

# Graph Theory: Homework #7

Due on February 18, 2015

*Professor McGinley MWF 9:15*

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## Problem 1

Determine all  $r, s$  such that  $K_{r,s}$  is planar.

### Solution

$K_{r,s}$  is planar whenever either  $r$  or  $s$  is less than or equal to 2. When either is 1, the graph is a tree, so it is planar. With either  $r$  or  $s$  equal to 2, the graph can be drawn with the larger set in the middle, with either of the pair on each side.

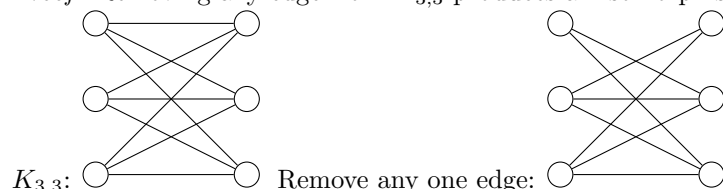
When both are greater than 2, the graph contains a subdivision of  $K_{3,3}$ , so it is not planar.

## Problem 2

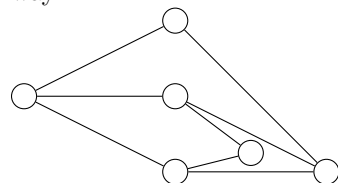
Show that the graph obtained by deleting one edge of  $K_{3,3}$  is planar.

### Solution

*Proof.* Removing any edge from  $K_{3,3}$  produces an isomorphism of the same graph.



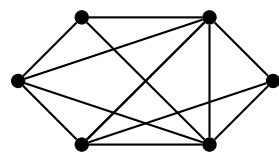
Moving the top and bottom vertices in the left set, we can embed the graph in the plane in the following way



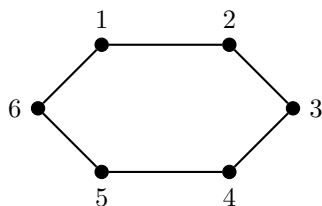
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## Problem 3

Determine (using the algorithm given in class) if the graph below is planar. If so, give its planar embedding.

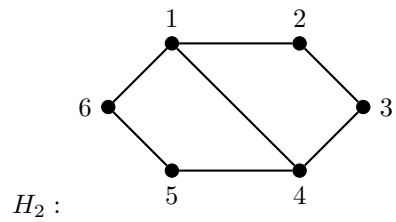


### Solution

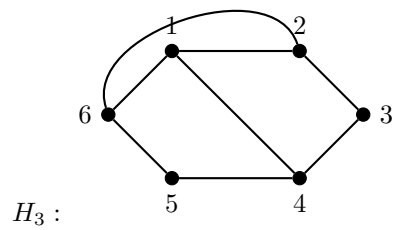


Choosing the exterior cycle  $H_1$  :

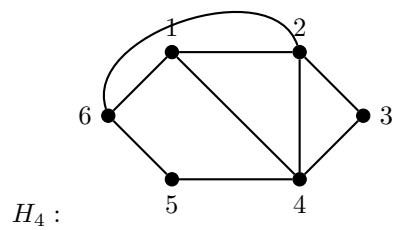
Any segment can be placed in either  $r_1$  or  $r_2$ , so choose  $s_1$  in  $r_1$ .



$R(s_1, H_2) = r_3$ . Choose  $s_1$  now because it can only be placed in one region.



Once again,  $s_1$  can only be placed in one region,  $r_4$ , so preemptively choose it for simplicity.

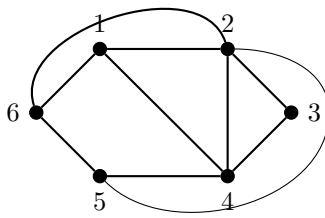


$$R(s_1, H_4) = r_5, r_1$$

$$R(s_2, H_4) = r_5$$

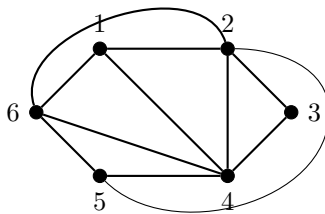
$$R(s_3, H_4) = r_5$$

Choose  $s_2$



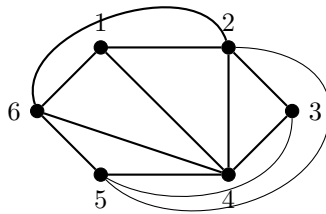
$H_5$  :

$s_1$  can be placed in  $r_1$ .



$H_6$  :

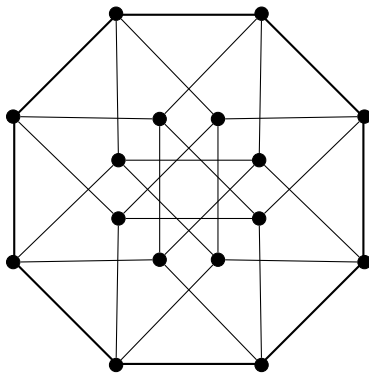
The final edge can be placed in  $r_5$ . Our final planar graph is



$G$  :

## Problem 4

Is the graph below planar? Use one of the theorems we talked about in class, not the algorithm.



**Solution**

*Proof.* According to a corollary of Euler's formula, for a graph to be planar,  $e(G) \leq 3n(G) - 6$ .

$$n(G) = 16$$

$$e(G) = 32$$

$$e(G) \leq 3n(G) - 6$$

$$32 \leq 48 - 6$$

$$32 \leq 42$$

For graphs with no triangular faces, it also must be true that  $e(G) \leq 2n(G) - 4$ . This graph has no triangular faces, so

$$e(G) \leq 2n(G) - 4$$

$$32 \leq 32 - 4$$

$$32 \leq 28$$

This is not true, so this graph is not planar.

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