The PSOATransRun 1.0 System for Object-Relational Reasoning in RuleML

The 6th Atlantic Workshop on Semantics and Services (AWoSS 2015) December 9. 2015

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- Introduction
- Normalization of PSOA Source
- Mapping the Normalized PSOA Source to Prolog
- 4 Conclusions

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Introduction – PSOA RuleML

Uses positional-slotted object-applicative (psoa) terms, permitting the application of a predicate (acting as a relation) to be [in an oidless/oidful dimension] without or with an Object IDentifier (OID) – typed by the predicate (acting as a class) – and the predicate's arguments to be [in an orthogonal dimension] positional, slotted, or combined

General case (multi-tuple), where "#" means "member of":

```
o # f([t_{1,1} ... t_{1,n_1}] ... [t_{m,1} ... t_{m,n_m}] p_1 -> v_1 ... p_k -> v_k)
```

Special cases (optional single-tuple brackets):

```
 \begin{array}{lll} \mbox{Relationship:} & \mbox{f([t_1 \dots t_n])} \\ \mbox{Shelf:} & \mbox{o\#f([t_1 \dots t_n])} \\ \mbox{Frame:} & \mbox{o\#f(} & \mbox{p_1->v_1 \dots p_k->v_k)} \\ \mbox{Shelframe:} & \mbox{o\#f([t_1 \dots t_n])} & \mbox{p_1->v_1 \dots p_k->v_k)} \\ \mbox{} \end{array}
```

- For an oidless psoa term, i.e. one without a 'user' OID, objectification will introduce a 'system'-generated OID
- Psoa terms without class typing are expressed by Top, specifying the root of class hierarchy

PSOA RuleML – Presentation Syntax (PS)

Integrates W3C RIF for relationships and frames (see above)

As in RIF PS:

- "oid is member of class" written as "oid#class" (see above)
- "class₁ is subclass of class₂" written as "class₁## class₂"
- Local constants prefixed by underscore ('_'); variables prefixed by question mark ('?')

Introduction – PSOATransRun 1.0

- Efficient reasoning in PSOA RuleML enabled
 - All forms of psoa terms, including relationships, shelves, pairships, and (shel)frames
 - (OID-)head-existential rules
 - Equality in the body, restricted to unification and external-function evaluation
 - Subclass formulas for 'ABox' reasoning only
 - Built-in arithmetic functions
- Released in Java source form and as an executable jar file. Downloadable from: http://wiki.ruleml.org/ index.php/PSOA RuleML#PSOATransRun



Introduction – PSOATransRun 1.0

- Includes a composition of translator PSOA2Prolog and well-known XSB Prolog engine
 - PSOA2Prolog translates Knowledge Bases (KBs) and queries in PSOA RuleML presentation syntax into a subset of the logic programming language ISO Prolog
 - XSB Prolog engine provides efficient query answering over translated Prolog KB
- PSOA2Prolog performs a multi-step source-to-source normalization followed by a mapping to the pure (Horn) subset of ISO Prolog

- Introduction
- Normalization of PSOA Source
- 3 Mapping the Normalized PSOA Source to Prolog
- 4 Conclusions

Normalization

- Objectification
- Skolemization
- Slotribution and Tupribution
- Flattening
- Rule Splitting

Static vs. Dynamic Objectification – Preview

```
KB:
_hire(_Ernie _Kate)
Query:
?O#_hire(?X ?Y)
```

Static: Generate fresh OID constant from '_1', '_2', ... (transform above **KB** ground atom, use **query** unchanged):

```
_1#_hire(_Ernie _Kate)
```

Dynamic: Virtualize with '_oidcons' function and equality '=' (keep above **KB** unchanged, transform **query** atom):

```
And (\_hire(?X ?Y) ?O = \_oidcons(\_hire ?X ?Y))
```

Objectification – From Static to Static/Dynamic

- Static: generate OIDs for all of the KB's oidless atoms
- Static/Dynamic (novel refinement)
 - Leave unchanged as many of the KB's oidless atoms as possible, instead dynamically constructing virtual OIDs as query variable bindings
 - Partition the set of KB predicates into two disjoint subsets: non-relational (at least one occurrence in a multi-tuple, oidful, or slotted atom, or in a subclass formula) and relational (no such occurrence)
 - Statically generate OIDs only for the KB's oidless psoa atoms with non-relational predicates

Static/Dynamic Objectification

- Perform OID virtualization for queries with OID variables corresponding to the KB's psoa atoms with relational predicates
 - Query atoms using the KB's relational predicates with OID variables are rewritten via equalities that unify an OID variable with a (Skolem-like) OID-constructor ('_oidcons') function application
 - Allow users to pose queries with OIDs regardless of whether the underlying KB clauses have OIDs or not
 - Make maximum use of the underlying Prolog engine for efficient inference on KB clauses with relational predicates

Dynamic Objectification by OID Virtualization

- Leave all KB atoms with relational predicates unchanged
- For each query atom ψ using a relational KB predicate f:
 - If ψ is a relationship, keep it unchanged
 - If ψ has a **non-variable** OID or a slot, rewrite it to explicit falsity, here encoded as Or ()
 - If ψ has an OID **variable** and m tuples, being of the form $?O\#f([t_1,_1 \ldots t_1,_{n_1}] \ldots [t_m,_1 \ldots t_m,_{n_m}])$, equivalent to a tupribution-like conjunction, copying ?O#f, And $(?O\#f(t_1,_1 \ldots t_1,_{n_1}) \ldots ?O\#f(t_m,_1 \ldots t_m,_{n_m})$, rewrite it to a relational conjunction using equality

Skolemization

- Specialized FOL Skolemization is employed to eliminate existentials in rule conclusions or facts, which are not allowed in Prolog
- Replace each existential formula <code>Exists</code> ?X (σ) in a rule conclusion or a fact with $\sigma[?X/_skolemk\ (?v_1\ldots ?v_m)]$, where each occurrence of ?X in σ becomes a Skolem function $_skolemk$ applied to all universally quantified variables $?v_1\ldots ?v_m$ from the clause's quantifier prefix
- New skolem function name $_skolem k$ (k = 1, 2, ...) generated for each existential variable

Slotribution and Tupribution

Rewrite each psoa atom

Flattening

- Extract each embedded interpreted function application as a separate equality
- Each atomic formula φ (in a rule premise or a query) that embeds an external function application ψ , which is not on the top level of an equality, is replaced with And $(?i=\psi \varphi[\psi/?i])$, where ?i is the first variable in ?1, ?2, ... that does not occur in the enclosing rule

Rule Splitting

- Remove conjunctions in rule conclusions, which are not supported in Prolog
- Each rule with an n-ary conjunction in the conclusion

Forall
$$v_1 \ldots v_m$$
 (And $(\varphi_1 \ldots \varphi_n) := \varphi'$) is split into n rules

Forall
$$v_1$$
 ... v_m $(\varphi_1 :- \varphi')$

•

Forall
$$?v_1 \dots ?v_m (\varphi_n :- \varphi')$$

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Simple Terms

Use recursive mapping function denoted by ρ_{psoa}

- Constants
 - If c is a number, ρ_{psoa}(c) is the corresponding Prolog number
 - If c is an arithmetic built-in, $\rho_{psoa}(c)$ is the corresponding Prolog built-in
 - Otherwise, $\rho_{psoa}(c)$ is the single-quoted version of c
- Variables
 For a '?'-prefixed variable ν, ρ_{psoa}(ν) replaces '?' with the upper-case letter 'Q' (Question mark)

Central PSOA Constructs

Mapping from PSOA/PS constructs to Prolog constructs

(To accommodate the relationships preserved by dynamic objectification, the 4th-row mapping of $f(t_1...t_k)$ has to be extended from functions to predicates f)

PSOA/PS Constructs	Prolog Constructs
$o # Top (t_1t_k)$	tupterm($ ho_{ extstyle{psoa}}$ (0), $ ho_{ extstyle{psoa}}$ (t $_1$) $ ho_{ extstyle{psoa}}$ (t $_k$))
o # Top (p -> v)	sloterm($ ho_{ extit{psoa}}$ (0), $ ho_{ extit{psoa}}$ (p), $ ho_{ extit{psoa}}$ (V))
o # c ()	memterm ($ ho_{psoa}$ (o) , $ ho_{psoa}$ (c))
$f(t_1t_k)$	$ ho_{ extstyle{psoa}}(extstyle{f}) \; (ho_{ extstyle{psoa}}(extstyle{t}_1) \; , \; \ldots \; , ho_{ extstyle{psoa}}(extstyle{t}_k) \;)$
And $(f_1 \ldots f_n)$	$(ho_{ extstyle{psoa}}(extstyle{f}_1)$, , $ ho_{ extstyle{psoa}}(extstyle{f}_n)$)
$Or(f_1 \dots f_n)$	$(\rho_{psoa}(f_1); \ldots; \rho_{psoa}(f_n))$
Exists $?v_1 \dots ?v_m (\varphi)$	$ ho_{ extsf{psoa}}(arphi)$
Forall $?v_1 \dots ?v_m (\varphi)$	$ ho_{ extsf{psoa}}(arphi)$
$\varphi:$ - ψ	$ ho_{ extsf{psoa}}(arphi):= ho_{ extsf{psoa}}(\psi)$.
$?v=External(f(t_1t_k))$	is $(\rho_{psoa}(?v), \rho_{psoa}(f)(\rho_{psoa}(t_1), \ldots, \rho_{psoa}(t_k)))$
c1 ## c2	memterm(X, ρ_{psoa} (c2)) :- memterm(X, ρ_{psoa} (c1)).

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Conclusions

- PSOATransRun 1.0 supports reasoning in PSOA RuleML by composing the efficient translator PSOA2Prolog and the efficient run-time engine XSB Prolog
- PSOA2Prolog performs a multi-step source-to-source normalization followed by a mapping to the pure (Horn) subset of ISO Prolog
- Normalization employs a novel static/dynamic approach for the objectification step, which makes optimal use of the underlying Prolog engine for efficient inference on KB clauses with relational predicates
- Future work includes adding support of other PSOA features, e.g. an expanded set of built-ins, and more detailed functionality and performance comparisons with realizations of other rule languages, e.g. Flora-2 and EYE