Using Discriminative Rule Mining to Discover Declarative Process Models with Non-Atomic Activities

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Outline

- Motivations
- 2 Backgrounds
- 3 The Approach
- 4 Case Study and Results
- **5** Conclusions and Future Work

Motivations

- Process discovery techniques are widely considered as critical for successful business process management and monitoring;
- The discovery of declarative models can be used in complex environments where process executions involve multiple alternatives and high flexibility is needed;
- Existing process discovery techniques for generating declarative specifications, do not take activity lifecycles and their characteristics into consideration.

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- Declare is a language for describing declarative process models consisting of a set of constraints applied to (atomic) activities.
- Constraints, in turn, are based on templates that are abstract parameterized patterns: constraints are their concrete instantiations on real activities.
- Template semantics can be formalized using different logics, the main one being Linear Temporal Logic over finite traces, making them verifiable and executable

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TEMPLATE	FORMALIZATION	NOTATION
existence(A)	$\Diamond A$	1* A
absence(A)	$\neg \Diamond A$	о А
choice(A,B)	$\lozenge A \vee \lozenge B$	A> B
exclusive choice(A,B)	$(\Diamond A \vee \Diamond B) \wedge \neg (\Diamond A \wedge \Diamond B)$	A -◆- B
responded existence(A,B)	$\Diamond A \rightarrow \Diamond B$	A B
response(A,B)	$\Box(A \to \Diamond B)$	A → B
precedence(A,B)	$\neg B \mathcal{W} A$	A → B
alternate response(A,B)	$\Box(A \to \operatorname{O}(\neg A \operatorname{\mathcal{U}} B))$	A → B
alternate precedence(A,B)	$(\neg B \mathcal{W} A) \wedge \Box (B \to \bigcirc (\neg B \mathcal{W} A))$	A → B
chain response(A,B)	$\Box(A \to \bigcirc B)$	A B
chain precedence(A,B)	$\Box(\bigcirc B \to A)$	A <u>→ B</u>
not resp. existence(A,B)	$\Diamond A \to \neg \Diamond B$	A
not response(A,B)	$\Box(A \to \neg \Diamond B)$	A → → B
not precedence(A,B)	$\Box(A \to \neg \Diamond B)$	A ▶ B

The Declare Language Activity Lifecycle Discriminative Mining

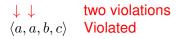
TEMPLATE	FORMALIZATION	NOTATION
response(A,B)	$\Box(A \to \Diamond B)$	A → B

 $\langle a,a,b,c \rangle$ Satisfied $\langle b,b,c,d \rangle$ Satisfied (vacuously) $\langle a,b,c,b \rangle$ Satisfied $\langle a,b,a,c \rangle$ Not Satisfied

TEMPLATE	FORMALIZATION	NOTATION
response(A,B)	$\Box(A \to \Diamond B)$	A → B

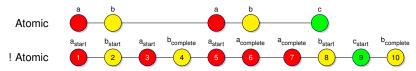
$$\begin{array}{ccc} \downarrow & & \text{one violation} \\ \langle a,a,b,c \rangle & & \text{Violated} \end{array}$$

TEMPLATE	FORMALIZATION	NOTATION
response(A,B)	$\Box(A \to \Diamond B)$	A → B



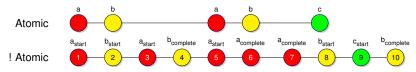
TEMPLATE	FORMALIZATION	NOTATION
response(A,B)	$\Box(A \to \Diamond B)$	A → B

$$\begin{array}{ccc} \downarrow & \downarrow & \downarrow & \text{2 fullfillments} \\ \langle a, a, b, c \rangle & \text{Satisfied} \end{array}$$



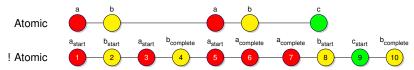
Activity Lifecycle

- In the reality, activities have a duration spanning over time intervals in which transactional states (a.k.a. event types) of the activity can occur;
- The sequences of event types that occur when an activity is executed, characterize the lifecycle of that activity;
- If available, event type data is very important and should be considered: it allows to understand the constraints between event types inside activities.



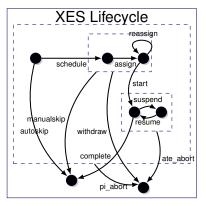
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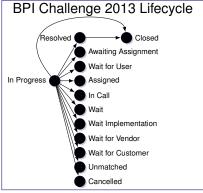
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- Discriminative mining was exploited to highlight patterns that are discriminative with respect to a given criterion from the existing set of constraints activations over an event logs;
- In order to mine the set of declarative rules that discriminate between fullfillmens and violations, decision tree supervised learning
- Decision tree learning uses a decision tree as a model to predict the value of a target variable based on input variables (features) and are built from a set of training dataset;
- Internal nodes of the tree are labeled with input features and edges with possible value ranges of the feature;
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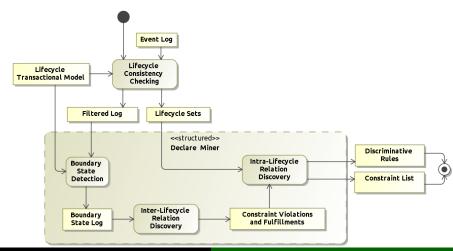
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The Process

Lifetime Consistency Checking Boundary State Detection Discovering inter-lifecycle relations Discovering intra-lifecycle relations

The Process



Lifetime Consistency Checking



The proposed technique aims at processing the input log to connect together activity states belonging to the same lifecycle; this can be done with 2 approaches:

- **FIFO** applying our FIFO-based algorithm, we can disambiguate the correlation using a conservative approach: $a_{complete}$ at pos. 6 can be connected to a_{start} at different positions but we connect it to the occurrence that occurred first (hence position 3) generating lifecycle 1-3-6.
- Data correlation if the event log contains (event) data attributes, it is possible to connect activity states that share some data values, e.g., an event id

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- This step is to replace activity lifecycles with placeholders marking the start and the end of each lifecycle in the log.
- The formulas directly follow by the corresponding formulas in standard Declare.
- The idea is that, for verifying constraints involving non-atomic activities, it is sufficient to take into account the boundary states of lifecycles, abstracting away from the lifecycle detail;
- From the Filtered Log and the Lifecycle specification this step generates a simplified log in which only boundary states are included.

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Discovering inter-lifecycle relations



- Boundary state log derived from the previous step is mined using the existing standard Declare mining approach presented in a previous work:
- Declare template semantics has been slightly modified for non-atomic activities (in the template the initial state is indicated with i and the final one with f):
- The outcome is a set of candidate Declare constraints connecting elements of different life-cycles (inter-lifecycle relations). For each candidate, we extract the set of fulfillments and the set of violations in the log; ad es:

 $b_{complete}(4)$ is a fulfillment for constraint $\Box(b_f \to \Diamond a_i)$ $b_{complete}(10)$ is a violation for the same constraint.

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- The problem is then riformulated as a supervised learning problem: the features of the lifecycle of activation a discriminating with respect to fullfillments/violations of constraint constr, are learned from a set L of sequences representing all the lifecycles of activation a in the log:
- Sequences **are classified in two sets** L_{ful} and L_{viol} according to whether the lifecycle corresponds to a fulfillment or to a violation of *constr*:
- Each lifecycle is encoded in terms of a set of Declare constraints as a vector of (boolean) values representing whether all these (intra-lifecycle) constraints are satisfied or not on that lifecycle. The decision tree is trained using the intra-lifecycle conditions as features and the classification of the lifecycle as part of L_{ful} or L_{viol} .

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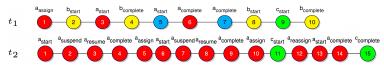
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Discovering intra-lifecycle relations



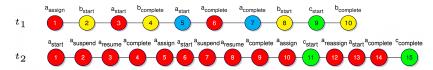
For example, consider constraint $\Box(a_f \to \Diamond b_i)$ over t_1 and t_2 : $a_1 = \langle a_{start}, a_{complete} \rangle$ $a_2 = \langle a_{assign}, a_{start}, a_{complete} \rangle$ in trace \mathbf{t}_1 correspond to **fulfillments** for the constraint (followed by b_i).

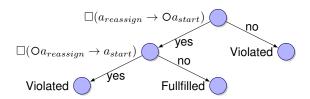
 $a_3 = \{\langle a_{start}, a_{suspend}, a_{resume}, a_{complete} \rangle, a_4 = \langle a_{assign}, a_{start}, a_{suspend}, a_{resume}, a_{complete} \rangle$, and $a_5 = \langle a_{assign}, a_{reassign}, a_{start}, a_{complete} \rangle$ in trace \mathbf{t}_2 correspond to **violations**.

Therefore $L_{ful} = \{a_1, a_2\}$ and $L_{viol} = \{a_3, a_4, a_5\}$ and hence it is likely that the constraint is verified when a_{start} is immediately followed by $a_{complete}$ but is not immediately preceded by $a_{reassign}$ in the lifecycle of a.

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The Case Study Setup

- We have implemented the approach as a plug-in of the process mining tool ProM.
- The implemented plug-in has been applied to a set of execution logs synthetically generated (to verify its capability to capture known discriminating behaviors);
- To assess the approach it was applied to a real-life log (logs provided by the BPI Challenge 2013).

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Synthetic Log 1 Synthetic Log 1 contains 1000 traces in which *Register* occurs with one of the following possible lifecycles:

```
- \ Register_start, Register_complete \\ - \ Register_abort \\ - \ Register_assign, Register_start, Register_complete \\ - \ Register_start, Register_abort \\ - \ Register_start, Register_supend, Register_abort \\ \end{array}
```

- Whenever Register is aborted (see second, fourth and fifth lifecycles in the list), the claimer is notified via phone;
- If the registration completes normally (see first and third lifecycle in the list), the e-mail notification is required;
- Hence whenever Register is aborted, the response constraint $\Box(Register \rightarrow \Diamond Notify_by_phone)$ is verified, otherwise this constraint does not hold; the intra-lifecycle analysis confirms this giving the constraint exactly(1, $Register_{abort}$) as discriminative constraints within lifecycle of Register.

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- \ Register_start, Register_bort \\
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- When Send_questionnaire is withdrawn or aborted (second, third and fourth), Skip response for skipping the response is executed;
- Whenever Send_questionnaire completes normally (first and fifth), Skip_response is not executed.
- Therefore, in the log, the response constraint $\Box(Send_questionnaire \rightarrow \Diamond(Skip_response))$ is verified only if $Send_questionnaire$ does not complete normally and in this case intra-lifecycle gives: exclusive-choice($Send-questionnaire_{abort}$, $Send-questionnaire_{withe}$

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Synthetic Log 2 Synthetic Log 2 contains 1000 traces in which the non-atomic activity *Send_questionnaire* occurs with one of the following possible lifecycles:

```
 \begin{array}{l} < Send.questionnaire_{start}, Send.questionnaire_{complete} \rangle \\ < Send.questionnaire_{withdraw} \rangle \\ < Send.questionnaire_{usisign}, Send.questionnaire_{withdraw} \rangle \\ < Send.questionnaire_{start}, Send.questionnaire_{suspend}, Send.questionnaire_{abort} \rangle \\ < Send.questionnaire_{complete} \rangle \\ < Send.questionnaire_{complete} \rangle \\ \end{aligned}
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```
    High.medical.history.check<sub>withdraw</sub> \)
    High.medical.history.check<sub>start</sub>, High.medical.history.check<sub>complete</sub> \)
    High.medical.history.check<sub>start</sub>, High.medical.history.check<sub>abort</sub> \)
    High.medical.history.check<sub>assign</sub>, High.medical.history.check<sub>autoskip</sub> \)
    High.medical.history.check<sub>assign</sub>, High.medical.history.check<sub>start</sub>,
    High.medical.history.check<sub>complete</sub> \)
```

- When *High_medical_history_check* does not complete normally (third and fourth), *Contact_hospital* is executed eventually.
- When verification procedure completes normally (second and fifth), there is no need to contact the hospital.
- Therefore, in the log, the response constraint $\Box(High_medical_history_check \rightarrow \Diamond Contact_hospital)$ holds if and only in High_medical_history_check fails.

```
    High.medical.history_check_withdraw \rangle
    High.medical.history_check_start, High.medical.history_check_complete \rangle
    High.medical.history_check_start, High.medical.history_check_abort \rangle
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```

- For this to be true the intra-lifecycle relation ¬ alternate response between High_medical_history_check_start and High_medical_history_check_complete must hold
- The result is in line with the log intent: whenever, in the lifecycle of High_medical_history_check, there is a High_medical_history_check_start that is not followed by High_medical_history_check_complete, the hospital has to be contacted.

```
    - ⟨ High.medical.history.check<sub>withdraw</sub> ⟩
    - ⟨ High.medical.history.check<sub>start</sub>, High.medical.history.check<sub>complete</sub>⟩
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```

- For this to be true the intra-lifecycle relation ¬ alternate response between High_medical_history_check_start and High_medical_history_check_complete must hold.
- The result is in line with the log intent: whenever, in the lifecycle of High_medical_history_check, there is a High_medical_history_check_start that is not followed by High_medical_history_check_complete, the hospital has to be contacted.

- $\langle Register_{start}, Register_{suspend}, Register_{resume}, Register_{complete} \rangle$
- $\langle Register_{start}, Register_{suspend}, Register_{abort} \rangle$
- (Register_{start}, Register_{suspend}, Register_{resume}, Register_{suspend}, Register_{resume}, Register_{complete})
- (Register_{start}, Register_{suspend}, Register_{resume}, Register_{suspend}, Register_{withdraw})
- When Register is suspended and ends with an abort or a withdraw (second and fourth lifecycles), or is suspended (and resumed) more than once (third lifecycle in the list), the claimer has to be notified via phone;
- if there is only one suspension correctly resumed (first lifecycle in the list) the claimer can be notified via e-mail
- In these cases the response constraint $\Box(Register \rightarrow \Diamond Notify_by_phone)$ is verified, otherwise not.
- The discriminative rule discovered for the verification of this constraint is the alternate succession between Register_{resume} and Register_{complete}.

- $\langle Register_{start}, Register_{suspend}, Register_{resume}, Register_{complete} \rangle$
- $\langle Register_{start}, Register_{suspend}, Register_{abort} \rangle$
- (Register_{start}, Register_{suspend}, Register_{resume}, Register_{suspend}, Register_{resume}, Redister_{complete})
- (Register_{start}, Register_{suspend}, Register_{resume}, Register_{suspend}, Register_{withdraw})
- When Register is suspended and ends with an abort or a withdraw (second and fourth lifecycles), or is suspended (and resumed) more than once (third lifecycle in the list), the claimer has to be notified via phone:
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- $\langle Register_{start}, Register_{suspend}, Register_{resume}, Register_{complete} \rangle$
- $\langle Register_{start}, Register_{suspend}, Register_{abort} \rangle$
- (Register_{start}, Register_{suspend}, Register_{resume}, Register_{suspend}, Register_{resume}, Redister_{complete})
- (Register_{start}, Register_{suspend}, Register_{resume}, Register_{suspend}, Register_{withdraw})
- When Register is suspended and ends with an abort or a withdraw (second and fourth lifecycles), or is suspended (and resumed) more than once (third lifecycle in the list), the claimer has to be notified via phone:
- if there is only one suspension correctly resumed (first lifecycle in the list), the claimer can be notified via e-mail.
- In these cases the response constraint $\Box(Register \rightarrow \Diamond Notify_by_phone)$ is verified, otherwise not.
- The discriminative rule discovered for the verification of this constraint is the alternate succession between Register_resume and Register_complete.

- $\langle Register_{start}, Register_{suspend}, Register_{resume}, Register_{complete} \rangle$
- $\langle Register_{start}, Register_{suspend}, Register_{abort} \rangle$
- (Register_{start}, Register_{suspend}, Register_{resume}, Register_{suspend}, Register_{resume}, Register_{complete})
- $\ \langle \ Register_{start}, Register_{suspend}, Register_{resume}, Register_{suspend}, Register_{withdraw} \ \rangle$
- When Register is suspended and ends with an abort or a withdraw (second and fourth lifecycles), or is suspended (and resumed) more than once (third lifecycle in the list), the claimer has to be notified via phone:
- if there is only one suspension correctly resumed (first lifecycle in the list), the claimer can be notified via e-mail.
- In these cases the response constraint $\Box(Register \rightarrow \Diamond Notify_by_phone)$ is verified, otherwise not.
- The discriminative rule discovered for the verification of this constraint is the alternate succession between Register_{resume} and Register_{complete}.

Table 4: BPI 2013 Results.

Tuble 4. Bi i 2015 Results.			
INTER-LIFECYCLE RELATION	INTRA-LIFECYCLE RELATION	CLASS PROB.	SUPPORT
(1) precedence (Accepted,Completed)		0.65	3711
(2) precedence (Queued, Accepted)	$\begin{array}{l} \text{init}(Accepted_{Assigned}) \vee \\ \text{init} \left(Accepted_{Wait_implementation}\right) \end{array}$	0.75	92
(3) precedence (Queued, Completed)		0.8	4551
(4) response (Completed, Queued)		0.98	7570
(5) responded existence (Completed, Accepted)		0.66	3771
(6) responded existence (Completed, Queued)		0.8	4595

■ The first row of the table (1) suggests as discriminating behavior for *Completed* to be preceded by *Accepted* is when either *Completed*_{Cancelled} or *Completed*_{Closed} occurs in the lifecycle of *Completed*.

Table 4: BPI 2013 Results.

Tuble 4. Bi i 2015 Results.			
INTER-LIFECYCLE RELATION	INTRA-LIFECYCLE RELATION	CLASS PROB.	SUPPORT
(1) precedence (Accepted,Completed)		0.65	3711
(2) precedence (Queued, Accepted)	$\begin{array}{l} \text{init}(Accepted_{Assigned}) \vee \\ \text{init} \left(Accepted_{Wait_implementation}\right) \end{array}$	0.75	92
(3) precedence (Queued, Completed)		0.8	4551
(4) response (Completed, Queued)		0.98	7570
(5) responded existence (Completed, Accepted)		0.66	3771
(6) responded existence (Completed, Queued)		0.8	4595

■ The same rule is also discriminating for responded existence between Completed and Accepted (expressing that whenever Completed occurs, then, Accepted has to occur in the future or has already occurred before).

Table 4: BPI 2013 Results.

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INTER-LIFECYCLE RELATION	INTRA-LIFECYCLE RELATION	CLASS PROB.	SUPPORT
(1) precedence (Accepted,Completed)		0.65	3711
(2) precedence (Queued, Accepted)	$\begin{array}{l} \operatorname{init}(Accepted_{Assigned}) \vee \\ \operatorname{init}\left(Accepted_{Wait.implementation}\right) \end{array}$	0.75	92
(3) precedence (Queued, Completed)		0.8	4551
(4) response (Completed, Queued)		0.98	7570
(5) responded existence (Completed, Accepted	$\left \begin{array}{l} \text{not responded existence} \\ (Completed_{Cancelled}, Completed_{Closed}) \end{array} \right $	0.66	3771
$ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ \ $		0.8	4595

- In the second row (2), whenever the lifecycle of Accepted starts with Accepted Assigned or with Accepted Wait_Implementation, Accepted is preceded by Queued. These results are the ones with the lowest class probability and support. (the support of (2) indicates that the cases in which the constraint and the discovered discriminative rule are both verified are 92).
- \blacksquare On the other hand, the remaining rules present both a reasonable class probability (> 0.8) and a good support (> 4500).

Table 4: BPI 2013 Results.

Table 4: BFI 2013 Results.			
INTER-LIFECYCLE RELATION	INTRA-LIFECYCLE RELATION	CLASS PROB.	SUPPORT
$(1) \bigg {\tt precedence} (Accepted, Completed)$		0.65	3711
$(2) \bigg \text{precedence } (Queued, Accepted)$	$\begin{array}{l} \text{init}(Accepted_{Assigned}) \lor \\ \text{init}\left(Accepted_{Wait_implementation}\right) \end{array}$	0.75	92
(3) precedence (Queued, Completed)		0.8	4551
(4) response (Completed, Queued)		0.98	7570
(5) responded existence (Completed, Accepted	$ \begin{array}{c} \text{not responded existence} \\ (Completed_{Cancelled}, Completed_{Closed}) \end{array}$	0.66	3771
$(6) \bigg {\it responded existence} \ (Completed, Queued)$		0.8	4595

■ In particular, the **co-occurrence** of *Completed* I_{n_Call} and *Completed* $C_{ancelled}$ discriminates both on the precedence (3) and on the responded existence (6) between *Queued* and *Completed*, i.e., whenever both I_{n_Call} and *Cancelled* occur in the lifecycle of *Completed*, it means that *Completed* is preceded by *Queued* or, more in general (with a slightly higher support), that *Queued* occurs at least once before or after *Completed*.

Table 4: BPI 2013 Results.

Table 4. BFI 2013 Results.			
INTER-LIFECYCLE RELATION	INTRA-LIFECYCLE RELATION	CLASS PROB.	SUPPORT
(1) precedence (Accepted,Completed)		0.65	3711
(2) precedence (Queued, Accepted)	$ \operatorname{init}(Accepted_{Assigned}) \lor \operatorname{init}(Accepted_{Wait_implementation}) $	0.75	92
(3) precedence (Queued, Completed)		0.8	4551
(4) response (Completed, Queued)		0.98	7570
(5) responded existence (Completed, Accepted)		0.66	3771
(6) responded existence (Completed, Queued)	$ $ co-existence $ $ $(Completed_{In_call}, Completed_{Cancelled})$	0.8	4595

- Finally, the discovered discriminating behavior with the highest class probability (almost 1) and support (more than 7000) is the one related to the response constraint between *Completed* and *Queued* (6):
- in this case Queued eventually follows Completed if and only if only one among Completed_{Resolved} and Completed_{Closed} occurs in the lifecycle of Completed.

Conclusions and Future Work

- This paper presents an approach for the discovery of declarative process models from logs containing non-atomic activities.
- Discriminative rule mining is used to characterize the lifecycle of each constraint activation and discriminate between fulfillments and violations of the constraint under examination.
- Our experiments show the effectiveness of the approach and its applicability in real-life scenarios.
- As future work, we will conduct a wider experimentation of the proposed framework on several case studies in real-life scenarios and different transactional models for activity lifecycles.

Thank you for listening!

Any questions?