

Algorithms, Week 2, Lec 1

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1 General Analysis Strategy

1. $T(n)$: Maximum line taken by algorithm to solve any input of size n
2. $T(n)$: Measure of goodness, how good or bad indicated by function
3. The function will indicate how good is a algorithm
4. Conservative Definition - *Worst Case*.

What is $T(n)$? It indicates the maximum time it would take for a machine to run that algorithm - the worst case scenario?

1. Form of $T(n)$ (independent from machine)
2. "Linear", "Cubic", "Quadratic" etc
3. Bounds of $T(n)$, upper bound, lower bound
4. Large n is important, as n becomes larger and larger which algorithm is better

2 Running Time Analysis

1. The running time depends upon the input size, e.g., n
2. Different inputs of the same size may result in different running time

3. Criteria for measuring running time
4. Worst-Case Time (Maximum running time over legal input of size n)
5. Criteria Worst-case time
6. Let I denote an input instance
7. Let $|I|$ denote its length
8. Let $T(I)$ denote the running time of algorithm on input I

Some measurements are as follows,

1. Average Case Time - the average running time over all inputs of size n
2. Let $P(I)$ denote the probability of seeing this input
3. Average case time is the weighted sum of running times with weights being the probabilities

2.1 Example: 2-Dimension Maxima

The car selection problem can be modelled this way: For each car we associate a (x, y) pair where,

1. x is the speed of the car
2. y is the negation of the price

2.1.1 Algorithm

```
MAXIMA(int n, Point P[1 ... n])
for i ← 1 to n do maximal ← true:
  for j ← 1 to n do
    if (i ≠ j) and (P[i].x ≤ P[j].x) and (P[i].y ≤ P[j].y) then
      maximal ← false;
  break;
if maximal == true continue;
```

2.1.2 Analysis

To do *Analysis*, just count the steps,