Sorting algorithms

Data Structures and Algorithms

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Overview

1. **Sorting concepts**
2. **Insertion Sort**
   - Straight Insertion Sort
   - Shell Sort
3. **Selection Sort**
   - Straight Selection Sort
4. **Exchange Sort**
   - Bubble Sort
5. **Devide-and-Conquer**
   - Quick Sort
   - Merge Sort

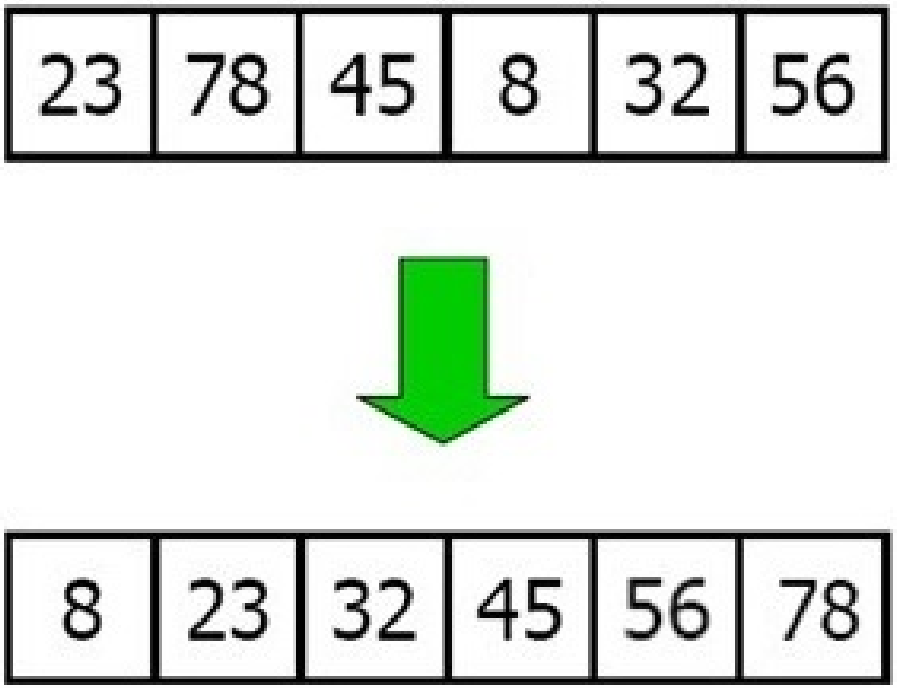
Course learning outcomes

L.O.1	Determine the complexity of simple algorithms (polynomial time - nested loop - no recursive)
L.O.1.1	Give definition of Big-O notation
L.O.1.2	Determine complexity of simple polynomial algorithms
L.O.2	Manipulate basic data structures such as list, tree and graph
L.O.2.1	Describe and present basic data structures such as: array, linked list, stack, queue, tree, and graph
L.O.2.2	Implement basic methods for each of basic data structures: array, linked list, stack, queue, tree, and graph
L.O.3	Implement basic sorting and searching algorithms
L.O.3.1	Illustrate how searching algorithms work on data structures: array, linked list, stack, queue, tree, and graph
L.O.3.2	Illustrate how sorting algorithms work on an array
L.O.3.3	Implement necessary methods and proposed algorithms on a given data structure for problem solving

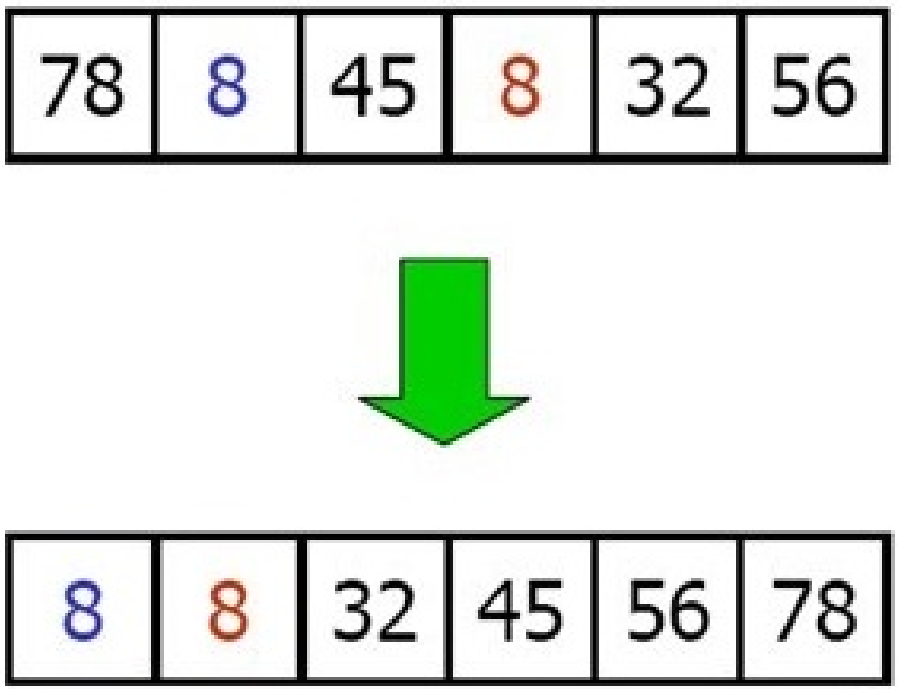
Sorting concepts

Sorting

One of the most important concepts and common applications in computing.

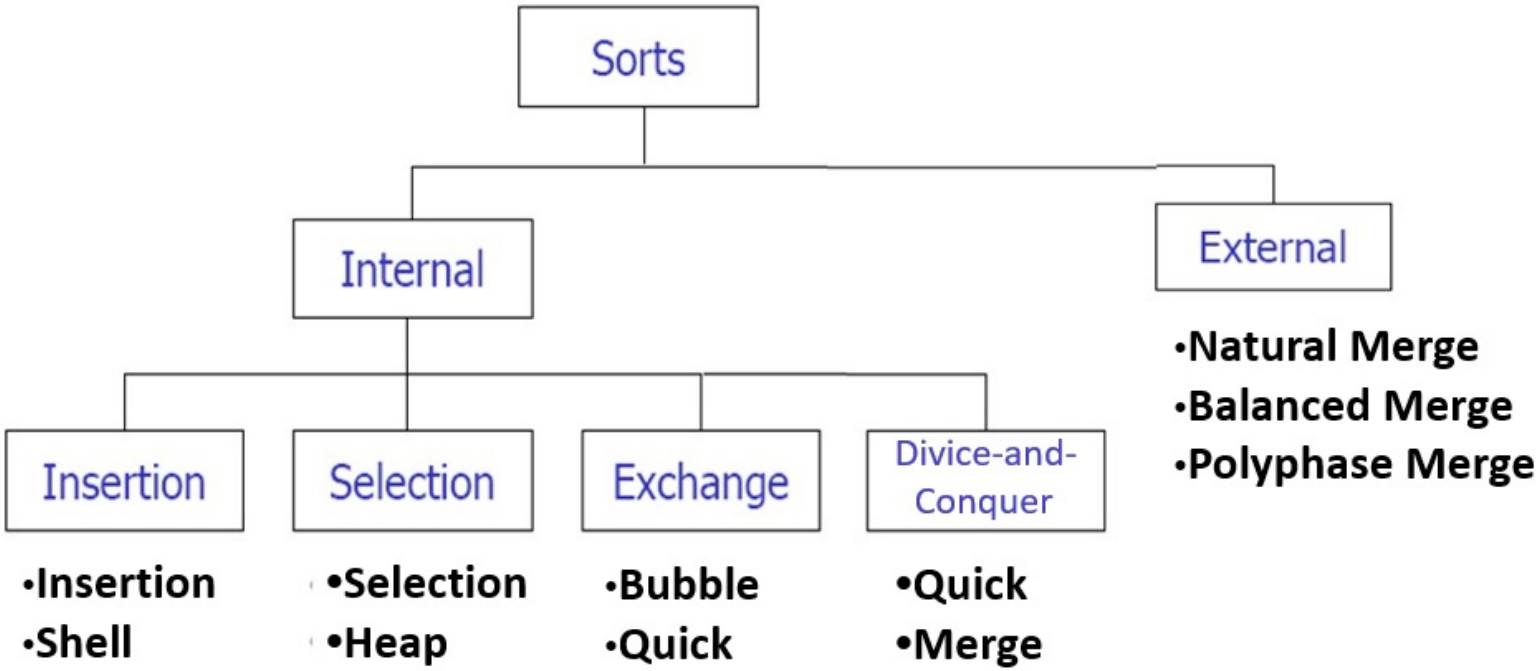


Sort stability: data with equal keys maintain their relative input order in the output.



Sort efficiency: a measure of the relative efficiency of a sort = number of comparisons + number of moves.

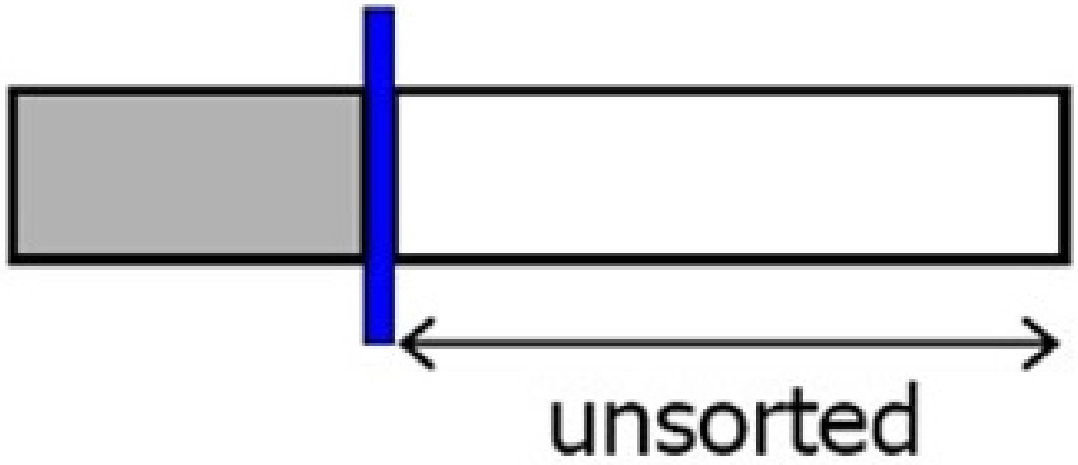
Sorting



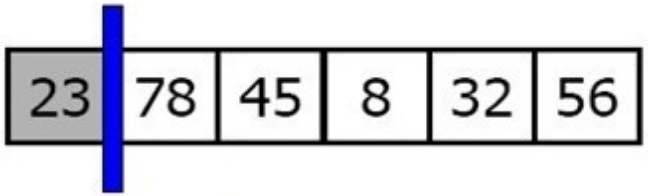
Insertion Sort

Straight Insertion Sort

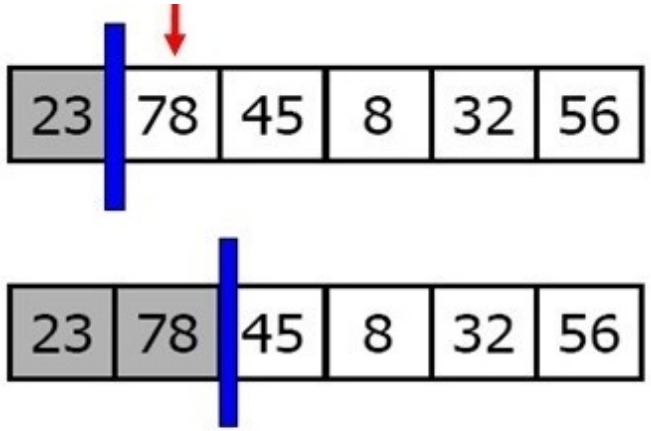
- The list is divided into two parts: sorted and unsorted.
- In each pass, the first element of the unsorted sublist is inserted into the sorted sublist.



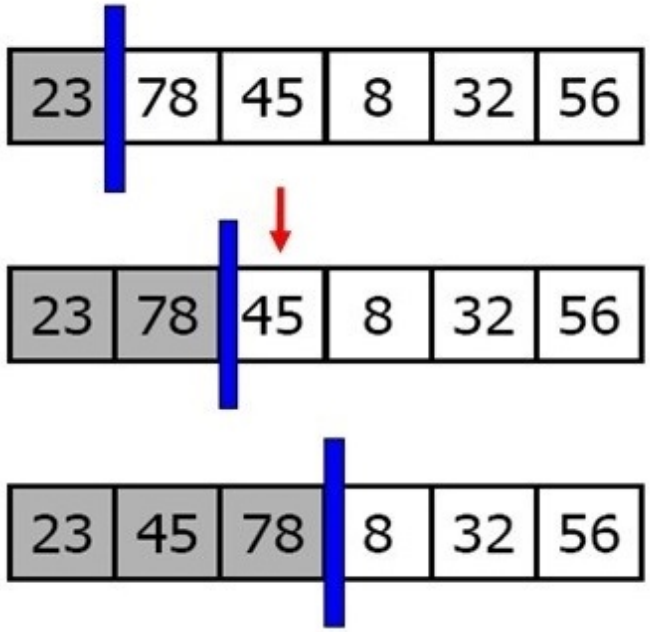
Straight Insertion Sort



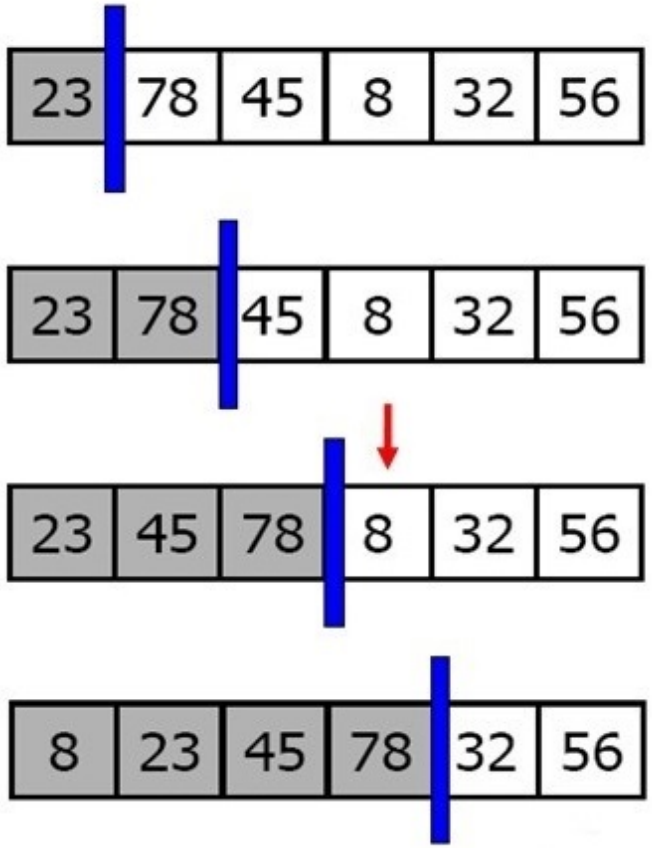
Straight Insertion Sort



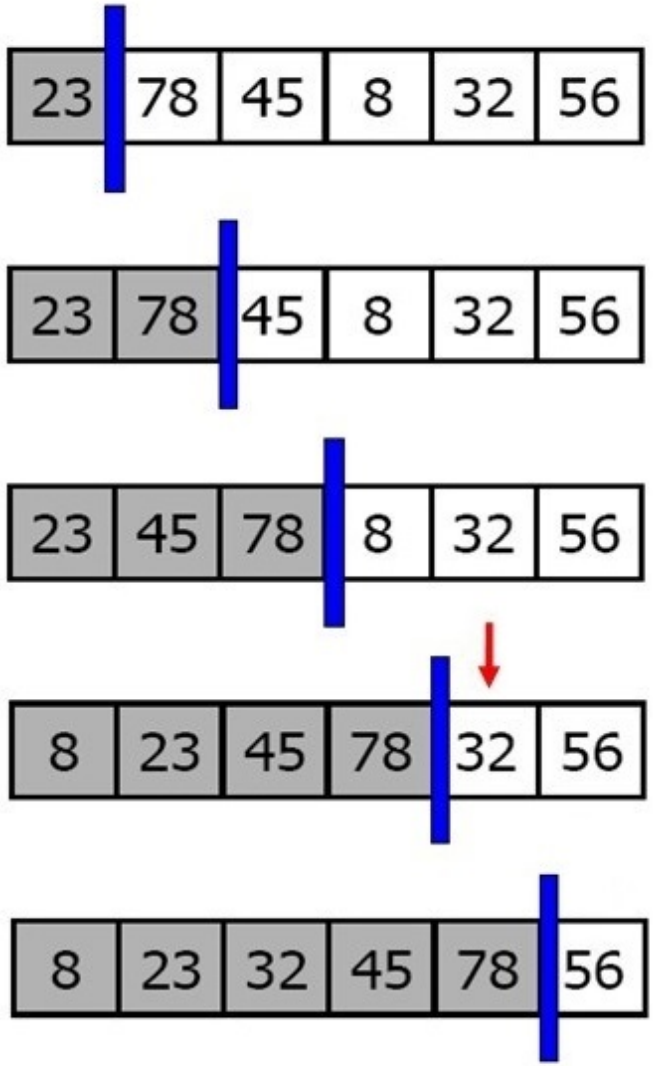
Straight Insertion Sort



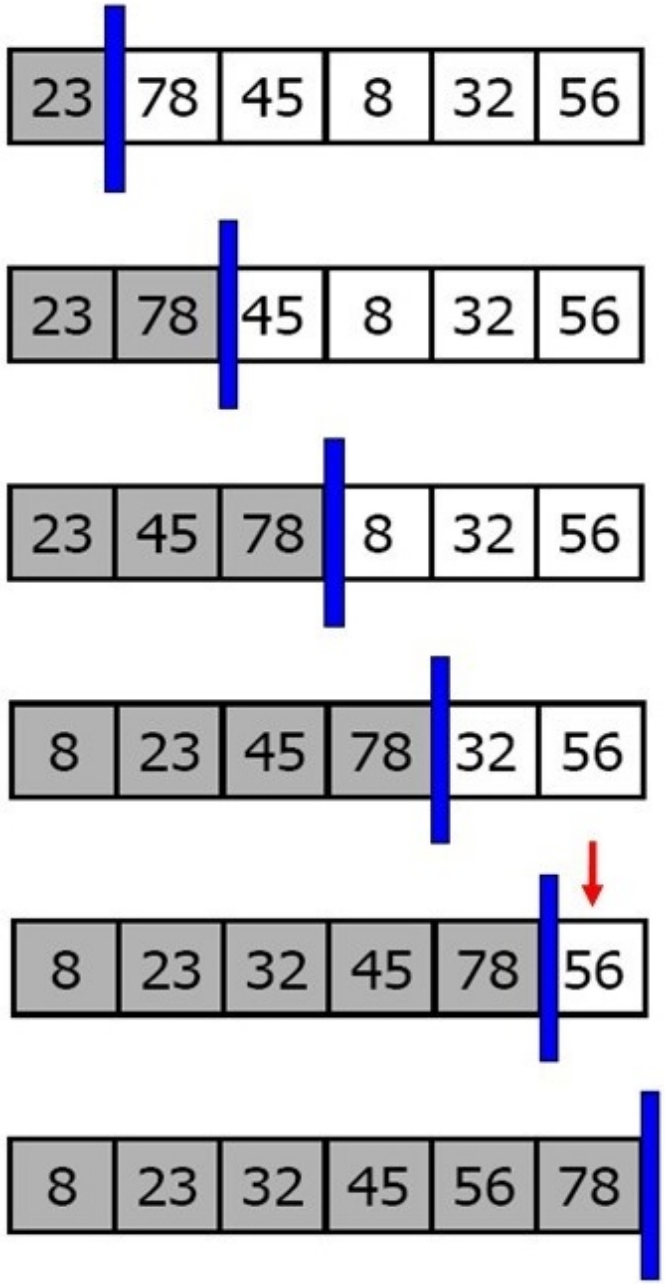
Straight Insertion Sort



Straight Insertion Sort



Straight Insertion Sort



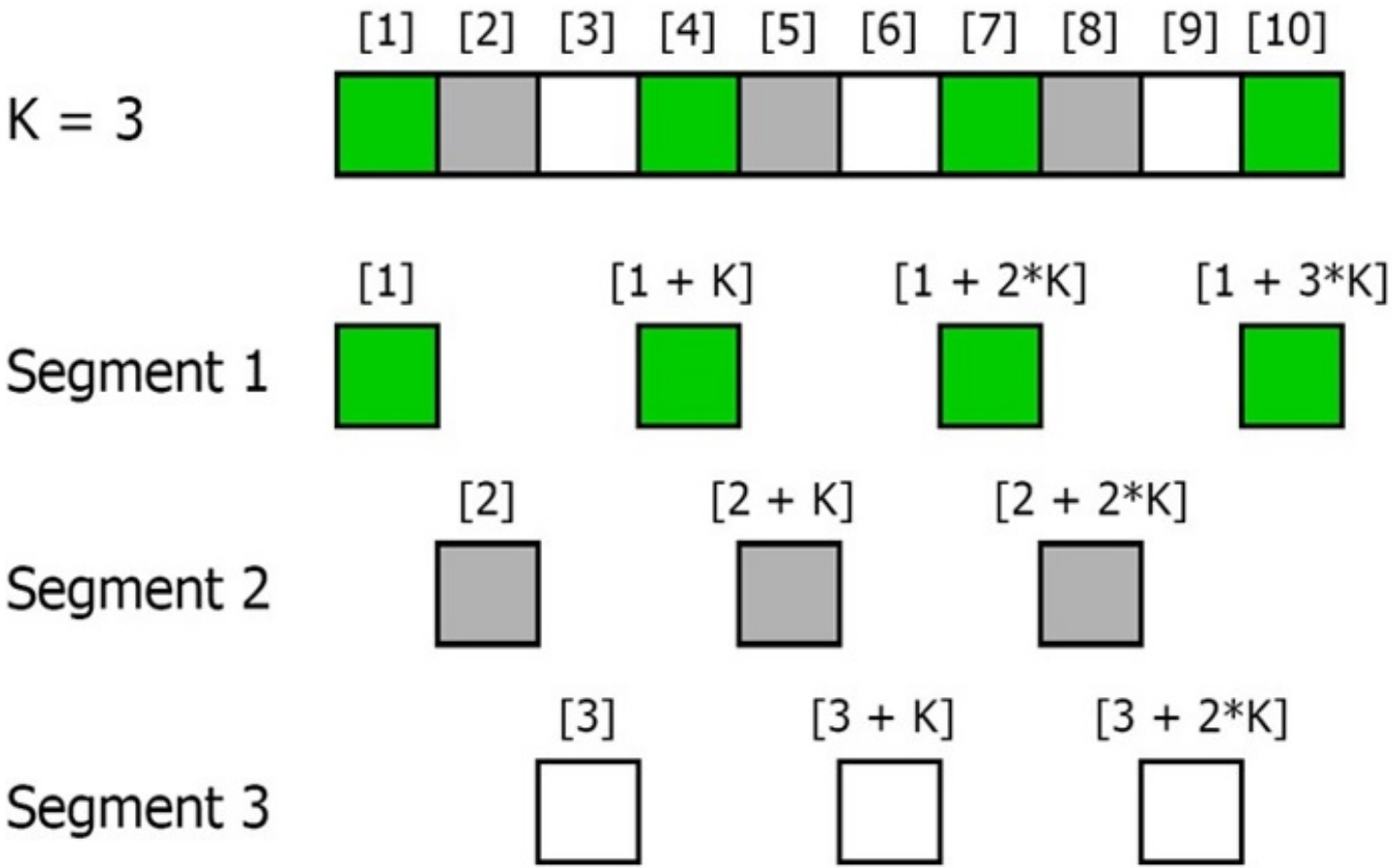
Straight Insertion Sort

```
1  Algorithm InsertionSort()
2  Sorts the contiguous list using straight insertion
     sort.

3  if count > 1 then
4      current = 1
5      while current < count do
6          temp = data[current]
7          walker = current - 1
8          while walker >= 0 AND temp.key <
             data[walker].key do
9              data[walker+1] = data[walker]
10             walker = walker - 1
11         end
12         data[walker+1] = temp
13         current = current + 1
14     end
15 end
```

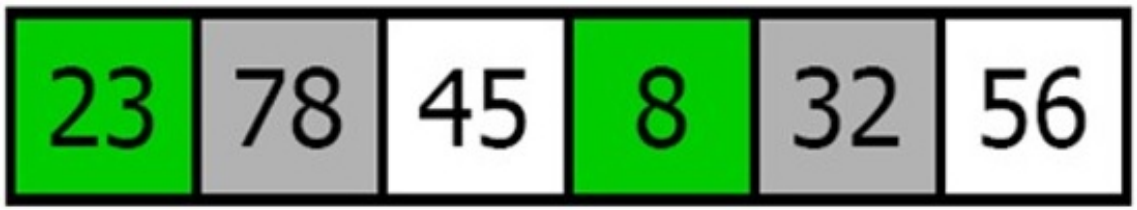

Shell Sort

- Named after its creator Donald L. Shell (1959).
- Given a list of N elements, the list is divided into K segments (K is called the increment).
- Each segment contains N/K or more elements.
- Segments are dispersed throughout the list.
- Also is called diminishing-increment sort.

Shell Sort

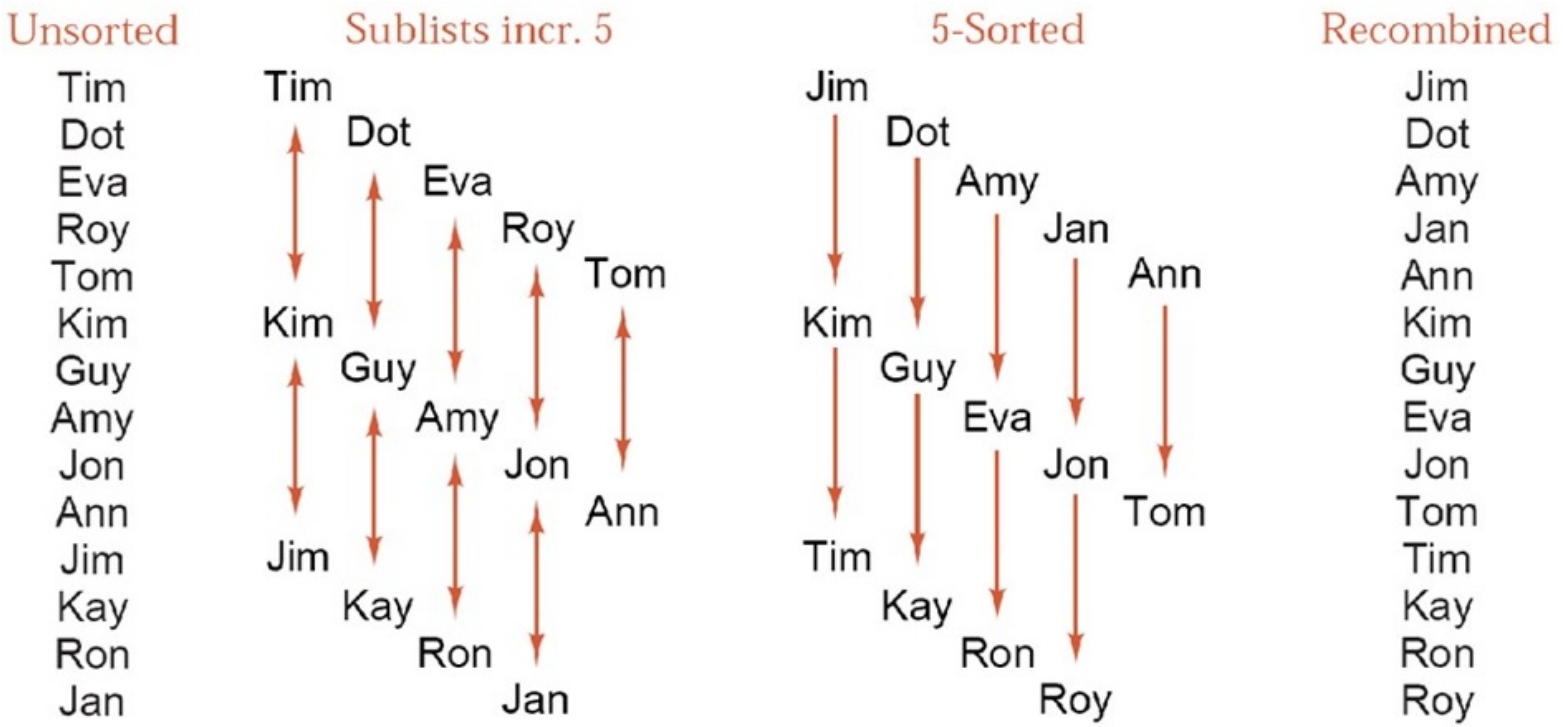


Shell Sort

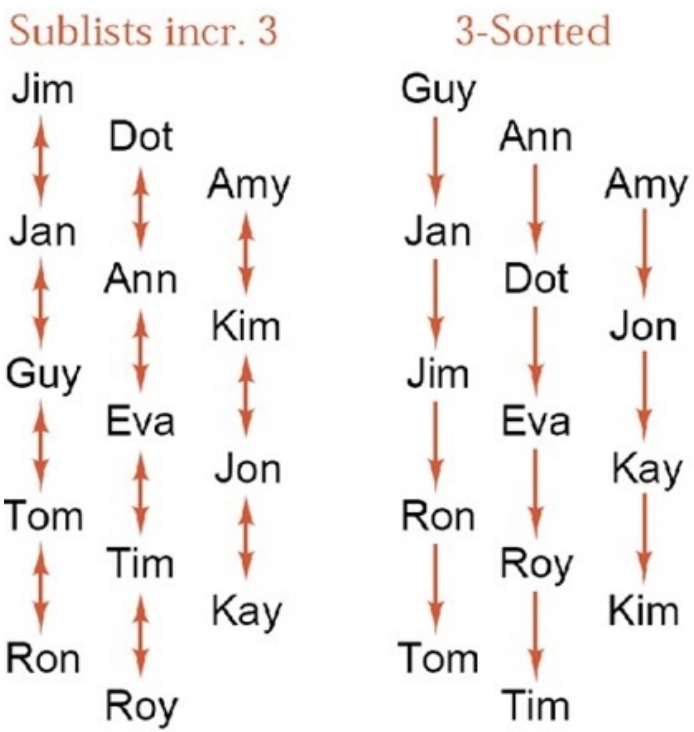


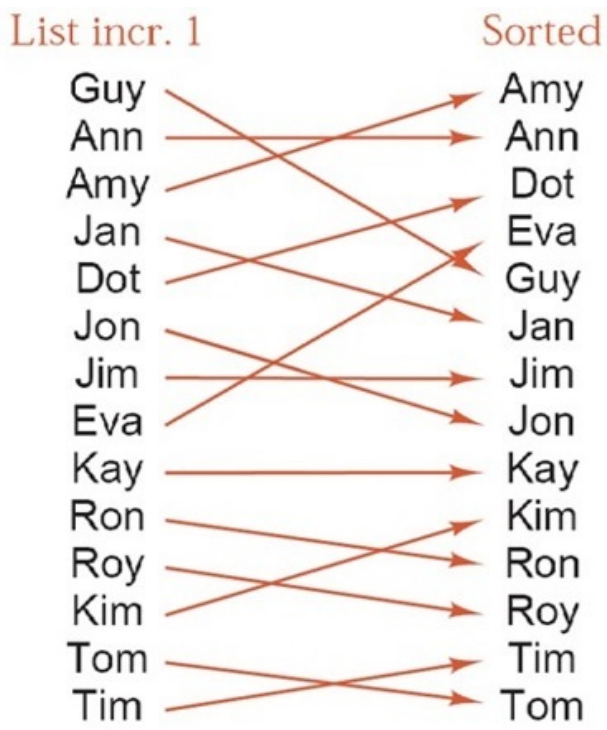
- For the value of K in each iteration, sort the K segments.
- After each iteration, K is reduced until it is 1 in the final iteration.

Example of Shell Sort



Example of Shell Sort





Choosing incremental values

- From more of the comparisons, it is better when we can receive more new information.
- Incremental values should not be multiples of each other, other wise, the same keys compared on one pass would be compared again at the next.
- The final incremental value must be 1.

Choosing incremental values

- Incremental values may be:
  $1, 4, 13, 40, 121, ...$
  $k_t = 1$
  $k_{i-1} = 3 * k_i + 1$
  $t = \lfloor \log_3 n \rfloor - 1$

- or:
  $1, 3, 7, 15, 31, ...$
  $k_t = 1$
  $k_{i-1} = 2 * k_i + 1$
  $t = \lfloor \log_2 n \rfloor - 1$

Shell Sort

```
1  Algorithm ShellSort()
2  Sorts the contiguous list using Shell sort.

3  k = first_incremental_value
4  while k >= 1 do
5  |   segment = 1
6  |   while segment <= k do
7  |   |   SortSegment(segment)
8  |   |   segment = segment + 1
9  |   end
10 |   k = next_incremental_value
11 end
12 End ShellSort
```


Shell Sort

```
1  Algorithm SortSegment(val segment <int>, val k
     <int>)
2  Sorts the segment beginning at segment using
     insertion sort, step between elements in the
     segment is k.

3  current = segment + k
4  while current < count do
5  |   temp = data[current]
6  |   walker = current - k
7  |   while walker >=0 AND temp.key <
   |     data[walker].key do
8  |   |   data[walker + k] = data[walker]
9  |   |   walker = walker - k
10 |   end
11 |   data[walker + k] = temp
12 |   current = current + k
13 end
```


Insertion Sort Efficiency

- Straight insertion sort:
  $f(n) = n(n+1)/2 = O(n^2)$
- Shell sort:
  $O(n^{1.25})$ (Empirical study)

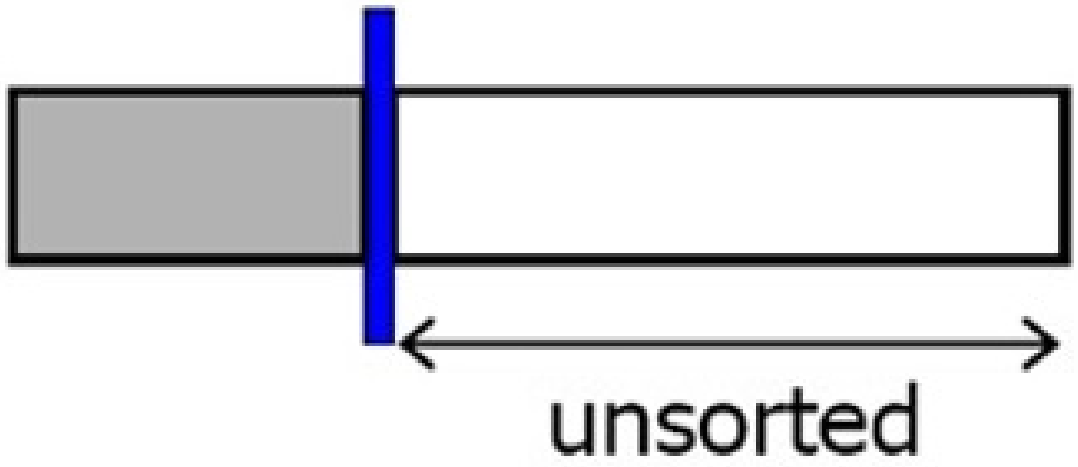
Selection Sort

Selection Sort

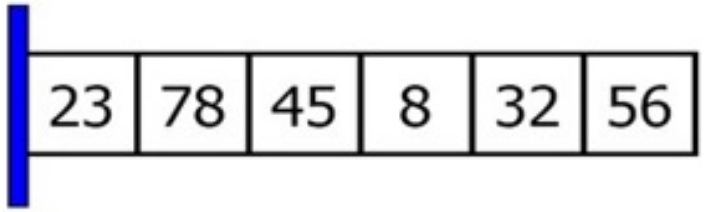
In each pass, the smallest/largest item is selected and placed in a sorted list.

Straight Selection Sort

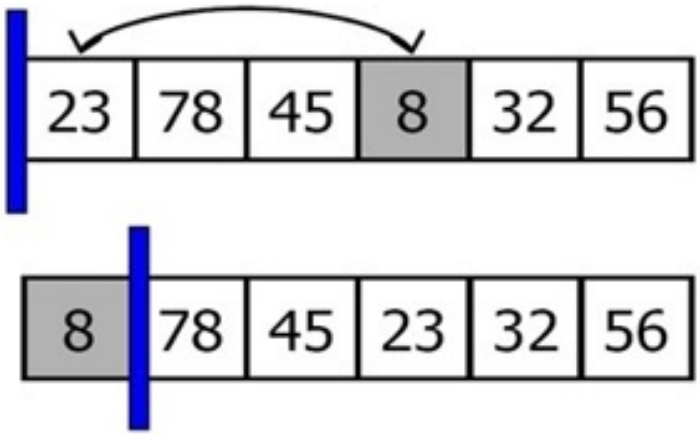
- The list is divided into two parts: sorted and unsorted.
- In each pass, in the unsorted sublist, the smallest element is selected and exchanged with the first element.



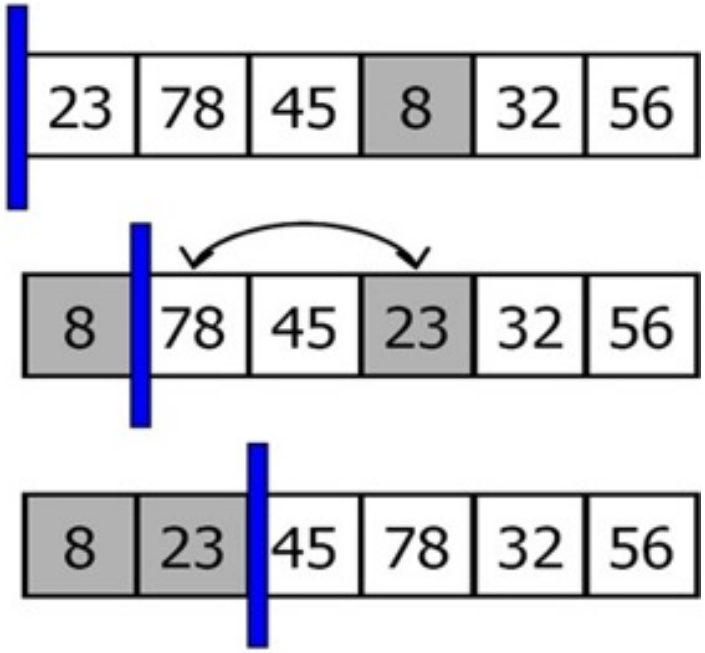
Straight Selection Sort



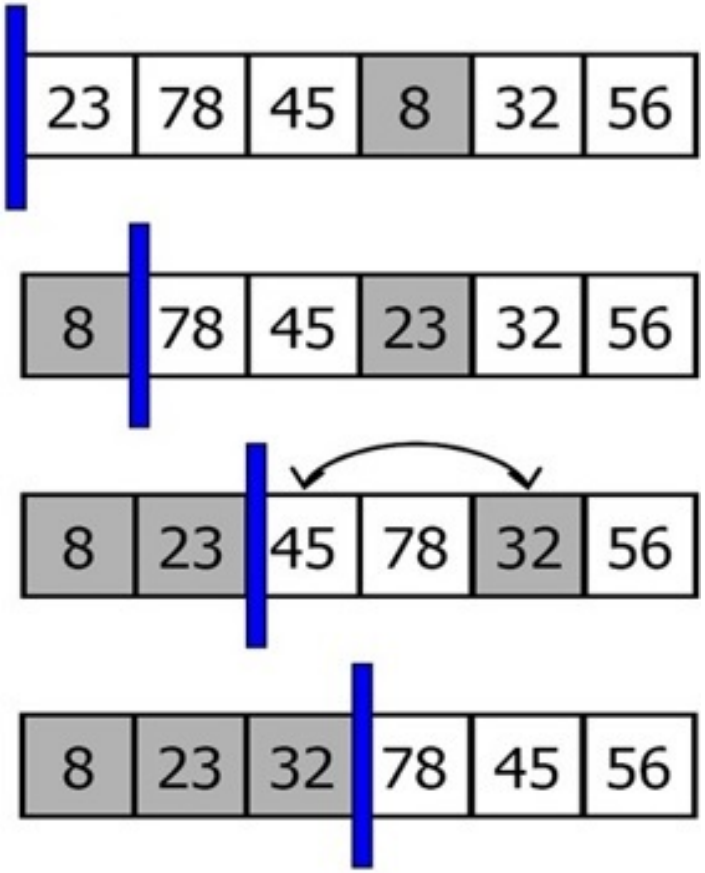
Straight Selection Sort



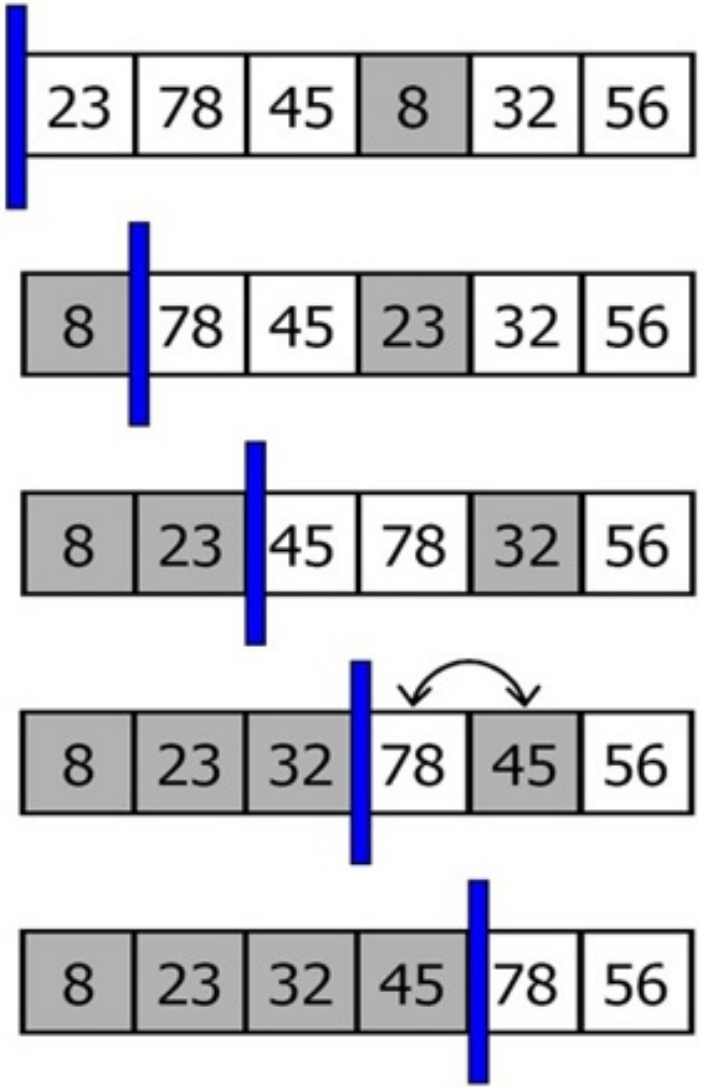
Straight Selection Sort



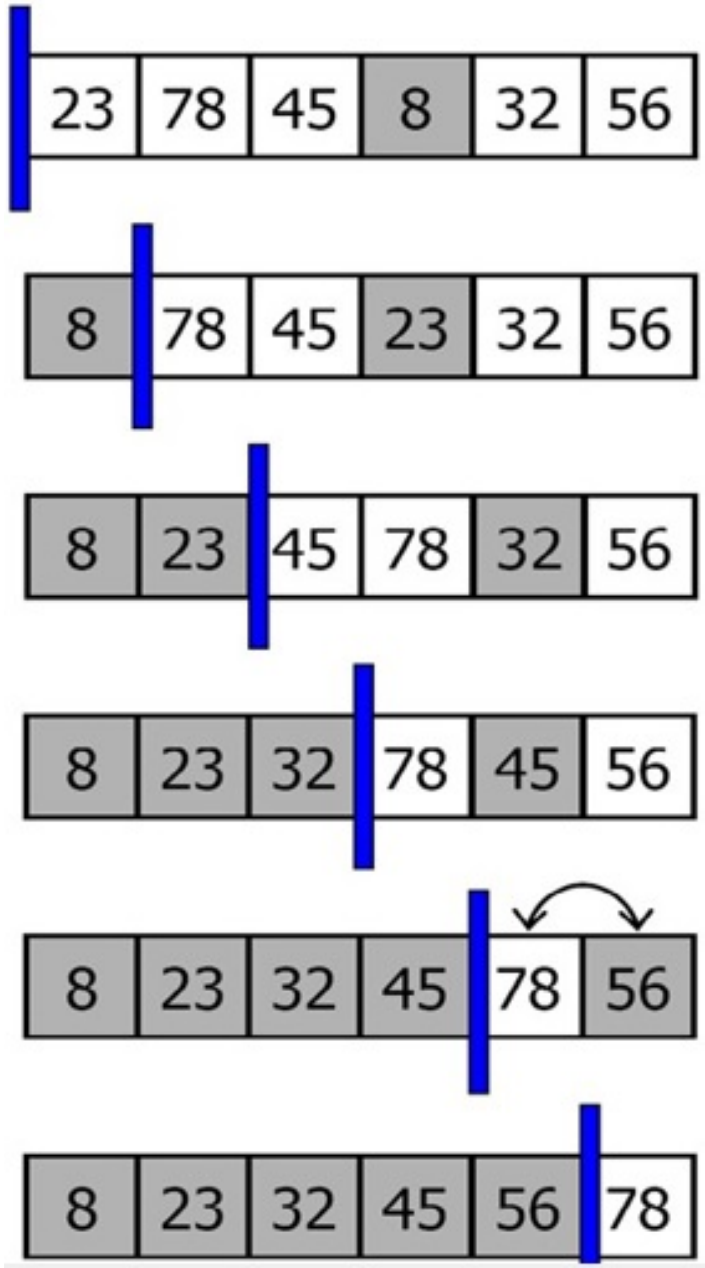
Straight Selection Sort



Straight Selection Sort



Straight Selection Sort



Straight Selection Sort

```
1  Algorithm SelectionSort()
2  Sorts the contiguous list using straight selection
     sort.

3  current = 0
4  while current < count - 1 do
5  |    smallest = current
6  |    walker = current + 1
7  |    while walker < count do
8  |    |    if data [walker].key < data [smallest].key
   |    |      then
9  |    |    |    smallest = walker
10 |    |    end
11 |    |    walker = walker + 1
12 |    end
13 |    swap(current, smallest)
14 |    current = current + 1
15 end
```


Selection Sort Efficiency

- Straight selection sort: $O(n^2)$

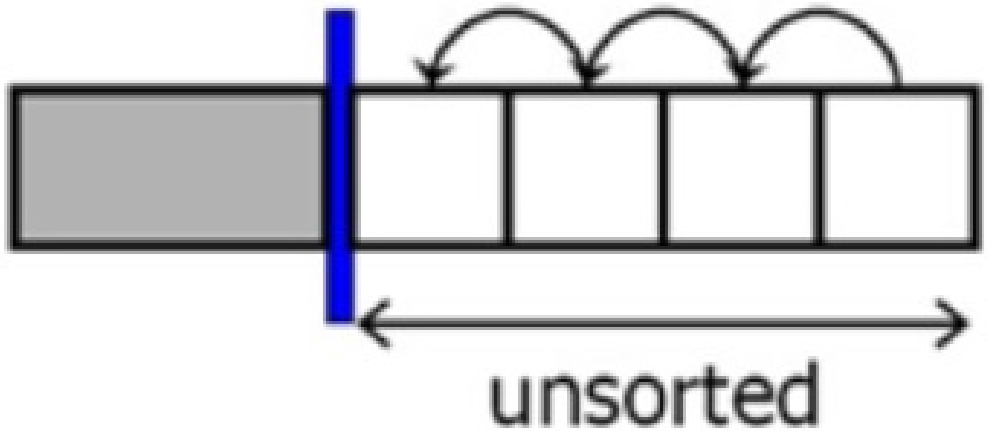
Exchange Sort

Exchange Sort

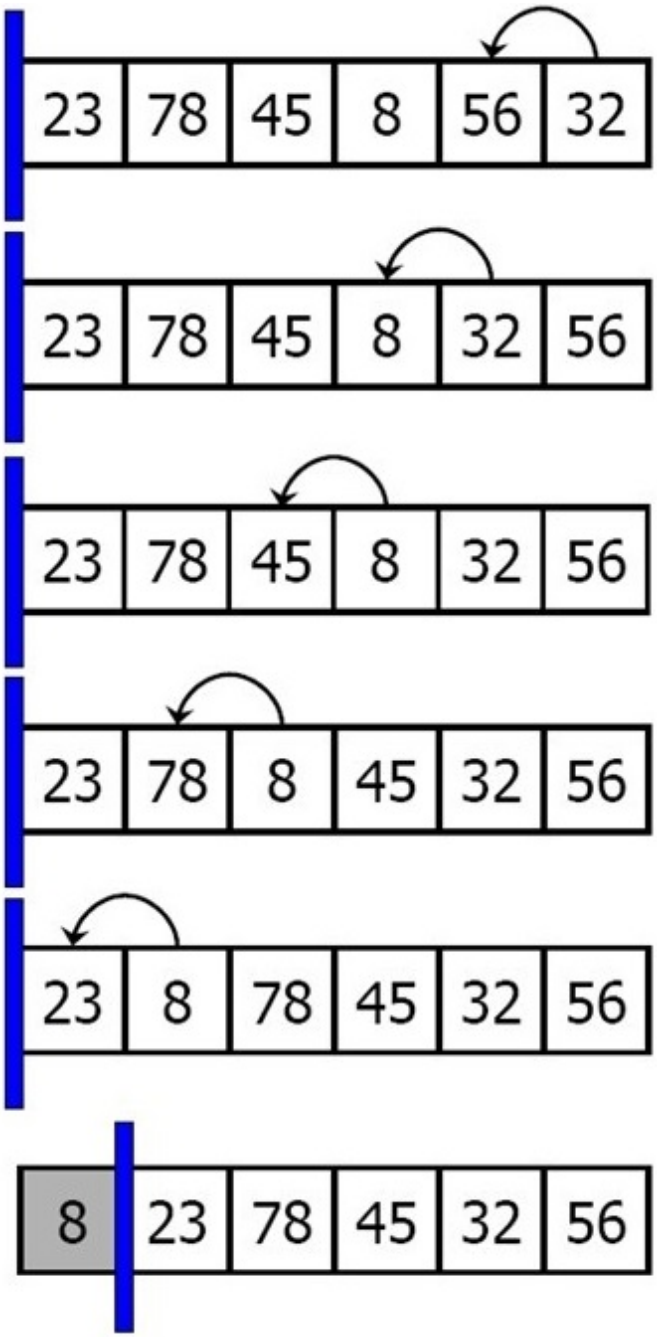
- In each pass, elements that are out of order are exchanged, until the entire list is sorted.
- Exchange is extensively used.

Bubble Sort

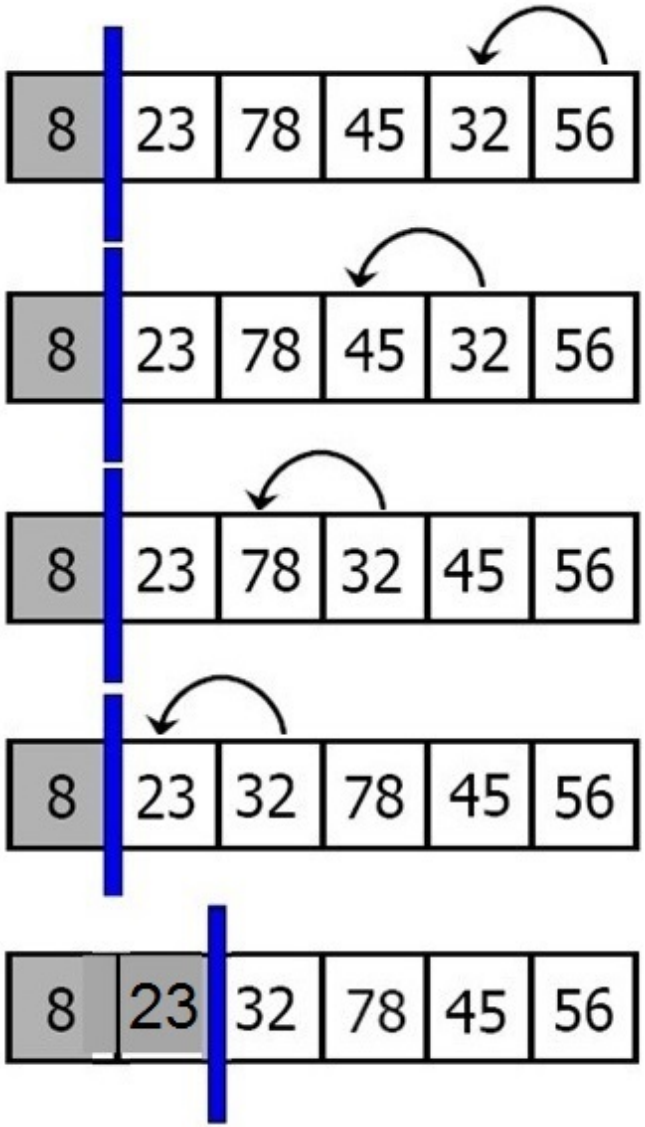
- The list is divided into two parts: sorted and unsorted.
- In each pass, the smallest element is bubbled from the unsorted sublist and moved to the sorted sublist.



Bubble Sort



Bubble Sort



Bubble Sort

```
1  Algorithm BubbleSort()
2  Sorts the contiguous list using bubble sort.

3  current = 0
4  flag = False
5  while current < count AND flag = False do
6  |   walker = count - 1
7  |   flag = True
8  |   while walker > current do
9  |   |   if data [walker].key < data [walker-1].key
   |   |     then
10 |   |   |   flag = False
11 |   |   |   swap(walker, walker - 1)
12 |   |   end
13 |   |   walker = walker - 1
14 |   end
15 |   current = current + 1
16 end
```


Exchange Sort Efficiency

- Bubble sort:
  $f(n) = n(n+1)/2 = O(n^2)$

Devide-and-Conquer

Devide-and-Conquer Sort

```
1 Algorithm DevideAndConquer()
2 if the list has length > 1 then
3 |   partition the list into lowlist and
  |     highlist
4 |   lowlist.DevideAndConquer()
5 |   highlist.DevideAndConquer()
6 |   combine(lowlist, highlist)
7 end
8 End DevideAndConquer
```


Devide-and-Conquer Sort

	Partition	Combine
Merge Sort	easy	hard
Quick Sort	hard	easy

Quick Sort

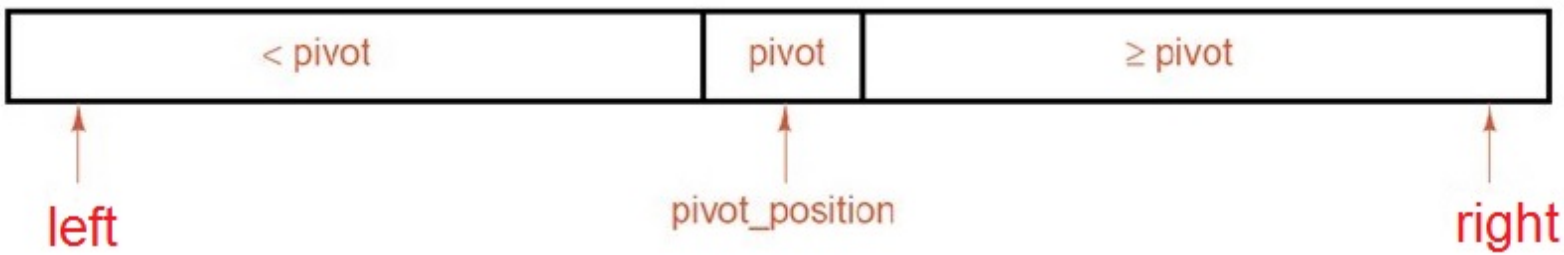
```
1 Algorithm QuickSort()
2 Sorts the contiguous list using quick sort.

3 recursiveQuickSort(0, count - 1)
4 End QuickSort
```


Quick Sort

```
1 Algorithm recursiveQuickSort(val left
    <int>, val right <int>)
2 Sorts the contiguous list using quick sort.
3 Pre: left and right are valid positions
    in the list
4 Post: list sorted

5 if left < right then
6 |   pivot_position = Partition(left, right)
7 |   recursiveQuickSort(left,
  |     pivot_position - 1)
8 |   recursiveQuickSort(pivot_position +
  |     1, right)
9 end
```


Quick Sort

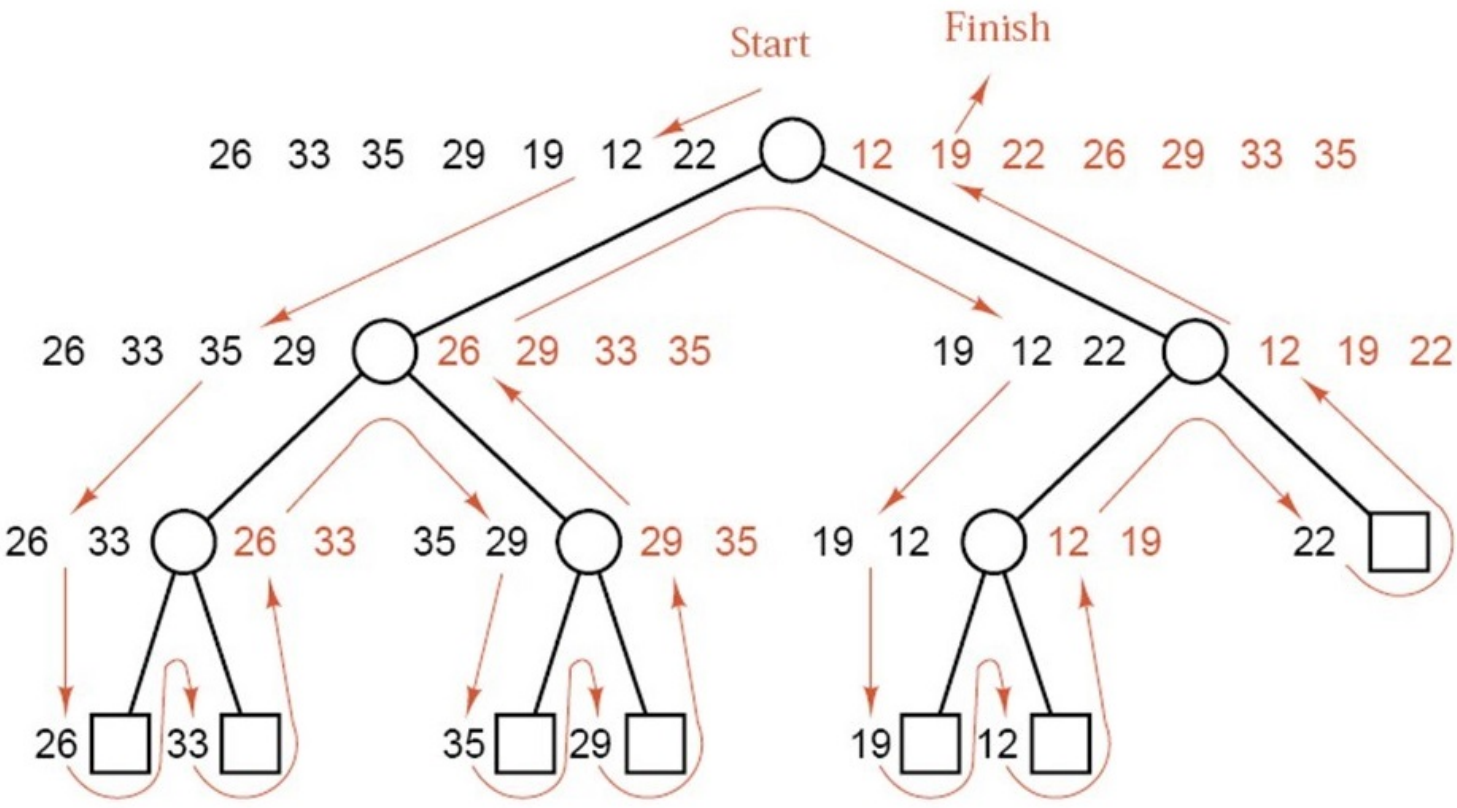
Given a pivot value, the partition rearranges the entries in the list as the following figure:



Quick Sort Efficiency

- Quick sort: $O(nlog_2n)$

Merge Sort



Merge Sort

```
1 Algorithm MergeSort()
2 Sorts the linked list using merge sort.

3 recursiveMergeSort(head)
4 End MergeSort
```

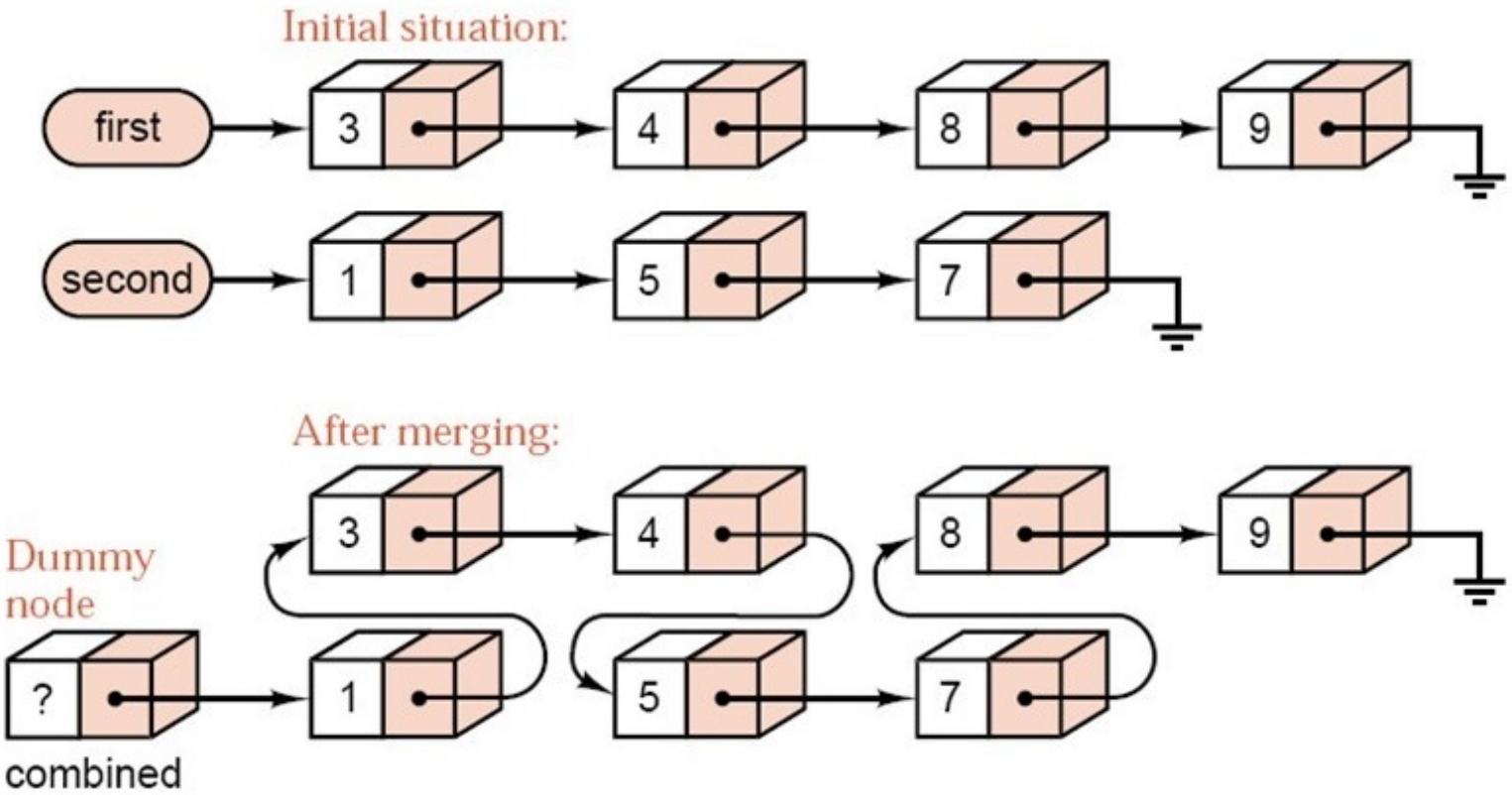

Merge Sort

```
1 Algorithm recursiveMergeSort(ref sublist <pointer>)
2 Sorts the linked list using recursive merge sort.

3 if sublist is not NULL AND sublist->link is not NULL then
4     Divide(sublist, second_list)
5     recursiveMergeSort(sublist)
6     recursiveMergeSort(second_list)
7     Merge(sublist, second_list)
8 end
9 End recursiveMergeSort
```


Merge Sort

```
1  Algorithm Divide(val sublist <pointer>, ref
     second_list <pointer>)
2  Divides the list into two halves.

3  midpoint = sublist
4  position = sublist->link
5  while position is not NULL do
6  |    position = position->link
7  |    if position is not NULL then
8  |    |    midpoint = midpoint->link
9  |    |    position = position->link
10 |    end
11 end
12 second_list = midpoint->link
13 midpoint->link = NULL
14 End Divide
```


Merge two sublists



Merge two sublists

```
1  Algorithm Merge(ref first <pointer>, ref second
     <pointer>)
2  Merges two sorted lists into a sorted list.

3  lastSorted = address of combined
4  while first is not NULL AND second is not NULL
     do
5  |    if first->data.key <= second->data.key then
6  |    |    lastSorted->link = first
7  |    |    lastSorted = first
8  |    |    first = first->link
9  |    else
10 |    |    lastSorted->link = second
11 |    |    lastSorted = second
12 |    |    second = second->link
13 |    end
14 end
```


Merge two sublists

```
1 // ...

2 if first is NULL then
3 |   lastSorted->link = second
4 |   second = NULL
5 else
6 |   lastSorted->link = first
7 end
8 first = combined.link
9 End Merge
```


THANK YOU.