

Refactoring and Verification in Rust
Expanding the REM toolchain
And providing a novel approach to verification

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Disclaimer

I hereby declare that the work undertaken in this R&D project has been conducted by me alone, except where indicated in the text. I conducted this work between September 2024 and July 2025, during which period I was an Engineering (R&D) student at the Australian National University. This report, in whole or any part of it, has not been submitted to this or any other university for a degree.

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Abstract

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Glossary

Term	Definition
AoL	Adventure of a Lifetime. ?.
GC	Garbage Collector.
REM	Rust Extraction Maestro, the refactoring toolchain implemented in Borrowing Without Sorrowing ?
AST	Abstract Syntax Tree.
CFG	Control Flow Graph.
IR	Intermediate Representation.
MIR	Mid-level Intermediate Representation.
LLBC	Low-Level Borrow Calculus
SMT	Satisfiability Modulo Theories.
CMBC	C Bounded Model Checker
LSP	Language Server Protocol.
rustc	The Rust compiler.
IDE	Integrated Development Environment.

Table 1: Glossary of Terms

Introduction

Rust is a modern systems programming language that aims to reconcile low-level control with strong safety guarantees. This is in stark contrast to previous norms, where a program was often forced to choose between low-level control (with languages such as C/C++) and memory safety (with garbage collected languages such as Java). However, as prior work like *Adventure of a Lifetime* (AoL) and the subsequent *Borrowing without Sorrowing* plugin have shown, these very features that make Rust robust also complicate automated refactoring - particularly “Extract Method” transformations that are straightforward in garbage collected (GC) languages. Building on these initial efforts, this paper extends the Rusty Extraction Maestro (REM) tool, removing dependencies on outdated IDEs, handling a broader array of Rust constructs (including `async/await` and generics), and introducing a more rigorous verification strategy to ensure that refactorings preserve semantics. By providing a standalone command-line interface and a proof-of-concept Visual Studio Code extension, this project lays the groundwork for seamlessly integrating safe, behavior-preserving “Extract Method” refactoring into mainstream Rust development workflows.

0.1 Motivation and Background

It is an inevitable fact of software development that the engineer will spend a substantial amount of their time modifying and improving existing codebases. Such modifications will range from small bug fixes to complete architecture overhauls. Regardless of how large (or small) the situation is, refactoring - systematic restructuring of the codebase without altering its external behaviour - plays a pivotal role in maintaining code quality and limiting code smell over time. Despite refactoring being commonplace in many systems programming languages like C++ or Java, implementing automated refactoring tools remains challenging, particularly in languages with advanced features and strict safety guarantees, such as Rust.

When it comes to memory management, two schools of thought have dominated for the last half-century. One tradition leans on automated techniques like garbage collection,

while the other relies on manual, explicit management. In this landscape, Rust stands out by enforcing memory safety and concurrency through its rigorous system of ownership, borrowing, and lifetimes. Yet, the very features that make Rust safe also introduce massive hurdles for automated refactoring. A naive, “copy and paste” implementation of extract method refactoring that works in more forgiving languages like Python and Java will hit the proverbial brick wall that is the Rust borrow checker. Whilst the borrow checker and other compilation checks will prevent unintended consequences, the programmer must now deal with a whole host of other issues.

Against this backdrop, *Adventure of a Lifetime* (AoL) ? introduced the conceptual framework for automated “Extract Method” refactoring in Rust. From there, *Borrowing without Sorrowing* ? implemented the concepts outlined in *Adventure of a Lifetime* as an IntelliJ IDEA plugin. This project now builds on those foundations, aiming to extend, improve and explore new frontiers for automated Rust refactoring.

0.1.1 Do People Actually Refactor?

Despite widespread agreement on the value of continuous, incremental code improvements, one might still wonder: do development teams truly refactor in practice, or is it just an idealized principle? A recent large-scale survey of 1,183 professional developers by the JetBrains team, *One Thousand and One Stories: A Large-Scale Survey of Software Refactoring* ?, sheds light on this question. The results show that refactoring is unquestionably a routine part of modern software engineering:

- **Frequent Refactoring:** Over 77% of surveyed developers reported refactoring code at least once a week, with 40.6% doing so “almost every day” and another 36.9% at least every week.
- **Lengthy Sessions:** Refactoring can be time-intensive; 66.3% of respondents indicated they had spent an hour or more in a single session refactoring their code.
- **Rare “Non-Refactorers”:** Only 2.5% stated that they “never” refactor code. Even within that group, many had at least renamed or relocated code elements—activities the survey authors classify as refactoring.

These findings confirm that refactoring is more than just a best-practice slogan: most teams regularly restructure and clean up their code. At the same time, the significant number of hour-plus refactoring sessions suggests that technical debt can build up to a point where larger-scale efforts become necessary. Understanding how, why, and when developers refactor can thus guide the design of more effective refactoring tools and strategies.

0.1.2 Importance of Refactoring in Software Engineering

In large projects, codebases can quickly become very unwieldy. Refactoring - the process of improving a programs internal structure without altering its external behaviour ? - is a crucial tool that helps engineers keep the technical debt at a manageable level. It ensures that system complexity remains manageable by promoting cleaner abstractions, fewer side effects, and much clearer data flows.

Additionally, whilst refactoring itself doesn't add new features, it often uncovers latent defects and unexpected behaviours or side effects. It can be used to bring clarity to confusing, hundred-plus line methods, and makes testing significantly more straightforward.

Finally, well-factored code is much easier to understand and thus modify. This, in turn, will speed up future iterations, feature implementations, and onboarding of new team members.

0.1.3 What is “Extract Method” Refactoring?

“Extract Method” is a common refactoring technique where a contiguous block of code (often a few lines or a loop body) is extracted into its own function. The original code is replaced with a call to this newly created function. By breaking down the large methods or blocks into smaller, aptly named functions, readers can far more quickly grasp the purpose of each piece of logic.

In Java or other garbage-collected languages, method extraction is straightforward: you collect parameters, create a new function, and replace the original lines with a function call. Rust, however, introduces additional layers of complexity due to ownership and borrowing. The example below (in Java) illustrates how this technique works at a superficial level.

```
1 // Before:
2 public void processNumbers(List<Integer> nums) {
3     int sum = 0;
4     for (int n : nums) {
5         sum += n;
6     }
7     System.out.println("Sum is: " + sum);
8 }
9
10 // After:
11 public void processNumbers(List<Integer> nums) {
12     int sum = calculateSum(nums);
13     System.out.println("Sum is: " + sum);
14 }
15
16 private int calculateSum(List<Integer> nums) {
17     int sum = 0;
18     for (int n : nums) {
19         sum += n;
```

```
20     }  
21     return sum;  
22 }
```

0.1.4 Why Rust Poses Unique Challenges

1. **Ownership:** Each value in Rust has a single owner that governs the value's lifetime. Once the owner goes out of scope, the value is dropped. This design eliminates many types of memory errors (such as use-after-free), but can complicate refactoring: relocating or splitting functionality into separate functions may transfer or divide ownership in unexpected ways.

```
1 fn main() {  
2     let s1 = String::from("Hello");  
3     let s2 = s1; // Ownership of the string data moves to s2  
4     // The following line would be invalid if uncommented:  
5     // println!("{}", s1);  
6     // ~ value borrowed here after move  
7     println!("{}", s2); // Prints "Hello"  
8 }
```

In the above example, `s1` is “moved” into `s2`, making `s1` no longer valid afterward. If a refactoring operation was to pass `s1` into a function that takes ownership, the original owner cannot access it unless ownership is returned. Alternatively, we could end up with code that refuses to compile, forcing the developer to manually debug and fix the refactored code.

2. **Borrowing:** Instead of copying or moving data, Rust encourages passing references. However, these references can be immutable (`&T`) or mutable (`&mut T`), and only one active mutable reference is allowed at a time. When extracting logic into a separate function, it is important to ensure that references remain valid and continue to adhere to Rust's borrowing rules. The examples below demonstrate how borrowing can be effectively handled. However, if multiple references with overlapping mutable access are introduced by refactoring, the compiler's rules may require parameter signatures or function boundaries to be revised accordingly.

Immutable Borrow	Mutable Borrow
<pre> fn print_message(message: &str) { println!("{}", message); } fn main() { let greeting = ↪ String::from("Hello, ↪ world!"); print_message(&greeting); // ↪ Borrows greeting immutably println!("Still own greeting: ↪ {}", greeting); } </pre>	<pre> fn add_exclamation(message: &mut ↪ String) { message.push('!'); } fn main() { let mut greeting = ↪ String::from("Hello"); add_exclamation(&mut greeting); ↪ // Borrows greeting mutably println!("{}", greeting); ↪ // Prints "Hello!" } </pre>

Table 1: Immutable and mutable borrowing in Rust

3. **Lifetimes and their Management:** The Rust compiler tracks how long references live to guarantee memory safety. Any function that takes (or returns) references must define or infer lifetime parameters. When we write Rust code normally, such as the example below, the compiler is capable of inferring these lifetimes for the programmer.

```
1 fn echo(message: &str) -> &str { message }
```

When the programmer writes `echo(message: &str) -> &str`, the compiler is quietly inferring a generic lifetime for both the parameter and the return value. In effect, the compiler treats this as if the programmer wrote:

```
1 fn echo<'a>(message: &'a str) -> &'a str { message }
```

So, behind the scenes, Rust sees that `message` is a reference that lives at least as long as `'a`, and it infers that the output must share that same `'a` lifetime. The programmer doesn't need to write `<'a>` or annotate the references explicitly, because the language automatically applies its lifetime inference rules whenever possible.

If engineers later extract code into new functions—especially if multiple reference parameters or more complex borrowing patterns are involved—the compiler might no longer be able to figure out the relationships between references, forcing programmers to specify such `'a` parameters explicitly.

4. **Non-local Control Flow** Rust allows `return`, `break`, and `continue` inside code blocks. Extracting a fragment of code containing, for example, an early `return`, into a function changes the flow of control unless handled explicitly.

0.1.5 Adventure of a Lifetime: Automated Refactoring for Rust

A foundational contribution toward automated refactoring for Rust was made by Sergey et al. in *Adventure of a Lifetime* ?. While “Extract Method” refactoring is relatively straightforward in languages like Java, the challenges discussed in the previous section demonstrate why a similar algorithm for Rust is far from trivial. To tackle these unique challenges, *Adventure of a Lifetime* introduced a specialized refactoring framework tailored to Rust’s complexities.

Their approach addressed Rust-specific issues using multiple specialized passes. For example, one pass systematically encodes **return**, **break**, and **continue** statements into an auxiliary enum, allowing the caller function to reconstruct the original control flow by matching on variants such as `Ok(...)`, `Return(...)`, etc. Another pass utilizes constraint-based ownership analysis to precisely determine and assign the correct ownership levels (owned or borrowed references) to variables involved in extracted code fragments.

The systematic handling of these issues ensures correctness with respect to ownership, borrowing, and lifetimes, thereby overcoming the primary hurdles faced when naively transferring techniques from languages like Java directly into Rust. Figure 1 from *Adventure of a Lifetime* outlines the complete refactoring process.

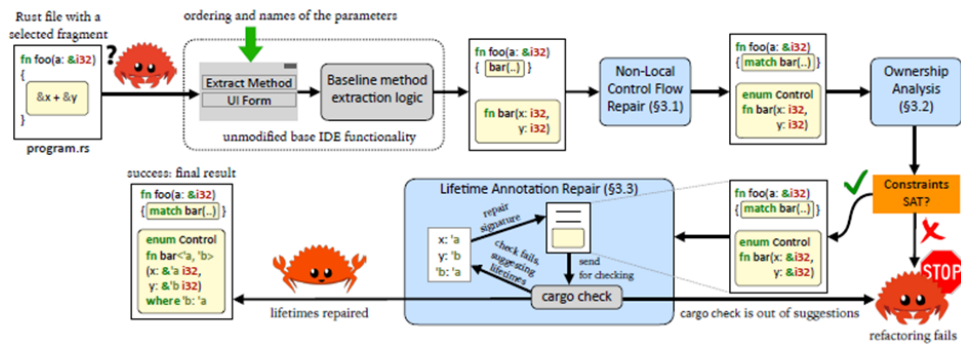


Figure 1: The refactoring process as outlined in *Adventure of a Lifetime*

The authors implemented this approach as *Rusty Extraction Maestro* (REM) on top of IntelliJ IDEA’s Rust plugin. Their empirical evaluation on multiple real-world Rust projects demonstrated that REM handles a wider range of extractions than existing IDE tools, including scenarios involving multiple borrows and nested references. Specifically, REM successfully managed 37 out of 40 tested cases, outperforming both IntelliJ IDEA alone and VSCode’s Rust Analyzer. The study also showed that REM-generated functions extracted quickly (on the order of one to two seconds), making it practical for everyday development. Most of the overhead arose from repeated calls to `cargo check`, although the delay remained acceptable for typical usage scenarios.

However, the evaluation also identified several limitations. REM struggles with code involving complex generics or elaborate trait bounds due to limited type inference capabilities in IntelliJ’s Rust support. Additionally, it lacks support for asynchronous programming constructs (`async/await`) and is unable to correctly extract functions that partially move struct fields. Macros beyond standard-library macros frequently disrupt refactoring attempts. Lastly, REM requires (and also assumes) a compilable project state, as it requires compilation checks to aid with its analysis. It often generates functions with overly verbose type signatures due to explicit lifetime and ownership annotations, despite the application a robust set of lifetime elision rules.

0.2 Problem Statement

0.2.1 Issues with Existing Automated Refactoring Tools

Rust’s unique ownership, borrowing, and lifetime system presents challenges that traditional refactoring tools don’t typically encounter. In languages like Java or Python, transformations such as “Extract Method” can be applied directly with minimal risk of breaking the code. However, Rust’s strict guarantees mean that naïvely lifting code into a new function often results in compilation errors or unintended changes in behavior due to ownership conflicts, borrowing rules, or lifetime mismatches. As a result, programmers often have to manually fix the refactored code—a process that can take significantly longer than rewriting it from scratch. While automated refactoring tools for Rust, such as RustRover and the IntelliJ Rust plugin, handle simple extractions well, they struggle with more complex constructs due to these inherent language constraints.

For instance, IntelliJ IDEA’s Rust plugin and Visual Studio Code’s Rust Analyzer both handle simpler scenarios—like extracting a small code block without complex borrows—but rely heavily on IntelliJ’s type inference or `rustc`. They often provide incomplete coverage for features such as `async/await`, generics with intricate trait bounds, or macros, leading the tools either to produce incorrect code or to refuse to recognize an extraction altogether. Moreover, these tools commonly struggle with non-local control flow, such as an early `return` from deep within a nested block.

By contrast, REM specifically targets these Rust-specific pain points. First, it applies a specialized ownership analysis algorithm that determines exactly how data should be passed—owned versus borrowed—in an extracted method. Second, it reifies non-local control flow operations like `return`, `break`, and `continue` by translating them into a custom enum and then patching the call-site logic. Finally, REM refines the resulting function’s lifetime annotations in a loop, using `cargo check` to repeatedly gather compiler feedback. This structured approach greatly expands the variety of Rust code patterns that can be safely refactored.

However, REM has several limitations. It heavily depends on the IntelliJ IDEA plugin for Rust, which is outdated and no longer supported. Additionally, it relies on numerous undocumented and rapidly changing `rustc` internals, leading to frequent build issues as third-party packages update and become incompatible. Another major drawback is that REM can only refactor code that compiles, as it relies on `cargo check`, which still requires successful compilation. Lastly, it passes multiple temporary files around on disk, making it incompatible with tools like `rust-analyzer`, which operate entirely in memory through the Language Server Protocol.

0.2.2 Need for Verification

An automated refactoring toolchain is only as good as its ability to preserve the program’s behavior. Rust refactoring poses particular challenges here due to its strict ownership and lifetime rules. A good refactoring must ensure that the transformed program maintains the exact observable behavior as before. Although comprehensive test coverage is the primary practical means of ensuring the preservation of functionality, tests are often slow and limited—they identify only *that* a problem exists, not precisely *where* it was introduced.

Verifying semantic preservation after transformations is notoriously challenging, particularly in Rust. Rust’s unique language features complicate semantic preservation proofs or automated verification techniques. Moreover, Rust code produced by REM may include constructs—such as enums for reified control flow or explicit lifetime annotations—that can appear unfamiliar or unnecessarily complex to human readers. While this complexity does not imply semantic changes, confirming correctness remains essential. Therefore, employing rigorous semantic validation (automated, partial, or manual) alongside refactoring is necessary to ensure transformations remain safe, accurate, and reliable.

0.3 Project Aims and Objectives

This project seeks to improve and extend the capabilities of the **Rusty Extraction Maestro** (REM) tool while making the refactoring process more accessible and verifiable. Specifically, the aims can be grouped into three key areas: enhancing REM’s standalone functionality, establishing a solid approach to verifying refactorings, and developing a proof-of-concept VSCode extension. A long-term ambition is to explore potential integration with `Rust Analyzer` if future architectural changes permit.

0.3.1 Improve and Extend REM

A principal goal is to remove REM’s existing dependencies on both `rustc` and IntelliJ IDEA, allowing the tool to function independently. By offering a standalone solution, REM would be simpler to incorporate into various workflows, free of the overhead or version constraints tied to external IDEs and compilers. This will be achieved through a single comprehensive CLI that links all stages of REM together.

To ensure broader adoption, another focus is to provide VSCode integration in the form of a user-friendly plugin. This plugin should install easily, require minimal setup, and offer automated `Extract Method` support for Rust directly in Visual Studio Code. In parallel, the project aims to expand REM’s capabilities to handle advanced Rust features, such as `async/await` constructs, macros, closures, `const` functions, and more intricate generics. Addressing these challenging language constructs will make REM suitable for a greater variety of real-world Rust codebases.

0.3.2 Verification of REM / Refactorings

Automated refactorings are only meaningful if they preserve the program’s original semantics. While typical “copy-and-paste” style transformations rarely alter program logic, REM’s deeper rewrites—which may include lifetime adjustments or reified control-flow enums—require stronger assurances that the resulting code truly behaves as intended. To achieve this, the project intends to develop a theoretical framework (or algorithm) for proving equivalence between the original and refactored code.

An emerging strategy involves functional translation—moving refactored Rust code into a simpler, more mathematically tractable language subset or an IR for proof purposes. Recent development to toolchains like CHARON? and AENEAS? can assist by allowing for a complete translation of a Rust program into a set of functional, proof oriented languages like F* and Coq. Though complete verification of all Rust code remains infeasible due to the language’s evolving semantics, partial or scenario-specific proofs can still bolster confidence in REM’s correctness. Demonstrating a handful of end-to-end examples—from extraction in Rust to verified equivalence in a proof assistant—will illustrate the viability of this approach. Moreover, the verification pipeline has the potential to be integrated into the main body of REM, providing the programmer with a guarantee that their code’s functionality hasn’t been affected or alternatively, the verification tool can be leveraged to execute a comprehensive series of case studies that further validate REM’s functionality.

0.3.3 Build a Proof-of-Concept Extension for VSCode

A strong emphasis is on producing a proof-of-concept **VSCode** extension that showcases **REM** in action within one of the most commonly used Rust development environments. This extension should seamlessly integrate into a developer's normal workflow, from code editing to debugging, without burdensome setup. By offering an easy-to-install plugin, the refactoring workflow becomes more accessible, and potential refinements can be validated in real-world coding sessions.

A core demonstration will involve extracting method fragments in a variety of Rust files and showing how the extension generates well-typed, behavior-preserving code. The key detail here is that **REM** relies on its new, standalone core, instead of a platform or program specific implementation.

0.3.4 Potential for the Long Term Goal of Merging **REM** with **Rust Analyzer**

Finally, looking beyond the immediate development schedule, an intriguing possibility is merging **REM**'s functionality with **Rust Analyzer**, the official language server for Rust. **Rust Analyzer** already converts source code into an intermediate representation to power syntax highlighting, code completion, and basic refactorings. If **REM** can leverage that IR, it might inject advanced lifetime analysis and partial verification steps directly into the analyzer pipeline.

However, **Rust Analyzer** is deliberately kept lightweight and incremental, while **REM** aspires to be a comprehensive solution that tackles complex rewriting tasks. Hence, if integration proves overly burdensome, a compromise could be to develop a stripped-down subset of **REM** for merging. This long-term exploration depends on how **Rust Analyzer** evolves and whether the Rust community adopts more sophisticated refactoring workflows within a single unified environment. The alternative solution of a standalone **VSCode** extension designed to be used alongside **Rust Analyzer** is also seen as acceptable in this case.

Literature Review

1.1 Introduction

Rust is a modern systems programming language that enforces memory safety through a strict ownership and borrowing discipline. This discipline, enforced by Rust’s borrow checker, ensures that well-typed Rust programs are free from data races and dangling pointers without requiring a GC ?. However, Rust’s unique type system — centered on ownership and lifetimes — poses new challenges for code transformation tools and formal reasoning. This literature review surveys academic work on automated refactoring techniques for Rust, approaches to verification of Rust programs (including formal methods and model checking), and efforts to formalize Rust’s semantics. We compare Rust’s refactoring and verification tooling to those of other languages, and we discuss integration with development environments. The goal is to highlight the state of the art in making Rust programs easier to evolve and prove correct, while pinpointing structural refinements for a clearer organization of this body of work.

1.2 Overview of Refactoring

As previously mentioned, refactoring is the process of restructuring existing code to improve its internal structure without changing its external behavior. It is a key practice for enhancing code maintainability and reducing technical debt in software projects ?. Modern IDEs, such as Eclipse and IntelliJ IDEA include a wide array of automated refactorings, providing developers with quick, semantics-preserving code transformations ?. In theory, these automated refactorings offer a safe way to restructure code (backed by well-defined preconditions to preserve behaviour ?) and should be widely used in practice.

However, empirical studies show a disparity between tool support and actual usage. Developers often remain reluctant to use automated refactoring features, preferring manual code changes ?. For example, even with “Rename” refactoring available at the click of a button, many programmers still rely on find-and-replace or other manual methods ?.

The survey also noted that many developers expressed a desire to better understand their IDE’s refactoring capabilities. This suggests that while the tools exist, there is a gap in awareness and useability of the tooling available. Additionally, as evidenced by the lack of adoption of even “Rename” refactorings, challenges still remain in getting developers to trust automated tools to perform complex manipulations.

Research into automated refactoring aims to bridge this gap by making tools more reliable and intelligent. For instance, Tip (2007) demonstrated that refactorings can be modeled with type constraints to systematically explore alternative valid program structures ?. Such approaches use static analyses to ensure that transformations maintain type correctness and behavioral equivalence, giving developers confidence in the automation. Nevertheless, the human factor (usability, understandability of tool behavior) remains critical.

1.3 Rust’s Ownership Model and Formal Semantics

Rust’s ownership model introduces compile-time enforced rules for aliasing and memory lifetime. Each value in Rust has a single owning scope, and references (borrows) must obey strict rules that prevent concurrent mutation and use-after-free ?,?. While this model provides strong safety guarantees, it complicates formal semantics and tool development. Early efforts to rigorously define Rust’s semantics culminated in *RustBelt*, which provided the first machine-checked safety proof for a realistic subset of Rust’s type system ?. *RustBelt* demonstrated that safe Rust code (even when using unsafe internals of standard libraries) is memory and thread safe, by modeling Rust in the Coq proof assistant ?. This foundational work established confidence in Rust’s core design and set the stage for verifying more complex properties. Subsequent research has extended Rust’s semantic foundation to cover aliasing in unsafe code (e.g. Stacked Borrows and Tree Borrows models for pointer aliasing) ?, and to connect Rust’s high-level rules with low-level memory models. These formal models provide a basis on which verification tools and refactoring tools can reason about Rust code behavior. In summary, Rust’s enforced aliasing discipline, while posing challenges, has inspired a rich line of work in formal semantics that underpins safe refactoring and verification.

1.4 Automated Refactoring in Rust vs Other Languages

Well established languages like Java and C# benefit from decades of IDE support for refactorings such as *Rename*, *Extract Method* and *Extract Module / Variable*. Rust, being newer and with far more complex compile-time rules, has lagged behind in automated refactoring support [1]. The ubiquitous *Extract Method* is “widely used in all major IDEs” for other languages, but implementing it for Rust is surprisingly non-trivial due to Rust’s ownership and lifetime constraints [2]. In contrast to a language like Java, where extracting a function involves mostly syntactic rearrangement, Rust’s refactoring tools must also infer where to borrow or clone data and how to introduce lifetime parameters to satisfy the borrow checker [3, 4]. Early comparisons noted that Rust’s lack of reflection and unstable compiler APIs made it harder to build refactoring tools as robust as those for Java / C#. Even in the two years since REM was released, compiler APIs have changed so much that it proved impossible to compile the original project without significant rewrites. On the positive side, Rust’s compiler errors can guide manual refactoring by pinpointing violations of borrowing rules, giving programmers immediate feedback; this means that if a refactoring compiles in Rust, it is likely behaviour-preserving [5]. Nonetheless, the consensus in both academia and industry is that first-class refactoring tools are needed to manage large Rust codebases with the same ease developers expect in other environments [6, 7].

1.5 Automated Refactoring Techniques for Rust

1.5.1 Renaming and Simple Refactorings

One of the first efforts to build a Rust refactoring tool was by G. Sam et al. (2017), who created a proof-of-concept refactoring framework utilizing the Rust compiler’s internal APIs. The team partnered with Mozilla Research to be among the first to implement the Rust specific refactorings of *Lifetime Elision* and *Lifetime Reification*. This allowed their program to introduce explicit lifetime parameters in instances where the compiler was implicitly inferring them - a refactoring that was brand new to Rust. The challenges encountered illustrated how Rust’s stricter scoping and shadowing rules required careful handling of name conflicts during renaming (e.g. avoiding situations where renaming a variable could unintentionally shadow another) [8]. Additionally, the 2017 study concluded that many refactorings are possible with Rust’s compiler infrastructure, but ensuring *behavioural preservation* (especially around ownership transfers) requires additional static analyses not needed in languages without Rust’s constraints. Their work on *Lifetime Elision* has since formed part of the REM toolchain, where the rich feedback from the

Rust compiler is leveraged to ensure that the final transformation is both valid Rust and as legible as possible ?, ?.

1.5.2 Extract Method (REM)

A major advance in Rust refactoring came with the Rust Extract Maestro (REM), with the theoretical background provided by the work of Costea et al. (2023) in *Adventure of a Lifetime: Extract Method Refactoring for Rust*. Thy et al. (2023) then provided a practical implementation of REM in *Borrowing Without Sorrowing*. REM is a tool that tackled notoriously complex Extract Method refactoring for Rust. REM decomposes the extract-method process into a sequence of transformations, each addressing a specific aspect of Rust’s type and borrow rules ?. The approach begins with a “naïve” function hoisting (moving a block of code into a new function), then applies a series of *automated program repairs* to make the code compile correctly. This process is guided by *oracles* - specialised analyses that resolve issues the compiler would flag. One such oracle is an intra-procedural **ownership analysis** which infers whether each value moved to the new function should be passed by value, by shared reference (using `&T`), or by mutable reference (using `&mut T`). Appendix 1 (6.1) provides a detailed example of what this looks like in practice.

Another oracle leverages Rust’s own compiler (`rustc`) as a sub-procedure: REM invokes Rust’s lifetime checker to identify lifetime inconsistencies, then introduces appropriate lifetime parameters to the extracted function until the borrow checker is satisfied. By iteratively repairing borrow errors in the extracted code, REM ensures the transformation preserves semantics and yields a well-typed program. The REM tool, implemented as an extension to the IntelliJ Rust plugin, was shown to handle complex extractions that involve non-local control flow (e.g. `return` or `break` inside the extracted fragment) and borrowing across function boundaries. In an evaluation on real-world Rust projects, REM could successfully perform extractions that developers had done manually (including cases with nested lifetimes), outperforming existing tools in the scope of code it can handle. Their work demonstrated that seemingly “unsafe” transformations (like cutting a function in half) can be automated in Rust by coupling transformations with verification from the Rust compiler’s own checks.

1.5.3 Automated Fixes for Ownership Errors

An alternative angle on Rust refactoring is automatically *fixing* code that violates ownership/borrow rules - essentially a refactoring from “non-compiling” to “compiling” code without changing behavior. **Rust-Lancet** (Yang et al., ICSE 2024) is a tool aimed at *automated ownership-rule-violation fixing with behaviour preservation*. Given a Rust

source file that fails to compile due to borrow checker errors, Rust-Lancet analyzes the Abstract Syntax Tree (AST) and applies code transformations to eliminate the error while preserving the program’s semantics ?. Under the hood, it uses Rust parsing libraries (`syn` and `quote`) to manipulate the AST and generate patched code. In a similar manner to REM, Rust-Lancet employs an iterative repair loop: it first applies a fix (for example, inserting a `clone()` on a value to appease ownership rules), then rechecks the compiler errors and repeats if further fixes are needed. The tool includes correctness checks to avoid over-fixing; a *behaviour preservation* module validates that the inserted changes do not alter program outputs for typical cases. In an evaluation on over 150 real cases of ownership rule violations (gathered from Rust’s test suite and prior research), Rust-Lancet was able to completely fix a large fraction of them with zero false positives (i.e. it never introduced an incorrect fix). It even outperformed suggestions given by the Rust compiler and baseline techniques using large language models, especially in scenarios requiring multiple coordinated edits. While Rust-Lancet is positioned as a bug-fixing tool, it overlaps with refactoring by automatically rewriting code in a semantics-preserving way. Its success highlights how Rust’s rigorous rules can be leveraged - the tool knows a patch is correct when the compiler’s checks pass and its behavior preservation validation succeeds ?.

1.5.4 Comparison and Summary

Overall, automated refactoring in Rust is characterized by taking full advantage of the rich and sophisticated output of the Rust compiler. Tools like REM explicitly invoke the borrow checker and type checker as so called oracles ?, whilst others, like Rust-Lancet, effectively automate a “fix until compile” loop with guarantees of behaviour preservation. This is in contrast to many refactoring tools for garbage-collected languages, which operate mostly on the AST or intermediate representation without needing to consider memory lifetimes. The research so far indicates that, with suitable abstractions (e.g. ownership analysis algorithms and lifetime inference strategies), even complex Rust refactorings can be automated soundly. As more of these techniques mature, we expect Rust’s refactoring support to begin matching the breadth available for languages like C++ and Java, albeit with different underlying algorithms tailored to Rust’s unique semantics.

1.6 Verification and Formal Methods for Rust

Rust’s appeal for building reliable software has spurred interest in applying formal verification techniques to Rust programs. Traditional *deductive verification* and *model checking* must be adapted to account for Rust’s ownership, lifetimes, and possibly unsafe code. In

this section, we review key verification tools and techniques, and the formal semantics advances enabling them.

1.6.1 RustBelt and Type-System Soundness

Jung et al. (2018) introduced **RustBelt**, a foundational framework for reasoning about Rust’s type system and memory safety. In their landmark paper, *RustBelt: Securing the Foundations of the Rust Programming Language*, they proved the soundness of Rust’s core type system. By embedding a Rust-like language (λ Rust) in the Iris framework (a higher-order concurrent separation logic), RustBelt verified that well-typed safe Rust code cannot exhibit undefined behaviour[?]. This was the first independent verification of Rust’s claims about memory and thread safety. Additionally, it provided a foundation on which other verification efforts could build: one can assume that safe code is memory-safe, allowing them to focus verification effort either on higher-level functional correctness or the remaining pitfalls of unsafe code.

1.6.2 Verification via Functional Translation (CHARON and AENEAS)

One prominent approach to verifying Rust programs is to reduce the problem to verifying *pure functional code*. **AENEAS** (Ho & Potzenko, ICFP 2022) is a toolchain that translates a substantial subset of Rust into a purely functional language, suitable for input to existing proof assistants or model checkers[?]. The key insight of Aeneas is that for many Rust programs (those not using interior mutability or unrestricted unsafe code), explicit memory reasoning can be eliminated. Aeneas introduces an intermediate representation called *Low-Level Borrow Calculus* (LLBC), inspired by Rust’s mid-level IR (MIR), which makes the ownership and borrowing structure of a program explicit. Crucially, LLBC’s semantics is *ownership-centric*: instead of a heap, it uses an abstraction of *loans* and *borrow*s to track aliasing, meaning individual values carry information about what they borrow.

Ho and Potzenko went on to define a set of pure, value-based semantics for LLBC - one with no pointer addresses - thereby “capturing the essence of the borrow mechanism” in a mathematical form[?]. They then translate LLBC into a pure lambda calculus¹, effectively converting Rust code into a functional program with the same behavior. The result is that one can verify properties of the Rust program by verifying the translated functional program using a theorem prover of choice. For example, one might translate Rust to Coq or Fstar and prove postconditions about the output, with Aeneas ensuring

¹See this excellent introduction to the Lambda Calculus for more information:
<https://plato.stanford.edu/entries/lambda-calculus/>

those properties carry back to the original Rust code.

To tackle tricky issues like termination of borrows across function boundaries, AENEAS introduced a notion of *backwards functions* to conservatively approximate the “borrow graph” of a program and ensure the translated code remains termination-friendly. The approach was shown to significantly reduce the burden on proof engineers, since they no longer needed to reason about low-level memory safety - Rust’s type system has already done the heavy lifting for them.

From there, Ho et al., (ICFP 2024) presented *Sound Borrow-Checking for Rust via Symbolic Semantics* ?, with the goal of addressing two open questions from AENEAS’s methodology. First, they established a link between Aeneas’s LLBC and a traditional heap-based semantics, by proving that LLBC’s admittedly unusual modelling is faithful to a standard memory model, ensuring that verifying programs in LLBC is not “proving the wrong thing”. Second, they introduced a set of formal *symbolic semantics* for Rust that act as a sound abstraction of concrete execution, and prove that these symbolic semantics are capable of correctly approximating Rust’s behavior. In essence, they show that if a program is accepted by their symbolic interpreter (which enforces borrow rules), then there exists a corresponding concrete execution in a real heap that is memory-safe. This result formally validates the idea that the borrow checker (or a tool mimicking it) prevents all illicit behaviors.

A very important outcome of this work is a proof that the **symbolic interpreter can server as a verified borrow-checker** for LLBC programs. This provides increased confidence that tools like Aeneas (which rely on symbolic execution as an intermediate) are built on solid theoretical ground, and it paves the way for future verified compilers or static analyzers for Rust that incorporate a proven-correct borrow checker. The broader significance is that Rust’s complex lifetime rules are now accompanied by a machine-checked proof of their soundness, something rare for industrial language features.

1.6.3 Static Analysers and Model Checkers

Static Analysis

Beyond full formal verification, several tools bring lightweight verification or bug-finding to Rust. **PRUSTI** (Astrauskas et al., OOPSLA 2019) is a static *deductive verifier* that uses Rust’s types to simplify program annotation and verification, ?. Prusti translates Rust code into the Viper verification language, using Rust’s lifetime and ownership information to implicitly encode permissions for verification ?. This means a developer can write specifications (pre/post-conditions, invariants) for Rust functions, and Prusti will prove them, assuming no undefined behavior in safe code. The advantage of leveraging Rust’s type system is evident - e.g., Prusti knows that two mutable references cannot

alias, which corresponds to a simple permission separation in the verification condition. Over the past few years, Prusti has grown to handle significant subsets of Rust, and it even provides a Visual Studio Code extension for interactive use ?.

In a similar vein, *RefinedRust: A Type System for High-Assurance Verification of Rust Programs* (Gäher et al., 2024) introduced a refinement type system for Rust that can verify both safe and certain unsafe code by embedding checking into the type system ?. RefinedRust’s types allow expressing rich safety and correctness properties, and a prototype tool translates annotated Rust into a Coq model for verification. This line of work brings Rust closer to languages like Dafny or Liquid Haskell in terms of having a path to verification integrated with the language’s types.

Model Checking

The most notable tool for *model checking* is **Kani**, a bounded model checker for Rust developed by AWS. Kani translates Rust programs (including those with `unsafe` code) into verification conditions that are checked by a SAT/SMT solver (built on the C Bounded Model Checker (CBMC) backend) ?. What sets Kani apart is that it operates on Rust’s MIR (Mid-level IR), leveraging information about traits and generics directly rather than reducing to LLVM bitcode. By working at the MIR level, Kani retains high-level semantics like trait object behavior, which allowed the team to verify complex properties involving dynamic trait objects (Rust’s form of dynamic dispatch). VanHattum et al. (ICSE 2023) describe how Kani handles dynamic trait dispatch soundly, making it the first model checker to support the full range of Rust trait features ?. In experiments, they found that using Rust-specific knowledge (MIR and trait info) improved verification performance by $5\times$ - $15\times$ compared to an equivalent analysis at the LLVM level. Kani is being used to verify critical systems components (for example, AWS has applied it to verify parts of their cryptographic libraries and virtualization code), focusing on memory safety and user-provided assertions in performance-critical Rust code. Other dynamic analyzers like MIRI (an interpreter that detects undefined behavior in Rust by executing code in a defined environment) complement these verification tools by checking one execution at a time but with a semantic model that catches unsafe violations; however, MIRI is more a testing aide than a formal verifier.

1.6.4 Summary

We have shown that the Rust verification ecosystem is incredibly rich and includes: foundational proofs of the language’s safety (RustBelt, symbolic semantics), translation-based verifiers (Aeneas, Prusti, RefinedRust) that build on the guarantees of Rust’s types, and model checkers (Kani) that leverage Rust’s own compiler internals for efficient exhaustive exploration. Each of these harnesses Rust’s key feature—its rich type system—as an

asset to simplify or optimize verification. This contrasts with verification for languages like C/C++, where the lack of a guaranteed memory safety baseline means much effort is spent just to mitigate wild pointers or data races. In Rust’s case, the baseline guarantees free the verifier to focus on higher-level properties or the remaining unsafe corners of the language. The research surveyed here shows a trend of *integrating verification with the language’s design*: from using the borrow checker in a proof assistant, to encoding Rust lifetimes as logical predicates, to running model checkers on MIR. This tight integration yields verifiers that are both sound (no false negatives on safety) and relatively precise, making formal assurance of Rust software increasingly attainable.

1.7 Integration with IDEs and Language Servers

The ultimate goal for refactoring and verification tools is seamless integration into developers’ workflows. For Rust, integration has meant working closely with the compiler or the emerging Language Server Protocol (LSP). However, `rustc` was not originally designed as a reusable library for IDE features, which led to the creation of separate projects to support code intelligence and analysis. The first generation Rust Language Server (RLS) attempted to use the compiler internals to power IDE features, but it struggled with performance and completeness as Rust evolved. Today, the community relies on **rust-analyzer**, a from-scratch implementation of a Rust language server that emphasizes speed and modularity ?. Rust-analyzer constructs its own parse tree and does its own type inference, enabling on-the-fly responses to editor queries without running the full compiler on every change. This architecture is conducive to implementing refactorings as lightweight code transformations called “assists.” In fact, rust-analyzer already supports basic refactorings like *Rename* and *Extract Variable*, and *Extract Method* as editor assists. These operate on rust-analyzer’s internal syntax tree and use its knowledge of ownership and types to ensure the edits are valid. For example, the Extract Function assist initially had limitations with Rust’s generics and lifetimes - extracting code from a generic function would produce errors because the assist didn’t propagate required type parameters or lifetime parameters to the new function. However, even with substantial development by the community, certain assists still fall significantly short of the capabilities of tools developed specifically for purpose, such as REM.

Integration into the popular IntelliJ IDEA (via the Rust plugin) has also seen some development. The REM tool was originally built atop the IntelliJ plugin, indicating that industrial IDEs can serve as a proving ground for research prototypes ?. In principle, once validated, these prototypes can be upstreamed, bringing advanced refactorings to everyday developers. A challenge here is maintaining the tools: Rust’s syntax and compiler behavior evolve, so refactoring algorithms must be kept in sync. One way this is being

addressed is by using Rust’s stable analysis APIs where possible. For instance, tools like Rust-Lancet chose to use the community-maintained `syn` crate for parsing so that they aren’t tied to internal compiler data structures that change nightly [?]. Similarly, projects to support large-scale refactoring use cases (such as Facebook’s code migration tools or the Rust-to-Rust transformations needed in the Rust 2018 edition uplift) relied on either compiler suggestions or external scripting. The Rust compiler team has introduced machine-applicable suggestions in compiler error messages, which tools like `rustfix` can apply automatically. This essentially allows the compiler to assist in simple refactors (e.g. renaming a deprecated syntax, adding a missing lifetime specifier) by emitting a structured suggestion. Whist not as general as a true refactoring engine, this has helped with batch changes across large codebases.

1.7.1 Compiler Internals for Ownership and Borrowing

It’s worth noting how Rust’s compiler architecture influences refactoring and verification tools. Rustc itself goes through a series of analyses: it parses into an AST (with desugaring), performs name resolution and type checking (including lifetime inference), and then produces MIR where borrow checking and other analyses occur. Because of this, tools integrated at various stages have different views of the program. For example, a refactoring tool working on an AST (like rust-analyzer’s assists or Rust-Lancet) must re-run portions of the compiler (or mimic them) to ensure that after transformation the code still passes borrow check. Meanwhile, a verification tool like Kani that operates on MIR has the advantage of piggybacking on the actual borrow-checker’s results - it can trust that MIR is already free of borrow rule violations and focus on exploring execution paths [?]. As research progresses, we may see more refactoring tools move to operate on MIR or use the borrow checker in the loop, because MIR provides a simpler, desugared view of the code with explicit borrow regions (e.g. the Polonius project, a next-generation borrow checker, could potentially expose an API for tools to query borrow relationships). For now, practical refactoring uses higher-level representations (AST plus perhaps some compiler query for type info) because it’s easier to map edits back to source code.

1.8 Related Work Beyond Rust: Refactoring in Other Languages and Paradigms

The challenges and advances in refactoring outlined throughout this review are not unique to Rust. Many other languages have inspired research and into automated refactoring and its many pitfalls. The easiest case study is Java, being one of the earliest targets of refactoring tools, and it has a rich literature. Beyond the generics case study mentioned earlier ([?]), researchers have looked at refactoring concurrent Java code, refactoring to use new

language features, and large-scale restructuring of legacy systems. Zhang et al. (2024) recently proposed an automated refactoring approach for Java’s asynchronous programming using the `CompletableFuture` API. Their tool, **ReFuture**, integrates static analysis (visitor patterns, alias analysis, etc.) to identify where an older async construct (like using raw threads or Futures) can be transformed into a modern `CompletableFuture` chain. Impressively, **ReFuture** was evaluated on 9 large Java projects, including *Hadoop* and *ActiveMQ*, and managed to refactor 639 out of 813 eligible code segments automatically, without introducing errors. This demonstrates that automated refactoring can manage to tackle performance or paradigm-migration tasks, not just cosmetic and maintenance changes.

In the mobile app domain, Lin et al. (2015) addressed refactoring for Android’s asynchronous constructs. Android apps often mis-use the `AsyncTask` class, leading to memory leaks or lost UI updates. Lin and his team developed **ASYNCDROID**, a tool to refactor improper `AsyncTask` usage into a sover structures like `IntentService`. Their formative study found that about 45% of the `AsyncTask` occurrences in real apps could be automatically refactored by **ASYNCDROID**’s rules, and an additional 10% could be handled with minor manual tweaks, resulting in more than half of problematic cases being fixed largely automatically. The tool was implemented as an Eclipse plugin and offered typical refactoring tool conveniences including a preview of the changes and a rollback option. Importantly, when the authors submitted some of these refactoring patches to open-source projects, the maintainers accepted them, validating that the transformations were indeed considered improvements with no behavioral regressions. This work illustrates how automated refactoring can be applied in a specialized context (mobile asynchronous UI code) by encoding domain-specific knowledge as transformation rules and preconditions.

Other notable refactoring research outside of Rust includes work on functional languages and dynamically-typed languages. In Haskell, for example, the **HaRe** tool (Li et al., 2005) brought refactoring support to a lazy, purely functional language, requiring careful handling of Haskell’s scoping, type classes, and purity concerns. In dynamic languages like JavaScript and Python, refactoring is complicated by the lack of static types; researchers have explored heuristic and runtime-analysis-based approaches to perform refactorings safely (e.g., refactoring Python’s modules or JavaScript’s callbacks to promises). While we will refrain from delving into specific studies for those here, it is worth noting that each language tends to spawn its own refactoring research to address unique features.

Finally, it is important that we note that refacotrning is very closely intertwined with other software engineering tasks like code smell detection, program repair and modernisation. Refactoring tools are sometimes used as building blocks for automated program repair, such as applying behavior-preserving transformations to enable a later bug-fixing change,

with one such case even mentioned in *Adventure of a Lifetime* ?. Tools like Facebook’s JScodeshift or Python’s Bowler allow developers to script custom refactorings (essentially semi-automated transformations), which blurs the line between a “human-driven” refactoring and a tool-driven migration. The increasing adoption of language servers means refactoring capabilities can be made available in a wide range of editors, not just heavyweight IDEs ?. This has been a positive development for languages like Rust and Go, where a language server can provide refactoring assists consistently across many development environments.

1.9 Conclusion

The landscape of Rust refactoring and verification has rapidly evolved, marrying the Rust compiler’s strengths with novel algorithms from the research community. Automated refactoring tools for Rust have progressed from simple renaming utilities to sophisticated systems that integrate program repair and static analysis (REM and Rust-Lancet) to handle ownership and lifetime challenges ?, ?. Verification tools leverage Rust’s unique guarantees to simplify proofs – whether by translating to pure functional models (Aeneas) ?, by encoding lifetime reasoning into logical frameworks (RustBelt, RefinedRust) ?, or by exploring executions with model checkers (Kani) that understand Rust-specific features ?. A unifying theme is that Rust’s ownership semantics, once seen purely as a hurdle, have become an enabler for both refactoring and verification: they localize the reasoning about memory, allowing tools to either confidently transform code or assert properties with the guarantee that certain classes of bugs are already ruled out.

In conclusion, Rust’s combination of strong static checks and low-level performance has fostered a unique blend of research. The advances in refactoring mean that developers can expect increasing assistance in restructuring Rust code without fear of breaking it. Meanwhile, the maturation of formal methods for Rust suggests that full verification of critical Rust code is becoming practical, supported by tools that integrate with Rust’s development ecosystem. By structuring the literature along clear axes (refactoring vs. verification, formal foundations vs. practical tools, compiler phases involved, safe vs. unsafe scope), future reviews can further improve clarity and coverage. The works cited here collectively show that Rust’s initially steep learning curve for both humans and tools can be mitigated by clever automation and rigorous semantics – turning Rust’s famed safety guarantees from a burden into a powerful ally for program analysis and transformation.

Expanding the Capabilities of REM

Automated refactoring in Rust has long been constrained by the languages strict ownership, borrowing and lifetime rules. Early prototypes of the Rusty Extraction Maestro demonstrated that *Extract Method* refactorings could be made effective, but only within relatively narrow limits. Complex language features, such as generic types and asynchronous code, were not supported. Moreover, the entire toolchain was reliant on a fragile integration with IntelliJ, which quickly became outdated. As a result, REM remained an academic proof-of-concept rather than a tool developers could realistically adopt.

This chapter details how REM has been transformed from a fragile prototype into a robust and flexible system. We introduce support for two language features that were previously out of reach - generic code and asynchronous functions - both of which are pervasive in real-world Rust codebases. We then describe the design of a standalone command-line interface (CLI), eliminating dependence on third-party IDEs and making REM portable across development environments. Finally, we present a prototype Visual Studio Code (VSCode) extension that integrates seamlessly into one of the most widely used Rust workflows. Together these advances greatly expand the scope and usability of REM, demonstrating that advanced refactoring can be brought to everyday developers without sacrificing continued improvement through research. This chapter thus sets the stage by showing how REM's capabilities now align with the realities of modern Rust programming.

2.1 Refactoring Generics

2.2 Supporting Asynchronous Code

2.3 Standalone, Editor Agnostic CLI

2.4 Integration with IDEs

2.5 Summary

Automatic Verification of Refactored Code

In the strictest sense, refactoring is defined as a behavior-preserving code transformation. In practice, refactoring is only valuable if it preserves the original program’s behaviour. In Rust, where ownership and lifetimes enforce strict invariants, the mere fact that refactored code compiles does not guarantee equivalence. Additionally, because REM performs complex, compiler-guided repairs as part of its extraction process [?], there is an increased risk of introducing subtle changes in program behaviour. Subtle shifts in aliasing or lifetime structure can produce programs that pass the compiler yet diverge semantically from their originals. For high-assurance domains, this risk is unacceptable: automated tools must not only generate compiling code but also provide evidence that transformations are correct.

This chapter introduces a novel verification pipeline that extends REM with automatic, annotation-free proofs of equivalence between original and refactored code. Our approach combines existing formal methods toolchains - CHARON ¹, which translates Rust into an ownership-explicit intermediate form, and AENEAS ², which then generates Coq code equivalent to the original Rust source through a pure λ -calculus based intermediary. This Coq code can then be used to formally prove that the refactored code is equivalent to the original. The result is an end-to-end system in which the Extract \rightarrow Repair cycle is followed by a Verify phase, discharging proofs in seconds with no additional burden on the developer. By embedding verification directly into the refactoring process, we move beyond compilation success to true semantic assurance, bridging the gap between theory and practical developer tools.

¹More information accessible from their GitHub, <https://github.com/AeneasVerif/charon>

²More information accessible from their GitHub, <https://github.com/AeneasVerif/aeneas>

3.1 Motivation: Why Verification is Essential

Despite decades of work on automated refactoring, developers continue to show reluctance to adopt such tools in practice. Empirical studies indicate that even simple, “safe” refactorings such as renaming are frequently performed manually, as programmers prefer to trust their own edits over opaque tool behaviour ?. This gap stems from usability and, more fundamentally, from trust: if a refactoring tool cannot convincingly guarantee that semantics are preserved, developers are unwilling to rely on it for non-trivial code transformations.

The challenge is particularly acute in Rust. Compared to languages like Java or C#, where refactorings operate largely at the syntactic or type level, Rust’s ownership and lifetime system makes every transformation potentially semantic. Prior work shows that Extract Method in Rust is not merely “cut and paste” but requires compiler-guided repairs such as introducing lifetime parameters, reifying control flow, and inferring ownership modes ?. Crucially, the final step to the original REM algorithm involves a series of non-trivial lifetime adjustments, guided by `rustc`, to ensure that the extracted code satisfies the borrow checker. These repairs enable REM to succeed where naïve tools fail, but they also heighten the risk of unintended behavioural changes. A transformation that compiles may nonetheless shift aliasing patterns, extend or shorten lifetimes, or alter when resources are released — all of which are semantically observable in Rust.

Traditional quality controls such as testing do not fill this gap. Tests, even when comprehensive, explore only a small subset of possible inputs and interleavings. In safety or performance critical domains, developers cannot afford to accept the residual risk that untested paths diverge after a refactoring. Verification, by contrast, offers a principled way to resolve the trust deficit: it can establish - on a program by program basis - that for all inputs respecting the refactored function’s preconditioned, that the refactored program is equivalent to the original. By integrating verification directly into REM’s pipeline, we address not only the technical challenge of making complex extractions compile, but also the socio-technical challenge of making developers confident that such transformations are behaviour-preserving.

3.2 Overview of the Verification Pipeline

The verification pipeline augments REM’s extract-and-repair workflow with an automatic equivalence check. Its design is guided by two principles: (1) verification must require no additional annotations or input from the developer, and (2) it must complete quickly enough to fit into interactive IDE workflows. To achieve this, the pipeline translates both the original program P and the refactored program P' into progressively more structured semantic representations, culminating in a machine-checked proof of observational equivalence.

At a high level, the pipeline consists of five stages: **Extraction**, **REM repairs**, **CHARON translation**, **AENEAS functionalisation**, and **Coq verification**. Each stage builds on the previous one, gradually exposing the semantics of the refactored code while preserving a clear correspondence to the original. Figure 3.1 provides an overview of the entire process.

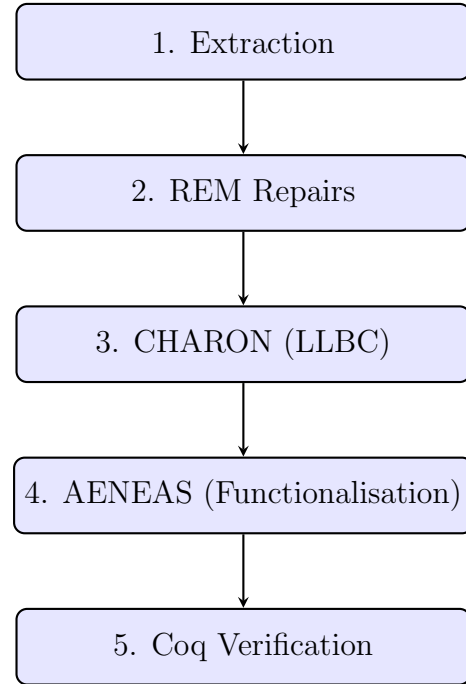


Figure 3.1: Overview of the verification pipeline: from source code extraction to proof of observational equivalence in Coq.

3.2.1 Extraction

At the first stage, the user identifies a region of code to extract into a new function. This can be done through REM’s CLI, however, the VSCode extension provides a much more user-friendly interface. The user highlights the code to be extracted and invokes the extraction command. This step mirrors the mechanics of traditional Extract Method refactorings: a block is lifted out of its enclosing function, and a corresponding call is inserted in its place. At this point the transformation is syntactic and only sometimes results in valid Rust. The below shows both cases:

- The first is a simple extraction that requires no further repairs. Strict typing makes this a trivial case, and REM would not be called.
- The second is a more complex extraction that requires REM’s repair phase to ensure

the refactored code compiles. It is adapted from a real commit in `GitOxide`³, and was also used in “Borrowing Without Sorrowing”⁴. Note how the extraction algorithm has failed to find the lifetime parameter 'a, and added an unnecessary borrow of `SectionID`. In this instance, REM would be called to fix the (non-compiling) code.

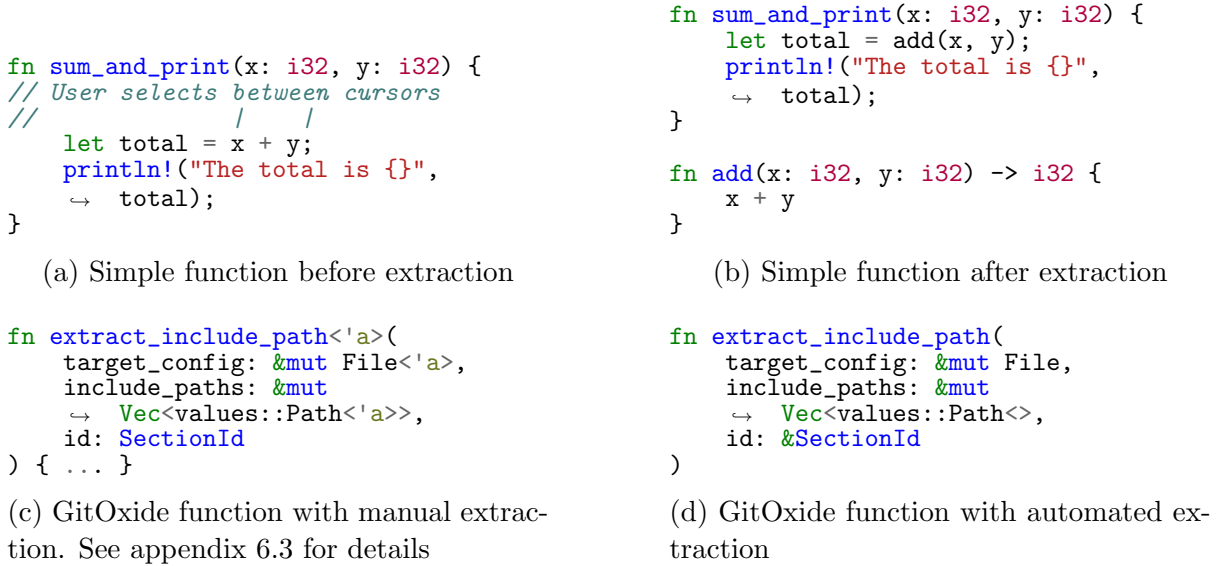


Figure 3.2: Two extraction examples: a trivial case (top) and a complex, failed case from real-world code (bottom).

3.2.2 REM

As seen in the above examples, a naïve extraction is often not going to compile. Free variables, non-local control flow and lifetimes make the extracted fragment ill-typed. the purpose of REM is to repair these issues automatically, producing a valid Rust program P' while preserving the intent of the original code. At a high level, REM applies three classes of fixes:

- (i) Reifying non-local control flow into a small result enum.
- (ii) Inferring ownership for free variables (by value, `&T`, or `&mut T`) and,
- (iii) Adjusting lifetime parameters in an iterative manner, guided by `rustc` until the borrow checker is satisfied.

These steps transform a purely syntactic extraction into a valid program P' that is much closer to what a developer would have written by hand.

³<https://github.com/GitoxideLabs/gitoxide/commit/c0786717c4979810002365a68d31abbf21d90f2d>

Example: Simple Case

In the toy example, extraction is trivial. REM observes that both `x` and `y` can be moved by value into the new function, so the inferred signature is simply:

```
1 fn add(x: i32, y: i32) -> i32
```

No further repairs are needed. This illustrates the best case: REM’s analysis immediately determines that the signature is valid Rust and human readable.

Example: Real World Case

By contrast, in the `extract_include_path` example (3.2d), naïve extraction fails. The plugin incorrectly infers that the loop variable `id` must be borrowed (`&SectionID`), and entirely misses the lifetime parameter `<'a>`, required to properly scope our mutable references to `File` and `Vec<path>` throughout the entire function.

REM corrects this in several ways:

1. It detects that `id` is consumed at the call site, so ownership can be transferred rather than borrowed.
2. It generalises lifetime parameters across both `File<'a>` and `Vec<Path<'a>>` to ensure consistent scoping.
3. It replays `cargo check` until no lifetime mismatches remain.

The resulting signature is significantly more natural (in this case it is identical to the manual extraction).

```
1 fn extract_include_path<'a>(  
2     target_config: &mut File<'a>,  
3     include_paths: &mut Vec<Path<'a>>,  
4     id: SectionId)
```

The shift from `&SectionId` to `SectionId` demonstrates REM’s key contribution: making semantically appropriate ownership decisions automatically, while still producing code that compiles under Rust’s strict type system.

3.2.3 CHARON

The first stage of our verification pipeline is translation from Rust into a lower-level, verification-friendly form. This role is performed by **CHARON**, a Rust to LLBC (Low-Level Borrow Calculus) translator. CHARON accepts Rust source code (or more precisely, its compiler intermediate representation MIR) ⁴. CHARON’s main task is to extract complex semantic information from `rustc` and produce a machine readable output containing said information.

⁴Even more precisely, the tool itself accepts Rust code, but its initial steps rely on a complex interaction with `rustc` to acquire the MIR

CHARON performs several non-trivial transformations:

Structured control flow. MIR is a control-flow graph with arbitrary `gotos`; CHARON reconstructs this into structured `if`, `match`, and `loop` constructs, or preserves the raw form with the `--ullbc` option.

Trait and type resolution. Explicitly records how trait bounds are proven, normalises default methods, and can transform associated types into explicit parameters (`--remove-associated-types`).

Lifetime handling. Hides the distinction between early- and late-bound variables and makes implied bounds explicit, simplifying reasoning about generic functions.

Closures and vtables. Represented as ordinary structs implementing the relevant `Fn*` traits, making them accessible to later verification steps.

Noise reduction. Options such as `--hide-marker-traits` remove built-in traits (`Sized`, `Send`, etc.) that would otherwise clutter the output.

The result is a cleaned, semantically faithful program in LLBC form. Unlike raw MIR, which is compiler-oriented, LLBC is designed for verification: ownership and borrowing are represented directly, trait resolution is explicit, and control flow is structured. A sample program and small portion of its MIR are shown below. However, the output of CHARON is not designed for human readability, and even the tiny program in Figure 3.3 produces over 140,000 characters of LLBC. This (very verbose) explicitness is what enables the subsequent translation into pure functional code for verification.

```

1 fn main() {
2     let v = vec![1, 2, 3];
3     let x = &v[0];
4     println!("{}", x);
5 }
```

Rust source

```

1 bb2: {
2     _9 = &_1;
3     _8 = <Vec<i32> as
        ↳ Index<usize>>::index(move _9,
        ↳ const 0_usize);
4 }
5
6 bb3: {
7     _7 = copy _8;
8     _16 = &_7;
9     _15 =
        ↳ core::fmt::rt::Argument:::<'_>::new_display
        ↳ _16);
10 }
```

Excerpt of MIR (`rustc`)

Figure 3.3: Illustration of CHARON’s role. While only Rust and MIR are shown here, CHARON translates the MIR into LLBC

3.2.4 AENEAS

With the program translated into LLBC by CHARON, the next stage of the pipeline is handled by **AENEAS**. AENEAS takes the borrow- and lifetime-explicit LLBC and

translates it into a *purely functional form*, stripping away low-level memory operations while preserving the program’s semantics. This functionalisation step is what makes the program suitable for downstream reasoning in Coq: instead of reasoning about references and loans, we can reason about values and functions. Currently AENEAS supports translation into *Coq*, *Lean*, *HOL4*, and *F**.

In practice, AENEAS acts as the bridge between compiler-level detail and proof-level abstraction. CHARON ensures that all the complexities of lifetimes, loads, and trait resolution are made explicit; AENEAS then packages these into a form where correctness properties can be stated and proved without touching the borrow checker or the stack/heap. For our pipeline, this is the point where Rust code “stops being Rust” and becomes a functional program whose equivalence can be checked in Coq.

The central insight behind AENEAS is that Rust’s ownership discipline makes much of low-level memory reasoning unnecessary.

3.2.5 Coq

AENEAS translates the LLBC into a purely functional form, in this case Coq. The final stage of the pipeline is thus to join the two Coq files (P_{coq} and P'_{coq}) with all of the auxiliary information needed to automatically verify equivalence between the original and refactored Rust programs.

What is Coq?

Coq is an interactive theorem prover and proof assistant built on a dependently typed functional language. It allows developers to define mathematical objects, programs, and logical propositions, and to construct machine-checked proofs about them. Every proof is validated by Coq’s small, trusted kernel, giving very high assurance of correctness. Coq has a long history of use in large-scale verification projects—such as CompCert (a verified C compiler) and formally verified operating systems—making it a natural fit for reasoning about the safety and correctness of Rust programs.

Why Coq?

Coq was selected over alternatives such as HOL4, Lean, or F* for three main reasons:

Mature ecosystem Decades of development have produced robust tooling, extensive libraries, and a stable kernel.

Direct integration AENEAS already targets Coq, whereas HOL4 and Lean would require substantial new backends for compatibility.

Automation Coq’s tactic language and SMT integrations allow many routine proofs—such as equivalence checks—to be discharged automatically.

Lean and HOL4 both provide highly expressive logics, but at the cost of significant engineering effort to integrate with AENEAS. F* is closer in spirit to AENEAS, but the platform is less mature and would duplicate existing functionality. Coq therefore offers the best combination of maturity, automation, and direct compatibility with our pipeline.

Verification Examples

To make the Coq stage of the pipeline concrete, we present two examples that have followed the entire process from Rust source through to Coq. Together, they illustrate both the best-case scenario and the additional complexity involved in real-world refactorings.

Before looking at the translations themselves, it is worth noting the supporting infrastructure. AENEAS generates not just the Coq files for individual functions, but also a small project scaffold. The `Primitives.v` file defines the Coq representations of Rust’s core datatypes (`i32`, `Vec`, `Option`, etc.), providing a foundation on which the translations depend. The `_CoqProject` file ties the entire build together, specifying load paths and dependencies so that the project can be automatically compiled and checked consistently.

1

1

Simple example: original translation

Simple example: refactored translation

Figure 3.4: Coq translations for the simple example.

Alongside these translations, the verifier produces the `EquivCheck.v` file. This module instantiates both the original and refactored definitions and generates an equivalence theorem. In practice, this means the Coq project is set up to automatically check that the two versions produce the same results for all inputs. We defer discussion of the proof obligations themselves to Section 3.3, but here the important point is that the generated Coq project is self-contained: once compiled, the equivalence is verified without further user intervention.

3.3 Proof Obligations

In the final stage of the verification pipeline, we discharge a formal *equivalence obligation* inside Coq. Intuitively, we must show that the refactored program behaves identically to the original program for all valid inputs. The central idea is functional equivalence: the observable outputs of the two programs must match.

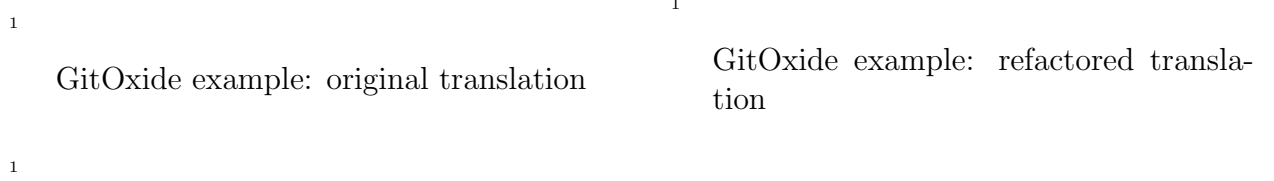


Figure 3.5: Coq translations and automatically generated equivalence check for the GitOxide example.

3.3.1 Simplified Obligation

In its simplest form, the obligation can be stated as:

$$\forall x. \text{tr}(f)(x) = \text{tr}(f')(x)$$

where f is the original function, f' is its refactored counterpart, and $\text{tr}(\cdot)$ denotes the translation into Coq performed by AENEAS. This expresses that for every input x , the translated functions return identical results. In practice, this is the form discharged automatically in our examples (see Section ??).

3.3.2 Whole-Program Context

For a more complete characterisation, we must acknowledge that functions are never executed in isolation, but as part of the enclosing program. Thus the obligation can be strengthened to whole-program equivalence:

$$\forall x. \text{tr}(P)(x) = \text{tr}(P')(x)$$

where P is the original program and P' is the refactored program, differing only in that f is replaced by f' . This captures the fact that equivalence must hold regardless of how the function is used internally.

An equivalent way to phrase this is through *contextual equivalence*:

$$\forall C, x. \text{tr}(C[f])(x) = \text{tr}(C[f'])(x)$$

where $C[\cdot]$ denotes an arbitrary program context. This emphasises that the replacement of f by f' preserves behaviour under any calling environment.

3.3.3 Relation to Prior Work

Earlier work on AENEAS stated equivalence in a deliberately general form, including preconditions $\Phi(x)$, effect traces τ , and projections $\text{obs}(\tau)$ onto observable behaviours.

These ingredients were necessary to account for impure features, partial programs, and richer effect semantics. In contrast, our pipeline deliberately restricts attention to a simpler setting. REM only produces well-typed, borrow-checked Rust fragments, and AENEAS is applied to safe Rust code without interior mutability or concurrency. Under these restrictions:

- Preconditions can be dropped, as REM ensures typing and borrow safety.
- Effect traces can be omitted, as our fragments are side-effect free modulo return values.
- Observable projections reduce to plain equality of results.

This simplification yields an obligation that is easier to state and automate, while still capturing the core requirement that refactoring preserves behaviour.

3.3.4 Practical Implications

The simplified obligation highlights a guiding design principle: the pipeline is intended to complement the extraction algorithm rather than over-approximate it. It is therefore preferable to *fail early* when unsupported constructs are encountered, rather than risk certifying programs incorrectly. AENEAS reports unsupported features (e.g. closures or nested borrows) as translation failures, while Coq distinguishes between an unprovable equivalence and an unsupported fragment. This separation ensures that developers never receive a false positive: code is either successfully verified, or it fails explicitly with a clear diagnostic.

This leads naturally to the question of coverage: which Rust features are currently supported by AENEAS, and which remain open problems? We address this in the following section.

3.4 Limitations of Current Verification

At present, the verification guarantees supported by our pipeline are constrained by the subset of Rust that AENEAS can translate. AENEAS is designed to operate on safe Rust; programs that make use of `unsafe` constructs fall outside its scope. While this excludes some low-level systems code, the restriction is deliberate: safe Rust already captures the majority of idiomatic Rust usage, and the long-term plan is to integrate AENEAS with complementary tools that target `unsafe` code. The current support landscape is summarised in Table 3.1.

Within safe Rust, several technical limitations remain. Loops are supported only in restricted forms (no nested loops), though this is under active development. Function

pointers and closures are not yet available, although trait-based abstractions are supported and ongoing work aims to extend this to first-class functions. Type parametricity is limited: currently, it is not possible to instantiate a generic type parameter with a type that itself contains a borrow. Similarly, function signatures cannot yet contain nested borrows, though this too is being addressed. Interior mutability is another open challenge; the current plan is to capture its effects using ghost states. Finally, concurrent execution is out of scope: AENEAS models sequential Rust programs only, with parallelism considered a long-term research goal.

Table 3.1: Current AENEAS support for Rust features

Feature	Status
Safe Rust	Fully supported
Unsafe Rust	Not supported (future integration planned)
Loops	Supported (no nested loops)
Function pointers / Closures	Not supported (work in progress)
Traits	Supported
Type parametricity	Limited: type parameters with borrows not supported
Nested borrows in function signatures	Not supported (work in progress)
Interior mutability (<code>Cell</code> , <code>RefCell</code> , etc.)	Not supported (planned via ghost states)
Concurrency	Out of scope (long-term research goal)

It is important to emphasise that the verification pipeline is designed to complement the extraction algorithm. As such, it is preferable for verification to *fail early* when encountering unsupported features rather than risk certifying code incorrectly. A failed verification due to unsupported language constructs is reported differently to the developer than a failed proof obligation. This distinction ensures that false positives—cases where unverifiable code would otherwise appear verified—are avoided, preserving the integrity of the guarantees provided by the pipeline.

In summary, the present pipeline is best understood as a verification framework for sequential, safe Rust programs without advanced features such as nested borrows, closures, or concurrency. These restrictions reflect the state of AENEAS rather than fundamental barriers, and ongoing work is progressively extending the supported subset. Broader limitations of the full REM–AENEAS–Coq pipeline are deferred to Section 5.2.

Evaluation and Experimental Results

4.1 Experimental Setup

4.2 Evaluation Criteria

4.3 Extraction: Results and Analysis

4.3.1 Complete Results Table

4.3.2 Analysis?

4.4 Verification: Results and Analysis

4.4.1 Complete Results Table

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Project (LOC)	Type	Size (LOC)		Verification		Time (s)		
		CLR	CLE	Attempted	Successful	LLBC Conversion	CoQ Conversion	Proof & V
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15								
16								
17								
18								
19								
20								

Table 4.2: Verification Results. See appendix ?? for details of each repository.

Conclusions and Future Work

5.1 Summary of Contributions

5.2 Limitations

5.3 Potential Improvements

5.4 Future Work

5.5 Broader Impact

Appendices

6.1 Appendix 1: Passing by Value and Reference

This appendix provides three short, self-contained examples of “naïve” function hoisting in Rust. In each example, we will define a function that is getting refactored, and we will demonstrate the case where the arguments to the function are passed by value and the case where they are passed by reference. To achieve this, we are imagining an automated refactoring tool that initially takes an inline block of code and moves it into a new function. The compilation errors (if any) guide a subsequent ownership analysis, which deduces how each variable should be passed: by value (`T`), by shared reference (`&T`), or by mutable reference (`&mut T`).

6.1.1 Example 1: Passing by Value

Original Code (inline)

```
1 fn main() {
2     let x = 42;
3     let y = x * 2;
4     println!("y is {}", y);
5 }
```

Naïve Hoisted Code (before fixes)

Suppose an automated refactoring tool decides to “hoist” the multiplication logic into a new function:

```
1 fn hoisted_block(x: i32) {
2     let y = x * 2;
3     println!("y is {}", y);
4 }
5
6 fn main() {
7     let x = 42;
8     hoisted_block(x);
9 }
```

In this example, no errors occur because after calling `hoisted_block(x)`, we do not need `x` again in `main`, so moving (passing ownership) is safe. The automated repair analysis detects that `x` is no longer needed in `main` and decides it can be passed by value.

6.1.2 Example 2: Passing by Shared Reference

Original Code (inline)

```

1 fn main() {
2     let text = String::from("Hello world!");
3     let length = text.len();
4
5     // Some inline block that only reads from `text`:
6     println!("First word is: {}", get_first_word(&text));
7     println!("Total length is: {}", length);
8 }
9
10 fn get_first_word(s: &str) -> &str {
11     s.split_whitespace().next().unwrap_or("")
12 }
```

Naïve Hoisted Code (before fixes)

Let's say the refactoring tool decides to hoist the usage of `get_first_word` (or a bit more logic around it) into a new helper function:

```

1 fn main() {
2     let text = String::from("Hello world!");
3     let length = text.len();
4
5     // Some inline block that only reads from `text`:
6     println!("First word is: {}", get_first_word(&text));
7     println!("Total length is: {}", length);
8 }
9
10 fn get_first_word(s: &str) -> &str {
11     s.split_whitespace().next().unwrap_or("")
12 }
```

This version won't compile as expected if the function `print_first_word_block` takes `text` by value. Once `text` is moved into `print_first_word_block`, we can no longer use `text` afterward in `main`. The compiler will complain that `text` is moved, thus invalidating the line `println!("Total length is: ", length);` if we needed `text` for anything else.

Automated Repair Yields

```

1 fn print_first_word_block(text: &String) {
2     // Only reads from `text`, so it only needs a shared reference.
3     println!("First word is: {}", get_first_word(text));
4 }
5
6 fn main() {
```

```

7     let text = String::from("Hello world!");
8     let length = text.len();
9
10    print_first_word_block(&text); // Pass a shared reference
11    println!("Total length is: {}", length);
12 }

```

What's going on?

- **Read-only usage:** The function `print_first_word_block` just needs to read the string (in order to print the first word).
- **Shared reference:** Because the string must remain valid afterward in `main`, the ownership analysis decides that passing by shared reference (`&String`) is correct.

6.1.3 Example 3: Passing by Mutable Reference

Original Code (inline)

```

1 fn main() {
2     let mut values = vec![1, 2, 3];
3
4     // Some inline block that mutates `values`:
5     values.push(4);
6     println!("Values are now: {:?}", values);
7 }

```

Naïve Hoisted Code (before fixes)

Hoisting the push operation into a new function could look like this:

```

1 fn push_block(values: Vec<i32>) {
2     let mut v = values;
3     v.push(4);
4     println!("Values are now: {:?}", v);
5 }
6
7 fn main() {
8     let mut values = vec![1, 2, 3];
9     push_block(values);
10    println!("Final values: {:?}", values); // error: `values` was moved
11 }

```

Here, we have a compilation error similar to Example 2: once `values` is passed by value, ownership is transferred, and we cannot use `values` again in `main`. However, in this case, we truly do need to mutate `values`, and we want `main` to see the updated vector.

Automated Repair Yields

```

1 fn push_block(values: &mut Vec<i32>) {
2     values.push(4);
3     println!("Values are now: {:?}", values);
4 }
5

```

```
6 fn main() {  
7     let mut values = vec![1, 2, 3];  
8     push_block(&mut values);  
9     println!("Final values: {:?}", values); // Now values is still in scope and updated  
10 }
```

What's going on?

- **Mutation required:** The code needs to modify the original vector.
- **Mutable reference:** By passing `&mut Vec<i32>`, the function can mutate `values` in place, and `main` retains ownership of `values` to use afterward.

6.2 Appendix 2: Details of each Repository in the Evaluation

6.3 Appendix 3: GitOxide Example in Detail

The following diff shows how the ‘if add_path’ block was refactored into the new helper method `extract_include_path`.

```

1 -         if add_path {
2 -             if let Some(body) = target_config.sections.get(&id) {
3 -                 let paths = body.values(&Key::from("path"));
4 -                 let paths = paths.iter().map(|path|
↪ values::Path::from(path.clone()));
5 -                 include_paths.extend(paths);
6 -             }
7 -         }
8         // Some specifics ommited for simplicity
9 +         if add_path {
10 +             extract_include_path(target_config, &mut include_paths, id)
11 +         }

```

And here is the new method in full:

```

1 fn extract_include_path<'a>(<
2     target_config: &mut File<'a>,
3     include_paths: &mut Vec<values::Path<'a>>,
4     id: SectionId
5 ) {
6     if let Some(body) = target_config.sections.get(&id) {
7         let paths = body.values(&Key::from("path"));
8         let paths = paths.iter().map(|path| values::Path::from(path.clone()));
9         include_paths.extend(paths);
10    }
11 }

```

6.4 Appendix 4: Complete MIR from ??

```

1 // WARNING: This output format is intended for human consumers only
2 // and is subject to change without notice. Knock yourself out.
3 fn main() -> () {
4     let mut _0: ();
5     let _1: std::vec::Vec<i32>;
6     let mut _2: std::boxed::Box<[i32]>;
7     let mut _3: usize;
8     let mut _4: usize;
9     let mut _5: *mut u8;
10    let mut _6: std::boxed::Box<[i32; 3]>;
11    let _8: &i32;
12    let mut _9: &std::vec::Vec<i32>;
13    let _10: ();
14    let mut _11: std::fmt::Arguments<'_>;
15    let _12: &[&str; 2];
16    let _13: &[core::fmt::rt::Argument<'_>; 1];
17    let _14: [core::fmt::rt::Argument<'_>; 1];
18    let mut _15: core::fmt::rt::Argument<'_>;
19    let _16: &&i32;
20    let mut _17: *const [i32; 3];
21    let mut _18: *const ();
22    let mut _19: usize;
23    let mut _20: usize;
24    let mut _21: usize;
25    let mut _22: usize;
26    let mut _23: bool;
27    scope 1 {
28        debug v => _1;
29        let _7: &i32;
30        scope 2 {
31            debug x => _7;
32        }
33    }
34
35    bb0: {
36        _3 = SizeOf([i32; 3]);
37        _4 = AlignOf([i32; 3]);
38        _5 = alloc::alloc::exchange_malloc(move _3, move _4) -> [return: bb1, unwind
39        ↪ continue];
40    }
41
42    bb1: {
43        _6 = ShallowInitBox(move _5, [i32; 3]);
44        _17 = copy (((_6.0: std::ptr::Unique<[i32; 3]>).0: std::ptr::NonNull<[i32;
45        ↪ 3]>).0: *const [i32; 3]);
46        _18 = copy _17 as *const () (PtrToPtr);
47        _19 = copy _18 as usize (Transmute);
48        _20 = AlignOf([i32; 3]);
49        _21 = Sub(copy _20, const 1_usize);
50        _22 = BitAnd(copy _19, copy _21);
51        _23 = Eq(copy _22, const 0_usize);
52        assert(copy _23, "misaligned pointer dereference: address must be a multiple
53        ↪ of {} but is {}", copy _20, copy _19) -> [success: bb10, unwind
54        ↪ unreachable];
55    }
56
57    bb2: {
58        _9 = &_1;
59        _8 = <Vec<i32> as Index<usize>>::index(move _9, const 0_usize) -> [return:
60        ↪ bb3, unwind: bb8];
61    }
62
63    bb3: {
64        _7 = copy _8;
65        _12 = const main::promoted[0];
66        _16 = &_7;
67        _15 = core::fmt::rt::Argument::<'_>::new_display::<&i32>(copy _16) -> [return:
68        ↪ bb4, unwind: bb8];

```

```

63     }
64
65     bb4: {
66         _14 = [move _15];
67         _13 = &_14;
68         _11 = Arguments:::<'>::new_v1::<2, 1>(copy _12, copy _13) -> [return: bb5,
        ↪ unwind: bb8];
69     }
70
71     bb5: {
72         _10 = _print(move _11) -> [return: bb6, unwind: bb8];
73     }
74
75     bb6: {
76         drop(_1) -> [return: bb7, unwind continue];
77     }
78
79     bb7: {
80         return;
81     }
82
83     bb8 (cleanup): {
84         drop(_1) -> [return: bb9, unwind terminate(cleanup)];
85     }
86
87     bb9 (cleanup): {
88         resume;
89     }
90
91     bb10: {
92         (*_17) = [const 1_i32, const 2_i32, const 3_i32];
93         _2 = move _6 as std::boxed::Box<[i32]> (PointerCoercion(Unsize, Implicit));
94         _1 = slice::<impl [i32]>::into_vec::<std::alloc::Global>(move _2) -> [return:
        ↪ bb2, unwind continue];
95     }
96 }
97
98 const main::promoted[0]: &[&str; 2] = {
99     let mut _0: &[&str; 2];
100     let mut _1: [&str; 2];
101
102     bb0: {
103         _1 = [const "", const "\n"];
104         _0 = &_1;
105         return;
106     }
107 }

```

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