Cryptography Engineering

- Lecture 3 (Nov 06, 2024)
- Today's notes:
 - Signed Diffie-Hellman Key Exchange (SigDH) Protocol
 - TLS handshake and HTTPS protocol
- Today's coding tasks (and homework):
 - Play with HKDF
 - Implement a toy example of TLS handshake



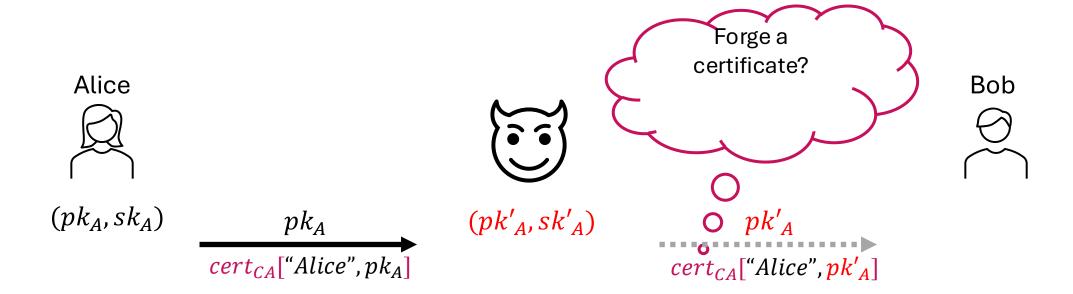
• Transporting (malicious) public keys



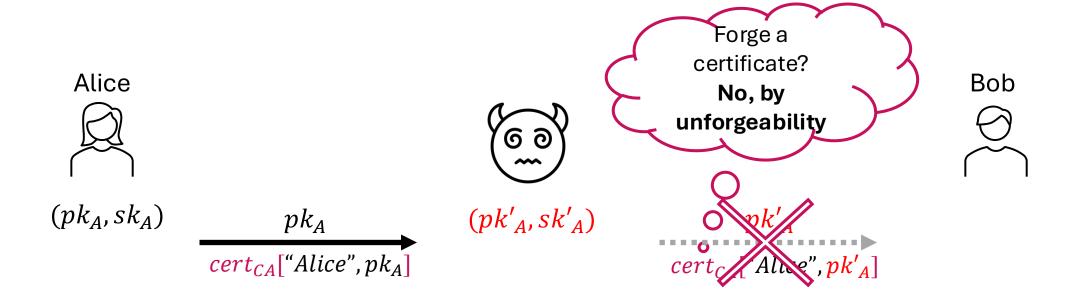
- Transporting (malicious) public keys (with signature/certificate)
 - (Note that a certificate binds a public key with the identity of owner)



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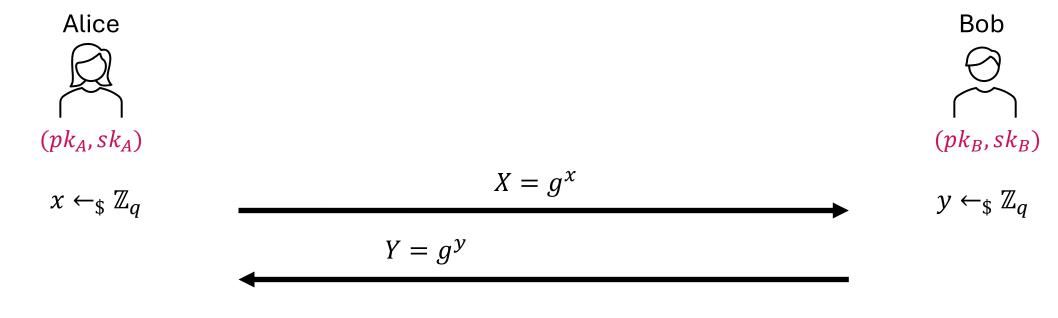


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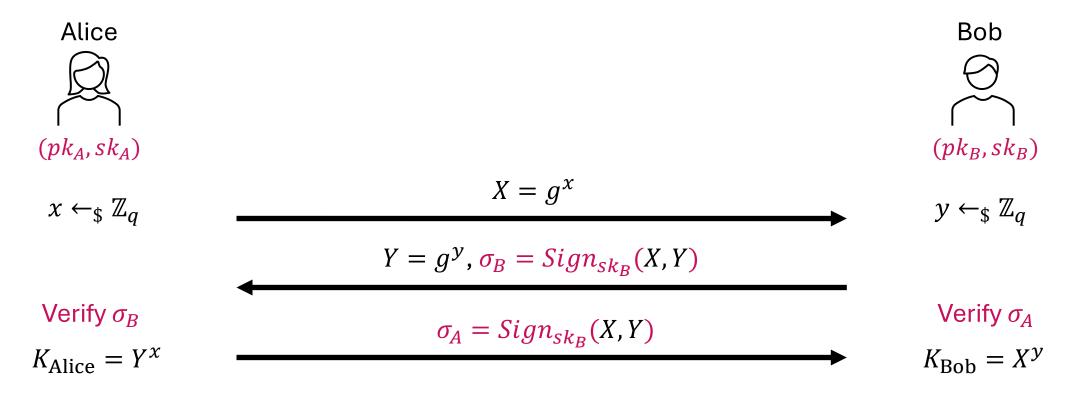
Signed Diffie-Hellman Protocol (Simplified)

• Use signature to avoid MitM attacks

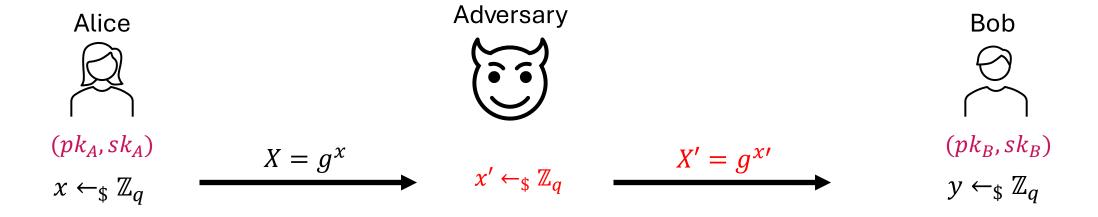


Signed Diffie-Hellman Protocol (Simplified)

Use signature to avoid MitM attacks

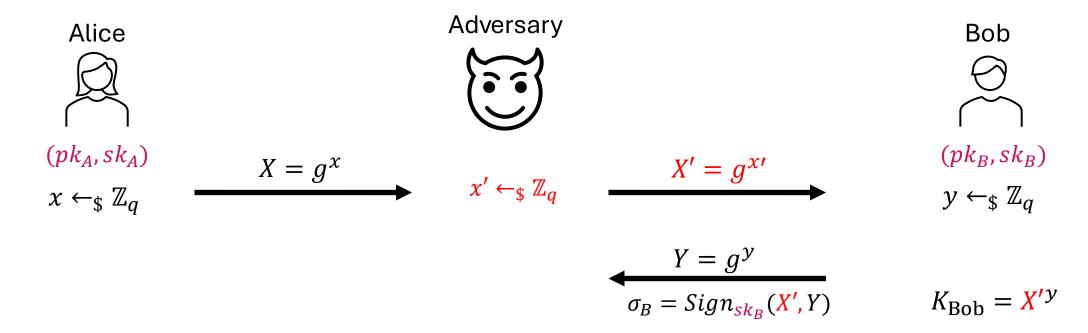


Can we launch a MitM attack on SigDH?

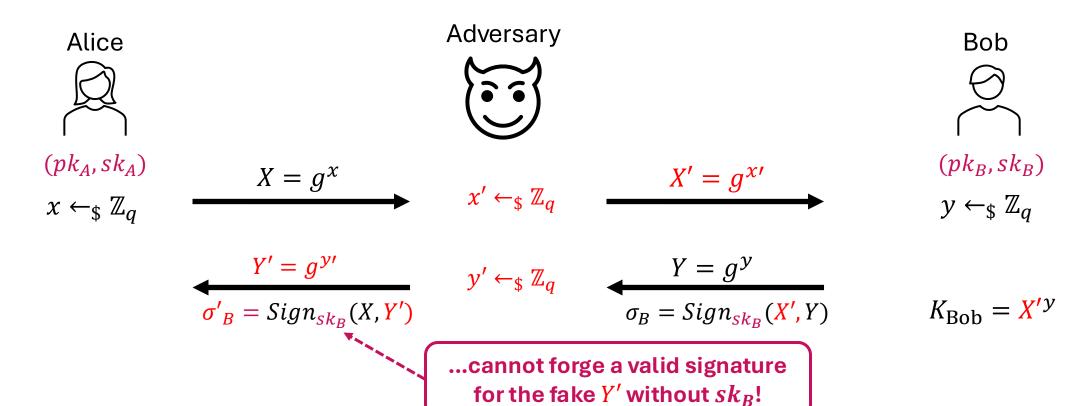


 $K_{\text{Bob}} = X'^{y}$

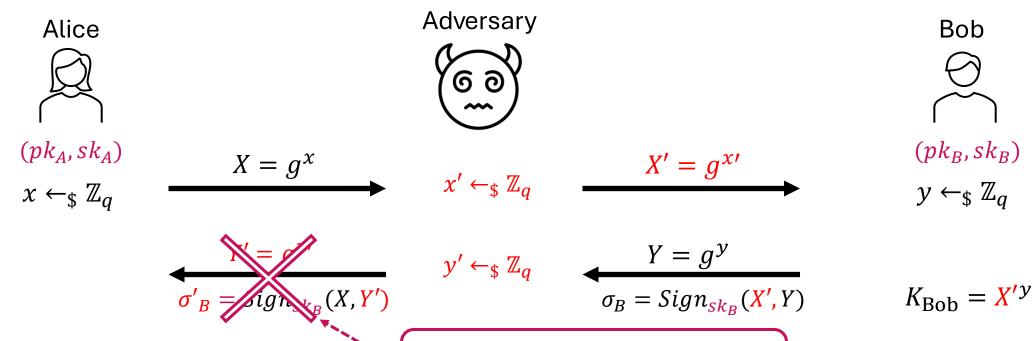
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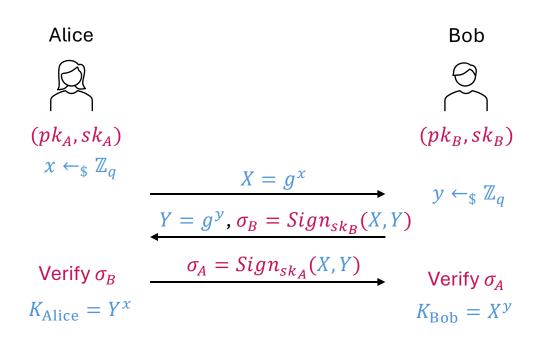


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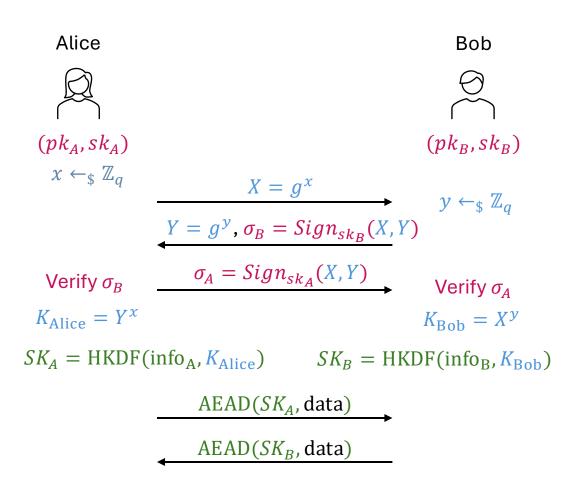


No, by unforgeability...

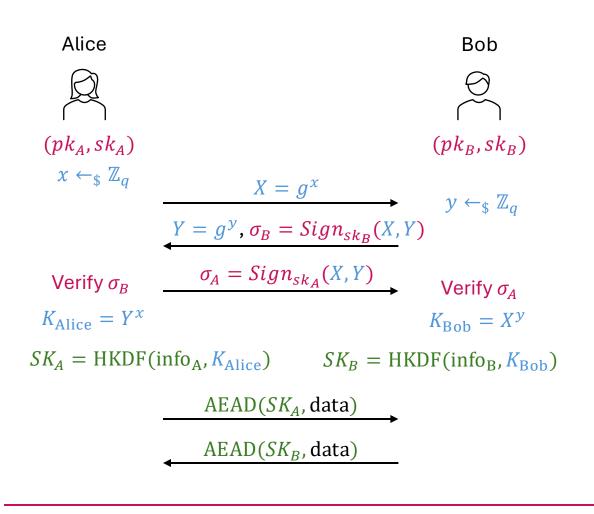
...cannot forge a valid signature for the fake Y' without $sk_B!$



- SigDH
 - Add signature to avoid MitM
 - Authenticated Key Exchange



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- In practice: ECDH + ECDSA + HKDF/HMAC + AEAD ...



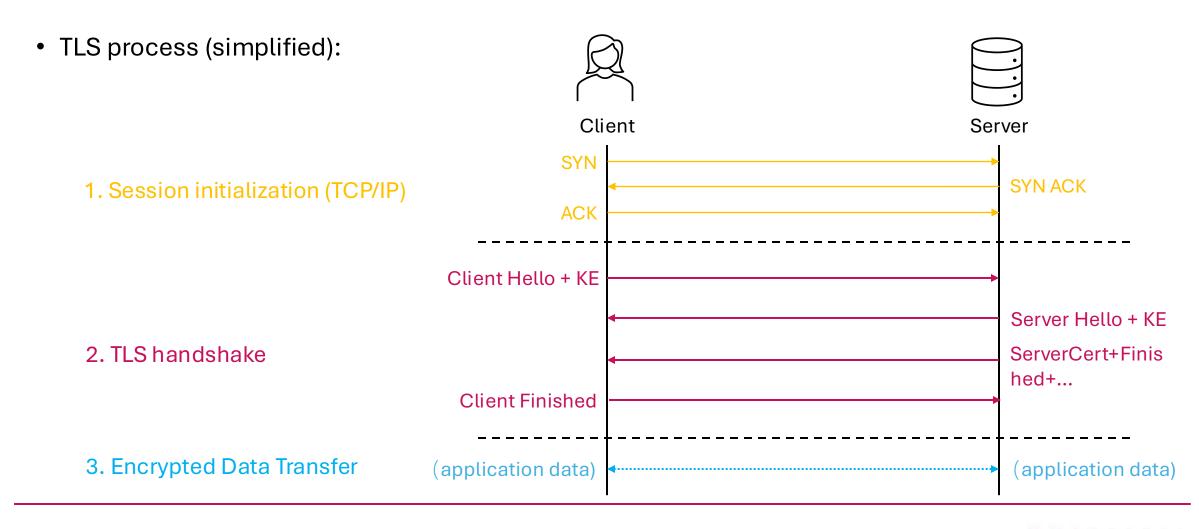
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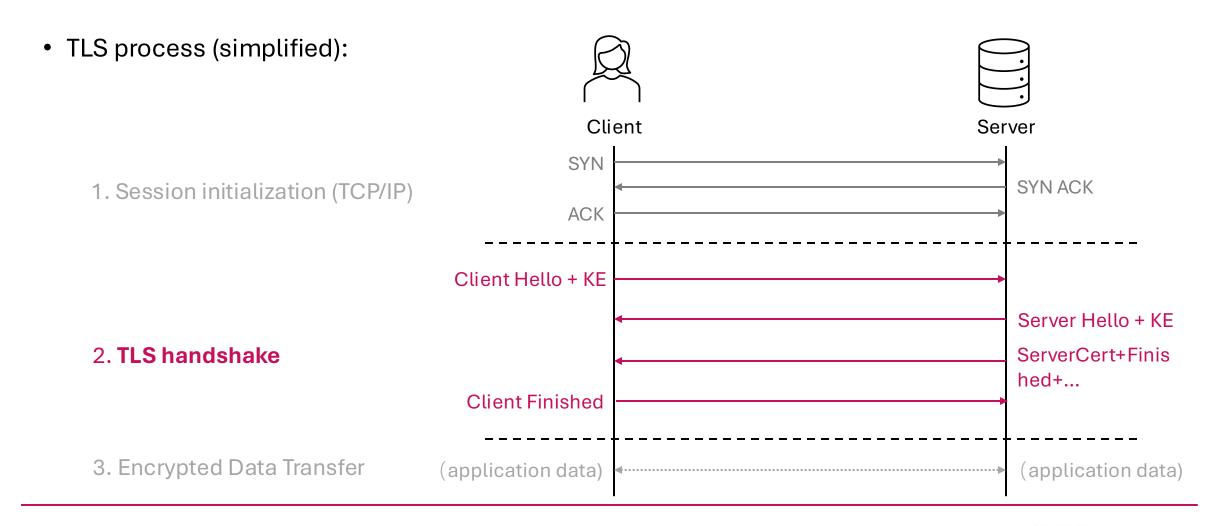
- Important Application: TLS handshake protocol...
 - Note: Cryptographic algorithms "in textbooks" often contrasts with their real-world implementation

- Transport Layer Security (TLS) Protocol
 - Designed to provide communications security over an open network
 - Used in HTTPS On the Latest of the Latest

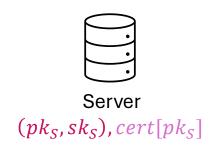
- Transport Layer Security (TLS) Protocol
 - Designed to provide communications security over an open network
 - Used in HTTPS On the Latest of the Latest
 - TLS process (simplified):
 - 1. Session initialization (TCP/IP)
 - 2. TLS handshake
 - 3. Encrypted Data Transfer
 - 4. Session end
 - In this lecture, we mainly consider the client-server setting
 - Server Authentication Only: A client normally does not have static public-private key pair and certificates



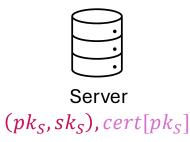












$$nonce_c, X = g^x$$

• TLS 1.3 handshake protocol (Simplified description, we ignore the TLS key schedule)



Server

 $(pk_S, sk_S), cert[pk_S]$

ClientHello + ClientKE Phase

$$nonce_c, X = g^x$$

 $nonce_S, Y = g^y$

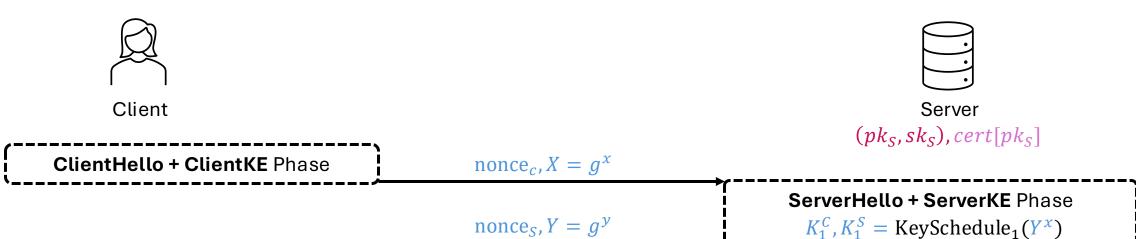
ServerHello + ServerKE Phase

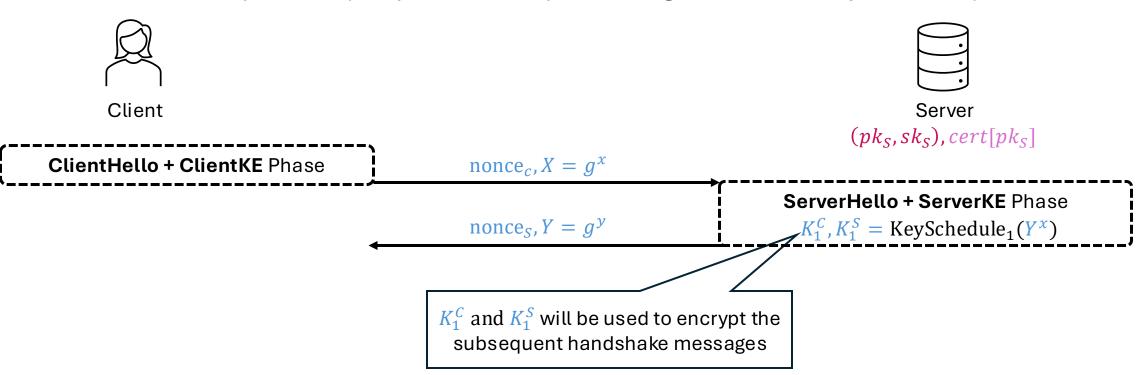
nonce_S
$$\leftarrow_{\$} \{0,1\}^{256}$$

 $y \leftarrow_{\$} \mathbb{Z}_q, Y = g^y$

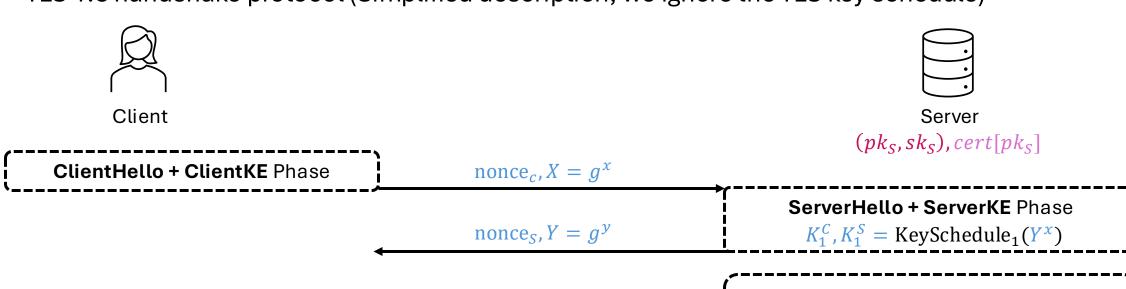
--- Derive keys (using some very --- complicated key schedule algorithm)

$$K_1^C$$
, K_1^S = KeySchedule₁ (Y^x)

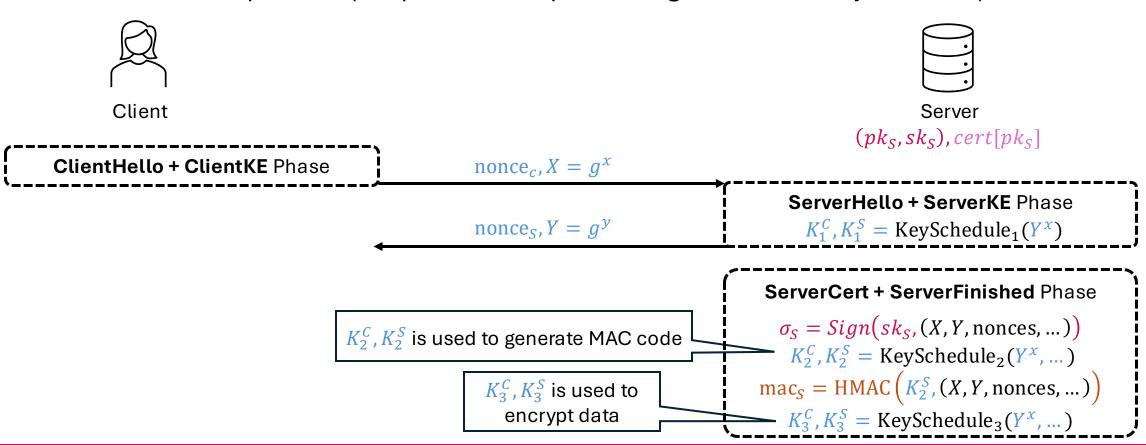


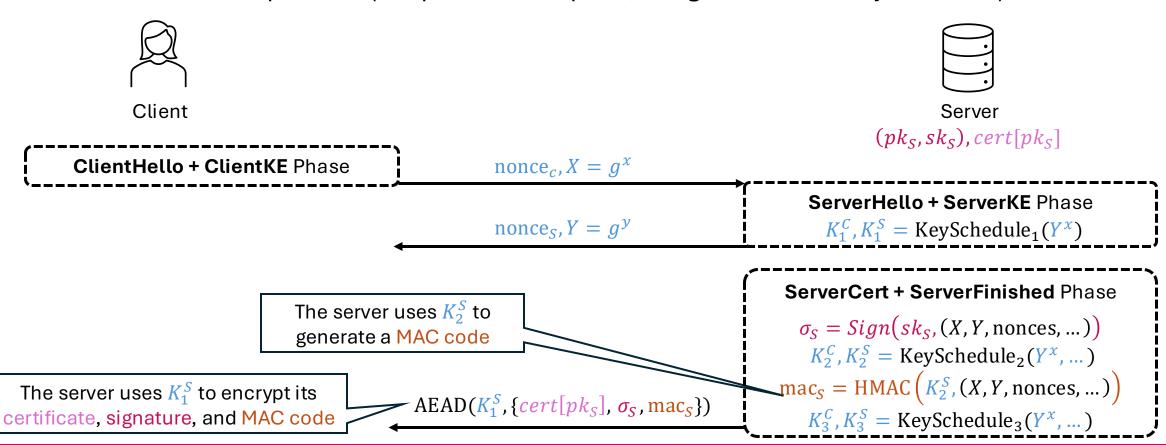


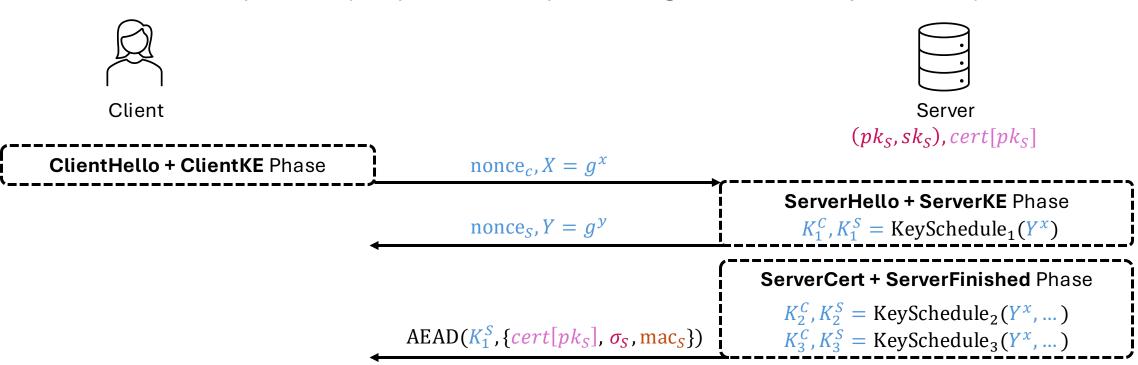
• TLS 1.3 handshake protocol (Simplified description, we ignore the TLS key schedule)



ServerCert + ServerFinished Phase $\sigma_{S} = Sign(sk_{S}, (X, Y, nonces, ...))$ $K_{2}^{C}, K_{2}^{S} = KeySchedule_{2}(Y^{x}, ...)$ $mac_{S} = HMAC(K_{2}^{S}, (X, Y, nonces, ...))$ $K_{3}^{C}, K_{3}^{S} = KeySchedule_{3}(Y^{x}, ...)$







• TLS 1.3 handshake protocol (Simplified description, we ignore the TLS key schedule)





 $(pk_S, sk_S), cert[pk_S]$

ClientHello + ClientKE Phase

$$nonce_c, X = g^x$$

 $nonce_S, Y = g^y$

ServerHello + ServerKE Phase

$$K_1^C, K_1^S = \text{KeySchedule}_1(Y^x)$$

ClientFinished Phase

$$K_1^C, K_1^S, K_2^C, K_2^S, K_3^C, K_3^S = \text{KeySchedule}(X^y, ...)$$

AEAD Decryption

Verify $cert[pk_S]$, σ_S , and mac_S $mac_C = HMAC(K_2^C, (X, Y, nonces, ...))$ $AEAD(K_1^S, \{cert[pk_S], \sigma_S, mac_S\})$

 $AEAD(K_1^C, \{mac_C\})$

ServerCert + ServerFinished Phase

 $K_2^C, K_2^S = \text{KeySchedule}_2(Y^x, ...)$

 K_3^C , K_3^S = KeySchedule₃(Y^x , ...)

• TLS 1.3 handshake protocol (Simplified description, we ignore the TLS key schedule)





ClientHello + ClientKE Phase

$$nonce_c, X = g^x$$

 $nonce_S, Y = g^y$

ServerHello + ServerKE Phase K_1^C , K_1^S = KeySchedule₁ (Y^x)

ClientFinished Phase

$$K_1^C, K_1^S, K_2^C, K_2^S, K_3^C, K_3^S = \text{KeySchedule}(X^y, ...)$$

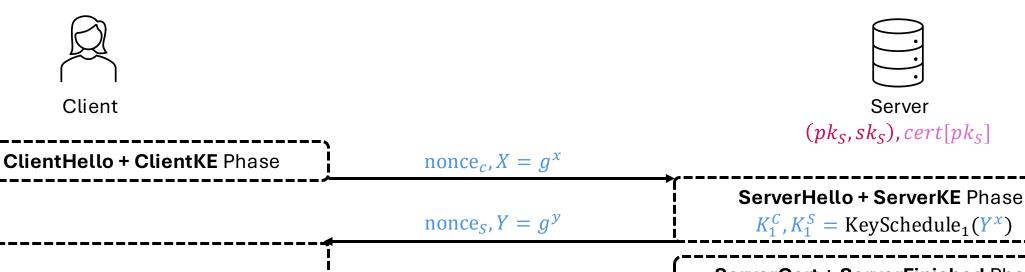
 $AEAD(K_1^S, \{cert[pk_S], \sigma_S, mac_S\})$

 $AEAD(K_1^C, \{mac_C\})$

ServerCert + ServerFinished Phase

$$K_2^C$$
, K_2^S = KeySchedule₂(Y^x , ...)
 K_3^C , K_3^S = KeySchedule₃(Y^x , ...)

• TLS 1.3 handshake protocol (Simplified description, we ignore the TLS key schedule)



ClientFinished Phase

$$K_1^C, K_1^S, K_2^C, K_2^S, K_3^C, K_3^S = \text{KeySchedule}(Y^x, ...)$$

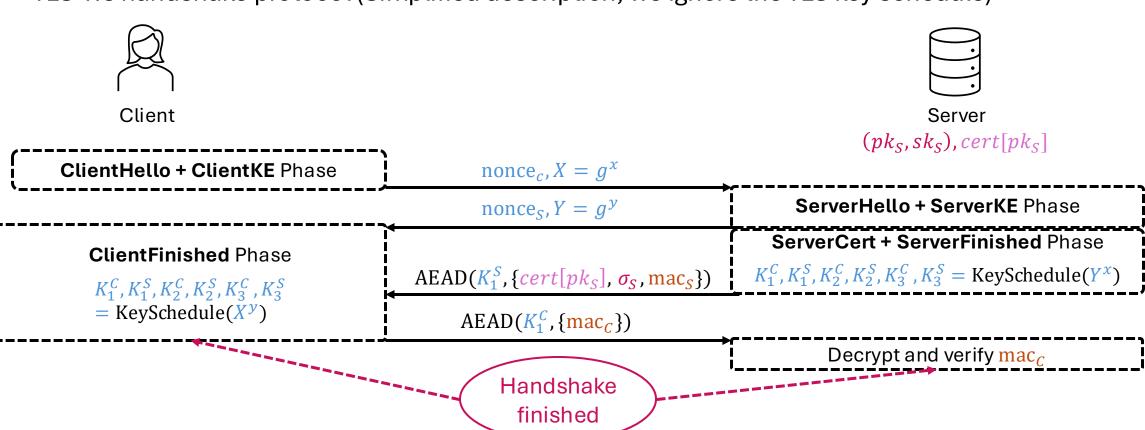
 $AEAD(K_1^S, \{cert[pk_S], \sigma_S, mac_S\})$

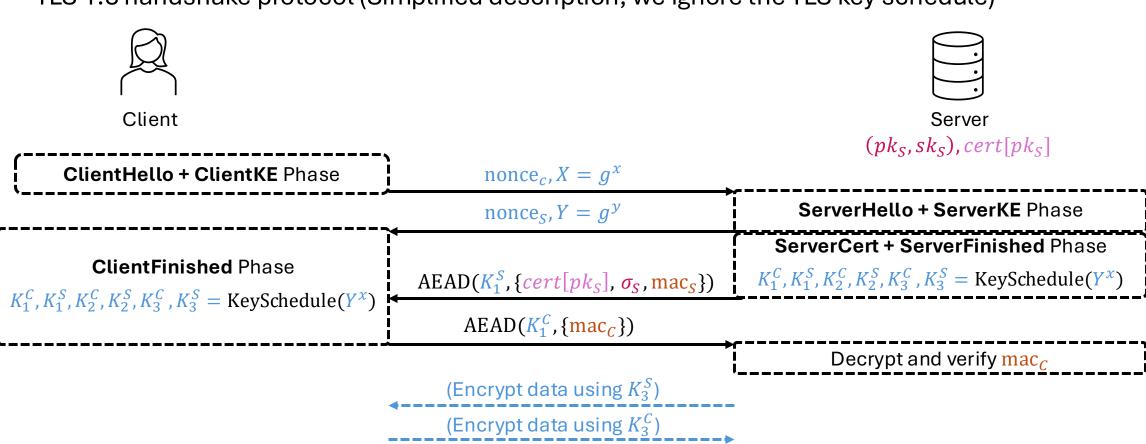
 $AEAD(K_1^C, \{mac_C\})$

ServerCert + ServerFinished Phase

$$K_2^C$$
, K_2^S = KeySchedule₂ $(Y^x, ...)$
 K_3^C , K_3^S = KeySchedule₃ $(Y^x, ...)$

Decrypt and verify mac_c





HTTP v.s. HTTPs



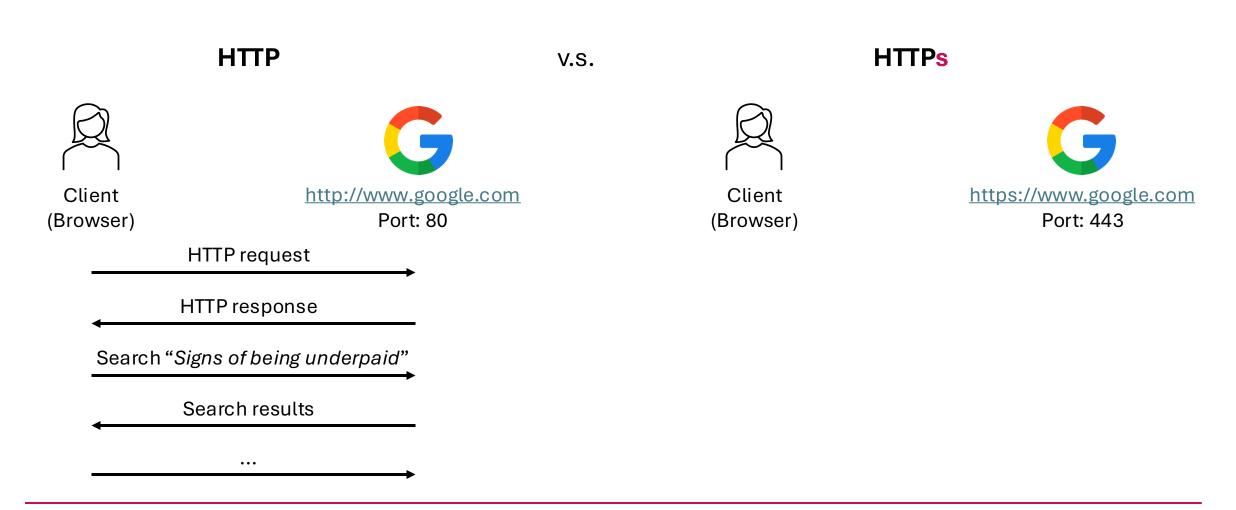


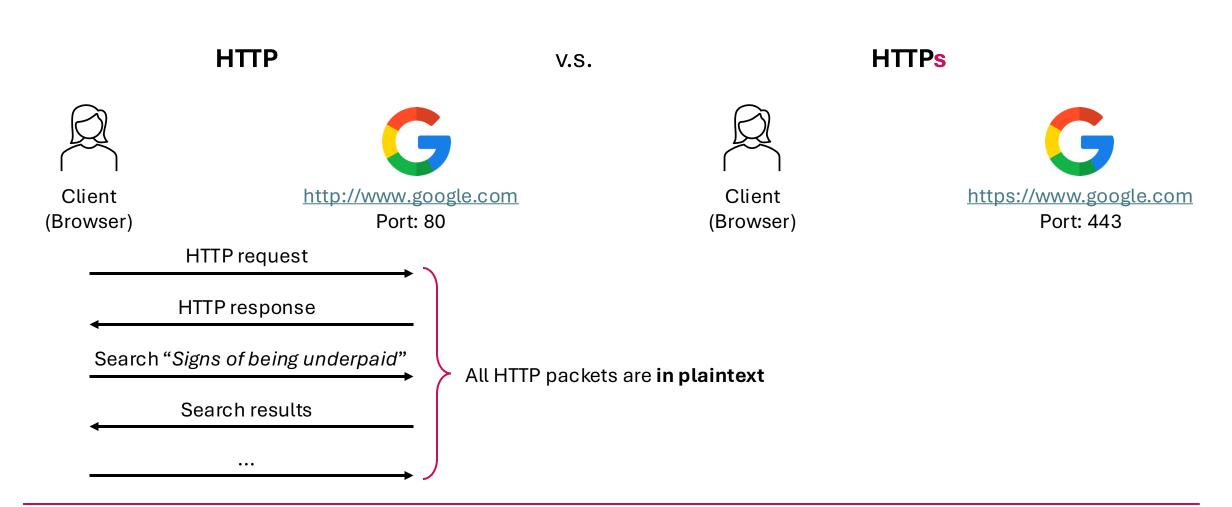


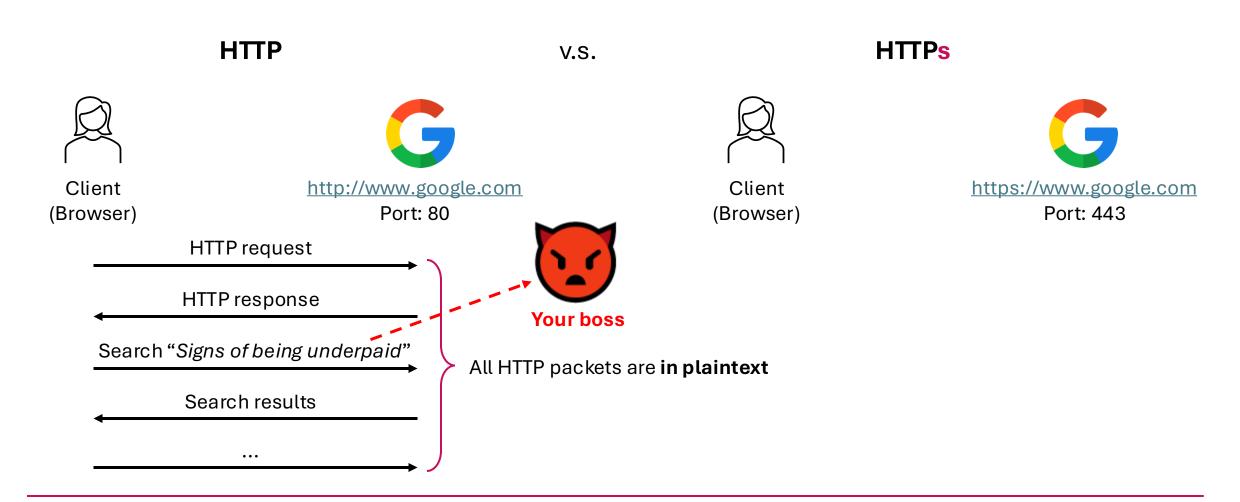


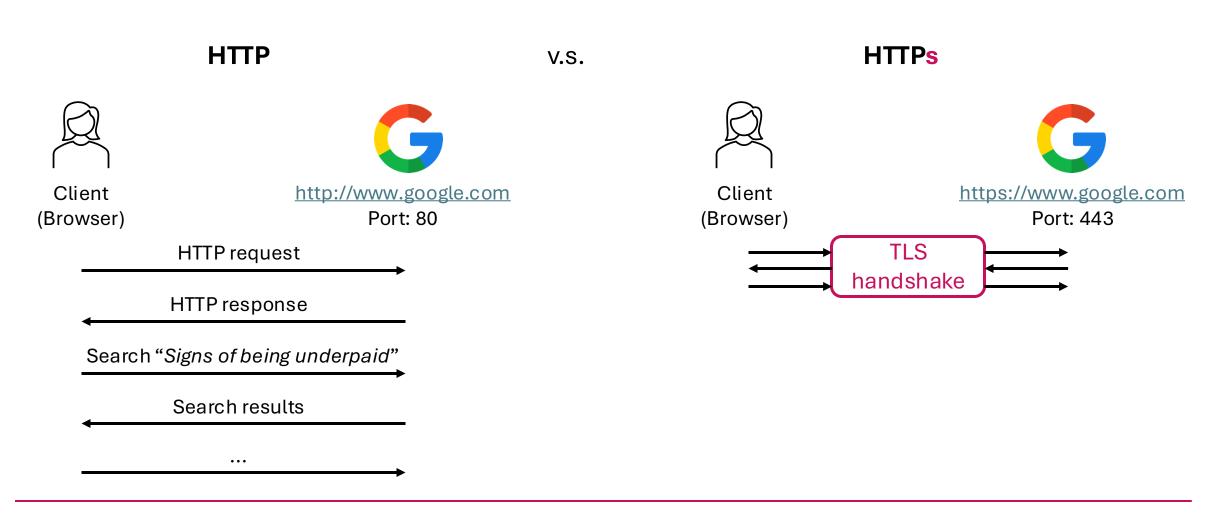
HTTP HTTPs V.S. Client http://www.google.com Client https://www.google.com (Browser) Port: 80 (Browser) Just an example. You probably cannot access http://www.google.com because your browser or Google enforces HTTPs connections.

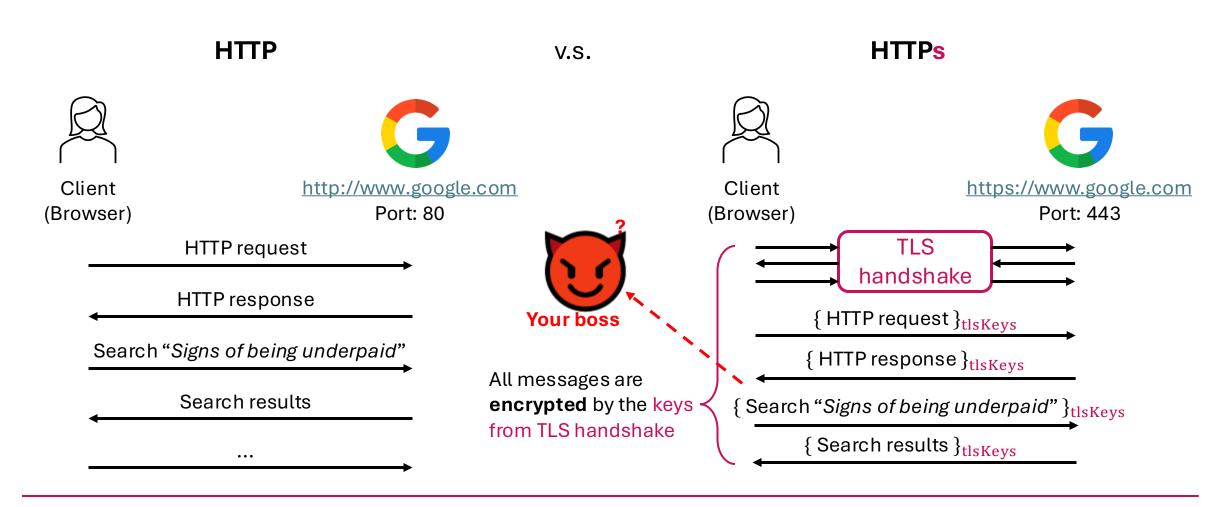
Port: 443





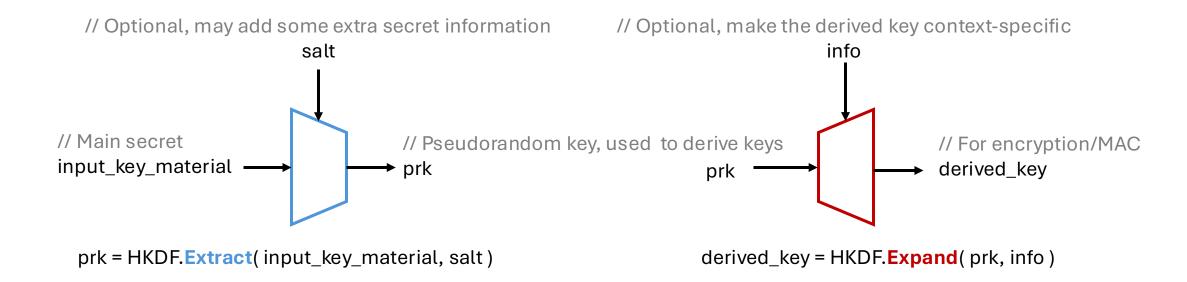






Coding Tasks

1. Run the example code "HKDF.py". Play with it and learn how to derive keys from a secret.



Homework

... (next page)

Warning: This key schedule scheme may not be secure. If you want to use TLS in real-world applications, please follow the TLS 1.3 standard

- Implement the tweaked TLS handshake protocol (in the Client-Server setting using sockets)
 - Use the simplified key schedule algorithm:

```
KeySchedule<sub>1</sub>(g^{xy}):
                                                                DeriveHS(q^{xy}):
1. HS = DeriveHS(g^{xy})
                                                                1. ES = HKDF.Extract(0, 0) // 0 = zeros (bytes) of length 32
2. K_1^C = HKDF.Expand(HS, SHA256("ClientKE"))
                                                                2. dES = HKDF.Expand(ES, SHA256("DerivedES"))
3. K_1^S = HKDF.Expand(HS, SHA256("ServerKE"))
                                                                3. HS = HKDF.Extract(dES, SHA256(q^{xy}))
4. return K_1^C, K_1^S
                                                                4. return HS
KeySchedule<sub>2</sub>(nonce<sub>C</sub>, X, nonce<sub>S</sub>, Y, g^{xy}):
1. HS = DeriveHS(g^{xy})
2. ClientKC = SHA256(nonce<sub>C</sub> || X || nonce<sub>S</sub> || Y || "ClientKC") // "||" is the concatenation operation
3. ServerKC = SHA256(nonce_C || X || nonce_S || Y || "ServerKC")
4. K_2^C = HKDF.Expand(HS, ClientKC)
5. K_2^S = HKDF.Expand(HS, ServerKC)
6. return K_2^C, K_2^S
```

Homework

Warning: This key schedule scheme may not be secure. If you want to use TLS in real-world applications, please follow the TLS 1.3 standard

```
KeySchedule<sub>3</sub>(nonce<sub>C</sub>, X, nonce<sub>S</sub>, Y, g^{xy}, \sigma, cert[pk_S], mac_S):

1. HS = DeriveHS(g^{xy})

2. dHS = HKDF.Expand(HS, SHA256("DerivedHS"))

3. MS = HKDF.Extract(dHS, 0) // 0 = zeros (bytes) of length 32

2. ClientSKH = SHA256(nonce<sub>C</sub> || X || nonce<sub>S</sub> || Y || \sigma || cert[pk_S] || mac_S|| "ClientEncK")

3. ServerSKH = SHA256(nonce<sub>C</sub> || X || nonce<sub>S</sub> || Y || \sigma || cert[pk_S] || mac_S|| "ServerEncK")

2. K_3^C = HKDF.Expand(MS, ClientSKH)

3. K_3^S = HKDF.Expand(MS, ServerSKH)

4. return K_3^C, K_3^S
```

How to compute the signature/MAC code:

```
For server: \sigma = Sign(sk_S, SHA256(nonce_C || X || nonce_S || Y || cert[pk_S])) // Use DSA with SHA256 and P256 For server: mac_S = HMAC(K_2^S, SHA256(nonce_C || X || nonce_S || Y || \sigma || cert[pk_S] || "ServerMAC")) For client: mac_C = HMAC(K_2^C, SHA256(nonce_C || X || nonce_S || Y || \sigma || cert[pk_S] || "ClientMAC"))
```

 How to verify HMAC: To verify if mac is the valid HMAC code of M with respect to the key K, Just check: mac =? HMAC(K, M)

Homework

• **Bonus:** Implement the same protocol, but this time use SHA3-512 as the hash function (for HKDF, HMAC, and the key schedule) and P-521 as the elliptic curve for key exchange. This should allow you to derive a key with 512 bits (64 bytes).

Further Reading

- RFC 8446 TLS 1.3: https://datatracker.ietf.org/doc/html/rfc8446
- RFC 2818 HTTP over TLS: https://datatracker.ietf.org/doc/html/rfc2818
- Felix Günther's lecture notes on TLS 1.3: https://www.felixguenther.info/teaching/2019-tls-seminar_02-21_TLS13-intro-MSKE-MKC.pdf
- Cryptography analysis of TLS 1.3 handshake: https://eprint.iacr.org/2020/1044

