Data Structure

資料結構

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什麼是資料結構

- 電腦組織、儲存資料的方式
- 對於不同用途選擇不同的資料結構(演算法、複雜度)
- 資料結構是一種工具
- 排序、分類、紀錄、分析

Abstract

- Computing Complexity
- Array
- Dynamic Array
- Linked List
- Stack
- Queue
- Tree

Computing Complexity

計算複雜度

- 如何表示電腦的計算速度?哪一種演算法比較好?
- 衡量程式的效率 (時間、空間)
- 做N次動作 花了N時間
- Ex: $N^3 + 2*N^2 + N \Rightarrow O(N^3)$
- O(n), O(n^2), O(log(N))表示計算量 常數用O(1)表示

Ex:如果輸出"讓我看看"的時間是k秒 那我就需要O(n^2)的時間 k*n^2秒

Array

陣列

- 一次儲存多個變數
- 優點:儲存快、拿取也快都是O(1)
- 缺點:大小被限制、刪除速度慢、插入速度慢 O(n)

```
int a[5] = {0,1,2,3,4};

a[0] = 99;

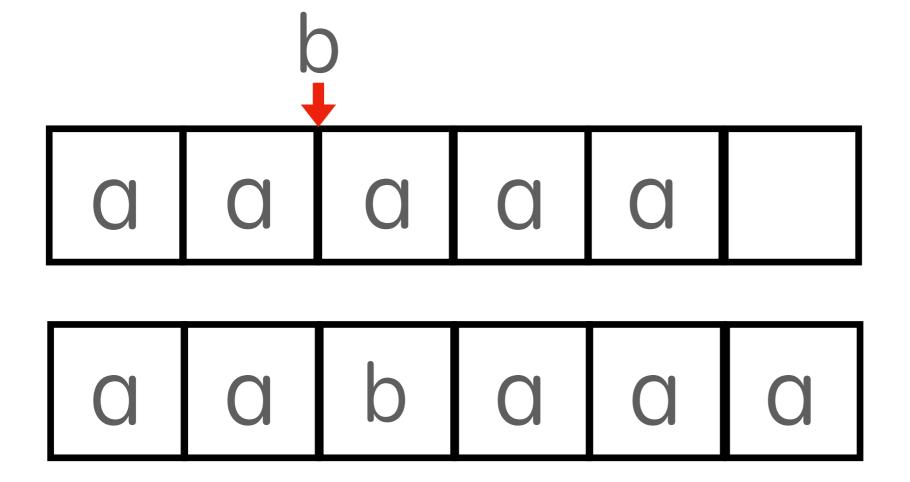
for(int i=0;i < 5;i++){
    cout<<a[i]<<" ";
}

#output: 99 1 2 3 4
```

Question

問題大了

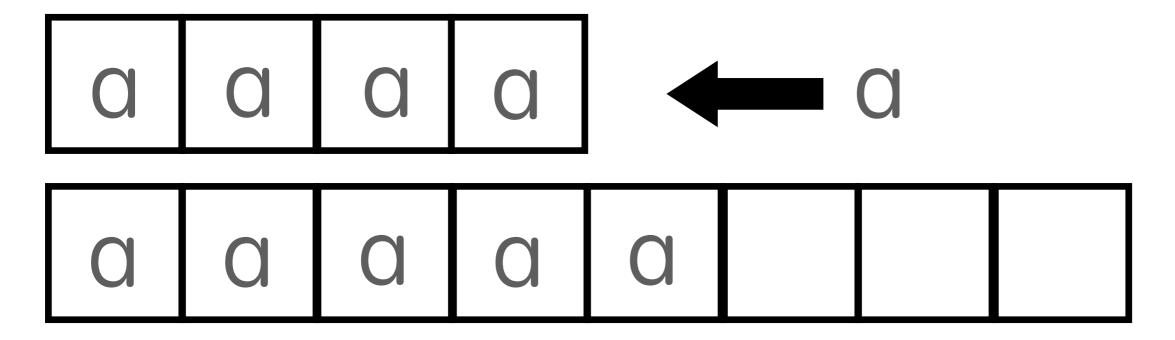
- 如何解決陣列大小被限制的問題?
- 如何解決插入速度很慢的問題?



Dynamic Array

動態陣列

- 可伸縮陣列
- 記憶體空間以2的冪次增加(不是無限增加)
- 優點:繼承陣列拿取快、存取快的優點,解決了大小限制問題
- 缺點:插入問題還是沒有解決



擴增兩倍!

Dynamic Array

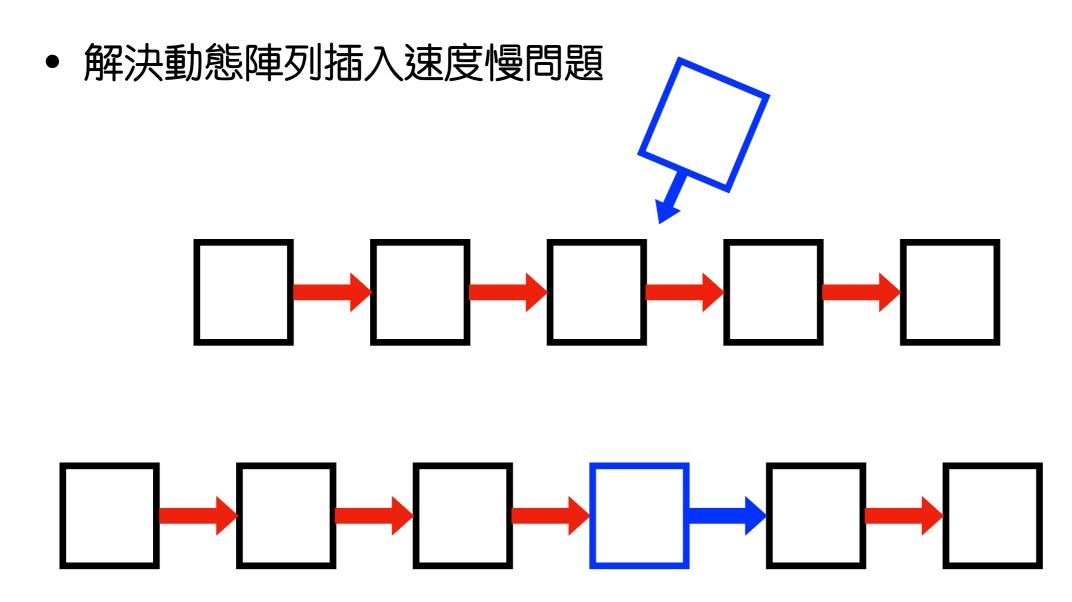
動態陣列

- member function
 - size()
 - capacity()
 - resize()
 - push()
 - pop()
 - del()
 - erase()
 - find()

Linked List

鏈結串列

• 以一個節點(Node)為單位,每個節點指向下一個節點



Linked List

Advantage and Disadvantage

- 插入、刪除 O(1)
- 但是取值、尋找 O(n) =>失去陣列結構的優點

```
1    struct node
2    {
3         int data;
4         node<int> *next;
5    };
6
7    list()
8    {
9         head = new node<T>;
10         tail = new node<T>;
11         tail->next = NULL;
12    }
```

沒有一個資料結構是完美的視不同的需求和算法就有不同的變化

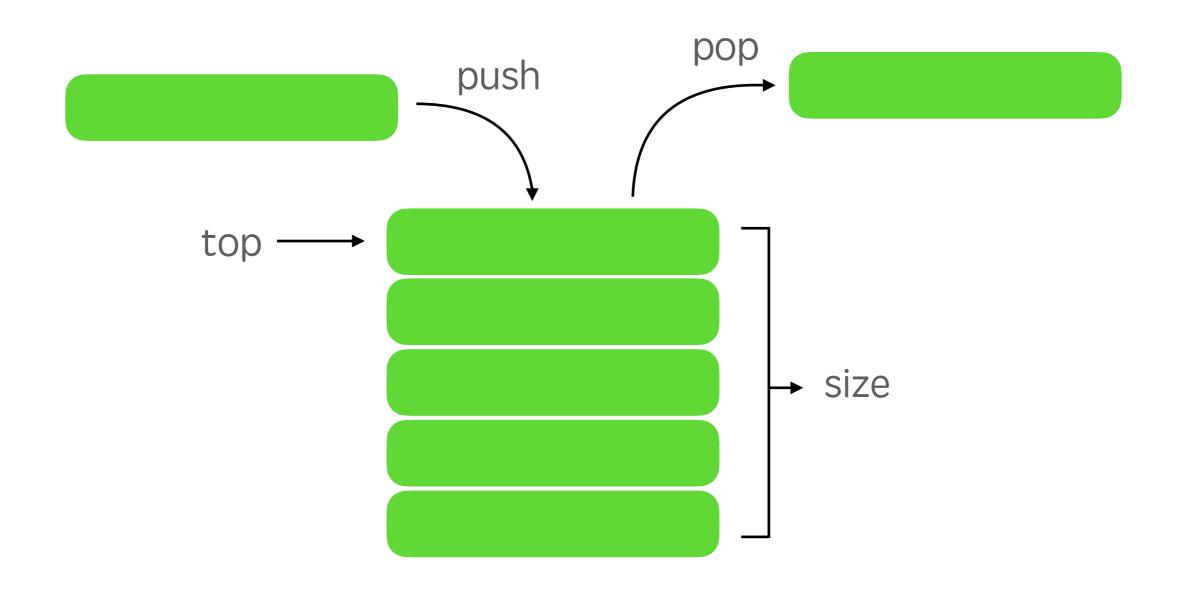
linked list

鏈結串列

- member function
 - begin()
 - end()
 - size()
 - empty()
 - push_front()
 - push_back()
 - pop_front()
 - pop_end()
 - insert()
 - del()
 - get()

Stack

堆疊



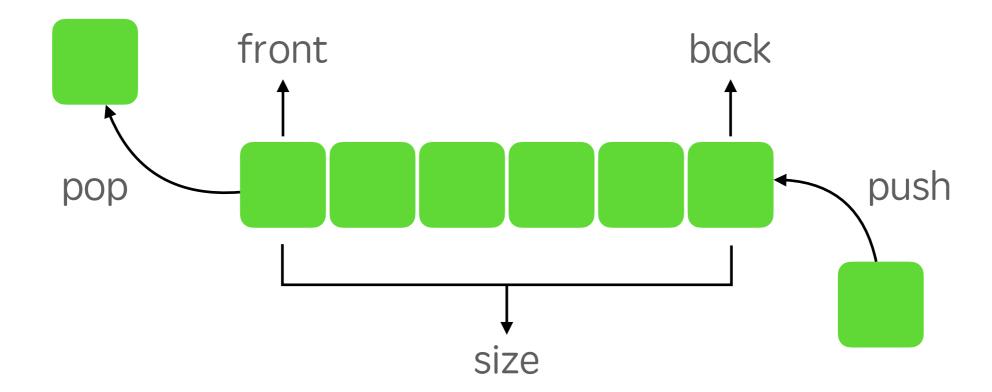
Stack

堆疊

- 後進先出
- 用 linked_list 實作
- member_function
 - size()
 - empty()
 - push()
 - pop()
 - top()

Queue

佇列

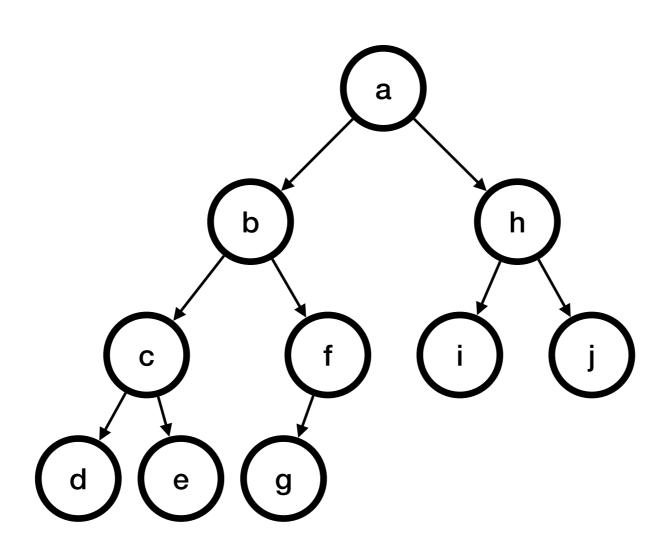


Queue 佇列

- 先進先出
- 用 linked_list 實作
- member function
 - size()
 - push()
 - pop()
 - empty()
 - front()
 - back()



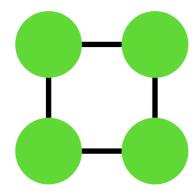
樹





樹

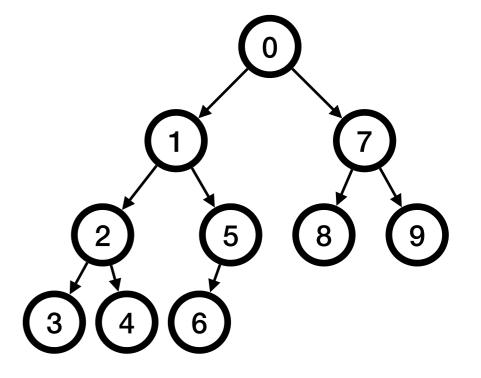
- 只有一個root,且不存在cycle
 - 對於所有的node有存在一條路到root
 - 每個node只有一個parent
- 檔案儲存方式是一種樹狀結構
- 族譜是一種樹狀結構
- 儲存方式
 - node
 - vector
 - matrix





儲存_matrix

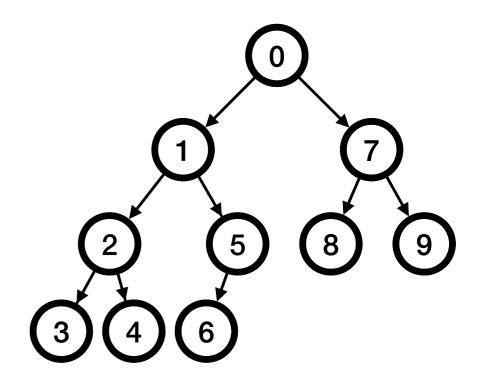
	0	1	2	3	4	5	6	7	8	9
0	F	Т	F	F	F	F	F	Т	F	F
1	F	F	Т	F	F	Т	F	F	F	F
2	F	F	F	Т	Т	F	F	F	F	F
3	F	F	F	F	F	F	F	F	F	F
4	F	F	F	F	F	F	F	F	F	F
5	F	F	F	F	F	F	Т	F	F	F
6	F	F	F	F	F	F	F	F	F	F
7	F	F	F	F	F	F	F	F	Т	Т
8	F	F	F	F	F	F	F	F	F	F
9	F	F	F	F	F	F	F	F	F	F





儲存_vector

data	child				
0	1	7			
1	2	5			
2	3	4			
3					
4					
5	6				
6					
7	8	9			
8					
9					





```
class node
        int *child_one;
3
        int *child_two;
        int data;
5
    };
```



儲存_node

```
class node
{
  int *child_one;
  int *child_two;
  int data;
};

class node
{
  int *child_one;
  int data;
}

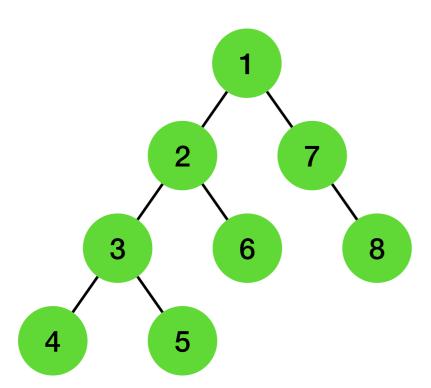
d
e
g
```



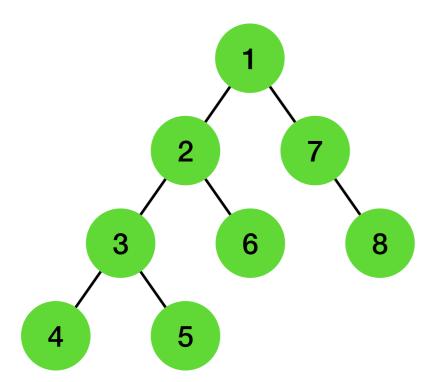
遍歷

- depth-first-search (connect with stack)
 - pre-order-traversal
 - in-order-traversal
 - post-order-traversal
- breadth-first-search (connect with queue)
 - level-order-traversal

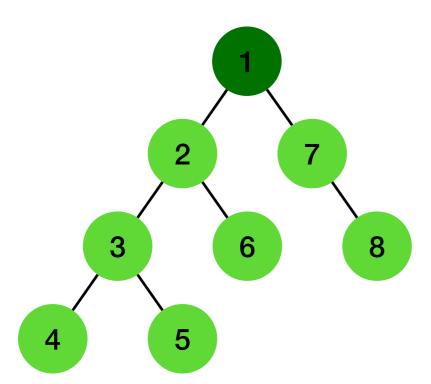




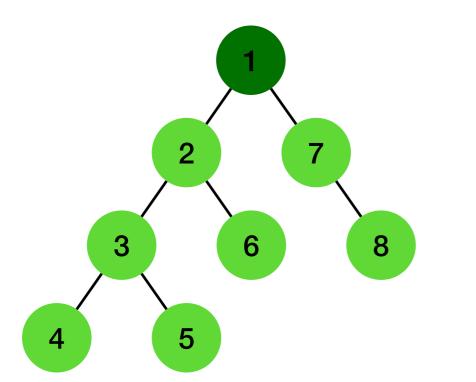




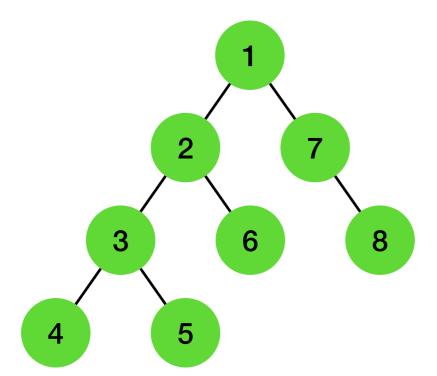




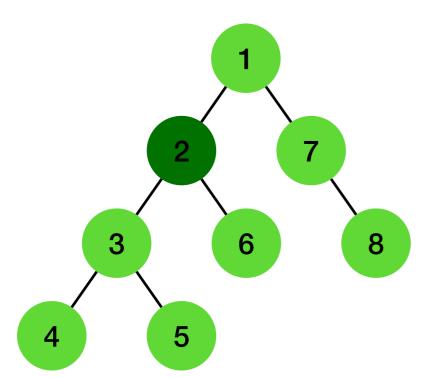




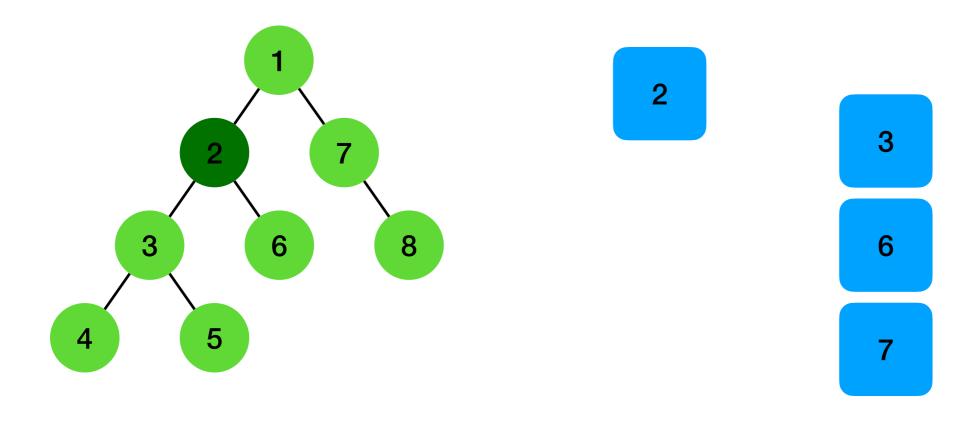




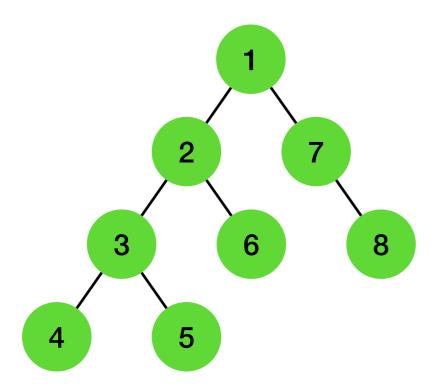




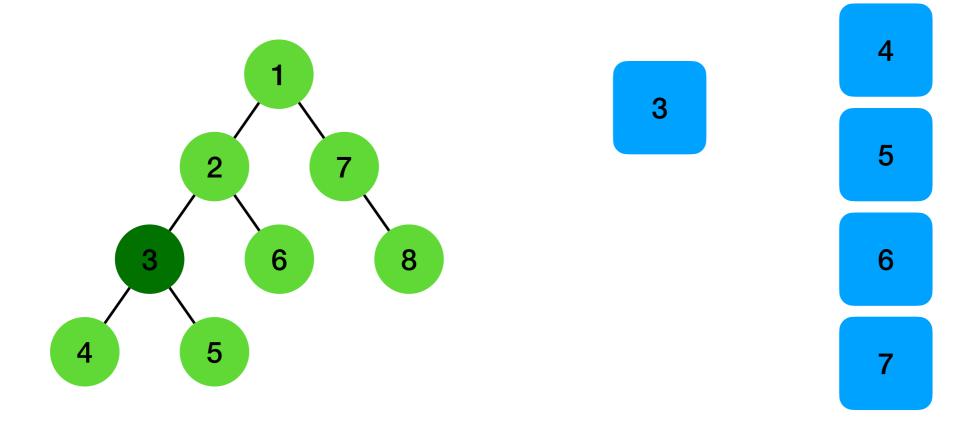




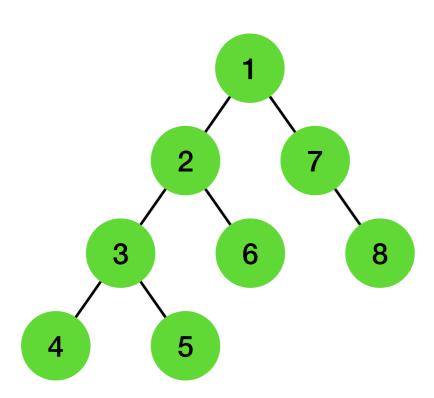




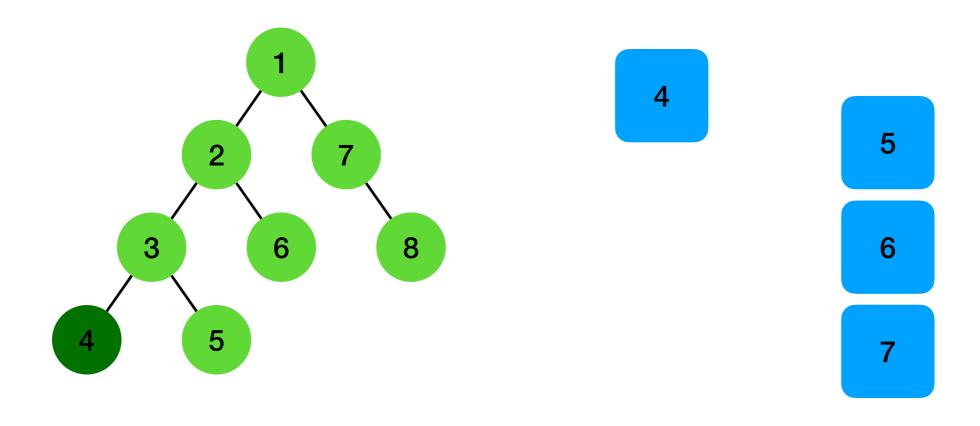




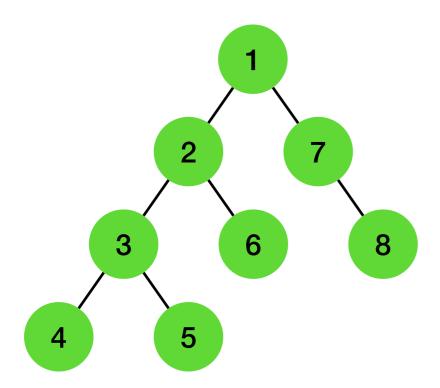




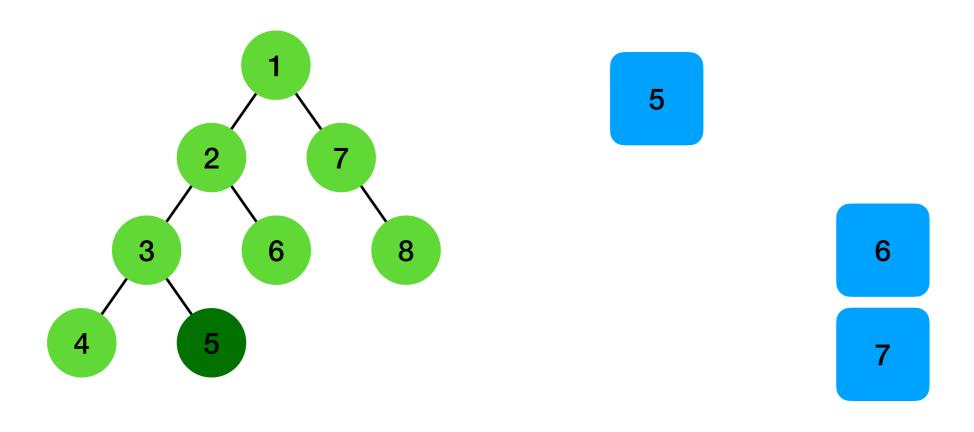




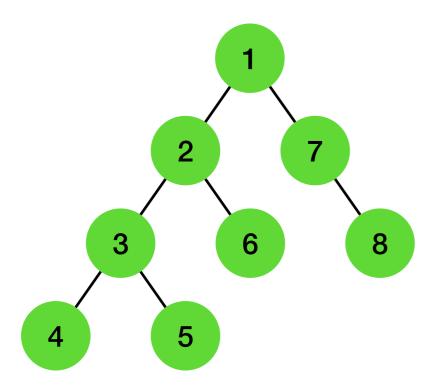




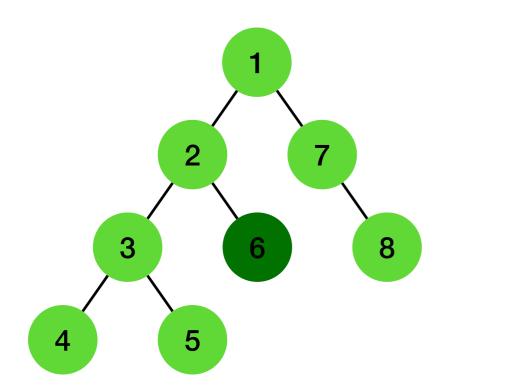




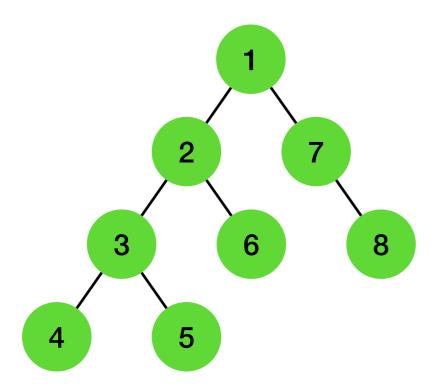




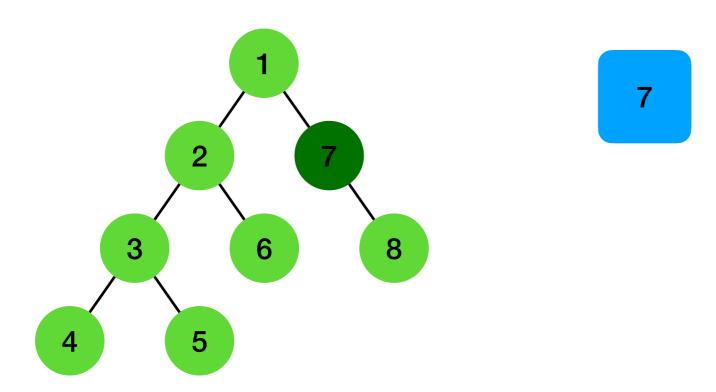




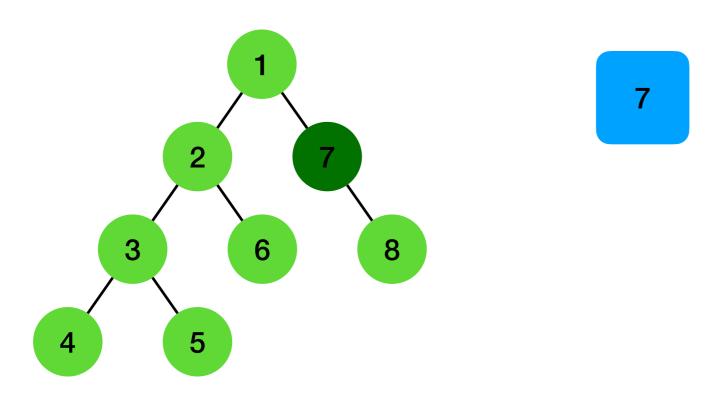






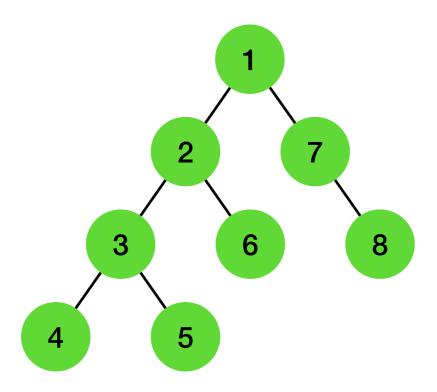




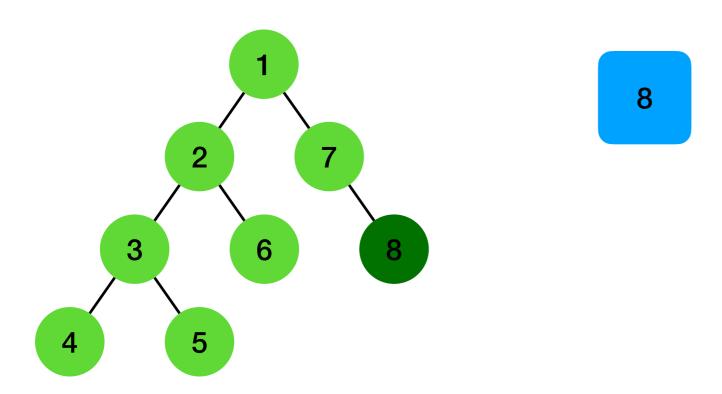


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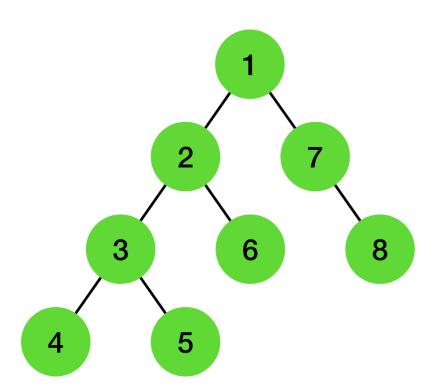




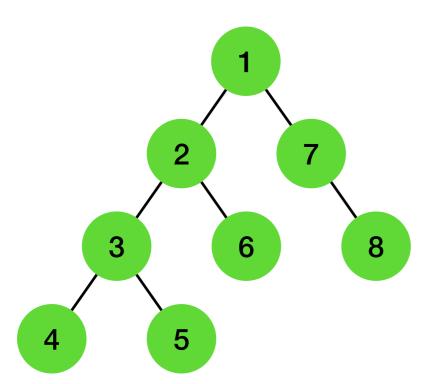




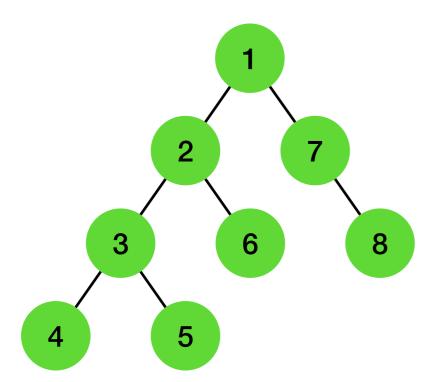




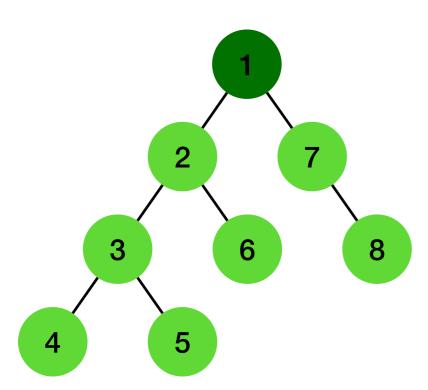




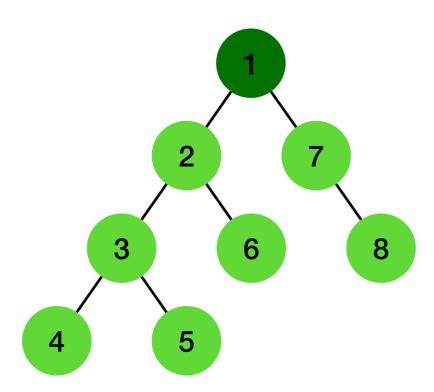




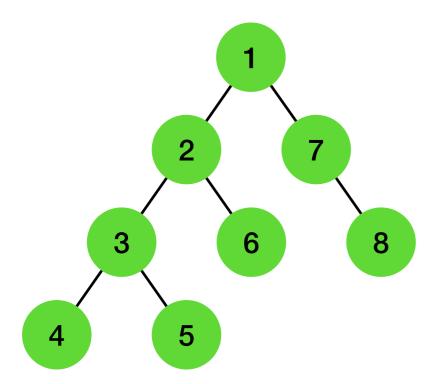




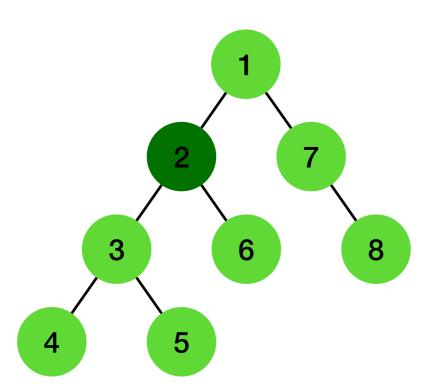




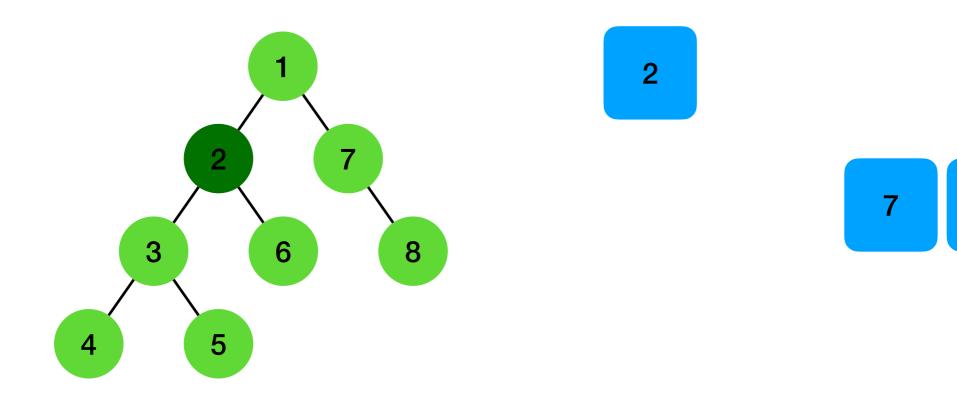




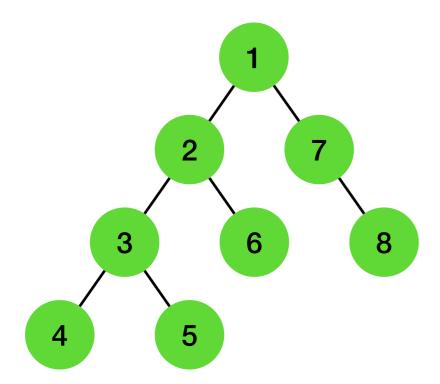






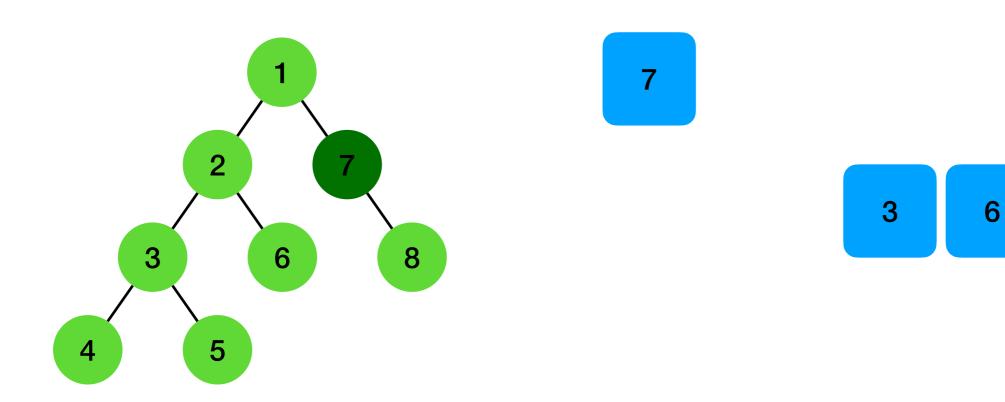




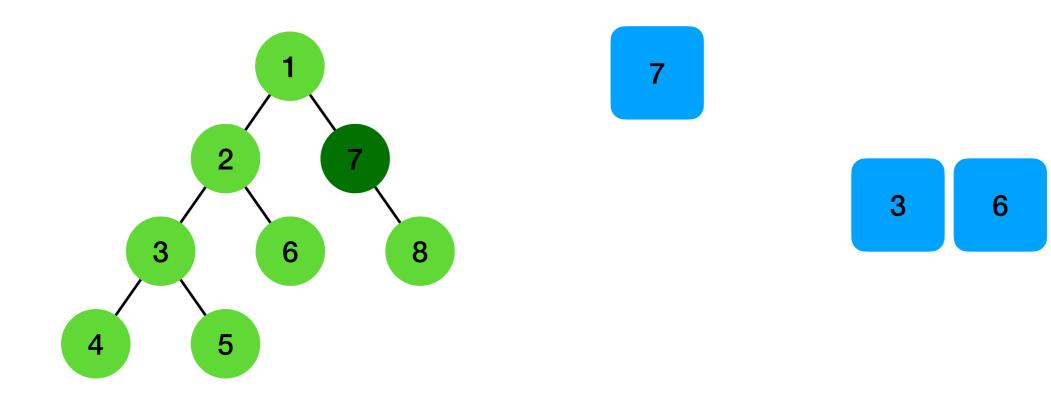


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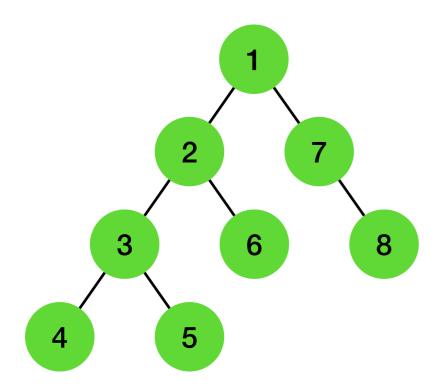






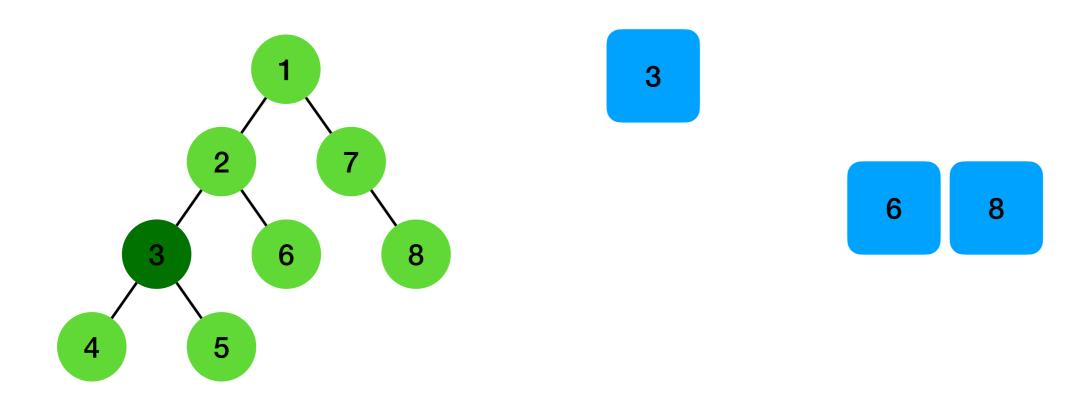




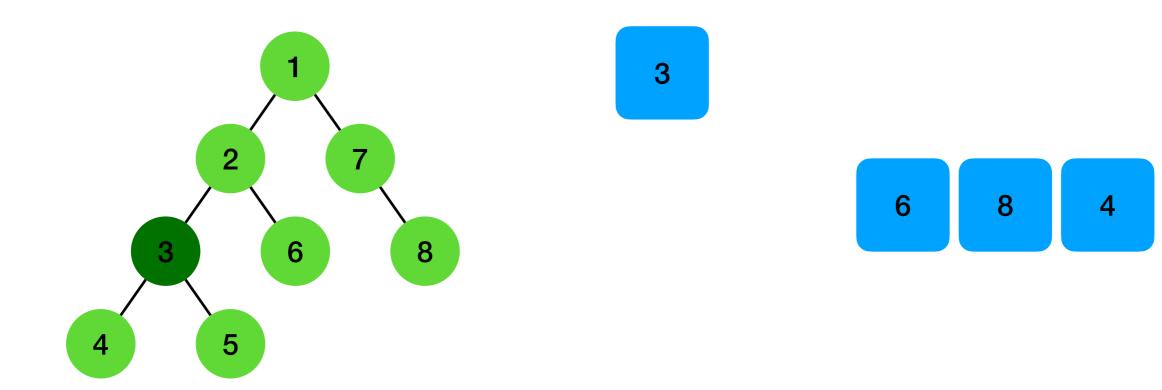


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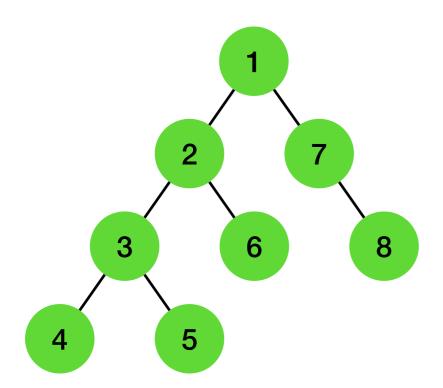






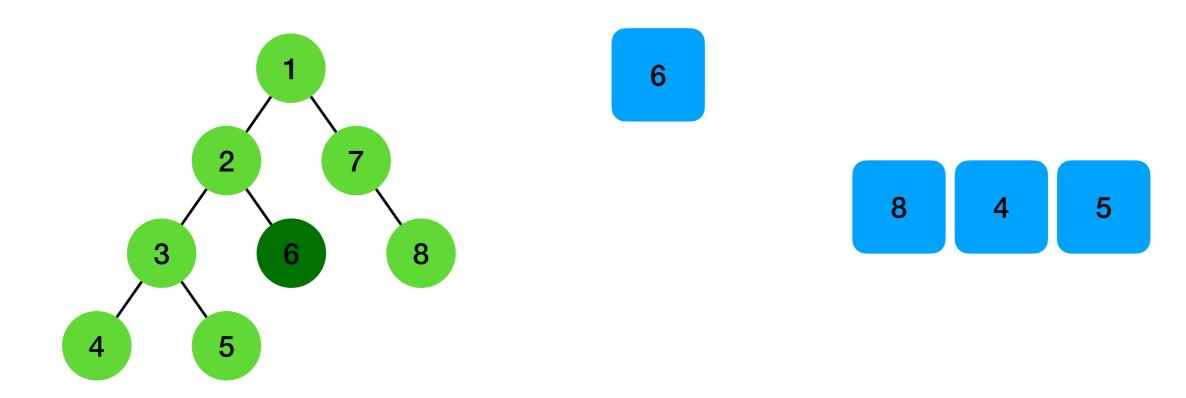




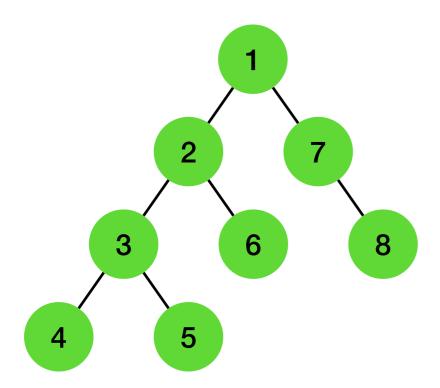


6 8 4 5



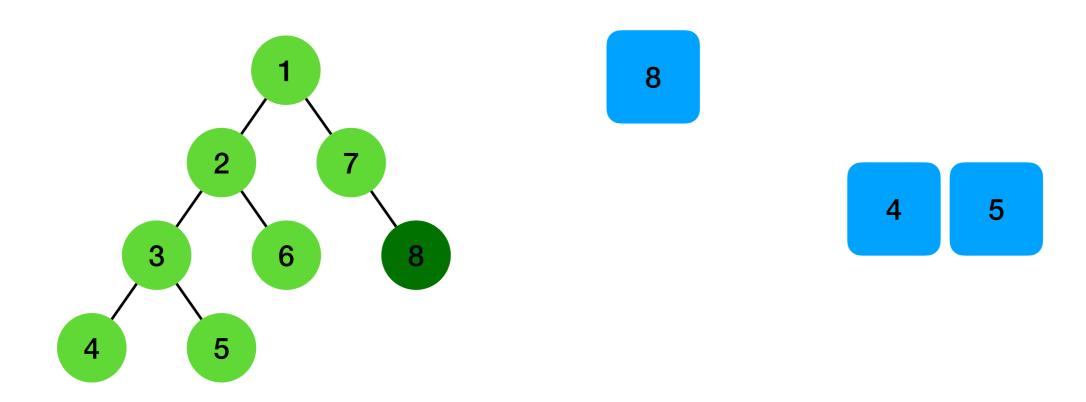




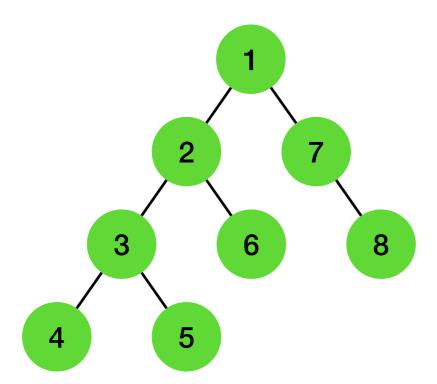


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