

# Stopping Rule for Secretary Problem

# Team Members

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# Problem Description

- Interview  $n$  candidates for a position
- Cannot recall
- After each can decide their rank
- Don't know the quality of rest

Goal: maximize the probability of choosing the best candidate.

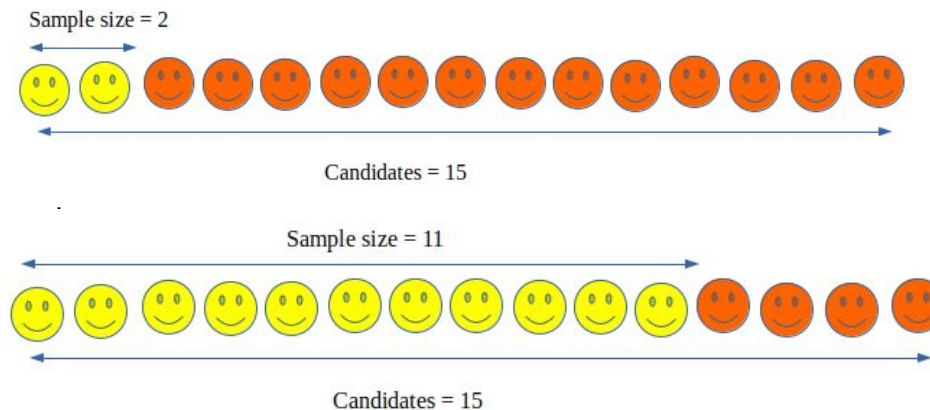


# Importance

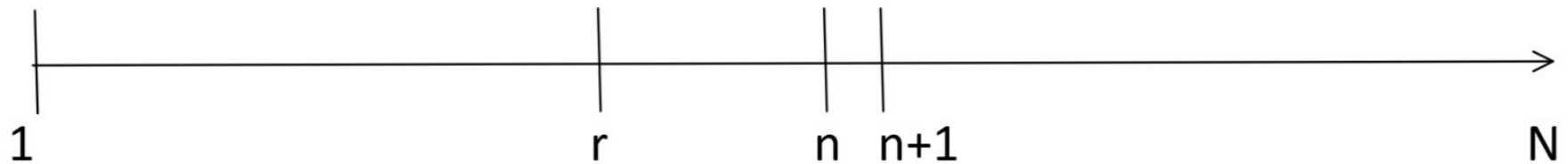
- A frequent problem companies face is how to choose the best possible job candidate
- The secretary problem ensures the best possible candidate can be chosen by forcing the hiring manager to wait
  - Can be used on any problem when making a choice among a group of options when you can't change your mind
  - Used by Mathematician Johannes Kepler in 1613 to successfully find his wife

# Difficulty

- Need an effective way of choosing a control group
  - Taking the first candidate will rarely result in the best
  - Need an effective way of finding the optimal control group



# Solution

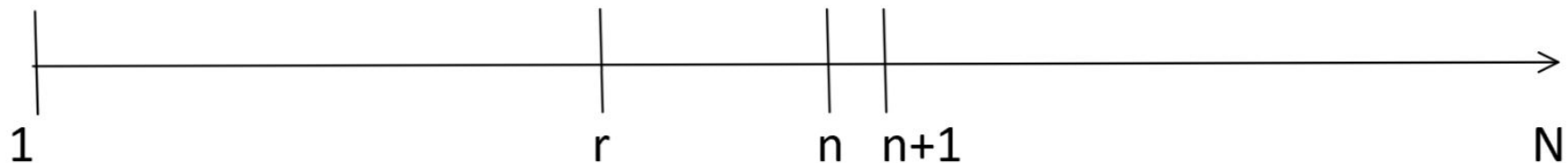


Interview and reject the first  $r$  applicants, for  $r < N$ .

Accept the very next applicant that is better than all the first  $r$  you interviewed.

$$P(\text{Success}) = P(r)$$

# Solution



$i^*$  (the best) occur at  $n+1$ . Do not pick it unless:

- (1)  $n \geq r$
- (2) The highest applicant in  $[1, n]$  is the same highest applicant in  $[1, r]$

The possibility of this is:  $\frac{r}{n} \frac{1}{N}$

The probability of  $i^*$  occurring at  $n + 1$  is  $1/N$

The probability of condition (2) is  $r/n$

# Solution

$$P(r) = \frac{1}{N} \left[ \frac{r}{r} + \frac{r}{r+1} + \frac{r}{r+2} + \dots + \frac{r}{N-1} \right] = \frac{r}{N} \sum_{n=r}^{N-1} \frac{1}{n}$$

$P(r)$  is a Riemann approximation to an integral.

$$P(r) = \lim_{N \rightarrow \infty} \frac{r}{N} \sum_{n=r}^{N-1} \frac{N}{n} \frac{1}{N} = x \int_x^1 \frac{1}{t} dt = -x \ln x$$

Solving  $P'(r) = 0$  for  $r$  gives us the optimal ratio and the probability of success  $P(r_{\text{optimal}})$ .



# Solution

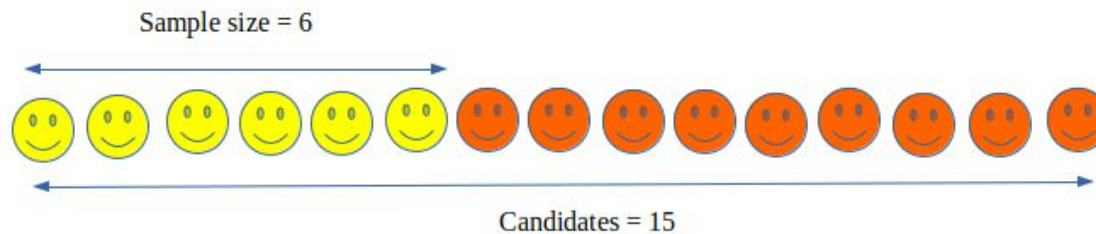
$$P'(r) = -\ln x - 1 = 0 \implies x = \frac{1}{e}$$

$$P\left(\frac{1}{e}\right) = \frac{1}{e}$$

$$\frac{1}{e} \approx .37$$

# Solution

- The solution shows that an optimal control group makes up a size of  $1/e$  or approximately 37%
- The first candidate better than the control group should be chosen
  - Results in the best probability of choosing the top candidate



# Time Complexity

- $O(n)$ 
  - In the worst case will look at  $n$  candidates choosing the last candidate
    - This will occur when the best candidate was found in the first  $1/e$  interviewed
- In most cases the best candidate will be found much earlier resulting in the code finishing in less than  $n$  iterations

# Experimental Result

Number of candidates	Probability of success	Running time (sec/10000 times)
20	0.376	0.550550891
50	0.379	1.140982333
100	0.3786	2.301025147
200	0.3629	3.678622066
500	0.3754	11.11833581
1000	0.3763	20.32150593
2000	0.364	41.23882053
5000	0.3774	104.4833268
10000	0.3647	214.7531664

# Conclusion

- The secretary problem shows that reviewing  $n/e$  candidates as part of a control group will maximize the chances of selecting the best
  - Does not guarantee the best is chosen
  - Reminds us to avoid rushing into a selection
- Can be used on any problem when a choice must be made without knowing all details.

# Citations

<https://www.geeksforgeeks.org/secretary-problem-optimal-stopping-problem/>

<https://thebryanhernandezgame.files.wordpress.com/2010/05/secretary-problem.pdf>

# Github

<https://github.com/RoyZhang7/secretary-problem>