Chapter 03

IS5306: Numerical Solution of System of Linear Equations

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Introduction

Numerical methods for solving systems of linear equations are divided into two categories:

- Direct methods
 Direct methods, such as Gaussian elimination and LU decomposition, provide exact solutions in a finite number of steps but may become impractical for very large systems due to computational and memory constraints.
- Iterative methods
 On the other hand, iterative methods, like the Jacobi and Gauss-Seidel methods, provide approximate solutions by successively refining an initial guess, making them suitable for large and sparse systems where direct methods may be too costly.

Linear System of Equations

A system of *n* linear equations with *n* variables can be represented as:

$$a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n = b_1,$$

 $a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n = b_2,$
 \vdots
 $a_{n1}x_1 + a_{n2}x_2 + \cdots + a_{nn}x_n = b_n,$

where:

- $ightharpoonup x_1, x_2, \dots, x_n$ are the *n* unknowns (variables),
- a_{ij} are the coefficients of the system, with *i* representing the equation number and *j* the variable's index,
- ▶ $b_1, b_2, ..., b_n$ are the constants on the right-hand side of the equations.

A linear system can be transformed into a new system that is easier to solve, while maintaining the same solutions.

During this process, it is not necessary to rewrite the full equations or track the variables x_1, x_2, \ldots, x_n , as long as they stay in the same columns. Only the coefficients and the constants on the right-hand side change.

This is why we often represent a linear system with a matrix, which provides all the necessary information in a compact and computer-friendly form.

Matrices and vectors

A matrix is a rectangular array of numbers or other mathematical objects arranged in rows and columns. It has the general form:

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\begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{m1} & a_{m2} & \cdots & a_{mn} \end{pmatrix}
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- The size of a matrix is defined by the number of rows and columns it contains. A matrix with *m* rows and *n* columns is called an *m* × *n* matrix, or *m*-by-*n* matrix, where *m* and *n* are called its dimensions.
- The individual items in a matrix are called its elements or entries.
- ▶ A square matrix is an $n \times n$ matrix. The main or principal diagonal of a square matrix consists of the elements a_{ii} , where i = 1, 2, ..., n, running from the upper left to the lower right.
- An $1 \times n$ matrix, such as $\mathbf{y} = (y_1 \ y_2 \ \cdots \ y_n)$, is called an n-dimensional row vector. Similarly, an $n \times 1$ matrix,

such as
$$\mathbf{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix}$$
, is called an *n*-dimensional column vector.

Special Types of Square Matrices

Diagonal Matrix:

$$D = \begin{pmatrix} d_1 & 0 & 0 & \cdots & 0 \\ 0 & d_2 & 0 & \cdots & 0 \\ 0 & 0 & d_3 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & d_n \end{pmatrix}$$

Identity Matrix:

$$I = \begin{pmatrix} 1 & 0 & 0 & \cdots & 0 \\ 0 & 1 & 0 & \cdots & 0 \\ 0 & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & 1 \end{pmatrix}$$

Upper Triangular Matrix:

$$A = \begin{pmatrix} a_{11} & a_{12} & a_{13} & \cdots & a_{1n} \\ 0 & a_{22} & a_{23} & \cdots & a_{2n} \\ 0 & 0 & a_{33} & \cdots & a_{3n} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & a_{nn} \end{pmatrix}$$

Lower Triangular Matrix:

$$L = \begin{pmatrix} a_{11} & 0 & 0 & \cdots & 0 \\ a_{21} & a_{22} & 0 & \cdots & 0 \\ a_{31} & a_{32} & a_{33} & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & a_{n3} & \cdots & a_{nn} \end{pmatrix}$$

Matrix form of a linear system

A system of *n* linear equations with *n* variables can be represented as:

$$a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n = b_1,$$

 $a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n = b_2,$
 \vdots
 $a_{n1}x_1 + a_{n2}x_2 + \cdots + a_{nn}x_n = b_n,$

In matrix form, this system can be written as:

$$A\mathbf{x} = \mathbf{b},$$

where:

▶ A is an $n \times n$ coefficient matrix:

$$A = \begin{pmatrix} a_{11} & a_{12} & \cdots & a_{1n} \\ a_{21} & a_{22} & \cdots & a_{2n} \\ \vdots & \vdots & \ddots & \vdots \\ a_{n1} & a_{n2} & \cdots & a_{nn} \end{pmatrix},$$

x is the vector of unknowns:

$$\mathbf{x} = \begin{pmatrix} x_1 \\ x_2 \\ \vdots \\ x_n \end{pmatrix},$$

b is the vector of constants:

$$\mathbf{b} = egin{pmatrix} b_1 \ b_2 \ dots \ b_n \end{pmatrix}.$$

The Elimination Method

The following three operations are called the elementary row operations:

1. Row Scaling: Multiplying any row R_i of a matrix by a non-zero number k. This operation is denoted by:

$$kR_i \rightarrow R_i$$

2. **Row Addition:** Adding a multiple of a row R_i to another row R_i . This operation is denoted by:

$$R_i + kR_i \rightarrow R_i$$

3. **Row Interchange:** Interchanging the order of two rows R_i and R_i . This operation is denoted by:

$$R_i \leftrightarrow R_j$$

Example 1: Let

$$A = \begin{pmatrix} 4 & 5 & 6 \\ 7 & 8 & 9 \\ 1 & 2 & 3 \end{pmatrix}$$

We can apply the following elementary row operations to the matrix *A*.

- 1. Multiply the second row of A by 2 and obtain a new matrix B.
- Add −2 times the first row to the third row and obtain a new matrix C.
- 3. Interchange the order of the second and third rows and obtain a new matrix *D*:

Solving linear systems of Equations

Numerical methods for the solution of systems of linear equations are of two types:

1. Direct Methods

- Gaussian Elimination Method with Backward Substitution
- LU Factorization

2. Iterative Methods

- Jacobi Method
- Gauss-Seidel Method

Direct Methods

Direct techniques are methods that gives the exact answers to the system in a finite number of steps.

Gaussian Elimination with back substitution

Gaussian elimination is a method to solve a linear system $A\mathbf{x} = \mathbf{b}$ by transforming the matrix A into an upper triangular matrix using elementary row operations on the augmented matrix $[A|\mathbf{b}]$.

Remark: Applying elementary row operations to the augmented matrix $[A|\mathbf{b}]$ does not change the solution of the linear system $A\mathbf{x} = \mathbf{b}$.

Let A be an upper triangular matrix. Back-substitution is the process of solving the linear system $A\mathbf{x} = \mathbf{b}$ by finding the unknown x_n from the last equation, then substituting this value into the (n-1)th equation to find the unknown x_{n-1} , and so on, until all unknowns are determined.

Example 1: Use Gauss Elimination and back-substitution to solve the linear system.

$$x_1 + x_2 + 3x_4 = 4$$

$$2x_1 + x_2 - x_3 + x_4 = 1$$

$$3x_1 - x_2 - x_3 + 2x_4 = -3$$

$$-x_1 + 2x_2 + 3x_3 - x_4 = 4$$

We can write the above linear system in the form $A\mathbf{x} = \mathbf{b}$ as follows:

$$\begin{pmatrix} 1 & 1 & 0 & 3 \\ 2 & 1 & -1 & 1 \\ 3 & -1 & -1 & 2 \\ -1 & 2 & 3 & -1 \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{pmatrix} = \begin{pmatrix} 4 \\ 1 \\ -3 \\ 4 \end{pmatrix}$$

Next, we consider the augmented matrix for the system $[A|\mathbf{b}]$:

$$[A|\mathbf{b}] = \begin{pmatrix} 1 & 1 & 0 & 3 & | & 4 \\ 2 & 1 & -1 & 1 & | & 1 \\ 3 & -1 & -1 & 2 & | & -3 \\ -1 & 2 & 3 & -1 & | & 4 \end{pmatrix}$$

We will apply elementary row operations to transform the matrix into an upper triangular form.

Step 1: Row Operations

1. Subtract 2R₁ from R₂:

$$R_2 - 2R_1 \rightarrow R_2 \implies R_2 = (2 - 2 \cdot 1, 1 - 2 \cdot 1, -1 - 2 \cdot 0, 1 - 2 \cdot 3 | 1 - 2 \cdot 4)$$

Resulting in:

$$\begin{pmatrix} 1 & 1 & 0 & 3 & | & 4 \\ 0 & -1 & -1 & -5 & | & -7 \\ 3 & -1 & -1 & 2 & | & -3 \\ -1 & 2 & 3 & -1 & | & 4 \end{pmatrix}$$

2. Subtract $3R_1$ from R_3 :

$$R_3-3R_1\to R_3 \implies R_3=(3-3\cdot 1,-1-3\cdot 1,-1-3\cdot 0,2-3\cdot 3|-3-3\cdot 4)$$
 Resulting in:

$$\begin{pmatrix}
1 & 1 & 0 & 3 & | & 4 \\
0 & -1 & -1 & -5 & | & -7 \\
0 & -4 & -1 & -7 & | & -15 \\
-1 & 2 & 3 & -1 & | & 4
\end{pmatrix}$$

3. Add R_1 to R_4 :

$$R_4 + R_1 \rightarrow R_4 \implies R_4 = (-1+1, 2+1, 3+0, -1+3|4+4)$$

Resulting in:

$$\begin{pmatrix} 1 & 1 & 0 & 3 & | & 4 \\ 0 & -1 & -1 & -5 & | & -7 \\ 0 & -4 & -1 & -7 & | & -15 \\ 0 & 3 & 3 & 2 & | & 8 \end{pmatrix}$$

4. Next, we simplify rows to achieve upper triangular form. From R_3 and R_4 , we can continue applying row operations:

$$R_3 + 4R_2 \rightarrow R_3$$

Resulting in:

$$\begin{pmatrix} 1 & 1 & 0 & 3 & | & 4 \\ 0 & -1 & -1 & -5 & | & -7 \\ 0 & 0 & 3 & 13 & | & 13 \\ 0 & 0 & 0 & -13 & | & -13 \end{pmatrix}$$

Step 2: Back-Substitution

Now that we have an upper triangular matrix, we can use back-substitution to solve for the variables.

1. From the last row:

$$-13x_4 = -13 \implies x_4 = 1$$

2. Substitute x_4 into the third row:

$$3x_3 + 13(1) = 13 \implies 3x_3 = 0 \implies x_3 = 0$$

3. Substitute x_3 and x_4 into the second row:

$$-x_2 - 0 - 5(1) = -7 \implies -x_2 - 5 = -7 \implies x_2 = 2$$

4. Finally, substitute x_2 , x_3 , and x_4 into the first row:

$$x_1 + 2 + 0 + 3(1) = 4 \implies x_1 + 2 + 3 = 4 \implies x_1 = -1$$

The solution to the system is:

$$x_1 = -1$$

$$x_2 = 2$$

$$x_3 = 0$$

$$x_4 = 1$$

Example 2: Consider the linear system represented by the following equations:

$$4x_1 + 5x_2 + 6x_3 = 7$$

 $7x_1 + 8x_2 + 9x_3 = 8$
 $1x_1 + 2x_2 + 3x_3 = 3$

Use Gauss Elimination and back-substitution to solve the above linear system.

Factorization method

LU Factorization

In this section, we factorize the coefficient matrix A into an upper triangular matrix U and a lower triangular matrix L, such that:

$$A = LU$$

By the factorization A = LU, we can rewrite the system of equations $A\mathbf{x} = \mathbf{b}$ as:

$$LU\mathbf{x} = \mathbf{b}$$

We introduce a new vector y, where:

$$U\mathbf{x} = \mathbf{y}$$

Thus, the system becomes:

$$L\mathbf{y} = \mathbf{b}$$

To solve for **x**:

- 1. First, solve $L\mathbf{y} = \mathbf{b}$ using forward substitution to find \mathbf{y} .
- 2. Then, solve $U\mathbf{x} = \mathbf{y}$ using backward substitution to find \mathbf{x} .

Therefore, solving $A\mathbf{x} = \mathbf{b}$ is reduced to solving two simpler triangular systems, which are more computationally efficient.

$$\mathbf{x} = U^{-1}L^{-1}\mathbf{b}$$

Algorithm

Step 1: Decomposition

Decompose the matrix *A* into *L* and *U*.Generally, the procedure is as follows:

1st row entries in U,

$$U_{ij} = \frac{a_{ij}}{l_{11}}$$
 where $i = 1, 2, .., n$

1st column entries in L,

$$\ell_{i1} = rac{a_{i1}}{u_{11}}$$
 where $i = 1, 2, .., n$
 $\ell_{ii}u_{ii} = a_{ii} - \sum_{k=1}^{i-1} \ell_{ik}u_{ki}$ where $i = 1, 2, ...n - 1$
If $i > j$, then $\ell_{ji} = rac{a_{ji} - \sum_{k=1}^{i-1} \ell_{jk}u_{ki}}{u_{ii}}$
If $j > i$, then $u_{ij} = rac{a_{ij} - \sum_{k=1}^{i-1} \ell_{ik}u_{kj}}{\ell_{ij}}$
 $\ell_{nn}u_{nn} = a_{nn} - \sum_{k=1}^{n-1} \ell_{nk}u_{kn}$

Step 2: Forward Substitution

Solve $L\mathbf{v} = \mathbf{b}$ for \mathbf{v} using forward substitution. This involves solving the system from top to bottom:

$$y_{1} = \frac{b_{1}}{\ell_{11}}$$

$$y_{2} = \frac{b_{2} - \ell_{21}y_{1}}{\ell_{22}}$$

$$y_{i} = \frac{b_{i} - \sum_{j=1}^{i-1} \ell_{ij}y_{j}}{\ell_{ii}}$$

Step 3: Backward Substitution

Solve $U\mathbf{x} = \mathbf{y}$ for \mathbf{x} using backward substitution. This involves solving the system from bottom to top:

$$x_{n} = \frac{y_{n}}{u_{nn}}$$

$$x_{n-1} = \frac{y_{n-1} - u_{n-1,n}x_{n}}{u_{n-1,n-1}}$$

$$x_{i} = \frac{y_{i} - \sum_{j=i+1}^{n} u_{ij}x_{j}}{u_{ii}}$$

$$u_{ij} = \frac{y_{i} - \sum_{j=i+1}^{n} u_{ij}x_{j}}{u_{ij}}$$

$$u_{ij} = \frac{y_{i} - \sum_{j=i+1}^{n} u_{ij}x_{j}}{u_{ij}}$$

- Suppose that $\ell_{ii} = 1$ for i = 1, 2, 3, ..., n then this method is called **Doolittle's LU Factorization**.
- Suppose $u_{ii} = 1$ for i = 1, 2, ..., n. In this case, the method is called **Crout's LU Factorization**.

Example:

Given a system of linear equations:

$$A\mathbf{x} = \mathbf{b}$$

where:

$$A = \begin{pmatrix} 2 & 3 & 1 \\ 4 & 7 & 1 \\ -2 & 4 & -3 \end{pmatrix}, \quad \mathbf{b} = \begin{pmatrix} 1 \\ 3 \\ -1 \end{pmatrix}$$

Using Doolittle's LU factorization method, decompose matrix A into L and U:

$$L = \begin{pmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -1 & 7 & 1 \end{pmatrix}, \quad U = \begin{pmatrix} 2 & 3 & 1 \\ 0 & 1 & -1 \\ 0 & 0 & 5 \end{pmatrix}$$

Now, solve $L\mathbf{y} = \mathbf{b}$:

$$L\begin{pmatrix} y_1 \\ y_2 \\ y_3 \end{pmatrix} = \begin{pmatrix} 1 \\ 3 \\ -1 \end{pmatrix}$$

Using forward substitution, solve for **y**:

$$y_1 = 1, \quad y_2 = 1, \quad y_3 = -7$$

Next, solve $U\mathbf{x} = \mathbf{y}$:

$$U\begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 1 \\ 1 \\ -7 \end{pmatrix}$$

Using backward substitution, solve for x:

$$x_3 = -7/5$$
, $x_2 = -2/5$, $x_1 = 9/5$

Thus, the solution is:

$$\mathbf{x} = \begin{pmatrix} -7/5 \\ -2/5 \\ 9/5 \end{pmatrix}$$

Importance of Factorization

- Efficiency:factorization reduces the computational effort when solving multiple systems with the same coefficient matrix.
- Versatility: It can be used for both square and rectangular matrices, making it widely applicable in numerical solutions of linear systems.
- **Reusability**: Once the matrix *A* is factored into *L* and *U*, the factorization can be reused to solve multiple systems with different right-hand side vectors **b**.

Importance of Direct Methods

- Exact Solutions: Direct methods provide exact solutions in a finite number of steps.
- Wide Applicability: These methods can be applied to a broad range of linear systems, including those arising from scientific computing, engineering simulations, and optimization problems.
- Computational Efficiency: For small to moderate-sized systems, direct methods often offer better computational efficiency compared to iterative methods.
- Foundational Techniques: Direct methods form the basis for many numerical algorithms and are often used to preprocess data for more complex numerical analysis techniques.
- 5. **Numerical Stability:** Well-implemented direct methods can exhibit favorable numerical stability properties.

Iterative Methods

Iterative methods provide an approximate solution to the system. An iterative technique to solve the system Ax = bstarts with an initial approximation $x^{(0)}$ and generates a sequence of vectors $\{x^{(k)}\}_{k=0}^{\infty}$ that converges to x. Iterative techniques reformulate the system Ax = b into an equivalent system of the form:

$$x = Tx + c$$

for some fixed matrix T and vector c. After the initial vector $x^{(0)}$ is selected, the sequence of approximate solution vectors is generated by computing:

$$x^{(k)} = Tx^{(k-1)} + c$$
 for $k = 1, 2, 3, ...$

We will cover two iterative methods:

- Jacobi Method
- Gauss-Seidel Method



Convergence Criteria

For the iterative methods to converge:

► The matrix A should be diagonally dominant for the Jacobi and Gauss-Seidel methods.

Diagonal Dominance

A matrix is said to be diagonally dominant if, for every row of the matrix, the magnitude of the diagonal entry in a row is larger than or equal to the sum of the magnitudes of all the off-diagonal elements. More precisely, for all *i*:

$$|a_{ii}| > \sum_{j \neq i} |a_{ij}|$$

Example:

$$x_1 + 3x_2 + 5x_3 = 8$$

 $3x_1 + x_2 + x_3 = 5$
 $x_1 + 8x_2 - 5x_3 = 3$

If a system of equations is rearranged to achieve diagonal dominance, the matrix can then be solved iteratively. Now, a_{ii} is the coefficient with the largest magnitude in the i-th equation.

Jacobi Method

The Jacobi method solves each equation for a particular variable and uses these values to approximate the solution.

Algorithm:

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- 1. Rearrange the system $A\mathbf{x} = \mathbf{b}$ so that the diagonal elements of A are isolated.
- 2. Start with an initial guess $\mathbf{x}^{(0)}$.
- 3. Update each $x_i^{(k)}$ according to:

$$x_i^{(k)} = \frac{1}{a_{ii}} \left(b_i - \sum_{j \neq i} a_{ij} x_j^{(k-1)} \right)$$

for each i = 1, 2, ..., n.

4. Repeat until convergence.

Example 1:

$$E_1: 10x_1-x_2+2x_3=6$$

$$E_2: -x_1+11x_2-x_3+3x_4=25$$

$$E_3: 2x_1-x_2+10x_3-x_4=-11$$

$$E_4: 3x_2-x_3+8x_4=15$$

This system is already in diagonal dominance form. To convert Ax = b to the form x = Tx + c, solve each equation E_i for x_i for i = 1, 2, 3, 4.

The first 10 iterations of the Jacobi method are given in the table below:

k	$X_1^{(k)}$	$x_{2}^{(k)}$	$X_3^{(k)}$	$X_4^{(k)}$
0	0.0000	0.0000	0.0000	0.0000
1	0.6000	2.2727	-1.1000	1.8750
2	1.0473	1.7159	-0.8052	0.8852
3	0.9326	2.0533	-1.0494	1.1309
4	1.0152	1.9537	-0.9681	0.9793
5	0.9890	2.0114	-1.0103	1.0214
6	1.0032	1.9922	-0.9945	0.9944
7	0.9981	2.0023	-1.0020	1.0036
8	1.0006	1.9987	-0.9940	0.9989
9	0.9997	2.0004	-1.0044	1.0006
10	1.0001	1.9998	-0.9998	0.9998

By considering the last two iterations ($x^{(9)}$ and $x^{(10)}$), we can see that convergence has occurred. Therefore, the required solution is:

Matrix Form of Jacobi Method

The Jacobi method can also be written in matrix form. We split the matrix A into its diagonal, lower, and upper triangular parts:

$$A = D - L - U$$

where D is the diagonal matrix, L is the strictly lower triangular part, and U is the strictly upper triangular part. If D^{-1} exist that $a_{ii} \neq 0$ for each i, this leads to the iterative form:

$$x^{(k)} = D^{-1}(L+U)x^{(k-1)} + D^{-1}b$$

In practice, this equation is used for computation, while the theoretical form is:

$$x^{(k)} = Tx^{(k-1)} + c$$

where $T = D^{-1}(L + U)$ and $c = D^{-1}b$.

Note

The **infinity norm** (also called the *maximum norm* or *supremum norm*) is a vector norm defined as:

$$\|x\|_{\infty} = \max_{1 \le i \le n} |x_i|$$

For a vector $x = (x_1, x_2, ..., x_n) \in \mathbb{R}^n$, the infinity norm is given by:

$$||x||_{\infty} = \max(|x_1|, |x_2|, \dots, |x_n|)$$

Example: For a vector x = (3, -7, 5), the infinity norm is:

$$||x||_{\infty} = \max(|3|, |-7|, |5|) = 7.$$

The Error of the n^{th} computation is defined as

$$e_n = \frac{\|x^{(n)} - x^{(n-1)}\|_{\infty}}{\|x^{(n)}\|_{\infty}}$$

in infinity norm(I_{∞}).

Example 2: Consider the following system of equations:

$$E_1: 10x_1-x_2+2x_3=6,$$

$$E_2: -x_1+11x_2-x_3+3x_4=25,$$

$$E_3: 2x_1-x_2+10x_3-x_4=-11,$$

$$E_4: 3x_2-x_3+8x_4=15.$$

Use the Jacobi method to solve this system, starting with an initial guess:

$$x^{(0)} = (0, 0, 0, 0).$$

Continue iterating until:

$$||x^{(k)} - x^{(k-1)}||_{\infty} < 10^{-3}.$$

Example 3: Use Jacobi method to solve the following linear system with Tolerance=TOL= 10^{-3} in the I_{∞} norm, that is

$$e_n = \frac{\|x^{(n)} - x^{(n-1)}\|_{\infty}}{\|x^{(n)}\|_{\infty}} < 10^{-3}$$

$$E_1: 10x_1 + 2x_2 + x_3 = 9,$$

$$E_2: \quad x_1+10x_2-x_3=-22,$$

$$E_3: -2x_1+3x_2+10x_3=22.$$

We initial guess:

$$x^{(0)} = (0, 0, 0).$$

The iterative equations can be derived as follows:

$$x_1 = \frac{9 - 2x_2 - x_3}{10},$$

$$x_2 = \frac{-22 - x_1 + x_3}{10},$$

$$x_3 = \frac{22 + 2x_1 - 3x_2}{10}.$$

Continue iterating until:

$$\frac{\|x^{(n)} - x^{(n-1)}\|_{\infty}}{\|x^{(n)}\|_{\infty}} < 10^{-3}$$

Now, we proceed with the iterations:

Iteration	$X_1^{(k)}$	$X_{2}^{(k)}$	$x_3^{(k)}$	$ x^{(k)}-x^{(k-1)} _{\infty}$	$Error(e_k)$
0	0.0000	0.0000	0.0000	_	_
1	0.9000	-2.2000	2.2000	2.2000	1.0000
2	1.1200	-2.0700	3.0400	0.8400	0.2763
3	1.0100	-2.0080	3.04500	0.1100	0.0361
4	0.9971	-1.9965	3.0044	0.0406	0.0135
5	0.9989	-1.9993	2.9984	0.0060	0.0020
6	1.0000	-2.0001	2.9996	0.0012	0.0004

After the 6th iteration, we can see that:

Convergence occurs:
$$\frac{\|x^{(6)} - x^{(5)}\|_{\infty}}{\|x^{(6)}\|_{\infty}} < 10^{-3}$$

Thus, the approximate solution is:

$$x_1 \approx 1$$
, $x_2 \approx -2$, $x_3 \approx 3$.



The Jacobi iterative method is a simple and effective technique for approximating the solution to a system of linear equations, especially when the matrix *A* is diagonally dominant. The method converges when the successive approximations are close enough, achieving the required accuracy.

Gauss-Seidel Method

The Gauss-Seidel method improves upon the Jacobi method by using the most recently computed values in the iterative process. In Jacobi's method, all components of $x^{(k)}$ are calculated using the values from $x^{(k-1)}$. In contrast, the Gauss-Seidel method updates each component as soon as its value is computed, leading to potentially faster convergence.

Algorithm:

- 1. Start with an initial guess $\mathbf{x}^{(0)}$.
- 2. Update each $x_i^{(k)}$ using:

$$x_i^{(k)} = \frac{1}{a_{ii}} \left(b_i - \sum_{j < i} a_{ij} x_j^{(k)} - \sum_{j > i} a_{ij} x_j^{(k-1)} \right)$$

3. Repeat until convergence.

Example 3:

Solve the system:

$$10x_1 - x_2 + 2x_3 = 6,$$

$$-x_1 + 11x_2 - x_3 + 3x_4 = 25,$$

$$2x_1 - x_2 + 10x_3 - x_4 = -11,$$

$$3x_2 - x_3 + 8x_4 = 15.$$

Start with:

$$x^{(0)} = (0, 0, 0, 0).$$

Continue iterating until:

$$\frac{\|x^{(k)}-x^{(k-1)}\|_{\infty}}{\|x^{(k)}\|_{\infty}}<10^{-3}.$$

k	$x_1^{(k)}$	$X_2^{(k)}$	$X_3^{(k)}$	$X_4^{(k)}$
0	0.0000	0.0000	0.0000	0.0000
1	0.6000	2.3227	-0.9873	0.8789
2	1.0302	2.0369	-1.0144	0.9844
3	1.0066	2.0035	-1.0025	0.9984
4	1.0009	2.0003	-1.0003	0.9999
5	1.0001	2.0000	-1.0000	1.0000

After several iterations, check:

$$\frac{\|x^{(k)}-x^{(k-1)}\|_{\infty}}{\|x^{(k)}\|_{\infty}}<10^{-3}.$$

Once this condition is met, $x^{(5)}$ is accepted as a reasonable approximation of the solution.

Matrix Form

The Gauss-Seidel method can also be represented in matrix form as:

$$(D-L)x^{(k)}=Ux^{(k-1)}+b$$

If $(D-L)^{-1}$ exists, then that

$$x^{(k)} = (D-L)^{-1}Ux^{(k-1)} + (D-L)^{-1}b$$

where D, L, and U are the diagonal, lower triangular, and upper triangular parts of the coefficient matrix, respectively. For the

lower triangular matrix (D-L) to be non-singular, it is necessary and sufficient that $a_{ii} \neq 0$, for each i = 1, 2, 3, ..., n.

Comparison Between Iterative Methods

- Convergence Speed: The Gauss-Seidel method converges more quickly than the Jacobi method due to its use of updated values immediately in calculations.
- Memory Efficiency: Iterative methods are typically more memory-efficient than direct methods, making them suitable for large-scale problems.
- ► Error Correction: Iterative methods benefit from inherent self-correcting properties, allowing for adjustments to errors made in previous iterations.

Comparison Between Direct and Iterative Methods

- Efficiency for Large Systems: For larger systems of equations, iterative methods are often faster than direct methods.
- ▶ Direct Methods for Specific Structures: Direct methods are particularly effective for solving tri-diagonal systems.
- Round-off Errors: Direct methods are prone to round-off errors.
- ➤ **Self-Correcting Nature:** Iterative methods correct errors automatically in subsequent iterations.
- Sparse Coefficient Matrices: When the coefficient matrix contains a significant number of zeros, iterative methods can leverage this sparsity for faster computations compared to direct methods.
- Scalability: Iterative methods can be more scalable for very large problems