

# Link List

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## # Disadvantage of Link List.

- 1) can't access random node from the list. we have to go sequentially for access that node.
- 2) Need extra space for pointers.

## # Referential class.

It is the class using to create specially in Link List or tree.

It is known as referential class because it declare a data member as a pointer to an object of the same class.

Ex -

node \*next;

## # main code:

#include <bits/stdc++.h>  
using namespace std;

class node {

public:

int data;

node \*next; // self pointer.

// constructor of node class

DATE

node (int data) {

    this->data = data;

    this->~~next~~ next = NULL;

}

};

void insertatstart (node\*& head, int data) {

    node \* new\_node = new node(data);

    new\_node->next = head;

    head = new\_node;

}

int main () {

    node \* node1 = new node(1);

    cout << node1->data << endl;

~~insertatstart~~

    node \* head = node1;

    insertatstart (head, 11);

    print (head);

    return 0;

}

❶ output

11 4

# ~~class~~ define <sup>element</sup> in class.

we define element in class because after that we don't need to define data & pointer again and again. It is useful when we build custom list.

# Link List

## # Overview.

Link List is a type of list which contain the data and the address of the next node.

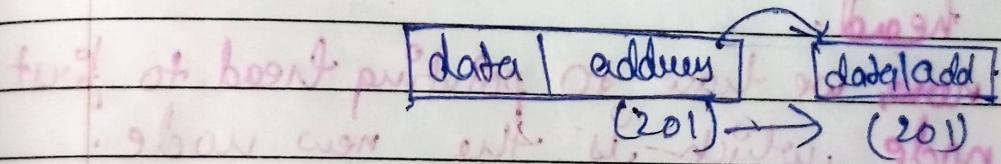
### → advantages.

- 1) No need to create new space
- 2) It is a dynamic memory.
- 3) It has no fixed size, we can add more nodes.
- 4) No shifting needed.

### → Types.

#### 1) Singly Link List

It contains single direction or one way direction with one data and one address section in a node.



## # operations on Link List

### 1) Create a node.

```

class node {
public:
    int data;
    node* next;
}
    
```

// define using class

```

node (int data) { // assign value then
    // store some data this -> data = data; // constr,
    // pointing to next this -> next = NULL;
    node.
}
}

```

## 2) Insert node at beginning.

```

void (int data, node *head) {
    ① — node *temp = new node(data);
    ② — temp -> next = head;
    ③ — head = temp;
}

```

- 1) create a new node space for inserting
- 2) ~~new node's~~ new node's next pointing to the existing node head;
- 3) now, we have to pointing head to first node, which is the new node.
- 4) Insert a node at the end.

```

void (int data, node *&tail) {

```

```

    node *new_data = new node(data);
    ① — tail -> next = new_data;
    ② — tail = new_data;
}

```

- 1) tail [which is last node], now tail's next pointing to the new node added.

2) And tail pointing the last node  
which is the new node [I just  
added].

4) print the node.on list

```
void print(*node * &head){ . . .
```

① — node \* temp = head;

② — while (temp != NULL) {

    cout << temp -> data << " and";

③ — temp = temp -> next;

}

    cout << endl;

}

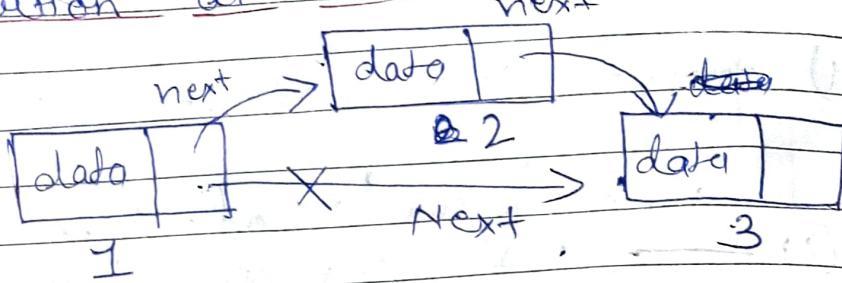
- 1) create a new pointer who point to the first node which indicated by "head".
- 2) run the loop while the temp is not equal to "NULL".  
[NULL represent the last node].
- 3) as we do in normal loop, increment the position by "1"; same we increment the temp value to the next node.

temp = temp -> next;

(5)

## ~~linked list insertion at mid node~~

### 5) Insertion at mid.



for inserting new node,

- the next of new node is pointing the node 3 which is pointed by the node 1.

And remove the next pointer of node 1 from node 3 to node 2.

# code at end of file

```
void insertatmid (&node * &tail, node *&head,
int position, int data) {
```

```
    node * temp = head;
```

```
    if (position == 1) {
```

```
        insertatstart (head, data)
        return;
```

```
(i) }
```

```
② — node * temp = head;
int count = 1;
```

while (~~count~~ < position - 1)

{

(3)

temp = temp  $\rightarrow$  next;

count ++;

}

// First Insert at end

if (tail  $\rightarrow$  next = NULL)

{

(4)

insert at last (tail, data);

}

node \* new\_node = new node(data);

(5)

new\_node  $\rightarrow$  next = temp  $\rightarrow$  next;

(6)

temp  $\rightarrow$  next = new\_data;

}

# 1) If the position is 1 than we call the insertatstart function for that.

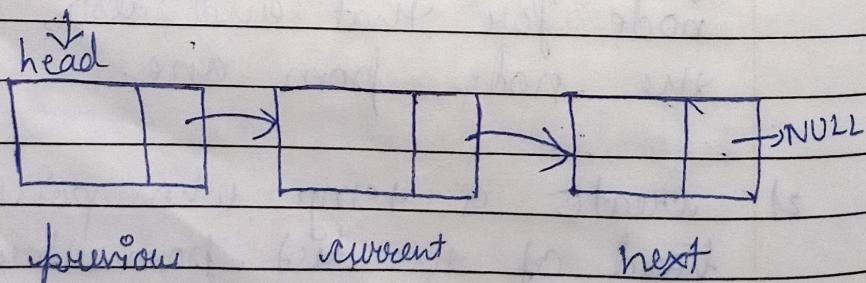
because it doesn't ~~start~~ insert at first position, it need a previous node for that and also we start the node from one.

2) Create a temp who point to the head of the list from starting point in purpose of traversing.

- 3) In this portion we traverse the list because we can't access the random node.
- 4) If in the case if the position is the last of the node. In this case we call the lastinsert function. It will work properly without this function but the tail is still not updated to the next node. for that we ~~will~~ need to call this function.
- 5) here we point the new node's next point to the node where previous next pointer point.
- 6) the previous pointer pointing to the ~~next~~ new node who is going to be added.

## # deletion of Link list

① deletion from mid and last node.

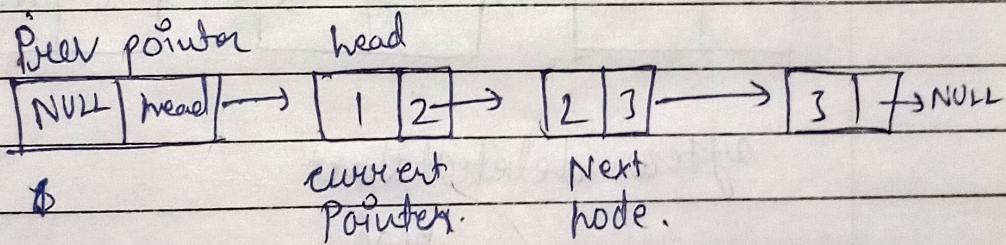


→ for deleting the current node firstly

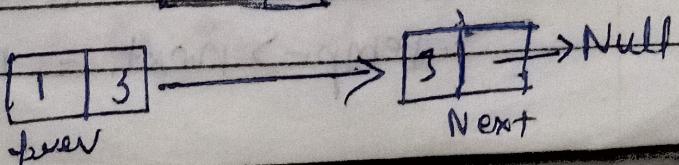
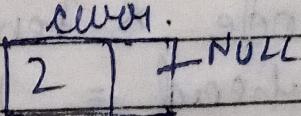
we have to pointing the previous's next to the next node and remove the current node add from it.

→ And for doing this pointer part first we have to initialize two pointer (previous, current) and traverse the list and increment the previous and current pointer till ( $\text{count} == \text{address}$ ).

\* for starting from first node, previous pointer should be null at starting than move forward and pointing to the next node which is Head.



→ when ( $\text{count} == \text{address}$ ) is arrived we do the first step remove curr add. from previous next pointer and locate to next node address and also NULL the ~~next~~ pointer of curr node.

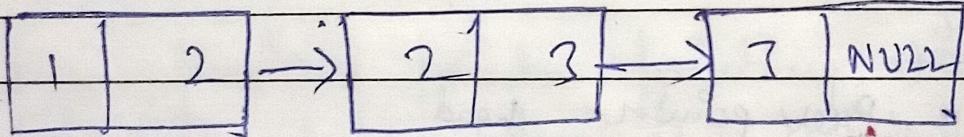


~~curr → Next = curr → Next;~~

$\text{prev} \rightarrow \text{Next} = \text{curr} \rightarrow \text{Next};$   
 $\text{curr} \rightarrow \text{Next} = \text{NULL};$   
delete curr;

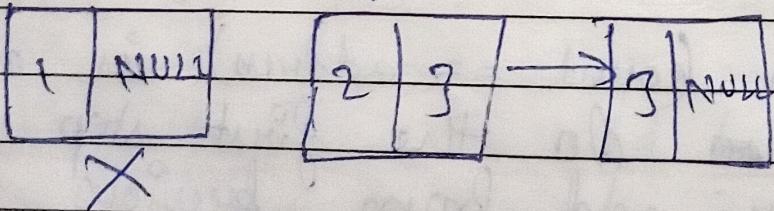
## \* Deletion at starting.

5) Deletion is not done by this process, because we need a previous node for that.  
That's why we call deleteAtFirst function where, we point the head the  $\text{head} \rightarrow \text{next}$  and make  $\text{head} \rightarrow \text{next}$  null  
head



after deleteAtStart

head



→ ~~node \* head~~

→  
 $\text{node} * \text{temp} = \text{head};$   
 $\text{head} = \text{head} \rightarrow \text{next};$   
 $\text{temp} \rightarrow \text{next} = \text{NULL};$

delete temp;

# main code.

void deletion (node \* &head, int add)

{

    if (add == 1)

    ②

        node \* temp = head;  
        head = head->next;  
        temp->next = NULL;  
        delete temp;

}

    else {

    ①

        node \* prev = NULL;  
        node \* curr = head;  
        int count = 1;

        while (count <= add)

{

            prev = curr;

            curr = curr->next;

            count++;

}

    }

    prev->next = curr->next;

    curr->next = NULL;

    delete curr;

}

→ Destructor code in class.

```
~ Node ()  
{
```

```
    if (this->next != NULL)
```

```
        delete next;
```

```
        next = NULL;
```

```
}
```

```
}
```

## # Doubly linked list

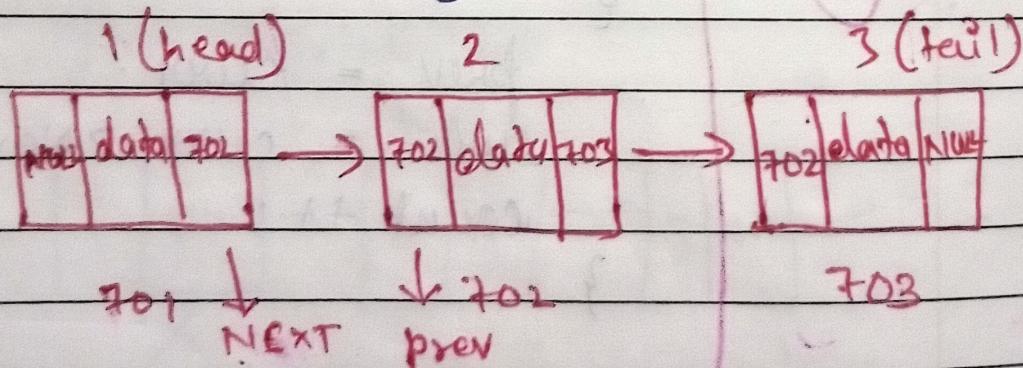
Q In double link list there is two pointer with every node

1) previous.

these pointer pointing to the previous node.

2) next

these pointing to the next node.



→ In this list the prev value of head is NULL.

and the next value of tail node is also NULL.

## # Insertion code:-

```
#include <bits/stdc++.h>
using namespace std;
```

```
class node {
```

```
    int data;
```

```
    node* prev;
```

```
    node* next;
```

```
    node(int data) {
```

```
        this->data = data;
```

```
        this->prev = NULL;
```

```
        this->next = NULL;
```

```
}
```

```
};
```

```
int main() {
```

```
    node *node1 = new node(5);
```

```
    node *head = node1;
```

```
    node *tail = node1 node1;
```

```
    print(head);
```

```
    return 0;
```

```
}
```

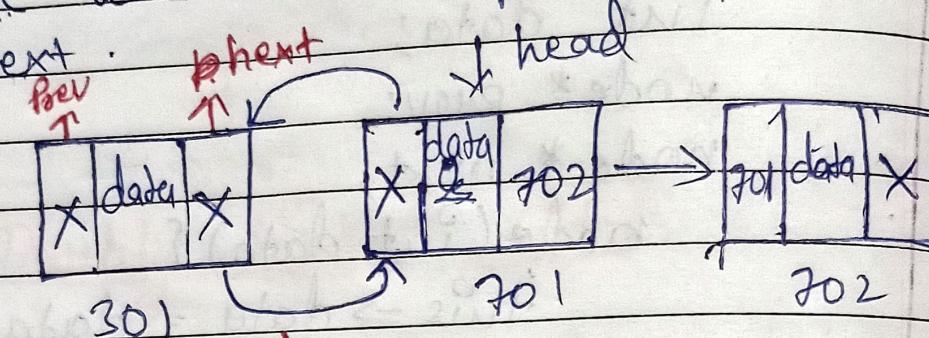
- ① It contains three items, ~~data~~ for store data, prev pointer for previous node, next pointer for next node.

# # operations

## 1) Insertion at start.

- Insert at start

We create two pointer prev and next.



(adding node)

new node's next should pointing to the head.

and the  $\text{head} \rightarrow \text{prev} = \text{new node}$ .

→ code.

```
void insertatstart (node * &head, int data)
```

① —  $\text{node} * \text{new\_node} = \text{new\_node}(\text{data})$

② —  $\text{new\_node} \rightarrow \text{next} = \text{head};$

③ —  $\text{head} \rightarrow \text{prev} = \text{new\_node};$

④ —  $\text{head} = \text{new\_node};$

```

void print(Node * &head)
{
    Node * temp = head;
    while (temp->next != NULL)
    {
        cout << temp->data << " ";
        temp = temp->next;
    }
    cout << endl;
}

```

```

int main()
{
    Node * node1 = new Node(4);
    Node * head = node1;
    insertAtStart(head, 40);
    insertAtStart(head, 64);
    print(head);
    return 0;
}

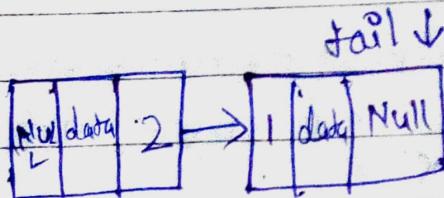
```

→ output

64 40 4

- 1) create new node space in heap.
- 2) new-node → next point to the head.
- 3) head move head to the newly added node.
- 4) head → prev pointing to the new node.

## 2) Insertion at end.

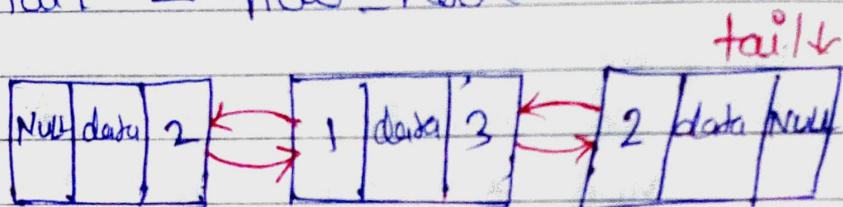


for inserting new node at end.

tail → next = new\_node

new\_node → prev = tail

tail = new\_node.



→ code.

```
void insertatend (node *ltail, int d)
```

```
{ int temp = new node(d);
```

```
tail → next = temp temp;
```

```
temp → prev = tail;
```

```
tail = temp;
```

```
}
```

```
int main() {
```

```
node * node1 = new node(4);
```

```
node * tail = node1;
```

```
insertatend(tail, 6);
```

```
insertatend(tail, 12);
```

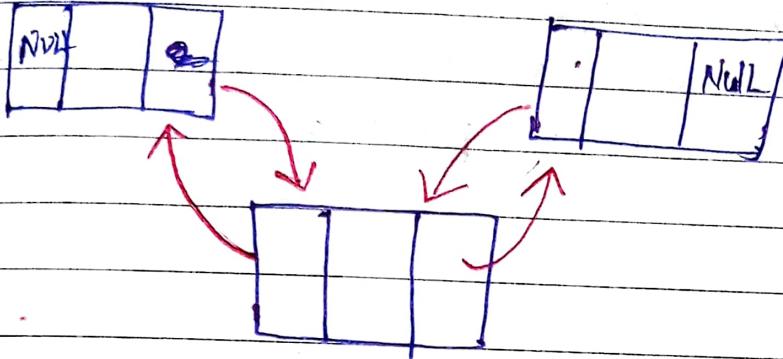
```
return 0;
```

output

4, 6, 12

3) Insertion at mid.

for inserting in mid.



→ Code for it

```
void insertatmid (node * &tail, node *&head, int position, int data)  
{
```

```
    if (position == 1) {  
        insertatstart (head, data);  
        return;  
    }
```

```
    node * temp = head;
```

```
    int count = 1;
```

```
    while (temp->next)
```

```
        while (count < position - 1)  
            {
```

```
                temp = temp->next;
```

```
                count++;
```

4

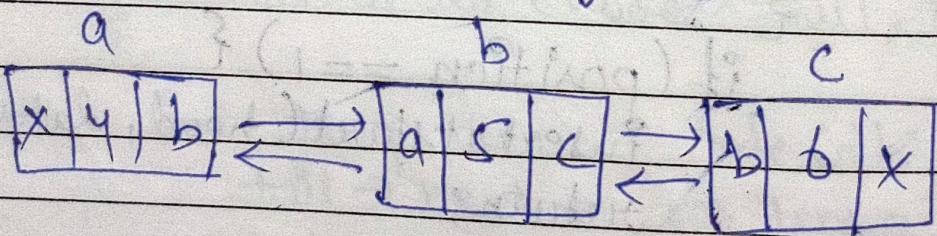
```
if (temp->next == NULL) {  
    insertatend (tail, data);  
    return;  
}
```

```
node * new_node = new node(data);  
new_node->next = temp->next;  
new_node->prev = temp;  
temp->next = new_node;  
temp->next->prev = new_node;
```

```
}
```

## # Deletion in Doubly Linked List

for deletion in doubly linked list let understand the logic first.

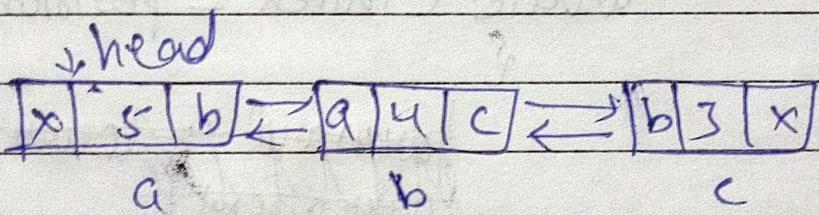


if we want to delete node b.  
what we have to do it.

- 1) the next of node a is equal to the node c
- 2) the prev of the node c is equal to the node a and delete the node b.

3) and before deletion you have to null the next and prev ~~next~~ pointer of the deletion node, otherwise it give segmentation fault.

→ exception case  
if it is first node in ~~in~~ list then  
simply null the next of node ~~of~~,  
and also null the ~~the~~ next node's  
prev pointer -  
and delete the node



→ code

~~void deletion(node \* &head, int position)~~

~~if (position == 1) {~~

~~node \* temp = head;~~

~~head -> next -> prev = NULL;~~

~~head -> next = NULL;~~

~~head = head -> next;~~

~~head -> next = NULL;~~

~~delete temp;~~

~~void deletion(node \* &head, int position)~~

~~if (position == 1) {~~

~~node \* temp = head;~~

```
temp->next->prev = NULL;  
head = temp->next;  
temp->next = NULL;  
Delete temp;
```

{

else {

node \* prev = NULL;

node \* curr = head;

~~node \* forward = curr->next;~~

int index = 1;

while (index &lt; position)

{

~~forward = curr->next;~~  
~~curr = forward;~~

prev = curr;

curr = curr-&gt;next;

index += 1;

}

curr-&gt;prev = NULL

prev-&gt;next = curr-&gt;next;

curr-&gt;next-&gt;prev = prev.

curr-&gt;next = NULL

delete curr;

}

3

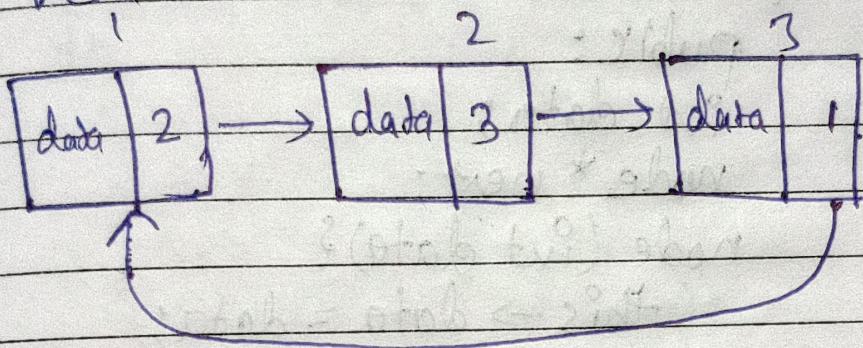
## # Circular Link List.

Circular link list is list in which the last node's pointer pointing to the first node.

and form a circular link list.

head

tail



### → Insertion in CLL

We use a tail pointer which point to the last node always and its pointer point to the head or first node.

→ It has two cases.

1) where the tail is NULL which means the empty list.

In that case we must create a node store data and represent it as tail. Its next pointer pointing to itself because it has none other node.

2) another case is non-empty list where new node store the data

represent as tail and pointing to  
the ~~last~~ first node.

→ code

class node {

public:

int data;

node \* next;

node (int data) {

this → data = data;

this → next = NULL;

}

void insertion(node \* tail, int element, int data) {

if (tail == NULL)

{

node \* new-node = new node (data);

tail = new-node;

new-node → next = new-node;

}

else {

node \* temp = tail;

while (temp → next != element)

{

temp = temp → next;

}

node \* new-node = new node (data);

create  
node

insert  
in  
list

new-node → next = ~~temp~~ temp → next;  
temp → next = new-node;

}  
}

void print(node \*tail) {

node \*temp = tail;

do

{

cout << tail → data << " ";

tail = tail → next;

} while (tail != temp);

cout << endl;

}

int main() {

node \*tail = NULL;

insertion(tail, 5, 10);

insertion(tail, 10, 22);

insertion(tail, 22, 14);

insertion(tail, 14, 20);

print(tail);

return 0;

}

Output

10 22 14 20

## → Deletion in CLL.

Create two pointers previous and current.  
We just have to do is pointing  
the previous pointer to the current  
next pointer.  
and delete that node.

code

node.

~~~node()~~

~~~  
node  
function~~

```
int value = this->data;  
if (this->data.next != NULL){  
    delete next;  
    this->next = NULL;  
}
```

Deletion

```
void deletion (node * &tail, int  
value) {  
    if (tail == NULL){  
        cout << "No element" << endl;  
    }  
    else {  
        node * prev = tail;  
        node * curr = prev->next;  
        while (curr->data != value)  
        {  
            prev = curr;  
            curr = curr->next;  
        }  
    }
```

$\text{prev} \rightarrow \text{next} = \text{curr} \rightarrow \text{next};$

$\text{if } (\text{prev} == \text{curr}) \quad // \text{single node}$

$\text{tail} = \text{NULL};$

}

$\text{if } (\text{tail} == \text{curr}) \quad // \text{last node}$

$\text{tail} = \text{prev};$

}

$\text{curr} \rightarrow \text{next} = \text{NULL};$

$\text{delete curr};$

3

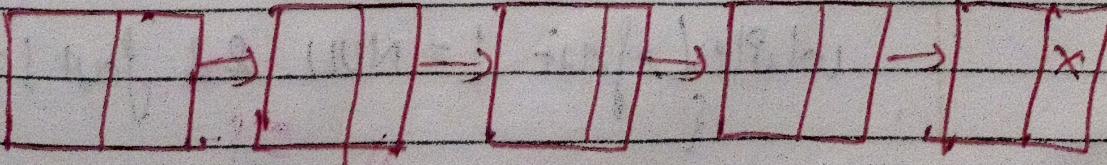
## # Floyd cycle Detection.

This algorithm is specially for detecting the loop in a link list.

where two pointers are used and both move with different speed. means, the other pointer is move twice time faster than the first one.

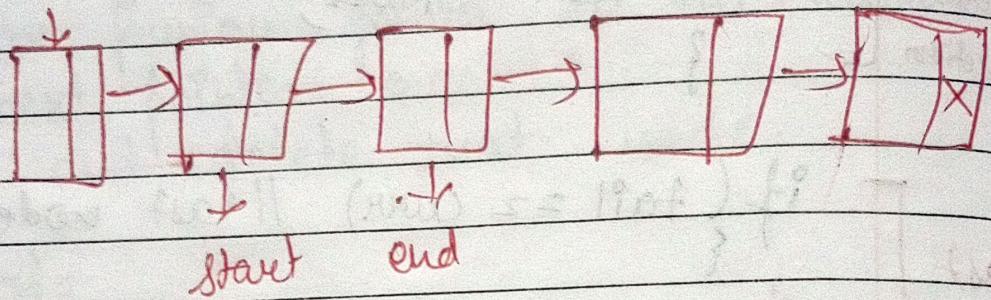
head

start  $\rightarrow$   
end  $\rightarrow$



now start move one step while  
end move twice or two step  
forward.

head



it traverse all the node in  
same time twice a time and  
if we find any loop or compare  
the node we can use this  
algorithm for that.

→ it use in many question.

1) find mid of a Link list.

2) detect loop in link list. etc so on.

→ example

find loop in Link list.

bool isLoop(Node \* head){

node \* slow = head; } - pointing  
node \* fast = head; } is head.

while (fast != NULL && fast->next != head)