

University of Glasgow
 School of Mathematics and Statistics
 Mathematics 5M Advanced Numerical Methods

Exercise Sheet 2

Computing questions 2

- C1 Write a MATLAB function which implements Gaussian Elimination without pivoting for computing the LU decomposition of a matrix $\mathbf{A} \in \mathbb{C}^{m \times m}$

```

U = A, L = I
for k = 1 to m - 1
    for j = k + 1 to m
        ljk = ujk/ukk
        uj,k:m = uj,k:m - ljkuk,k:m
    
```

- C2 Write a MATLAB function which solves the matrix problem $\mathbf{Ax} = \mathbf{b}$, for $\mathbf{x} \in \mathbb{C}^m$ where $\mathbf{A} \in \mathbb{C}^{m \times m}$ and $\mathbf{b} \in \mathbb{C}^m$ are given.

- C3 Use your algorithm to solve the four examples in T3 below. What do you notice?

- C4 Write a MATLAB function which implements Gaussian Elimination with partial pivoting for computing the LU decomposition of a matrix $\mathbf{A} \in \mathbb{C}^{m \times m}$.

```

U = A, L = I, P = I
for k = 1 to m - 1
    Select i ≥ k to maximise |uik|
    uk,k:m ↔ ui,k:m (interchange two rows)
    lk,1:k-1 ↔ li,1:k-1
    pk,:  ↔ pi,: 
    for j = k + 1 to m
        ljk = ujk/ukk
        uj,k:m = uj,k:m - ljkuk,k:m
    
```

- C5 Use this modified algorithm to solve two examples in T3 below. Note how the problem identified in C3(iv) has now been resolved.

C6 Write a MATLAB function which implements Cholesky Factorisation of a square matrix $\mathbf{A} \in \mathbb{C}^{m \times m}$

```

R = A
for k = 1 to m
    for j = k + 1 to m
         $R_{j,j:m} = R_{j,j:m} - R_{k,j:m} \overline{R_{kj}} / R_{kk}$ 
    end for
     $R_{k,k:m} = R_{k,k:m} / \sqrt{R_{kk}}$ 
end for

```

C7 Compute the Cholesky Factorisation of the hermitian positive definite matrix

$$\mathbf{A} = \begin{bmatrix} 2 & -1 & 0 \\ -1 & 2 & -1 \\ 0 & -1 & 2 \end{bmatrix}.$$

Tutorial questions 1

T1 Use LU decomposition to solve the following systems $\mathbf{Ax} = \mathbf{b}$,

$$(i) \mathbf{A} = \begin{bmatrix} 2 & -1 \\ 1 & 3 \end{bmatrix}, \quad \mathbf{b} = \begin{bmatrix} 3 \\ 5 \end{bmatrix}, \quad (ii) \begin{bmatrix} 2 & 3 & -1 \\ 1 & 0 & 1 \\ -1 & 2 & -2 \end{bmatrix}, \quad \mathbf{b} = \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix},$$

$$(iii) \mathbf{A} = \begin{bmatrix} 2 & 1 & 5 \\ 5 & 1 & -5 \\ 1 & 2 & 0 \end{bmatrix}, \quad \mathbf{b} = \begin{bmatrix} 7 \\ 9 \\ 4 \end{bmatrix}, \quad (iv) \begin{bmatrix} 0 & 6 & 4 \\ -2 & 2 & 2 \\ 2 & 1 & -1 \end{bmatrix}, \quad \mathbf{b} = \begin{bmatrix} 4 \\ 2 \\ 1 \end{bmatrix}.$$

Consider the general matrix system $\mathbf{Ax} = \mathbf{b}$, $\mathbf{A} \in \mathbb{C}^{m \times m}$, $\mathbf{b} \in \mathbb{C}^m$.

T2 Show that the number of flops required for forward substitution and backward substitution is $O(m^2)$ respectively.

T3 How many flops would it require to compute the pivot at each step of Gaussian elimination if every available element needed to be tested against some criterion?

T4 Show that the number of flops required for Gaussian elimination with partial pivoting (algorithm above) is $O(\frac{2}{3}m^3)$

T5 Show that the number of flops required for Cholesky factorisation of a hermitian positive definite matrix is half as many as Gaussian elimination.