# Privilege separation

## - UNIX security model -

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# Privilege separation

Modern computers are

1. multi-users
2. multi-tasking

Goal: Prevent potentially misbehaving users and/or applications from harming the rest of the system

Permissions system: mechanisms for achieving separation between components

# Central question

"Who is allowed to access what and how?"

The subject (who) - *eg.* user, application, process

The object (what) - protected resource, *eg.* hardware device, network socket, memory, files, directories, *etc.*

The access operation (how) - *eg.* read, write, execute

# Key assumptions for separation

1. The system know who the user is - user has authenticated, *e.g.* using username / password
2. Complete mediation - all requests are mediated - all requests go to the reference monitor that enforces specified access control policies


The **reference monitor** grants permission to users to apply certain operations to a given resource

# Users

Two types of accounts each with a unique identifier, the user ID (`uid`):

1. User accounts - associated with humans
2. Service accounts - associated with background processes

```
marapini@myrto-thinkpad:~$ more /etc/passwd
root:x:0:0:root:/root:/bin/bash
daemon:x:1:1:daemon:/usr/sbin:/usr/sbin/nologin
bin:x:2:2:bin:/bin:/usr/sbin/nologin
sys:x:3:3:sys:/dev:/usr/sbin/nologin
sync:x:4:65534:sync:/bin:/bin/sync
games:x:5:60:games:/usr/games:/usr/sbin/nologin
man:x:6:12:man:/var/cache/man:/usr/sbin/nologin
lp:x:7:7:lp:/var/spool/lpd:/usr/sbin/nologin
mail:x:8:8:mail:/var/mail:/usr/sbin/nologin
news:x:9:9:news:/var/spool/news:/usr/sbin/nologin
uucp:x:10:10:uucp:/var/spool/uucp:/usr/sbin/nologin
proxy:x:13:13:proxy:/bin:/usr/sbin/nologin
```

- One entry in the `/etc/passwd` per account with the fields:

  `username:password:uid:gid:uid_info:home:shell`

- `uid 0` - user root `uid`

# Groups

- Groups are sets of users that share resources
- Every group has a name and a unique identifier, the group ID (`gid`)
- Allow for easier users management and monitoring

```
marapini@myrto-thinkpad:~$ more /etc/group
root:x:0:
daemon:x:1:
bin:x:2:
sys:x:3:
adm:x:4:syslog,marapini
tty:x:5:
disk:x:6:
lp:x:7:
mail:x:8:
```

- One entry in the /etc/group per group with the fields:

`group_name:password:gid:group_list`

# File permissions

- All resources (sockets, directories, files) are managed as files
- 3 defined permissions: read (r), write (w), execute (x)
- Permissions are defined for the owner, the owner's group, and other users
- Root and owner can change file permissions
- Only root can change file ownership

```
marapini@myrto-thinkpad:~/Documents/Work/Teaching/INFR10067-ComputerSecurity/2021/Lectures/L18.AccessControl$ ls -l
total 352
drwxrwxr-x 2 marapini marapini   4096 Feb 23 17:27 Images
-rw-r--r-- 1 marapini marapini   1839 Feb 23 17:31 L18.AccessControl.aux
-rw-r--r-- 1 marapini marapini  47121 Feb 23 17:31 L18.AccessControl.log
-rw-r--r-- 1 marapini marapini    835 Feb 23 17:31 L18.AccessControl.nav
-rw-r--r-- 1 marapini marapini      0 Feb 23 17:31 L18.AccessControl.out
-rw-r--r-- 1 marapini marapini 258111 Feb 23 17:31 L18.AccessControl.pdf
-rw-r--r-- 1 marapini marapini      0 Feb 23 17:31 L18.AccessControl.snm
-rw-rw-r-- 1 marapini marapini   9769 Feb 23 18:05 L18.AccessControl.tex
-rw-rw-r-- 1 marapini marapini  23638 Feb 20 01:28 L18.AccessControl.tex~
-rw-r--r-- 1 marapini marapini      0 Feb 23 17:31 L18.AccessControl.toc
```

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Quizz: Imagine you have the following groups:

- `infr10067` - for any user involved with the Computer Security course
- `tas` - for all Informatics TAs

How can you have a folder only for Computer Security TAs?

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- `infr10067` - for any user involved with the Computer Security course
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How can you have a folder only for Computer Security TAs?

```
marapini@myrto-thinkpad:~/Documents/Work/Teaching/INFR10067-ComputerSecurity/2021/Lectures/L18.AccessControl/conjunction$ ls -l
total 4
drwxr-xr-- 3 marapini tas 4096 Feb 23 22:50 only_for_tas
marapini@myrto-thinkpad:~/Documents/Work/Teaching/INFR10067-ComputerSecurity/2021/Lectures/L18.AccessControl/conjunction$ ls -l only_for_tas/
total 4
drwxr-xr-- 2 marapini infr10067 4096 Feb 23 22:50 only_for_infr10067_tas
marapini@myrto-thinkpad:~/Documents/Work/Teaching/INFR10067-ComputerSecurity/2021/Lectures/L18.AccessControl/conjunction$
```

# Processes

- Each process has a unique identifier, the process ID (`pid`)
- Each process is associated with the user that spanned it

```
marapini@myrto-thinkpad:~$ ps -ef
UID          PID  PPID  C STIME TTY          TIME CMD
root           1     0  0 Feb22 ?        00:00:43 /sbin/init splash
root           2     0  0 Feb22 ?        00:00:00 [kthreadd]
root           4     2  0 Feb22 ?        00:00:00 [kworker/0:0H]
root           6     2  0 Feb22 ?        00:00:00 [mm_percpu_wq]
root           7     2  0 Feb22 ?        00:00:00 [ksoftirqd/0]
root           8     2  0 Feb22 ?        00:00:12 [rcu_sched]
root           9     2  0 Feb22 ?        00:00:00 [rcu_bh]
root          10     2  0 Feb22 ?        00:00:00 [migration/0]
root          11     2  0 Feb22 ?        00:00:00 [watchdog/0]
root          12     2  0 Feb22 ?        00:00:00 [cpuhp/0]
root          13     2  0 Feb22 ?        00:00:00 [cpuhp/1]
root          14     2  0 Feb22 ?        00:00:00 [watchdog/1]
```

- When a user runs a process, it runs with that user's privileges, *i.e.* they can access any resource that user has permissions for
- By default, a child process inherits its parent's privileges
- Processes are isolated in memory

# Process user IDs

Every process has:

- Real user ID (`uid`) - the user ID that started that process
- Effective user ID (`euid`) - the user ID that determines the process' privileges
- Saved user ID (`suid`) - the effective user ID before the last modification

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Users can change a process' IDs:

```
setuid(x)          seteuid(x)
  uid ← x            uid ← uid
  euid ← x           euid ← x
  suid ← x           suid ← suid
```

- Root can change `euid`/`uid` to arbitrary values `x`:
- Unprivileged users can only change `euid` to `uid` or `suid`:

# Dropping privileges with `setuid`

Imagine a program that runs as root and wants to fork a process with lower privileges using the following code:

```
if (auth(uid, pwd) == SUCCESS) {
  if (fork() == 0) {
    seteuid(uid);
    exec("/bin/bash");
  }
}
```

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    exec("/bin/bash");
  }
}
```

← the user can call `seteuid(0)` and become root!!

# Dropping privileges with `setuid`

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```
if (auth(uid, pwd) == SUCCESS) {
  if (fork() == 0) {
    setuid(uid);
    exec("/bin/bash");                    ← the user cannot
  }                                         change uid
}
```

# Elevating privileges - `setuid` programs

- An executable file can have the `set-user-ID` property (`setuid`) enabled
- If A executes a `setuid` file owned by B, then the `euid` of the process is B and not A
- Writing secure `setuid` programs is tricky because vulnerabilities may be exploited by malicious user actions

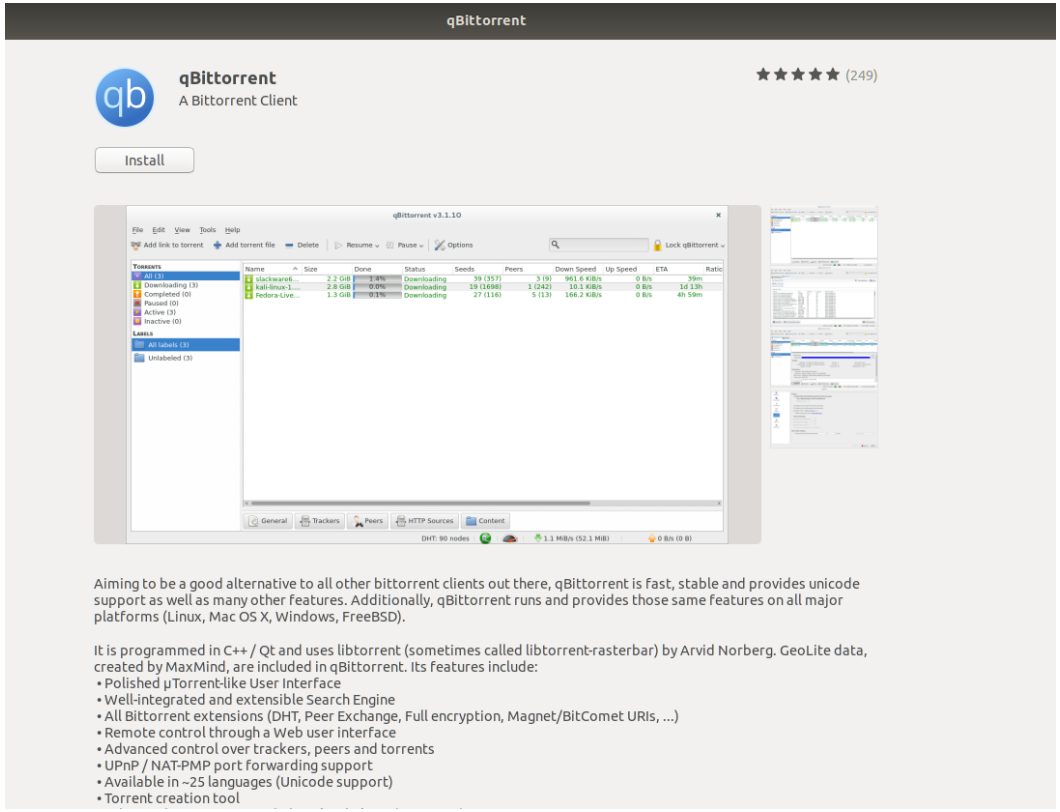
- Some programs that access system resources are owned by root and have the `setuid` bit set (`setuid` programs)

```
marapini@myrto-thinkpad:/usr/bin$ ls -l | grep passwd
-rwsr-xr-x 1 root root       75824 Jan 25  2018 gpasswd
-rwxr-xr-x 1 root root      249976 Feb  7 23:20 grub-mkpasswd-pbkdf2
-rwsr-xr-x 1 root root       59640 Jan 25  2018 passwd
marapini@myrto-thinkpad:/usr/bin$
```

```
marapini@myrto-thinkpad:/etc$ ls -l | grep shadow
-rw-r-----  1 root shadow     860 Feb 27 10:11 gshadow
-rw-r-----  1 root shadow     845 Sep 21 15:46 gshadow-
-rw-r-----  1 root shadow    1373 Feb 27 10:11 shadow
-rw-r-----  1 root shadow    1373 Feb 27 10:11 shadow-
marapini@myrto-thinkpad:/etc$
```

# UNIX permissions are too coarse-grained

All application installed by a single user account have the same privileges!



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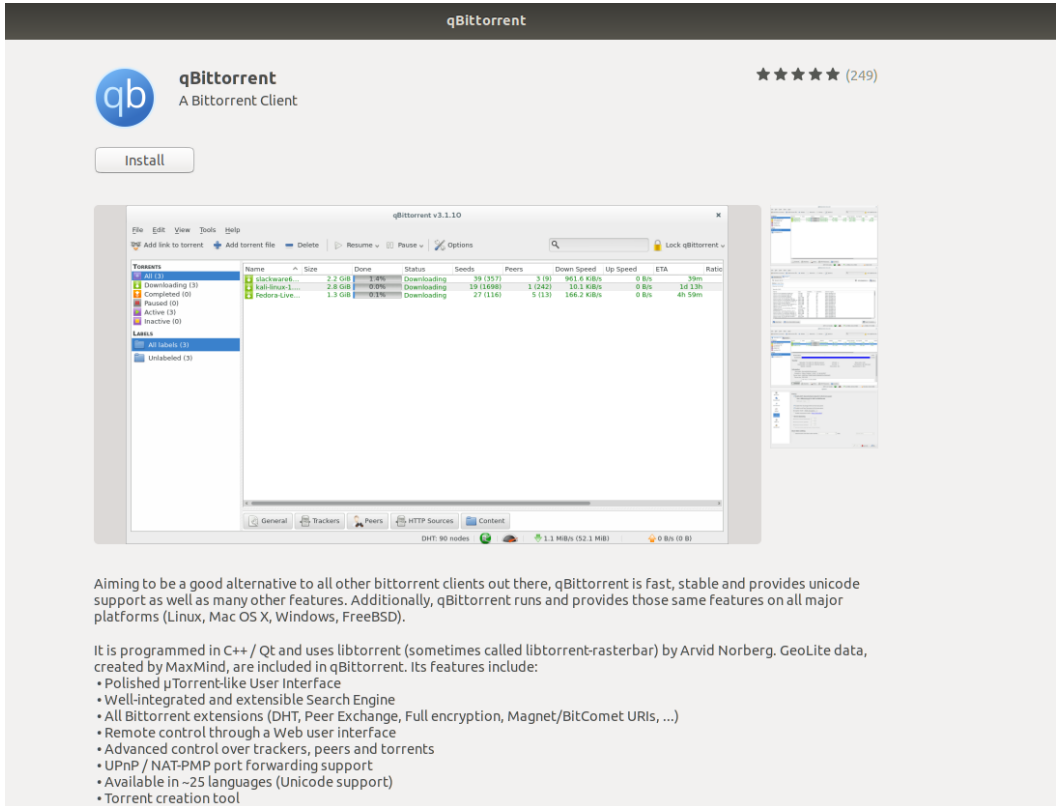


!!? What if qBittorent is malware ?!!

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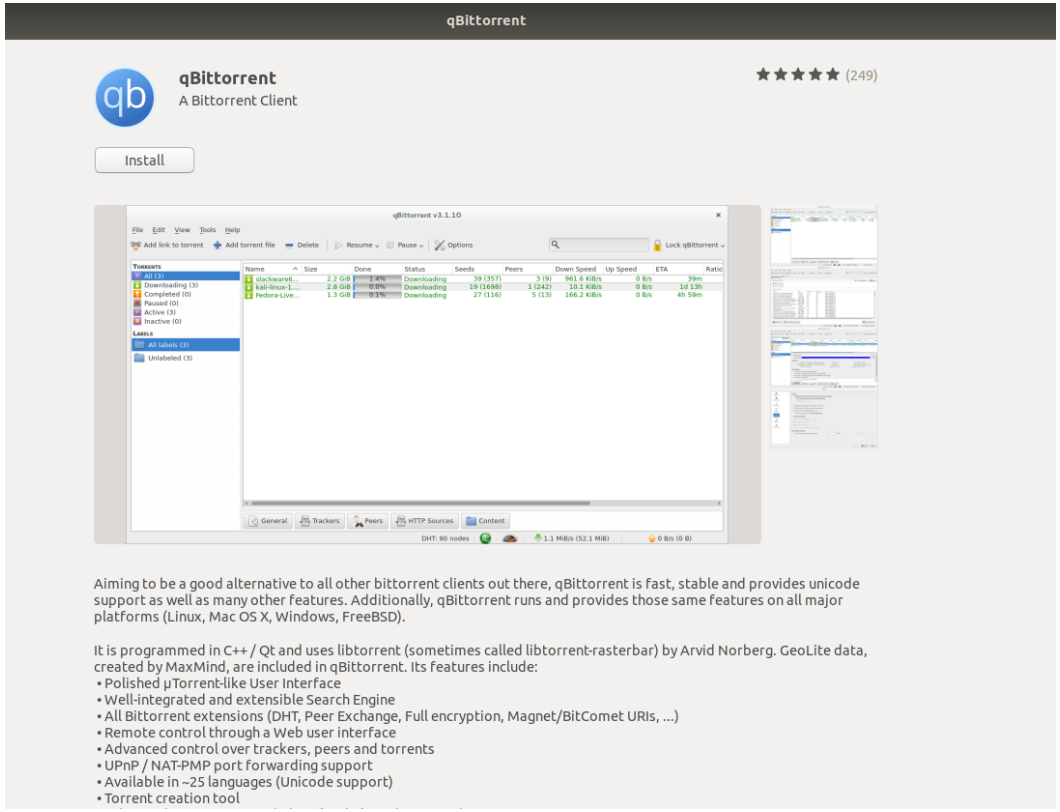
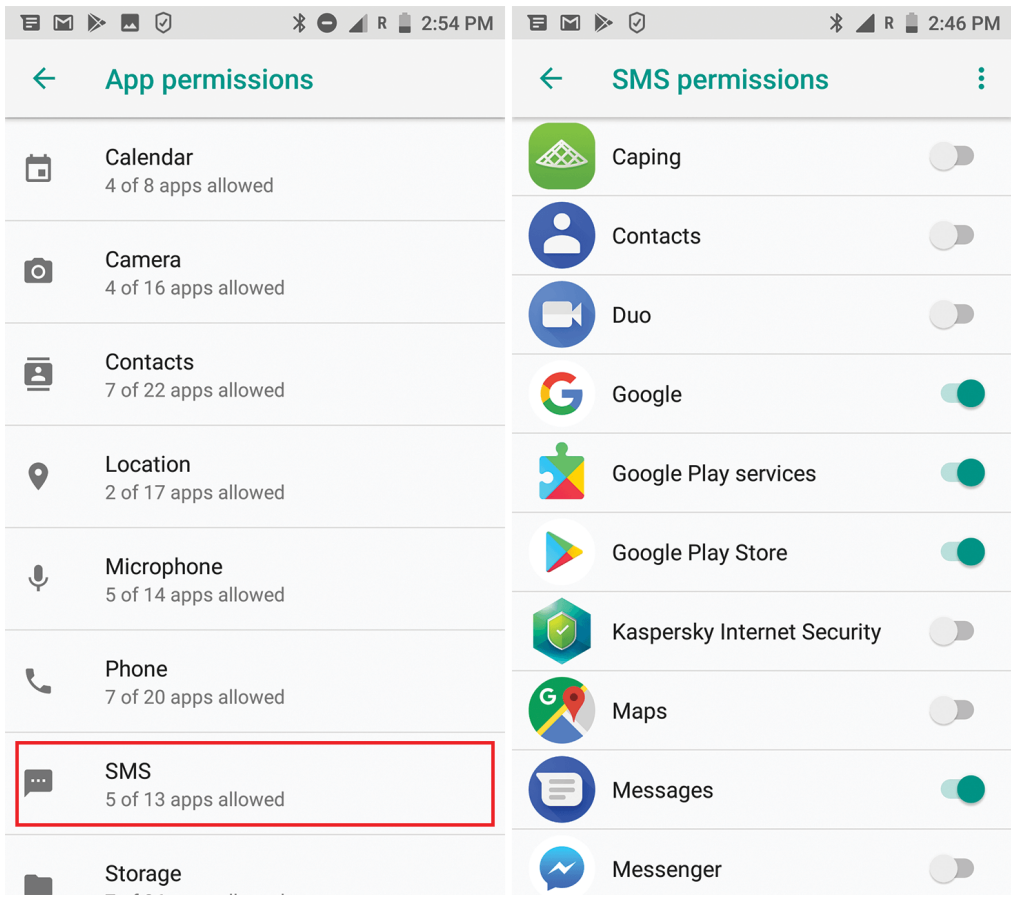



!!? What if qBittorent is malware ?!!

- Better delegate capabilities associated with specific root powers

# Android permissions

- Each app runs with a different user ID
- Apps do not interact
- Permissions are set per app



# Take aways

The UNIX security model provides a simple and flexible model, **but permissions are too coarse-grained**:

- same permissions for all applications ran under a single user account
- many utilities have the `setuid` bit enabled

→ many opportunities for privilege escalation attacks

→ better use capabilities when delegating privileges