a one-time pad using key K .

XOR table:

$1 \text{ xor } 1 = 0$

$1 \text{ xor } 0 = 1$

$0 \text{ xor } 0 = 0$

$0 \text{ xor } 1 = 1$

Answer:

$C1 = M1 \text{ xor } K$

$M1 = 10011101$

$K = 01011000$

$C1 = 11000101$

b) Eve is eavesdropping on the channel and obtains the ciphertext $C1$. What can she do to find the secret message $M1$?

Answer:

The best she can do is directly guessing the message $M1$. Knowing the ciphertext $C1$ does not give her any information about the message, i.e. the one-time pad is perfectly secure, irrespective of the amount of computing resources and time an attacker has.

c) Now let's assume Alice wants to send another secret message $M2$ to Bob. The problem is that Alice and Bob have run out of secret keys to use, so Alice decides to reuse key K for the new message $M2$. The resulting ciphertext is $C2 = M2 \text{ xor } K$.

$M2 = 11100010$

$K = 01011000$

$C2 = 10111010$

Eve has been eavesdropping all the time and she has observed both $C1$ and $C2$. Through some other means (e.g. social engineering or guessing), she managed to obtain the message $M1$. With her knowledge of $M1$, $C1$ and $C2$, she can easily find $M2$. How?

You might use the following properties of xor:

$A \text{ xor } B = B \text{ xor } A$ (Commutativity)

$A \text{ xor } (B \text{ xor } C) = (A \text{ xor } B) \text{ xor } C$ (Associativity)

$A \text{ xor } A = 0$

$A \text{ xor } 0 = A$

Answer:

$$M1 = C1 \text{ xor } K$$

$$K = M1 \text{ xor } C1$$

$$M1 = 10011101$$

$$C1 = 11000101$$

$$K = 01011000$$

Now Eve has the secret key K which she can use to decrypt M2.

$$M2 = C2 \text{ xor } K$$

$$C2 = 10111010$$

$$K = 01011000$$

$$M2 = 11100010$$

Alternatively, Eve can use the following relationship:

$$M2 = M1 \text{ xor } (C1 \text{ xor } C2)$$

$$C1 = 11000101$$

$$C2 = 10111010$$

$$C1 \text{ xor } C2 = 01111111$$

$$C1 \text{ xor } C2 = 01111111$$

$$M1 = 10011101$$

$$M1 \text{ xor } C1 \text{ xor } C2 = 11100010 \text{ which is indeed } = M2$$

This example illustrates the danger of reusing keys in a one-time pad.

Q6) Given is a Feistel Cipher with the following parameters:

8 bit blocks

$$K_i = 1010$$

$$F(R_{i-1}, K_i) = 1111 = \text{const}$$

The output of F is always 1111, no matter what the input. This does not make sense for a real cipher, but it allows us to demonstrate some important characteristics of Feistel ciphers.

Plaintext block: 1001 1100

Calculate the ciphertext using two rounds.

Decrypt the ciphertext by applying two Feistel rounds to the ciphertext. You should get the original plaintext.

Remember, at the end of each encryption and decryption process, the left and the right half are swapped one more time.

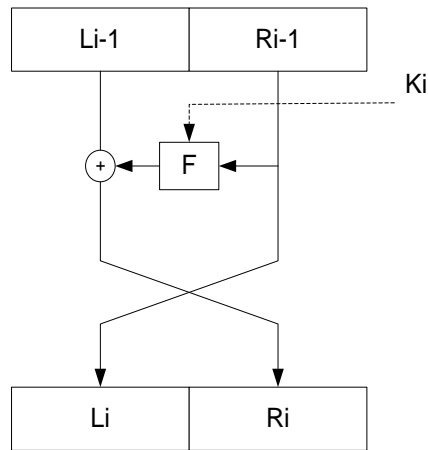
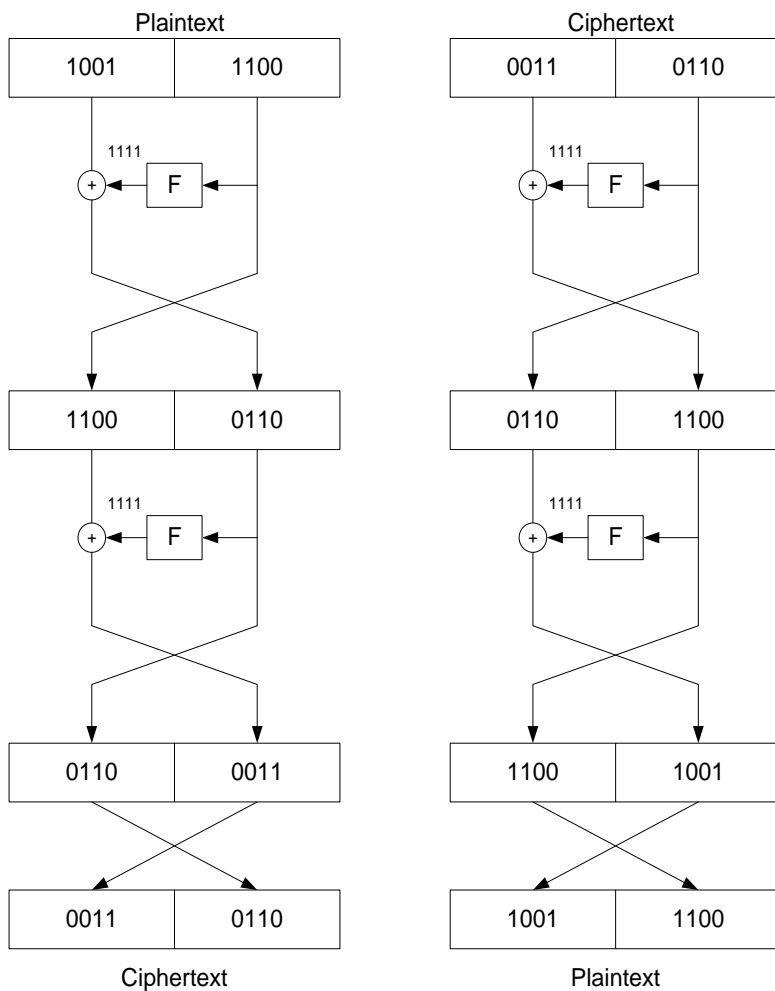


Figure 1 A Feistel Round

Answer:



The example shows that a Feistel cipher is reversible, even if F is not a reversible function. The security of a Feistel cipher relies on the properties of F . The constant function in our example is obviously not a very secure choice.