

A Datalog+ RuleML 1.01 Architecture for Rule-Based Data Access in Ecosystem Research

(Long version: cs.unb.ca/~boley/talks/RulesOBDA.pdf)

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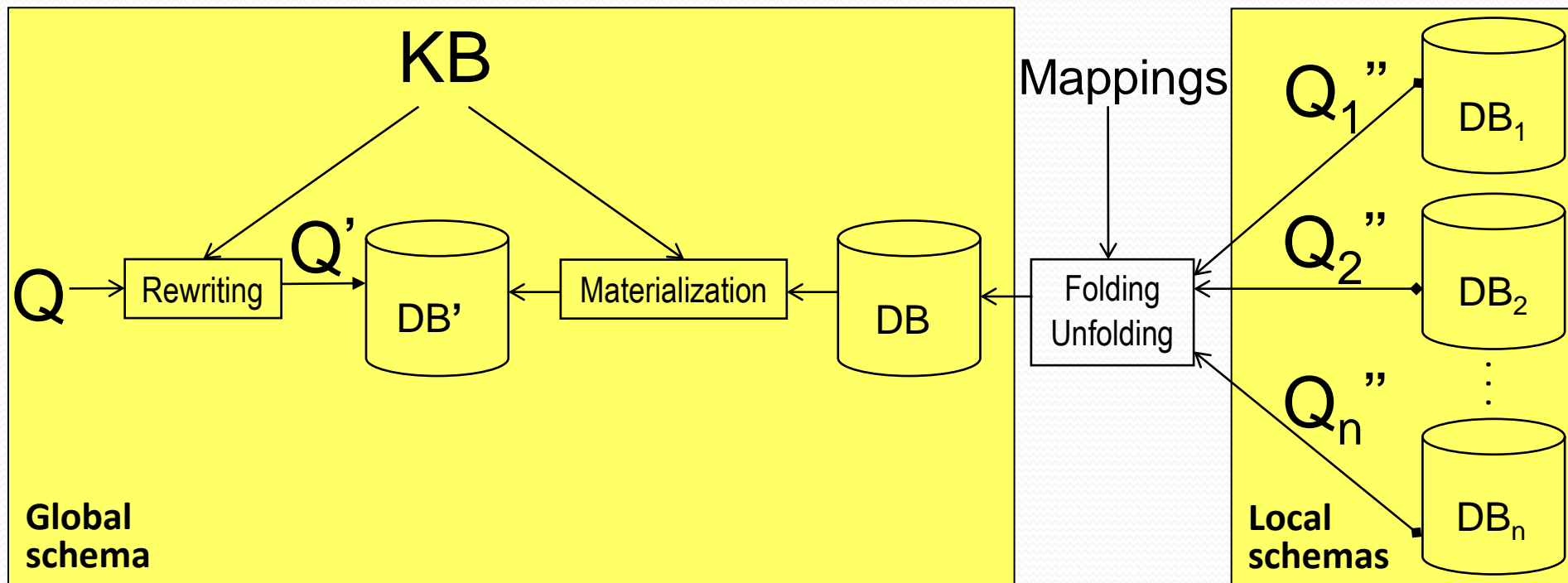
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What is Knowledge-Based Data Access?

- KBDA applies AI / Semantic Technologies to databases
- Founded on **(logical) knowledge** representation for:
 - **Ontology-Based Data Access (OBDA)**, e.g. data integration/federation and query optimization
 - **Rule-Based Data Access (RBDA)**, e.g. Datalog / deductive databases and query answering
- **Knowledge Base** as generalized **global** schema for data in **local** (e.g., relational or graph) DBs
- KB module amplifies data storage & query execution of **distributed, heterogeneous** (No)SQL DBs
- Provides **multi-purpose** knowledge level for data

Preview: Unified Architecture



Why Knowledge-Based Data Access?

- Domain knowledge utilized to deal with data torrent
 - Domain experts conceptually *fold* data / *unfold* queries via **Mappings** defined with IT (SQL, SPARQL, ...) experts
 - User concepts are captured in **Knowledge Base** for *domain-enriched* database materialization / querying without IT experts
 - **Engines** use KB to deduce answers implicit in DBs
 - **Analytics** enabled by queries exploring hypotheses
- KB as major organizational resource also for, e.g.:
 - Data validation (consistency, completeness, ...)
 - Schema-level query answering (even without DBs)

RBDA Realizes Uniform KBDA – 1 of 3: Queries as Rules

1. a) A conjunctive **query** is a special Datalog **rule** whose body can be *rewritten* (see 2.) and *unfolded* (see 3.), and whose head instantiates the distinguished answer variables of the body
b) KBDA ontologies beyond RDF Schema (RDFS) often permit **Boolean** conjunctive queries corresponding to **integrity rules**
2. ...
3. ...

RBDA Realizes Uniform KBDA – 2 of 3: KBs as Rules

1. ...
2. KBDA **KB** supports, e.g., query *rewriting* through global-schema-level reasoning, including with RDFS **taxonomies** or Datalog **rule** axioms, and DL-Lite (OWL 2 QL) or (head-)existential **rules**; KBDA **rules** also permit Description Logic Programs (OWL 2 RL), Datalog[±], and Disjunctive Datalog. [Semantics of **ontology languages** customizable for expressivity and efficiency requirements by adding/deleting **rules** (SPIN)]
3. ...

RBDA Realizes Uniform KBDA – 3 of 3: Mappings as Rules

1. ...
2. ...
3. KBDA data integration is centered on Global-As-View (GAV) *mappings*, which are Datalog **rules** for, e.g., *unfolding* each global head predicate to (a join, i.e. conjunction, of) local body predicates

Example: Forest/Orchard Knowledge



EntityWithTree KB: Named Root Class (1)

Subsumption axioms (in higher-order rule syntax):

EntityContainingAtLeastOneTree ← Forest.

EntityContainingAtLeastOneTree ← Orchard.

Forest ← Woodland.

“←” is
taxonomy-style
'subsumes' infix

EntityContainingAtLeastOneTree

Root of taxonomy tree of
tree-containing entities to
see the forest for the trees

Forest

Orchard

Woodland

EntityWithTree KB: Named Root Class (2)

Subsumption axioms (in higher-order rule syntax):

EntityContainingAtLeastOneTree :- Forest.

EntityContainingAtLeastOneTree :- Orchard.

Forest :- Woodland.

“:-” is
rule-style
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EntityContainingAtLeastOneTree

Root of taxonomy tree of
tree-containing entities to
see the forest for the trees

Forest

Orchard

Woodland

EntityWithTree KB: Constructed Root Class

Subsumption axioms (in higher-order rule syntax):

$\exists \text{contains.Tree} \text{ :- Forest.}$

$\exists \text{contains.Tree} \text{ :- Orchard.}$

$\text{Forest :- Woodland.}$

Cf. ontology-style (description logic) axioms:

$\exists \text{contains.Tree} \equiv \text{Forest}$

$\exists \text{contains.Tree} \equiv \text{Orchard}$

$\exists \text{contains.Tree}$

Entities each having a
contains property with
at least one value in class Tree

Forest

Orchard

Woodland

Three Dimensions of KBDA_s : R, Q, m

R B D Q mediator

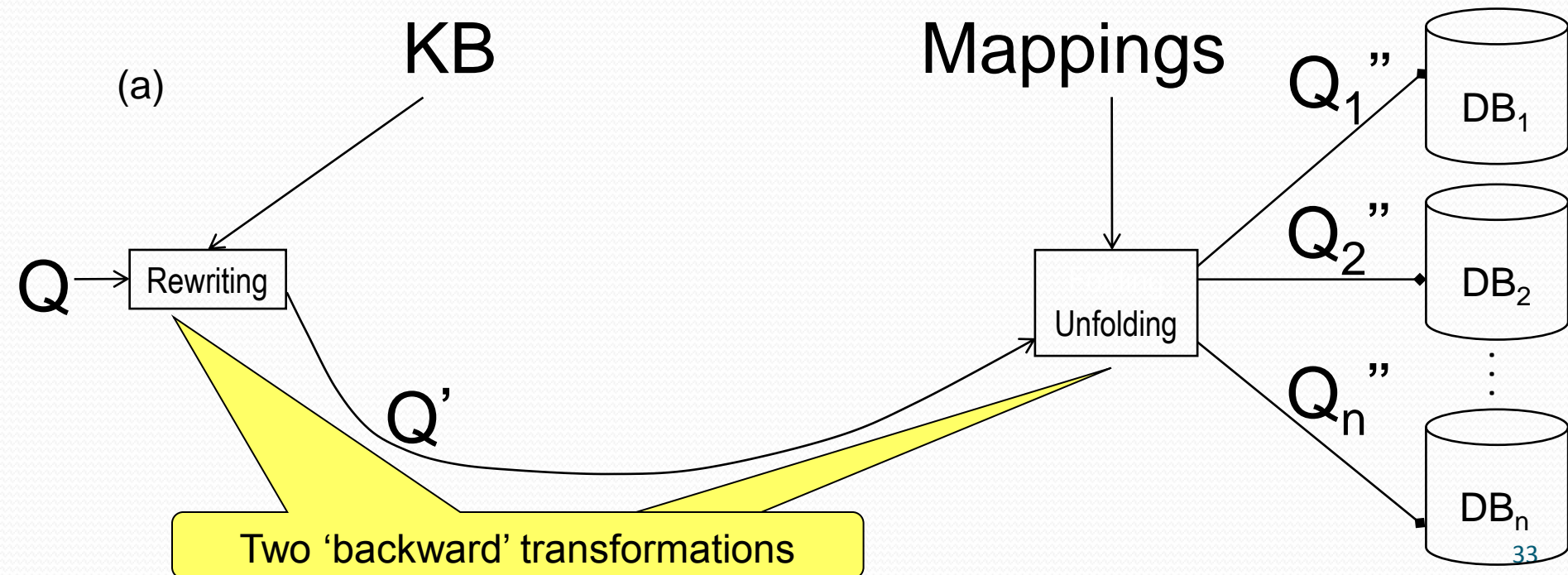
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e -Based Data Querying

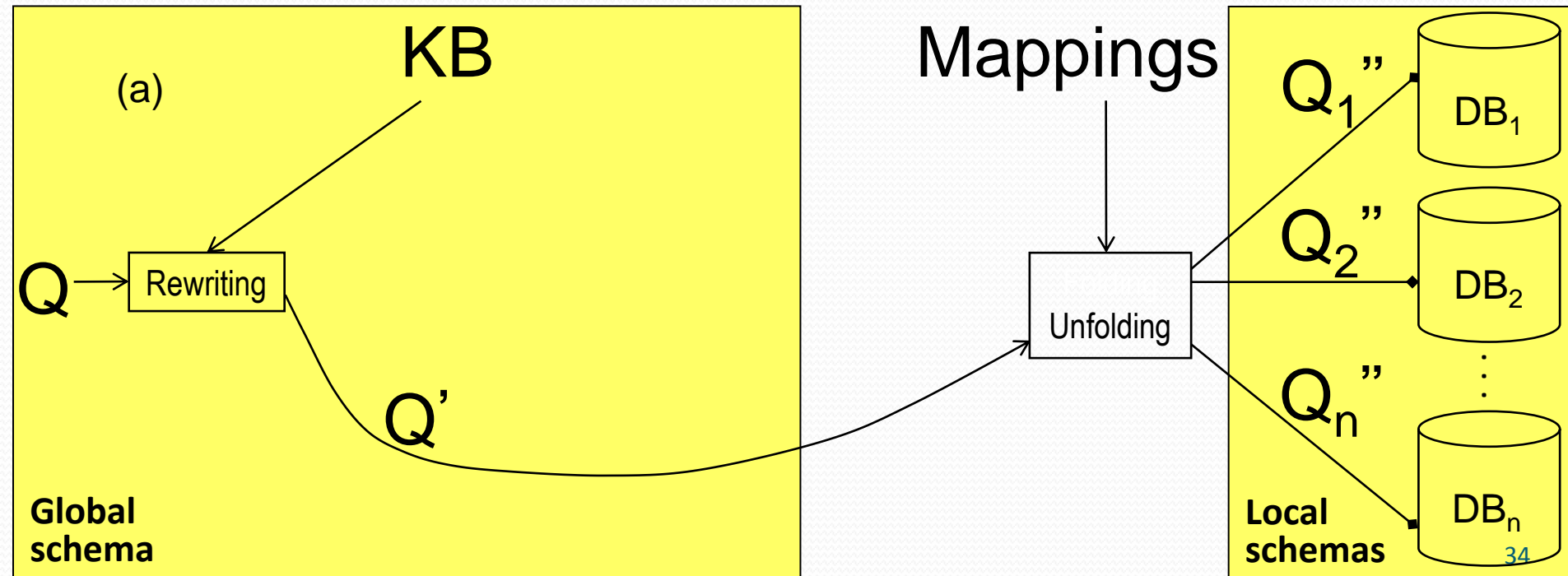
Query Rewriting and Unfolding

- ❖ Mediator strategy uses:
KB to *rewrite* Q to Q' and Mappings (**rules**) to *unfold* Q' to Q_i''
- ❖ KB can be **ontology**, e.g. in OWL 2 QL (DL-Lite), or **rules**
- ❖ Abstract (relational/graph/...) queries Q_i'' ♦-grounded (to SQL/SPARQL/...) for DB_i
- ❖ Each (relational/graph/...) database DB_i left as original; answers at ♦



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Query Rewriting Use of KB

- In Information Retrieval:

Query expansion

Intuitive idea of **query rewriting**

- With increased *recall*
- Without loss of *precision*

- From Logic Programming (for Horn expressivity):

Can use resolution method

for KB-enrichment of a given (conjunctive) query
with expanded (conjunctive) queries

so that, for any DB,

the answers to the enriched queries no longer using the KB
are the same as the answers to the original query using the KB

Ontology-to-Rule Clausification of EntityWithTree KB Omitting Orchard

Description Logic subsumptions (as higher-order rules):

$\exists \text{contains.Tree} \text{ :- Forest.}$
 $\text{Forest :- Woodland.}$

Higher-order syntactic sugar for existential rule:
 $\exists ?y(\text{contains}(?x ?y) \wedge \text{Tree}(?y)) \text{ :- Forest}(?x).$

Horn Logic rules (the first with conjunctive head):

$(\text{contains}(?x s(?x)) \wedge \text{Tree}(s(?x))) \text{ :- Forest}(?x).$
 $\text{Forest}(?x) \text{ :- Woodland}(?x).$

s is fresh Skolem function symbol

Horn Logic rules (the first head split into two conjuncts):

$\text{contains}(?x s(?x)) \text{ :- Forest}(?x).$
 $\text{Tree}(s(?x)) \text{ :- Forest}(?x).$
 $\text{Forest}(?x) \text{ :- Woodland}(?x).$

KB Rules Perform Rewriting of Given Query

KB with Horn rules (from above) for rewriting of query rules:

$\text{contains}(\text{?x } s(\text{?x})) \text{ :- Forest}(\text{?x}).$

$\text{Tree}(s(\text{?x})) \text{ :- Forest}(\text{?x}).$

$\text{Forest}(\text{?x}) \text{ :- Woodland}(\text{?x}).$

Rewriting Datalog query rule to obtain extra query rules:

Q: Given

$q(\text{?z}) \text{ :- contains}(\text{?z } \text{?y}) \wedge \text{Tree}(\text{?y}).$

$q(\text{?z}) \text{ :- Forest}(\text{?z}) \wedge \text{Tree}(s(\text{?z})) \quad q(\text{?z}) \text{ :- contains}(\text{?z } s(\text{?x})) \wedge \text{Forest}(\text{?x})$

$q(\text{?z}) \text{ :- Forest}(\text{?z}) \wedge \text{Forest}(\text{?z}).$

$q(\text{?z}) \text{ :- Forest}(\text{?z}).$

Expansion

$q(\text{?z}) \text{ :- Woodland}(\text{?z}).$

Q': Given \cup Expansion

Query Unfolding Use of Mappings to Original Database Sources

- Datalog rules **bridging** between:

- KB
- Distributed DBs

Intuitive idea of **query unfolding**

- Use partial deduction-like unfolding (and simplification) of (conjunctive) KB queries to (conjunctive) *abstract DB queries*
 - Abstract relational queries grounded to SQL, abstract graph queries grounded to SPARQL, etc.
 - Lower-level optimization and execution by SQL, SPARQL, etc. engines
- Generated queries distributed over multiple DBs as indicated by “*source.*” name prefixes

Sample Mapping Rules to Three Local Data Sources

Map KB predicates to locDB/regionDB tables for geo data:

`contains(?x ?y) :- locDB.cnt(?x ?kind ?y).`

`contains(?x ?y) :- regionDB.sub(?x ?r) \wedge locDB.cnt(?r ?kind ?y).`

Map KB predicates to locDB/ecoDB tables for forestry data:

`Tree(?t) :- locDB.cnt(?plot "tree" ?t).`

`Tree(?t) :- ecoDB.Plant(?plot "tree" ?size ?t).`

`Forest(?x) :- ecoDB.Habitat(?plot "forest" ?size ?x).`

`Woodland(?x) :- ecoDB.Habitat(?plot "wood" ?size ?x).`

Mapping Rules Perform Unfolding of Rewritten Queries

Union of conjunctive queries as Datalog rules (rewritten):

$q(?z) :- \text{contains}(?z ?y) \wedge \text{Tree}(?y).$

$q(?z) :- \text{Forest}(?z).$

$q(?z) :- \text{Woodland}(?z).$

Q': Given \cup Expansion

Unfolding above queries via mappings from previous slide:

$q(?z) :- \text{locDB.cnt}(?z ?kind ?y) \wedge \text{locDB.cnt}(?plot \text{"tree"} ?y).$

$q(?z) :- \text{locDB.cnt}(?z \text{"tree"} ?y).$

Simplification: Integrity rules

$q(?z) :- \text{locDB.cnt}(?z ?kind ?y) \wedge \text{ecoDB.Plant}(?plot \text{"tree"} ?size ?y).$

$q(?z) :- \text{regionDB.sub}(?z ?r) \wedge \text{locDB.cnt}(?r ?kind ?y) \wedge \text{locDB.cnt}(?plot \text{"tree"} ?y).$

$q(?z) :- \text{regionDB.sub}(?z ?r) \wedge \text{locDB.cnt}(?r \text{"tree"} ?y).$

$q(?z) :- \text{regionDB.sub}(?z ?r) \wedge \text{locDB.cnt}(?r ?kind ?y) \wedge \text{ecoDB.Plant}(?plot \text{"tree"} ?size ?y).$

$q(?z) :- \text{ecoDB.Habitat}(?plot \text{"forest"} ?size ?z).$

$q(?z) :- \text{ecoDB.Habitat}(?plot \text{"wood"} ?size ?z).$

Q_i': Bold-faced
($1 \leq i \leq 3$)

Three Dimensions of $KBD A_s$: R, Q, w

R B D Q warehouse

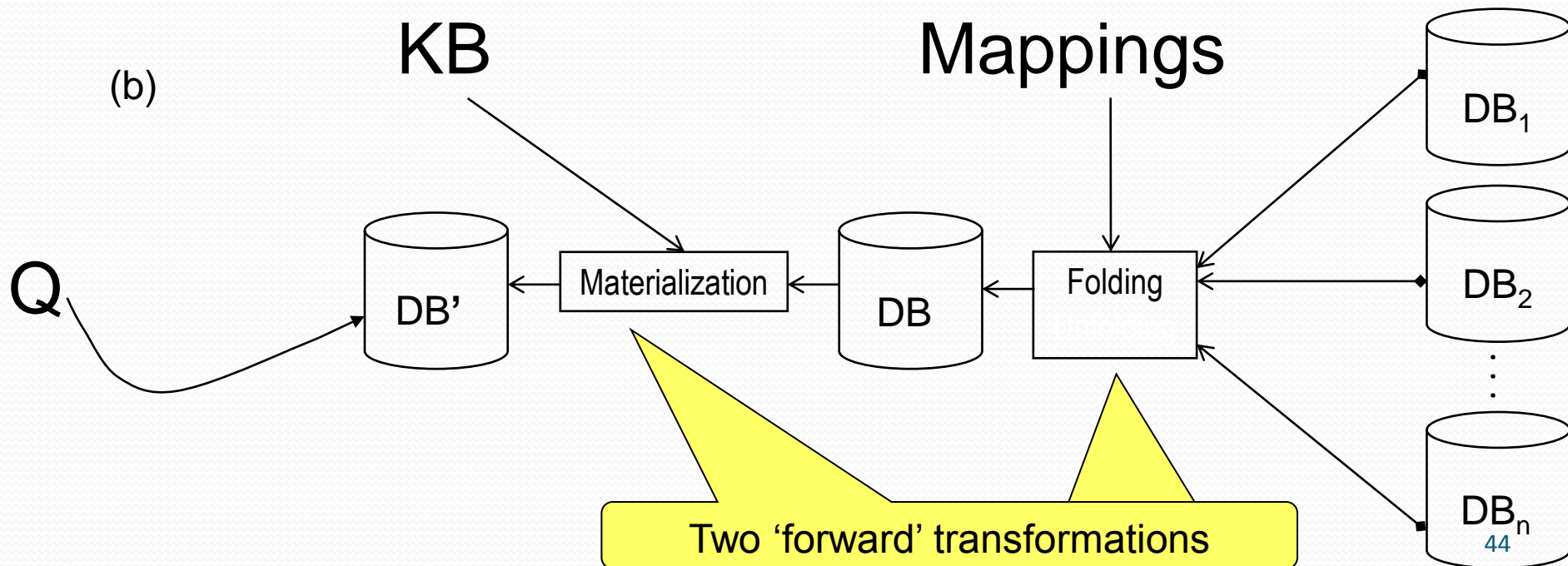
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e -Based Data Querying

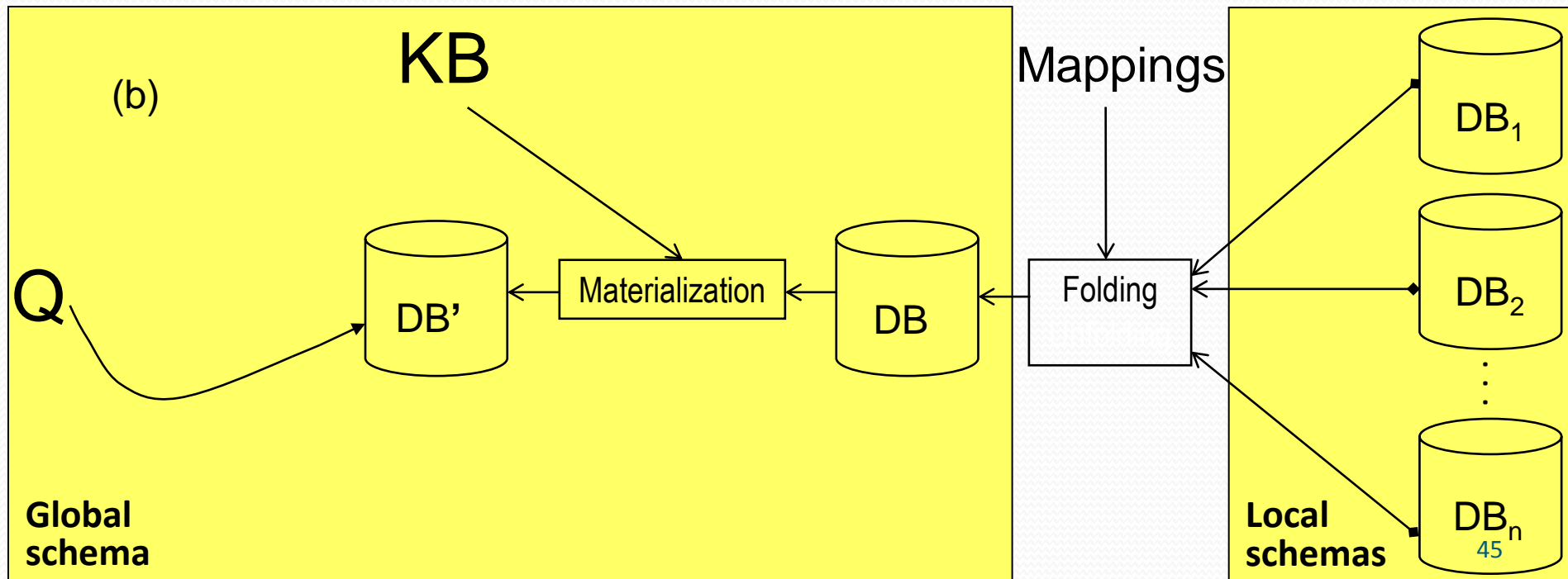
Database Materialization after Folding

- Warehouse strategy uses:
Mappings (same **rules** as for *unfolding*) to *fold* DB_i to DB and KB to *materialize* Database DB to DB'
- KB can be **ontology**, e.g. in OWL 2 RL (DLP), or **rules**
- Query is left as original; answers at solid triangular arrow head



Database Materialization after Folding

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EWT KB and DB: 'Populated' KB

Subsumptions and Data (as higher-order rules and facts):

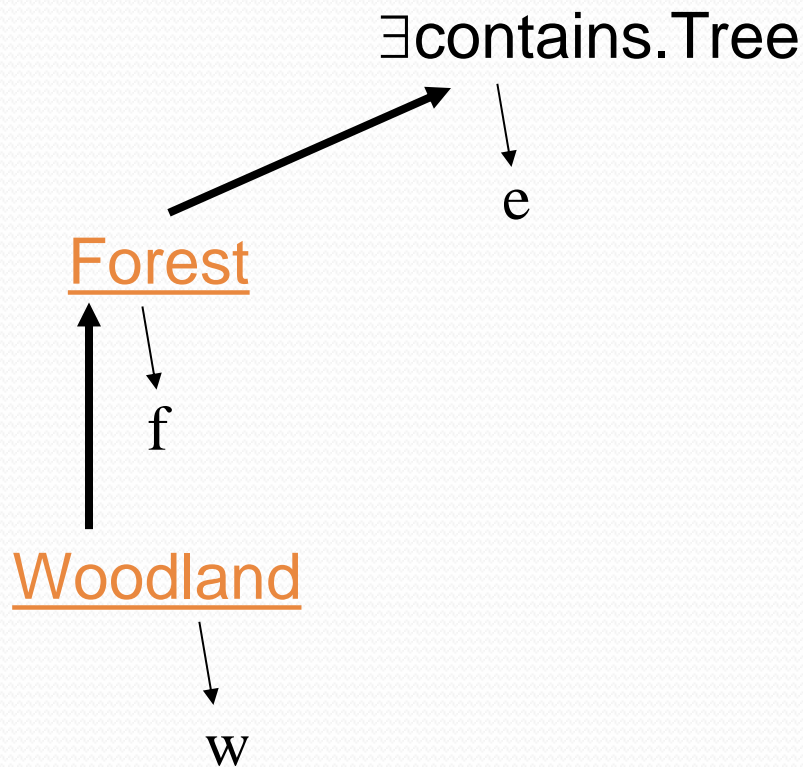
$\exists \text{contains.Tree} \text{ :- Forest.}$

$\text{Forest} \text{ :- Woodland.}$

$\exists \text{contains.Tree}(e).$

$\text{Forest}(f).$

$\text{Woodland}(w).$



Applying All KB Rules to All DB Facts (1)

Subsumptions and Data (as higher-order rules and facts):

$\exists \text{contains.Tree} \text{ :- Forest.}$

$\text{Forest} \text{ :- Woodland.}$

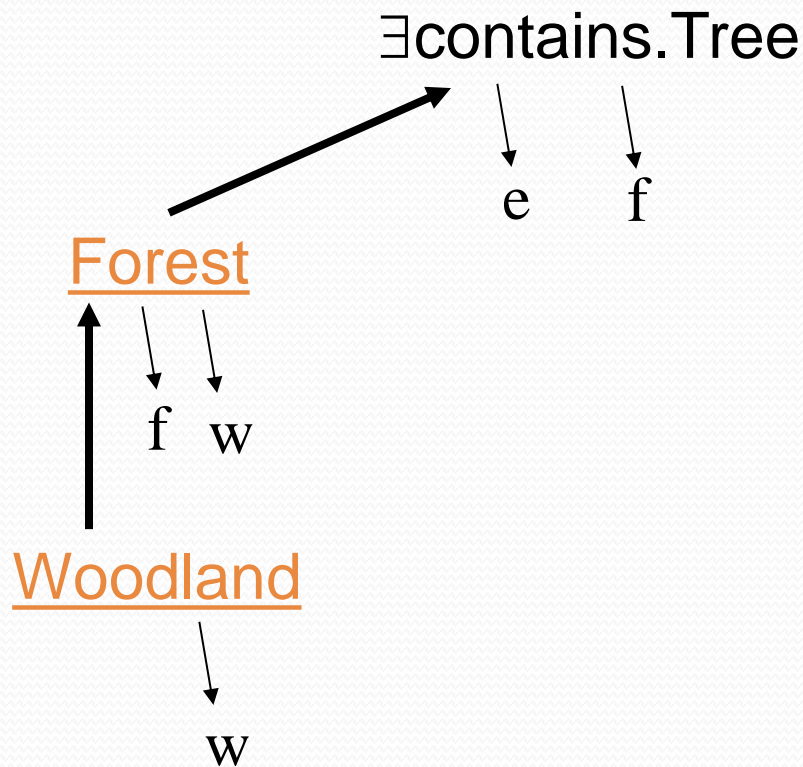
$\exists \text{contains.Tree}(e).$

$\exists \text{contains.Tree}(f).$

$\text{Forest}(f).$

$\text{Forest}(w).$

$\text{Woodland}(w).$



Applying All KB Rules to All DB Facts (2)

Subsumptions and Data (as higher-order rules and facts):

$$\exists \text{contains.Tree} :- \text{Forest.}$$

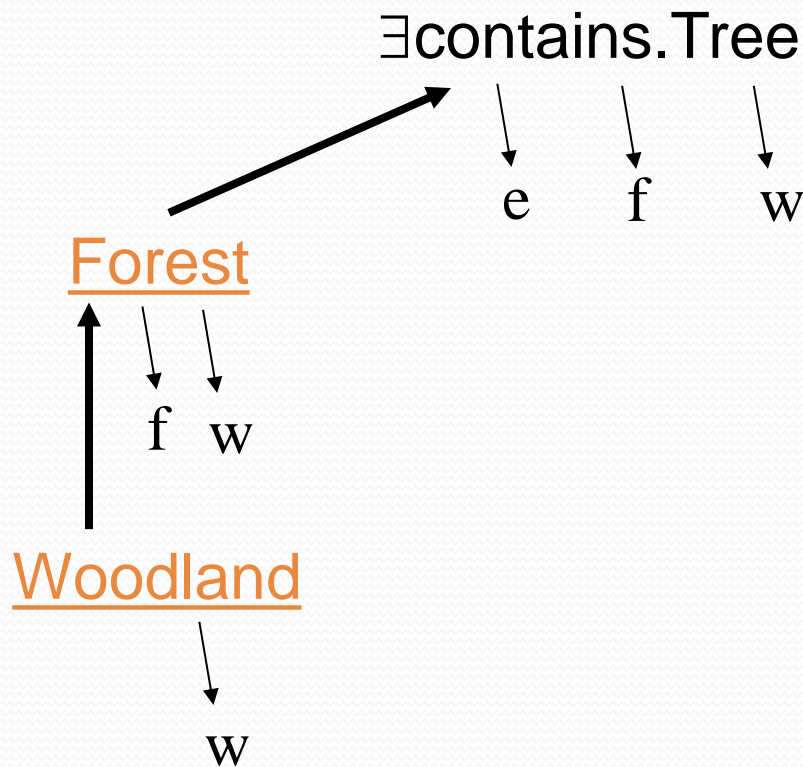
Forest :- Woodland.

$$\exists \text{contains.Tree}(e).$$
$$\exists \text{contains.Tree}(f).$$
$$\exists \text{contains.Tree}(w).$$

Forest(f).

Forest(w).

Woodland(w).



Fixpoint with No New Rule-Derived Facts

Subsumptions and Data (as higher-order rules and facts):

$\exists \text{contains.Tree} \text{ :- Forest.}$

$\text{Forest} \text{ :- Woodland.}$

$\exists \text{contains.Tree}(e).$

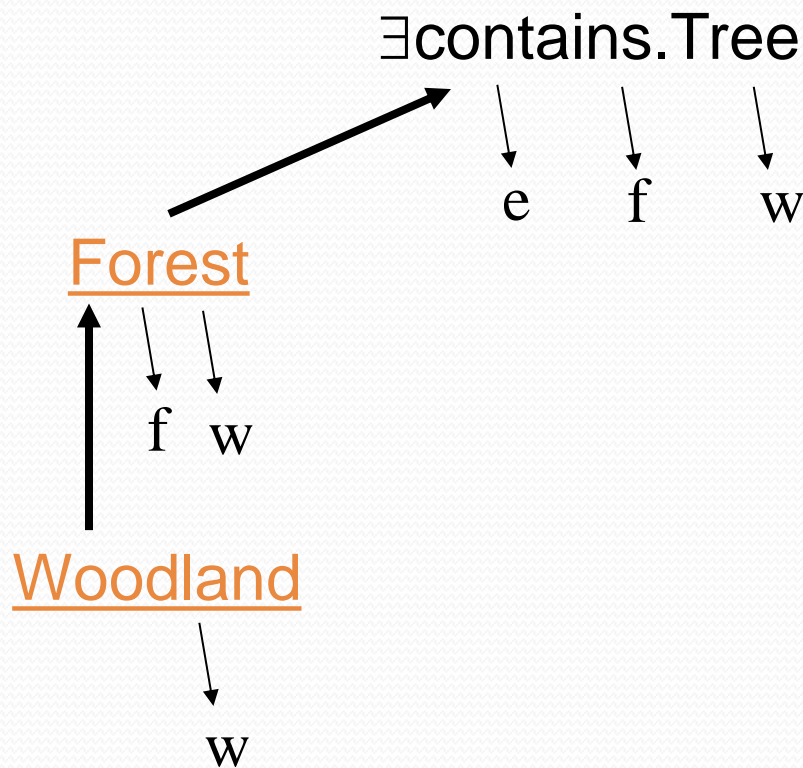
$\exists \text{contains.Tree}(f).$

$\exists \text{contains.Tree}(w).$

$\text{Forest}(f).$

$\text{Forest}(w).$

$\text{Woodland}(w).$



Three Dimensions of KBDA_s

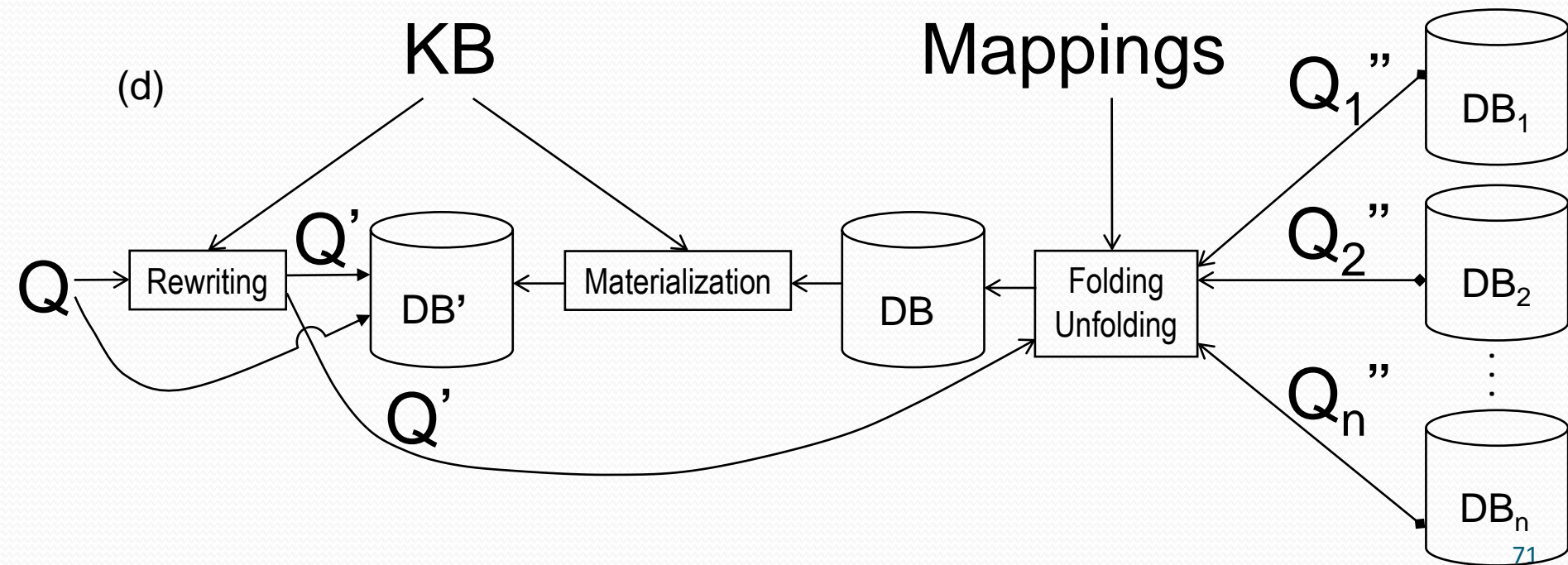
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KBDA_s strategy

-Based Data Access

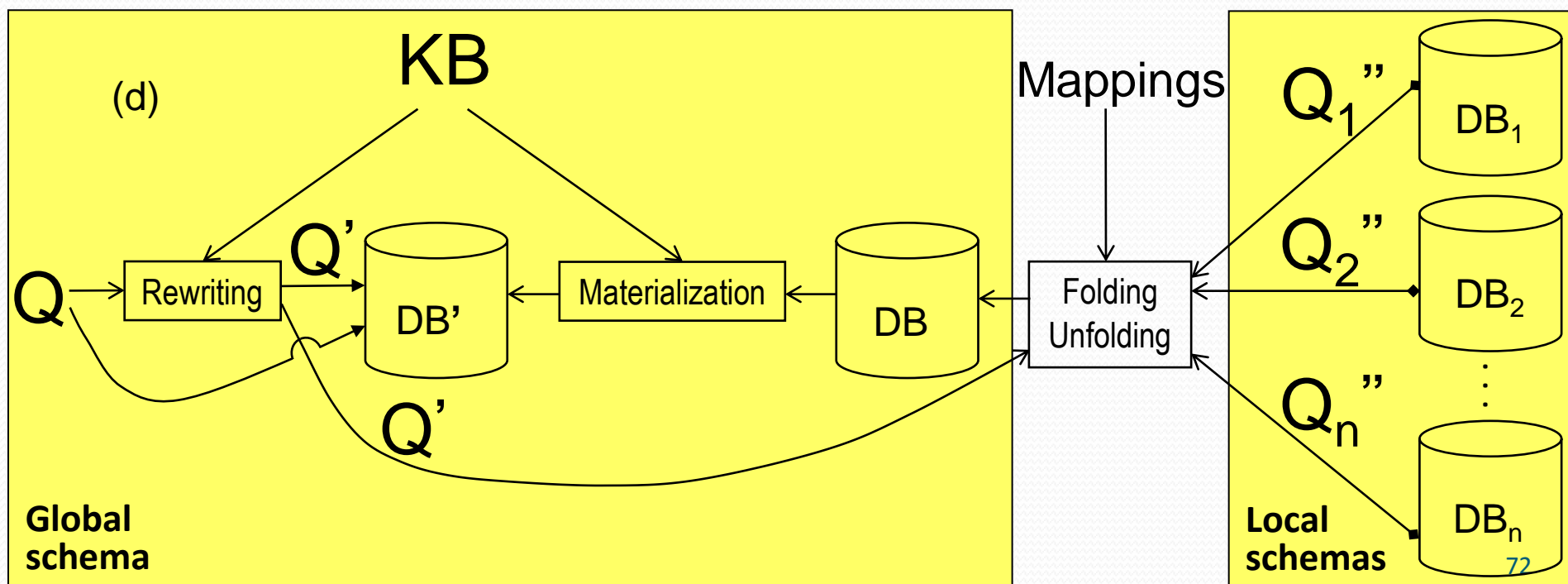
Unified Architecture

- Combines strategies (a)-(c) of earlier slides
- Meets the needs of Δ Forest case study



Unified Architecture

- Combines strategies (a)-(c) of earlier slides
- Meets the needs of Δ Forest case study



RBDA_s -Style KBDA_s Architecture: Expressivity of Rule Systems

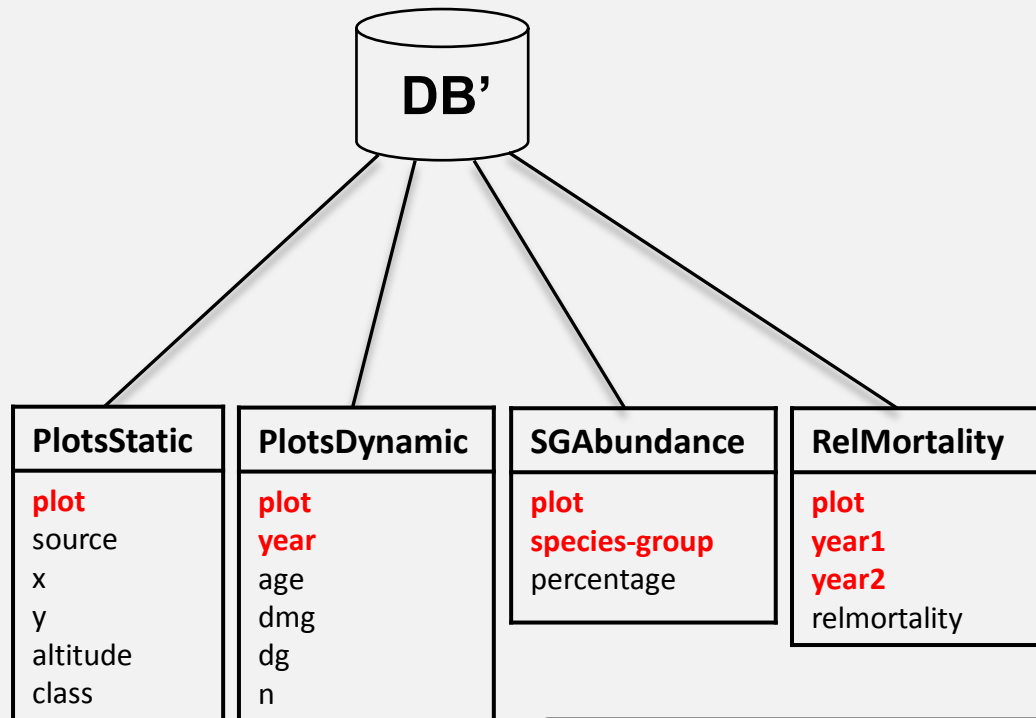
- The language of the global schema can be generalized from unary/binary (OBDA_s) to n-ary predicates (RBDA_s)
- When decidability of querying is not required, RBDA_s expressivity can be extended from Datalog, Datalog[±], and description logic to Datalog⁺, Horn logic, and FOL, as enabled by [Deliberation RuleML 1.01](#),
 - Features customizable with the [MYNG 1.01](#) GUI
- Moreover, [Reaction RuleML 1.0](#) can express updates, as needed for KBDM_s ([Ontology-based Data Management](#))

RBDAs-Style KBDA_s Architecture: Uniformity via Rule Systems

- Rule-based style of Unified Architecture (earlier slide)
- Presentation syntax (“:-”), serialization (“<RuleML>”), and semantics approach (model theory) uniform from *queries (+ integrity constraints)* to *KBs* to *mappings* to *abstract DBs*
- Division of labor between KB rules and mapping rules can be modified *without crossing paradigm boundaries*
 - Allows KB- and mapping-directed normal forms
- Assumptions (unique-name and closed-world) of DBs accommodated by *default assumptions of rule systems*

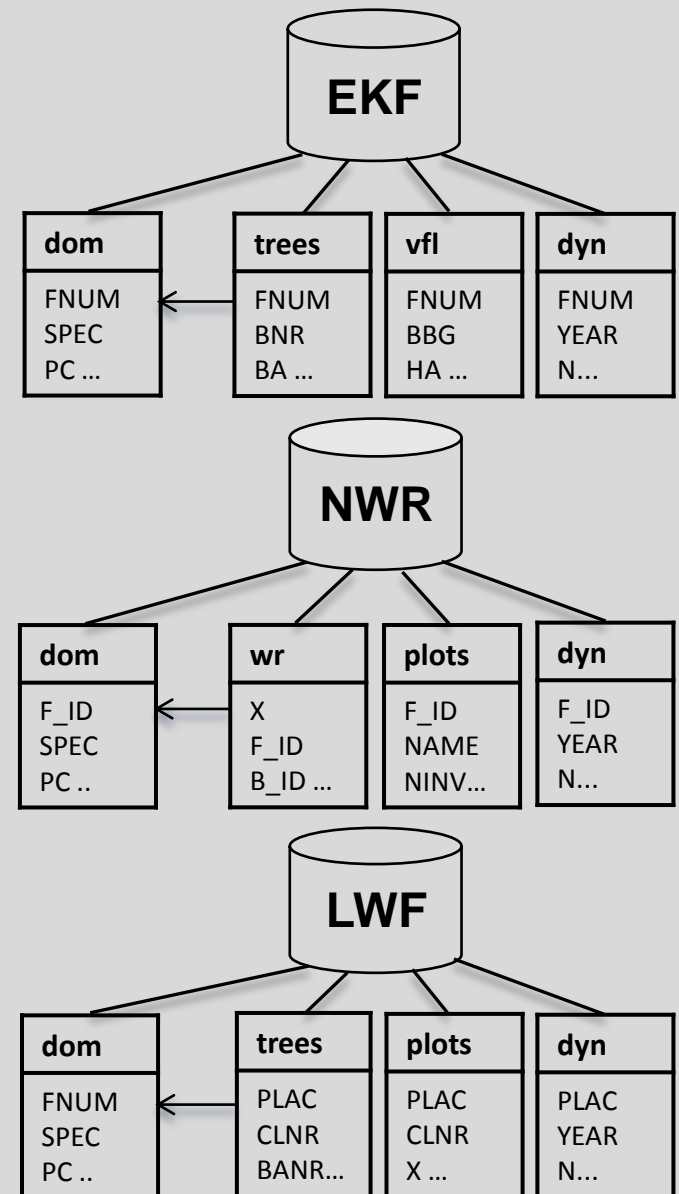
Global schema

ΔForest: Schemas for Architecture (a)-(d)



Keys – three composite – in **bold red**

Local schemas



Questions Addressed

1. Are there sufficiently many eligible plots in order to perform an analysis per main tree species?
2. Are there sufficiently many eligible plots in order to perform an analysis per main tree species **and** climatic region?
3. Which eligible plots represent pure tree stands and which eligible plots represent mixed tree stands?

Query Rewriting

**q(?plot) :- .EligiblePlot(?plot)
 .TreeStandKey(?id ?plot "oak")
 .TreeStandAbundance(?id ?pct)
 ?pct >= 15.**

**Exists ?id (.TreeStandKey(?id ?plot ?sp) .TreeStandAbundance(?id ?pct)) :-
 .TreeStandMerged(?plot ?sp ?pct).**

▶ **q(?plot) :- .EligiblePlot(?plot)
 .TreeStandMerged(?plot "oak" ?pct)
 ?pct >= 15.**

Query

KB Rule

Rewritten

Input

Output

Query Unfolding

**q2(?plot) :- .TreeStandMerged(?plot "oak" ?pct)
 ?pct **>= 15.****

**.TreeStandMerged(?plot "oak" ?pct) :-
 EKF.dom(?plot "Quercus petraea" ?pct1)
 EKF.dom(?plot "Quercus robur" ?pct2)
 ?pct = ?pct1 + ?pct2.**

▶ **q2(?plot) :- EKF.dom(?plot "Quercus petraea" ?pct1)
 EKF.dom(?plot "Quercus robur" ?pct2)
 ?pct1+?pct2 **>= 15.****

Query 2

Mapping
Rule

Unfolded

Input

Output

Conclusions

- Ontology-Based Data Access (OBDA) founded on three kinds of rules: *Query rules* (including integrity rules), *KB rules* (for query rewriting and DB materialization), as well as *mapping rules* (for query unfolding and DB folding)
- OBDA complemented by Rule-Based Data Access (RBDA) and generalized to Knowledge-Based Data Access (KBDA)
- Specified an RBDA-uniform KBDA_s architecture with unified mediator, warehouse, and bidirectional strategies
- RuleML used for XML-serialized rules, MYNG-customized rule expressivity, and platform-independent RBDA
- Introduced Δ Forest specialization of RBDA architecture for statistical data analysis in ecosystem research at WSL

Future Work (1)

- Translate simplified presentation syntax into released XML serialization of Deliberation RuleML 1.01 / MYNG 1.01
- Support implementations of specified architecture reusing (open source) KBDA technology (cf. RBDA wiki page)
- For high-precision RBDQ_s language support, complement current techniques of Datalog⁺ RuleML for Datalog[±] RuleML using context-sensitive/semantic validators for “-” constraints
- Evaluate (mediator/warehouse, relational/graph, ...) trade-offs for KBDQ_s in PSOA RuleML as executed in PSOATransRun
- Develop Δ Forest study at WSL for extended and new data sources of big volume, variety, and velocity (e.g., about climate change)
- Augment geospatial KBDQ_s mappings with Optique mapping (bootstrapping, repair, ...) techniques

Future Work (2)

- Compare engines for OBDQ_s and RBDQ_s , including HYDRA and RDFox, w.r.t. expressivity and efficiency
- Adapt *Semantic Automated Discovery and Integration* (SADI) test cases for KBDQ_s experiments in PSOA RuleML querying
- Evaluate *Abstract Logic-based Architecture Storage systems & Knowledge base Analysis* (ALASKA) for RBDQ_s
- Extend the KBDQ_s architecture with semantic annotation rules for (Deep) Web data extraction (Deep Web Mediator)
- Use Grailog for KBDQ_s data and knowledge visualization
- Explore synergies between the logical KBDQ_s approach with statistical approaches, e.g. from Statistical Relational AI