

# Chapter 8

## Security

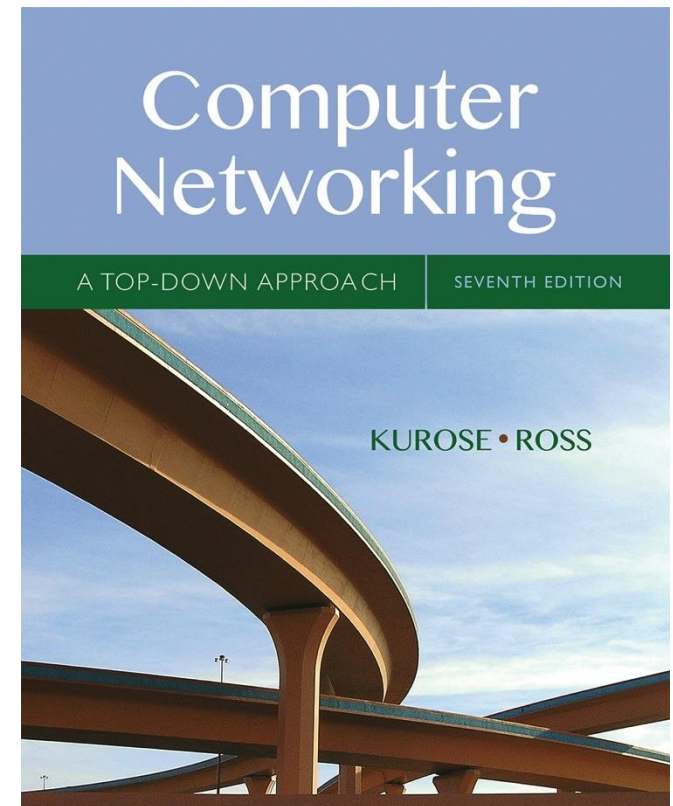
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## Computer Networking: A Top Down Approach

7<sup>th</sup> edition

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# Chapter 8: Network Security

## *Chapter goals:*

- understand principles of network security:
  - cryptography and its *many* uses beyond “confidentiality” (保密)
  - authentication (真实)
  - message integrity (完整)
- security in practice:
  - firewalls and intrusion detection systems
  - security in application, transport, network, link layers

# Chapter 8 roadmap

*8.1 What is network security?*

8.2 Principles of cryptography

8.3 Message integrity, authentication

8.4 Securing e-mail

8.5 Securing TCP connections: SSL

8.6 Network layer security: IPsec

8.7 Securing wireless LANs

8.8 Operational security: firewalls and IDS

# What is network security?

**Confidentiality:** only sender, intended receiver should “understand” message contents

- sender encrypts message
- receiver decrypts message

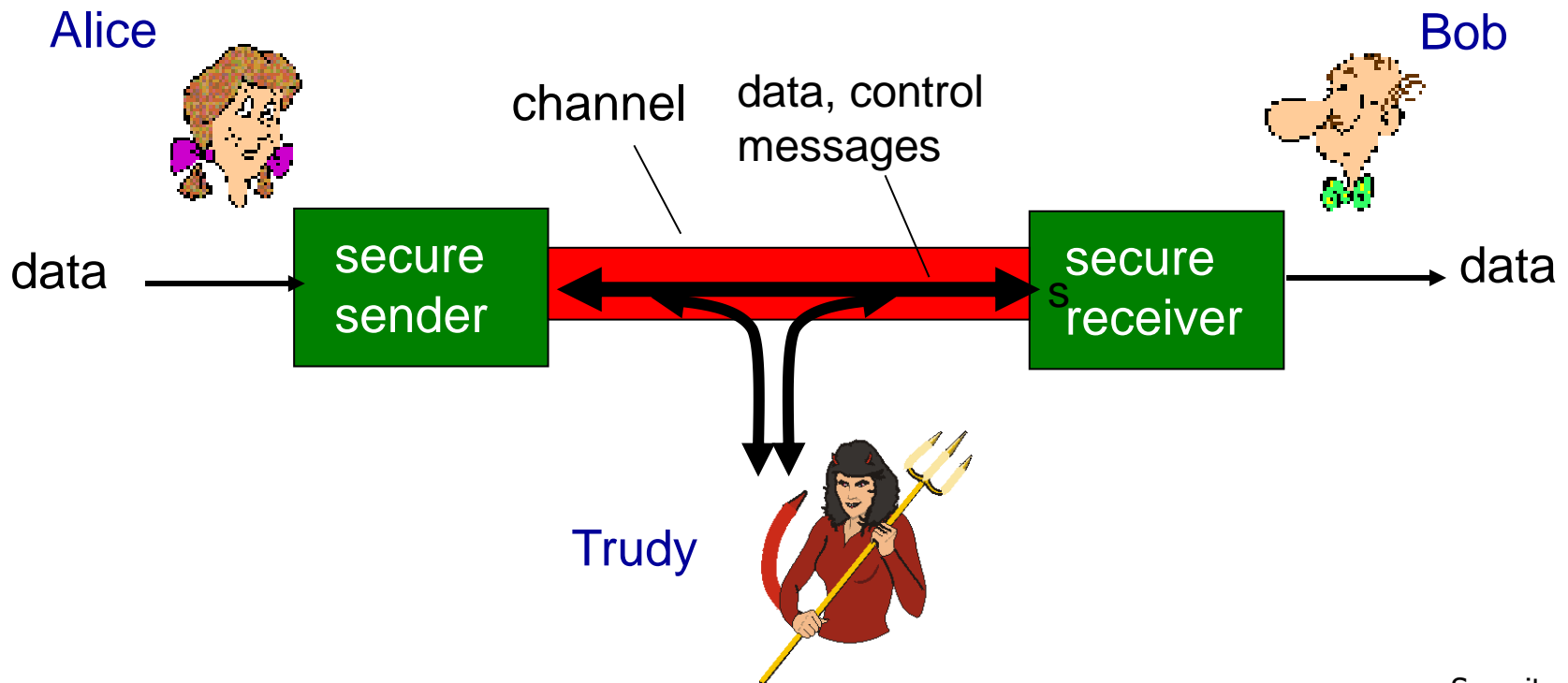
**authentication:** sender, receiver want to confirm identity of each other

**message integrity:** sender, receiver want to ensure message not altered (in transit, or afterwards) without detection

**access and availability:** services must be accessible and available to users

# Friends and enemies: Alice, Bob, Trudy

- well-known in network security world
- Bob, Alice (lovers!) want to communicate “securely”
- Trudy (intruder) may intercept, delete, add messages



# Who might Bob, Alice be?

- ... well, *real-life* Bobs and Alices!
- Web browser/server for electronic transactions (e.g., on-line purchases)
- on-line banking client/server
- DNS servers
- routers exchanging routing table updates
- other examples?

# There are bad guys (and girls) out there!

Q: What can a “bad guy” do?

A: A lot! See section 1.6

- *eavesdrop*: intercept messages
- actively *insert* messages into connection
- *impersonation*: can fake (spoof) source address in packet (or any field in packet)
- *hijacking*: “take over” ongoing connection by removing sender or receiver, inserting himself in place
- *denial of service*: prevent service from being used by others (e.g., by overloading resources)

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8.2 *Principles of cryptography*

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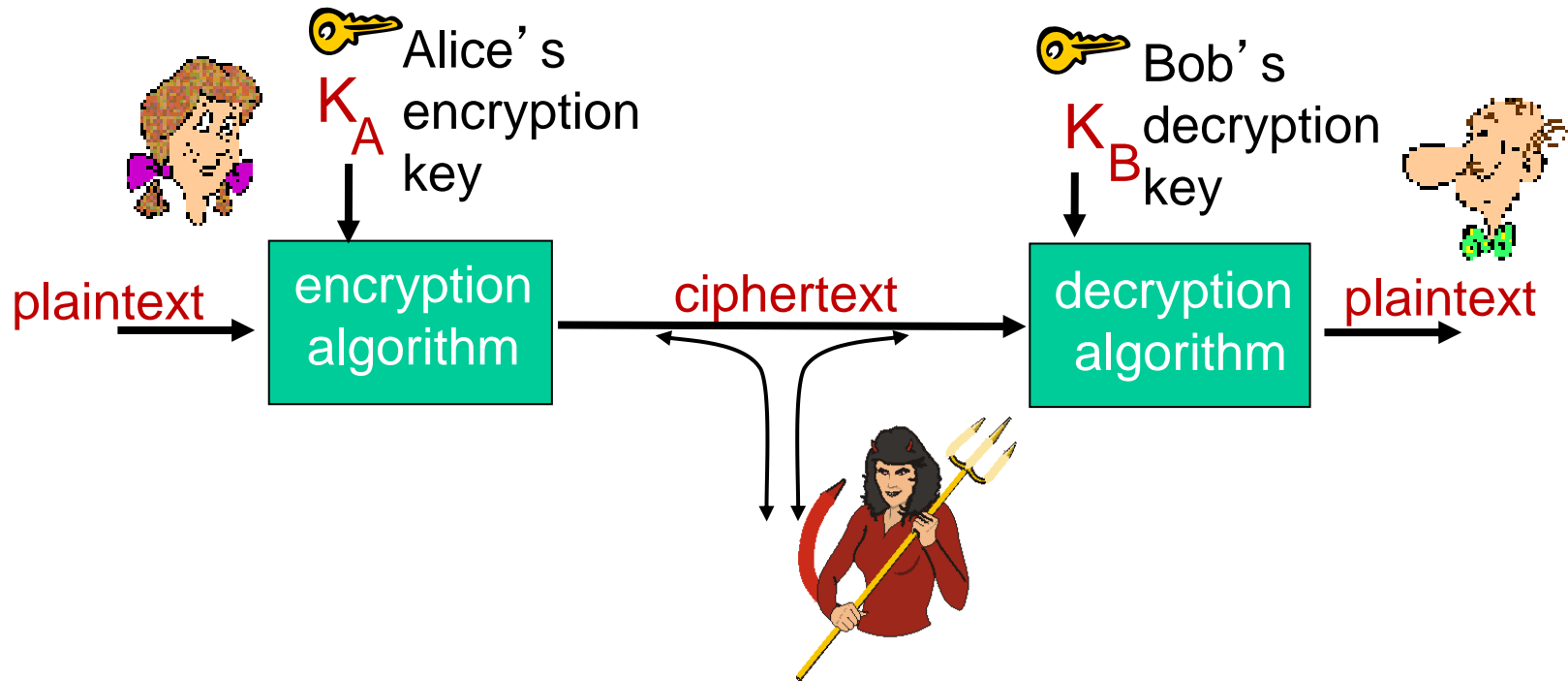
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# The language of cryptography



$m$  plaintext message

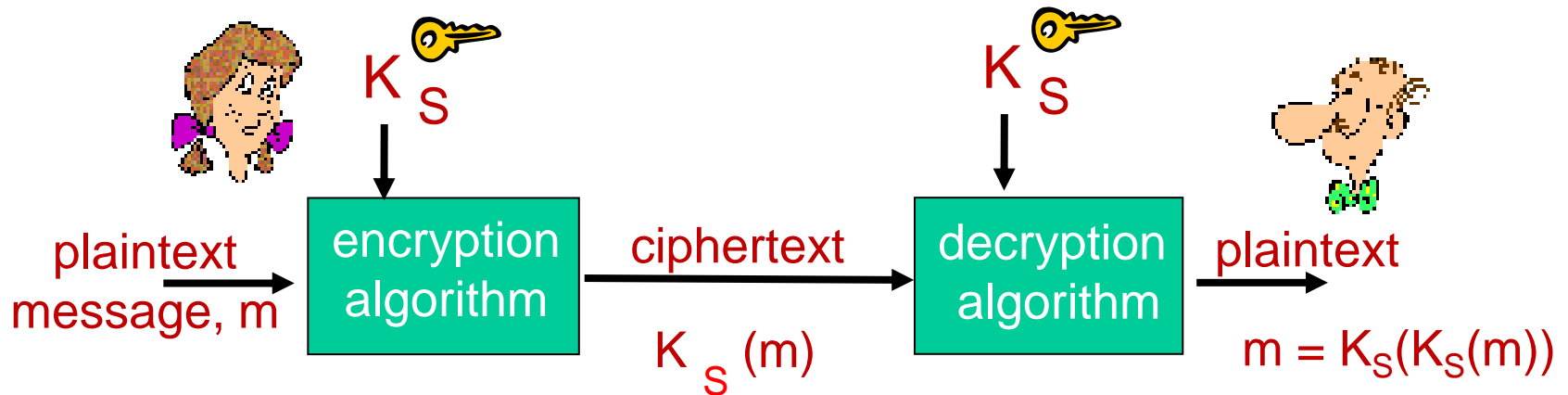
$K_A(m)$  ciphertext, encrypted with key  $K_A$

$m = K_B(K_A(m))$

# Breaking an encryption scheme

- **cipher-text only attack:**  
Trudy has ciphertext she can analyze
- **two approaches:**
  - brute force: search through all keys
  - statistical analysis
- **known-plaintext attack:**  
Trudy has plaintext corresponding to ciphertext
  - e.g., in monoalphabetic cipher, Trudy determines pairings for a,l,i,c,e,b,o,
- **chosen-plaintext attack:**  
Trudy can get ciphertext for chosen plaintext

# Symmetric key cryptography



**symmetric key crypto:** Bob and Alice share same (symmetric) key:  $K_S$

- e.g., key is knowing substitution pattern in mono alphabetic substitution cipher

Q: how do Bob and Alice agree on key value?

# Simple encryption scheme

*substitution cipher*: substituting one thing for another

- monoalphabetic cipher: substitute one letter for another (Caesar cipher)

plaintext:   abcdefghijklmnopqrstuvwxyz

letter after 12 letters

ciphertext: **mnbvcxz**asdfghjklpoiuytrewq

e.g.: Plaintext: bob. i love you. alice

**ciphertext:** nkn. s gktc wky. mgsbc

🔑 **Encryption key:** mapping from set of 26 letters to set of 26 letters

# A more sophisticated encryption approach

- n substitution ciphers,  $M_1, M_2, \dots, M_n$
- cycling pattern:
  - e.g.,  $n=4$ :  $M_1, M_3, M_4, M_3, M_2$ ;  $M_1, M_3, M_4, M_3, M_2$ ; ..
- for each new plaintext symbol, use subsequent substitution pattern in cyclic pattern
  - dog: d from  $M_1$ , o from  $M_3$ , g from  $M_4$

*Encryption key:* n substitution ciphers, and cyclic pattern



- key need not be just n-bit pattern

# Polyalphabetic encryption

Plaintext letter:	a b c d e f g h i j k l m n o p q r s t u v w x y z
$C_1(k = 5)$ :	f g h i j k l m n o p q r s t u v w x y z a b c d e
$C_2(k = 19)$ :	t u v w x y z a b c d e f g h i j k l m n o p q r s

Using two Caesar ciphers

Repeating pattern:  $C_1, C_2, C_2, C_1, C_2$

**Plaintext:** bob. i love you.

**ciphertext:** ghu. n etox dhz.

# Block cipher

- ❖ Message is encrypted in blocks of  $k$  bits.

$k=3$

input	output	input	output
000	110	100	011
001	111	101	010
010	101	110	000
011	100	111	001

- ❖ Large  $k$  values against brute-force attack
  - ❖  $2^k$  possible mapping
  - ❖ E.g,  $k=1024$
- ❖ Problem: maintaining a table of  $2^k$  input values
  - ❖ Hard to maintain and update

# Block cipher

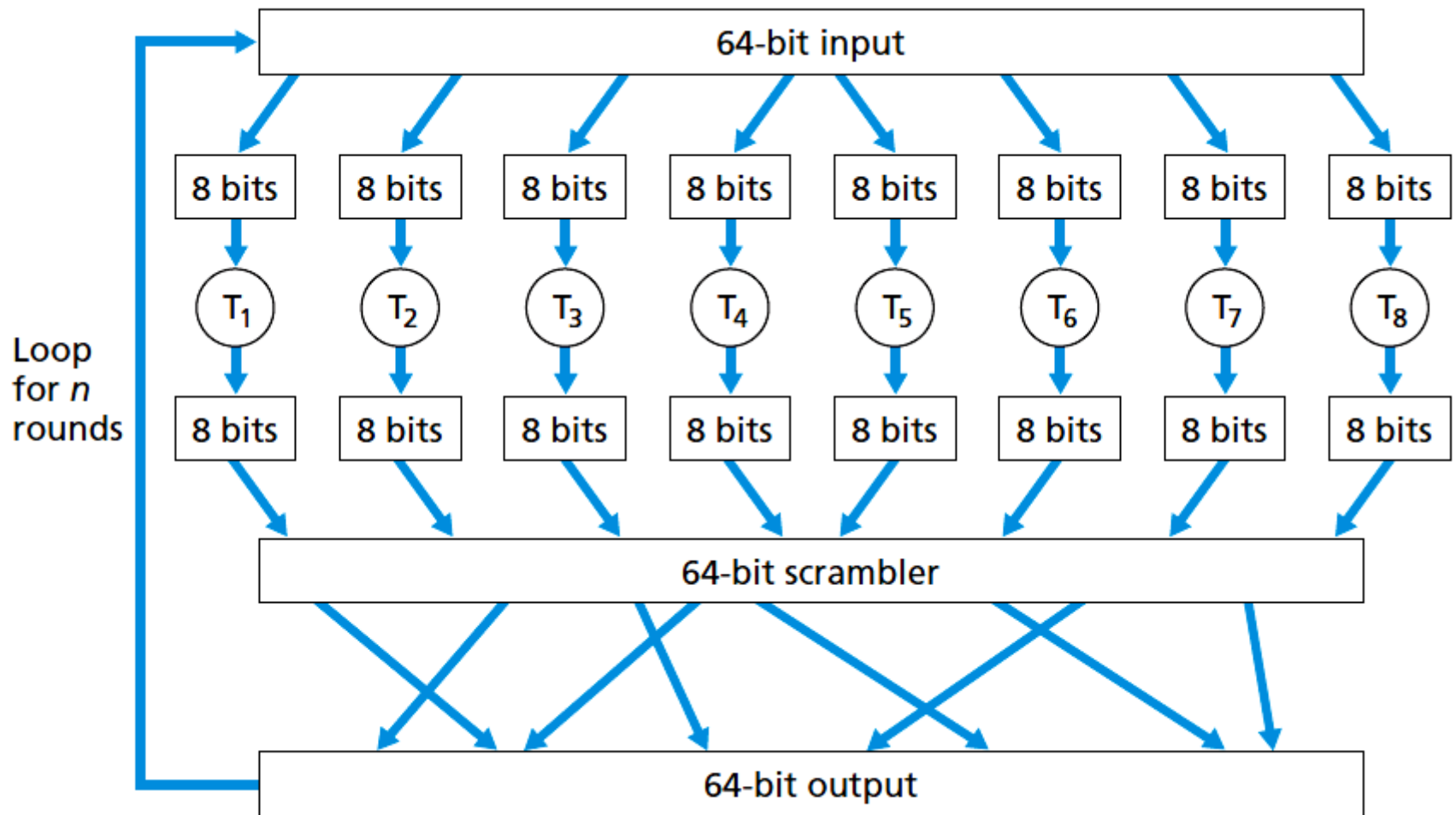
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- ❖ Example:  $K=64$

1. Break a 64-bit block into 8 chunks, each 8-bit chunk is processed by an 8-bit to 8-bit table
  2. The 8 output chunks are reassembled into a 64-bit block.
  3. The positions of the 64 bits in the block are then scrambled (permuted) to produce a 64-bit output.
  4. This output is fed back to the 64-bit input, where another cycle begins.
  5. After  $n$  cycles, the function provides a 64-bit block of ciphertext.
- ❖ The key for this block cipher algorithm would be the eight permutation tables
    - ❖ scramble function is publicly known



# Block cipher



# Cipher-block chaining

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- ❖ Problem of block cipher: two or more of the cleartext blocks can be identical.
  - ❖ E.g., 200 OK
- ❖ The sender creates a random  $k$ -bit number  $r(i)$  for the  $i$ th block and calculates  $c(i) = K_s(m(i) \oplus r(i))$ . ( $\oplus$  means exclusive or (异或))
  - ❖ A new  $k$ -bit random number is chosen for each block.
- ❖ Receiver receives  $c(i)$  and  $r(i)$ , it can recover each block of the plaintext by computing  $m(i) = K_d(c(i)) \oplus r(i)$ .
- ❖ Example:
  - ❖  $k=3$ , plain text 010 010 010
  - ❖  $r(1)=001$   $r(2)=111$   $r(3)=100$
  - ❖ Produce different ciphertext for same plaintext

# Cipher-block chaining

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- ❖ Problem: need to transmit  $r(i)$  from sender to receiver
  - ❖ Transmit twice as many bits as before
- ❖ Cipher-block chaining (CBC)
  - ❖ Sender generates a  $k$ -bit string, **Initialization Vector** (IV), denoted as  $c(0)$ , sent IV to receiver.
  - ❖ Runs block cipher algorithm  $c(1)=K_s(m(1)\oplus c(0))$ , send  $c(1)$  to receiver.
  - ❖  $c(i)=K_s(m(i)\oplus c(i-1))$
  - ❖ Receiver: knows  $c(i-1)$ , obtain  $m(i)=K_s(c(i))\oplus c(i-1)$

# Cipher-block chaining: example

input	output	input	output
000	110	100	011
001	111	101	010
010	101	110	000
011	100	111	001

- ❖ plain text 010 010 010
- ❖  $r(0)=001$
- ❖  $c(1)=K_s(010 \oplus 001)=K_s(011)=100$
- ❖  $c(2)=K_s(010 \oplus 100)=K_s(110)=000$
- ❖  $c(3)=K_s(010 \oplus 000)=K_s(010)=101$

# Block cipher

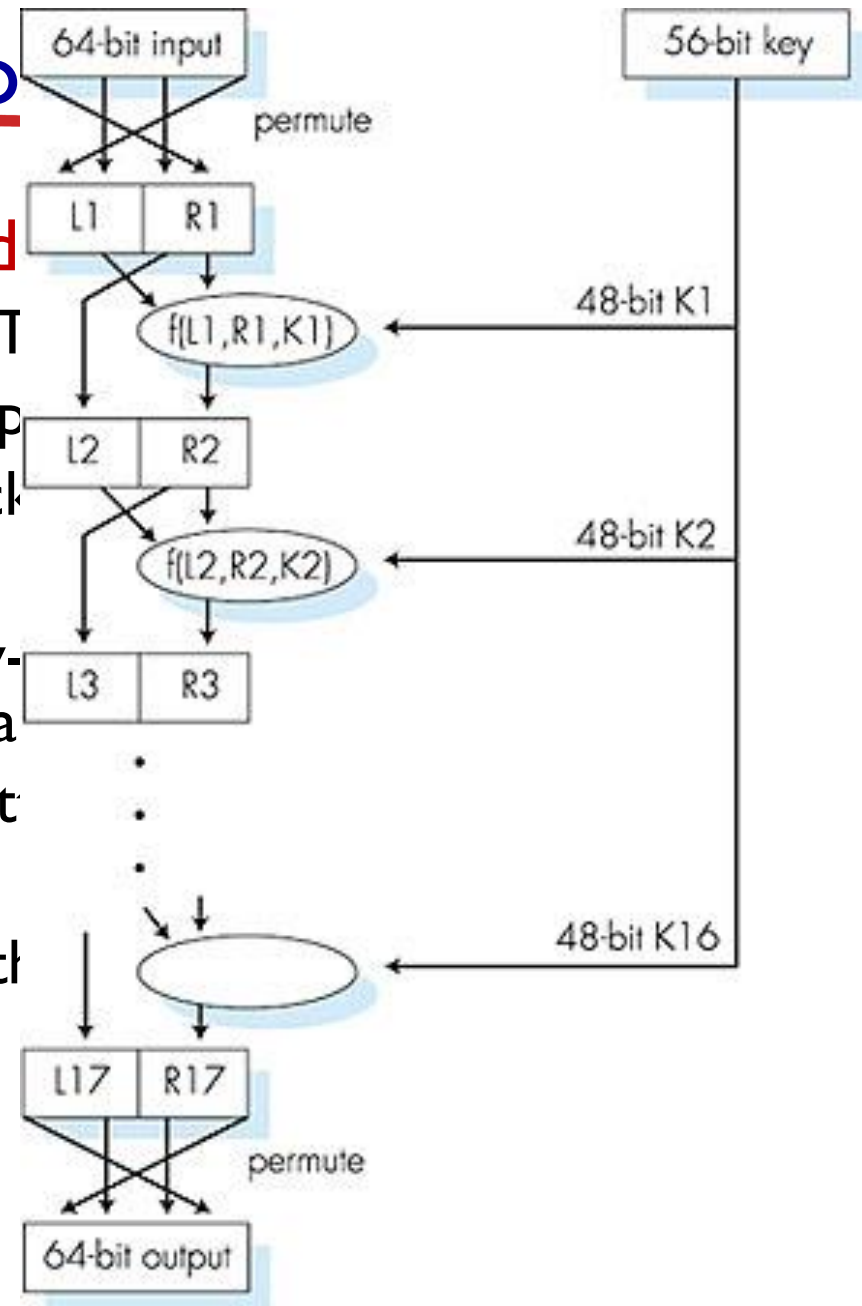
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- ❖ Standards
  - ❖ Use functions rather than tables
  - ❖ Use a string of bits as key: determines table mappings, permutations, and loops.
- ❖ DES uses 64-bit blocks with a 56-bit key.
- ❖ AES uses 128-bit blocks and can operate with keys that are 128, 192, and 256 bits long.

# Symmetric key crypto

## DES: Data Encryption Standard

- US encryption standard [NIST]
- 56-bit symmetric key, 64-bit plaintext block
- block cipher with cipher block size of 64 bits
- how secure is DES?
  - DES Challenge: 56-bit-key - (brute force) in less than a year
  - no known good analytic attack
- making DES more secure:
  - 3DES: encrypt 3 times with different keys



# AES: Advanced Encryption Standard

- symmetric-key NIST standard, replaced DES (Nov 2001)
- processes data in 128 bit blocks
- 128, 192, or 256 bit keys
- brute force decryption (try each key) taking 1 sec on DES, takes 149 trillion years for AES

# Public Key Cryptography



## *symmetric key crypto*

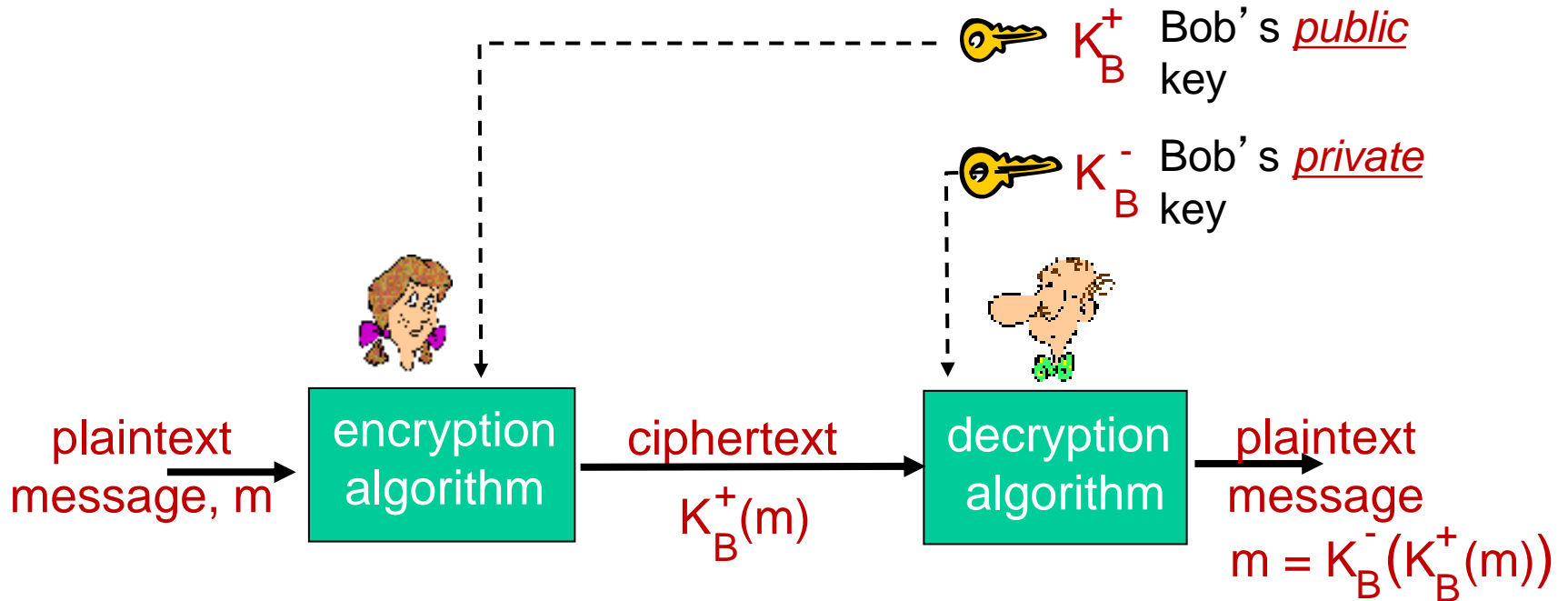
- requires sender, receiver know shared secret key
- Q: how to agree on key in first place (particularly if never “met”)?

## *public key crypto*

- radically different approach [Diffie-Hellman76, RSA78]
- sender, receiver do *not* share secret key
- *public* encryption key known to *all*
- *private* decryption key known only to receiver



# Public key cryptography



# Public key encryption algorithms

requirements:

- ① need  $K_B^+(\cdot)$  and  $K_B^-(\cdot)$  such that

$$K_B^-(K_B^+(m)) = m$$

- ② given public key  $K_B^+$ , it should be impossible to compute private key  $K_B^-$

**RSA:** Rivest, Shamir, Adelson algorithm

# Prerequisite: modular arithmetic

- $x \bmod n$  = remainder of  $x$  when divide by  $n$

- facts:

$$[(a \bmod n) + (b \bmod n)] \bmod n = (a+b) \bmod n$$

$$[(a \bmod n) - (b \bmod n)] \bmod n = (a-b) \bmod n$$

$$[(a \bmod n) * (b \bmod n)] \bmod n = (a*b) \bmod n$$

- thus

$$(a \bmod n)^d \bmod n = a^d \bmod n$$

- example:  $x=14$ ,  $n=10$ ,  $d=2$ :

$$(x \bmod n)^d \bmod n = 4^2 \bmod 10 = 6$$

$$x^d = 14^2 = 196 \quad x^d \bmod 10 = 6$$

# RSA: getting ready

- message: just a bit pattern
- bit pattern can be uniquely represented by an integer number
- thus, encrypting a message is equivalent to encrypting a number

## *example:*

- $m = 10010001$ . This message is uniquely represented by the decimal number 145.
- to encrypt  $m$ , we encrypt the corresponding number, which gives a new number (the ciphertext).

# RSA: Creating public/private key pair

1. choose two large prime numbers (质数)  $p, q$ .  
(e.g., 1024 bits each)
2. compute  $n = pq$ ,  $z = (p-1)(q-1)$
3. choose  $e$  (with  $e < n$ ) that has no common factors with  $z$  ( $e, z$  are “relatively prime (互质)”).
4. choose  $d$  such that  $ed-1$  is exactly divisible by  $z$ .  
(in other words:  $ed \bmod z = 1$ ).
5. public key is  $\underbrace{(n, e)}_{K_B^+}$ . private key is  $\underbrace{(n, d)}_{K_B^-}$ .

# RSA: encryption, decryption

0. given  $(n,e)$  and  $(n,d)$  as computed above

1. to encrypt message  $m$  ( $<n$ ), compute

$$c = m^e \bmod n$$

2. to decrypt received bit pattern,  $c$ , compute

$$m = c^d \bmod n$$

*magic happens!*

$$m = \underbrace{(m^e \bmod n)}_c^d \bmod n$$

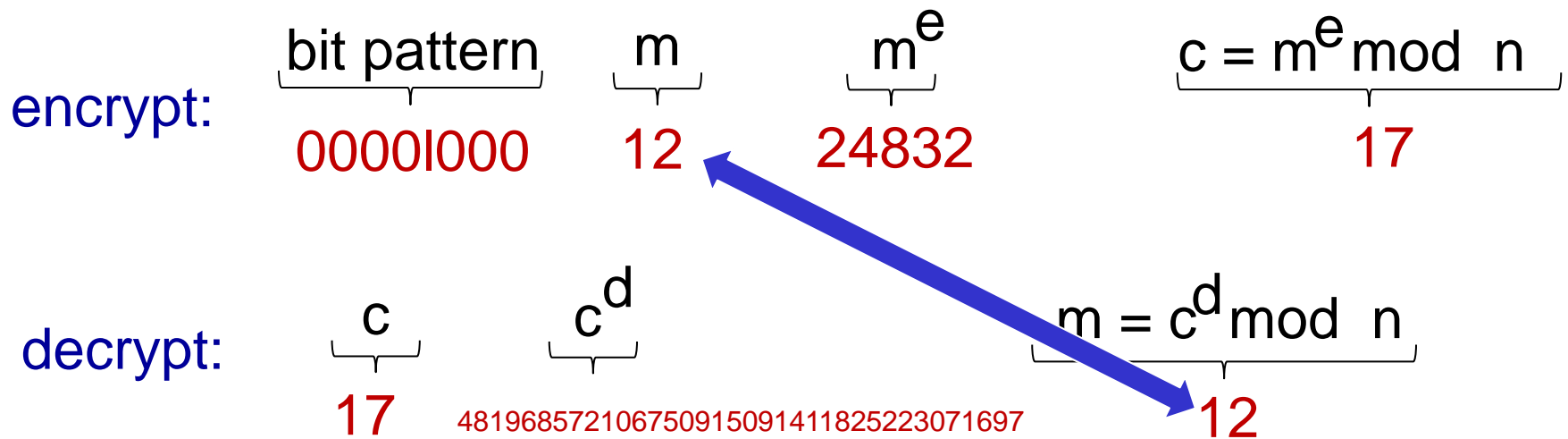
# RSA example:

Bob chooses  $p=5$ ,  $q=7$ . Then  $n=35$ ,  $z=24$ .

$e=5$  (so  $e$ ,  $z$  relatively prime).

$d=29$  (so  $ed-1$  exactly divisible by  $z$ ).

encrypting 8-bit messages.



# Why does RSA work?

- must show that  $c^d \bmod n = m$   
where  $c = m^e \bmod n$
- fact: for any  $x$  and  $y$ :  $x^y \bmod n = x^{(y \bmod z)} \bmod n$ 
  - where  $n = pq$  and  $z = (p-1)(q-1)$
- thus,
$$\begin{aligned}c^d \bmod n &= (m^e \bmod n)^d \bmod n \\&= m^{ed} \bmod n \\&= m^{(ed \bmod z)} \bmod n \\&= m^1 \bmod n \\&= m\end{aligned}$$



# RSA: another important property

The following property will be *very* useful later:

$$\underbrace{K_B^-(K_B^+(m))}_{\text{use public key first, followed by private key}} = m = \underbrace{K_B^+(K_B^-(m))}_{\text{use private key first, followed by public key}}$$

use public key first,  
followed by  
private key

use private key  
first, followed by  
public key

*result is the same!*

**Why**  $K_B^-(K_B^+(m)) = m = K_B^+(K_B^-(m))$  ?

---

follows directly from modular arithmetic:

$$\begin{aligned}(m^e \bmod n)^d \bmod n &= m^{ed} \bmod n \\ &= m^{de} \bmod n \\ &= (m^d \bmod n)^e \bmod n\end{aligned}$$

# Why is RSA secure?

- suppose you know Bob's public key  $(n,e)$ . How hard is it to determine  $d$ ?
- essentially need to find factors of  $n$  without knowing the two factors  $p$  and  $q$ 
  - If  $p$  and  $q$  is known, can get  $z=(p-1)(q-1)$
  - given  $e$ , can find  $d=e^{-1} \bmod z$
  - fact: factoring a big number is hard

# RSA in practice: session keys

- exponentiation in RSA is computationally intensive
- DES is at least 100 times faster than RSA
- use public key crypto to establish secure connection, then establish second key – symmetric session key – for encrypting data

## *session key, $K_S$*

- Bob and Alice use RSA to exchange a symmetric key  $K_S$ 
  - Sender computes  $c = (K_S)^e \bmod n$ , receiver obtains  $K_S$
- once both have  $K_S$ , they use symmetric key cryptography

# Chapter 8 roadmap

- 8.1 What is network security?
- 8.2 Principles of cryptography
- 8.3 Message integrity, *authentication*
- 8.4 Securing e-mail
- 8.5 Securing TCP connections: SSL
- 8.6 Network layer security: IPsec
- 8.7 Securing wireless LANs
- 8.8 Operational security: firewalls and IDS

# Authentication

*Goal:* Bob wants Alice to “prove” her identity to him

*Protocol ap1.0:* Alice says “I am Alice”



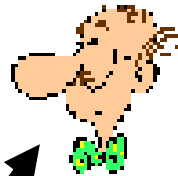
Failure scenario??



# Authentication

*Goal:* Bob wants Alice to “prove” her identity to him

*Protocol ap1.0:* Alice says “I am Alice”

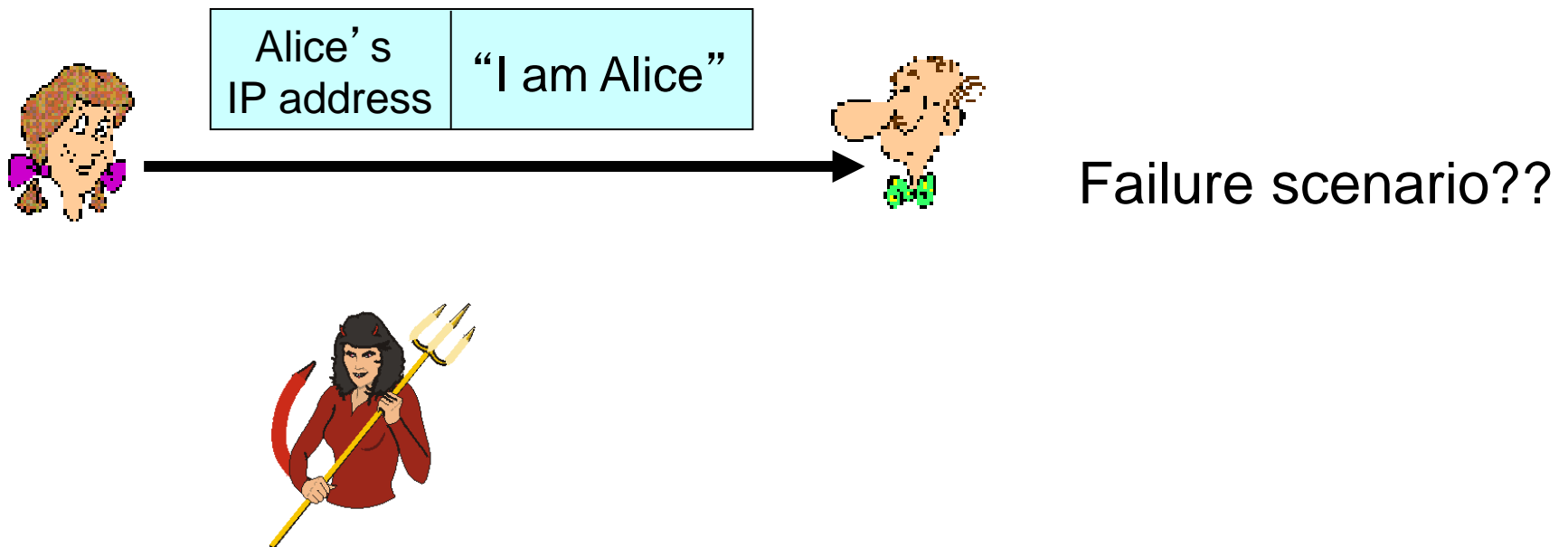


“I am Alice”

in a network,  
Bob can not “see” Alice,  
so Trudy simply declares  
herself to be Alice

# Authentication: another try

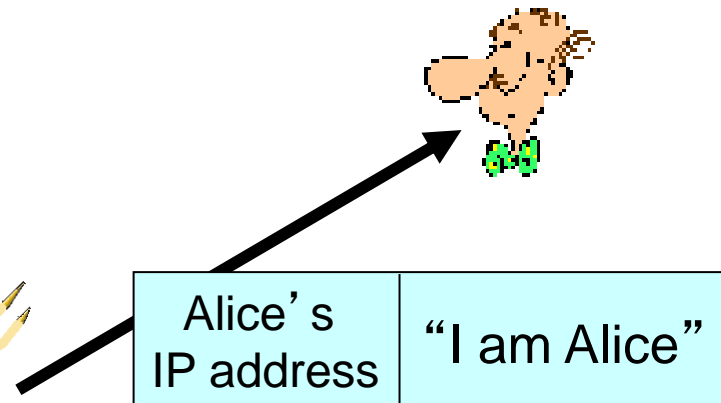
*Protocol ap2.0:* Alice says “I am Alice” in an IP packet containing her source IP address





# Authentication: another try

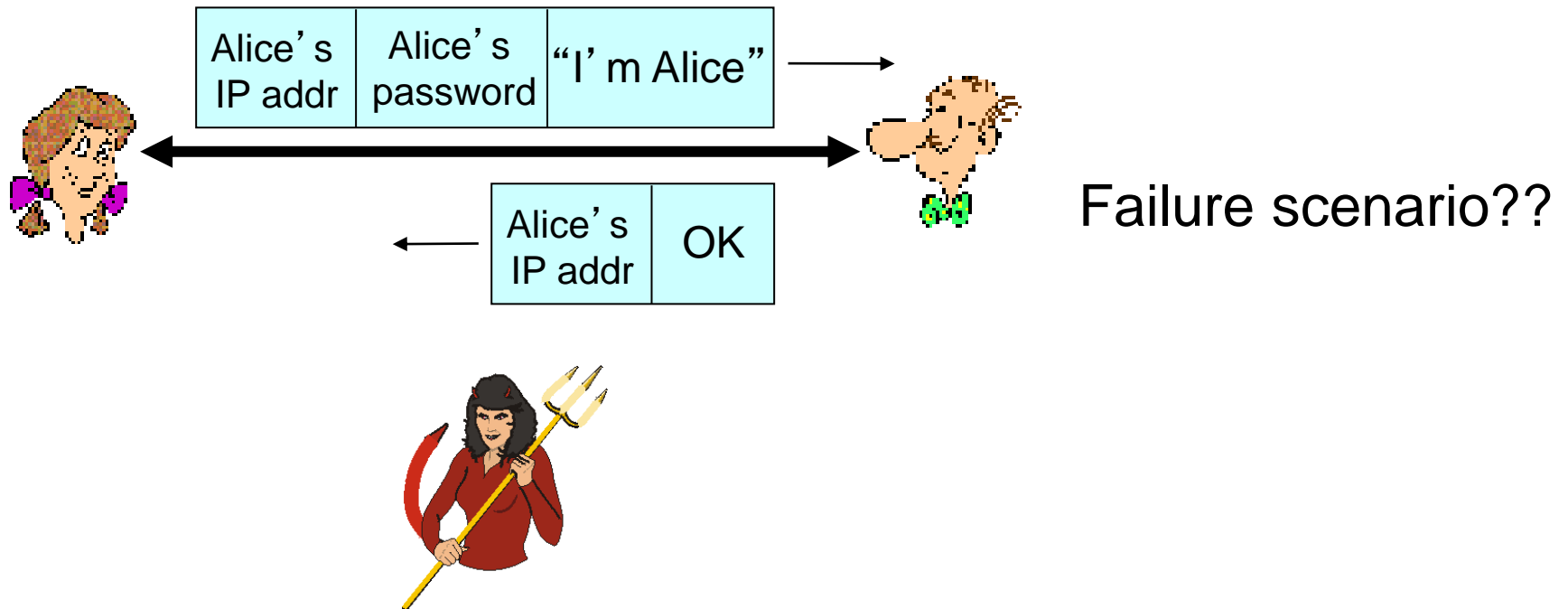
*Protocol ap2.0:* Alice says “I am Alice” in an IP packet containing her source IP address



Trudy can create  
a packet  
“spoofing”  
Alice’s address

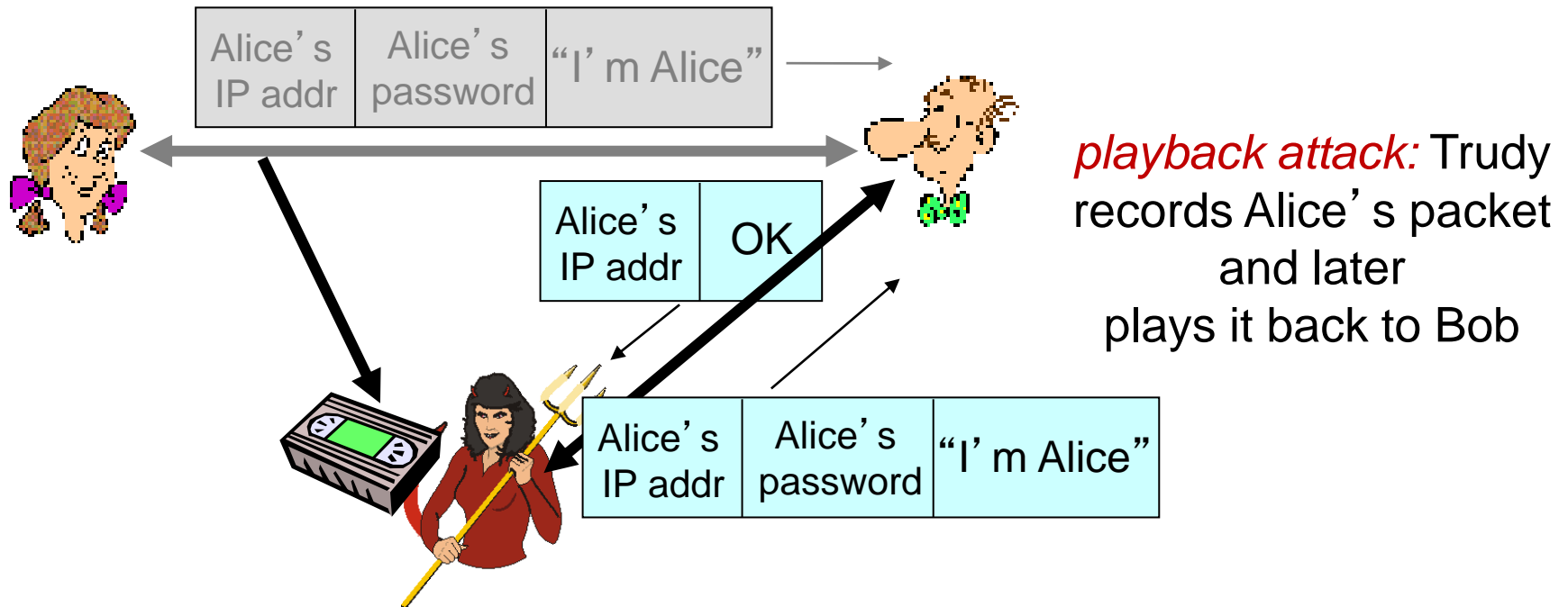
# Authentication: another try

*Protocol ap3.0:* Alice says “I am Alice” and sends her secret password to “prove” it.



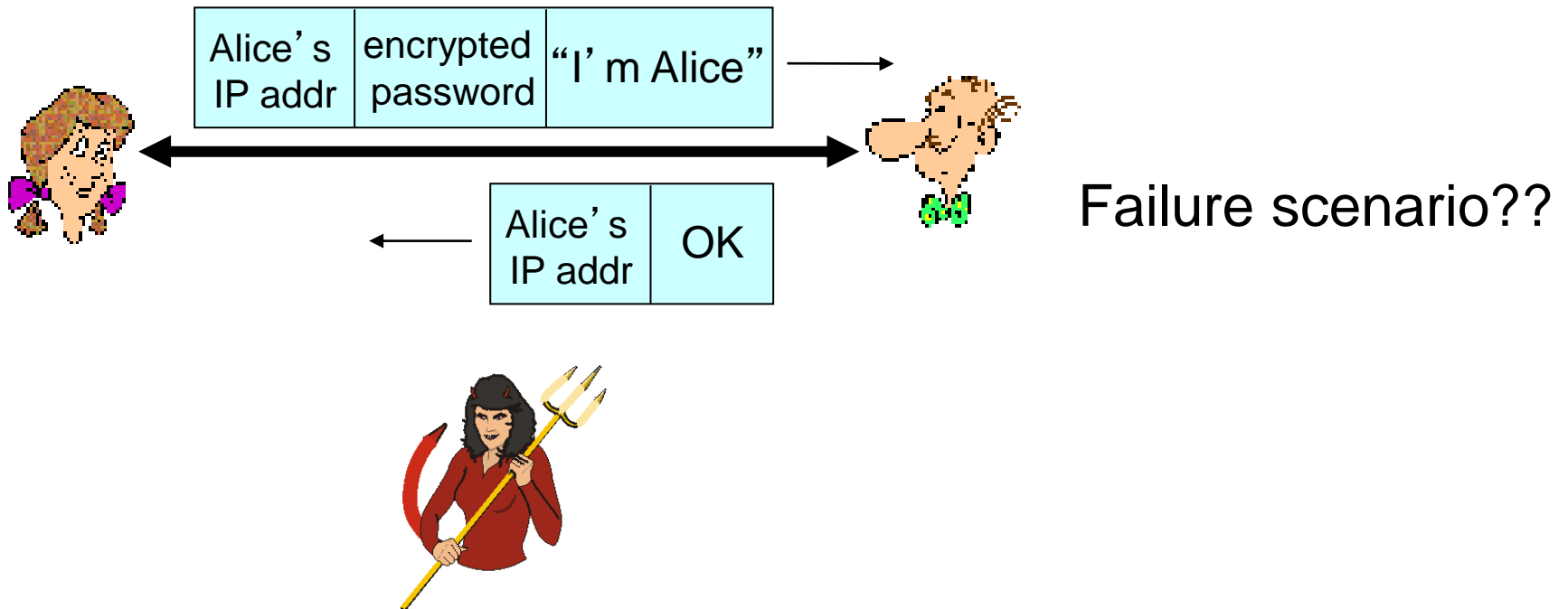
# Authentication: another try

*Protocol ap3.0:* Alice says “I am Alice” and sends her secret password to “prove” it.



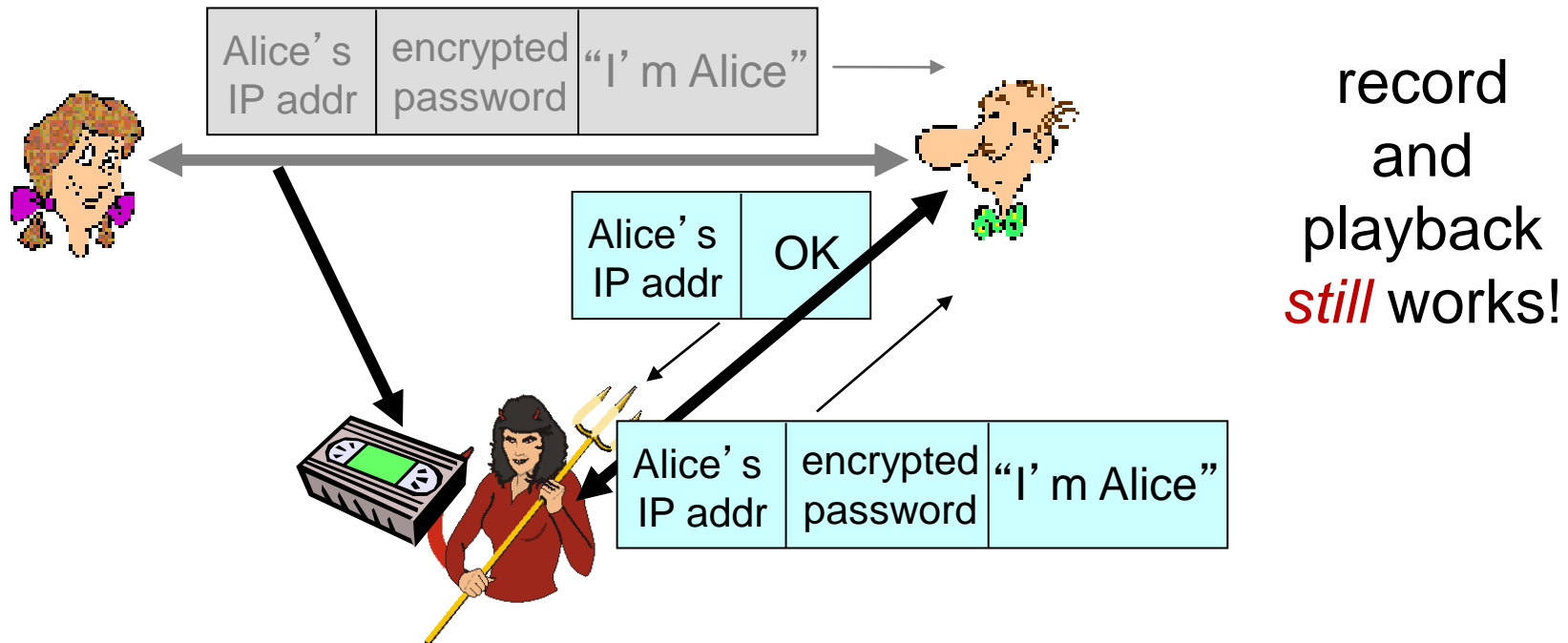
# Authentication: yet another try

*Protocol ap3.1:* Alice says “I am Alice” and sends her *encrypted* secret password to “prove” it.



# Authentication: yet another try

*Protocol ap3.1:* Alice says “I am Alice” and sends her *encrypted* secret password to “prove” it.

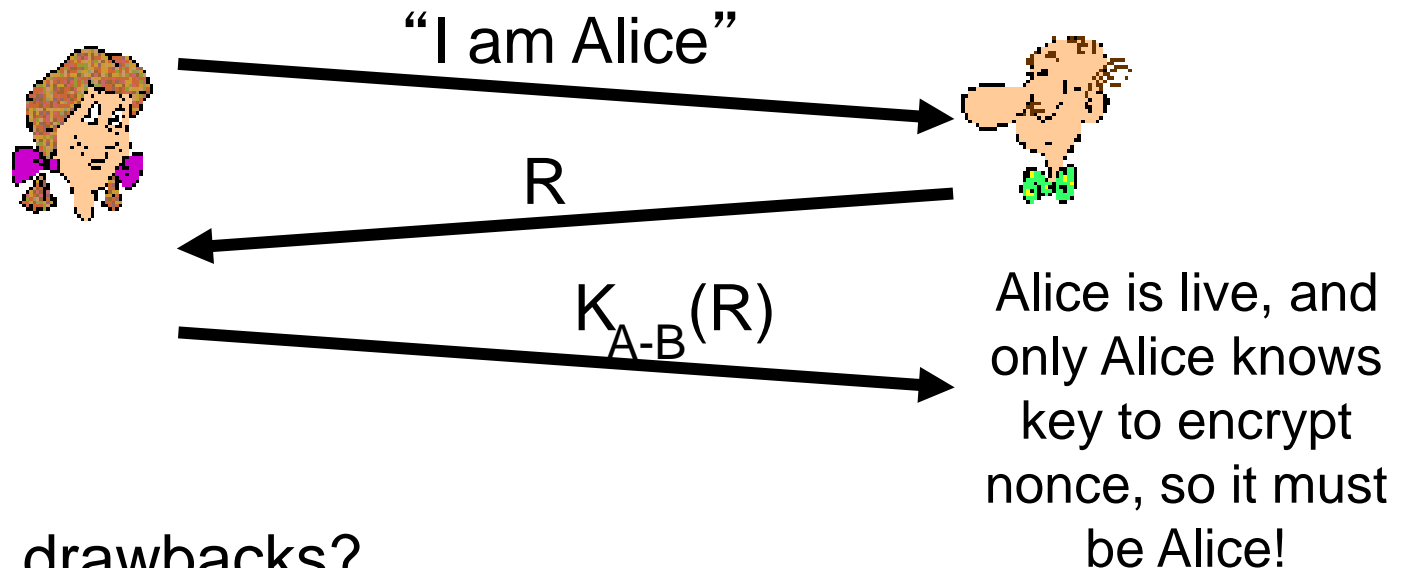


# Authentication: yet another try

**Goal:** avoid playback attack

**nonce:** number (R) used only *once-in-a-lifetime*

**ap4.0:** to prove Alice “live”, Bob sends Alice **nonce**, R. Alice must return R, encrypted with shared secret key



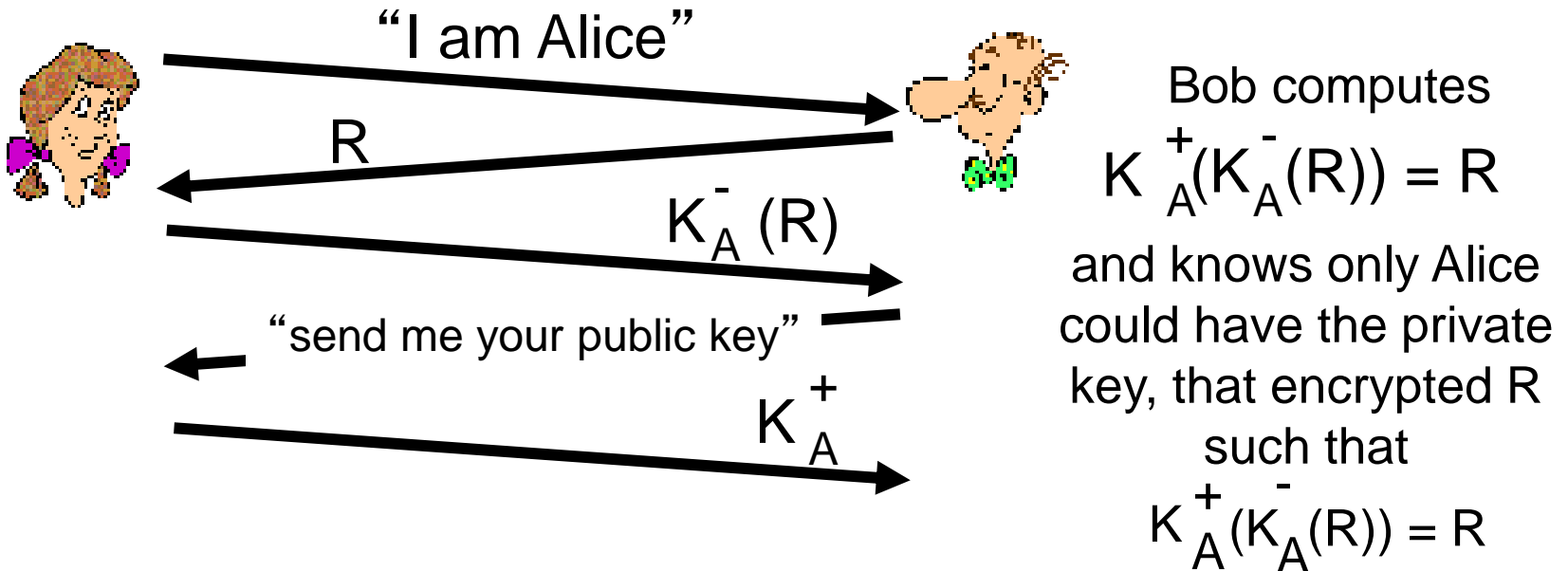
Failures, drawbacks?

# Authentication: ap5.0

ap4.0 requires shared symmetric key

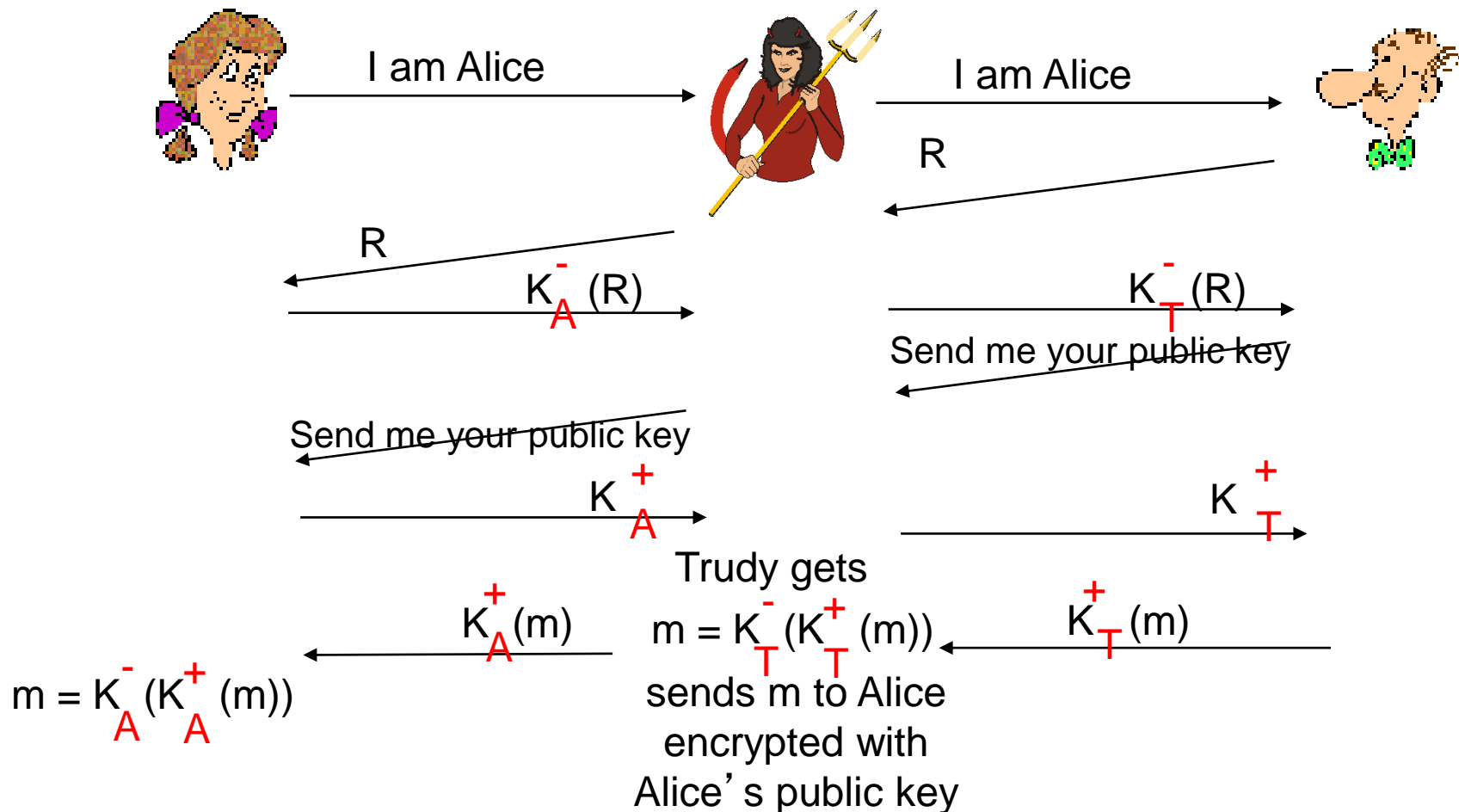
- can we authenticate using public key techniques?

*ap5.0*: use nonce, public key cryptography



# ap5.0: security hole

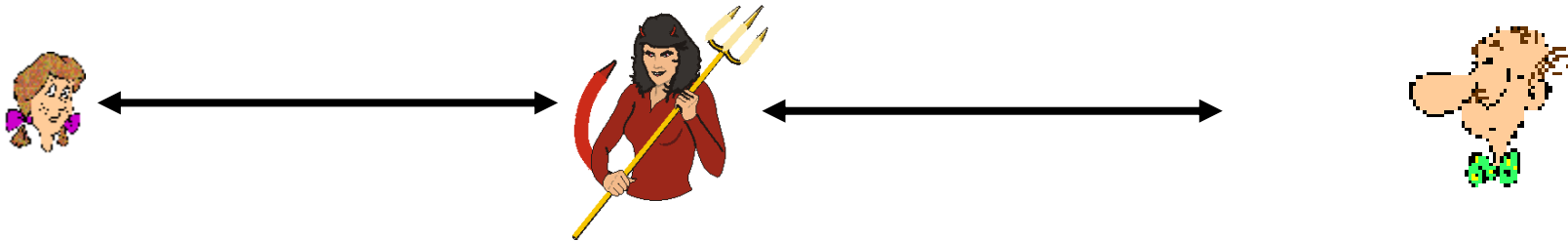
*man (or woman) in the middle attack:* Trudy poses as Alice (to Bob) and as Bob (to Alice)





# ap5.0: security hole

*man (or woman) in the middle attack:* Trudy poses as Alice (to Bob) and as Bob (to Alice)



difficult to detect:

- Bob receives everything that Alice sends, and vice versa. (e.g., so Bob, Alice can meet one week later and recall conversation!)
- problem is that Trudy receives all messages as well!

未完待续

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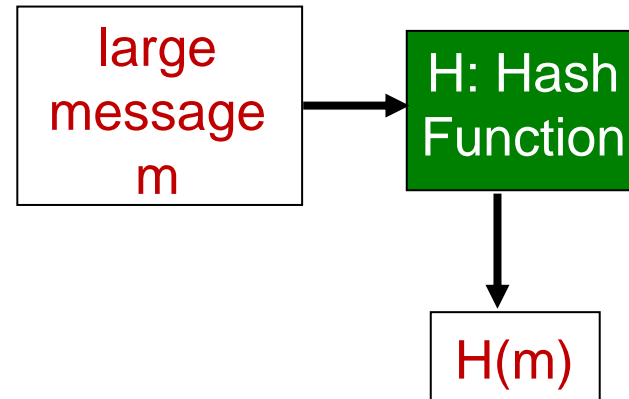
8.8 Operational security: firewalls and IDS

# Message digests

computationally  
expensive to public-key-  
encrypt long messages

**goal:** fixed-length, easy-  
to-compute digital  
“fingerprint”

- apply hash function  $H$  to  $m$ , get fixed size **message digest** (报文摘要),  $H(m)$ .



**Hash function properties:**

- many-to-1
- produces fixed-size msg digest (fingerprint)
- given message digest  $x$ , computationally infeasible to find  $m$  such that  $x = H(m)$

# Internet checksum: poor crypto hash function

Internet checksum has some properties of hash function:

- produces fixed length digest (16-bit sum) of message
- is many-to-one

But given message with given hash value, it is easy to find another message with same hash value:

<u>message</u>	<u>ASCII format</u>		<u>message</u>	<u>ASCII format</u>
I O U 1	49 4F 55 31		I O U <u>9</u>	49 4F 55 <u>39</u>
0 0 . 9	30 30 2E 39		0 0 . <u>1</u>	30 30 2E <u>31</u>
9 B O B	39 42 D2 42		9 B O B	39 42 D2 42
	<u>B2 C1 D2 AC</u>	different messages but identical checksums!		<u>B2 C1 D2 AC</u>

# Hash function algorithms

- **MD5 hash function widely used (RFC 1321)**
  - computes 128-bit message digest in 4-step process.
  - arbitrary 128-bit string  $x$ , appears difficult to construct msg  $m$  whose MD5 hash is equal to  $x$
- **SHA-1 is also used**
  - US standard [NIST, FIPS PUB 180-1]
  - 160-bit message digest

# Message authentication code

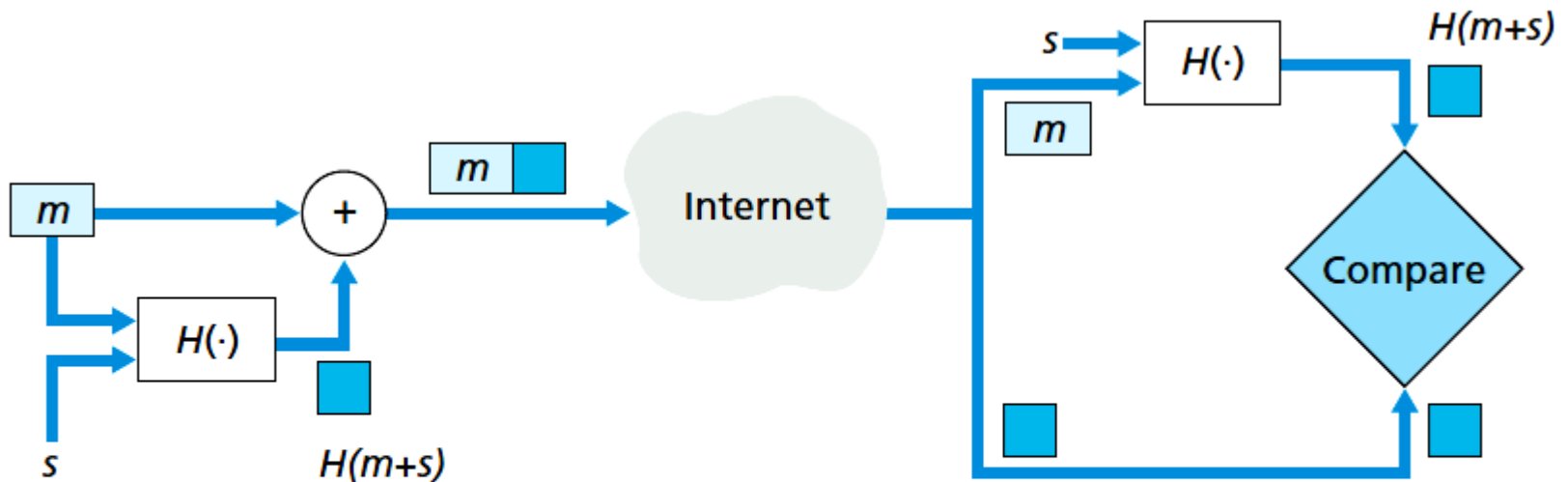
- A flawed approach for protecting message integrity
  1. Alice creates message  $m$  and calculates the hash  $H(m)$  (for example with SHA-1).
  2. Alice then appends  $H(m)$  to the message  $m$ , creating an extended message  $(m, H(m))$ , and sends the extended message to Bob.
  3. Bob receives an extended message  $(m, h)$  and calculates  $H(m)$ . If  $H(m) = h$ ,
  4. Bob concludes that everything is fine.

Trudy can create a bogus message  $m'$  in which she says she is Alice, calculate  $H(m')$ , and send Bob  $(m', H(m'))$ .

# Message authentication code

- Alice and Bob will need a shared secret  $s$ , called the *authentication key*.
  1. Alice concatenates  $s$  with  $m$  to create  $m + s$ , and calculates the hash  $H(m + s)$  (for example with SHA-1).
    - $H(m + s)$  is called the **message authentication code (MAC)**.
  2. Alice then appends the MAC to the message  $m$ , creating an extended message  $(m, H(m + s))$ , and sends to Bob.
  3. Bob receives an extended message  $(m, h)$  and knowing  $s$ , calculates the MAC  $H(m + s)$ .
    - If  $H(m + s) = h$ , Bob concludes that everything is fine.

# Message authentication code



Key:

$m$  = Message  
 $s$  = Shared secret

- Question: how to distributed authentication key?
  - Manual config
  - Public key



# Digital signatures

cryptographic technique analogous to hand-written signatures:

- sender (Bob) digitally signs document, establishing he is document owner/creator.
- *verifiable, nonforgeable*: recipient (Alice) can prove to someone that Bob, and no one else (including Alice), must have signed document


# Digital signatures

simple digital signature for message  $m$ :

- Bob signs  $m$  by encrypting with his private key  $K_B^-$ , creating “signed” message,  $K_B^-(m)$

Bob's message,  $m$

Dear Alice  
Oh, how I have missed  
you. I think of you all the  
time! ... (blah blah blah)  
Bob

  $K_B^-$  Bob's private  
key

Public key  
encryption  
algorithm

$m, K_B^-(m)$

Bob's message,  
 $m$ , signed  
(encrypted) with  
his private key

# Digital signatures

- suppose Alice receives msg  $m$ , with signature:  $m, K_B^-(m)$
- Alice verifies  $m$  signed by Bob by applying Bob's public key  $K_B^+$  to  $K_B^-(m)$  then checks  $K_B^+(K_B^-(m)) = m$ .
- If  $K_B^+(K_B^-(m)) = m$ , whoever signed  $m$  must have used Bob's private key.

Alice thus verifies that:

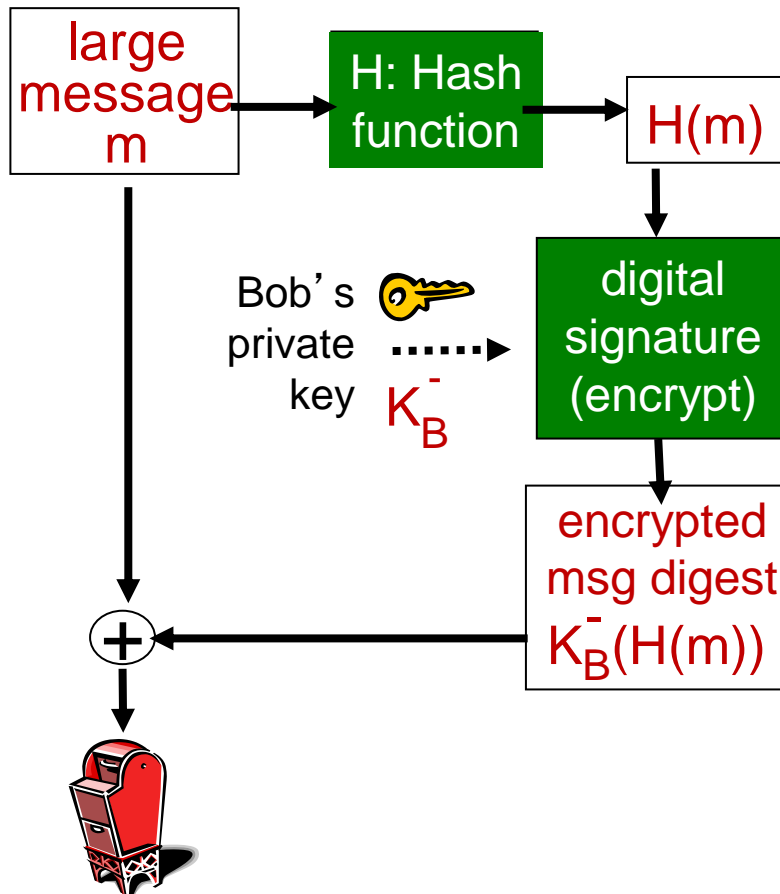
- Bob signed  $m$
- no one else signed  $m$
- Bob signed  $m$  and not  $m'$

non-deniable:

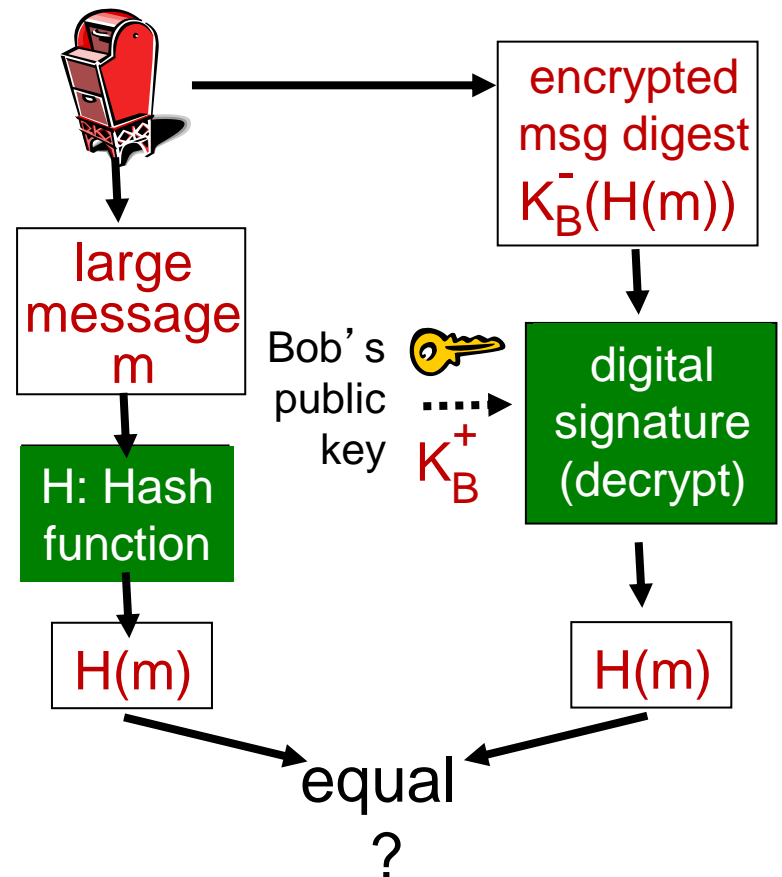
- ✓ Alice can take  $m$ , and signature  $K_B^-(m)$  to court and prove that Bob signed  $m$

# Digital signature = signed message digest

Bob sends digitally signed message:

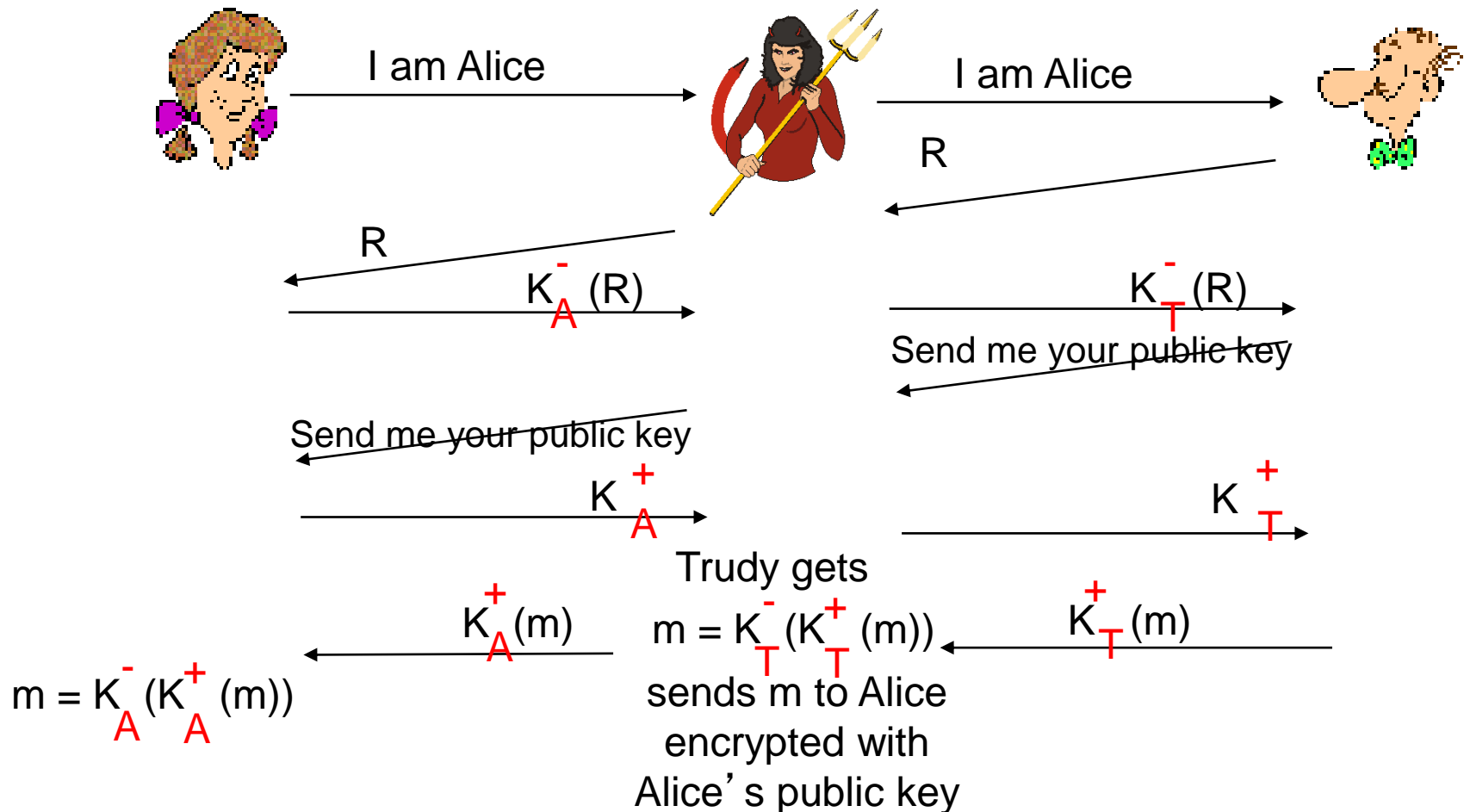


Alice verifies signature, integrity of digitally signed message:



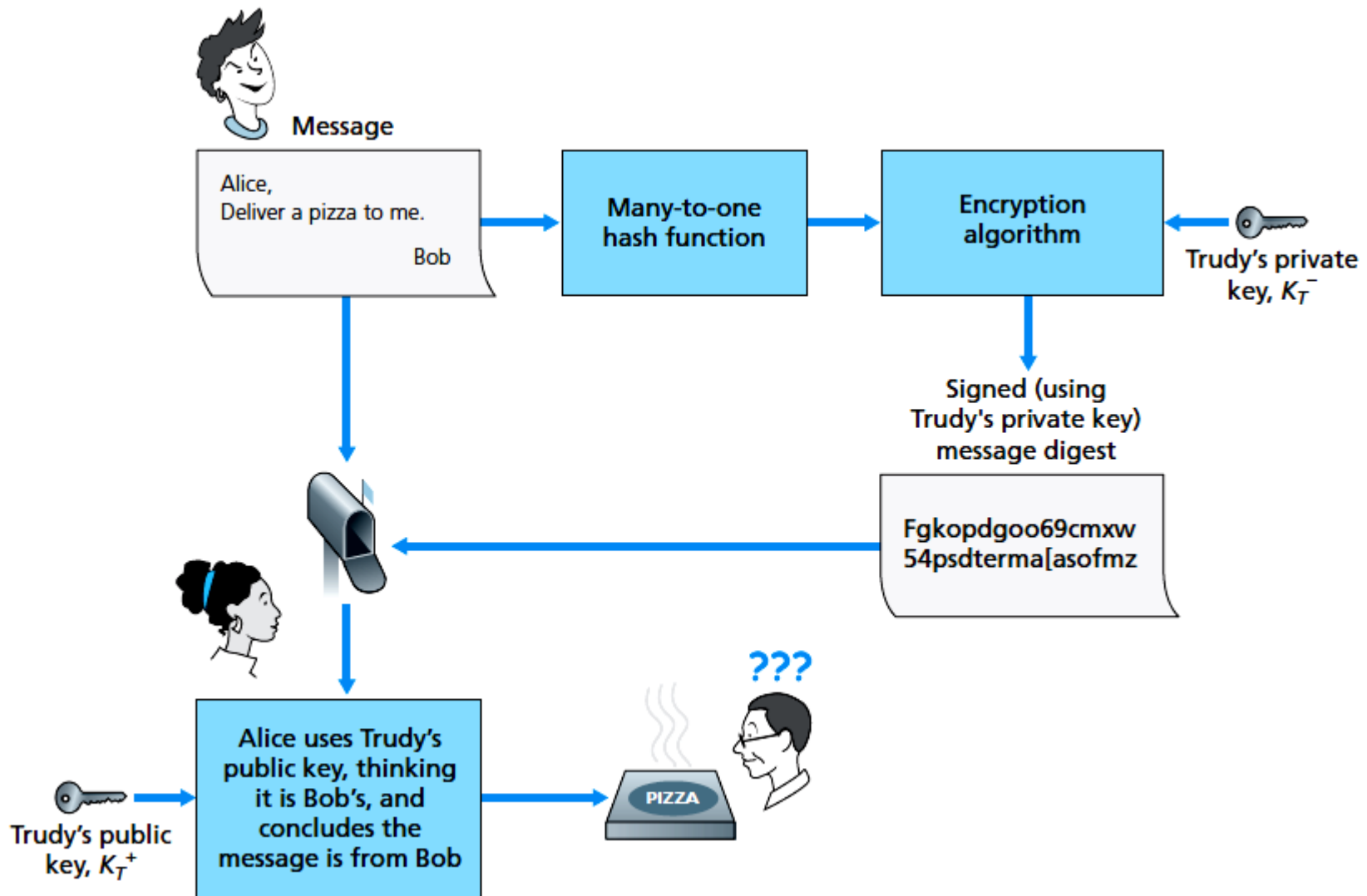
# Recall ap5.0: security hole

*man (or woman) in the middle attack:* Trudy poses as Alice (to Bob) and as Bob (to Alice)



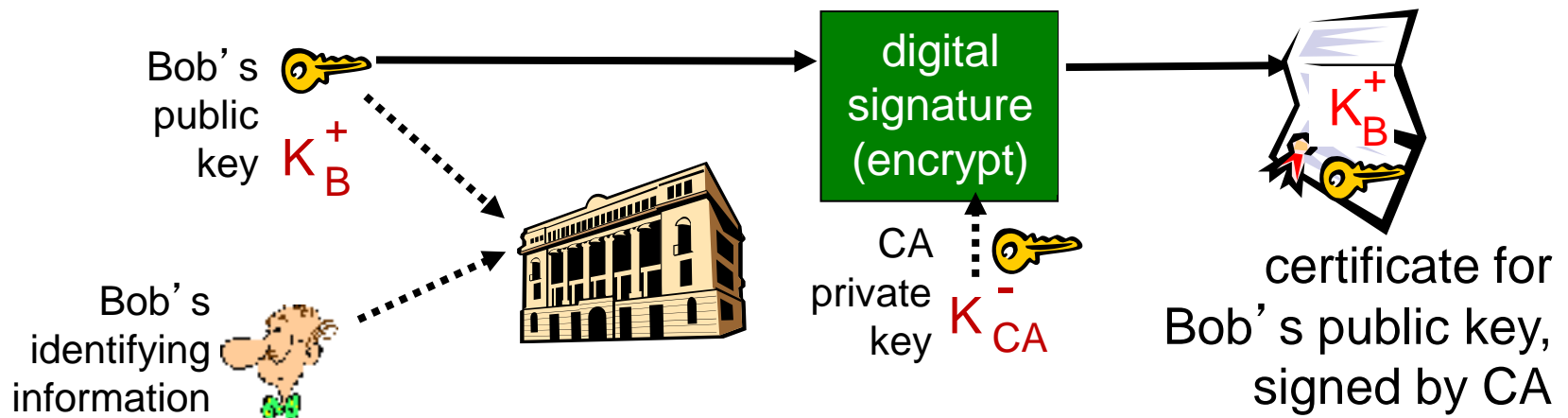
# Public-key certification

- motivation: Trudy plays pizza prank on Bob
  - Trudy creates e-mail order:  
*Dear Pizza Store, Please deliver to me four pepperoni pizzas. Thank you, Bob*
  - Trudy signs order with her private key
  - Trudy sends order to Pizza Store
  - Trudy sends to Pizza Store her public key, but says it's Bob's public key
  - Pizza Store verifies signature; then delivers four pepperoni pizzas to Bob
  - Bob doesn't even like pepperoni



# Certification authorities

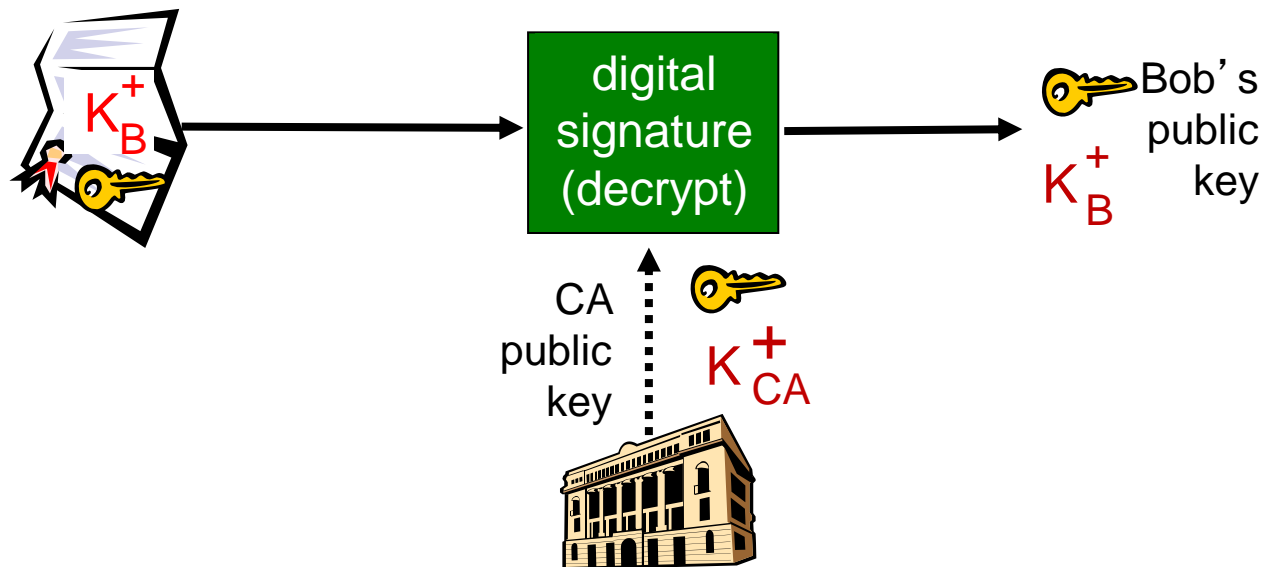
- **certification authority (CA):** binds public key to particular entity, E.
- E (person, router) registers its public key with CA.
  - E provides “proof of identity” to CA.
  - CA creates certificate binding E to its public key.
  - certificate containing E’s public key digitally signed by CA – CA says “this is E’s public key”





# Certification authorities

- when Alice wants Bob's public key:
  - gets Bob's certificate (Bob or elsewhere).
  - apply CA's public key to Bob's certificate, get Bob's public key



# RSA in Linux

---

## ❖ *Create RSA key pair*

```
root@server1:~# ssh-keygen -t rsa
Generating public/private rsa key pair.
Enter file in which to save the key (/root/.ssh/id_rsa):
Enter passphrase (empty for no passphrase):
Enter same passphrase again:
Your identification has been saved in /root/.ssh/id_rsa.
```

## ❖ *Upload the public key on a server*

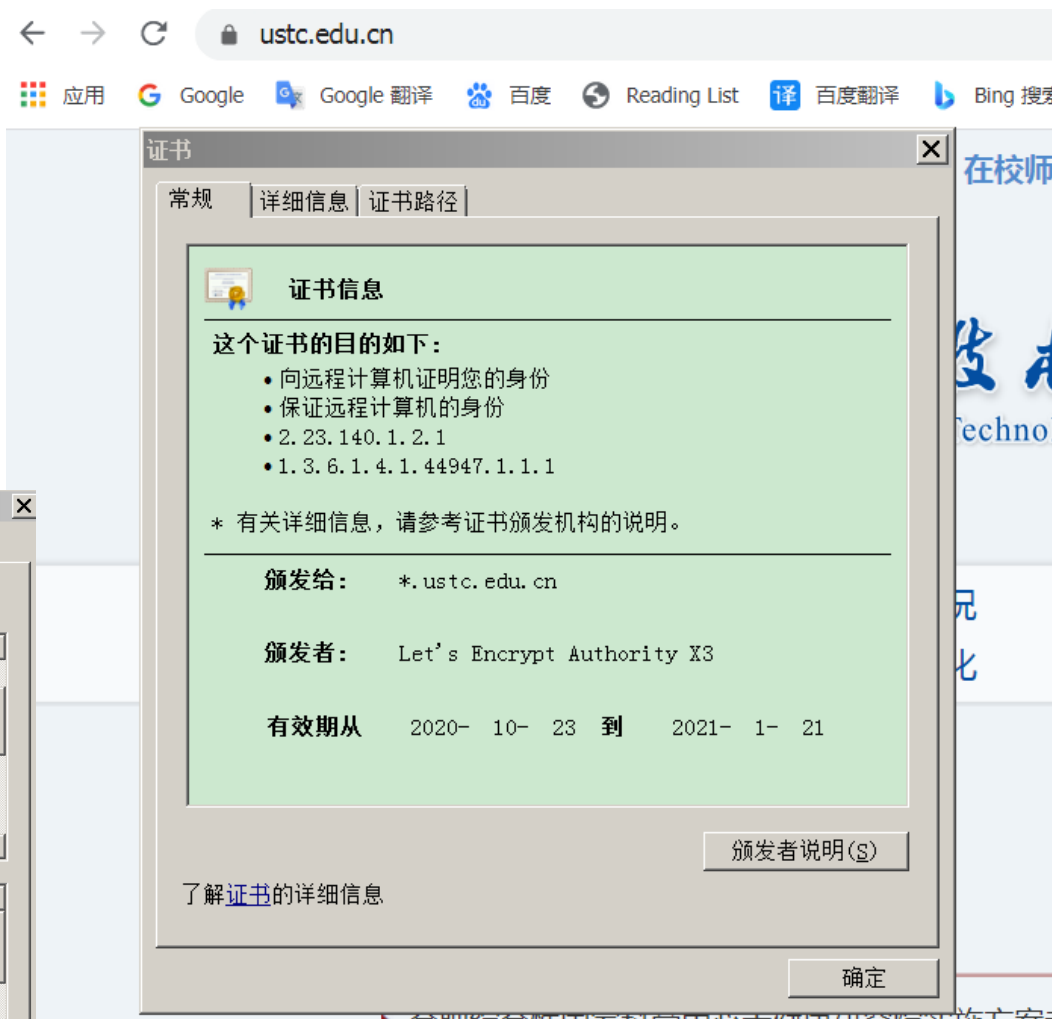
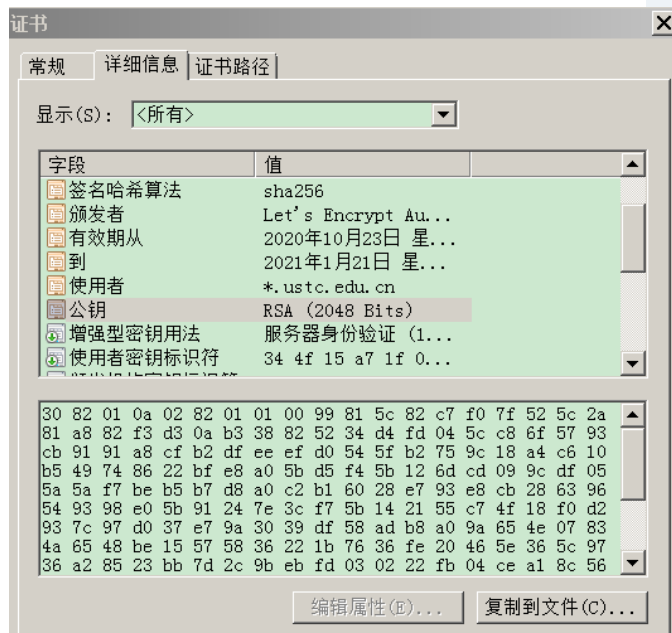
```
ssh-copy-id user@192.168.0.100
```

*or*

```
cat ~/.ssh/id_rsa.pub | ssh user@192.168.0.100 "cat >>
~/.ssh/authorized_keys"
```

# CA in practice

Let's Encrypt作为CA颁发给ustc.edu.cn的证书



# Summary

---

- ❖ Objective: message encryption
  - ❖ Solution: block cipher,  
 $m = K_S(K_S(m))$
  - ❖ cipher block chaining  $c(i) = K_S(m(i) \oplus c(i-1))$   $m(i) = K_S(c(i)) \oplus c(i-1)$
  - ❖ Require: shared secret key (symmetric key)
- ❖ Requirement: negotiate shared secret key
  - ❖ Solution: public key cryptography  
 $K_B^-(K_B^+(m)) = m$
  - ❖ Require: how to know someone's public key?

# Summary

---

- ❖ Objective: authentication (Alice proves she is Alice to Bob)

- ❖ Solution: nonce, CA

$$K_A^+(K_A^-(R)) = R$$

$$K_{CA}^+(K_{CA}^-(K_A^+))$$

- ❖ Objective: Integrity

- ❖ Solution: MAC using hash function and authentication key

$$H(m + s)$$

- ❖ Digital signature:

$$K_A^+(K_A^-(H(m)))$$

# Chapter 8 roadmap

8.1 What is network security?

8.2 Principles of cryptography

8.3 Message integrity, authentication

*8.4 Securing e-mail*

8.5 Securing TCP connections: SSL

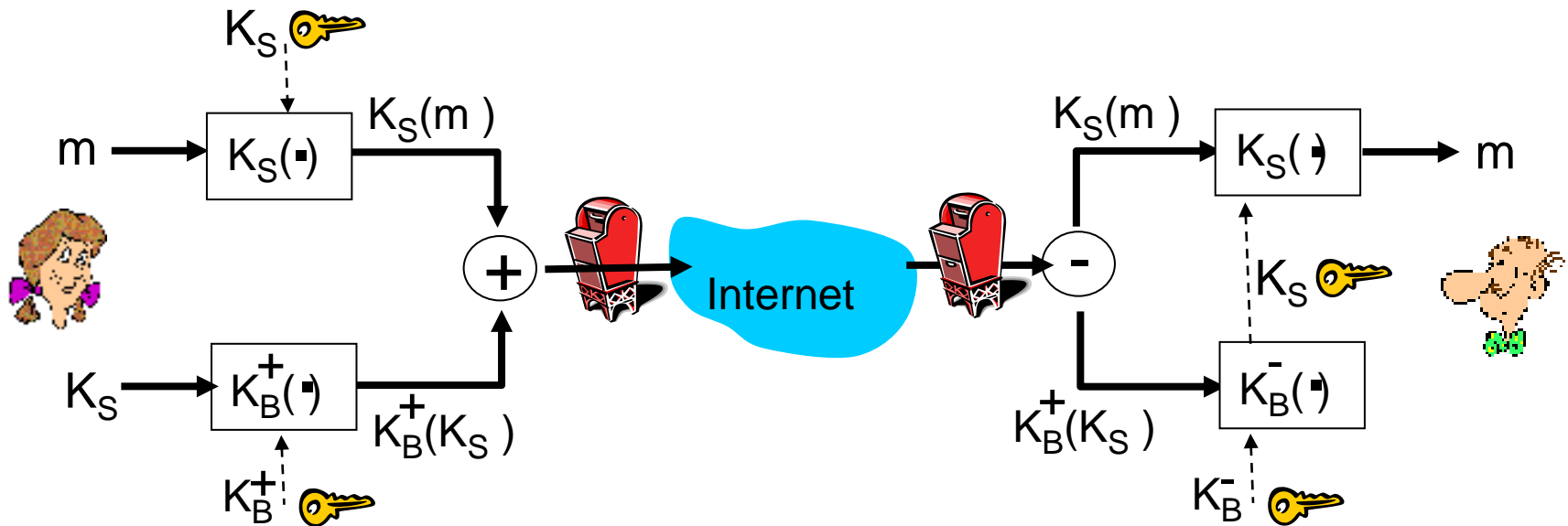
8.6 Network layer security: IPsec

8.7 Securing wireless LANs

8.8 Operational security: firewalls and IDS

# Secure e-mail

Alice wants to send confidential e-mail,  $m$ , to Bob.

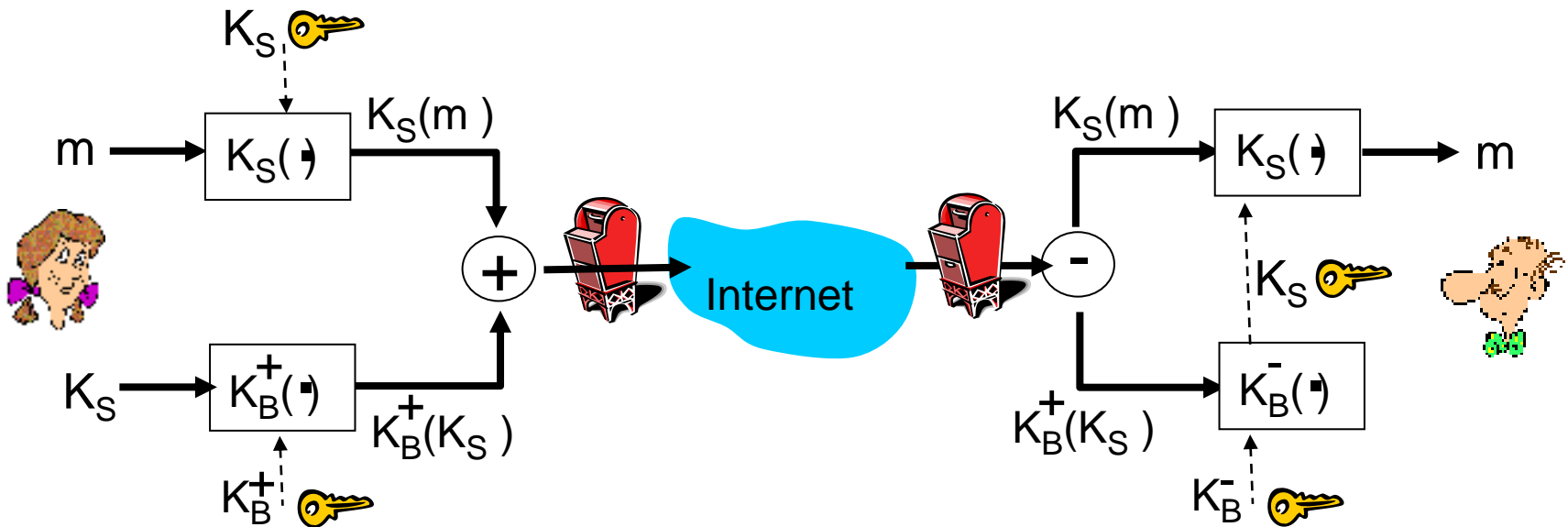


*Alice:*

- generates random *symmetric* private key,  $K_S$
- encrypts message with  $K_S$  (for efficiency)
- also encrypts  $K_S$  with Bob's public key
- sends both  $K_S(m)$  and  $K_B(K_S)$  to Bob

# Secure e-mail

Alice wants to send confidential e-mail,  $m$ , to Bob.



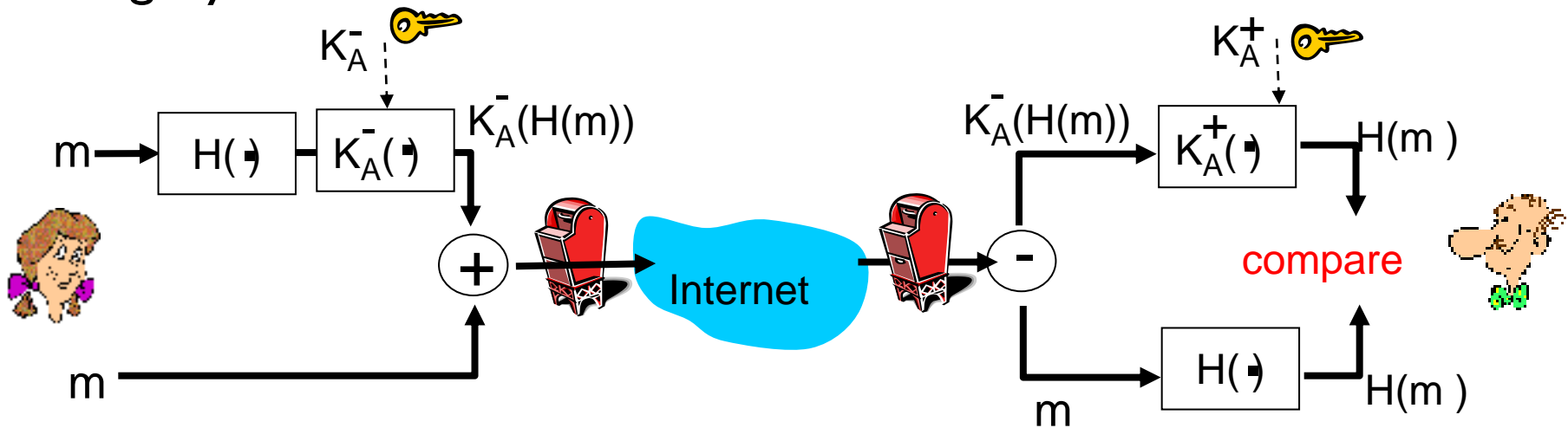
**Bob:**

- uses his private key to decrypt and recover  $K_S$
- uses  $K_S$  to decrypt  $K_S(m)$  to recover  $m$



# Secure e-mail (continued)

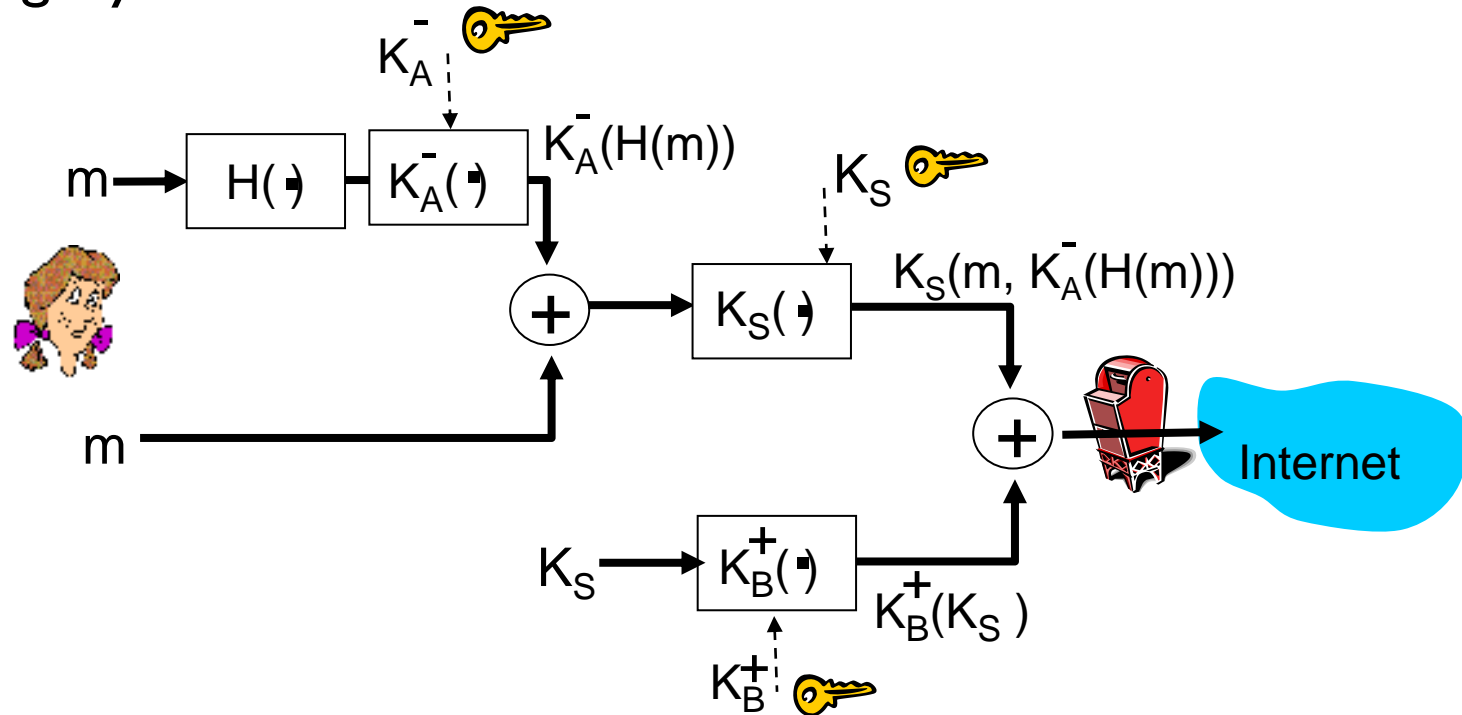
Alice wants to provide sender authentication and message integrity



- Alice digitally signs message
- sends both message (in the clear) and digital signature

# Secure e-mail (continued)

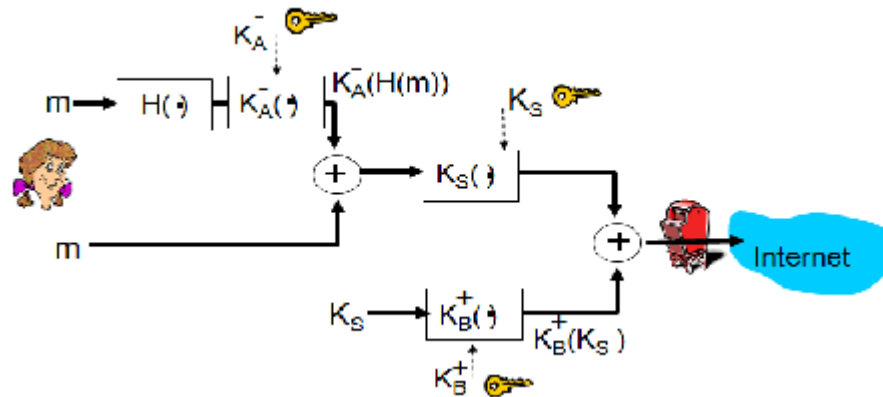
Alice wants to provide secrecy, sender authentication, message integrity.



*Alice uses three keys:* her private key, Bob's public key, newly created symmetric key

# Secure e-mail (continued)

- ❖ PGP (Pretty Good Privacy)
  - ❖ De facto standard
  - ❖ Essentially same as the previous slide



- ❖ Software generate public key pair, choose MD5 or SHA for hashing

# Chapter 8 roadmap

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*8.5 Securing TCP connections: SSL*

8.6 Network layer security: IPsec

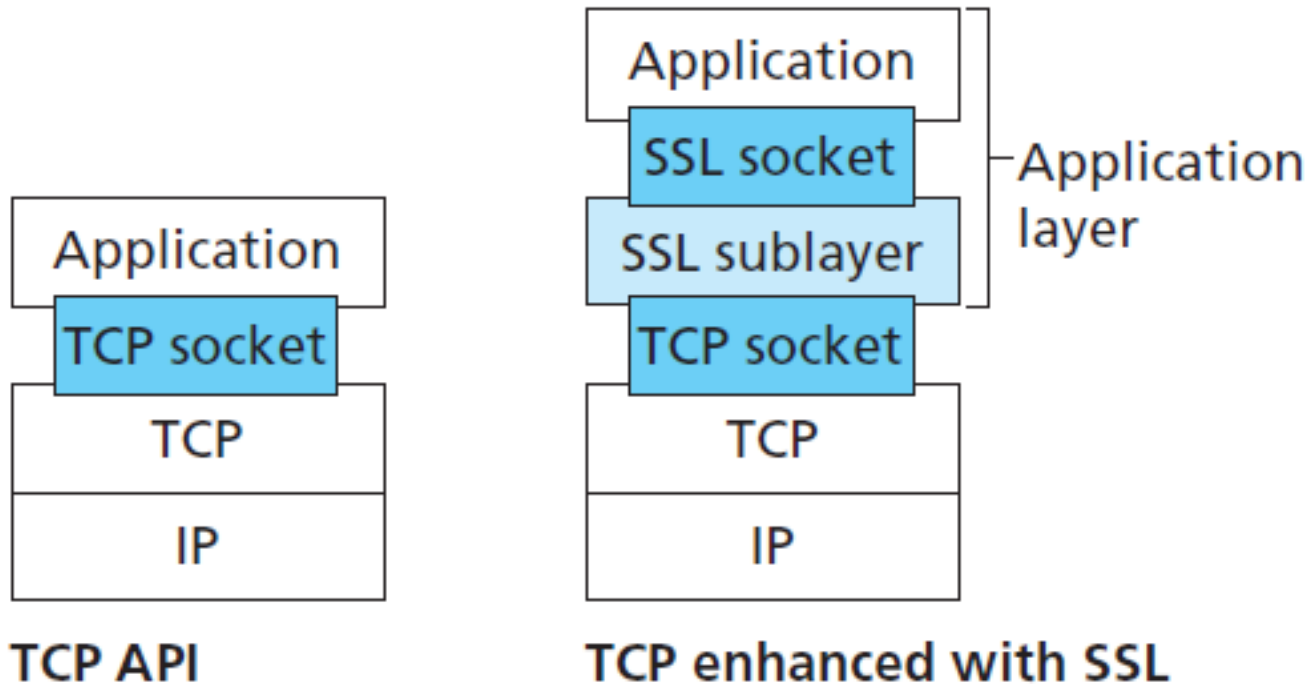
8.7 Securing wireless LANs

8.8 Operational security: firewalls and IDS

# SSL: Secure Sockets Layer

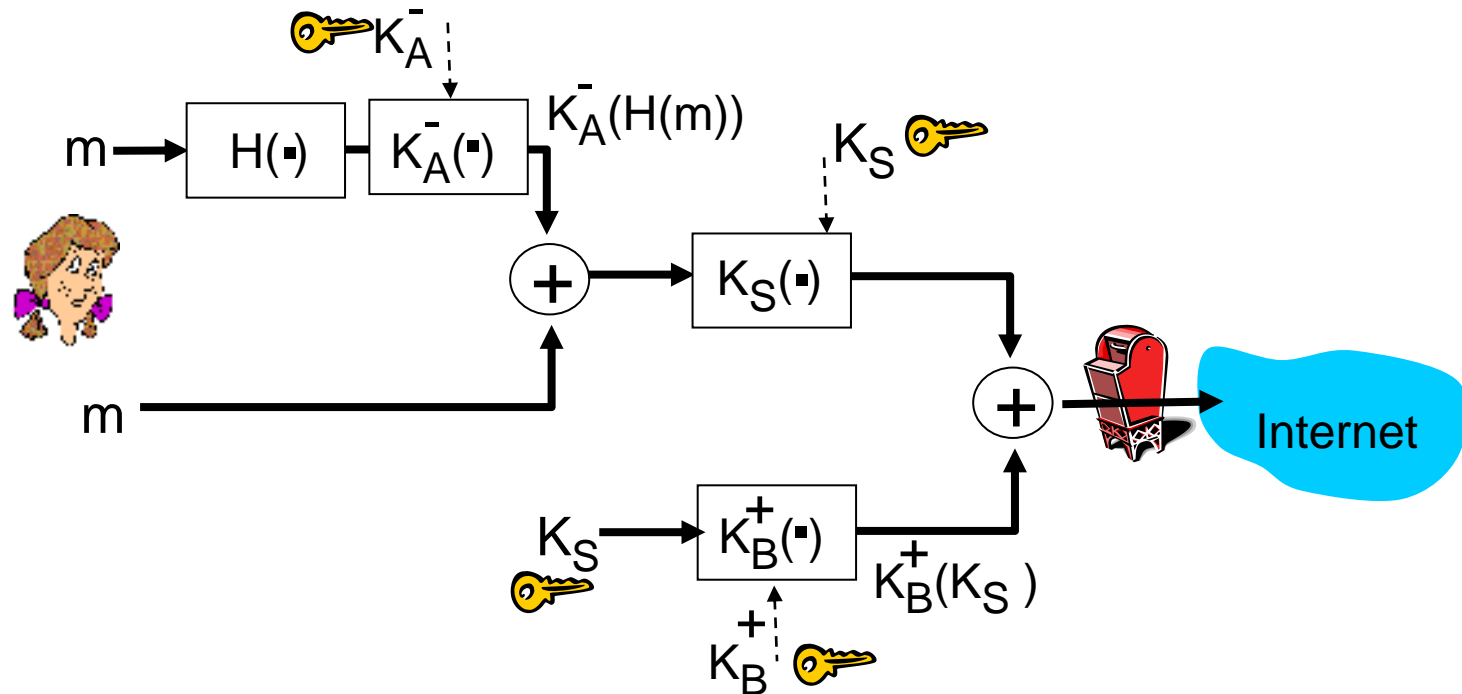
- widely deployed security protocol
  - supported by almost all browsers, web servers
  - https
  - billions \$/year over SSL
- mechanisms: [Woo 1994], implementation: Netscape
- variation -TLS: transport layer security, RFC 2246
- provides
  - *confidentiality*
  - *integrity*
  - *authentication*
- original goals:
  - Web e-commerce transactions
  - encryption (especially credit-card numbers)
  - Web-server authentication
  - optional client authentication
  - minimum hassle in doing business with new merchant
- available to all TCP applications
  - secure socket interface

# SSL and TCP/IP



- SSL provides application programming interface (API) to applications
- C and Java SSL libraries/classes readily available

# Could do something like PGP:



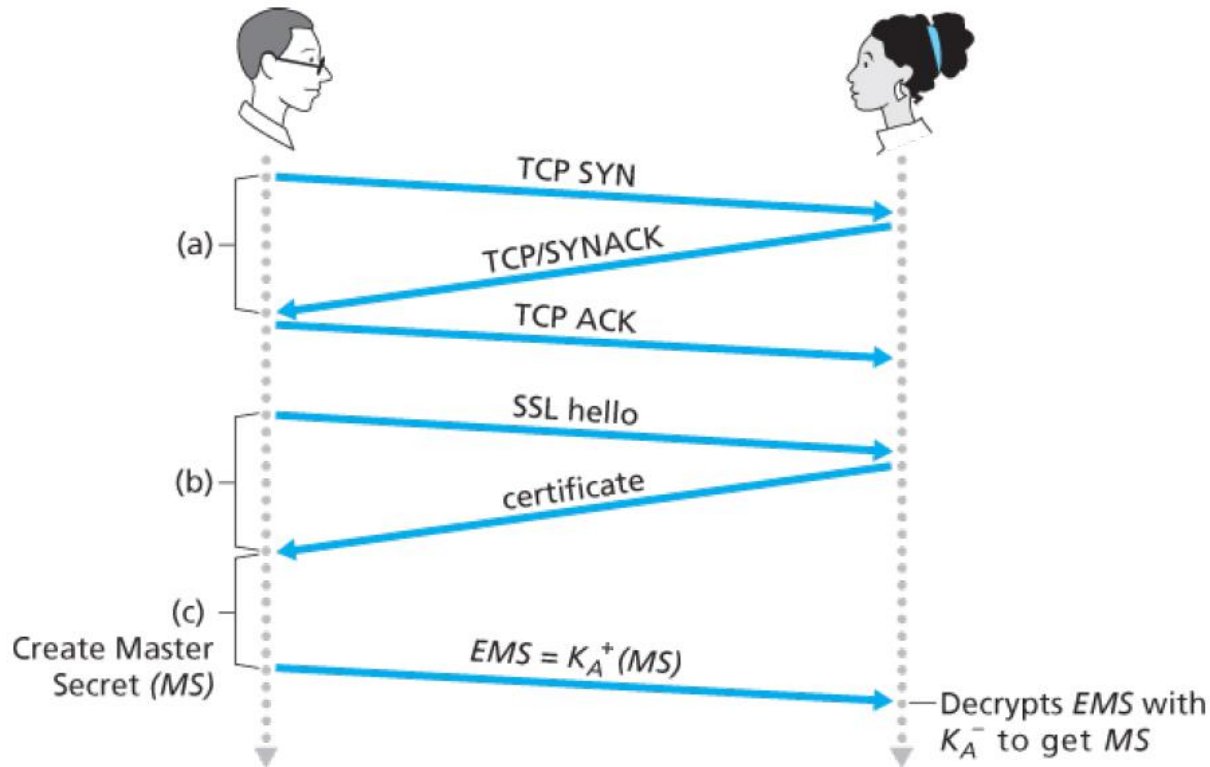
- but want to send byte streams & interactive data
- want set of secret keys for entire connection
- want certificate exchange as part of protocol: handshake phase

# Toy SSL: a simple secure channel

- *handshake*: Alice and Bob use their certificates, private keys to authenticate each other and exchange shared secret
- *key derivation*: Alice and Bob use shared secret to derive set of keys
- *data transfer*: data to be transferred is broken up into series of records
- *connection closure*: special messages to securely close connection



# Toy: a simple handshake



**MS:** master secret

**EMS:** encrypted master secret

# Toy: key derivation

- considered bad to use same key for more than one cryptographic operation
  - use different keys for message authentication code (MAC) and encryption
- four keys:
  - $E_B$  = encryption key for data sent from Bob to Alice
  - $M_B$  = MAC key for data sent from Bob to Alice
  - $E_A$  = encryption key for data sent from Alice to Bob
  - $M_A$  = MAC key for data sent from Alice to Bob
- keys derived from the MS, using key derivation function (KDF)
  - Example: simply slicing the MS into four keys

# Toy: data records

- why not encrypt data in constant stream as we write it to TCP?
  - where would we put the MAC? If at end, no message integrity until all data processed.
  - e.g., with instant messaging, how can we do integrity check over all bytes sent before displaying?
- instead, break stream in series of records
  - each record carries a MAC
  - receiver can act on each record as it arrives
- issue: in record, receiver needs to distinguish MAC from data
  - want to use variable-length records



# Toy: sequence numbers

- *problem:* attacker can capture and replay record or re-order records
- *solution:* put sequence number into MAC:
  - $\text{MAC} = \text{Hash}(M_x, \text{sequence}||\text{data})$
  - note: send maintains a sequence number locally, no sequence number field in the record
- *problem:* attacker could replay all records
- *solution:* use nonce

# Toy: control information

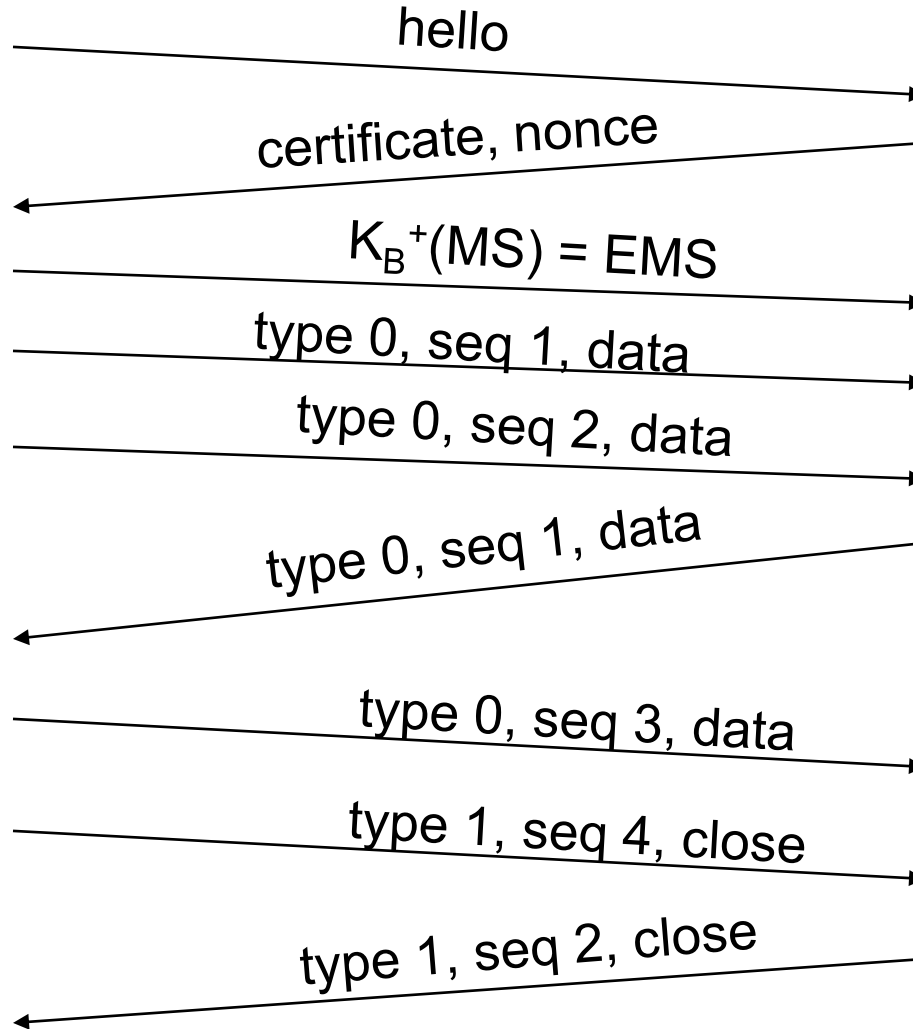
- *problem:* truncation attack:
  - attacker forges TCP connection close segment
  - one or both sides thinks there is less data than there actually is.
- *solution:* record types, with one type for closure
  - type 0 for data; type 1 for closure
- $MAC = Hash(M_x, \text{sequence} || \text{type} || \text{data})$



# Toy SSL: summary



*encrypted*



bob.com

# Toy SSL isn't complete

- how long are fields?
- which encryption protocols?
- want negotiation?
  - allow client and server to support different encryption algorithms
  - allow client and server to choose together specific algorithm before data transfer

# SSL cipher suite

- cipher suite
  - public-key algorithm
  - symmetric encryption algorithm
  - MAC algorithm
- SSL supports several cipher suites
- negotiation: client, server agree on cipher suite
  - client offers choice
  - server picks one

## common SSL symmetric ciphers

- DES – Data Encryption  
Standard: block
- 3DES – Triple strength: block
- RC2 – Rivest Cipher 2: block
- RC4 – Rivest Cipher 4: stream

## SSL Public key encryption

- RSA



# Real SSL: handshake (I)

## *Purpose*

1. server authentication
2. negotiation: agree on crypto algorithms
3. establish keys
4. client authentication (optional)

# Real SSL: handshake (2)

1. client sends list of algorithms it supports, along with client nonce
2. server chooses algorithms from list; sends back: choice + certificate + server nonce
3. client verifies certificate, extracts server's public key, generates pre\_master\_secret, encrypts with server's public key, sends to server
4. client and server independently compute encryption and MAC keys from pre\_master\_secret and nonces
5. client sends a MAC of all the handshake messages
6. server sends a MAC of all the handshake messages

# Real SSL: handshaking (3)

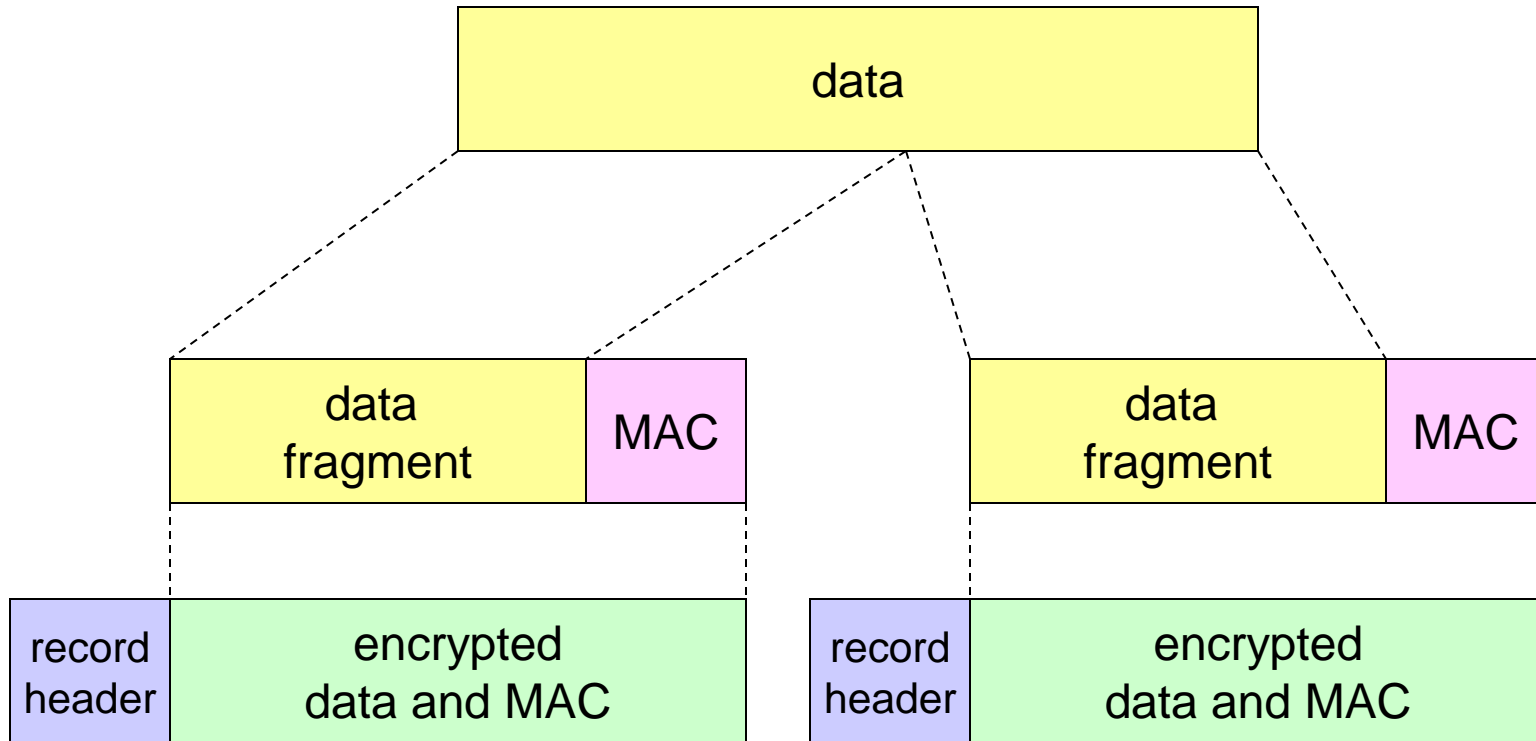
last 2 steps protect handshake from tampering

- client typically offers range of algorithms, some strong, some weak
- man-in-the middle could delete stronger algorithms from list
- last 2 steps prevent this
  - last two messages are encrypted

# Real SSL: handshaking (4)

- why two random nonces?
- suppose Trudy sniffs all messages between Alice & Bob
- next day, Trudy sets up TCP connection with Bob, sends exact same sequence of records
  - Bob (Amazon) thinks Alice made two separate orders for the same thing
  - solution: Bob sends different random nonce for each connection. This causes encryption keys to be different on the two days
  - Trudy's messages will fail Bob's integrity check

# SSL record protocol

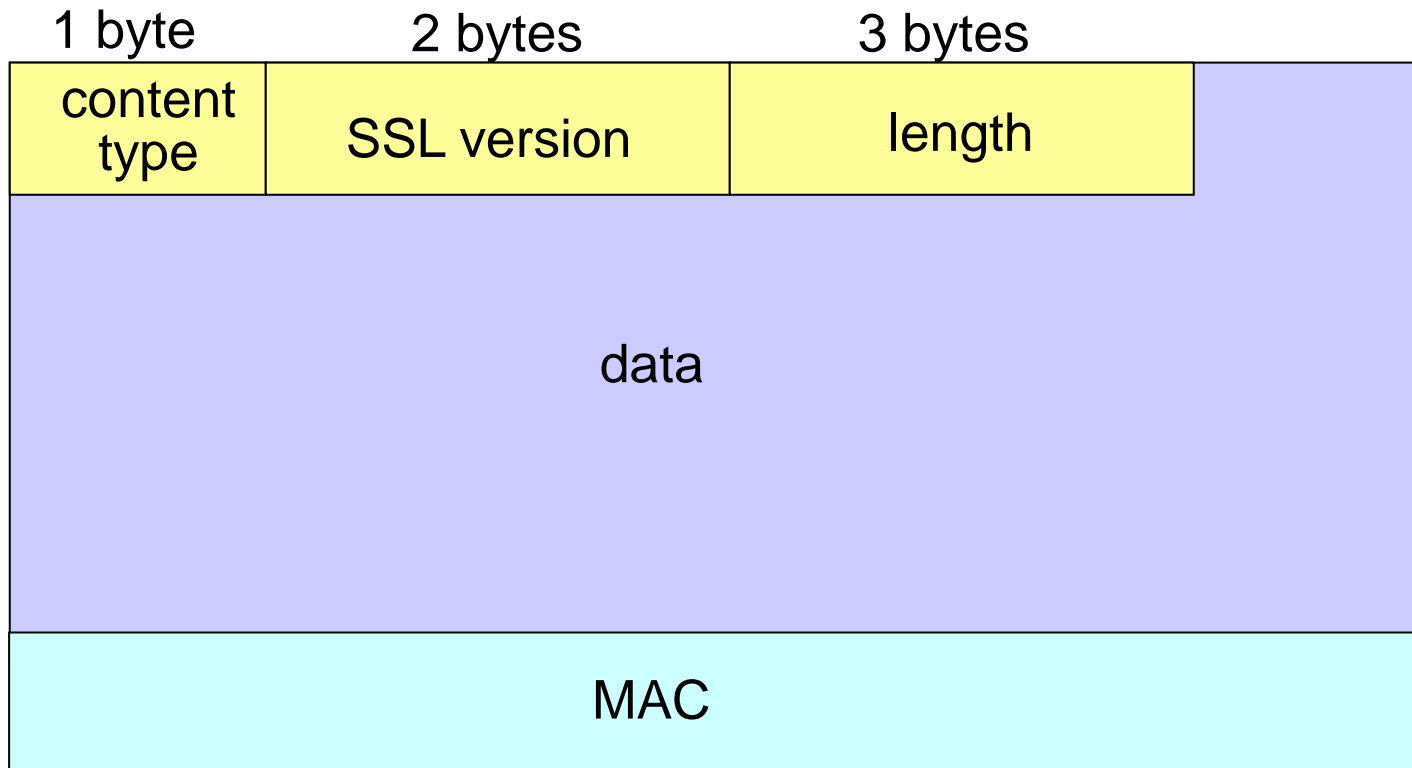


*record header:* content type; version; length

*MAC:* includes sequence number, MAC key  $M_x$

*fragment:* each SSL fragment  $2^{14}$  bytes (~16 Kbytes)

# SSL record format



data and MAC encrypted (symmetric algorithm)

# Chapter 8 roadmap

8.1 What is network security?

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8.3 Message integrity

8.4 Securing e-mail

8.5 Securing TCP connections: SSL

*8.6 Network layer security: IPsec*

8.7 Securing wireless LANs

8.8 Operational security: firewalls and IDS

# What is network-layer confidentiality ?

*between two network entities:*

- sending entity encrypts datagram payload, payload could be:
  - TCP or UDP segment, ICMP message, OSPF message ....
- all data sent from one entity to other would be hidden:
  - web pages, e-mail, P2P file transfers, TCP SYN packets
  - ...
- “blanket coverage”

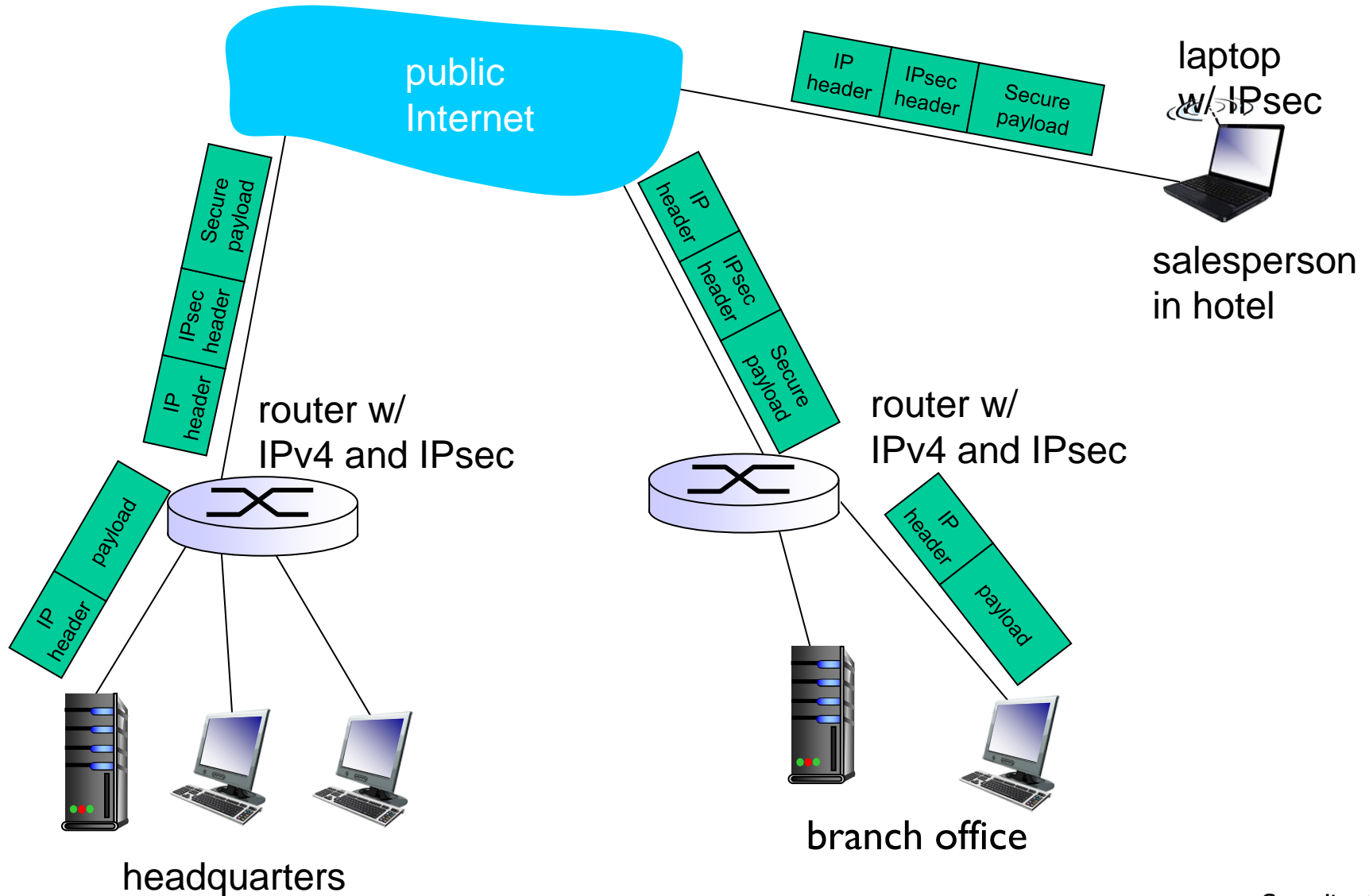


# Virtual Private Networks (VPNs)

## *motivation:*

- institutions often want private networks for security.
  - costly: separate routers, links, DNS infrastructure.
- VPN: institution's inter-office traffic is sent over public Internet instead
  - encrypted before entering public Internet
  - logically separate from other traffic
- USTC VPN
  - Connect to USTC router from outside USTC
  - Access resources (e.g., digital library) as if the host is within the USTC network.

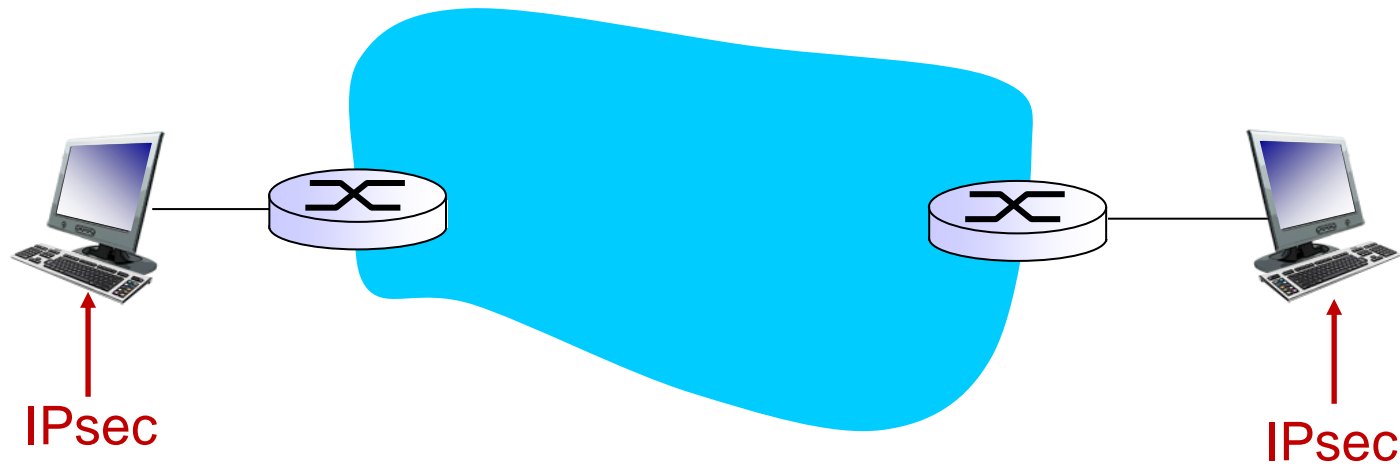
# Virtual Private Networks (VPNs)



# IPsec services

- data integrity
- origin authentication
- replay attack prevention
- confidentiality
  
- two protocols providing different service models:
  - Authentication Header (AH) protocol: provides source authentication & data integrity but *not* confidentiality
  - Encapsulation Security Protocol (ESP): provides source authentication, data integrity, *and confidentiality*
    - more widely used than AH

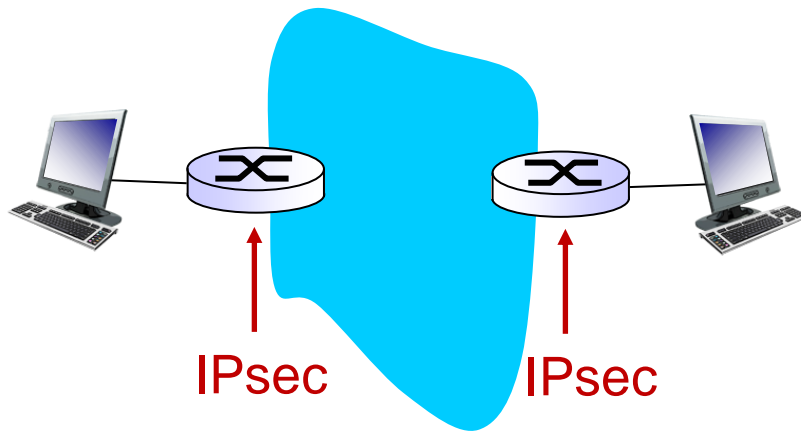
# IPsec transport mode



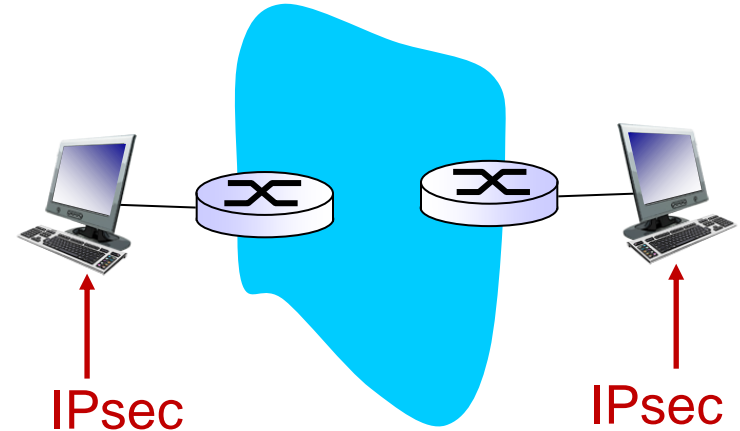
- IPsec datagram emitted and received by end-system
- protects upper level protocols

Retain the original IP header unprotected, only protect the payload.

# IPsec – tunneling mode



- edge routers IPsec-aware



- hosts IPsec-aware
- Host connect to router

Protect the entire original IP packet by encapsulate it as payload of a new IP packet

# Four combinations are possible!

Transport mode with AH	Transport mode with ESP
Tunnel mode with AH	Tunnel mode with ESP

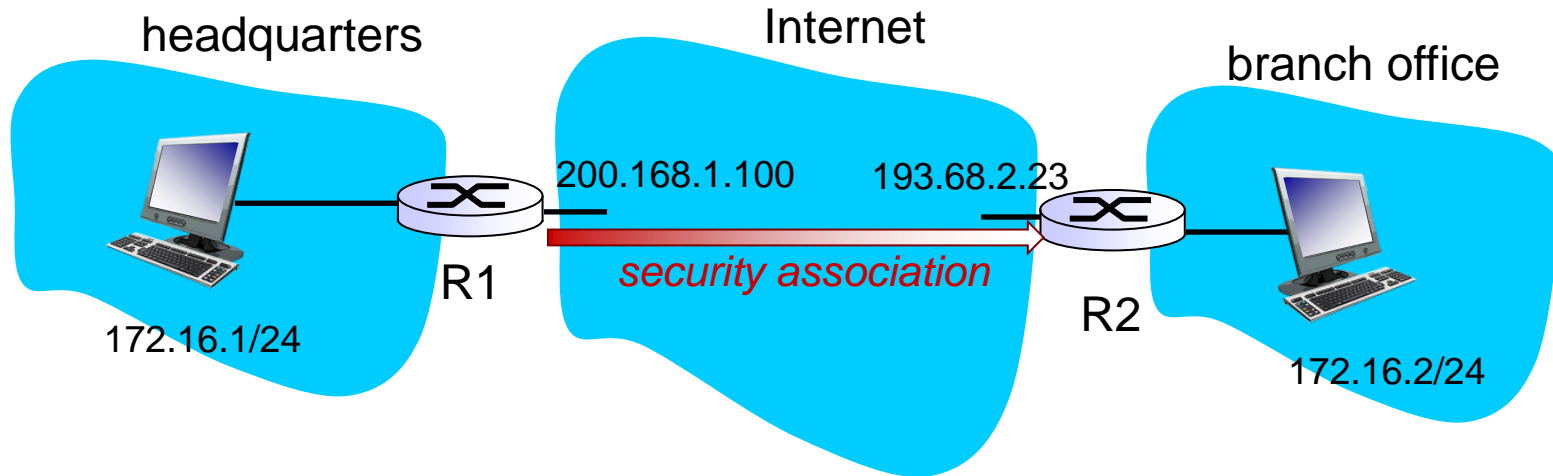


most common and  
most important

# Security associations (SAs)

- before sending data, set up a logical connection called “**security association (SA)**” from sending to receiving entity
  - SAs are simplex(单向): for only one direction
- ending, receiving entitles maintain *state information* about SA
  - recall: TCP endpoints also maintain state info
  - IP is connectionless; IPsec is connection-oriented!
- how many SAs in VPN w/ headquarters, branch office, and n traveling salespeople?

# Example SA from R1 to R2



## *R1 stores for SA:*

- 32-bit SA identifier: *Security Parameter Index (SPI)*
- origin SA interface (200.168.1.100)
- destination SA interface (193.68.2.23)
- type of encryption used (e.g., 3DES with CBC)
- encryption key
- type of integrity check used (e.g., HMAC with MD5)
- authentication key

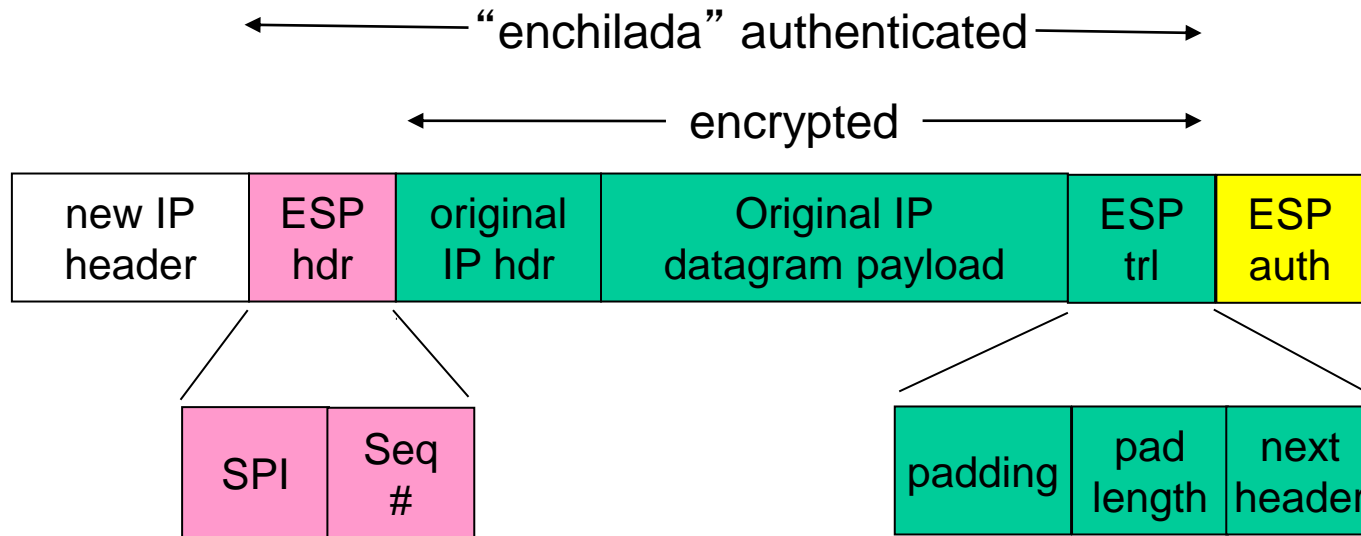


# Security Association Database (SAD)

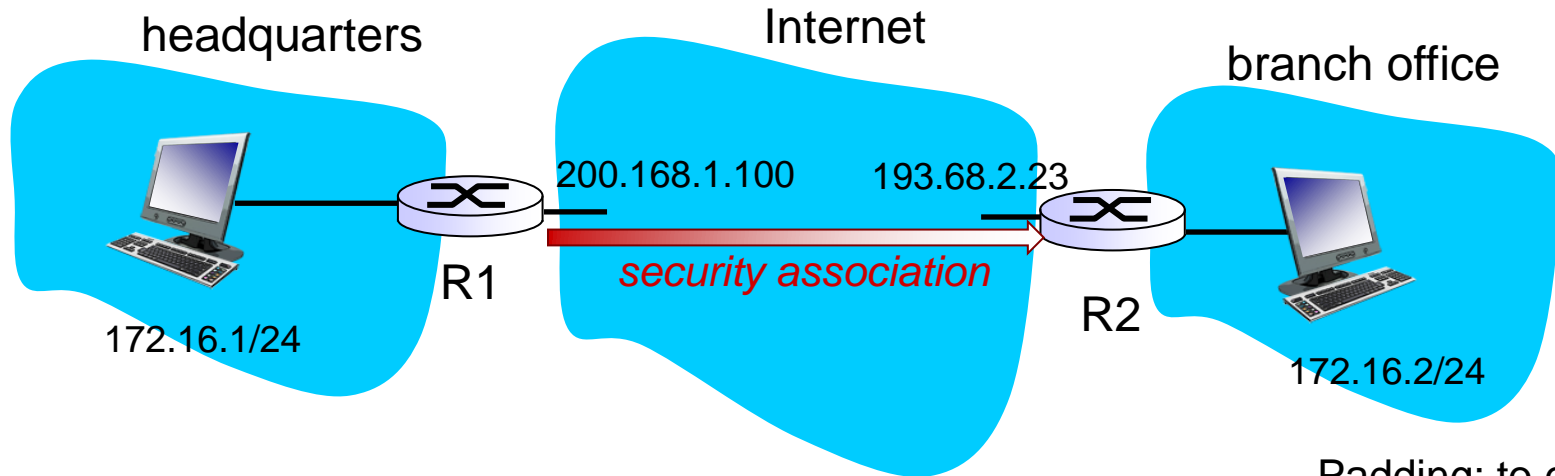
- endpoint holds SA state in *security association database (SAD)*, where it can locate them during processing.
- with  $n$  salespersons,  $2 + 2n$  SAs in R1's SAD
  - $2n$  for  $n$  salespersons and 2 for the branch.
- when sending IPsec datagram, R1 accesses SAD to determine how to process datagram.
- when IPsec datagram arrives to R2, R2 examines SPI in IPsec datagram, indexes SAD with SPI, and processes datagram accordingly.

# IPsec datagram

focus for now on tunnel mode with ESP

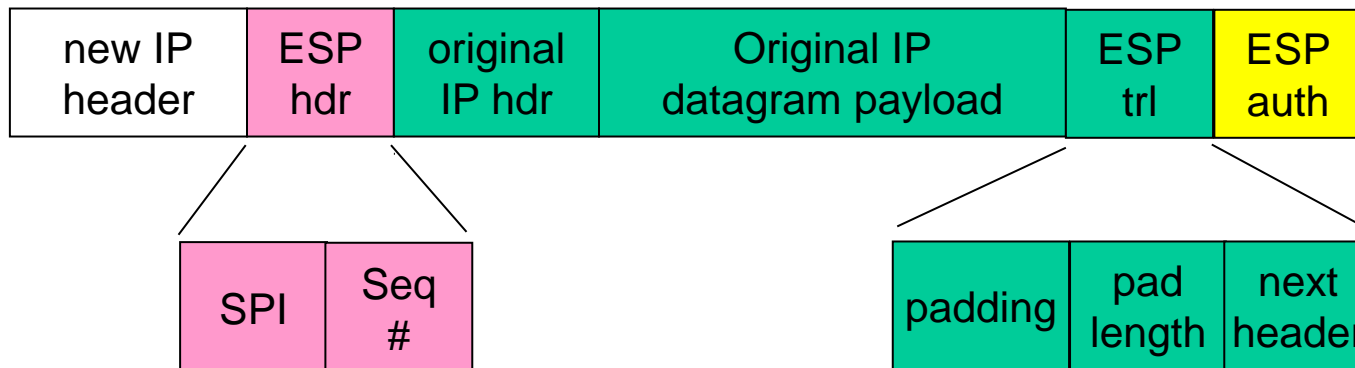


# What happens?



← "enchilada" authenticated →  
← encrypted →

Padding: to ensure  
the data for  
encryption is times of  
block size

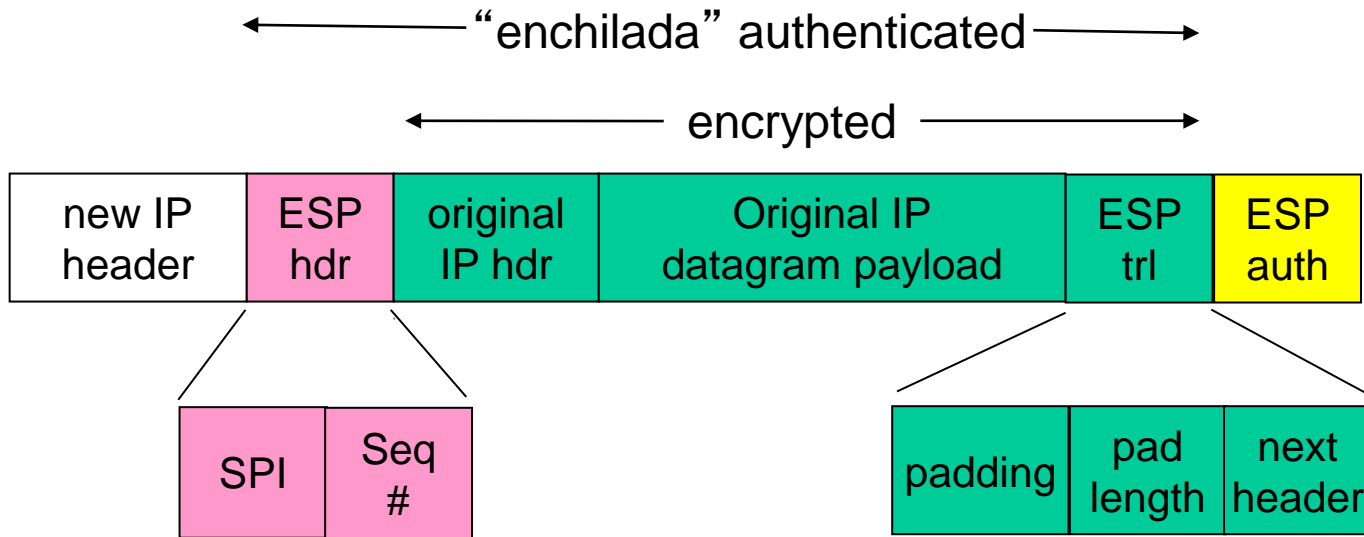


next header: UDP or TCP?

# RI: convert original datagram to IPsec datagram

- appends to back of original datagram (which includes original header fields!) an “ESP trailer” field.
- encrypts result using algorithm & key specified by SA.
- appends to front of this encrypted quantity the “ESP header”, creating “enchilada”.
- creates authentication MAC over the *whole enchilada*, using algorithm and key specified in SA;
- appends MAC to back of enchilada, forming *payload*;
- creates brand new IP header, with all the classic IPv4 header fields, which it appends before payload

# Inside the enchilada:



- ESP trailer: Padding for block ciphers
- ESP header:
  - SPI, so receiving entity knows what to do
  - Sequence number, to thwart replay attacks
- MAC in ESP auth field is created with shared authentication key

# IPsec sequence numbers

- for new SA, sender initializes seq. # to 0
- each time datagram is sent on SA:
  - sender increments seq # counter
  - places value in seq # field
- goal:
  - prevent attacker from sniffing and replaying a packet
  - receipt of duplicate, authenticated IP packets may disrupt service

# Security Policy Database (SPD)

- policy: For a given datagram, sending entity needs to know if it should use IPsec (or direct send it out)
- needs also to know which SA to use
  - may use: source and destination IP address; protocol number
- info in SPD indicates “what” to do with arriving datagram
- info in SAD indicates “how” to do it

# Summary: IPsec services



new IP header	ESP hdr	original IP hdr	Original IP datagram payload	ESP trl	ESP auth
---------------	---------	-----------------	------------------------------	---------	----------

- suppose Trudy sits somewhere between R1 and R2. she doesn't know the keys.
  - will Trudy be able to see original contents of datagram? How about source, dest IP address, transport protocol, application port? **Not be able to see encrypted part.**
  - flip bits without detection? **Fail integrity check using MAC**
  - masquerade as R1 using R1's IP address? **Fail integrity check at R2**
  - replay a datagram? **Fail with wrong sequence number**



# IKE: Internet Key Exchange

- *previous examples:* manual establishment of IPsec SAs in IPsec endpoints:

*Example SA*

SPI: 12345

Source IP: 200.168.1.100

Dest IP: 193.68.2.23

Protocol: ESP

Encryption algorithm: 3DES-cbc

HMAC algorithm: MD5

Encryption key: 0x7aeaca...

HMAC key:0xc0291f...

- manual keying is impractical for VPN with 100s of endpoints
- instead use *IPsec IKE (Internet Key Exchange)*

# IKE phases

- IKE has two phases
  - *phase 1*: establish bi-directional IKE SA
    - note: IKE SA different from IPsec SA
    - aka ISAKMP security association
  - *phase 2*: ISAKMP is used to securely negotiate IPsec pair of SAs
- phase 1 has two modes: aggressive mode and main mode
  - aggressive mode uses fewer messages
  - main mode provides identity protection and is more flexible

# IPsec summary

- IKE message exchange for algorithms, secret keys, SPI numbers
- either AH or ESP protocol (or both)
  - AH provides integrity, source authentication
  - ESP protocol (with AH) additionally provides encryption
- IPsec peers can be two end systems, two routers/firewalls, or a router/firewall and an end system

# Chapter 8 roadmap

- 8.1 What is network security?
- 8.2 Principles of cryptography
- 8.3 Message integrity
- 8.4 Securing e-mail
- 8.5 Securing TCP connections: SSL
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- 8.7 Securing wireless LANs*
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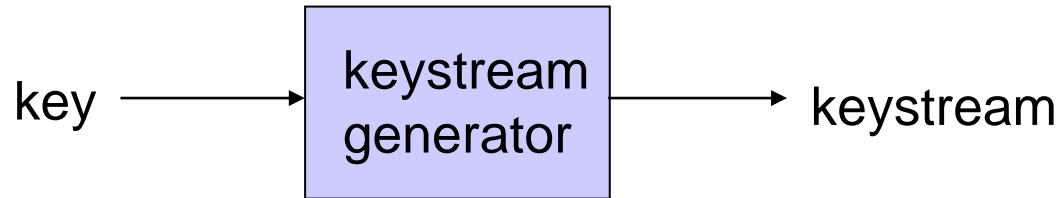
# WEP design goals



- symmetric key crypto
  - confidentiality
  - end host authorization
  - data integrity
- self-synchronizing: each packet separately encrypted
  - given encrypted packet and key, can decrypt; can continue to decrypt packets when preceding packet was lost (unlike Cipher Block Chaining (CBC) in block ciphers)
- Efficient
  - implementable in hardware or software



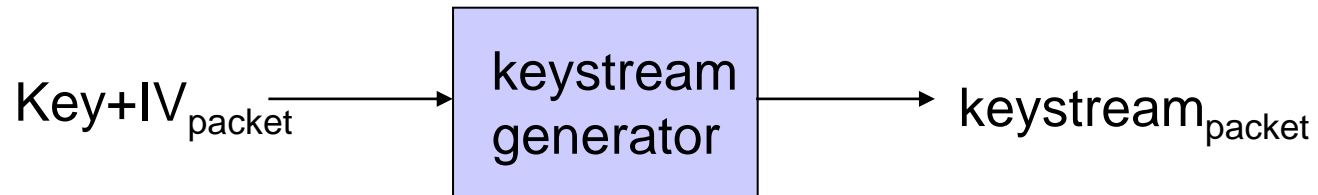
# Review: symmetric stream ciphers



- *combine each byte of keystream with byte of plaintext to get ciphertext:*
  - $m(i)$  = ith unit of message
  - $ks(i)$  = ith unit of keystream
  - $c(i)$  = ith unit of ciphertext
  - $c(i) = ks(i) \oplus m(i)$  ( $\oplus$  = exclusive or)
  - $m(i) = ks(i) \oplus c(i)$
- WEP uses RC4

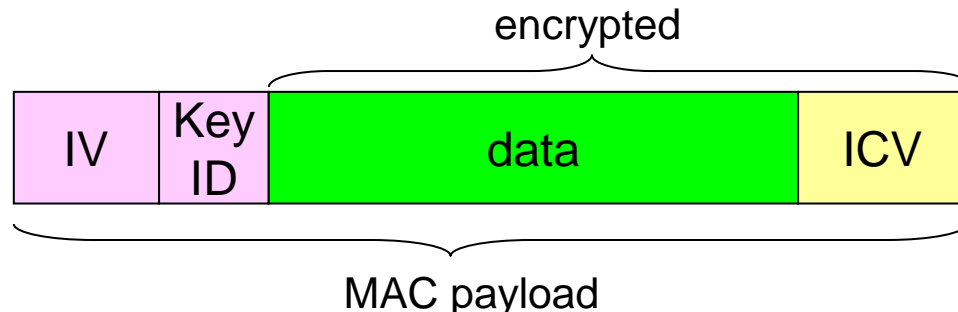
# Stream cipher and packet independence

- recall design goal: each packet separately encrypted
- if for frame  $n+1$ , use keystream from where we left off for frame  $n$ , then each frame is not separately encrypted
  - need to know where we left off for packet  $n$
- WEP approach: initialize keystream with key + new IV for each packet:



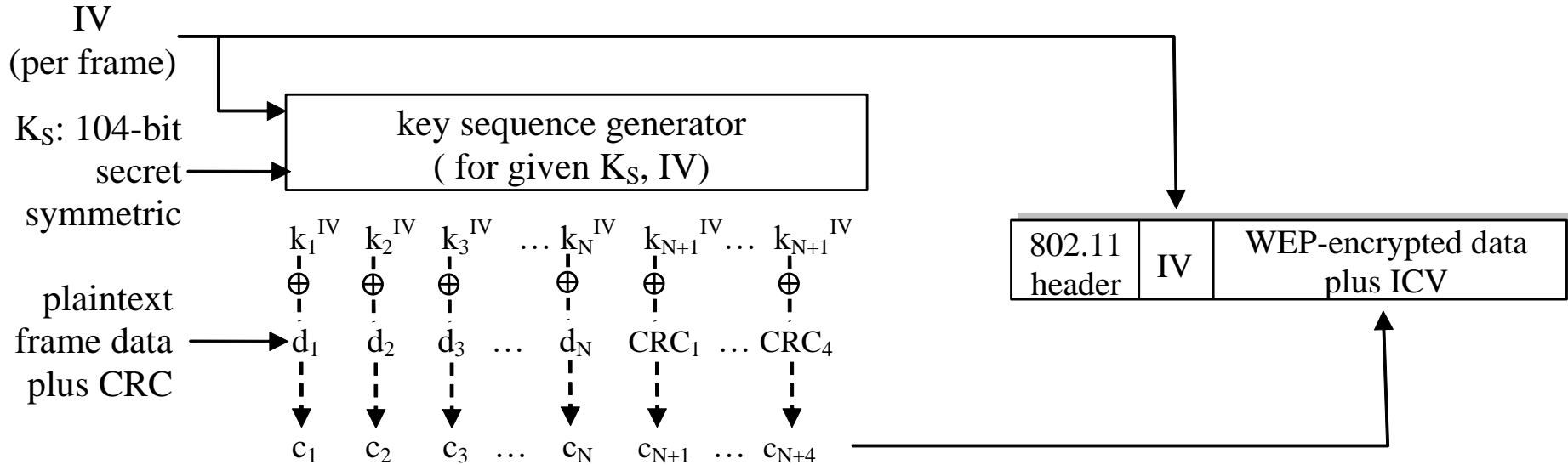
# WEP encryption (I)

- sender calculates Integrity Check Value (ICV, four-byte hash/CRC over data)
- each side has 104-bit shared key
- sender creates 24-bit initialization vector (IV), appends to key: gives 128-bit key
- sender also appends keyID (in 8-bit field)
- 128-bit key inputted into pseudo random number generator to get keystream
- data in frame + ICV is encrypted with RC4:
  - bytes of keystream are XORed with bytes of data & ICV
  - IV & keyID are appended to encrypted data to create payload
  - payload inserted into 802.11 frame



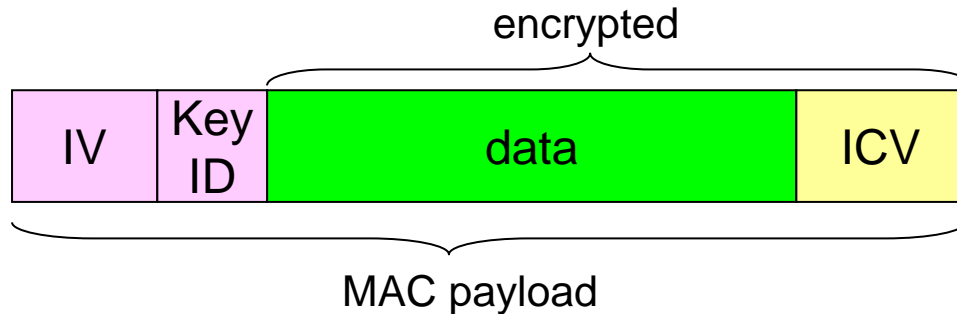


# WEP encryption (2)



*new IV for each frame*

# WEP decryption overview



- receiver extracts IV
- inputs IV, shared secret key into pseudo random generator, gets keystream
- XORs keystream with encrypted data to decrypt data + ICV
- verifies integrity of data with ICV
  - note: message integrity approach used here is different from MAC (message authentication code) and signatures (using PKI).

# Breaking 802.11 WEP encryption

## *security hole:*

- 24-bit IV, one IV per frame, -> IV's eventually reused
- IV transmitted in plaintext -> IV reuse detected

## *attack:*

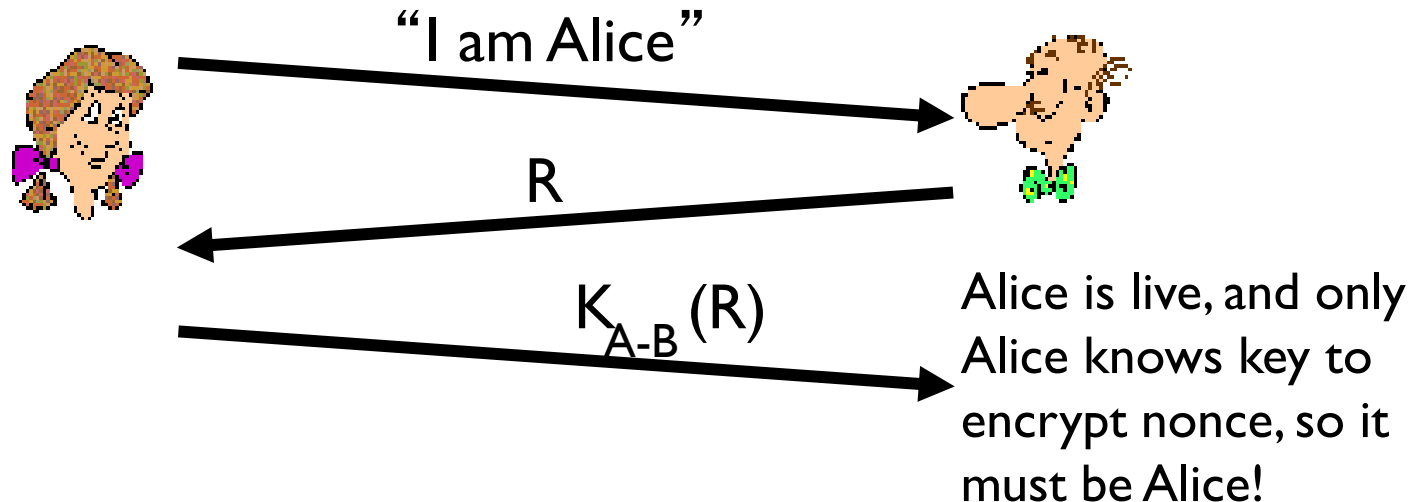
- Trudy causes Alice to encrypt known plaintext  $d_1 d_2 d_3 d_4 \dots$
- Trudy sees:  $c_i = d_i \text{ XOR } k_i^{\text{IV}}$
- Trudy knows  $c_i d_i$ , so can compute  $k_i^{\text{IV}}$
- Trudy knows encrypting key sequence  $k_1^{\text{IV}} k_2^{\text{IV}} k_3^{\text{IV}} \dots$
- Next time Trudy finds that IV is used, Trudy can decrypt with  $k_1^{\text{IV}} k_2^{\text{IV}} k_3^{\text{IV}} \dots$

WEP is replaced by WPA and WPA2.

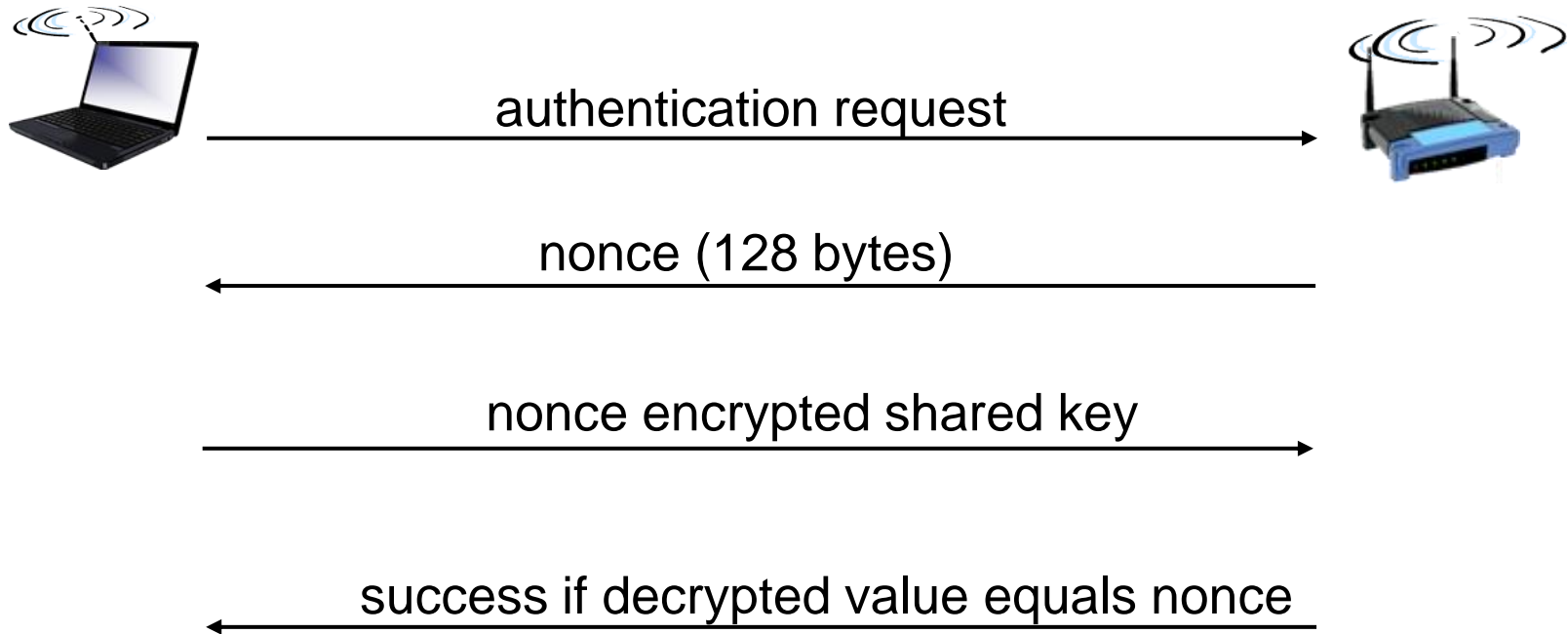
# End-point authentication w/ nonce

**Nonce:** number (R) used only *once* –*in-a-lifetime*

**How to prove Alice “live”:** Bob sends Alice **nonce**, R. Alice must return R, encrypted with shared secret key



# WEP authentication



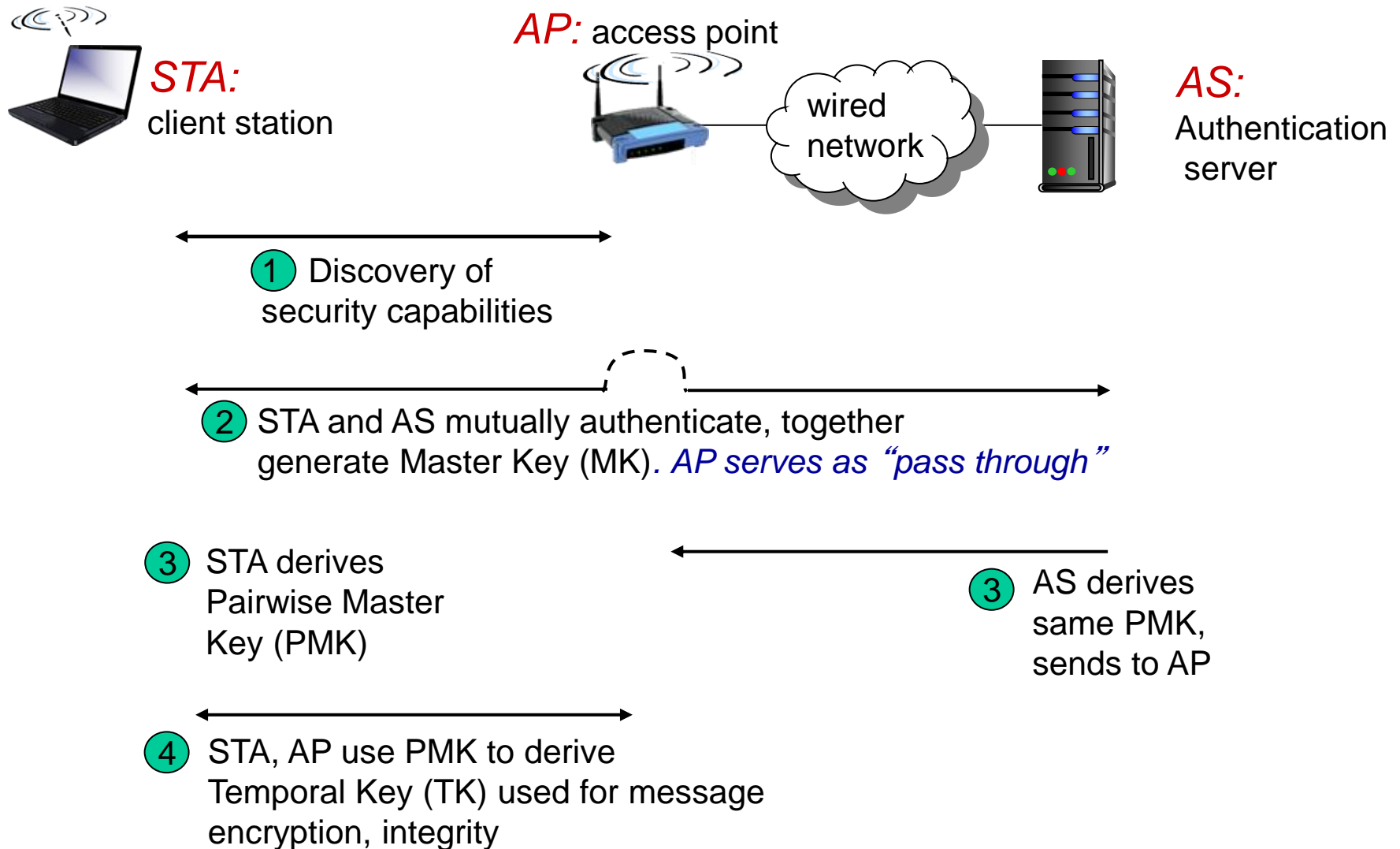
## *Notes:*

- not all APs do it, even if WEP is being used
- AP indicates if authentication is necessary in beacon frame
- done before association

# 802.11i: improved security

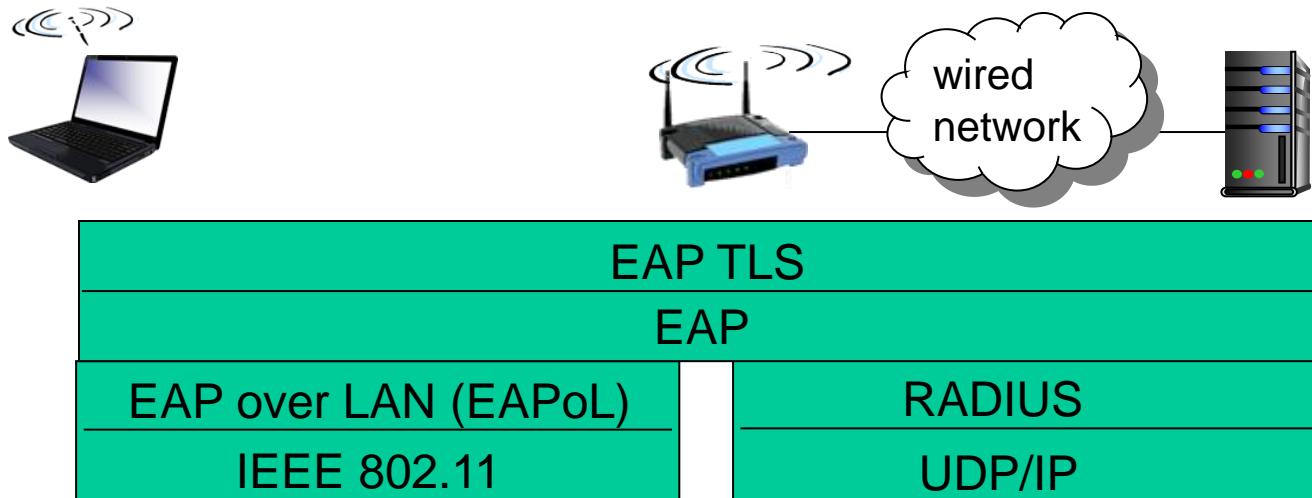
- numerous (stronger) forms of encryption possible
- provides key distribution
- uses authentication server separate from access point

# 802.11i: four phases of operation



# EAP: extensible authentication protocol

- EAP: end-end client (mobile) to authentication server protocol
- EAP sent over separate “links”
  - mobile-to-AP (EAP over LAN)
  - AP to authentication server (RADIUS over UDP)





# Chapter 8 roadmap

8.1 What is network security?

8.2 Principles of cryptography

8.3 Message integrity

8.4 Securing e-mail

8.5 Securing TCP connections: SSL

8.6 Network layer security: IPsec

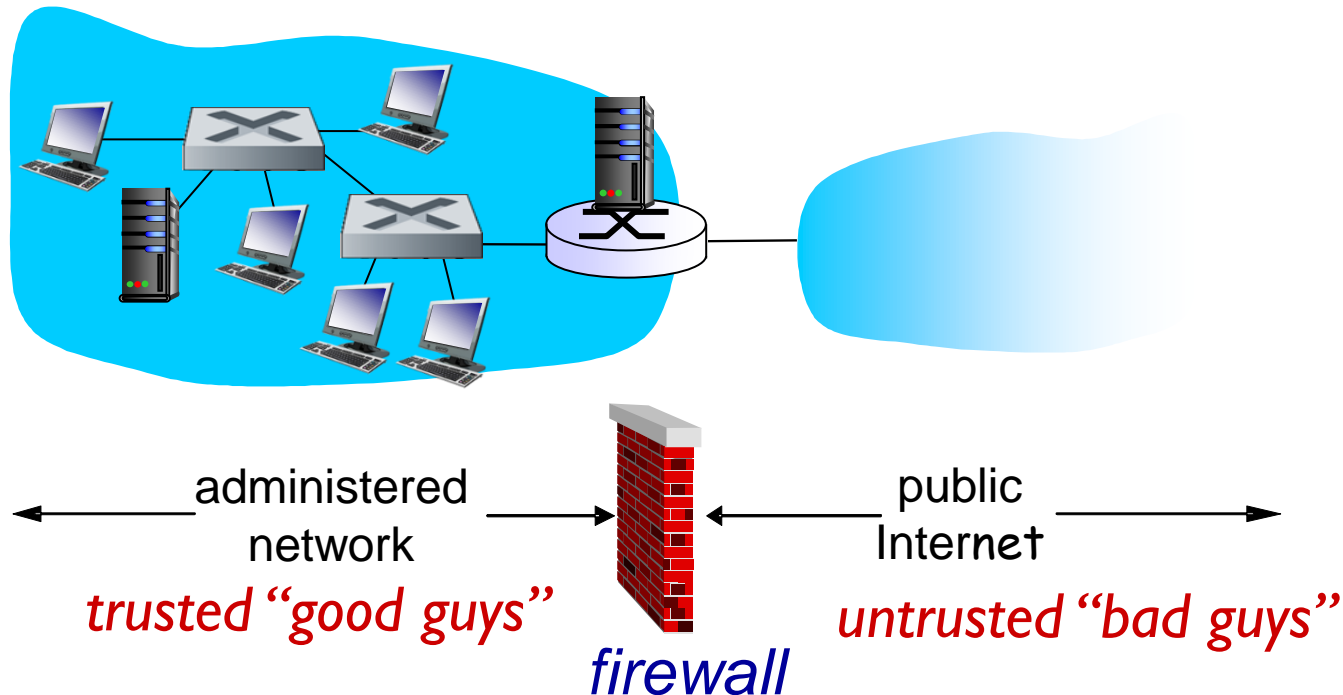
8.7 Securing wireless LANs

*8.8 Operational security: firewalls and IDS*

# Firewalls

## *firewall*

isolates organization's internal net from larger Internet, allowing some packets to pass, blocking others



# Firewalls: why

prevent denial of service attacks:

- SYN flooding: attacker establishes many bogus TCP connections, no resources left for “real” connections

prevent illegal modification/access of internal data

- e.g., attacker replaces CIA’s homepage with something else

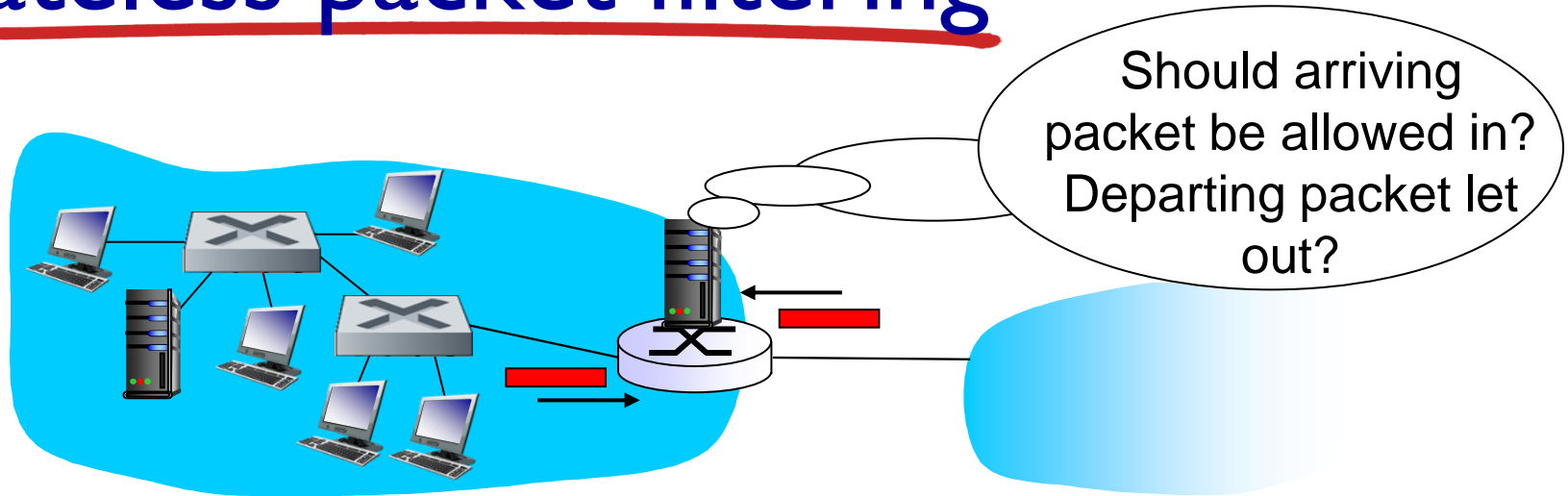
allow only authorized access to inside network

- set of authenticated users/hosts

three types of firewalls:

- stateless packet filters
- stateful packet filters
- application gateways

# Stateless packet filtering



- internal network connected to Internet via *router firewall*
- router *filters packet-by-packet*, decision to forward/drop packet based on:
  - source IP address, destination IP address
  - TCP/UDP source and destination port numbers
  - ICMP message type
  - TCP SYN and ACK bits

# Stateless packet filtering: example

- *example 1*: block incoming and outgoing datagrams with IP protocol field = 17 and with either source or dest port = 23
  - *result*: all incoming, outgoing UDP flows and telnet connections are blocked
- *example 2*: block inbound TCP segments with ACK=0.
  - *result*: prevents external clients from making TCP connections with internal clients, but allows internal clients to connect to outside.

# Stateless packet filtering: more examples

<i>Policy</i>	<i>Firewall Setting</i>
No outside Web access.	Drop all outgoing packets to any IP address, port 80
No incoming TCP connections, except those for institution's public Web server only.	Drop all incoming TCP SYN packets to any IP except 130.207.244.203, port 80
Prevent Web-radios from eating up the available bandwidth.	Drop all incoming UDP packets - except DNS and router broadcasts.
Prevent your network from being used for a smurf DoS attack.	Drop all ICMP packets going to a "broadcast" address (e.g. 130.207.255.255).
Prevent your network from being tracerouted	Drop all outgoing ICMP TTL expired traffic

# Access Control Lists

**ACL:** table of rules, applied top to bottom to incoming packets:  
(action, condition) pairs: looks like OpenFlow forwarding (Ch. 4)!

action	source address	dest address	protocol	source port	dest port	flag bit
allow	222.22/16	outside of 222.22/16	TCP	> 1023	80	any
allow	outside of 222.22/16	222.22/16	TCP	80	> 1023	ACK
allow	222.22/16	outside of 222.22/16	UDP	> 1023	53	---
allow	outside of 222.22/16	222.22/16	UDP	53	> 1023	----
deny	all	all	all	all	all	all

# Stateful packet filtering

- *stateless packet filter*: heavy handed tool
  - admits packets that “make no sense,” e.g., dest port = 80, ACK bit set, even though no TCP connection established:

action	source address	dest address	protocol	source port	dest port	flag bit
allow	outside of 222.22/16	222.22/16	TCP	80	> 1023	ACK

- *stateful packet filter*: track status of every TCP connection
  - track connection setup (SYN), teardown (FIN): determine whether incoming, outgoing packets “makes sense”
  - timeout inactive connections at firewall: no longer admit packets



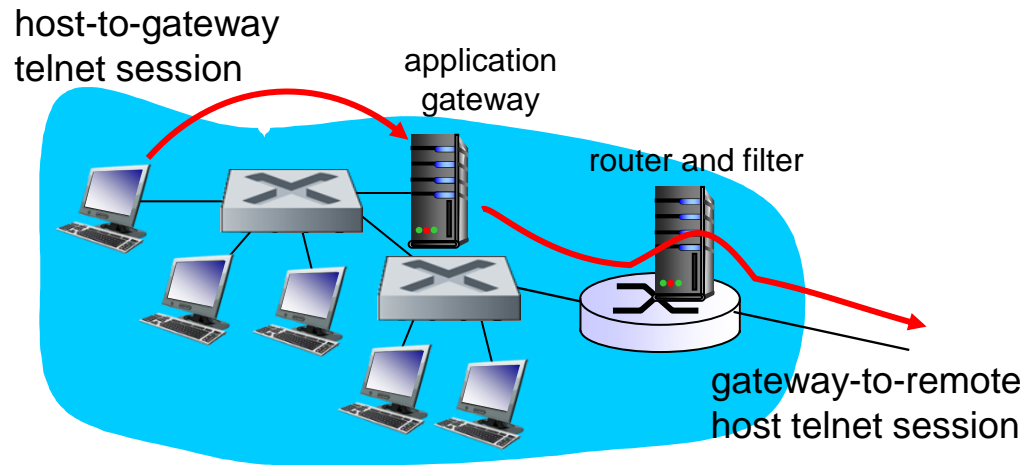
# Stateful packet filtering

ACL augmented to indicate need to check connection state table before admitting packet

action	source address	dest address	proto	source port	dest port	flag bit	check conxion
allow	222.22/16	outside of 222.22/16	TCP	> 1023	80	any	
allow	outside of 222.22/16	222.22/16	TCP	80	> 1023	ACK	<b>X</b> 是否有连接？
allow	222.22/16	outside of 222.22/16	UDP	> 1023	53	---	
allow	outside of 222.22/16	222.22/16	UDP	53	> 1023	----	<b>X</b> 是否有DNS query？
deny	all	all	all	all	all	all	

# Application gateways

- filter packets on application data as well as on IP/TCP/UDP fields.
- *example:* allow select internal users to telnet outside



1. require all telnet users to telnet through gateway.
2. for authorized users, gateway sets up telnet connection to dest host. Gateway relays data between 2 connections
3. router filter blocks all telnet connections not originating from gateway.

# Limitations of firewalls, gateways

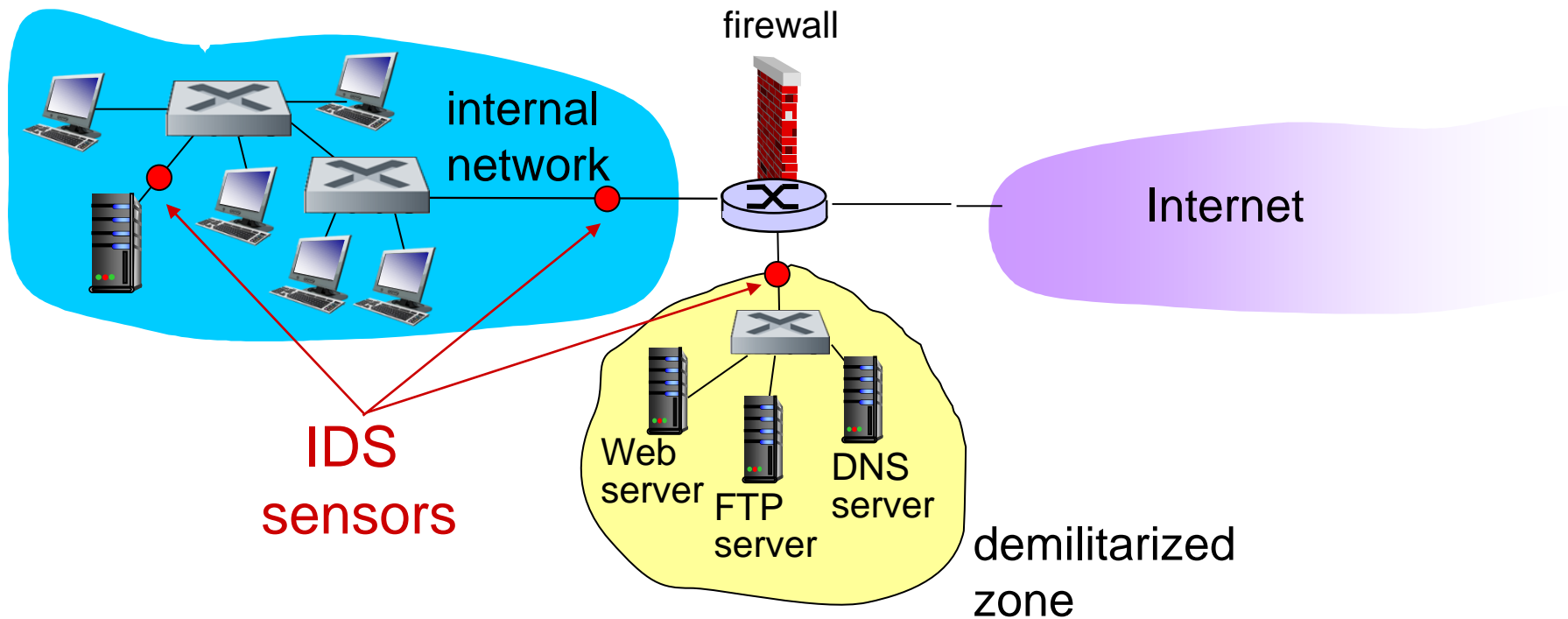
- *IP spoofing*: router can't know if data “really” comes from claimed source
- if multiple app's. need special treatment, each has own app. gateway
- client software must know how to contact gateway.
  - e.g., must set IP address of proxy in Web browser
- filters often use all or nothing policy for UDP
- *tradeoff*: degree of communication with outside world, level of security
- many highly protected sites still suffer from attacks

# Intrusion detection systems

- packet filtering:
  - operates on TCP/IP headers only
  - no correlation check among sessions
- *IDS: intrusion detection system*
  - *deep packet inspection*: look at packet contents (e.g., check character strings in packet against database of known virus, attack strings)
  - *examine correlation* among multiple packets
    - port scanning
    - network mapping
    - DoS attack

# Intrusion detection systems

multiple IDSs: different types of checking at different locations



# Network Security (summary)

## basic techniques.....

- cryptography (symmetric and public)
- message integrity
- end-point authentication

## .... used in many different security scenarios

- secure email
- secure transport (SSL)
- IP sec
- 802.11

## operational security: firewalls and IDS