

## Experiment No. 4

**Aim:-** Write a program to find first and follow sets for the given grammar. Program should accept the grammar from user and output the first and follow sets for each of the grammar symbol.

### Theory:-

$\text{First}(\alpha)$  is set of terminals that begins strings derived from  $\alpha$ . If  $\alpha \Rightarrow \epsilon$  then  $\epsilon$  is also in  $\text{First}(\epsilon)$ .

In predictive parsing when we have  $A \rightarrow \alpha | \beta$ , if  $\text{First}(\alpha)$  and  $\text{First}(\beta)$  are disjoint sets then we can select appropriate A-production by looking at the next input.

$\text{Follow}(A)$ , for any non terminal A, is set of terminals a that can appear immediately after A in some sentential form

If we have  $S \Rightarrow \alpha A a \beta$  for some  $\alpha$  and  $\beta$  then a is in  $\text{Follow}(A)$

If A can be the rightmost symbol in some sentential form, then \$ is in  $\text{Follow}(A)$

### Finding First:-

**First(x)** for all grammar symbols X

Apply following rules:

1. If X is terminal,  $\text{FIRST}(X) = \{X\}$ .
2. If  $X \rightarrow \epsilon$  is a production, then add  $\epsilon$  to  $\text{FIRST}(X)$ .
3. If X is a non-terminal, and  $X \rightarrow Y_1 Y_2 \dots Y_k$  is a production, and  $\epsilon$  is in all of  $\text{FIRST}(Y_1), \dots, \text{FIRST}(Y_k)$ , then add  $\epsilon$  to  $\text{FIRST}(X)$ .
4. If X is a non-terminal, and  $X \rightarrow Y_1 Y_2 \dots Y_k$  is a production, then add a to  $\text{FIRST}(X)$  if for some i, a is in  $\text{FIRST}(Y_i)$ , and  $\epsilon$  is in all of  $\text{FIRST}(Y_1), \dots, \text{FIRST}(Y_{i-1})$ .

Applying rules 1 and 2 is obvious. Applying rules 3 and 4 for  $\text{FIRST}(Y_1 Y_2 \dots Y_k)$  can be done as follows:

Add all the non- $\epsilon$  symbols of  $\text{FIRST}(Y_1)$  to  $\text{FIRST}(Y_1 Y_2 \dots Y_k)$ . If  $\epsilon \in \text{FIRST}(Y_1)$ , add all the non- $\epsilon$  symbols of  $\text{FIRST}(Y_2)$ . If  $\epsilon \in \text{FIRST}(Y_1)$  and  $\epsilon \in \text{FIRST}(Y_2)$ , add all the non- $\epsilon$  symbols of  $\text{FIRST}(Y_3)$ , and so on. Finally, add  $\epsilon$  to  $\text{FIRST}(Y_1 Y_2 \dots Y_k)$  if  $\epsilon \in \text{FIRST}(Y_i)$ , for all  $1 \leq i \leq k$ .

### **Example:**

Consider the following grammar.

$E \rightarrow E + T \mid T$

$T \rightarrow T * F \mid F$

$F \rightarrow (E) \mid \text{id}$

Grammar after removing left recursion:

$E \rightarrow TX$

$X \rightarrow +TX \mid \epsilon$

$T \rightarrow FY$

$Y \rightarrow *FY \mid \varepsilon$

$F \rightarrow (E) \mid id$

For the above grammar, following the above rules, the FIRST sets could be computed as follows:

$FIRST(E) = FIRST(T) = FIRST(F) = \{ (, id \}$

$FIRST(X) = \{ +, \varepsilon \}$

$FIRST(Y) = \{ *, \varepsilon \}$

### **Finding Follow:-**

To compute Follow(A) for all nonterminals A, apply following rules until nothing can be added to any follow set:

1. Place \$ in Follow(S) where S is the start symbol
2. If there is a production  $A \rightarrow \alpha B \beta$  then everything in First( $\beta$ ) except  $\varepsilon$  is in Follow(B).
3. If there is a production  $A \rightarrow B$  or a production  $A \rightarrow \alpha B \beta$  where First( $\beta$ ) contains  $\varepsilon$ , then everything in Follow(A) is in Follow(B)

### **Example:-**

$E \rightarrow TX$

$X \rightarrow +TX \mid \varepsilon$

$T \rightarrow FY$

$Y \rightarrow *FY \mid \varepsilon$

$F \rightarrow (E) \mid id$

$FOLLOW(E) = \{ \$, ) \}$

$FOLLOW(X) = \{ \$, ) \}$

$FOLLOW(T) = \{ +, \$, ) \}$

$FOLLOW(Y) = \{ +, \$, ) \}$

$FOLLOW(F) = \{ *, +, \$, ) \}$

### **Input Specifications:**

1. Students are supposed to take grammar productions(non-left recursive) as input.
2. They can implement the given program in any language like C, C++, Java or Python, etc.

```
# First and Follow
```

```
import re
```

```
no_of_rules=int(input("Enter the number of rules : "))
```

```
list_of_rules=[]
```

```
i=0
```

```
while i<no_of_rules :
```

```
    list_of_rules.append(input("Rule "+str(i+1)+" : "))
```

```
    i=i+1
```

```
print(list_of_rules)
```

```
old_rules=[]
```

```

rule_status=[]
rule_id=[]
first_list=[]
i=0
while i<no_of_rules :
    old_rules.append(re.split('->\\|',list_of_rules[i]))
    rule_status.append(False)
    rule_id.append(old_rules[i][0].strip())
    i=i+1
print(old_rules)
# Rule 1 and Rule 2
i=0
while i<len(old_rules) :
    j=1
    temp_list=[]
    value=True
    while j<len(old_rules[i]):
        temp=old_rules[i][j].strip()
        if temp != "abs" and not temp[0].islower():
            value=False
            break
        j+=1
    rule_status[i]=value
    j=1
    while j<len(old_rules[i]):
        temp=old_rules[i][j].strip()
        if temp == "abs" :
            temp_list.append(temp)
        elif temp[0].islower() or temp.islower() :
            temp_list.append(temp[0])
        j+=1
    first_list.append(temp_list)
    i+=1
print(rule_status)
print(rule_id)
print(first_list)
print("-----")
print("-----")

while False in rule_status:
    i=0
    while i<len(old_rules):
        if(rule_status[i]==True):
            i=i+1
            continue

```

```

j=1
temp_status=[False]*(len(old_rules[i])-1)
m=1
while m<(len(old_rules[i])):
    if old_rules[i][m].strip() == "abs" or
(old_rules[i][m].strip())[0].islower():
        temp_status[m-1]=True
        m+=1

while j<len(old_rules[i]) :
    if old_rules[i][j].strip() == "abs" or
(old_rules[i][j].strip())[0].islower():
        j+=1
        continue
    temp=old_rules[i][j].strip()

y = [False] * len(temp)
k=0
for a in temp:
    if a.islower() or rule_status[rule_id.index(a)] :
        y[k]=True
        k+=1

if False in y:
    break
else:
    for a in temp:

        if not a.islower() and not "abs" in
first_list[rule_id.index(a)]:
            first_list[i] = list(set(first_list[i] +
first_list[rule_id.index(a)]))
            break
        elif not a.islower() and "abs" in
first_list[rule_id.index(a)] and temp[-1]!=a:
            temp1=list(first_list[rule_id.index(a)])
            temp1.remove("abs")
            first_list[i] = list(set(first_list[i]
+temp1))

            elif not a.islower() and "abs" in
first_list[rule_id.index(a)] and temp[-1]==a:
            temp1=list(first_list[rule_id.index(a)])
            first_list[i] = list(set(first_list[i]
+temp1))

```

```

        elif a.islower():
            if not a in first_list[i]:
                first_list[i].append(a)
            break
        else:
            None
            print(first_list)
            temp_status[j-1]=True
        j+=1
    if not False in temp_status:
        rule_status[i]=True
    i+=1

print("First : ")
print(first_list)
print("Follow :")
start=input("Enter the start symbol : ")
follow_list=[]

i=0
while i< no_of_rules:
    temp_list=[]
    if old_rules[i][0] == start :
        temp_list.append("$")
    follow_list.append(temp_list)
    i+=1
print(follow_list)
i=0
while i<no_of_rules:
    character=old_rules[i][0]
    j=0
    while j<no_of_rules:
        temp=''.join(old_rules[j][1:])
        if not character in temp:
            j=j+1
            continue
        follow=[]
        list_contain_char=[]
        for word in old_rules[j][1:]:
            if character in word:
                list_contain_char.append(word)
        for word in list_contain_char:
            start_index=word.index(character)+1

```

```

        while start_index<len(word):
            if word[start_index].islower():
                if not word[start_index] in follow:
                    follow.append(word[start_index])
                break
            elif start_index!=(len(word)-1) and not "abs" in
first_list[rule_id.index(word[start_index])]:
                follow = list(set(follow +
first_list[rule_id.index(word[start_index])]))
                break
            elif start_index!=(len(word)-1) and "abs" in
first_list[rule_id.index(word[start_index])]:

temp1=list(first_list[rule_id.index(word[start_index])])
                temp1.remove("abs")
                follow = list(set(follow + temp1))
            elif start_index==(len(word)-1) and not "abs" in
first_list[rule_id.index(word[start_index])]:
                follow = list(set(follow +
first_list[rule_id.index(word[start_index])]))
                break
            elif start_index==(len(word)-1) and "abs" in
first_list[rule_id.index(word[start_index])]:

temp1=list(first_list[rule_id.index(word[start_index])])
                temp1.remove("abs")
                follow = list(set(follow + temp1+follow_list[j]))

                start_index+=1
            if start_index==(len(word)):
                follow = list(set(follow + follow_list[j]))
                follow_list[i]=list(set(follow_list[i]+follow))
        j+=1
        i+=1
print(follow_list)

```

Output:

```

# root@coder:/coder/mnt/Rstar/SPCC_PY_CODES# python first_follow.py
# Enter the number of rules : 5
# Rule 1 : S->ABCD
# Rule 2 : A->a|abs
# Rule 3 : B->CD|b

```

```

# Rule 4 : C->c|abs
# Rule 5 : D->Aa|d|abs
# ['S->ABCD', 'A->a|abs', 'B->CD|b', 'C->c|abs', 'D->Aa|d|abs']
# [['S', 'ABCD'], ['A', 'a', 'abs'], ['B', 'CD', 'b'], ['C', 'c', 'abs'], ['D', 'Aa', 'd', 'abs']]
# [False, True, False, True, False]
# ['S', 'A', 'B', 'C', 'D']
# [[], ['a', 'abs'], ['b'], ['c', 'abs'], ['d', 'abs']]
# -----
--
# [[], ['a', 'abs'], ['b'], ['c', 'abs'], ['d', 'a', 'abs']]
# [[], ['a', 'abs'], ['abs', 'd', 'b', 'a', 'c'], ['c', 'abs'], ['d', 'a', 'abs']]
# [['abs', 'b', 'd', 'a', 'c'], ['a', 'abs'], ['abs', 'd', 'b', 'a', 'c'], ['c', 'abs'], ['d', 'a', 'abs']]
# First :
# [['abs', 'b', 'd', 'a', 'c'], ['a', 'abs'], ['abs', 'd', 'b', 'a', 'c'], ['c', 'abs'], ['d', 'a', 'abs']]
# Follow :
# Enter the start symbol : S
# [['$'], [], [], [], []]
# [['$'], ['$ ', 'd', 'b', 'a', 'c'], ['d', 'a', '$ ', 'c'], ['d', 'a', '$ ', 'c'], ['d', 'a', '$ ', 'c']]
# root@coder:/coder/mnt/Rstar/SPCC_PY_CODES#

```

### **Post Lab Assignment:-**

**Q1. Which symbols are used as synchronizing tokens in predictive parsing?**

A. Follow symbols    B. First symbols    C. First and Follow symbols    D. None of these

### **Answer:**

First and Follow symbols are used as synchronizing tokens in predictive parsing i.e. **Option C.**

**Q2. Find the First and Follow sets for the grammar**

```

E -> TP
P -> Eps | |TP
T -> FQ
Q -> Eps | FQ
F -> (E) R | iR
R -> Eps | *R

```

**Answer:**

	<b>First</b>	<b>Follow</b>
E -> TP	{ ( , i }	{ \$ , ) }
P -> Eps    TP	{ Eps ,   }	{ \$ , ) }
T -> FQ	{ ( , i }	{   , \$ , ) }
Q -> Eps   FQ	{ Eps , ( , i }	{   , \$ , ) }
F -> (E) R   iR	{ ( , i }	{   , \$ , ( , ) , i }
R -> Eps   *R	{ Eps , * }	{   , \$ , ( , ) , i }

**Conclusion:** From this experiment I understood about the computation of first and follow. This concept will be further used in next experiment