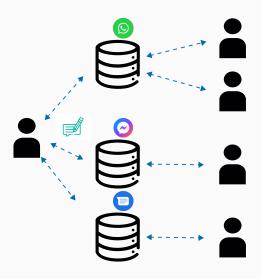
Revisiting Link Prioritization for Efficient Traversal in Structured Decentralized Environments

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The Need for Decentralized Personal Data Storage



- Each application has its own data
- Stiffles innovation
- Causes vendor lock-in

The Need for Decentralized Personal Data Storage



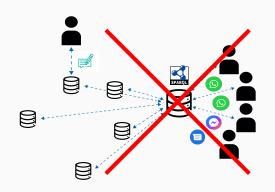
- Each application uses common data storage
- Easy to switch vendors
- Promotes innovation

The Problem with Centrally Aggregating and Querying



- Why not aggregate data and query it?
- Impossible in case of personal data due to privacy concerns

The Problem with Centrally Aggregating and Querying

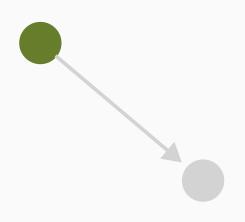


- Why not aggregate data and query it?
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Link Traversal-based Query Processing

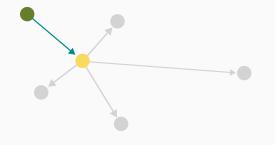
- Link Traversal iteratively dereferences data to query over
- Continuously produces results
- Can enforce fine-grained (document-level) access-control

Link Traversal: Seed Document



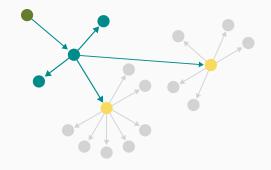
- Link Traversal starts from seed documents (URIs)
- These are provided by the user or in the query.

Link Traversal: Traversal



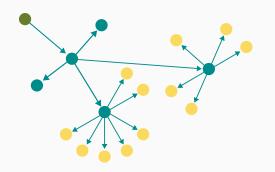
- New URIs are extracted from the seed document
- URIs are extracted in accordance with reachability criterions

Link Traversal: Traversal



 New URIs are dereferenced and the process is repeated

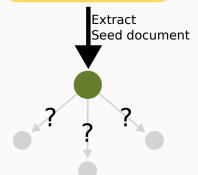
Link Traversal: Termination



 This continues untill all links are dereferenced

Query Optimization for Link Traversal

```
SELECT * WHERE {
    <seedUri> <ex:p1> ?o1.
    <seedUri> <ex:p2> ?o2.
    ...
}
```

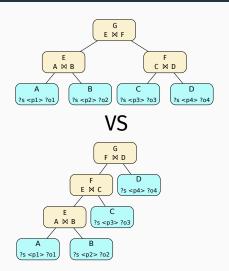


The query (partly) determines:

- The queried data
- The topology of the queried data
- The query-relevant documents

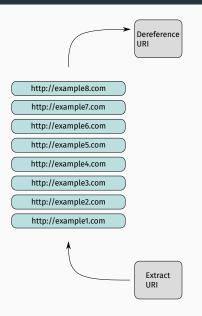
Result: limited prior knowledge for query optimization

Query Optimization for Link Traversal: Traditional Query Planning



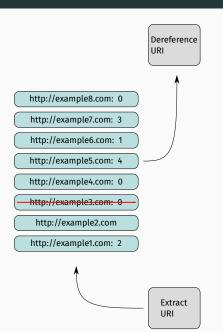
 Query optimization for link traversal involves traditional (zero knowledge) query planning

Query Optimization for Link Traversal: Traversal optimization



- URIs are put into a link queue
- Link traversal uses a FiFo queue by default

Query Optimization for Link Traversal: Traversal optimization



Optimizations:

- Prioritize query-relevant URIs
- Prune irrelevant URIs
- Pruning can lead to missing results without prior knowledge

Problem Statement

Investigate the performance of link prioritization algorithms in literature in new structured decentralized environments

Problem Statement

- These algorithms are implemented in an engine that is no longer maintained
- The baseline performance depends on design choices orthogonal to the prioritization algorithm
- The baseline should measure marginal prioritization performance

R^3 metric: Motivation

Context

During link traversal, the engine traverses a fixed topology entirely

Goal

The goal of prioritization is to find query-relevant (where-provenance) documents as soon as possible.

Challenge

Find the **optimal traversal order** of a given query in hindsight to compare the traversal order taken by the engine to.

R^3 metric: Steiner trees

- The problem of finding the optimal traversal order can be translated to solving the directed graph steiner tree problem.
- The traversed topology admits directed graph G = (V, A) with root r
- Query-relevant documents D_T serve as terminals $T \subseteq V$
- Find minimum cost sub-graph X = (V', A') starting at root r and spanning all vertexes T.
- With cost: $C(X) = \sum_{a \in A'} c(a)$,, with c(a) the cost of an edge in the topology

R^3 metric

Definition

$$R^3 = \frac{C(X)}{C(O_T)}$$

- X: the optimal traversal order (minimal-cost path)
- O_T : traversal order produced by the link prioritization algorithm
- $C(\cdot)$: The total cost of all arcs in the traversal path

Interpretation

- $lacksquare R^3=1
 ightarrow {
 m perfect}$ match with the optimal traversal
- lacksquare $R^3 < 1
 ightarrow$ deviation from optimal performance
- Higher values indicate better prioritization quality.

R³ Metric Illustration

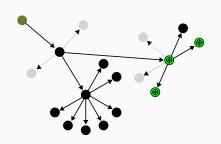


Figure 1: Actual traversal

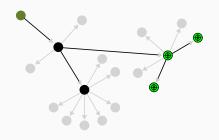


Figure 2: Optimal traversal: minimal cost path

$$R^3 = \frac{5}{13} = 0.38$$

Prioritization algorithms

Non-adaptive

Breadth-first (default), depth-first, random prioritization

Graph-based

In-degree, PageRank score

Result-based

- Uses result contribution count (RCC) of each node (URI) in the topology
- Priority set based on the sum or count of non-zero RCC of the 1 or 2-hop in-neighbours of a node
- Called rcc-1, rcc-2, rel-1, rel-2 respectively.

Prioritization algorithms

Intermediate-results

- Uses intermediate solutions in the engine
- Priority of URI is determined by the largest intermediate result with that URI bound
- IS sets initial priorities to 0, while ISdcr sets priority to the priority of the parent node - 1

Hybrid

- Multiply intermediate and full result scoring functions
- is-rcc1, is-rcc2, is-rel1, is-rel2

TypeIndex

- TypeIndex points to location for resource of specific type
- Prioritize TypeIndex

Prioritization algorithms

Oracle

- Compute RCC in hindsight
- Scores are propegated through the shortest path
- Serves as optimal performance oracle