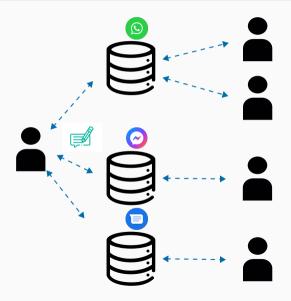
# Revisiting Link Prioritization for Efficient Traversal in Structured Decentralized Environments

Ruben Eschauzier, Ruben Taelman, Ruben Verborgh

October 27, 2025

Department of Electronics and Information Systems, Ghent University Imec

# The Need for Decentralized Personal Data Storage



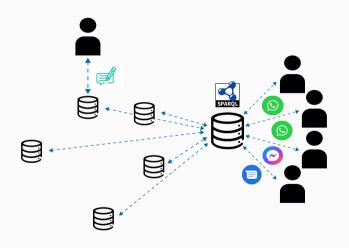
- Each application has its own data
- Stifles innovation
- Causes vendor lock-in

# The Need for Decentralized Personal Data Storage



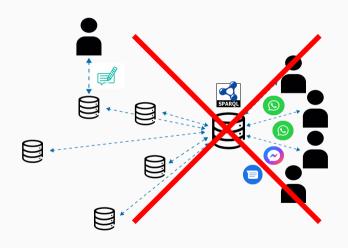
- Each application uses common data storage
- Easy to switch vendors
- Promotes innovation

# The Problem with Centrally Aggregating and Querying



- Why not aggregate data and query it?
- Difficult in case of personal data due to privacy concerns

# The Problem with Centrally Aggregating and Querying



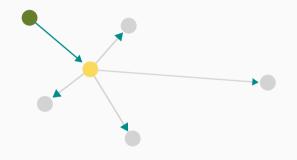
- Why not aggregate data and query it?
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#### Link Traversal: Seed Document



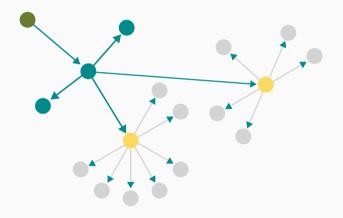
- Link Traversal starts from seed documents (URIs)
- These are provided by the user or in the query.

#### Link Traversal: Traversal



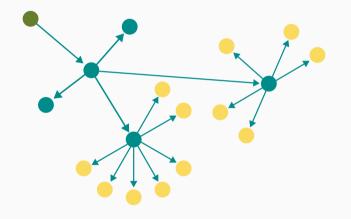
- New URIs are extracted from the seed document
- URIs are extracted in accordance with reachability criterions

## Link Traversal: Traversal



 New URIs are dereferenced and the process is repeated

#### **Link Traversal: Termination**



- This continues until all links are dereferenced
- Continuously produces results
- Can enforce fine-grained (document-level)
   access-control

# Query Optimization for Link Traversal

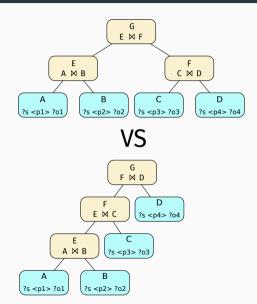
```
SELECT * WHERE {
 <seedUri> <ex:p1> ?o1.
 <seedUri> <ex:p2> ?o2.
 ...
            Extract
            Seed document
```

The query (partly) determines:

- The queried data
- The topology of the queried data
- The query-relevant documents

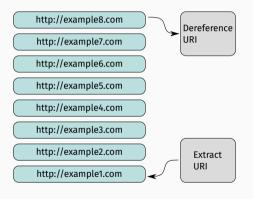
Result: limited prior knowledge for query optimization

# Query Optimization for Link Traversal: Traditional Query Planning



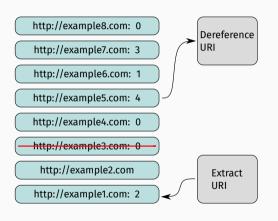
- Query optimization for link traversal involves traditional (zero knowledge) query planning
- Assume an optimize-then-execute

# Query Optimization for Link Traversal: Traversal optimization



- URIs are put into a link queue
- Link traversal uses a FiFo queue by default

# Query Optimization for Link Traversal: Traversal optimization



#### Optimizations:

- Prioritize query-relevant URIs
- Prune irrelevant URIs
- Pruning can lead to missing results without prior knowledge

#### **Problem Statement**

Investigate the performance of link prioritization algorithms in literature in new structured decentralized environments

#### **Problem Statement**

- These algorithms are implemented in an engine that is no longer maintained
- The baseline performance depends on design choices orthogonal to the prioritization algorithm
- The baseline should measure marginal prioritization performance

# R<sup>3</sup> Metric: Introduction

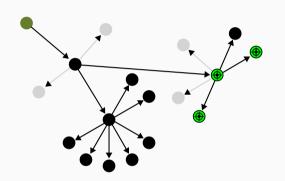
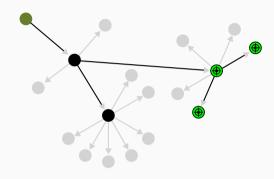


Figure 1: Actual traversal



**Figure 2:** Optimal traversal: minimal cost path

$$R^3 = \frac{5}{13} = 0.38$$

## $R^3$ metric: Definition

#### **Definition**

$$R^3 = \frac{C(X)}{C(O_T)}$$

- X: the **optimal traversal order** (minimal-cost path)
- ullet  $O_{\mathcal{T}}$ : traversal order produced by the link prioritization algorithm
- $C(\cdot)$ : The total cost of all arcs in the traversal path

#### Interpretation

- $R^3 = 1$  perfect match with the optimal traversal
- $R^3 < 1$  deviation from optimal performance
- Higher values indicate better prioritization quality.

## $R^3$ metric: Steiner Trees

- Directed graph steiner tree over traversed directed graph G = (V, A) with root r
- lacksquare Query-relevant documents  $D_{\mathcal{T}}$  serve as terminals  $\mathcal{T}\subseteq V$
- Find minimum cost sub-graph X = (V', A') starting at root r and spanning all vertexes T.
- With cost:  $C(X) = \sum_{a \in A'} c(a)$ , with c(a) the cost of an edge in the topology

# **Investigated Structured Decentralized Environment**

#### Simulated Solid Environment: SolidBench (Taelman and Verborgh 2023)

- Data is stored in vaults
- Vaults use document-centric structure
- Vaults simulate a social media application

#### **Data attributes**

- Generates 158,233 RDF files
- Across 1,531 data vaults
- containing a total of 3,556,159 triples.

#### **Analysed queries**

Use discover queries

# Prioritization algorithms (Hartig and Özsu 2016)

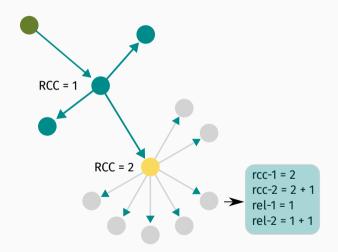
#### Non-adaptive

Breadth-first (default), depth-first, random prioritization

## **Graph-based**

■ In-degree, PageRank score

# Prioritization algorithms: Result Contribution-Based (RCC)



- Nodes adaptively scored according to the number of results they contribute to
- Called *rcc-1*, *rcc-2*, *rel-1*, *rel-2*.

# Prioritization algorithms: Intermediate Results-based

```
SELECT ?person ?project ?topic
WHERE {
 ?person ex:worksFor ?company.
 ?person ex:worksOn ?project .
 ?project ex:hasTopic ?topic .
                                   ?person: p1
                                   ?project: pri1
       http://t1.com
                       0
                                   ?topic: http://t3.com
       http://t2.com
                       0
       http://t3.com
                        2
                                   Dereference
                                   URI
       http://t4.com
                       0
```

- Nodes adaptively scored according to their contribution to intermediate results
- IS with initial priorities of 0, while ISdcr sets priority to the priority of the parent node - 1

# Prioritization algorithms

## Hybrid

- Multiply intermediate result and RCC-based scoring functions
- *is-rcc1*, *is-rcc2*, *is-rel1*, *is-rel2*

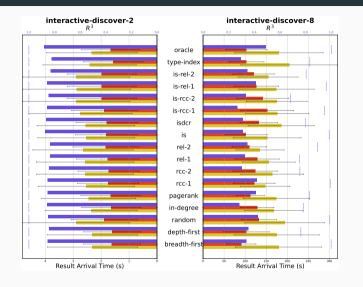
## **TypeIndex**

- TypeIndex points to location for resource of specific type
- Prioritize TypeIndex

#### **Oracle**

- Compute RCC in hindsight
- Scores are propegated through the shortest path
- Serves as optimal performance oracle

# Prioritization Has Limited Impact on Result Arrival Time



- Breadth-first outperforms most algorithms
- Oracle has significantly better R<sup>3</sup>, but not execution time

# Prioritization Has Limited Impact on Result Arrival Time

	1st		Cmpl		$R^3$	
	better	worse	better	worse	better	worse
depth-first	<u>15.7</u>	<u>17.1</u>	<u>14.3</u>	17.1	7.5	12.5
random	4.3	68.6	5.7	74.3	<u>15.0</u>	15.0
is-rcc-1	4.3	57.1	5.7	60.0	7.5	5.0
type-index	<u>15.7</u>	18.6	10.0	<u>15.7</u>	12.5	20.0
oracle	<u>18.6</u>	<u>11.4</u>	<u>18.6</u>	12.9	<u>25.0</u>	2.5

**Table 1:** Percentage of queries for which an algorithm performs at least 10% better or worse than the breadth-first baseline.

• No algorithm improves performance, and the oracle only marginally improves performance

## Link Prioritization in Structured Decentralized Environments

- Link prioritization in Solid environment will not significantly improve query performance
- Research should instead focus on query planning Hanski et al. (2025) or traversal pruning Tam et al. (2024)

#### References i

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- Tam, B.-E. et al. (2024). "Opportunities for Shape-based Optimization of Link Traversal Queries". In: arXiv preprint arXiv:2407.00998.