- 1. What is Disjoint get? What are its applications
 - A Disjoint Set (Union-Find) is a data Structure that keeps track of elements partitioned into disjoint (non-overdapping) sets.
 - · Applicautions:
 - · kruskal's algorithm for minimum Spanning Tree
 - · Network Connectivity
 - · Image Processing Connected Componente)
 - · Cycle detection in grouphs.
- 9. What is a Red-Black Trace? What area to its
 Propercises?
- · A Red Black Tree is a balanced binary Gearch tree where each node has a colore treed on black) ensuring balance through Prioperatios:
 - · Each node is tred at block. ! [[[]]]
 - . Root is always black. (iii) : 2000 1239.
 - . No two red nodes are adjacent.
- · Every path from a node to 9ts descendant NULL nodes here the same numbers of black nodes
- 8. What 1s divide and conquer paradigm for problem Golving.
 - · Divide and conquer involves:
 - . Dividue Breeak the Problem into Subproblems
- recursively. Solve the gularcobkern into
 - final ressult.

4-What is Art Trace? An AVL True is a Self. balancing binarry Search tree where the difference between heights of left and reight subtraces cannot be more than I for any node how tong the Applications: Section-B · Kruskal's · Olgorathin for mintagen span 1. Algorithm: selection Sorch (int arcre[], int n) ? void Common force Cent 1 = Di i Kn-1; itt) F int onin-ida = 1; for (inf j = 1+1; j < n; j++) { if (ann [i] & ann [min-ida]) min_idn= j; Swap (arer [min-idn], arer [i]); stated there should bear they properly Prisade malact prisade Capall no bord Time Complenity: (EDITANIGOUS · Best cose: O(m) . Assid appeals it had . · Average case: Q(n2) times words coise: Q(n) Anti- single trees the come simplere of process andice 2. Write an Algorithm: Void BFS (int Start) \$ University cooldons Queue 9; visited [stard] = true;

9. enqueire (stard);

while (!9. is finpty ())? int node = q. dequeue (); for (each neighbour of node)? if (! visited [neighbouri])?

4/3/ted [neighbour] = treve;

```
9. enqueue (neighborn);

P

Brs traversal for given graph Starcting of A:

Order: A >B > C > D > F > F
```

Section-c

1. Heap: A complete binary frue where every Parcent node is greatien (man-Heap) on Smaller (min-Heap) than its children.

· max- HEAPIFY:

Not described for area [], int n, int i) ?

int largest = i;
int left = 2*i + 1;

int left - 2*i + 2;

if (left < n & area [left] > area [larges])

largest = left;

if (reight < n & area [reight] > area [largest])

largest - reight;

if (largest - reight;

swap (area [i], area [largest]);

mantleapity (area, n, baregest);

?

Heap Sort Steps on 8:

Afton building man heap: [84,22,19,10,5,17,6,3]

Sorted antroy after Heap Sort:

[3,5,6,10,17,19,22,84]

· Time complemity:
· Q(n log n)

```
2.
   Algorathm:
       void quick Soret (int arcre[], int low, int high) {
           : & (low < high ) {
               Int Pi = Parctition (arcre, low, high);
               ·quick Sout (aren, los , Pi-1);
                quickSorer (arcr, Pi+1, high);
   Partition
        int Pourtition (int arcrity, int low, int high) of
          ; nt Pluot = arcrathigh ];
                             POHOCKE , nodes to green
          int i = low - 1;
          for (int i = low; ix= high -1; itt)
            if Carcre [j] < Pivo+ ) ?
        swap (arere [i] arere [i]);
         Swap (arm [i+1], arm [high]);
return 1+1;
                     1 1701 - 123 Bugl
   Time Complenity:
       · Best case: O (n log n)
       · Average casio: 0 (n log n)
        · Worst cose: (Cn2)
  · Quick Sort Steps for A:
     Sorched ourceay : [3,5,6,10,17,19,22,849
10, F1, 2,01, PT ( EC. P.S. ] . Sep. 1 20001 . Dellard . Dellard . Dellard .
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Counted control offers head south 1 12,55 pt; Ft, of, d 2, 81 }

: Philosophias Dail