

APPENDIX A
PROOF OF LEMMA 1

Lemma 1. *Given any feasible task replication and power allocation, the Lyapunov drift-plus-penalty function B in our formulation (32) is upper bounded by the following expression.*

$$B \geq \frac{1}{2}E \left[\left(\sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) \right)^2 + (\mu_b t_a)^2 \right] + \frac{1}{2}E \left[\left(\sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) \right)^2 + (1_{Q(t)>0} \mu_c t_a)^2 \right],$$

where $R_{i,b}(t)$ denotes the total amount of computing from TaV i in the time slot t , $1_{Q(t)>0}$ represents a characteristic function, which only takes the value of 1 when $Q(t) > 0$.

Proof. By squaring both sides of the equation, we can get the expression below.

$$\begin{aligned} Q(t+1)^2 &= \left(\max \left\{ Q(t) + \sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) - \mu_b t_a, 0 \right\} \right)^2 \\ &\leq Q(t)^2 + (\mu_b t_a)^2 + \left(\sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) \right)^2 \\ &\quad - 2\mu_b t_a Q(t) - 2\mu_b t_a \sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) \\ &\quad + 2 \sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) \mu_b t_a. \end{aligned}$$

The following inequality can be obtained as follows.

$$\begin{aligned} \frac{1}{2}[Q(t+1)^2 - Q(t)^2] &\leq \frac{1}{2} \left[\left(\sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) \right)^2 \right. \\ &\quad \left. + (\mu_b t_a)^2 + 2Q(t) \left(\sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) - \mu_b t_a \right) \right]. \end{aligned}$$

Similarly, we can obtain the following expression.

$$\begin{aligned} \frac{1}{2}[H(t+1)^2 - H(t)^2] &\leq \frac{1}{2} \left[\left(\sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) \right)^2 \right. \\ &\quad \left. + (1_{Q(t)>0} \mu_c t_a)^2 + 2 \sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) H(t) \right. \\ &\quad \left. - 2 \sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) 1_{Q(t)>0} \mu_c t_a - 2H(t) 1_{Q(t)>0} \mu_c t_a \right] \\ &\leq \frac{1}{2} \left[(1_{Q(t)>0} \mu_c t_a)^2 + \left(\sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) \right)^2 \right. \\ &\quad \left. + 2H(t) \left(\sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) - 1_{Q(t)>0} \mu_c t_a \right) \right]. \end{aligned}$$

According to the definition of Lyapunov drift function, we can get the expression below.

$$L(\Theta(t)) = \frac{1}{2} (Q(t)^2 + H(t)^2),$$

$$\Delta L(t) = E[L(\Theta(t+1)) - L(\Theta(t))].$$

Then, we can get a bound on $\Delta L(t)$.

$$\begin{aligned} \Delta L(t) &\leq B + \left[Q(t) \left(\sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) - \mu_b t_a \right) \right. \\ &\quad \left. + H(t) \left(\sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) - 1_{Q(t)>0} \mu_c t_a \right) \right], \end{aligned}$$

where B is the upper bound by the expression as below.

$$\begin{aligned} B &\geq \frac{1}{2}E \left[\left(\sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) \right)^2 + (\mu_b t_a)^2 \right] \\ &\quad + \frac{1}{2}E \left[\left(\sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) \right)^2 + (1_{Q(t)>0} \mu_c t_a)^2 \right], \end{aligned}$$

□

APPENDIX B
PROOF OF LEMMA 2

Lemma 2. *An algorithm that can effectively control $Q(t)$ and $H(t)$ within prescribed limits ($Q(t) \leq Q_{\max}$, $H(t) \leq H_{\max}$) will ensure a capped worst-case delay. The worst-case delay for the task transmission to the server queue is bounded by the constant T_{\max} defined as follows:*

$$T_{\max} = \lceil \frac{Q_{\max} + H_{\max}}{\mu_b t_a} \rceil,$$

where $\lceil \cdot \rceil$ denotes the smallest integer that is greater than or equal to the value.

Proof. The task arrives at slot t can depart the queue at the slot $t+1$ earliest. We prove that these tasks can depart the queue by time $t+T_{\max}$ by contradiction. It must be that $Q(t_x) > 0$ for all $t_x \in \{t+1, \dots, t+T_{\max}\}$. We have:

$$H(t_x+1) = \max[H(t_x) + \sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) - \mu_b t_a, 0],$$

and hence for all $t_x \in \{t+1, \dots, t+T_{\max}\}$:

$$H(t_x+1) \geq H(t_x) + \sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) - \mu_b t_a.$$

Summing the above over $t_x \in \{t+1, \dots, t+T_{\max}\}$ yields:

$$\begin{aligned} H(t+T_{\max}+1) - H(t+1) &\geq \mu_b t_a T_{\max} - \sum_{t_x=t+1}^{t+T_{\max}} \sum_{i=1}^I x_{i,b}(t_x) R_{i,b}(t_x). \end{aligned}$$

Using the fact that $H(t+T_{\max}+1) \leq T_{\max}$ and $H(t+1) \leq 0$, we have:

$$\mu_b t_a T_{\max} - H_{\max} \leq \sum_{t_x=t+1}^{t+W_{\max}} \sum_{i=1}^I x_{i,b}(t_x) R_{i,b}(t_x).$$

Because we have assumed that not all of the tasks arrive at t slot depart by time $t + W_{\max}$, we must have:

$$\sum_{t_x=t+1}^{t+W_{\max}} \sum_{i=1}^I x_{i,b}(t_x) R_{i,b}(t_x) < Q(t) \leq Q_{\max}.$$

Combining above equations, we can get:

$$\mu_b t_a T_{\max} - H_{\max} \leq Q_{\max}.$$

Therefore:

$$T_{\max} \leq \frac{Q_{\max} + H_{\max}}{\mu_b t_a}.$$

This contradicts the definition we have made before. \square

APPENDIX C PROOF OF LEMMA 3

Lemma 3. The BER $\Omega_{i,b}(t)$ between TaV i and the edge cloud and the BER $\Omega_{i,j}(t)$ between TaV i and SeV j under Rayleigh fading conditions have the approximation as below.

$$\Omega_{i,b}(t) \approx \frac{(e^{R_{i,b}(t)} - 1) \left(\sum_{k=i+1}^I |h_{k,b}(t)|^2 \alpha_i^n(t) p_{k,b}(t) + \sigma^2 \right)}{p_{i,b}(t) |h_{i,b}(t)|^2},$$

$$\Omega_{i,j}(t) \approx \frac{(e^{\frac{K_{i,j}}{L_{i,j}}} - 1) \sigma^2}{p_{i,j}(t) |h_{i,j}(t)|^2}.$$

Proof. According to Taylor's Expanded Form, we can get the approximations of BER as below.

$$\Omega_{i,b}(t) \approx 1 - k_1 \gamma_{i,b}(t) \left(1 + \frac{k_2}{\gamma_{i,b}(t)} + \frac{k_2^2}{2\gamma_{i,b}^2(t)} - 1 - \frac{k_3}{\gamma_{i,b}(t)} - \frac{k_3^2}{2\gamma_{i,b}^2(t)} \right),$$

$$\text{where } k_1 = \sqrt{\frac{L_{i,b}}{2\pi(e^{\frac{K_{i,b}}{L_{i,b}}} - 1)}}, k_2 = 1 - e^{\frac{K_{i,b}}{L_{i,b}}} + \sqrt{\frac{\pi(e^{\frac{K_{i,b}}{L_{i,b}}} - 1)}{2L_{i,b}}},$$

$$\text{and } k_3 = 1 - e^{\frac{K_{i,b}}{L_{i,b}}} - \sqrt{\frac{\pi(e^{\frac{K_{i,b}}{L_{i,b}}} - 1)}{2L_{i,b}}}.$$

$$\Omega_{i,j}(t) \approx 1 - g_1 \gamma_{i,j}(t) \left(1 + \frac{g_2}{\gamma_{i,j}(t)} + \frac{g_2^2}{2\gamma_{i,j}^2(t)} - 1 - \frac{g_3}{\gamma_{i,j}(t)} - \frac{g_3^2}{2\gamma_{i,j}^2(t)} \right),$$

$$\text{where } g_1 = \sqrt{\frac{L_{i,j}}{2\pi(e^{\frac{K_{i,j}}{L_{i,j}}} - 1)}}, g_2 = 1 - e^{\frac{K_{i,j}}{L_{i,j}}} + \sqrt{\frac{\pi(e^{\frac{K_{i,j}}{L_{i,j}}} - 1)}{2L_{i,j}}},$$

$$\text{and } g_3 = 1 - e^{\frac{K_{i,j}}{L_{i,j}}} - \sqrt{\frac{\pi(e^{\frac{K_{i,j}}{L_{i,j}}} - 1)}{2L_{i,j}}}.$$

After simplification, we can obtain the below expressions.

$$\Omega_{i,b}(t) \approx \frac{(e^{\frac{K_{i,b}}{L_{i,b}}} - 1) \left(\sum_{k=i+1}^I |h_{k,b}(t)|^2 p_{k,b}(t) + \sigma^2 \right)}{p_{i,b}(t) |h_{i,b}(t)|^2},$$

$$\Omega_{i,j}(t) \approx \frac{(e^{\frac{K_{i,j}}{L_{i,j}}} - 1) \sigma^2}{p_{i,j}(t) |h_{i,j}(t)|^2}.$$

\square

APPENDIX D PROOF OF LEMMA 4

Lemma 4. The expected service reliability incurred by PATRLF could achieve the effect that has a gap with the maximal service reliability, which can be described below.

$$\lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=0}^{T-1} E\{\Upsilon(t_x) | \Theta(t_x)\} \leq \Upsilon^{\sup}(t) - \frac{B}{V},$$

where $\Upsilon^{\sup}(t)$ is the supremum of the average service reliability achieved by any joint strategy under the required constraints, and B is mentioned in **Lemma 1**.

Proof. From **Lemma 2**, we can get the below expression.

$$\begin{aligned} \Delta L(t) - VE\{\Upsilon(t) | \Theta(t)\} &\leq B + Q(t) E\left\{ \sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) \right. \\ &\quad \left. - \mu_b t_a | \Theta(t) \right\} + H(t) E\left\{ \sum_{i=1}^I x_{i,b}(t) R_{i,b}(t) \right. \\ &\quad \left. - 1_{Q(t) > 0} \mu_c t_a | \Theta(t) \right\} - VE\{\Upsilon(t) | \Theta(t)\}. \end{aligned}$$

B is the upper bound, which is mentioned in **Lemma 1**. From **Section 3.1.4** in the study [16], we have the following inequality.

$$B - V\Upsilon^{\sup}(t) \leq \Delta L(t) - VE\{\Upsilon(t) | \Theta(t)\}.$$

By adding $t = 0, 1, \dots, T-1$ cumulatively, we get:

$$\begin{aligned} TB - TV\Upsilon^{\sup}(t) &\leq E\{L(\Theta(t)) - L(\Theta(0))\} \\ &\quad - V \sum_{t=0}^{T-1} E\{\Upsilon(t) | \Theta(t)\}. \end{aligned}$$

Since $L(\Theta(t)) \geq 0, L(\Theta(0)) = 0$, both sides divide by the same $-TV$, when $t \rightarrow +\infty$, we can obtain that:

$$\limsup_{t \rightarrow +\infty} \frac{1}{t} \sum_{t=0}^{T-1} E\{\Upsilon(t)\} \leq \Upsilon^{\sup}(t) - \frac{B}{V}.$$

With Jensen's inequality, we can get:

$$\lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=0}^{T-1} E\{\Upsilon(t)\} \leq \Upsilon^{\sup}(t) - \frac{B}{V}.$$

\square

APPENDIX E
PROOF OF LEMMA 5

Lemma 5. *The virtual task queue in the edge cloud and the virtual delay queue for all TaVs in NOMA-enabled VEC are stable, which can be presented as follows.*

$$\lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t_x=0}^{T-1} (E\{Q(t_x)\} + E\{H(t_x)\}) \leq \frac{V(\hat{\Upsilon}(t) - \check{\Upsilon}(t)) - B}{\varsigma},$$

where $\hat{\Upsilon}(t)$ and $\check{\Upsilon}(t)$ are defined as the maximum and minimum service reliability among all strategies. B is mentioned in **Lemma 1**.

Proof. We now prove that $Q(t)$ and $G_i(t)$ are mean rate stable under Algorithm 3.

Denote $\Phi(t) = TV\{\hat{\Upsilon}(t)\} - V \sum_{t=0}^{T-1} E\{\Upsilon(t)|\Theta(t)\} - TB + E\{L(\Theta(0))\}$.

Based on **Lemma 3**, we can get $E\{L(\Theta(t))\} \leq \Phi(t)$, i.e., $E\{Q(t)^2\} \leq 2\Phi(t)$ and $E\{H(t)^2\} \leq 2\Phi(t)$. Thus, we have the following inequality.

$$\begin{aligned} 0 &\leq E\{Q(t)\} \leq \sqrt{2\Phi(t)}. \\ 0 &\leq E\{H(t)\} \leq \sqrt{2\Phi(t)}. \end{aligned}$$

Since $\lim_{T \rightarrow \infty} \frac{\sqrt{2\Phi(t)}}{T} = 0$, we can also acknowledge that

$$\begin{aligned} \lim_{T \rightarrow \infty} \frac{E\{Q(t)\}}{T} &= 0. \\ \lim_{T \rightarrow \infty} \frac{E\{H(t)\}}{T} &= 0. \end{aligned}$$

To find the upper bound of the average sum queue length of $E\{Q(t)\}$ and $E\{H(t)\}$, we can obtain the below expression in the proof of **Lemma 3**,

$$\begin{aligned} L(\Theta(t+1)) - L(\Theta(t)) - VE\{\Upsilon(t)|\Theta(t)\} &\geq \\ B - V\Upsilon^{\sup}(t) + \left[Q(t) \left(\sum_{i=1}^I x_{i,b}(t)R_{i,b}(t) - \mu_b t_a \right) \right. \\ &\quad \left. + H(t) \left(\sum_{i=1}^I x_{i,b}(t)R_{i,b}(t) - 1_{Q(t)>0} \mu_c t_a \right) \right], \end{aligned}$$

where $\Upsilon^{\sup}(t)$ is the supremum of the average service reliability achieved by any joint strategy under the required constraints.

Then, summing over $t \in \{0, 1, 2, \dots, T-1\}$, we can obtain the expression below.

$$\begin{aligned} L(\Theta(T)) - L(\Theta(0)) - \sum_{t=0}^{T-1} VE\{\Upsilon(t)|\Theta(t)\} &\geq \\ TB - TV\Upsilon^{\sup}(t) + \varsigma(E\{H(t)\} + E\{Q(t)\}), \end{aligned}$$

where $\varsigma > 0$, and after dividing T and ς by both sides of the above inequality, we can finally get.

$$\begin{aligned} \frac{1}{T} (E\{H(t)\} + E\{Q(t)\}) &\leq \frac{E(L(T))}{\varsigma T} - \frac{E(L(0))}{\varsigma T} \\ &\quad - \frac{B - V\Upsilon^{\sup}(t) + \frac{1}{T} \sum_{t=0}^{T-1} VE\{\Upsilon(t)|\Theta(t)\}}{\varsigma}, \end{aligned}$$

$$\Upsilon^{\sup}(t) - \frac{1}{T} \sum_{t=0}^{T-1} E\{\Upsilon(t)|\Theta(t)\} \leq \hat{\Upsilon}(t) - \check{\Upsilon}(t).$$

From the predefined value $L(0)$, we can get the conclusion.

$$\begin{aligned} \lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t_x=0}^{T-1} (E\{H(t_x)\} + E\{Q(t_x)\}) \\ \leq \frac{V(\hat{\Upsilon}(t) - \check{\Upsilon}(t)) - B}{\varsigma}. \end{aligned}$$

□