

# Lecture 21

March 29, 2018 9:00 AM

## Pages

- Physical memory is divided into equal-sized frames
- Divide logical memory into equal size pages

Frame size = Page size

- To run a program w/ a size of  $n$  pages, we can do it if there are  $n$  free frames
- Load the  $n$  pages into the  $n$  frames
- Create a **page table** to translate page numbers to frame numbers

### Page Logical Memory

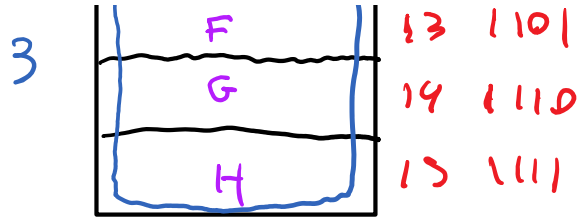
0	A	0	0000
	B	1	0001
	C	2	0010
	D	3	0011
1	E	4	0100
	F	5	0101
	G	6	0110
	H	7	0111
2	I	8	1000
	J	9	1001
	K	10	1010
	L	11	1011

CP4

### Frame Physical Memory

0		0	0000
		1	0001
		2	0010
		3	0011
1	I	4	0100
	J	5	0101
	K	6	0110
	L	7	0111
2	A	8	1000
	B	9	1001
	C	10	1010
	D	11	1011
3	E	12	1100
	F	13	1101
	G	14	1110
		15	1111

CP4



Page      Frame

0 → 2  
2 → 1  
1 → 3

Page Table

00	0	2	10
01	1	3	11
10	2	1	01

CPU  
want

0101  
↑      ↑  
page    offset  
number

Replace page #  
with frame # → 1101

Consider 8-bit Addresses

⇒  $2^8 = 256$  byte addresses



- How big is a page? Ans  $2^5 = 32$  bytes
- How many pages are there? Ans  $2^3 = 8$  pages

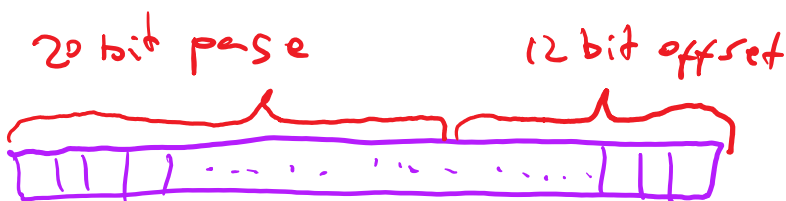
- How many entries are in the page table? Ans  $2^3 = 8$
- How big is the page table?  $8 * \text{storage needed for a number}$

Consider a 32-bit Address space



- If the page specifies 10 bits, How many bits are in the offset? Ans 22
- How big is a frame? Ans  $2^{32}$  bytes  $\sim 4,096,000$
- How many pages are there? Ans  $2^{10} = 1024$
- How many entries in page table? Ans  $2^{10} = 1024$
- How big is the page table? Ans: Need more than 1 byte to count up to 1024. Need 2 bytes for each integer in the page table  $\Rightarrow 1024 * 2 \text{ bytes} = 2 \text{ KB}$

Consider 32-bit Addresses



- How big is a page? Ans  $2^{12} \text{ bytes} = 4 \text{ KB}$

- How many pages are there? Ans  $2^{20} \sim 1,000,000$
- How many entries are in the page table? Ans  $2^{20}$
- How big is the page table?  $2^{20} \times 4 \text{ bytes (size of int)}$   
 $= 4 \text{ MB}$

But, Page table is in memory, which means every memory access requires 2 memory access  $\Rightarrow$  System is slower

Partial Solution

- Associative memory  
 (a cache that is electronically near the CPU)

$\Rightarrow$  Effective Access Time (EAT)

$$EAT = \alpha \overset{\text{Hit}}{(1 + \epsilon)} + (1 - \alpha) \overset{\text{Miss}}{(2 + \epsilon)}$$

$$(0 \leq \alpha \leq 1)$$

$\alpha = 0.99 \Rightarrow 99\%$  of the time there is a hit  
 $1\%$  of the time a miss

