

Exercise 3 - Efficient Algorithms

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Ex. 1

Let X be a random variable describing the number of 1-sequences with $l = 1 + \log_2 n$. Let X_i be a indicator random variable describing whether an 1-sequence begins at index/toss i or not.

Since the coin gets tossed $n + \log_2 n$ times and an 1-sequence is $1 + \log_2 n$ long, the highest index/toss, where an 1-sequence can start is n . Therefore

$$X = \sum_{i=1}^n X_i \quad \text{and} \quad E(X) = \sum_{i=1}^n E(X_i)$$
$$= \sum_{i=1}^n \Pr(X_i)$$

The probability to get L consecutive of either heads or tails is

$$2 \cdot \left(\frac{1}{2}\right)^L = 2 \cdot \left(\frac{1}{2}\right)^{1+\log_2 n} = \left(\frac{1}{2}\right)^{\log_2 n} = \frac{1}{n}$$

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head- or
tail-sequence

$$\Rightarrow E(X) = \sum_{i=1}^n \frac{1}{n} = n \cdot \frac{1}{n} = 1$$

Ex. 2

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a)

Let $m \in \mathbb{N}$ be the constant desired size of S .
Let $\Pr: P \mapsto \mathbb{R}$ and $\omega: P \mapsto \mathbb{R}$ be functions.

1 weighted Set($P, d, C, OPT, m, \Pr, \omega$)

2 $S = \emptyset$

3 for every $p \in P$ do

4 $\Pr(p) = d(p, C) / OPT$

5 while $|S| < m$ do

6 $s =$ Select a $p \in P$ with probability
7 $\Pr(p)$, without replacement

8 $\omega(s) = OPT / (m \cdot d(s, C))$

9 $S = S \cup \{s\}$

10 return S, ω

- b)
- Line 1-2: $O(1)$
 - Line 3-4: $O(n)$
 - Line 5-9: $O(m)$ with constant $m \Rightarrow O(1)$
 - Line 10: $O(1)$

\Rightarrow running time is $O(n)$

- c) The bigger the distance $d(p, c)$ of a $p \in P$, the higher the probability

$$Pr(p) = \frac{d(p, c)}{\sum_{q \in P} d(q, c)} = \frac{d(p, c)}{OPT}$$

of being selected. The weight of a selected $s \in S$ behaves antiproportional with $d(s, c)$

$$\omega(s) = \frac{\sum_{p \in P} d(p, c)}{m \cdot d(s, c)} = \frac{OPT}{m \cdot d(s, c)}$$

Therefore the expected value of the contribution to the sum of a single point is

$$\mathbb{E}(\omega(s)d(s, c)) = \mathbb{E}\left(\frac{OPT}{|S| \cdot d(s, c)} \cdot d(s, c)\right)$$

$$= \mathbb{E}\left(\frac{OPT}{|S|}\right)$$

Since there is no more random variable contained

$$\mathbb{E}\left(\frac{\text{OPT}}{|S|}\right) = \frac{\text{OPT}}{|S|}$$

As a result we get

$$\begin{aligned}\mathbb{E}\left(\sum_{s \in S} w(s)d(s,c)\right) &= \sum_{s \in S} \mathbb{E}(w(s)d(s,c)) \\ &= \sum_{s \in S} \frac{\text{OPT}}{|S|} \\ &= \text{OPT}\end{aligned}$$

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