Combining ECA Rules with Process Algebras for the Semantic Web

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Motivation and Goals

(Semantic) Web:

- XML: bridge the heterogeneity of data models and languages
- RDF, OWL provide a computer-understandable semantics
- ... same goals for describing behavior:
 - description of behavior in the Semantic Web
 - semantic description of behavior

Event-Condition-Action Rules are suitable for both goals:

- operational semantics
- ontology of rules, events, actions

ECA Rules

"On Event check Condition and then do Action"

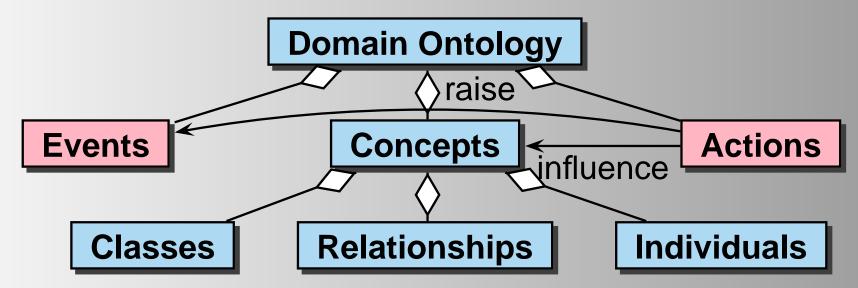
- paradigm of Event-Driven Behavior,
- modular, declarative specification in terms of the domain ontology
- sublanguages for specifying Events, Conditions, Actions
- global ECA rules that act "in the Web"

Requirements

- Ontology of behavior aspects
- modular markup definition,
- implement an operational and executable semantics

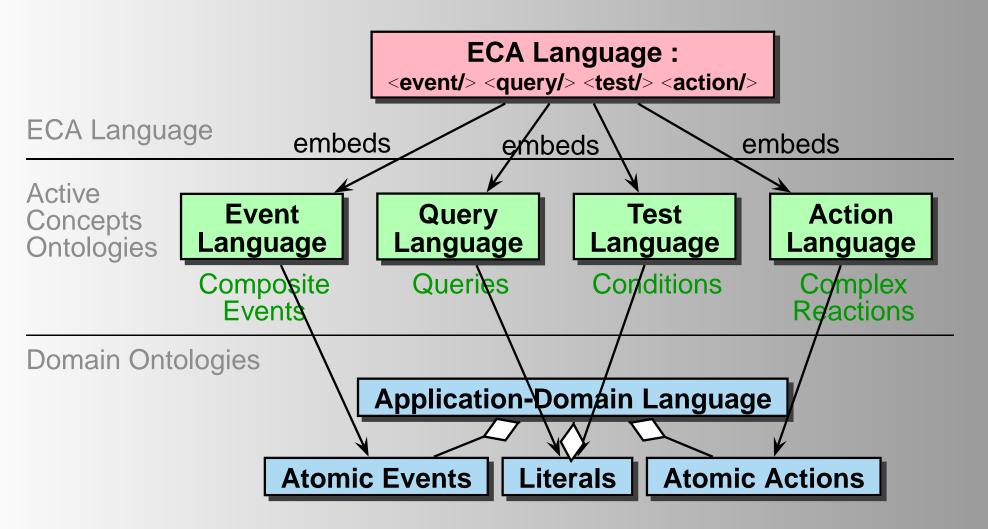
Events and Actions in the Semantic Web

- applications do not only have an ontology that describes static notions
 - cities, airlines, flights, etc., relations between them ...
- but also an ontology of events and actions
 - cancelling a flight, cancelling a (hotel, flight) booking,
- Domain languages also describe behavior:



Embedding of Languages

... there are not only atomic events and actions.

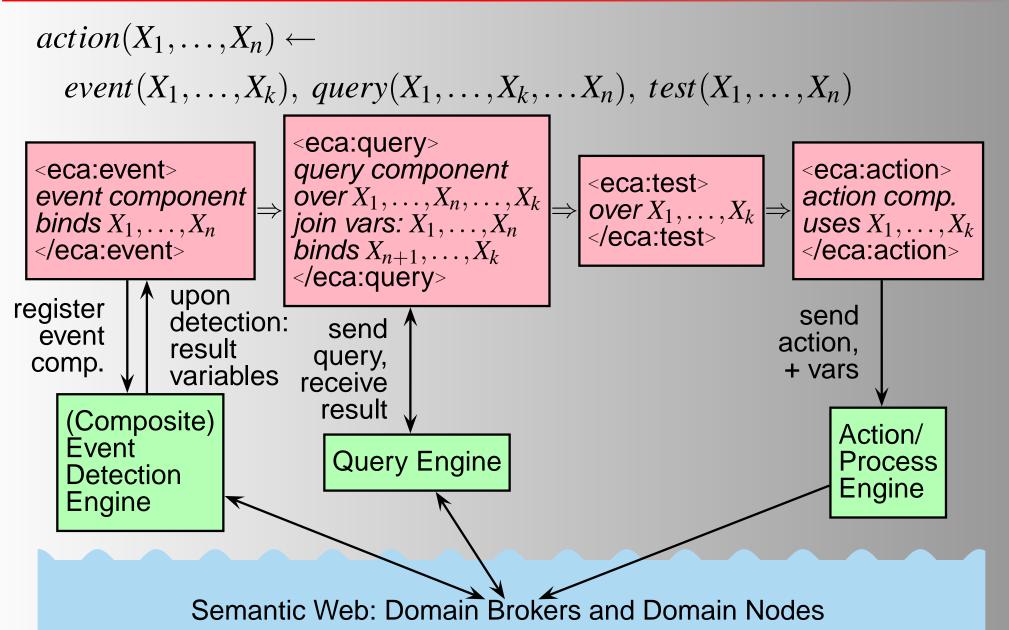


Rule Markup: ECA-ML

</eca:rule>

```
<!ELEMENT rule (event,query*,test?,action<sup>+</sup>) >
<eca:rule rule-specific attributes>
 <eca:event identification of the language >
  event specification, probably binding variables
 </eca:event>
 <eca:query identification of the language > <!-- there may be several queries -->
  query specification; using variables, binding others
 </eca:query>
 <eca:test identification of the language >
  condition specification, using variables
 </eca:test>
 <eca:action identification of the language > <!-- there may be several actions -->
  action specification, using variables, probably binding local ones
 </eca:action>
```

Binding and Use of Variables in ECA Rules



Rule Markup: Example (Stripped)

```
<!ELEMENT rule (event,query*,test?,action+) >
<eca:rule xmlns:travel="http://www.travel.de">
 <eca:event xmlns:snoop="http://www.snoop.org">
  <snoop:seq> <travel:delayed-flight flight="{$flight}"/>
     <travel:canceled-flight flight="{\$flight}"/> </snoop:seq>
 </eca:event>
 <eca:query>
  <eca:variable name="email">
    <eca:opaque lang="http://www.w3.org/xpath">
     doc("http://xml.lufthansa.de")/flights[code="{$flight}"]/passenger/@e-mail
    </eca:opaque> </eca:variable> </eca:query>
 <eca:action xmlns:smtp="...">
  <smtp:send-mail to="$email" text="..."/>
 </eca:action>
</eca:rule>
```

Active Concepts Ontologies

Domains specify atomic events, actions and static concepts

Composite [Algebraic] Active Concepts

- Event algebras: composite events
- Process algebras (e.g. CCS)
- consist of composers/operators to define composite events/processes,
- leaves of the terms are atomic domain-level events/actions,
- as operator trees: "standard" XML markup of terms
- RDF markup as languages,
- every expression can be associated with its language.

Composite Actions: Process Algebras

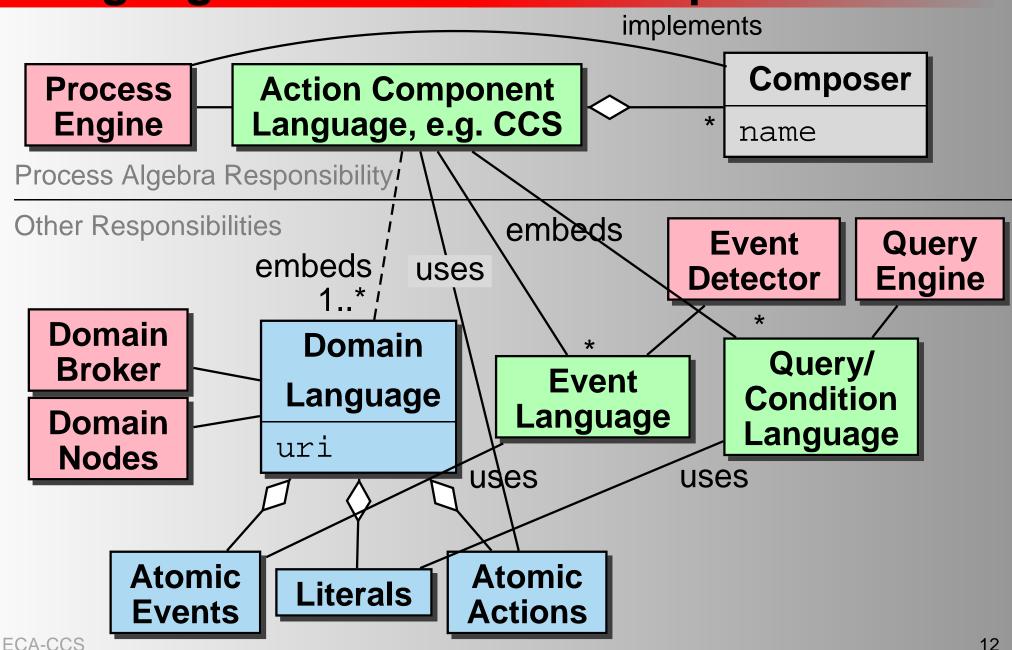
- e.g., CCS Calculus of Communicating Systems [Milner'80]
- operational semantics defined by transition rules, e.g.
 - a sequence of actions to be executed,
 - a process that includes "receiving" actions,
 - guarded (i.e., conditional) execution alternatives,
 - the start of a fixpoint (i.e., iteration or even infinite processes), and
 - a family of communicating, concurrent processes.
- Originally only over atomic processes/actions
- reading and writing simulated by communication a (send), \bar{a} (receive) "match" as communication
- ... extend this to the (Semantic) Web environment with autonomous nodes.

Adaptation of Process Algebras

Goal: specification of reactions in ECA rules

- liberal asynchronous variant of CCS: go on when possible, waiting and delaying possible
- extend with variable bindings semantics
- input variables come bound to values/URIs
- additional variables can be bound by "communication"
- queries as atomic actions: to be executed, contribute to the variable bindings
- event subexpressions as atomic actions: like waiting for \bar{a} communication
- ⇒ subexpressions in other kinds of component languages

Languages in the Action Component



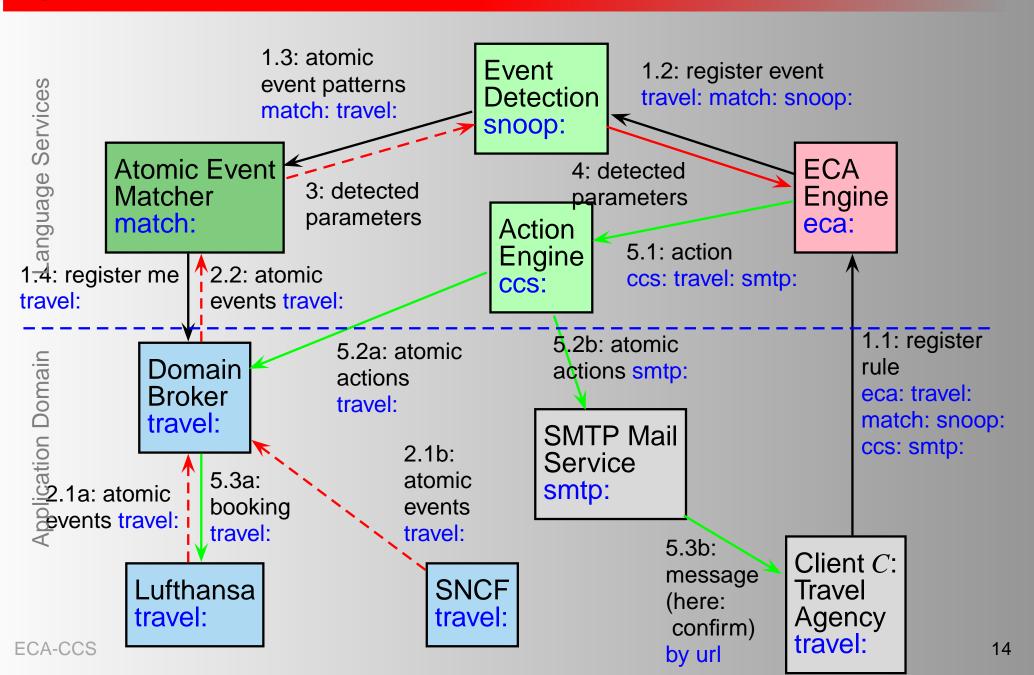
CCS Markup

- <ccs:sequence>CCS subexpressions </ccs:sequence></ccs:alternative>CCS subexpressions </ccs:alternative></ccs:concurrent>CCS subexpressions </cs:concurrent>
- <ccs:fixpoint variables="X1 X2 ... Xn" index="i" // "my" index localvars="..."> n subexpressions </ccs:fixpoint>
- <ccs:atomic-action>domain-level action </ccs:atomic-action></ccs:event xmlns:ev-ns="uri">event expression </ccs:event></ccs:query xmlns:q-ns="uri">query expression </ccs:query></ccs:test xmlns:t-ns="uri">test expression </ccs:test>

Embedding Mechanisms: Same as in ECA-ML

- communication by logical variables
- namespaces for identifying languages of subexpressions

Service-Based Architecture



Comparison

- CCS (extended with events and queries) strictly more expressive than ECA rules alone:
 ECA pattern in CCS: event:condition:action,
- many ECA rules have much simpler actions and do not need CCS,
- useful to have CCS as an option for the action part.

Summary

- RDF/OWL as integrating semantic model in the Semantic Web
- describe events and actions of an application within its RDF/OWL model
- languages of different expressiveness/complexity available
- ECA rules
 - components
 - application-level atomic events and atomic actions
 - specific languages (event algebras, process algebras)
- Architecture: functionality provided by specialized nodes

Thank You Questions ??

Further information and publications: http://dbis.informatik.uni-goettingen.de/eca/

Complementing Slides

Action Component: Process Algebras

- example: CCS (Calculus of Communicating Systems, Milner 1980)
- describes the execution of processes as a transition system:
 - (only the asynchronous transitions are listed)

$$a: P \xrightarrow{a} P$$
 , $\frac{P_i \xrightarrow{a} P}{\sum_{i \in I} P_i \xrightarrow{a} P}$ (for $i \in I$)
$$\frac{P \xrightarrow{a} P'}{P|Q \xrightarrow{a} P'|Q}$$
 , $\frac{Q \xrightarrow{a} Q'}{P|Q \xrightarrow{a} P|Q'}$

$$\frac{P_i \{ \text{fix } \vec{X} \vec{P} / \vec{X} \} \xrightarrow{a} P'}{\text{fix:} \vec{X} \vec{P} \xrightarrow{a} P'}$$

Atomic Event Specifications

```
Sample Event: | <travel:canceled-flight flight="LH123">
                     <travel:reason>bad weather</travel:reason>
                   </travel:canceled-flight>
```

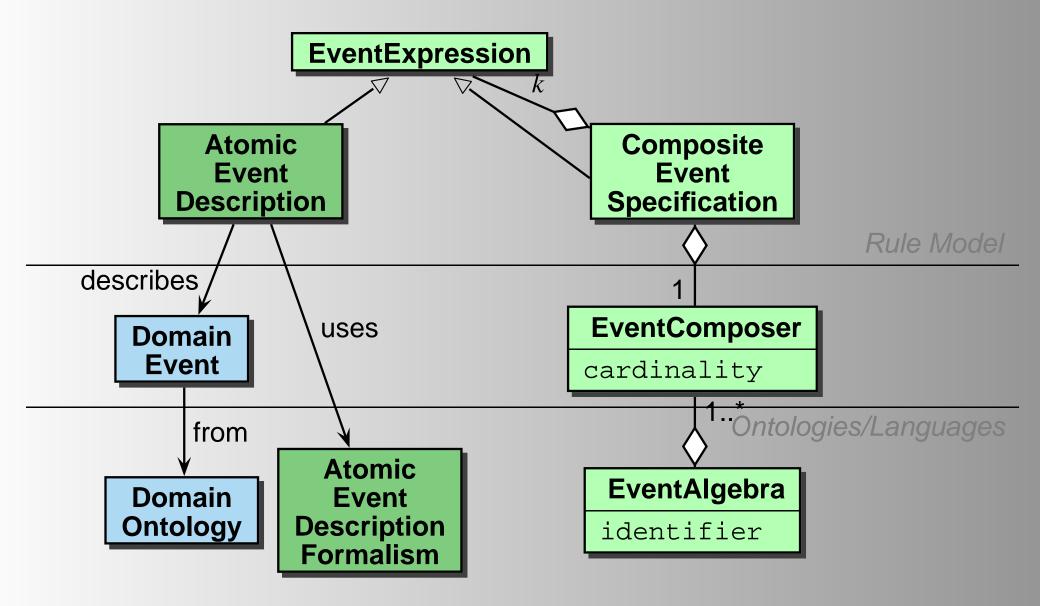
Event expressions require an auxiliary formalism for specifying relevant events:

- type of event ("travel:canceled-flight"),
- constraints ("must have a travel:reason subelement"),
- extract data from events ("bind @flight to variable

Sample: XML-QL-style matching

```
<atomic-event language="match">
 <travel:canceled-flight flight="{$flight}"><travel:reason/></travel:canceled-flight>
</atomic-event>
```

Event Expressions: Languages



Sample Markup (Event Component)

```
<eca:rule xmlns:travel="...">
 <eca:variable name="theSeq">
  <eca:event xmlns:snoop="...">
  <snoop:sequence>
     <snoop:atomic-event language="match">
      <travel:delayed-flight flight="{$Flight}" minutes="{$Minutes}"/>
     </snoop:atomic-event>
     <snoop:atomic-event language="match">
      <travel:canceled-flight flight="{$Flight}"/>
     </snoop:atomic-event>
   </snoop:sequence>
  </eca:event>
                        binds variables:
 </eca:variable>
```

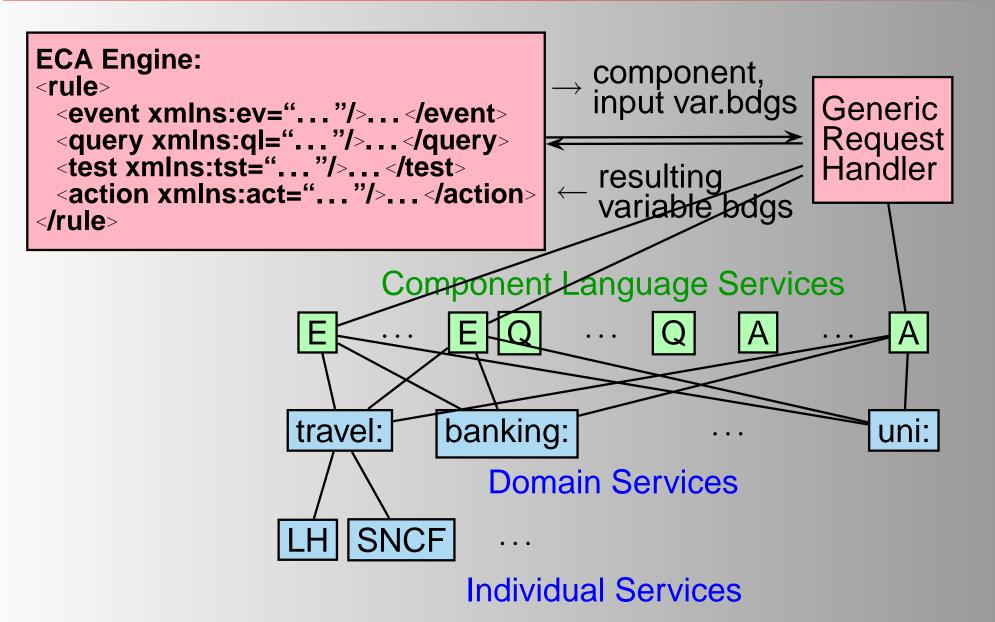
</eca:rule>

- Flight, Minutes: by matching
- theSeq is bound to the sequence of events that matched the pattern

Tasks

- ECA Engine: Rule Semantics
 - Control flow: registering event component, receiving "firing" answer, continuing with queries etc.
 - Variable Bindings, Join Semantics
- Generic Request Handler: Mediator with Component Engines
 - depending on Service Descriptions
- Component Engines: dedicated to certain Event Algebras,
 Query Languages, Action Languages
- Domain Services (Portals): atomic events, queries, atomic actions

ECA Architecture



Communication of Variable Bindings

XML markup for communication of variable bindings:

Communication ECA → **GRH**

- the component to be processed
- bindings of all relevant variables

```
[Sample: a query component]

<eca:query xmlns:ql="url"

rule="rule-id" component="component-id">

<!-- query component -->

<eca:query>
<log:variable-bindings>
 <log:tuple> ... </log:tuple>
  :

<log:tuple> ... </log:tuple>
<log:variable-bindings>
```

- url is the namespace used by the event language
- identifies appropriate service

Communication

ECA engine sends component to be processed together with bindings of all relevant variables to GRH.

Generic Request Handler (GRH)

- Submits component (with relevant input/used variable bindings) to appropriate service (determined by namespace/language used in the component)
- if necessary: does some wrapping tasks (for non-framework-aware services)
- receives results and transforms them into flat variable bindings and sends them back to the ECA engine ...
- ... where they are joined with the existing tuples ...
- ... and the next component is processed.

Communication Component Engine → **GRH**

result-bindings-pairs (semantics of expression)

```
log:answers rule="rule-id" component="component-id">
< log:answer>
 <log:result>
  <!-- functional result -->
 log:variable-bindings>
  <log:tuple> . . . 
  <log:tuple> . . . </log:tuple>
 /log:variable-bindings>
log:answer> ... 
log:answer> ... 
/log:answers>
```

Communication GRH → **ECA**

- set of tuples of variable bindings
 (i.e., input/used variables and output/result variables)
- is then joined with tuples in ECA engine
- ... and next component is processed