

4. Home Exercises

Efficient Algorithms

Ex. 1 (2 + 2 points)

Suppose you shuffle n pieces of paper labeled with $1, \dots, n$ in an urn and draw one after another without replacement.

- a) What is the probability that the number drawn in the i -th round is larger than all previous numbers? Explain your answer. (*Hint: You can also use backward analysis. Fix the largest number to be the one drawn in the i -th iteration. Then, calculate the probability that all the numbers drawn before it are smaller than this.*)
- b) Compute the expectation of the number of rounds in which you draw a number that is larger than all previous numbers.

Ex. 2 (4 points)

In the analysis of the edge contraction algorithm it was used that

$$\Pr\left(\bigcap_{j=1}^i \{e_j \notin C\}\right) = \prod_{j=1}^i \Pr(e_j \notin C \mid C \subseteq E_{j-1}),$$

see Section 2.4.3. This is a consequence of the chain rule of probability theory, which states that for random variables X_1, X_2, \dots, X_n we have that

$$\Pr(X_1 \cap X_2 \cap \dots \cap X_n) = \Pr(X_1) \Pr(X_2 \mid X_1) \cdots \Pr(X_n \mid X_1, \dots, X_{n-1}).$$

Prove the chain rule by induction.

Ex. 3 (4 points)

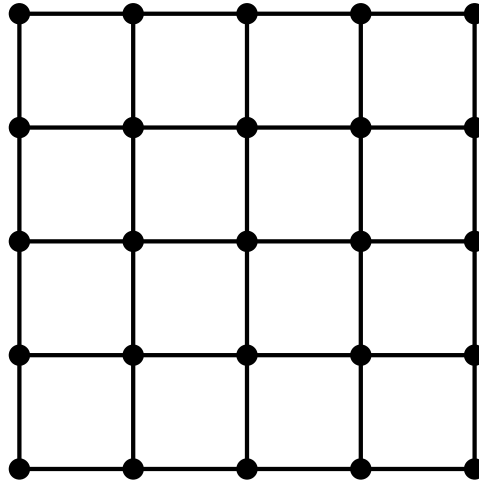
Recall that a spanning tree of a graph $G = (V, E)$ is a subgraph without cycles that connects all vertices of G . In this exercise, we call a spanning tree α -good if it is an α -spanner of G .

Give an example of a graph for which there exists no α -good spanning tree with $\alpha \leq |V|/2$. Explain your answer.

Ex. 4 ($4 + 2 + 2$ points)

In this exercise we will look at *grid graphs*. The vertex set of a grid graph G of size n^2 is the set of points in the plane with coordinates (i, j) for $1 \leq i \leq n$ and $1 \leq j \leq n$. The edges of G are given by $\{(i, j), (i + 1, j)\}$ and $\{(i, j), (i, j + 1)\}$. Note that the number of edges of a grid graph with n^2 vertices is $(2 - o(1))n^2$, where the $o(1)$ factor is due to the vertices in the outermost layer of the graph having degree less than 4.

An example of a grid graph with 5^2 vertices is given in the figure below.



- a) Design an algorithm that computes a spanner of a grid graph. The number of edges of the spanner should be at most $(\frac{3}{2} + o(1))n^2$ (even fewer is of course also okay) and the stretch factor should be constant, i.e., it should not depend on n .
- b) Show that the number of edges of your spanner is indeed at most $(\frac{3}{2} + o(1))n^2$.
- c) Show that the stretch factor of your spanner is constant.