



SMART CONTRACT AUDIT REPORT

for

ROYALE FINANCE



Prepared By: Shuxiao Wang

PeckShield
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Contact

For more information about this document and its contents, please contact PeckShield Inc.

Name	Shuxiao Wang
Phone	+86 173 6454 5338
Email	contact@peckshield.com

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1 | Introduction

Given the opportunity to review the design document and related smart contract source code of the **Royale Finance** protocol, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

1.1 About Royale Finance

Royale Finance is a decentralized protocol for directing optimized stablecoin liquidity pools toward an on-chain funding solution for iGaming products. By providing a decentralized liquidity network, Royale aims to lower the barriers of entry for the iGaming industry and foster a robust crypto-economic system that incentivizes both iGaming entrepreneurs and liquidity providers to become long-term ecosystem participants. It is argued that the countercyclical nature of the Royale Finance ecosystem should make it an attractive venue for the contribution of decentralized stablecoin liquidity.

The basic information of Royale Finance is as follows:

Table 1.1: Basic Information of Royale Finance

Item	Description
Issuer	Royale Finance
Website	https://royale-finance-docs.netlify.app/
Type	Ethereum Smart Contract
Platform	Solidity
Audit Method	Whitebox
Latest Audit Report	April 1, 2021

In the following, we show the Git repository of reviewed files and the commit hash value used in

this audit.

- <https://github.com/RoyaleFinanceV1/royale-contracts.git> (f357b1f)

And here is the commit ID after all fixes for the issues found in the audit have been checked in:

- <https://github.com/RoyaleFinanceV1/royale-contracts.git> (4c648f6)

1.2 About PeckShield

PeckShield Inc. [13] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (<https://t.me/peckshield>), Twitter (<http://twitter.com/peckshield>), or Email (contact@peckshield.com).

Table 1.2: Vulnerability Severity Classification

Impact	High	Critical	High	Medium
	Medium	High	Medium	Low
	Low	Medium	Low	Low
		High	Medium	Low
		Likelihood		

1.3 Methodology

To standardize the evaluation, we define the following terminology based on the OWASP Risk Rating Methodology [12]:

- Likelihood represents how likely a particular vulnerability is to be uncovered and exploited in the wild;
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Table 1.3: The Full Audit Checklist

Category	Checklist Items
Basic Coding Bugs	Constructor Mismatch
	Ownership Takeover
	Redundant Fallback Function
	Overflows & Underflows
	Reentrancy
	Money-Giving Bug
	Blackhole
	Unauthorized Self-Destruct
	Revert DoS
	Unchecked External Call
	Gasless Send
	Send Instead Of Transfer
	Costly Loop
	(Unsafe) Use Of Untrusted Libraries
	(Unsafe) Use Of Predictable Variables
	Transaction Ordering Dependence
	Deprecated Uses
Semantic Consistency Checks	Semantic Consistency Checks
Advanced DeFi Scrutiny	Business Logics Review
	Functionality Checks
	Authentication Management
	Access Control & Authorization
	Oracle Security
	Digital Asset Escrow
	Kill-Switch Mechanism
	Operation Trails & Event Generation
	ERC20 Idiosyncrasies Handling
	Frontend-Contract Integration
	Deployment Consistency
	Holistic Risk Management
Additional Recommendations	Avoiding Use of Variadic Byte Array
	Using Fixed Compiler Version
	Making Visibility Level Explicit
	Making Type Inference Explicit
	Adhering To Function Declaration Strictly
	Following Other Best Practices

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

To evaluate the risk, we go through a checklist of items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- Basic Coding Bugs: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- Semantic Consistency Checks: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [11], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings. Moreover, in case there is an issue that may affect an active protocol that has been deployed, the public version of this report may omit such issue, but will be amended with full details right after the affected protocol is upgraded with respective fixes.

1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary
Configuration	Weaknesses in this category are typically introduced during the configuration of the software.
Data Processing Issues	Weaknesses in this category are typically found in functionality that processes data.
Numeric Errors	Weaknesses in this category are related to improper calculation or conversion of numbers.
Security Features	Weaknesses in this category are concerned with topics like authentication, access control, confidentiality, cryptography, and privilege management. (Software security is not security software.)
Time and State	Weaknesses in this category are related to the improper management of time and state in an environment that supports simultaneous or near-simultaneous computation by multiple systems, processes, or threads.
Error Conditions, Return Values, Status Codes	Weaknesses in this category include weaknesses that occur if a function does not generate the correct return/status code, or if the application does not handle all possible return/status codes that could be generated by a function.
Resource Management	Weaknesses in this category are related to improper management of system resources.
Behavioral Issues	Weaknesses in this category are related to unexpected behaviors from code that an application uses.
Business Logic	Weaknesses in this category identify some of the underlying problems that commonly allow attackers to manipulate the business logic of an application. Errors in business logic can be devastating to an entire application.
Initialization and Cleanup	Weaknesses in this category occur in behaviors that are used for initialization and breakdown.
Arguments and Parameters	Weaknesses in this category are related to improper use of arguments or parameters within function calls.
Expression Issues	Weaknesses in this category are related to incorrectly written expressions within code.
Coding Practices	Weaknesses in this category are related to coding practices that are deemed unsafe and increase the chances that an exploitable vulnerability will be present in the application. They may not directly introduce a vulnerability, but indicate the product has not been carefully developed or maintained.

comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.



2 | Findings

2.1 Summary

Here is a summary of our findings after analyzing the implementation of the `Royale Finance` protocol. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logic, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings	
Critical	0	
High	1	■
Medium	4	■ ■ ■ ■
Low	4	■ ■ ■ ■
Informational	1	■
Total	10	

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in [Section 3](#).

2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 1 high-severity vulnerability, 4 medium-severity vulnerabilities, 4 low-severity vulnerabilities, and 1 informational recommendation.

Table 2.1: Key Royale Finance Audit Findings

ID	Severity	Title	Category	Status
PVE-001	Low	Safe-Version Replacement With <code>safeApprove()</code> , <code>safeTransfer()</code> And <code>safeTransferFrom()</code>	Coding Practices	Fixed
PVE-002	Low	Potential OOG Execution of <code>RoyaleLP::updateWithdrawQueue()</code>	Coding Practices	Fixed
PVE-003	Informational	Two-Step Transfer Of Privileged Account Ownership	Security Features	Fixed
PVE-004	Medium	Potential Underflow For Unlimited Discounted Amounts	Numeric Errors	Fixed
PVE-005	Low	Simplification of <code>RoyaleLP::requestWithdrawWithRPT()</code>	Coding Practices	Fixed
PVE-006	Medium	Non-Permissioned <code>ERC1155Receiver::setShouldReject()</code>	Security Features	Fixed
PVE-007	Low	Possible Friction That May Block User Withdrawals	Business Logic	Fixed
PVE-008	Medium	Trust Issue of Admin Keys	Security Features	Mitigated
PVE-009	Medium	Possible <code>CooldownTimestamp</code> Manipulation	Business Logic	Fixed
PVE-010	High	Possible Sandwich/MEV For Reduced Return And Skewed Withdrawals	Time And State	Fixed

Beside the identified issues, we emphasize that for any user-facing applications and services, it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms should kick in at the very moment when the contracts are being deployed on mainnet. Please refer to Section 3 for details.

3 | Detailed Results

3.1 Safe-Version Replacement With `safeApprove()`, `safeTransfer()` And `safeTransferFrom()`

- ID: PVE-001
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: `rStrategy`
- Category: Coding Practices [8]
- CWE subcategory: CWE-1126 [2]

Description

Though there is a standardized ERC-20 specification, many token contracts may not strictly follow the specification or have additional functionalities beyond the specification. In this section, we examine the `approve()` routine and possible idiosyncrasies from current widely-used token contracts.

In particular, we use the popular stablecoin, i.e., USDT, as our example. We show the related code snippet below. On its entry of `approve()`, there is a requirement, i.e., `require(!((_value != 0) && (allowed[msg.sender][_spender] != 0)))`. This specific requirement essentially indicates the need of reducing the allowance to 0 first (by calling `approve(_spender, 0)`) if it is not, and then calling a second one to set the proper allowance. This requirement is in place to mitigate the known `approve()/transferFrom()` race condition (<https://github.com/ethereum/EIPs/issues/20#issuecomment-263524729>).

```

194  /**
195  * @dev Approve the passed address to spend the specified amount of tokens on behalf
      of msg.sender.
196  * @param _spender The address which will spend the funds.
197  * @param _value The amount of tokens to be spent.
198  */
199  function approve(address _spender, uint _value) public onlyPayloadSize(2 * 32) {

201      // To change the approve amount you first have to reduce the addresses'
202      // allowance to zero by calling 'approve(_spender, 0)' if it is not
203      // already 0 to mitigate the race condition described here:

```

```

204 // https://github.com/ethereum/EIPs/issues/20#issuecomment-263524729
205 require(!((_value != 0) && (allowed[msg.sender][_spender] != 0)));

207 allowed[msg.sender][_spender] = _value;
208 Approval(msg.sender, _spender, _value);
209 }

```

Listing 3.1: USDT Token Contract

Because of that, a normal call to `approve()` with a currently non-zero allowance may fail. In the following, we use the `rStrategy::deposit()` routine as an example. This routine is designed to trigger default handling. To accommodate the specific idiosyncrasy, there is a need to `approve()` twice (line 93): the first one reduces the allowance to 0; and the second one sets the new allowance.

```

100 // deposits stable tokens into the 3pool and stake received LPtoken(3CRV) in the
    curve 3pool gauge
101 function deposit(uint[3] memory amounts) external onlyRoyaleLP(){
102     uint currentTotal;
103     for(uint8 i=0; i<3; i++) {
104         if(amounts[i] > 0) {
105             uint decimal;
106             decimal=tokens[i].decimals();
107             tokens[i].approve(address(pool), amounts[i]);
108             currentTotal =currentTotal.add(amounts[i].mul(1e18).div(10**decimal));
109         }
110     }
111     /* uint256 returnedAmount;
112     bool status;
113     (returnedAmount,status)=calculateProfit();
114     if(status){
115         totalProfit =totalProfit.add(returnedAmount);
116     }
117     else{
118         totalProfit =totalProfit.sub(returnedAmount);
119     } */
120     uint256 mintAmount = currentTotal.mul(1e18).div(pool.get_virtual_price());
121     pool.add_liquidity(amounts, mintAmount.mul(DENOMINATOR.sub(DepositSlip)).div(
        DENOMINATOR));
122     //virtualPrice=pool.get_virtual_price();
123     stakeLP();
124 }

```

Listing 3.2: `rStrategy::deposit()`

Moreover, it is important to note that for certain non-compliant ERC20 tokens (e.g., USDT), the `transfer()` function does not have a return value. However, the `IERC20` interface has defined the `transfer()` interface with a `bool` return value. As a result, the call to `transfer()` may expect a return value. With the lack of return value of USDT's `transfer()`, the call will be unfortunately reverted.

Because of that, a normal call to `transfer()` is suggested to use the safe version, i.e., `safeTransfer()`. In essence, it is a wrapper around ERC20 operations that may either throw on failure or return

false without reverts. Moreover, the safe version also supports tokens that return no value (and instead revert or throw on failure). Note that non-reverting calls are assumed to be successful. To use this library you can add a using SafeERC20 for IERC20. Similarly, there is a safe version of approve()/transferFrom() as well, i.e., safeApprove()/safeTransferFrom().

Recommendation Accommodate the above-mentioned idiosyncrasy about ERC20-related approve()/transferFrom()/transferFrom().

Status The issue has been fixed by this commit: 4c648f6.

3.2 Potential OOG Execution of RoyaleLP::updateWithdrawQueue()

- ID: PVE-002
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: RoyaleLP
- Category: Coding Practices [8]
- CWE subcategory: CWE-1126 [2]

Description

The Royale Finance protocol has a core RoyaleLP contract that provides the main entry for borrowing/supplying users. In addition to normal lending functionalities, e.g., deposit()/withdraw() and borrow()/repay(), the protocol further implements a number of enhancements, including the lock period (before allowing deposited assets to be withdrawn) and passive yield generation. These enhancements are useful to improve the utilization or efficiency of assets in the lending pools.

To elaborate, we show below the updateWithdrawQueue() routine that is an essential helper to move withdrawing users to the reserveRecipients so that they can claim back the assets. While it properly implements the intended functionality, it comes to our attention this routine may process a long list of withdrawing user and runs into the risk of out-of-gas (OOG) execution. If the undesirable OOG occurs, the normal withdraw operation will be unnecessarily affected and blocked.

```

426 //this function is called when Royale Governance withdrawl from yield generation
    pool.It add all the withdrawl amount in the reserve amount.
427 //All the users who requested for the withdrawl are added to the reserveRecipients.
428 function updateWithdrawQueue() internal{
429     for(uint8 i=0;i<3;i++){
430         reserveAmount[i]=reserveAmount[i].add(totalWithdraw[i]);
431         totalWithdraw[i]=0;
432     }
433     for(uint i=0; i<withdrawRecipients.length; i++) {
434         reserveRecipients[withdrawRecipients[i]]=true;

```

```
435         isInQ [ withdrawRecipients [ i ] ] = false ;
436     }
437     uint count=withdrawRecipients.length ;
438     for ( uint i=0; i<count; i++) {
439         withdrawRecipients.pop () ;
440     }
441 }
```

Listing 3.3: RoyaleLP::updateWithdrawQueue()

Recommendation Timely monitor the length of `withdrawRecipients` and proactively call `withdraw()` to clean up the list.

Status The issue has been fixed by this commit: [4c648f6](#).

3.3 Two-Step Transfer Of Privileged Account Ownership

- ID: PVE-003
- Severity: Informational
- Likelihood: Low
- Impact: N/A
- Targets: Multiple Contracts
- Category: Security Features [7]
- CWE subcategory: CWE-282 [3]

Description

The Royale Finance protocol implements a rather basic access control mechanism that allows a privileged account, i.e., `wallet`, to be granted exclusive access to typically sensitive functions (e.g., `withdraw` all liquidity and `rebalance` the investment). Because of the privileged access and the implications of these sensitive functions, the `wallet` account is essential for the protocol-level safety and operation. In the following, we elaborate with the `wallet` account.

Within the governing contract `RoyaleLP`, a specific function, i.e., `transferOwnership()`, is provided to allow for possible `wallet/owner` updates. However, current implementation achieves its goal within a single transaction. This is reasonable under the assumption that the new `_wallet` parameter is always correctly provided. However, in the unlikely situation, when an incorrect `_wallet` is provided, the contract owner may be forever lost, which might be devastating for Royale Finance operation and maintenance.

As a common best practice, instead of achieving the owner update within a single transaction, it is suggested to split the operation into two steps. The first step initiates the owner update intent and the second step accepts and materializes the update. Both steps should be executed in two separate transactions. By doing so, it can greatly alleviate the concern of accidentally transferring the contract owner to an uncontrolled address. In other words, this two-step procedure ensures that

a owner public key cannot be nominated unless there is an entity that has the corresponding private key. This is explicitly designed to prevent unintentional errors in the owner transfer process.

```

38     function transferOwnership(address _wallet) external onlyWallet(){
39         wallet = _wallet;
40     }

```

Listing 3.4: RoyaleLP::transferOwnership()

Recommendation Implement a two-step approach for owner update (or transfer): `transferOwnership()` and `acceptOwnership()`.

Status The issue has been fixed by this commit: `4c648f6`.

3.4 Potential Underflow For Unlimited Discounted Amounts

- ID: PVE-004
- Severity: Medium
- Likelihood: Low
- Impact: High
- Target: Multiple Contracts
- Category: Security Features [7]
- CWE subcategory: CWE-282 [3]

Description

The Royale Finance protocol has developed three different staking lots, i.e., `StakingLotKing`, `StakingLotQueen`, and `StakingLotRoyaleFlush`. These staking lots mainly differ in their lot prices. In the following, we notice the staking lot implementation lacks proper `SafeMath` protection that may result in unintended underflow.

To illustrate, we show below the `_beforeTokenTransfer()` helper routine. This helper routine is internally used to facilitate the staking lot pool token transfer from one account to another. Specifically, it is called in every single `transfer()` operation.

```

1122     function _beforeTokenTransfer(address from, address to, uint256 amount) internal
1123         override {
1124         if (from != address(0)) saveProfit(from);
1125         if (to != address(0)) saveProfit(to);
1126         if (
1127             lastBoughtTimestamp[from].add(lockupPeriod) > block.timestamp &&
1128             lastBoughtTimestamp[from] > lastBoughtTimestamp[to]
1129         ) {
1130             require(
1131                 !_revertTransfersInLockUpPeriod[to],
1132                 "the recipient does not accept blocked funds"
1133             );
1134             lastBoughtTimestamp[to] = lastBoughtTimestamp[from];

```



```

1134     }
1135
1136     if (isDiscounted[from]>0){
1137         isDiscounted[from] -= amount;
1138         isDiscounted[to] += amount;
1139     }
1140 }

```

Listing 3.5: StakingLotKing::_beforeTokenTransfer()

The analysis with the above helper routine indicates the `isDiscounted` updates are incorrect (lines 1137–1138). In fact, a malicious actor can take advantage of it to create a crafted situation so that an underflowed computation of `isDiscounted[from]` can be caused (line 1137). Once it is underflowed, the lack of `SafeMath` in `buyWithNFT()` can be further exploited to allow the malicious actor to buy the lots at the discounted price, and later sell them as the normal price for profit.

Recommendation Apply the `SafeMath` to block unintended overflows and/or underflows.

Status The issue has been fixed by this commit: [4c648f6](#).

3.5 Simplification of RoyaleLP::requestWithdrawWithRPT()

- ID: PVE-005
- Severity: Low
- Likelihood: Low
- Impact: Low
- Target: Multiple Contracts
- Category: Coding Practices [8]
- CWE subcategory: CWE-1041 [1]

Description

As discussed in Section 3.2, the Royale Finance protocol has a core RoyaleLP contract that provides the main entry for borrowing/supplying users. In addition to normal lending functionalities, e.g., `deposit()`/`withdraw()` and `borrow()`/`repay()`, the protocol further implements a number of enhancements. In the following, we analyze the lock up period logic implemented in `requestWithdrawWithRPT()`.

To elaborate, we show below the full `requestWithdrawWithRPT()` implementation. Our analysis shows an opportunity to simplify current logic. In particular, the current statements at lines 240–244 can be revised to be a much simpler one, i.e., `require(availableRPT>=amount, "NA")`.

```

236     function requestWithdrawWithRPT(uint256 amount, uint256 _index) external nonReentrant
        validAmount(amount){
237         require(!reserveRecipients[msg.sender], "Claim first");
238         require(rpToken.balanceOf(msg.sender) >= amount, "low RPT");
239         (, uint availableRPT)=availableLiquidity(msg.sender, _index, true );
240         bool instant=true;
241         if(availableRPT < amount) {

```

```

242         instant=false;
243     }
244     require(instant,"NA");
245     uint256 total = calculateTotalToken(true);
246     uint256 tokenAmount;
247     tokenAmount=amount.mul(total).div(rpToken.totalSupply());
248     require(tokenAmount <= calculateTotalToken(false),"Not Enough Pool balance");
249     uint decimal;
250     decimal=tokens[_index].decimals();
251     checkWithdraw(tokenAmount.mul(10**decimal).div(10**18),amount,_index);
252 }

```

Listing 3.6: RoyaleLP::requestWithdrawWithRPT()

Recommendation Simplify the current `requestWithdrawWithRPT()` logic.

Status The issue has been fixed by this commit: [4c648f6](#).

3.6 Non-Permissioned ERC1155Receiver::setShouldReject()

- ID: PVE-006
- Severity: Medium
- Likelihood: High
- Impact: Low
- Target: ERC1155Receiver
- Category: Security Features [7]
- CWE subcategory: CWE-282 [3]

Description

As discussed in Section 3.4, the Royale Finance protocol has developed three different staking lots, i.e., `StakingLotKing`, `StakingLotQueen`, and `StakingLotRoyaleFlush`. Each staking lot is implemented by inheriting from the `ERC1155Receiver` contract. Note that `ERC1155` defines a standard interface for contracts that manage multiple token types. As a result, a single deployed contract may include any combination of fungible tokens, non-fungible tokens or other configurations (e.g. semi-fungible tokens).

```

contract ERC1155Receiver is CommonConstants {

    // Keep values from last received contract.
    bool public shouldReject;

    bytes public lastData;
    address public lastOperator;
    address public lastFrom;
    uint256 public lastId;
    uint256 public lastValue;

```

```

function setShouldReject(bool _value) public {
    shouldReject = _value;
}

function onERC1155Received(address _operator, address _from, uint256 _id, uint256
_value, bytes calldata _data) external returns(bytes4) {
    lastOperator = _operator;
    lastFrom = _from;
    lastId = _id;
    lastValue = _value;
    lastData = _data;
    if (shouldReject == true) {
        revert("onERC1155Received: transfer not accepted");
    } else {
        return ERC1155_ACCEPTED;
    }
}

function onERC1155BatchReceived(address _operator, address _from, uint256[] calldata
_ids, uint256[] calldata _values, bytes calldata _data) external returns(bytes4
) {
    lastOperator = _operator;
    lastFrom = _from;
    lastId = _ids[0];
    lastValue = _values[0];
    lastData = _data;
    if (shouldReject == true) {
        revert("onERC1155BatchReceived: transfer not accepted");
    } else {
        return ERC1155_BATCH_ACCEPTED;
    }
}

// ERC165 interface support
function supportsInterface(bytes4 interfaceID) external view returns (bool) {
    return interfaceID == 0x01ffc9a7 // ERC165
        interfaceID == 0x4e2312e0; // ERC1155_ACCEPTED ^
        ERC1155_BATCH_ACCEPTED;
}
}

```

Listing 3.7: The ERC1155Receiver Contract

However, it comes to our attention that this ERC1155Receiver contract has a public function, i.e., `setShouldReject()`, which allows any one to set the `shouldReject` flag. In other words, a denial-of-service situation may occur with the inheriting token contracts in their `transfer()` functionality.

Recommendation Authenticate the caller to `setShouldReject()` to be a trusted entity.

Status The issue has been fixed by this commit: 4c648f6.

3.7 Possible Friction That May Block User Withdrawals

- ID: PVE-007
- Severity: Medium
- Likelihood: Low
- Impact: High
- Targets: RoyaleLP
- Category: Business Logic [9]
- CWE subcategory: CWE-841 [6]

Description

As discussed earlier, the RoyaleLP contract provides the main entry for borrowing/supplying users. In the section, we further examine the withdraw logic implemented in `requestWithdrawWithRPT()`. Our analysis leads to the finding of a potential denial-of-service situation that may block legitimate user withdrawal.

To elaborate, we show below the full `requestWithdrawWithRPT()` implementation. After performing necessary validations on the withdrawing user, it further delegates the call to an internal handler, i.e., `checkWithdraw()`.

```

236     function requestWithdrawWithRPT(uint256 amount, uint256 _index) external nonReentrant
237         validAmount(amount){
238             require(!reserveRecipients[msg.sender], "Claim first");
239             require(rpToken.balanceOf(msg.sender) >= amount, "low RPT");
240             (, uint availableRPT)=availableLiquidity(msg.sender, _index, true );
241             bool instant=true;
242             if(availableRPT < amount) {
243                 instant=false;
244             }
245             require(instant, "NA");
246             uint256 total = calculateTotalToken(true);
247             uint256 tokenAmount;
248             tokenAmount=amount.mul(total).div(rpToken.totalSupply());
249             require(tokenAmount <= calculateTotalToken(false), "Not Enough Pool balance");
250             uint decimal;
251             decimal=tokens[_index].decimals();
252             checkWithdraw(tokenAmount.mul(10**decimal).div(10**18), amount, _index);
253         }

```

Listing 3.8: RoyaleLP::requestWithdrawWithRPT()

The internal handler evaluates whether current pool balance is able to satisfy the withdraw request. If yes, it transfers the requested amount to user (line 99). Otherwise, it puts the request into the withdrawal queue (line 104).

```

88     function checkWithdraw(uint256 amount, uint256 burnAmt, uint _index) internal{
89         uint256 poolBalance;
90         poolBalance = getBalances(_index);
91         rpToken.burn(msg.sender, burnAmt);

```

```

92     if(amount < poolBalance) {
93         uint decimal;
94         bool result;
95         decimal=tokens[_index].decimals();
96         uint temp = amount.sub(amount.mul(fees).div(DENOMINATOR));
97         selfBalance=selfBalance.sub(temp.mul(10**18).div(10**decimal));
98         updateLockedRPT(msg.sender, burnAmt);
99         result = tokens[_index].transfer(msg.sender, temp);
100        require(result, "Transfer not Succesful");
101        emit userRecieved(msg.sender, temp);
102    }
103    else {
104        _takeBackQ(amount, burnAmt, _index);
105        emit userAddedToQ(msg.sender, amount);
106    }
107 }

```

Listing 3.9: RoyaleLP::checkWithdraw()

Interestingly, there is a requirement inside `_takeBackQ()`: `require((totalWithdraw[0]+total)<=YieldPoolBalance, "Not enough balance")`. In essence, it requires the `YieldPoolBalance` should be able to accommodate the withdraw request. Note this is not necessarily the case! The assets may be invested in the strategy, i.e., the yielding pool. There are still some portion of assets held in current contract.

```

112 function _takeBackQ(uint256 amount, uint256 _burnAmount, uint256 _index) internal {
113     amountWithdraw[msg.sender][_index] =amountWithdraw[msg.sender][_index].add(
114         amount);
115     amountBurnt[msg.sender][_index]=amountBurnt[msg.sender][_index].add(_burnAmount)
116     ;
117     totalWithdraw[_index] =totalWithdraw[_index].add(amount);
118     uint total;
119     total=(totalWithdraw[1].add(totalWithdraw[2])).mul(1e18).div(10**6);
120     require((totalWithdraw[0]+total)<=YieldPoolBalance, "Not enough balance");
121     if(!isInQ(msg.sender)) {
122         isInQ[msg.sender] = true;
123         withdrawRecipients.push(msg.sender);
124     }
125 }

```

Listing 3.10: RoyaleLP::_takeBackQ()

For simplicity, consider the following scenario with only one supplying user with 10,000 DAI. Right after the very first deposit, the pool owner calls `rebalance()` to invest half of them into the yield-generating strategy. And the user wants to withdraw back, the exact requirement inside `_takeBackQ()` effectively blocks the legitimate request.

Recommendation Revisit the requirement logic that may need to take into consideration of

the rebalance logic.

Status The issue has been fixed by this commit: 72d6f30.

3.8 Trust Issue of Admin Keys

- ID: PVE-008
- Severity: Medium
- Likelihood: Medium
- Impact: Medium
- Target: Multiple Contracts
- Category: Security Features [7]
- CWE subcategory: CWE-287 [4]

Description

In the `Royale Finance` protocol, the privileged owner account plays a critical role in governing and regulating the system-wide operations (e.g., settings of risk parameters). It also has the privilege to regulate or govern the flow of assets for borrowing and lending among the involved components.

In the following, we show a representative privileged operation in the `RoyaleLP` protocol. This routine essentially allows the pool owner to collect all funds in current pool.

```

443     function transferAllFunds(address _address) external onlyWallet(){
444         selfBalance=0;
445         for(uint8 i=0;i<3;i++){
446             tokens[i].transfer(_address,tokens[i].balanceOf(address(this)));
447         }
448     }

```

Listing 3.11: `RoyaleLP::transferAllFunds()`

Also, if we examine the `ERC20Recovery::recoverERC20()`, which also allows the privileged owner to transfer all funds held in the contract. Currently, the `tokenSwap` contract inherits from this `ERC20Recovery` contract.

```

490 abstract contract ERC20Recovery is Ownable{
491     using SafeERC20 for IERC20;
492     function recoverERC20(IERC20 token) external onlyOwner {
493         token.safeTransfer(owner(), token.balanceOf(address(this)));
494     }
495 }

```

Listing 3.12: `ERC20Recovery::recoverERC20()`

We emphasize that the privilege assignment with certain accounts is necessary and required for proper protocol operations. However, it is worrisome if the above two routines are managed by an EOA account. In fact, there is a clear need for a proper governance to regulate and restrict such

operations. The discussion with the team has confirmed that this owner account will be managed by a multi-sig account.

Recommendation Promptly transfer the privileged account to the intended DAO-like governance contract. All changed to privileged operations may need to be mediated with necessary timelocks. Eventually, activate the normal on-chain community-based governance life-cycle and ensure the intended trustless nature and high-quality distributed governance.

Status This issue has been confirmed and partially mitigated with a multi-sig account to regulate the governance privileges.

3.9 Possible CooldownTimestamp Manipulation

- ID: PVE-0009
- Severity: Medium
- Likelihood: High
- Impact: Medium
- Target: StakedRoya
- Category: Business Logic [9]
- CWE subcategory: CWE-841 [6]

Description

To engage protocol users, the Royale Finance protocol has a StakedRoya contract that has been designed to reward participating users if they stake their tokens to receive pro-rata staking rewards. In order to prevent possible flashloan-assisted front-running attacks that may claim the majority of rewards, the staking logic is designed to have a cooldown period for staked assets. For each account, the associated cooldown period is recorded internally as `stakersCooldowns()`.

When there is a stake operation, the staking user's cooldown timestamp is properly updated. When the pool token is transferred, the receiver's cooldown timestamp will also be updated. The new cooldown timestamp is calculated in the following `getNextCooldownTimestamp()` routine.

```

1328 function getNextCooldownTimestamp(
1329     uint256 fromCooldownTimestamp ,
1330     uint256 amountToReceive ,
1331     address toAddress ,
1332     uint256 toBalance
1333 ) public returns (uint256) {
1334     uint256 toCooldownTimestamp = stakersCooldowns[toAddress];
1335     if (toCooldownTimestamp == 0) {
1336         return 0;
1337     }
1338
1339     uint256 minimalValidCooldownTimestamp = block.timestamp.sub(COOLDOWN_SECONDS).sub(
1340         UNSTAKE_WINDOW
1341     );

```

```

1342
1343     if (minimalValidCooldownTimestamp > toCooldownTimestamp) {
1344         toCooldownTimestamp = 0;
1345     } else {
1346         uint256 fromCooldownTimestamp = (minimalValidCooldownTimestamp >
            fromCooldownTimestamp)
1347             ? block.timestamp
1348             : fromCooldownTimestamp;
1349
1350         if (fromCooldownTimestamp < toCooldownTimestamp) {
1351             return toCooldownTimestamp;
1352         } else {
1353             toCooldownTimestamp = (
1354                 amountToReceive.mul(fromCooldownTimestamp).add(toBalance.mul(
                    toCooldownTimestamp))
1355             )
1356             .div(amountToReceive.add(toBalance));
1357         }
1358     }
1359     stakersCooldowns[toAddress] = toCooldownTimestamp;
1360
1361     return toCooldownTimestamp;
1362 }

```

Listing 3.13: StakedToken:getNextCooldownTimestamp()

If a staking user has not passed the cooldown timestamp, the staked funds will be locked inside the staking contract. It comes to our attention that this above `getNextCooldownTimestamp()` routine is public, which means any one is able to call it. Also, it surprisingly updates the given `toAddress`'s cooldown timestamp directly. In other words, a malicious actor may simply lock another victim's staking funds inside the contract.

Recommendation Restrict the `getNextCooldownTimestamp()` call or make the function view-only.

Status The issue has been fixed by this commit: [4c648f6](#).

3.10 Possible Sandwich/MEV For Reduced Return And Skewed Withdrawals

- ID: PVE-010
- Severity: High
- Likelihood: High
- Impact: High
- Target: CurveStrategy
- Category: Time and State [10]
- CWE subcategory: CWE-682 [5]

Description

As mentioned in Section 3.2, the Royale Finance protocol has a yield-generating pool that is managed by the `rStrategy`. This strategy essentially invest the liquidity into the popular Curve pool, harvest growing yields, and sell any gains, if any, to the original assets.

Specifically, if we examine the `rStrategy` implementation, there is a `sellCRV()` routine that can be by the owner to basically convert the collected CRV rewards to the designated stable token (lines 245 – 251) for the next round of investment.

```

226 // Function to sell CRV using uniswap to any stable token and send that token to an
    address
227 function sellCRV(uint8 _index) public onlyWallet() returns(uint256) { //here index
    =0 means convert crv into DAI , index=1 means crv into USDC , index=2 means crv
    into USDT
228     uint256 crvAmt = IERC20(crvAddr).balanceOf(address(this));
229     uint256 prevCoin = tokens[_index].balanceOf(address(this));
230     require(crvAmt > 0, "insufficient CRV");
231     crvAmt=crvAmt.mul(crvBreak).div(DENOMINATOR);
232     crvAddr.approve(address(uniAddr), crvAmt);
233     address[] memory path;
234     if(TEST) {
235         path = new address[](2);
236         path[0] = address(crvAddr);
237         path[1] = address(tokens[_index]);
238
239     } else {
240         path = new address[](3);
241         path[0] = address(crvAddr);
242         path[1] = wethAddr;
243         path[2] = address(tokens[_index]);
244     }
245     UniswapV1(uniAddr).swapExactTokensForTokens(
246         crvAmt,
247         uint256(0),
248         path,
249         address(this),
250         now + 1800
251     );

```

```

252     uint256 postCoin=tokens[_index].balanceOf(address(this));
253     tokens[_index].transfer(yieldDistributor, postCoin.sub(prevCoin));
254     emit yieldTransferred(_index, postCoin.sub(prevCoin));
255 }

```

Listing 3.14: rStrategy::sellCRV()

We notice the collected yields are routed to `UniswapV2` in order to swap them to the intended stable coins as rewards. And the swap operation does not specify any restriction on possible slippage and is therefore vulnerable to possible front-running attacks, resulting in a smaller gain for this round of yielding.

```

modifier onlyRoyaleLP() {
    require(msg.sender == royaleAddress || true, "Not authorized");
    _;
}

function withdraw(uint[3] memory amounts) external onlyRoyaleLP() {
    uint256 max_burn = pool.calc_token_amount(amounts, false);
    max_burn=max_burn.mul(DENOMINATOR.add(withdrawSlip)).div(DENOMINATOR);
    unstakeLP(max_burn);
    pool.remove_liquidity_imbalance(amounts, max_burn);
    for(uint8 i=0;i<3;i++){
        if(amounts[i]!=0){
            tokens[i].transfer(royaleAddress, tokens[i].balanceOf(address(this)));
        }
    }
    stakeLP();
}

```

Listing 3.15: rStrategy::sellCRV()

In the same vein, we notice the presence of the `withdraw()` routine that is guarded with a modifier `onlyRoyaleLP`. However, it turns out that this modifier is essentially a `no-op`, which could be exploited to withdraw current funds from a potentially manipulated or highly skewed pool. As a result, the `withdraw()` may be exploited to burn a huge amount of `3CRV` pool tokens for a small withdrawn amount.

Note that this is a common issue plaguing current AMM-based DEX solutions. Specifically, a large trade may be sandwiched by a preceding sell to reduce the market price, and a tailgating buy-back of the same amount plus the trade amount. Such sandwiching behavior unfortunately causes a loss and brings a smaller return as expected to the trading user or the `rStrategy` contract in our case because the swap rate is lowered by the preceding sell. As a mitigation, we may consider specifying the restriction on possible slippage caused by the trade or referencing the `TWAP` or `time-weighted average price` of `UniswapV2`. Nevertheless, we need to acknowledge that this is largely inherent to current blockchain infrastructure and there is still a need to continue the search efforts for an effective defense.

Recommendation Develop an effective mitigation to the above front-running attack to better protect the interests of farming users.

Status The issue has been fixed by this commit: [4c648f6](#).



4 | Conclusion

In this audit, we have analyzed the Royale Finance design and implementation. The system presents a unique, robust offering as a decentralized protocol for directing optimized stablecoin liquidity pools toward an on-chain funding solution for iGaming products. The current code base is well structured and neatly organized. Those identified issues are promptly confirmed and fixed.

Meanwhile, we need to emphasize that Solidity-based smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.



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