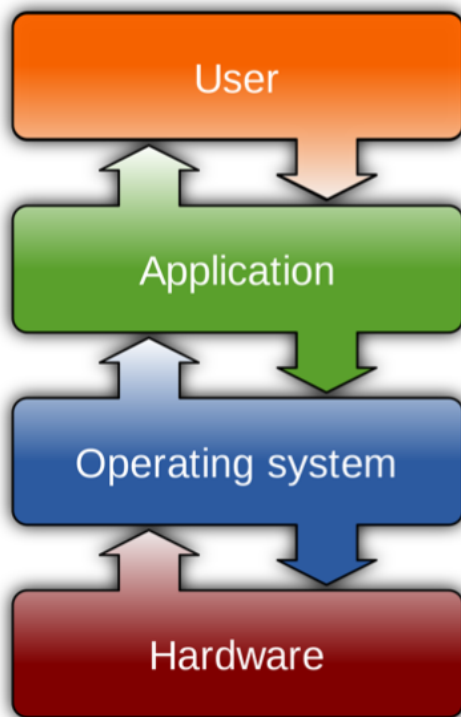


CS Revision Lecture 17, 18

- **Lecture 17 - Operating Systems : Key concepts & security principles**

- **Operating Systems**

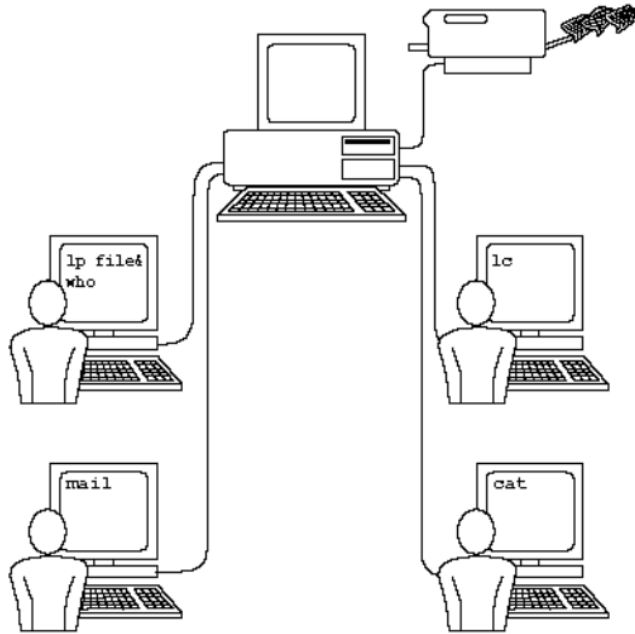
- An OS provides the interface between the users of a computer and that computer's hardware
- The OS handles the management of low-level hardware resources:



- disk drives
- CPU
- RAM
- I/O devices
- Network interfaces

- **Multi-users**

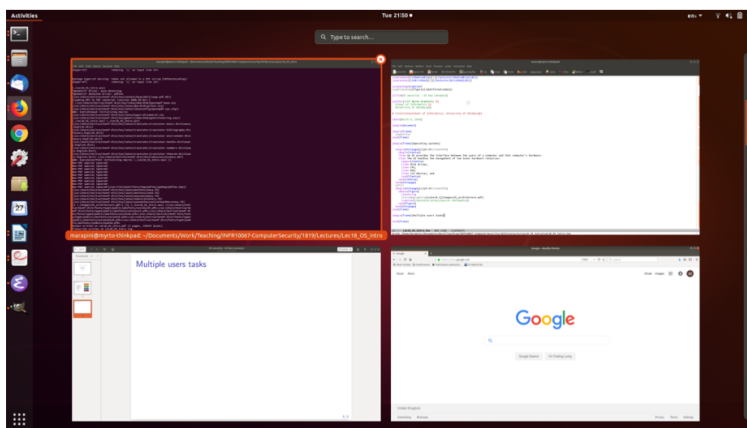
- OS must allow for multiple users with potentially different levels of access to the same computer



- The OS needs to have in place mechanisms to isolate different users

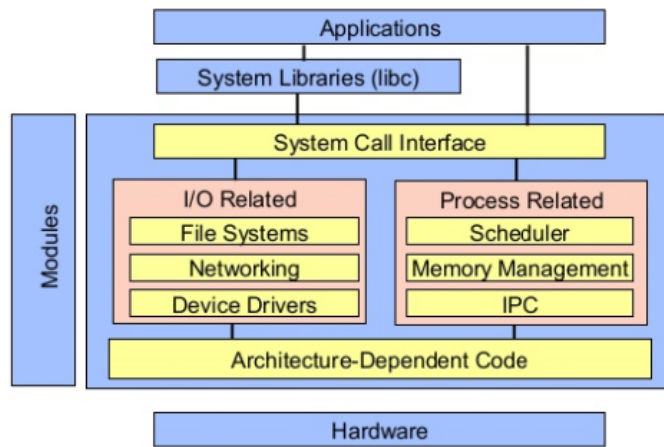
Multi-tasking

- OS must allow multiple application programs to run at the same time



- The OS needs to have in place mechanisms to **isolate different applications running**

Essential Unix architecture



- **Kernel**
 - key component of the OS
 - supports secure sharing of low-level resources between users/applications
 - kernel limits how applications access computer resources
- **Execution modes**
 - User mode - access to resources through syscall to kernel
 - Kernel mode - direct access to resources
- **System calls** are usually contained in a collection of programs, e.g. a library such as the C library libc:
 - `open()` `close()` `read()` `write()`
 - `wait()` `fork()` `exec()` `exit()`
- **Processes and process management**

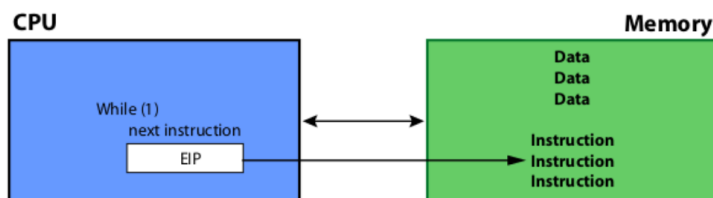
```

marapin@myrto-thinkpad: ~/Documents/Work/Teaching/INFR10067-ComputerSecurity/1819/Lectures/Lec18_05_Intro$ ps -ef
File Edit View Search Terminal Help
marapin@myrto-thinkpad: ~/Documents/Work/Teaching/INFR10067-ComputerSecurity/1819/Lectures/Lec18_05_Intro$ ps -ef
UID        PID     PPID  C  STIME TTY          TIME CMD
root         1         0  0  Mar03 ?        00:00:29 /sbin/init splash
root         2         0  0  Mar03 ?        00:00:00 [kthreadd]
root         4         2  0  Mar03 ?        00:00:00 [kworker/0:0]
root         6         2  0  Mar03 ?        00:00:00 [mm_percpu_wq]
root         7         2  0  Mar03 ?        00:00:01 [ksoftirqd/0]
root         8         2  0  Mar03 ?        00:00:41 [rcu_sched]
root         9         2  0  Mar03 ?        00:00:00 [rcu_bh]
root        10         2  0  Mar03 ?        00:00:00 [nigrtion/0]
root        11         2  0  Mar03 ?        00:00:00 [watchdog/0]
root        12         2  0  Mar03 ?        00:00:00 [cpuhp/0]
root        13         2  0  Mar03 ?        00:00:00 [cpuhp/1]
root        14         2  0  Mar03 ?        00:00:00 [watchdog/1]
root        15         2  0  Mar03 ?        00:00:00 [nigrtion/1]
root        16         2  0  Mar03 ?        00:00:01 [ksoftirqd/1]
root        18         2  0  Mar03 ?        00:00:00 [kworker/1:0]
root        19         2  0  Mar03 ?        00:00:00 [cpuhp/2]
root        20         2  0  Mar03 ?        00:00:00 [watchdog/2]
root        21         2  0  Mar03 ?        00:00:00 [nigrtion/2]
root        22         2  0  Mar03 ?        00:00:01 [ksoftirqd/2]
root        24         2  0  Mar03 ?        00:00:00 [kworker/2:0]
root        25         2  0  Mar03 ?        00:00:00 [cpuhp/3]
root        26         2  0  Mar03 ?        00:00:00 [watchdog/3]
root        27         2  0  Mar03 ?        00:00:00 [nigrtion/3]
root        28         2  0  Mar03 ?        00:00:01 [ksoftirqd/3]
root        30         2  0  Mar03 ?        00:00:00 [kworker/3:0]
root        31         2  0  Mar03 ?        00:00:00 [kdevtmpfs]
root        32         2  0  Mar03 ?        00:00:00 [netns]
root        33         2  0  Mar03 ?        00:00:00 [rcu_tasks_kthre]
root        34         2  0  Mar03 ?        00:00:00 [kauditd]
root        38         2  0  Mar03 ?        00:00:00 [khungtaskd]
root        39         2  0  Mar03 ?        00:00:00 [oom_reaper]
root        40         2  0  Mar03 ?        00:00:00 [writeback]
root        41         2  0  Mar03 ?        00:00:00 [kcompactd0]
root        42         2  0  Mar03 ?        00:00:00 [ksmd]
root        43         2  0  Mar03 ?        00:00:00 [khugepaged]
root        44         2  0  Mar03 ?        00:00:00 [crypto]
root        45         2  0  Mar03 ?        00:00:00 [kintegrityd]
root        46         2  0  Mar03 ?        00:00:00 [kblockd]
root        48         2  0  Mar03 ?        00:00:00 [ata_sff]
root        49         2  0  Mar03 ?        00:00:00 [nd]
root        50         2  0  Mar03 ?        00:00:00 [edac-poller]
root        51         2  0  Mar03 ?        00:00:00 [devfreq_wq]
root        52         2  0  Mar03 ?        00:00:00 [watchdogd]
root        55         2  0  Mar03 ?        00:00:00 [kswapd0]
root        56         2  0  Mar03 ?        00:00:00 [ecryptfs-kthrea]
root        98         2  0  Mar03 ?        00:00:00 [kthrotld]
root        99         2  0  Mar03 ?        00:00:00 [acpi_thermal_pm]
root       103         2  0  Mar03 ?        00:00:00 [ipvs_addrconf]
root       112         2  0  Mar03 ?        00:00:00 [kstrp]
root       129         2  0  Mar03 ?        00:00:00 [charger_manager]
root       176         2  0  Mar03 ?        00:00:00 [nme-wq]
root       180         2  0  Mar03 ?        00:00:10 [/915/signal:0]
root       181         2  0  Mar03 ?        00:00:00 [/915/signal:1]

```

- A process is an instance of a program that is currently executing
- To actually be executed the program must be loaded into RAM and uniquely identified
- Each process running is identified by a **unique process ID (pid)**
- To a pid, we can associate its CPU time, memory usage, userID (uid), program name, etc
- A process might control other processes (fork)
- Child process inherits context from parent process

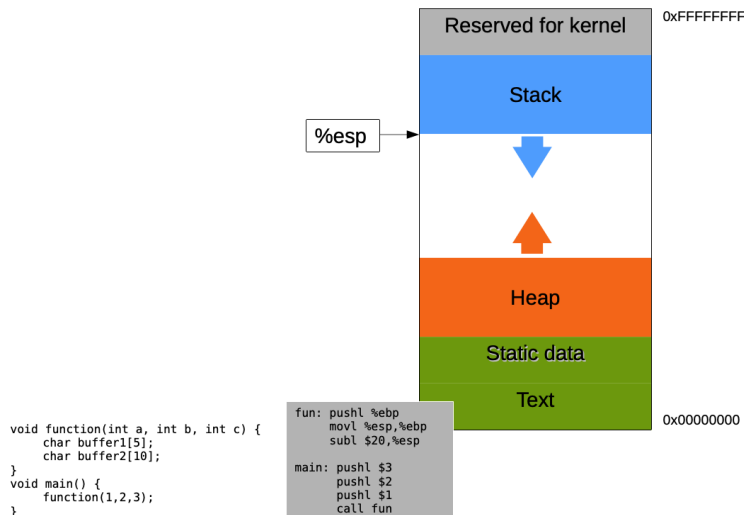
x86 CPU/Memory



- To actually be executed the program must be loaded into RAM and uniquely identified

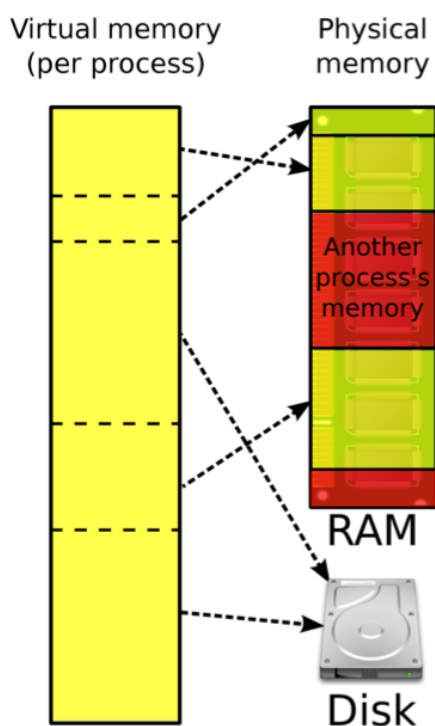
- The RAM memory allocated to a process is its address space
- It contains both the code for the running program, its input data and its working memory
- CPU interprets instructions - **%eip points to next instruction**

x86 process memory layout(simplified)



- Stack: from top to bottom
- Heap: from bottom to top

Virtual memory



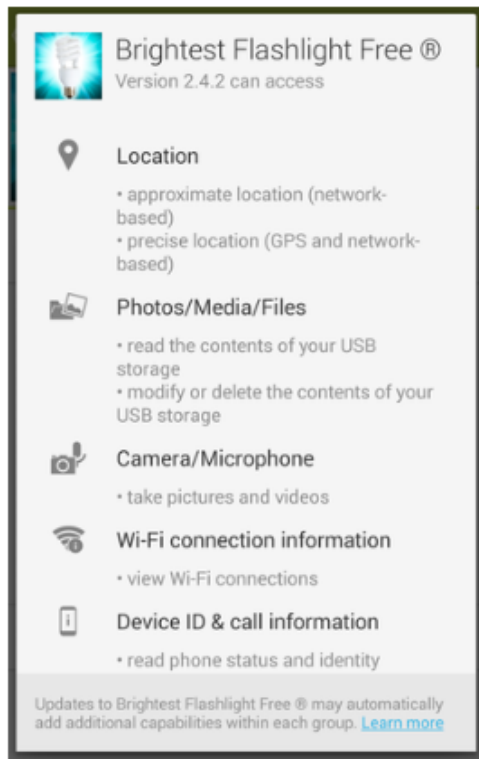
- Common technique used in a computer's OS
- Sometimes **available RAM is not enough to run several programs at one time**. This is where virtual memory comes in
- A system using virtual memory uses a section of the hard drive to emulate RAM - secondary memory treated as though it were main memory
- A **memory management unit (MMU)** maps a logical address space to a corresponding physical address
- **Live CD attacks on memory**
 - The attack:
 - Attacker with physical access to computer powers off the computer (without properly shutting down)
 - Attacker boots to different OS via external media
 - Attacker retrieves the *Pagefile.sys*, *Swapfile.sys*, *Hiberfile.sys* files
 - Attacker gains access to passwords and sensitive information that were stored in memory
 - Mitigation:
 - Hard disk encryption must be used!
- **Security principles**
 - **Defence-in-depth**



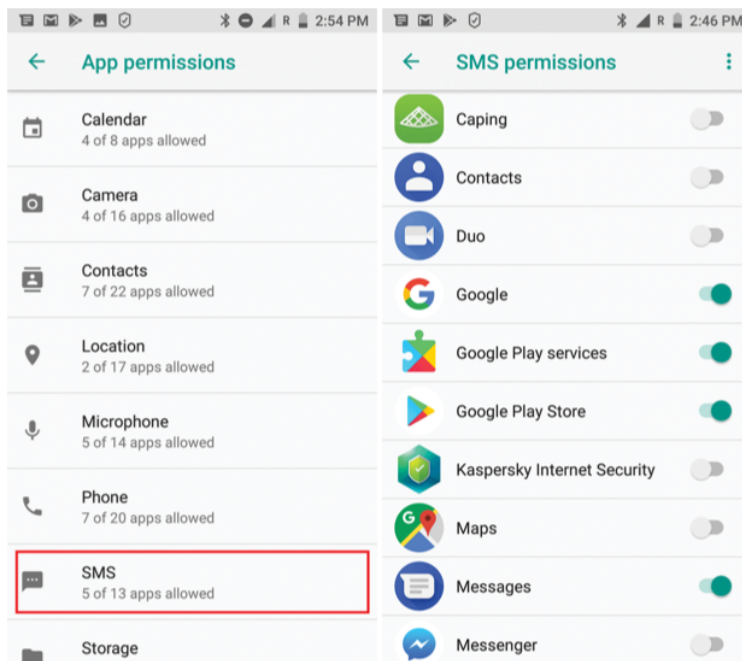
- Security protections built in multiple layers of the system: if one mechanism fails, another steps up immediately behind to thwart attacks
- Firewalls, intrusion detection and protection system, network segmentation, anti-virus, least privilege, strong password, path management
- **Least privilege**
- Users and programs should only access the data and resources required to perform its function



- A torch application does not need access to your location, photos, camera, microphone, wifi, device id, to perform its intended task!



Privilege separation



- Segment the system into components to which we can limit access
- Will limit the damage caused by a security break of any individual component

Open design

- The security of a mechanism should not depend on its secrecy

- The design and implementation details always get leaked
- **Economy of mechanism**
 - When designing a security mechanism keep it simple
 - It will facilitate the job of security researchers and allow verification
 - It will facilitate the task of developers and avoid bugs
 - It will facilitate the life of users and avoid misuses
- **Fail-safe defaults**
 - default configuration should be conservative, eg. new user should be granted least privileges by default
- **Complete mediation**
 - every access to a resource must be checked for compliance with security policy
- **Usable security**
 - UIs and security mechanisms should be designed with the ordinary user in mind – the users should be supported in interacting in a secure way with the system – you can't blame users !
- **What we learned today**
 - Many tasks handled by the OS relate to fundamental security problems
 - **1. OS concepts**
 - basic tasks of the OS
 - security concerns arise from multiple users and multiple processes
 - processes and process management
 - x86 runtime memory
 - **2. Security design principle**
 - Defence-in-depth
 - Least privilege

- Privilege separation
- Open design
- Economy of mechanism

- **Lecture 18 - Usable Security and Phishing**

- **Outline**

- What is "usable security"?
- An explanation of phishing
- Authentication in brief
- Passive vs. active indication

- **Usable security is challenging because:**

- Users are unmotivated to care about security over their current task
- Complex configurations make sense to computer scientists but not to end users
- Good feedback/advice is very hard to give to users
- Barn door - once security or privacy is lost, its gone. There is no undo

- **Phishing**

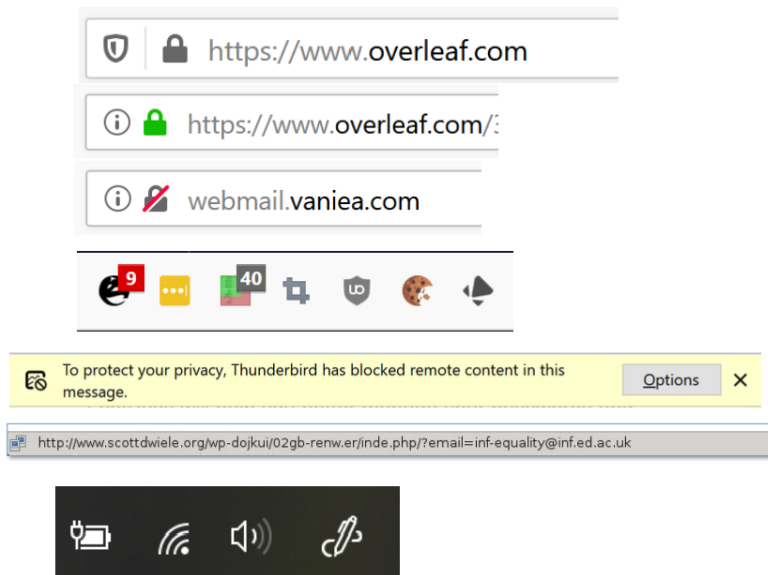
- Phishing is very common and very disruptive to UK businesses
- Also, it really annoys those of us who are just trying to get our work done.

- **Phishing Ecosystem**

- Limit the damage of credential sharing to one transaction
 - Let users authenticate websites
- Why does phishing work? Authentication is very broken
 - Authentication is how Entity A proves their identity to Entity B
 - We normally think of authentication as one directional
 - But it is actually two directional
 - The user must first make sure they are interacting with the “correct” website. Then the website must make sure that they are interacting with the “correct” user.

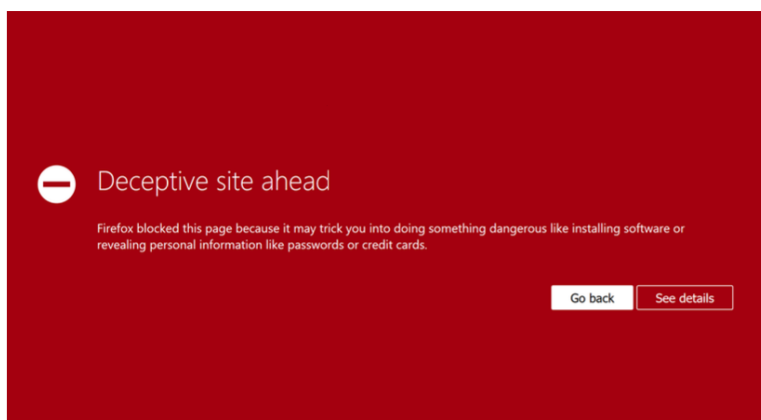
- Phishing Support(a history lesson)

- Passive indicator



- A UI element that provides information, but the user is not forced to look at or interact with.
- Phishing moves to email
 - Phishing moved off AOL and onto the less secure email. Directing people to fake sites, particularly fake financial sites.
 - Massive rise in identity theft
 - Financial loss skyrocketing

- Low conviction rate with "1-in-700 chance of escaping capture"
- Burdon falling on consumers
- **Recommend:**
 - Businesses should take security seriously
 - Financial organizations should auto identify fraudulent applications
 - Reduce impact on consumers
- Passive security indicators were not working at the level researchers wanted.
- So indicators started getting more obvious and intrusive.
- **Active Indicator**



- A UI element that interrupts the user's activity and demands a response.
- Active indicators work better than passive ones in terms of helping people avoid phishing.
- **Click through rates**
 - "Click through" – when a user sees a warning and chooses to proceed anyway.
 - Willingness to use Linux or use nightly builds of browsers indicates users are more willing to click through warnings.
 - A huge downside of active indicators is "habituation" where the user starts learning that the

warnings always happen and starts ignoring them.

- **Active indicators are alive and well in 2022**

- Screenshot of Santander payment page
- Asks payment purpose, then gives specific advice based on answer
- More customized to user needs, but still likely ignored

- **Where are users learning about security?**

- Users follow advice from sources they trust
- Users self-evaluate advice they feel they understand, like password advice
- Marketing material in the advice results in less trust
- High socioeconomic users get phishing training at work and tend to follow it
- Low socioeconomic users tend to get and follow advice from friends and service providers (ISPs)

- **Many different types of advice given in guidance**

- **Lookout for:**

- Requests for sensitive data
- Poor grammar
- Unusual senders
- Use of an alarming tone
- Has a link to a website
- Content is account related
- Content too good to be true

- **Protection actions:**

- Check the URL in the address bar
- Check for HTTPS
- Type the URL yourself
- Check for a lock icon
- Bookmark sensitive websites

- Inclusion
 - Usable security
 - Harder than it looks
 - Phishing only requires one side of the two way authentication to fail
 - Passive indicators
 - Show information but do not block user tasks
 - Are easily ignored by users
 - Active indicators
 - Block the user till they interact with the dialog in some way
 - Much more effective than passive indicators
 - Lead to habituation if a user sees the warning frequently, they stop reading it

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