

Part I.

Data-structures.scm

First, we need to add the data structures, we need to add the nested procedure:

```
(nested-procedure
  (bvar symbol?)
  (count number?)
  (name symbol?)
  (body expression?)
  (env environment?))
```

Then the extension of the environment:

```
(extend-env-rec-nested
  (id symbol?)
  (bvar symbol?)
  (count number?)
  (body expression?)
  (saved-env environment?))
```

Environments.scm

First of all we need to extend out environment with our count variable which will be n in my case:

```
(extend-env  
  'n (num-val 0))
```

Then we need to handle the values in apply-env:

```
(cases expval val  
  (proc-val (procval)  
    (cases proc procval  
      (nested-procedure (bvar count name body env)  
        (proc-val (nested-procedure bvar count var body env)))  
      (else procval)))  
  (else val))
```

And now we can handle the new environment extension:

```
(extend-env-rec-nested (id bvar count body saved-env)  
  (if (eqv? search-sym id)  
    (proc-val (nested-procedure bvar count id body env))  
    (apply-env saved-env search-sym)))
```

Interp.scm

Now we can start handling the interpreter

```
(proc-nested-exp (var count name body)
  (proc-val (nested-procedure var (expval->num (value-of (var-exp count)
env)) name body env)))

(call-nested-exp (rator rand count)
  (let ((rnd (value-of rand env))
        (procedure (expval->proc (value-of rator env)))
        (newcount (expval->num (value-of count env))))
    (apply-procedure
      (cases proc procedure
        (nested-procedure (var count name body env)
          (nested-procedure var newcount name body env))
        (else procedure))
      (rnd)))

(letrec-nested-exp (p-name b-var b-count p-body letrec-body)
  (let ((count (expval->num (value-of (var-exp b-count) env)))
        (value-of letrec-body
          (extend-env-rec-nested p-name b-var count p-body env))))
```

And we also need to add the nested-procedure to apply-procedure:

```
(nested-procedure (var count name body saved-env)
                  (begin
                     (recursive-displayer name count)
                     (value-of body (extend-env 'n (num-val count)
                                                (extend-env var arg saved-env)))
                   ))
```

Lang.scm

Now we can add them to our grammar:

```
(expression
  ("proc-nested" "(" identifier "," identifier "," identifier ")" expression)
  proc-nested-exp)

(expression
  ("call-nested" "(" expression expression "," expression ")")
  call-nested-exp)

(expression
  ("letrec-nested" identifier "(" identifier "," identifier ")" "=" expression "in"
expression)
  letrec-nested-exp)
```

Translator.scm

Now we need to translate all of the newly added expressions

```
(proc-exp (var body)
  (proc-nested-exp var 'n 'anonym
    (translation-of body env))
)

(call-exp (rator rand)
  (let* ((operator (translation-of rator env))
    (operand (translation-of rand env))
    (count (cases expression operator
      (var-exp (var) (diff-exp
        (var-exp 'n)
        (const-exp -1)))
      (else (const-exp 1)))))
    (call-nested-exp operator operand count))
)

(letrec-exp (p-name b-var p-body letrec-body)
  (letrec-nested-exp p-name b-var 'n
    (translation-of p-body env)
    (translation-of letrec-body env))
) (letrec-exp (p-name b-var p-body letrec-body)
  (letrec-nested-exp p-name b-var 'n
    (translation-of p-body env)
    (translation-of letrec-body env))
)
```

Tests.scm

Extra test cases:

Test 1: here we evaluate a function based on the output of another function. Notice how b is unused, it is there only to check how let handles letrec in its body

```
(double-letrec-1
  "let b = 1 in letrec func1(y) = -(10,y) in letrec func2(z) = if zero?(z) then 1 else 2 in
(func1 (func2 0))"
  9)
```

Test 2: here we have an anonymous function which gets constructed on the fly and its result is passed to the letrec which will get called a max. of 2 times.

```
(letrec-forever
  "letrec l(y) = if zero?(y) then 1 else (l 0) in (l (proc(x) -(x,-1) 4))"
  1)
```

Test 3: here a letrec constructs different anonymous functions on the fly depending on its input.

```
(letrec-notletrec-letrec
  "letrec fun1(r) = if zero?(r) then
    (proc(x) if zero?(r) then -(r,-(0,r)) else 0 r)
  else
    (proc(x) if zero?(-(r, 1)) then 1 else 0 r)
  in (fun1 1)"
  1)
```

Part II.

Translator.scm

The translation of var-exp:

```
(let ((count (apply-senv-number senv var)))
      (if (> count 0) (var-exp (string->symbol
                                (string-append (symbol->string var) (number->string
count))))
                    (eopl:error 'translation-of "unbound variable in code: ~s" var)))
```

Translation of let:

```
(let ((count (apply-senv-number senv var)))
      (if (> count 0) (var-exp (string->symbol
                                (string-append (symbol->string var) (number->string
count))))
                    (eopl:error 'translation-of "unbound variable in code: ~s" var)))
```

Translation of proc:

```
(let* ((count (apply-senv-number senv var))
        (var-string (symbol->string var))
        (old-var
         (string-append var-string (number->string count)))
        (new-var
         (string-append var-string (number->string (+ 1 count)))))
      (message (if (> count 0)
                    (string-append var-string " has been reinitialized. " new-
var " is created and shadows " old-var ".")
                    ""))
      (var-field (string->symbol (string-append new-var " " message))))
(proc-exp var-field
          (translation-of body (extend-senv var senv))))
```

Senv:

```
(lambda (senv var)
  (cond
    ((null? senv) 0)
    ((eqv? var (car senv))
     (+ 1 (apply-senv-number (cdr senv) var)))
    (else
     (+ 0 (apply-senv-number (cdr senv) var)))))
```