# Chapter 7: Advanced Frequent Pattern Mining

- Pattern Mining: A Road Map
- Pattern Mining in Multi-Level, Multi-Dimensional Space
- Constraint-Based Frequent Pattern Mining
- Mining High-Dimensional Data and Colossal Patterns
- Mining Compressed or Approximate Patterns
- Sequential Pattern Mining
- Graph Pattern Mining
- Summary

## Graph Pattern Mining

- Frequent subgraphs
  - ■A (sub)graph is *frequent* if its *support* (occurrence frequency) in a given dataset is no less than a *minimum support* threshold
- ■Applications of graph pattern mining
  - ■Mining biochemical structures
  - ■Program control flow analysis
  - ■Mining XML structures or Web communities
  - ■Building blocks for graph classification, clustering, compression, comparison, and correlation analysis

# Example: Frequent Subgraphs

#### **GRAPH DATASET**

$$(A) \qquad (B) \qquad (C)$$

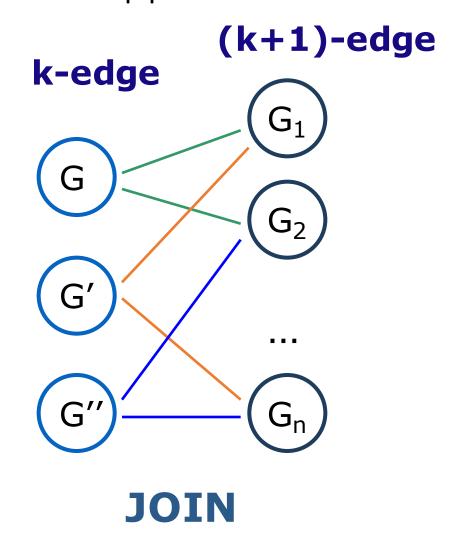
# FREQUENT PATTERNS (MIN SUPPORT IS 2)

$$(1) \qquad (2) \qquad \sqrt[N]{}$$

# Properties of Graph Mining Algorithms

- ■Search order
  - ■breadth vs. depth
- ■Generation of candidate subgraphs
  - ■apriori vs. pattern growth
- ■Elimination of duplicate subgraphs
  - ■passive vs. active
- ■Support calculation
  - ■embedding store or not
- ■Discover order of patterns
  - $\blacksquare$ path → tree → graph

# Apriori-Based Approach

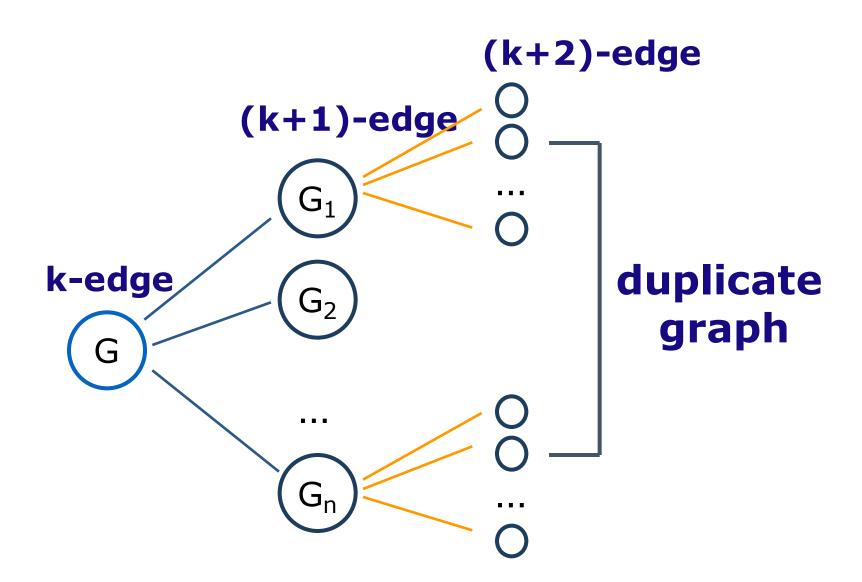


# Apriori-Based, Breadth-First Search

- Methodology: breadth-search, joining two graphs
- AGM (Inokuchi, et al. PKDD'00)
  - generates new graphs with one more node

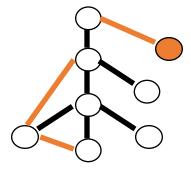
- FSG (Kuramochi and Karypis ICDM'01)
  - generates new graphs with one more edge

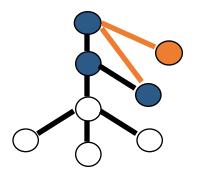
## Pattern Growth Method



# GSPAN (Yan and Han ICDM'02)

### **Right-Most Extension**





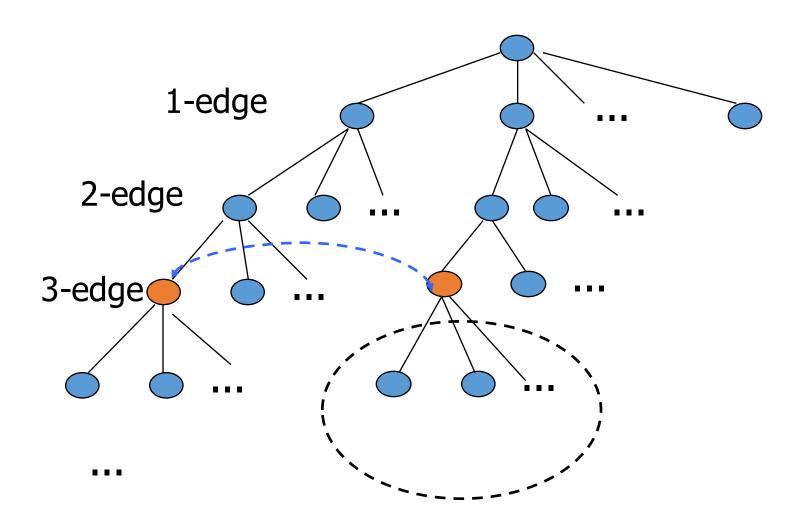
# **Theorem: Completeness**

The Enumeration of Graphs using Right-most Extension is COMPLETE

# Lexicographic Ordering in Graph

- It can tell us whether the graph has been discovered.
- And more, the most important, if a graph has been discovered, all its children nodes in the hierarchy must have been discovered.

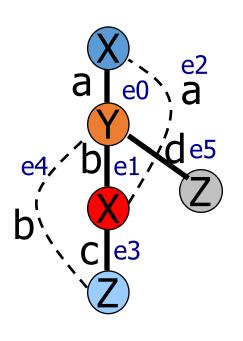
# Lexicographic Ordering in Graph

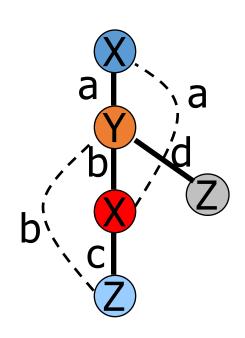


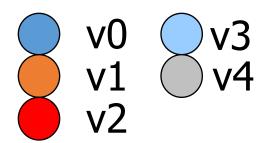
### DFS code and Minimum DFS code

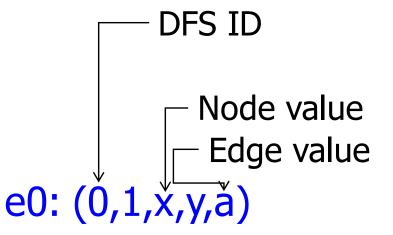
- Edge: (v<sub>i</sub>, v<sub>i</sub>, I(v<sub>i</sub>), I(v<sub>i</sub>), I(v<sub>i</sub>,v<sub>i</sub>))
  - v<sub>i</sub>: node id by DFS
  - I(v<sub>i</sub>): node value (e.g. Oxygen, Carbon, Person name)
  - $I(v_i, v_i)$ : edge value (e.g. covalent bond, relationship)
- Graph: sequence of edges of following order:
  - Extend one *new* node, add the *forward edge* that connect one node from old nodes to the new node.
  - Add all backward edges that connect this new node to other old nodes
  - repeat this procedure.

## DFS code









e1: (1,2,y,x,b)

e2: (2,0,x,x,a)

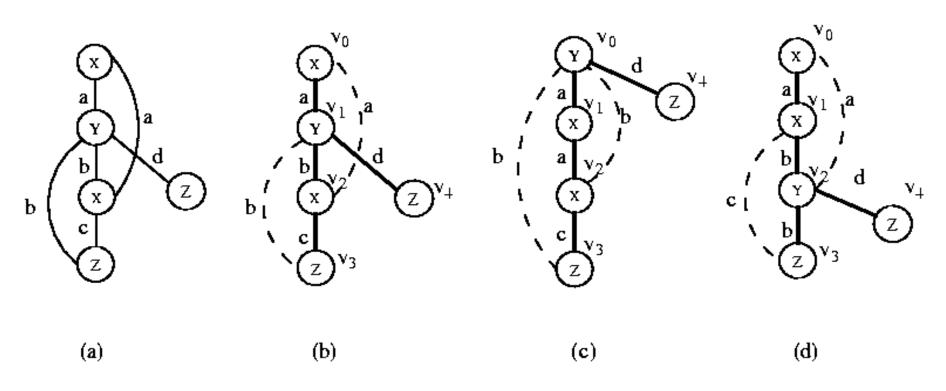
e3: (2,3,x,z,c)

e4: (3,1,z,y,b)

e5: (1,4,y,z,d)

Forward expansion Backward Linking

### DFS code and Minimum DFS code

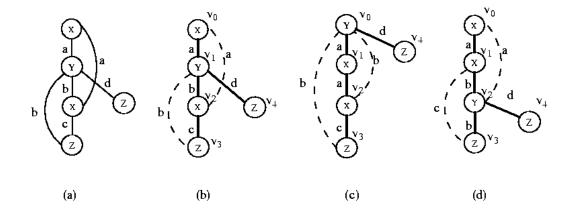


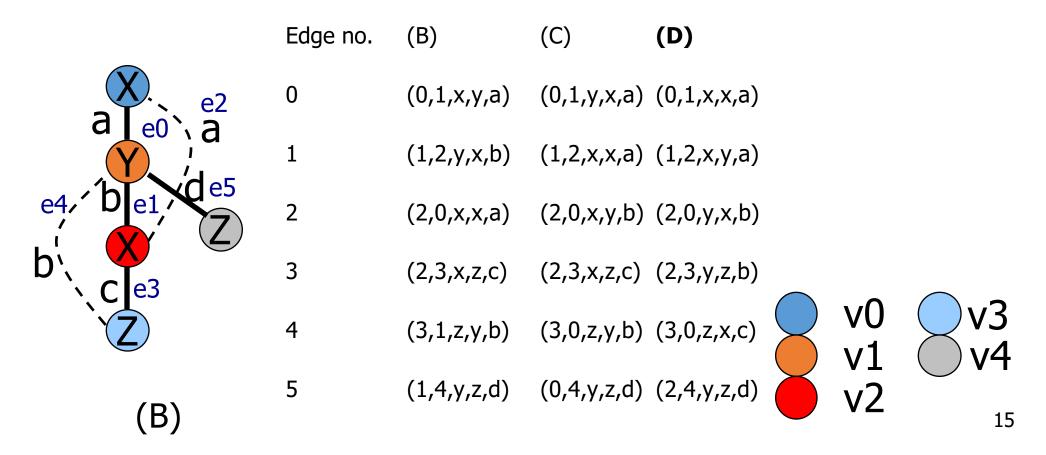
- For same graph, it might have different DFS codes due to its order of nodes and edges visited.
  - (b), (c), (d) are isomorphic graphs of (a)
- Use minimum DFS code to cancel the representation uncertainty.

## Minimum DFS code

Each Graph may have lots of DFS code (why?):

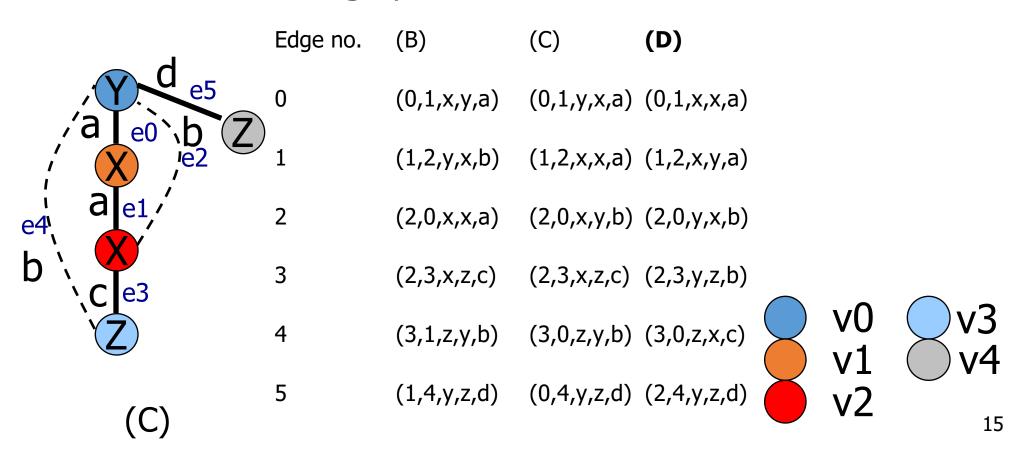
one smallest lexicographic one is its Minimum DFS Code





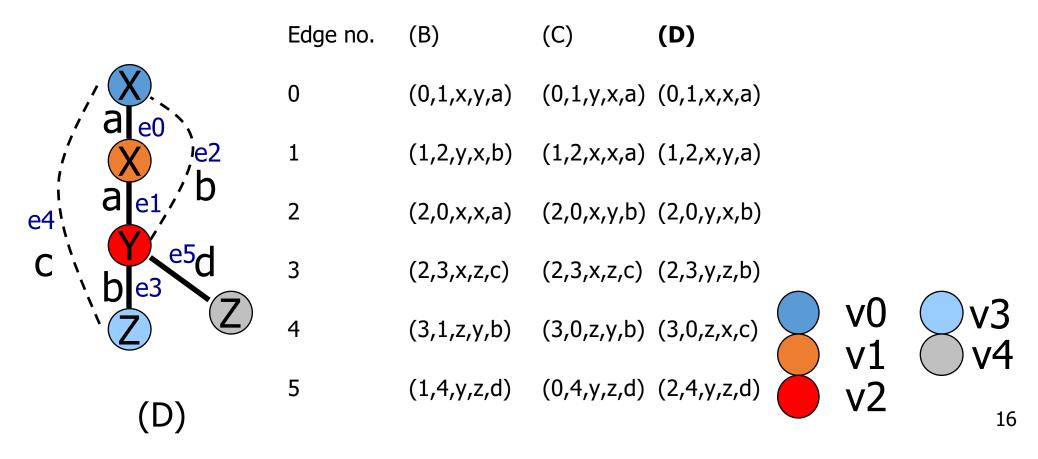
## Minimum DFS code

# Each Graph may have lots of DFS code (why?): one smallest lexicographic one is its Minimum DFS Code



#### Minimum DFS code

# Each Graph may have lots of DFS code (why?): one smallest lexicographic one is its Minimum DFS Code



# Minimum DFS Code

### DFS code:

List of tuples

# Lexicographic order:

Compare tuples in sequence alphabetically

C > B > D

Minimum DFS: D

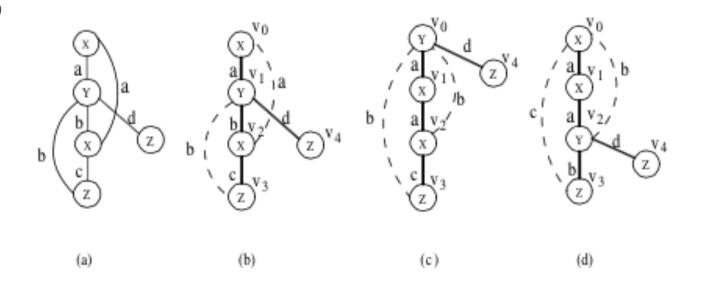


Figure 1. Depth-First Search Tree

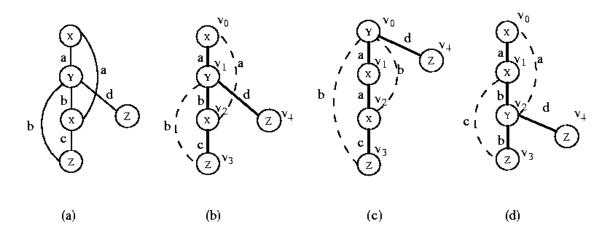
edge	(Fig 1b) α	(Fig 1c) β	(Fig 1d) γ
0	(0, 1, X, a, Y)	(0, 1, Y, a, X)	(0, 1, X, a, X)
1	(1, 2, Y, b, X)	(1, 2, X, a, X)	(1, 2, X, a, Y)
2	(2,0,X,a,X)	(2, 0, X, b, Y)	(2, 0, Y, b, X)
3	(2, 3, X, c, Z)	(2, 3, X, c, Z)	(2, 3, Y, b, Z)
4	(3, 1, Z, b, Y)	(3, 0, Z, b, Y)	(3, 0, Z, c, X)
5	(1, 4, Y, d, Z)	(0, 4, Y, d, Z)	(2, 4, Y, d, Z)

Table 1. DFS codes for Fig. 1(b)-(d)

### How to Generate Minimum DFS code

- Given node labels ordered (e.g., alphabetically)
- Selecting the minimum DFS code
  - select the nodes with the minimum label as the candidate roots
  - construct the DFS spanning tree from each root, always visit edges with smaller labels first; and if two edges with same label, visit the one whose second node has the smaller label
  - insert the backward edges

# Forward/Backward Edge Set



Forward Edge and Backward Edge. Given  $G_T$ , the forward edge (tree edge [3]) set contains all the edges in the DFS tree, and the backward edge (back edge [3]) set contains all the edges which are not in the DFS tree. For simplicity, (i, j) is an ordered pair to represent an edge. If i < j, it is a forward edge; otherwise, a backward edge. A linear order,  $\prec_T$  is built among all the edges in G by the following rules (assume  $e_1 = (i_1, j_1), e_2 = (i_2, j_2)$ ): (i) if  $i_1 = i_2$  and  $j_1 < j_2$ ,  $e_1 \prec_T e_2$ ; (ii) if  $i_1 < j_1$  and  $j_1 = i_2$ ,  $e_1 \prec_T e_2$ ; and (iii) if  $e_1 \prec_T e_2$  and  $e_2 \prec_T e_3$ ,  $e_1 \prec_T e_3$ .

min\_sup = 2 -> Patterns?

Only number of transaction matters, Not occurrences

Alternatively,

$$C \stackrel{a}{-} S \stackrel{a}{-} C \stackrel{a}{-} C$$

$$|b|$$

$$N \stackrel{a}{-} O$$

$$g3$$

min\_sup = 2 -> Patterns?

S a C Sup: 3
 C a C Sup: 3

### Alternatively,

$$C \xrightarrow{a} S \xrightarrow{a} C \xrightarrow{a} C$$

$$N \xrightarrow{a} O$$

$$g3$$

$$S \stackrel{a}{-} C$$

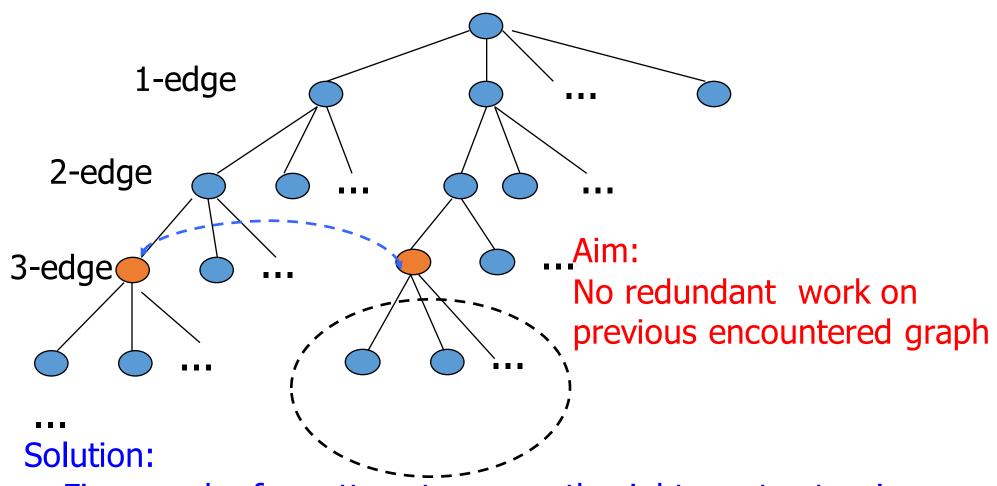
$$S \stackrel{a}{\longrightarrow} C$$
 Sup: 3  $\xrightarrow{Pattern growth}$   $S \stackrel{a}{\longrightarrow} C \stackrel{a}{\longrightarrow} C$  Sup: 3

min\_sup = 2 -> Patterns?

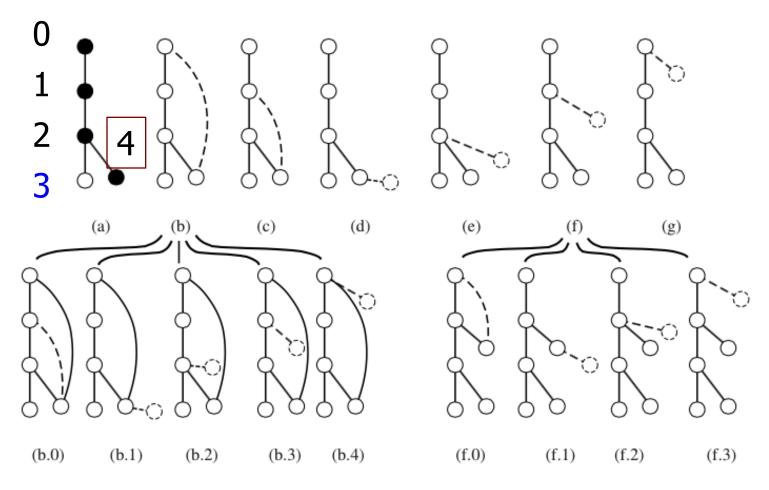
- Generate candidate growth patterns
- Remove duplication and find frequent ones
- Repeat

generating the same pattern Multiple times is inefficient

## Rightmost extension

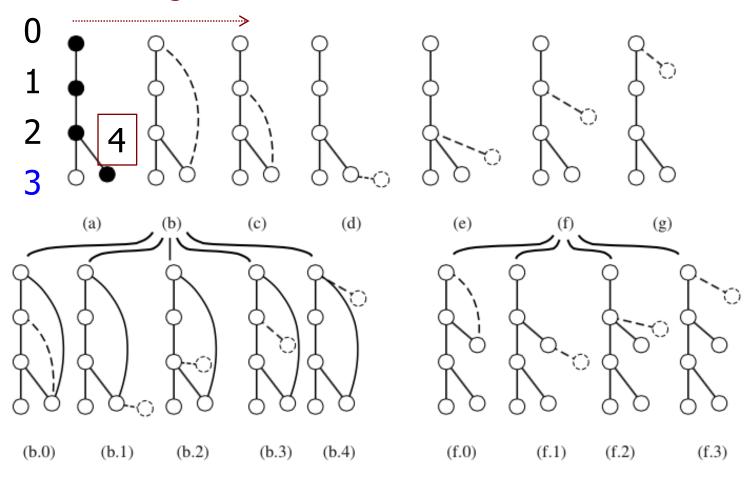


- Fix an order for pattern tree growth: right-most extension
- Test identical patterns: minimum DFS code



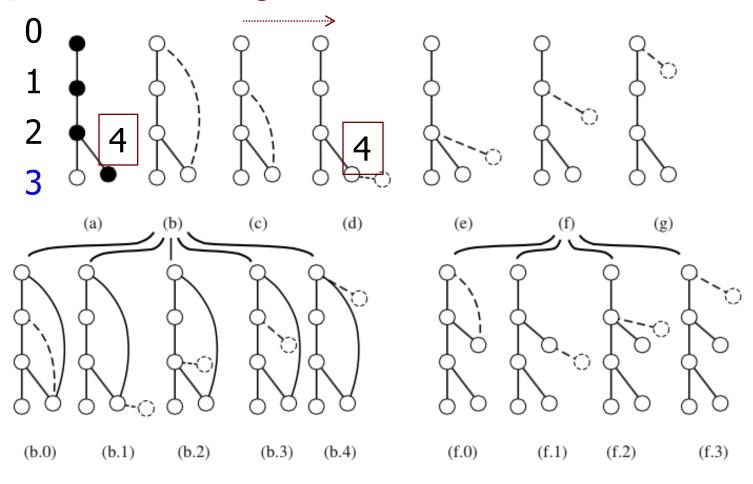
Rightmost Path: 0, 1, 2, 4

First, link backward edges from current RM node to old nodes in order of nodes



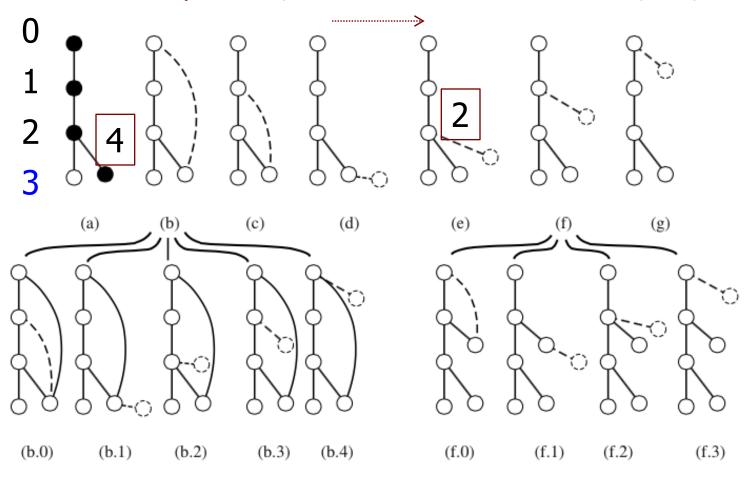
Rightmost Path: 0, 1, 2, 4

Then, extend forward edge from RM node



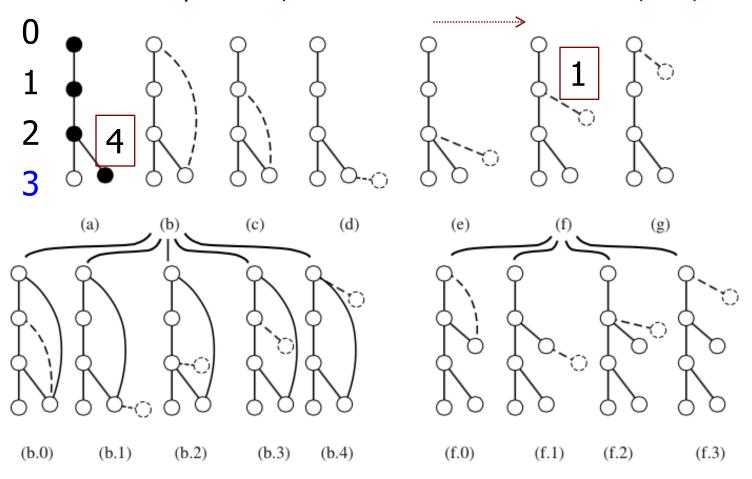
Rightmost Path: 0, 1, 2, 4

If not possible, use next available RM node, i.e., node 2



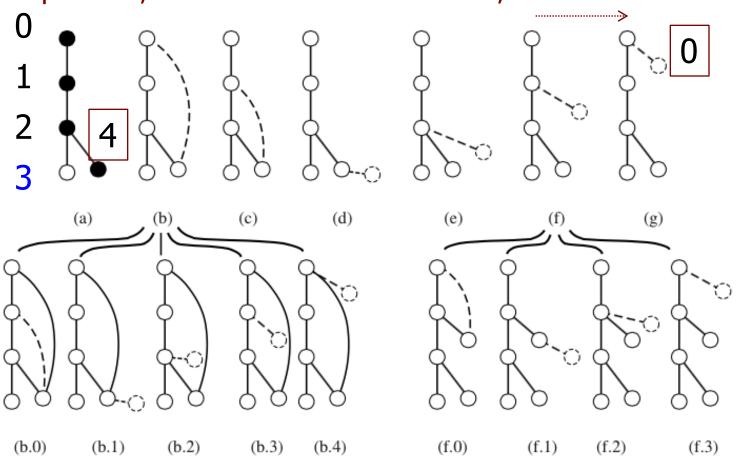
Rightmost Path: 0, 1, 2

If not possible, use next available RM node, i.e., node 1



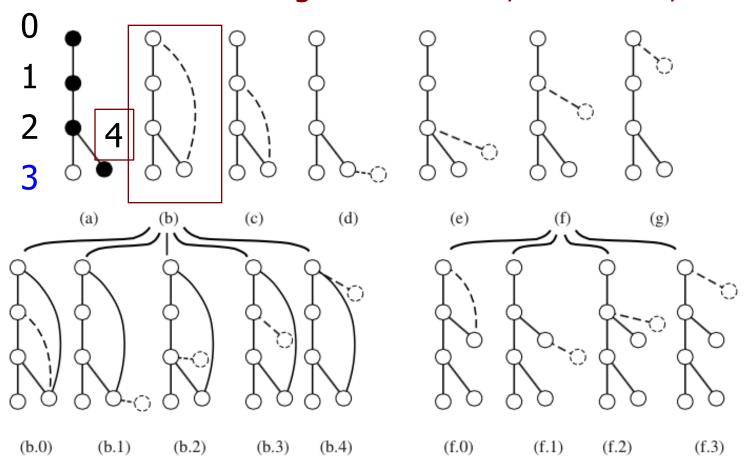
Rightmost Path: 0, 1

If not possible, use next available RM node, until extension terminates



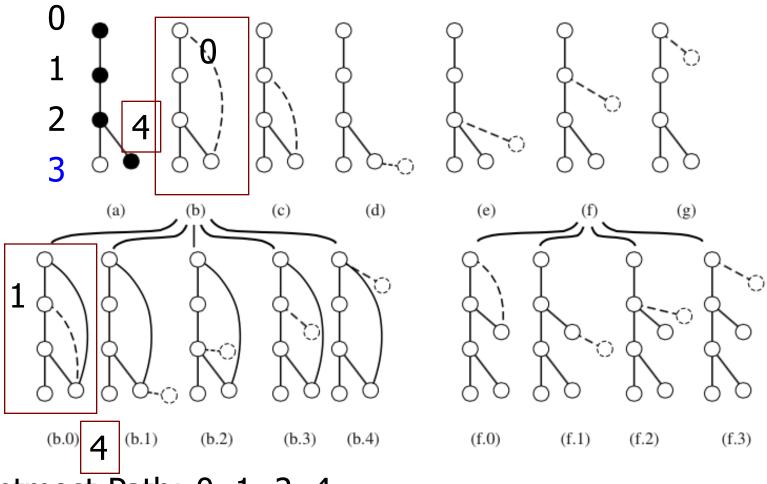
Rightmost Path: 0

If backward edge 4->0 exists, add 4->0, and repeat



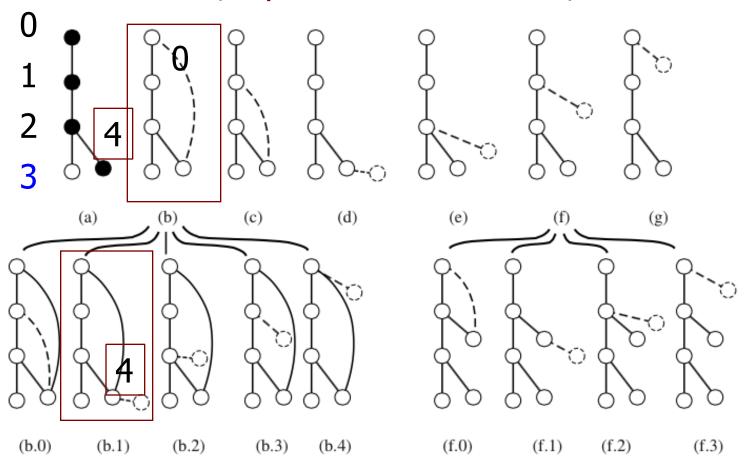
Rightmost Path: 0, 1, 2, 4

First, if backward 4->1 exists, link it



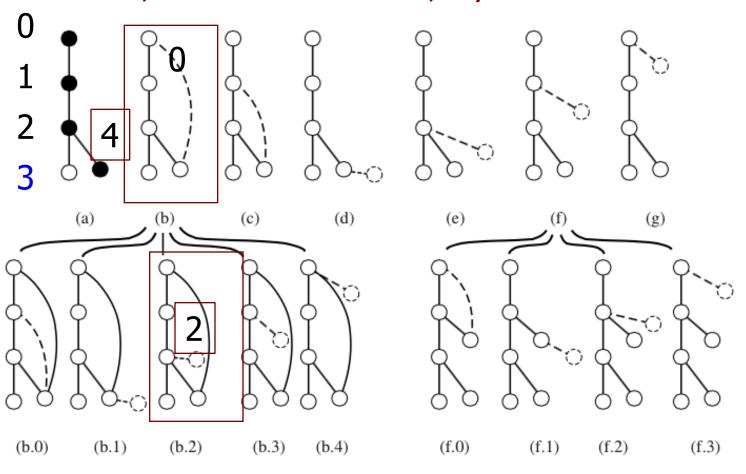
Rightmost Path: 0, 1, 2, 4

Otherwise, try extend RM Node 4,



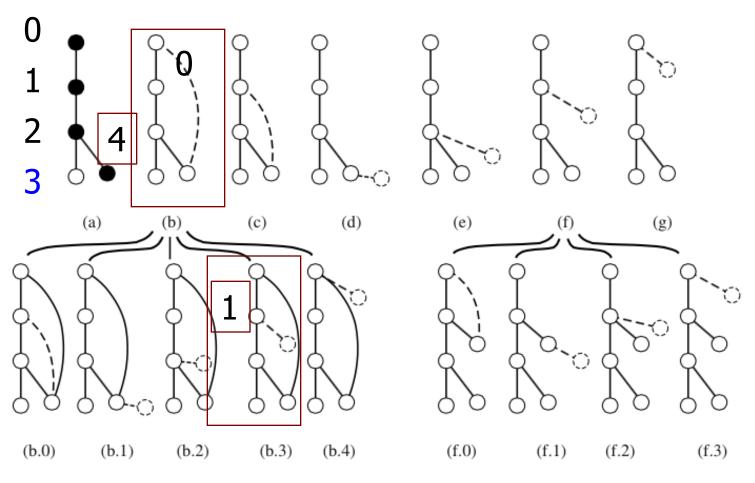
Rightmost Path: 0, 1, 2, 4

Then, use 2 as RM node, try forward extension

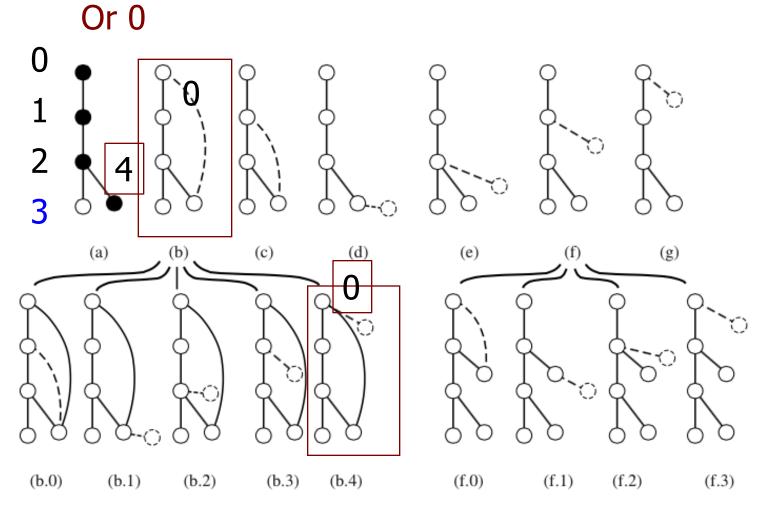


Rightmost Path: 0, 1, 2

#### Or 1 as RM Node



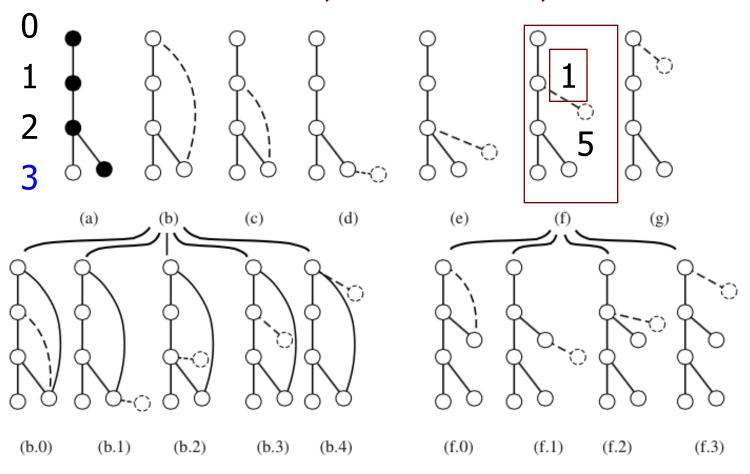
Rightmost Path: 0, 1



Rightmost Path: 0 Rightmost Node: 0

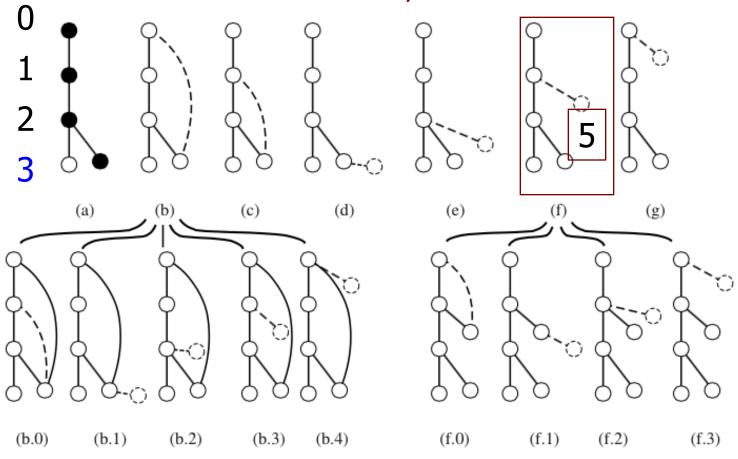
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Another situation, RM node is 1, forward 1->5 exists



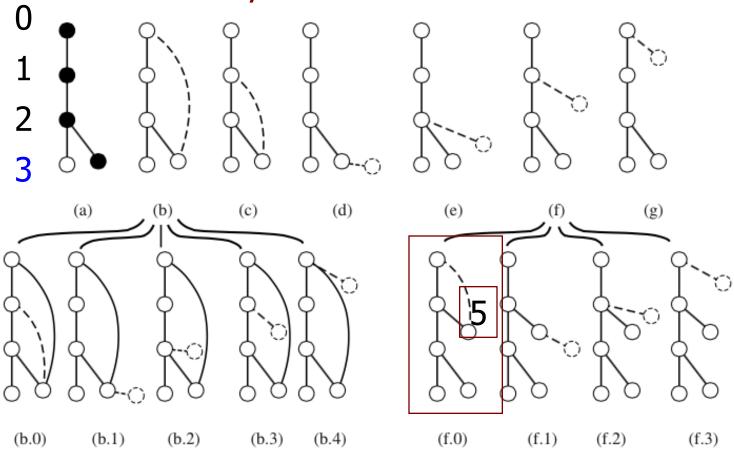
Rightmost Path: 0, 1

#### Extend 1->5, 5 is now RM node



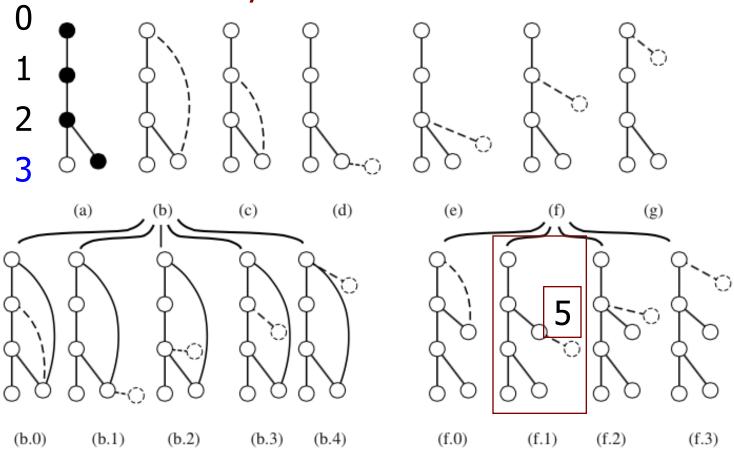
Rightmost Path: 0, 1, 5

#### Try backward 5->0



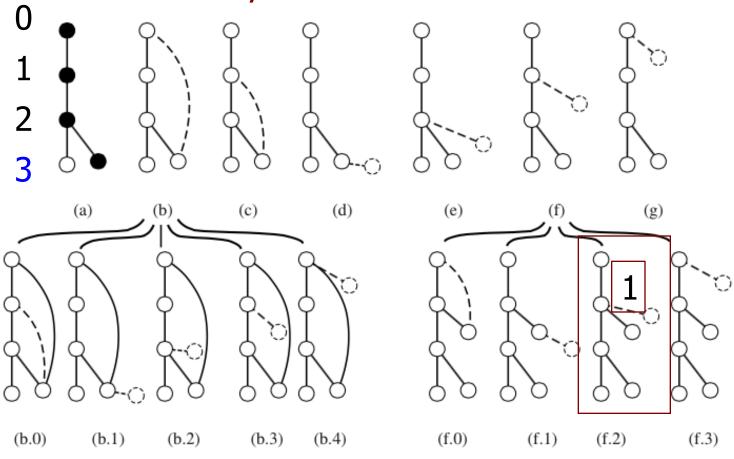
Rightmost Path: 0, 1, 5

#### Try extend RM node 5



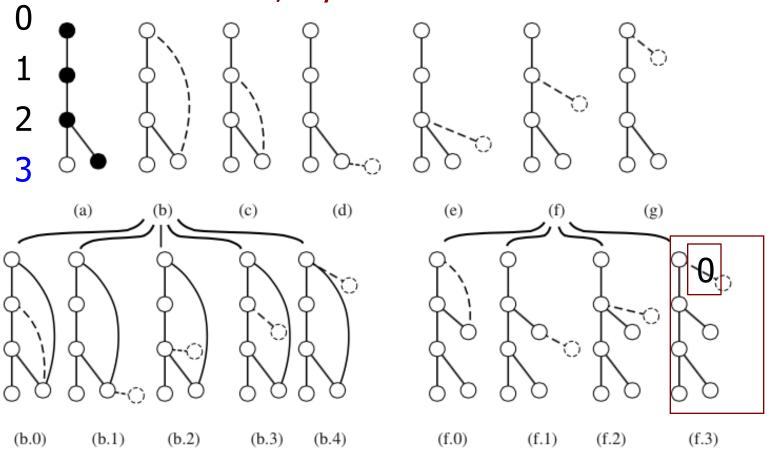
Rightmost Path: 0, 1, 5

#### Try extend RM node 1



Rightmost Path: 0, 1

#### Last, try extend RM node 0



Rightmost Path: 0

Frequency:

C: 3, N: 3, O: 3, S: 3, H: 1

Grow from C, N, O, S

#### Frequency:

C: 3, N: 3, O: 3, S: 3, H: 1

Grow from C (extension)

 $C \stackrel{a}{-} C$  sup: 3

 $C \stackrel{a}{\longrightarrow} N$  sup: 2

c <u>b</u> o sup: 2

 $C \stackrel{a}{\longrightarrow} S$  sup: 3

$$S \stackrel{H}{=} C \stackrel{O}{=} V \stackrel{|b}{=} V \stackrel{|c}{=} V \stackrel{|b}{=} V \stackrel{|c}{=} V \stackrel{|c}{=}$$

Frequency:

C: 3, N: 3, O: 3, S: 3, H: 1

Grow from N, S similarly (extension)

$$S \stackrel{H}{=} C \stackrel{O}{=} C \stackrel{|b}{=} N \qquad C \stackrel{|b}{=} C \stackrel{|b}{=} N \stackrel{a}{=} C \qquad C \stackrel{a}{=} C \stackrel{a}{=} C \qquad C \stackrel{a}{=} C \stackrel{b}{=} C \qquad C \stackrel{b}{=} S \stackrel{a}{=} C \stackrel{a}{=} C \qquad C \stackrel{b}{=} S \stackrel{a}{=} C \stackrel{a}{=} C \qquad C \stackrel{b}{=} S \stackrel{a}{=} C \stackrel{a}{=} C \stackrel{a}{=} C \stackrel{b}{=} C \stackrel{b}{=} C \stackrel{a}{=} C \stackrel{a}{$$

Frequency:

C: 3, N: 3, O: 3, S: 3, H: 1

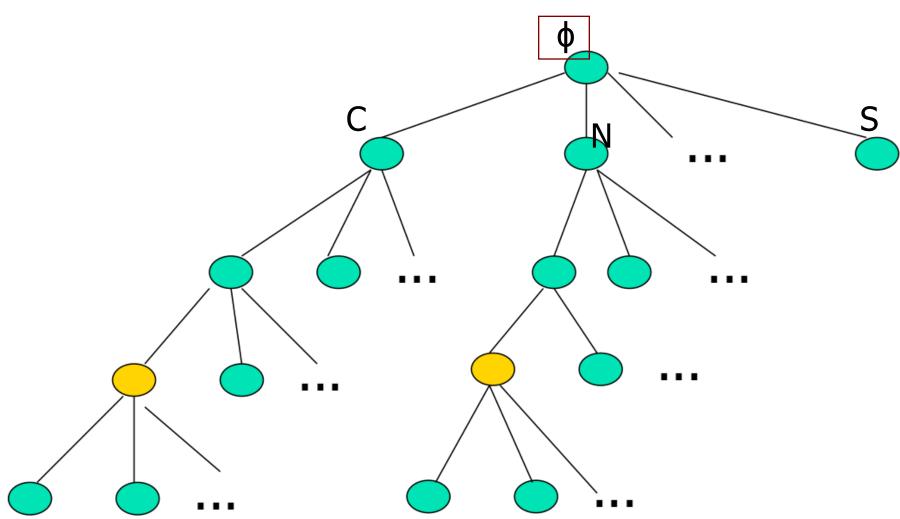
Grow from  $C \stackrel{a}{\longrightarrow} C$  (extension)

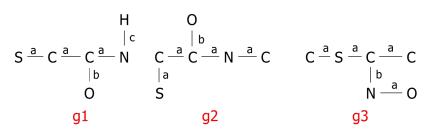
Frequency:

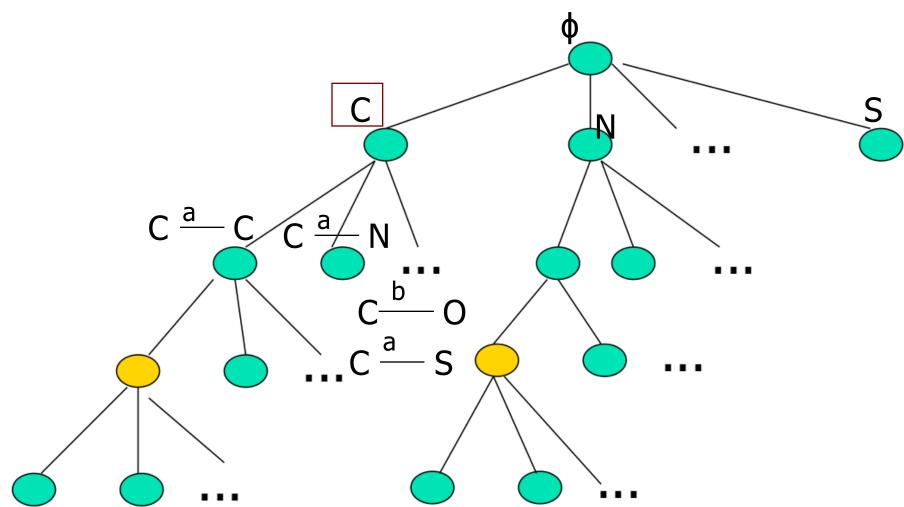
C: 3, N: 3, O: 3, S: 3, H: 1

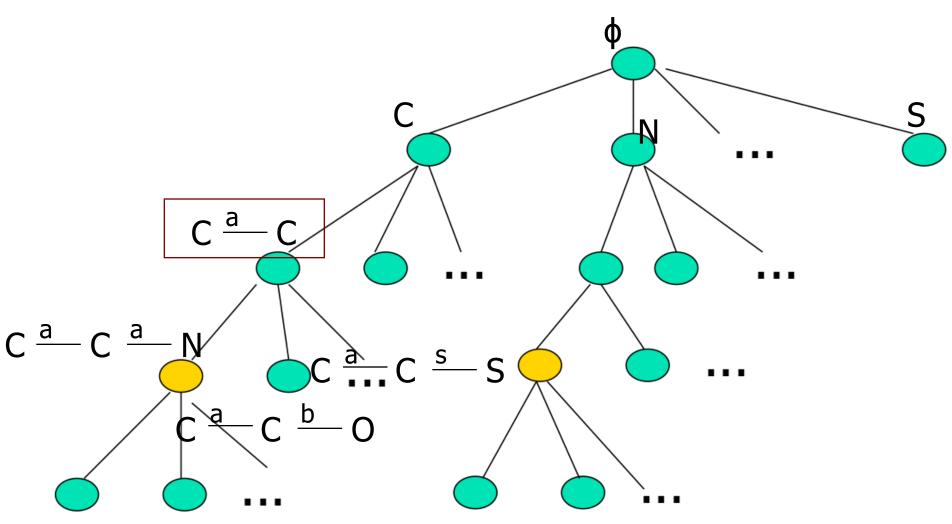
Grow from  $C \stackrel{a}{\longrightarrow} N$  (extension)

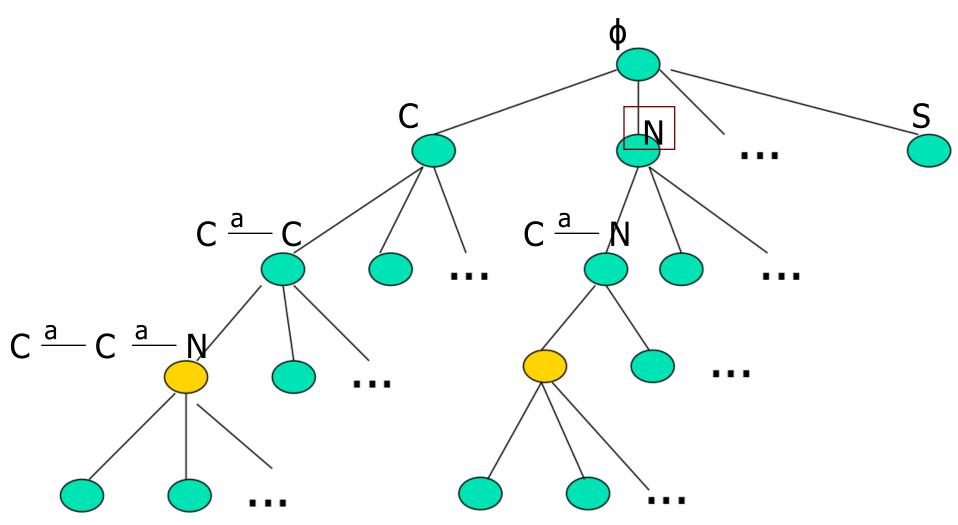
Once  $C \stackrel{a}{=} C \stackrel{a}{=} N$  sup: 2 is encountered, by minimum DFS code, it is identified redundant pattern, and thus removed.

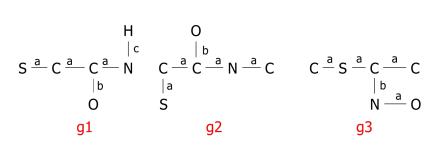


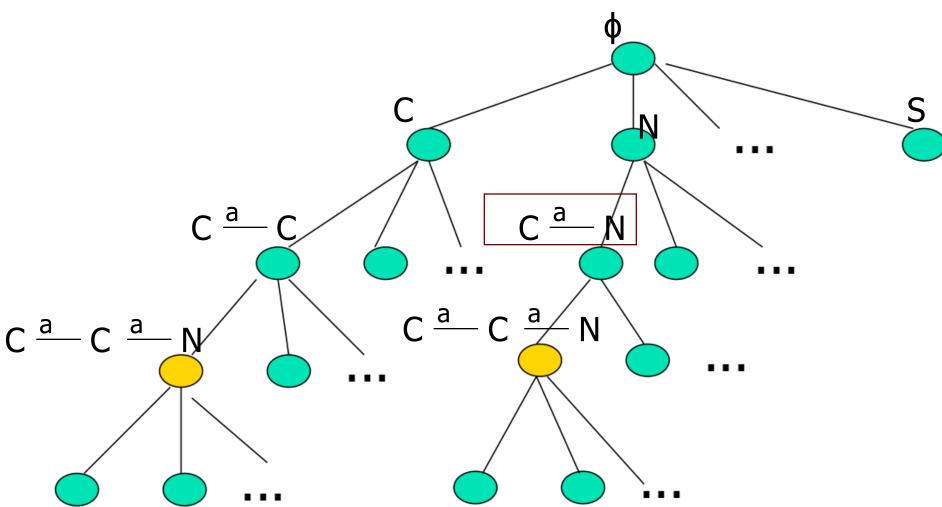


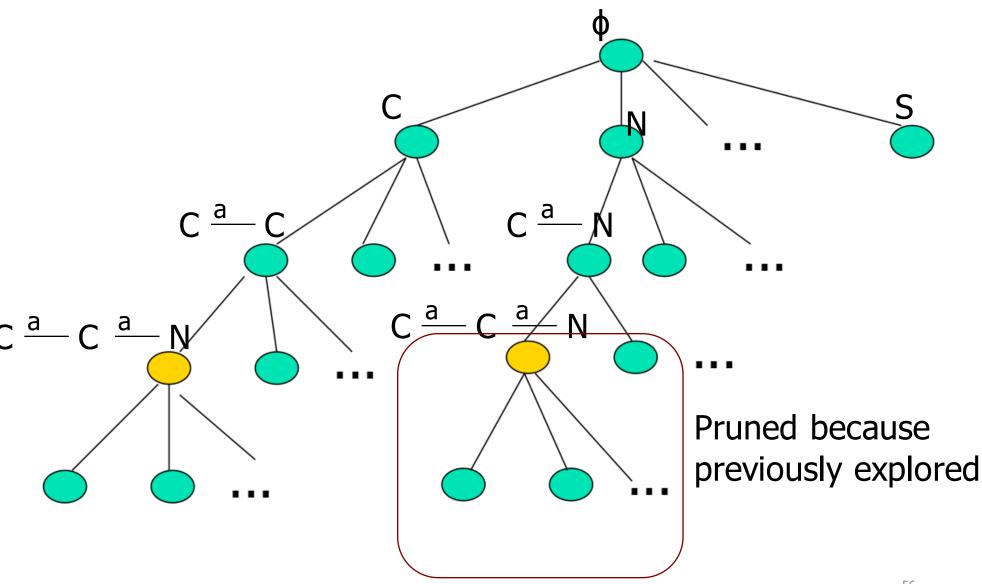




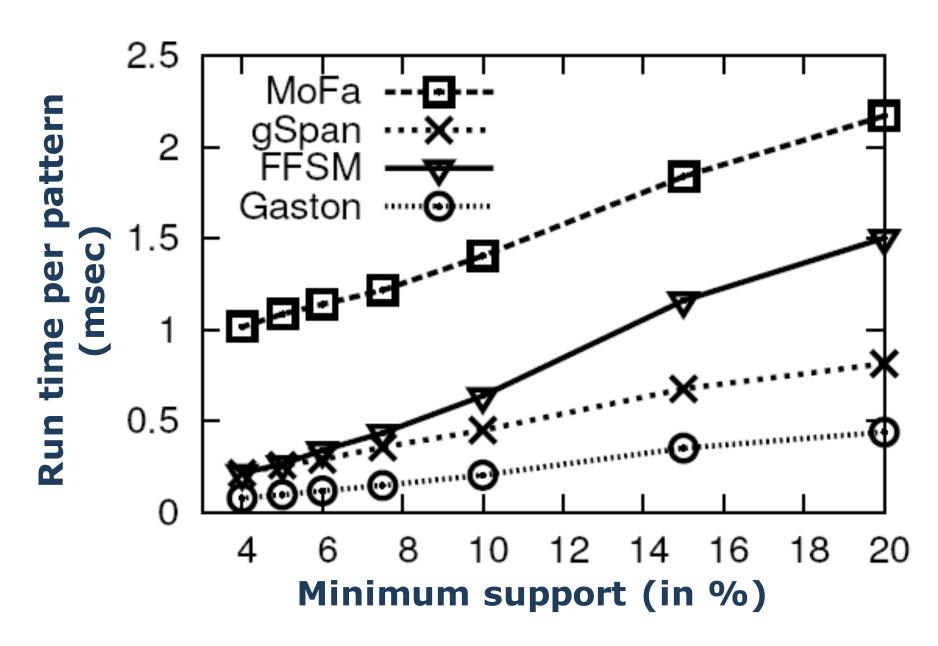




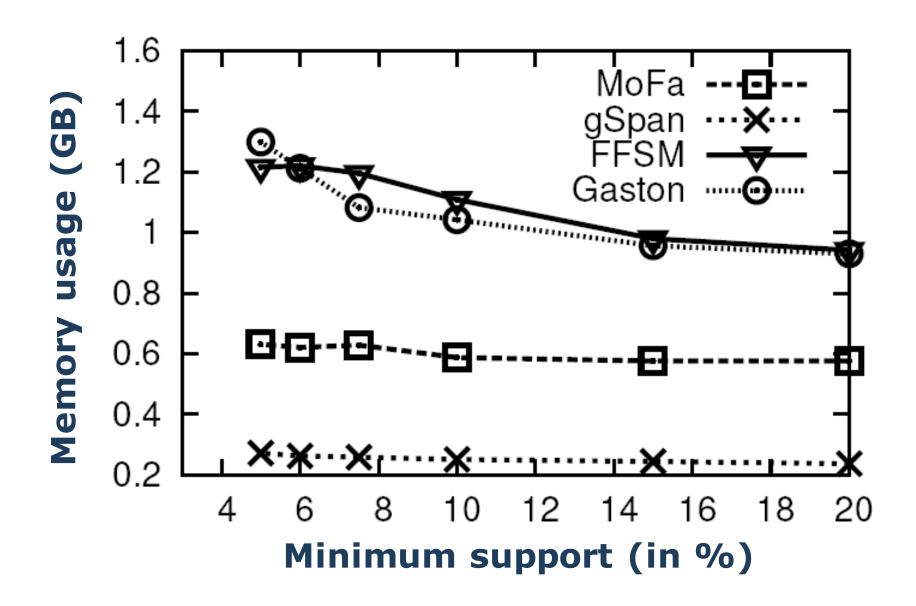




### Performance (1): Run Time



### Performance (2): Memory Usage



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## Summary

- Roadmap: Many aspects & extensions on pattern mining
- Mining patterns in multi-level, multi dimensional space
- Mining rare and negative patterns
- Constraint-based pattern mining
- Specialized methods for mining high-dimensional data and colossal patterns
- Mining compressed or approximate patterns
- Pattern exploration and understanding: Semantic annotation of frequent patterns