

SMART CONTRACT AUDIT REPORT

for

RULER PROTOCOL

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1 Introduction

Given the opportunity to review the design document and related smart contract source code of the Ruler protocol, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

1.1 About Ruler Protocol

The Ruler protocol is a lending platform where users can borrow their preferred cryptocurrency with any other cryptocurrency. It aims to fill the gap by enabling the following goals: 1) supply and demand decide interest rate; 2) no liquidation as long as borrowers payback on time; 3) fungible loans that can be traded anytime. At the core of Ruler Protocol is the notion of ruler pairs and each pair consists of four elements: collateral token, paired token, expiry, and mint ratio. In essence, the protocol enables a market driven lending platform that provides non-liquidatable and fungible loans.

The basic information of the Ruler protocol is as follows:

Table 1.1: Basic Information of The Ruler Protocol

ltem	Description
lssuer	Ruler Protocol
Website	https://rulerprotocol.com/
Туре	Ethereum Smart Contract
Platform	Solidity
Audit Method	Whitebox
Latest Audit Report	February 19, 2021

In the following, we show the Git repository of reviewed files and the commit hash value used in

this audit.

• https://github.com/Ruler-Protocol/ruler-core.git (b90e7a6)

And this is the commit ID after all fixes for the issues found in the audit have been checked in:

https://github.com/Ruler-Protocol/ruler-core.git (bd0ca68)

1.2 About PeckShield

PeckShield Inc. [10] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (https://t.me/peckshield), Twitter (http://twitter.com/peckshield), or Email (contact@peckshield.com).

High Critical High Medium

High Medium

Low

High Low

High Medium

Low

High Medium

Low

Likelihood

Table 1.2: Vulnerability Severity Classification

1.3 Methodology

To standardize the evaluation, we define the following terminology based on OWASP Risk Rating Methodology [9]:

- <u>Likelihood</u> represents how likely a particular vulnerability is to be uncovered and exploited in the wild;
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Table 1.3: The Full List of Check Items

Category	Check Item
	Constructor Mismatch
	Ownership Takeover
	Redundant Fallback Function
	Overflows & Underflows
	Reentrancy
	Money-Giving Bug
	Blackhole
	Unauthorized Self-Destruct
Basic Coding Bugs	Revert DoS
Dasic Coung Dugs	Unchecked External Call
	Gasless Send
	Send Instead Of Transfer
-	Costly Loop
	(Unsafe) Use Of Untrusted Libraries
	(Unsafe) Use Of Predictable Variables
	Transaction Ordering Dependence
	Deprecated Uses
Semantic Consistency Checks	Semantic Consistency Checks
	Business Logics Review
	Functionality Checks
	Authentication Management
	Access Control & Authorization
	Oracle Security
Advanced DeFi Scrutiny	Digital Asset Escrow
Advanced Berr Scruting	Kill-Switch Mechanism
	Operation Trails & Event Generation
	ERC20 Idiosyncrasies Handling
Additional Recommendations	Frontend-Contract Integration
	Deployment Consistency
	Holistic Risk Management
	Avoiding Use of Variadic Byte Array
	Using Fixed Compiler Version
	Making Visibility Level Explicit
	Making Type Inference Explicit
	Adhering To Function Declaration Strictly
	Following Other Best Practices

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

To evaluate the risk, we go through a list of check items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- Basic Coding Bugs: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- <u>Semantic Consistency Checks</u>: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [8], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings.

1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary
Configuration	Weaknesses in this category are typically introduced during
	the configuration of the software.
Data Processing Issues	Weaknesses in this category are typically found in functional-
	ity that processes data.
Numeric Errors	Weaknesses in this category are related to improper calcula-
	tion or conversion of numbers.
Security Features	Weaknesses in this category are concerned with topics like
	authentication, access control, confidentiality, cryptography,
	and privilege management. (Software security is not security
	software.)
Time and State	Weaknesses in this category are related to the improper man-
	agement of time and state in an environment that supports
	simultaneous or near-simultaneous computation by multiple
	systems, processes, or threads.
Error Conditions,	Weaknesses in this category include weaknesses that occur if
Return Values,	a function does not generate the correct return/status code,
Status Codes	or if the application does not handle all possible return/status
	codes that could be generated by a function.
Resource Management	Weaknesses in this category are related to improper manage-
	ment of system resources.
Behavioral Issues	Weaknesses in this category are related to unexpected behav-
	iors from code that an application uses.
Business Logics	Weaknesses in this category identify some of the underlying
	problems that commonly allow attackers to manipulate the
	business logic of an application. Errors in business logic can
	be devastating to an entire application.
Initialization and Cleanup	Weaknesses in this category occur in behaviors that are used
	for initialization and breakdown.
Arguments and Parameters	Weaknesses in this category are related to improper use of
	arguments or parameters within function calls.
Expression Issues	Weaknesses in this category are related to incorrectly written
	expressions within code.
Coding Practices	Weaknesses in this category are related to coding practices
	that are deemed unsafe and increase the chances that an ex-
	ploitable vulnerability will be present in the application. They
	may not directly introduce a vulnerability, but indicate the
	product has not been carefully developed or maintained.

2 | Findings

2.1 Summary

Here is a summary of our findings after analyzing the Ruler Protocol implementation. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings
Critical	0
High	0
Medium	1
Low	3
Informational	1
Total	5

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities that need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in Section 3.

2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 1 medium-severity vulnerability, 3 low-severity vulnerabilities, and 1 informational recommendation.

ID Title Severity Category Status PVE-001 Medium Safe-Version Replacement With safeAp-Coding Practices Fixed prove(), safeTransfer() And safeTransfer-From() **PVE-002** Low Accommodation of approve() Idiosyncrasies **Business Logic** Fixed PVE-003 Fixed Informational Improved Precision By Multiplication And Di-Numeric Errors vision Reordering PVE-004 Improved Sanity Checks Of System/Function Coding Practices Confirmed Low **Parameters PVE-005** Front-Running For Nonce Invalidation Time and State Low Confirmed

Table 2.1: Key Ruler Protocol Audit Findings

Besides recommending specific countermeasures to mitigate these issues, based on the fact that compiler upgrades might bring unexpected compatibility or inter-version consistencies, it is always suggested to use fixed compiler versions whenever possible. As an example, we highly encourage to explicitly indicate the Solidity compiler version, e.g., pragma solidity 0.8.0 instead of specifying a range, e.g., pragma solidity ^0.8.0.

In addition, we also emphasize that it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms need to kick in at the very moment when the contracts are being deployed in mainnet. Please refer to Section 3 for details.

3 Detailed Results

3.1 Safe-Version Replacement With safeApprove(), safeTransfer() And safeTransferFrom()

• ID: PVE-001

• Severity: Medium

Likelihood: Medium

• Impact: Medium

• Target: Multiple Contracts

• Category: Coding Practices [4]

• CWE subcategory: CWE-1126 [1]

Description

Though there is a standardized ERC-20 specification, many token contracts may not strictly follow the specification or have additional functionalities beyond the specification. In this section, we examine the transfer() routine and possible idiosyncrasies from current widely-used token contracts. In particular, we use the popular stablecoin, i.e., USDT, as our example. We show the related code snippet below.

```
121
122
         * @dev transfer token for a specified address
123
         * @param _to The address to transfer to.
124
         * @param _value The amount to be transferred.
125
         function transfer(address _to, uint _value) public onlyPayloadSize(2 * 32) {
126
127
             uint fee = ( value.mul(basisPointsRate)).div(10000);
128
             if (fee > maximumFee) {
129
                 fee = maximumFee;
130
131
             uint sendAmount = value.sub(fee);
132
             balances [msg.sender] = balances [msg.sender].sub( value);
133
             balances [ to] = balances [ to].add(sendAmount);
134
             if (fee > 0) {
135
                 balances[owner] = balances[owner].add(fee);
136
                 Transfer(msg.sender, owner, fee);
137
```

Listing 3.1: USDT Token Contract

It is important to note the transfer() function does not have a return value. However, the IERC20 interface has defined the following transfer() interface with a bool return value: function transfer(address recipient, uint256 amount)external returns (bool). As a result, the call to transfer() may expect a return value. With the lack of return value of USDT's transfer(), the call will be unfortunately reverted.

Because of that, a normal call to transfer() is suggested to use the safe version, i.e., safeTransfer (), In essence, it is a wrapper around ERC20 operations that may either throw on failure or return false without reverts. Moreover, the safe version also supports tokens that return no value (and instead revert or throw on failure). Note that non-reverting calls are assumed to be successful. To use this library you can add a using SafeERC20 for IERC20. Similarly, there is a safe version of approve()/transferFrom() as well, i.e., safeApprove()/safeTransferFrom().

In the following, we show the collectDust() routine in the BonusRewards contract. If the USDT token is given as the routine's argument, i.e., _token, the unsafe version of IERC20(_token).transfer (owner(), balance) (line 236) may revert as there is no return value in the USDT token contract's transfer() implementation (but the IERC20 interface expects a return value)!

```
219
      /// @notice collect bonus token dust to treasury
220
      function collectDust(address _token, address _lpToken, uint256 _poolBonusId) external
          override onlyOwner {
221
        require(pools[ token].lastUpdatedAt == 0, "BonusRewards: lpToken, not allowed");
223
        if (\_token == address(0)) \{ // token address(0) = ETH
          payable(owner()).transfer(address(this).balance);
224
225
        } else {
          uint256 balance = IERC20( token).balanceOf(address(this));
226
227
           if (bonusTokenAddrMap[ token] == 1) {
228
             // bonus token
229
             Bonus memory bonus = pools [ lpToken]. bonuses [ poolBonusId];
230
             require(bonus.bonusTokenAddr == token, "BonusRewards: wrong pool");
231
             require(bonus.endTime + WEEK < block.timestamp, "BonusRewards: not ready");</pre>
232
             balance = bonus.remBonus;
233
             pools[ lpToken].bonuses[ poolBonusId].remBonus = 0;
234
          }
236
          IERC20( token).transfer(owner(), balance);
237
```

Listing 3.2: BonusRewards::collectDust()

Note that the same issue exists in the _depositAndAddLiquidity() routine from the RulerZap contract, which reverts related liquidation additions. Also, the _approve() helper from the same contract shares the same issue, which may revert a number of calling routines.

Recommendation Accommodate the above-mentioned idiosyncrasy about ERC20-related approve()/transfer()/transferFrom().

Status The issue has been fixed by this commit: 4e753ce.

3.2 Accommodation of approve() Idiosyncrasies

ID: PVE-002

Severity: Low

Likelihood: medium

Impact: Low

• Target: RulerZap

• Category: Business Logic [5]

CWE subcategory: N/A

Description

In Section 3.1, we have examined certain non-compliant ERC20 tokens that may exhibit specific idiosyncrasies in their transfer() and transferFrom() implementations. In this section, we examine the approve() routine and possible another idiosyncrasy from current widely-used token contracts.

In particular, we use the popular stablecoin, i.e., USDT, as our example. We show the related code snippet below. On its entry of approve(), there is a requirement, i.e., require(!((_value != 0) && (allowed[msg.sender][_spender] != 0))). This specific requirement essentially indicates the need of reducing the allowance to 0 first (by calling approve(_spender, 0)) if it is not, and then calling a second one to set the proper allowance. This requirement is in place to mitigate the known approve()/transferFrom() race condition (https://github.com/ethereum/EIPs/issues/20#issuecomment-263524729).

```
194
195
        * @dev Approve the passed address to spend the specified amount of tokens on behalf
            of msg.sender.
196
        * @param _spender The address which will spend the funds.
197
        * @param _value The amount of tokens to be spent.
198
199
        function approve(address spender, uint value) public onlyPayloadSize(2 * 32) {
201
            // To change the approve amount you first have to reduce the addresses '
202
            // allowance to zero by calling 'approve(_spender, 0)' if it is not
203
            // already 0 to mitigate the race condition described here:
204
            // https://github.com/ethereum/EIPs/issues/20#issuecomment-263524729
205
            require(!(( value != 0) && (allowed[msg.sender][ spender] != 0)));
207
            allowed [msg.sender] [ _spender] = _value;
208
            Approval (msg. sender, spender, value);
209
```

Listing 3.3: USDT Token Contract

Because of that, a normal call to approve() with a currently non-zero allowance may fail. In the following, we use as an example the RulerZap contract that is designed to facilitate the interaction with the Ruler Core contract. To accommodate the specific idiosyncrasy, there is a need to approve() twice: the first one reduces the allowance to 0; and the second one sets the new allowance.

```
310
        function _deposit(
311
             address _col,
312
             address paired,
313
             uint48 expiry,
             uint256 mintRatio,
314
315
             uint256 colAmt
316
        ) internal returns (address rcTokenAddr, uint256 rcTokenReceived, uint256
             rcTokenBalBefore) {
317
             ( , , , IRERC20 rcToken , IRERC20 rrToken , , , ) = core.pairs ( col , paired ,
                 expiry, mintRatio);
318
             // receive collateral from sender
319
             IERC20 collateral = IERC20( col);
320
             uint256 colBalBefore = collateral.balanceOf(address(this));
321
             collateral.safeTransferFrom(msg.sender, address(this), colAmt);
322
             uint256 received = collateral.balanceOf(address(this)) - colBalBefore;
323
             require(received > 0, "RulerZap: col transfer failed");
325
             // deposit collateral to Ruler
326
             rcTokenBalBefore = rcToken.balanceOf(address(this));
327
             uint256 rrTokenBalBefore = rrToken.balanceOf(address(this));
328
             _approve(collateral, address(core), received);
329
             core.deposit(\_col,\_paired,\_expiry,\_mintRatio, received);\\
331
             // send rrToken back to sender, and record received rcTokens
332
             rrToken.transfer(msg.sender, rrToken.balanceOf(address(this)) - rrTokenBalBefore
333
             rcTokenReceived = rcToken.balanceOf(address(this)) - rcTokenBalBefore;
334
             rcTokenAddr = address(rcToken);
335
        }
337
        function approve(IERC20 token, address spender, uint256 amount) internal {
338
             if (\_token.allowance(address(this), \_spender) < \_amount) {
339
                 _token.approve(_spender, type(uint256).max);
340
             }
341
```

Listing 3.4: RulerZap:: deposit()

Recommendation Accommodate the above-mentioned idiosyncrasy of approve().

Status The issue has been fixed by this commit: 4e753ce.

3.3 Improved Precision By Multiplication And Division Reordering

• ID: PVE-003

• Severity: Informational

Likelihood: N/A

• Impact: N/A

• Target: BonusRewards

• Category: Numeric Errors [7]

• CWE subcategory: CWE-190 [2]

Description

In the Ruler protocol, there is a BonusRewards contract that allows for rewarding multiple bonus tokens for participating users. The reward logic enforces pro-rata claims of disseminated rewards.

For illustration, we show below the updateBonus() routine that is used to update an active rewarding pool. It implements a rather straightforward logic by firstly performing necessary sanity checks on the input arguments, and then updating the related bonus entry with the specified startTime and weeklyRewards (lines 154 - 164).

```
136
      /// @notice called by authorizers only, update weeklyRewards (if not ended), or update
           startTime (only if rewards not started, 0 is ignored)
137
      function updateBonus(
138
        address _IpToken,
139
        address bonusTokenAddr,
140
        uint256 _weeklyRewards,
141
        uint48 _startTime
142
      ) external override nonReentrant notPaused {
        require( isAuthorized(allowedTokenAuthorizers[_lpToken][_bonusTokenAddr]), "
143
            BonusRewards: not authorized caller");
144
        the past");
146
        // make sure the pool is in the right state (exist with no active bonus at the
           moment) to add new bonus tokens
147
        Pool memory pool = pools[ lpToken];
148
        require(pool.lastUpdatedAt > 0, "BonusRewards: pool does not exist");
149
        Bonus[] memory bonuses = pool.bonuses;
150
        for (uint256 i = 0; i < bonuses.length; i++) {
151
          if (bonuses[i].bonusTokenAddr == bonusTokenAddr && bonuses[i].endTime > block.
152
            Bonus storage bonus = pools[ lpToken].bonuses[i];
153
            updatePool( IpToken); // update pool with old weeklyReward to this block
154
            if (bonus.startTime >= block.timestamp) {
155
              // only honor new start time, if program has not started
156
              if ( startTime >= block.timestamp) {
157
                bonus.startTime = startTime;
158
```

```
159
               bonus.endTime = uint48 (bonus.remBonus * WEEK / weeklyRewards + bonus.
                   startTime);
160
             } else {
               uint256 remBonusToDistribute = (bonus.endTime - block.timestamp) * bonus.
161
                   weeklyRewards / WEEK;
               bonus.endTime = uint48(remBonusToDistribute * WEEK / weeklyRewards + block.
162
                   timestamp);
163
             bonus.weeklyRewards = weeklyRewards;
164
165
          }
166
        }
167
```

Listing 3.5: BonusRewards::updateBonus()

The update of startTime and weeklyRewards naturally leads to the re-calculation of the endTime. It comes to our attention that the current endTime is computed as follows: uint48(remBonusToDistribute * WEEK / _weeklyRewards + block.timestamp), where remBonusToDistribute = (bonus.endTime - block.timestamp)* bonus.weeklyRewards / WEEK.

It is important to emphasize that the lack of float support in Solidity may introduce subtle, but troublesome issue: precision loss. One possible precision loss stems from the computation when both multiplication (mul) and division (div) are involved. Specifically, the computation at lines 161 — 162 can be performed as follows: uint48((bonus.endTime - block.timestamp)* bonus.weeklyRewards / _weeklyRewards + block.timestamp).

A better approach is the one that can always avoid or reduce any precision loss. In other words, the computation of the form A / B * C can be converted into A * C / B under the condition that A * C does not introduce any overflow.

Recommendation Avoid unnecessary precision loss due to the lack of floating support in Solidity. An example revision to the above endTime is shown as: remBonusToDistribute = (bonus.endTime - block.timestamp)* bonus.weeklyRewards / WEEK.

Status The issue has been fixed by this commit: 40f65e4.

3.4 Improved Sanity Checks For System/Function Parameters

• ID: PVE-004

• Severity: Low

• Likelihood: Low

• Impact: Low

Target: RulerCore

• Category: Coding Practices [4]

• CWE subcategory: CWE-1126 [1]

Description

DeFi protocols typically have a number of system-wide parameters that can be dynamically configured on demand. The Ruler Protocol protocol is no exception. Specifically, if we examine the RulerCore contract, it has defined a number of protocol-wide risk parameters: flashLoanRate and minColRatioMap. In the following, we show the corresponding routine that allows for their changes.

```
307
      function setFlashLoanRate(uint256 newRate) external override onlyOwner {
        emit FlashLoanRateUpdated(flashLoanRate, newRate);
308
309
        flashLoanRate = _newRate;
310
      }
311
312
      function setMinColRatio(address _col, uint256 _minColRatio) external override
        require(minColRatioMap[_col] > 0, "Ruler: collateral not listed");
313
314
        require( minColRatio >= 0.5 ether, "Ruler: min colRatio < 50%");</pre>
315
        emit MinColRatioUpdated( col, minColRatioMap[ col], minColRatio);
316
        minColRatioMap[ col] = minColRatio;
317
      }
```

Listing 3.6: RulerCore :: setFlashLoanRate() and RulerCore :: setMinColRatio()

These parameters define various aspects of the protocol operation and maintenance and need to exercise extra care when configuring or updating them. Our analysis shows the update logic on these parameters can be improved by applying more rigorous sanity checks. Based on the current implementation, certain corner cases may lead to an undesirable consequence. For example, an unlikely mis-configuration of flashLoanRate may charge unreasonably high fee in the flashLoan() operation, hence incurring cost to participating users or hurting the adoption of the flashloans.

Recommendation Validate any changes regarding these system-wide parameters to ensure they fall in an appropriate range. If necessary, also consider emitting relevant events for their changes.

Status This issue has been confirmed and by design it is allowed to set any flashloan rate of choice.

3.5 Possible Front-Running For Nonce Invalidation

• ID: PVE-005

Severity: Low

Likelihood: Low

Impact: Low

• Target: Multiple Contracts

• Category: Time and State [6]

• CWE subcategory: CWE-663 [3]

Description

In the Ruler protocol, both RulerCore and RulerZap contracts support permit-style function variants that allow users to authorize actions using an off-chain signature. The intention is to support metatransactions such that an user can simply sign an intended transaction offline and then send the signed transaction to a relayer. The replayer will take care of submitting the transaction for mining by paying required transaction fee. The signature correctness will be verified on chain via a helper routine, i.e., permit().

In the following, we elaborate the related execution logic. Particularly, the depositWithPermit() function verifies the signature using the internal _permit() helper (line 126) before executing the intended deposit().

```
/// @notice deposit collateral to a Ruler Pair, sender receives rTokens
 92
 93
       function deposit (
 94
         address col,
 95
         address _ paired,
 96
         uint48 _expiry ,
 97
         uint256 _ mintRatio ,
 98
         uint256 colAmt
 99
      ) public override onlyNotPaused nonReentrant {
100
         Pair memory pair = pairs [ col][ paired][ expiry][ mintRatio];
101
         validateDepositInputs( col, pair);
103
         // receive collateral
104
         IERC20 collateral = IERC20( col);
105
         uint256 colBalBefore = collateral.balanceOf(address(this));
106
         collateral.safeTransferFrom(msg.sender, address(this), colAmt);
107
         uint256 received = collateral.balanceOf(address(this)) - colBalBefore;
108
         require(received > 0, "Ruler: transfer failed");
109
         pairs [_col][_paired][_expiry][_mintRatio]. colTotal = pair.colTotal + received;
111
         // mint rTokens for reveiced collateral
112
         uint256 mintAmount = getRTokenAmtFromColAmt(received, col, paired, pair.mintRatio
113
         pair.rcToken.mint(msg.sender, mintAmount);
114
         pair.rrToken.mint(msg.sender, mintAmount);
115
         emit Deposit(msg.sender, _col, _paired, _mintRatio, received);
116
```

```
118
       function depositWithPermit(
119
         address col,
120
         address _paired,
         uint48 _expiry,
121
122
         uint256 mintRatio,
123
         uint256 colAmt,
124
         Permit calldata colPermit
125
       ) external override {
126
         _permit(_col, _colPermit);
127
         deposit(_col, _paired, _expiry, _mintRatio, _colAmt);
128
```

Listing 3.7: RulerCore :: depositWithPermit()

Within the _permit() helper, we observe that it basically invokes the built-in public permit() function in the token contract. This public function takes the normal procedure, i.e., ecrecover(), to retrieve the signer information. If validated, it advances the nonce by 1, i.e., nonce_[owner]++. Since this function is defined as public, any one could call this function to verify the signature but with the side-effect of advancing the nonce (if successfully verified)!

```
39
        function permit(address owner, address spender, uint256 amount, uint256 deadline,
            uint8 v, bytes32 r, bytes32 s) public virtual override {
40
            // solhint-disable-next-line not-rely-on-time
41
            require(block.timestamp <= deadline, "ERC20Permit: expired deadline");</pre>
43
            bytes32 structHash = keccak256(
44
                abi.encode(
45
                     PERMIT TYPEHASH,
46
                    owner,
47
                    spender,
48
                    amount,
49
                     nonces [owner],
50
                    deadline
                )
51
52
            );
            bytes32 hash = hashTypedDataV4(structHash);
54
56
            address signer = ECDSA.recover(hash, v, r, s);
57
            require(signer == owner, "ERC20Permit: invalid signature");
59
            nonces[owner]++;
60
            _approve(owner, spender, amount);
61
```

Listing 3.8: ERC20Permit::permit()

Here comes the problem: when an user invokes depositWithPermit() to perform specified actions by signing the transaction offline, but before the transaction is mined, it is possible for a malicious actor to observe it (by closely monitoring the transaction pool) and then possibly front-runs it by

crafting a new transaction and offering a higher gas fee for block inclusion. The new transaction may perform a fresh permit() call. If the front-running is successful, the crafted transaction essentially advances the nonce by 1, effectively invalidating the user transaction that is being front-run.

Recommendation This is a common issue inherent in current blockchain infrastructure. However, its impact is rather limited in causing frictions without actual damages.

Status This issue has been confirmed.



4 Conclusion

In this audit, we have analyzed the design and implementation of the Ruler protocol that is a market driven lending platform that provides non-liquidatable and fungible loans. During the audit, we notice that the current code base is well organized and those identified issues are promptly confirmed and fixed.

Meanwhile, we need to emphasize that smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.



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