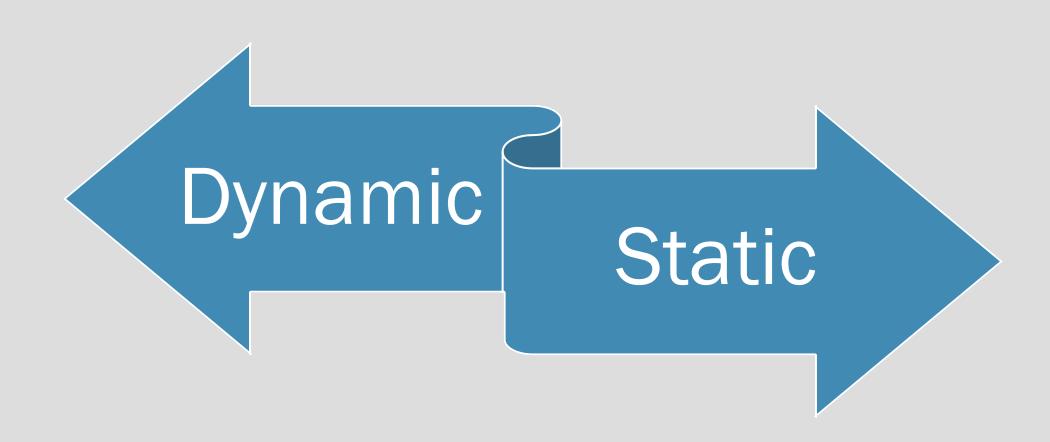
WHY TO LOCALIZE NON-LOCAL CONTROL

Ross Tate

Enforcement Strategies



Lexical Switching

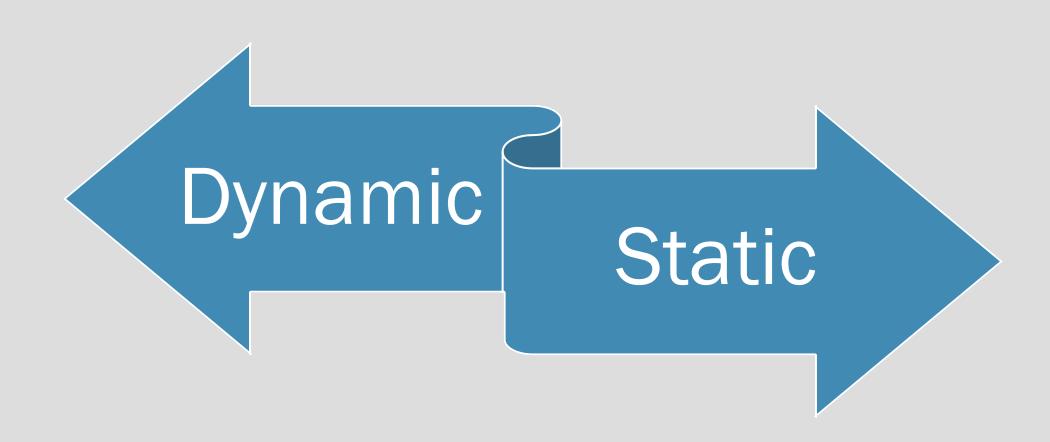
foo1 foon executing bar stacklet (foo-bar connection) bar1 contains host frame? barn • executing baz stacklet (bar-baz connection) baz1 • contains host frame? bazn • yields control to foo using direct reference to foo-bar connection

What to do when there is a host frame?

- Only suspend up to the host frame?
 - Performance: every call from host into WebAssembly allocates a stack
 - Migration: wasm programs can come to depend on this functionality
 - Difficult to change platform-specific host code to cross-platform wasm code
- Trap?
 - Migration: difficult to change wasm host code to platform-specialized code
- Corner cases?
 - Interface types intertwines code from two sources (including host)
 - Different behavior for JavaScript frames versus C/C++/assembly frames?

Host frames migrate between being foreign versus wasm

Enforcement Strategies



Effect Signatures

- (func \$foo (param ti*) (result to*) (effect ...))
 - Specifics of ... depends on other design choices (e.g. lexical vs. dynamic scope)
 - ... at least indicates whether or not \$foo can suspend
- Caller of suspending \$foo must either
 - set up infrastructure for \$foo to be able to suspend
 - E.g. installs an effect handler or prompt or calls \$foo on a new stacklet
 - or declare itself to also be suspending (under same conditions as \$foo)
- No surprise suspends when calling effect-free wasm functions! (i.e. "opt-in" design)
 - Make only effect-free wasm functions accessible to host language (e.g. C or JS)
 - Guarantees host frames will not be captured

Faulty Comparison

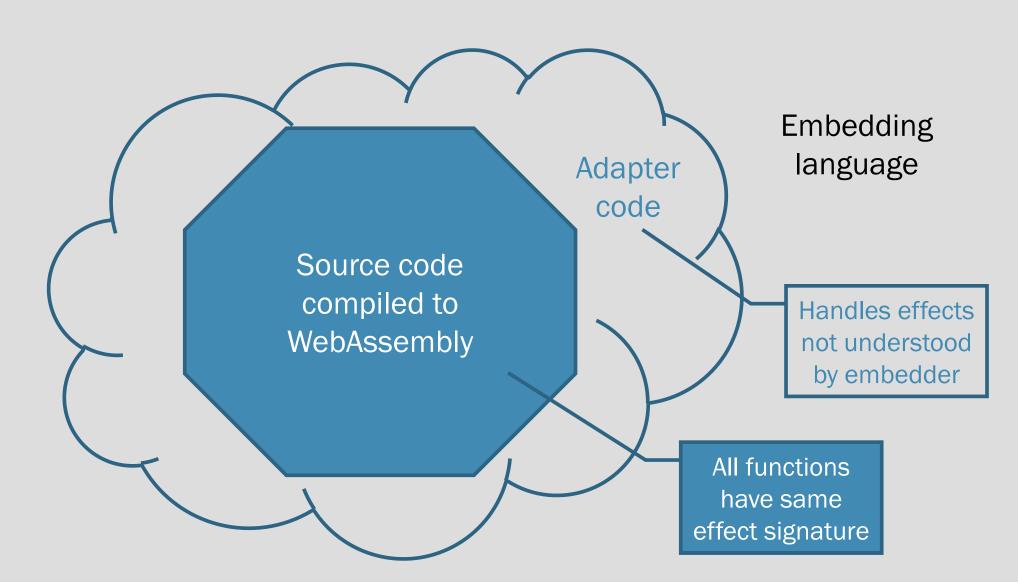
Java Exceptions

- Many exception types within module
- No guarantees (RuntimeException)
- Written by hand

Wasm Effect Signatures

- One effect type in typical module
- Strong guarantees
- Procedurally generated

Easy to generate (preliminary estimate)



Simpler in the long run

No unknown/foreign effects

- No need for handle_all (i.e. catch_all)
- No need for eventref (i.e. exnref)
- No need for dynamic-wind (i.e. unwind)

Program invariants automatically respected

- Program can easily keep track of all its stack frames (for linear-memory GC root marking)
- Stacks cannot be unexpectedly duplicated inside critical regions
- Shadow stacks stay aligned with wasm stacks

Enforcement Strategies

