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Daily Coding Problem #149

Problem

This problem was asked by Goldman Sachs.

Given a list of numbers L , implement a method `sum(i, j)` which returns the sum from the sublist $L[i:j]$ (including i , excluding j).

For example, given $L = [1, 2, 3, 4, 5]$, `sum(1, 3)` should return `sum([2, 3])`, which is 5.

You can assume that you can do some pre-processing. `sum()` should be optimized over the pre-processing step.

Solution

This problem has many different possible tradeoffs you could make:

- Have no pre-processing step and compute sum on the fly.

- Pre-process the sum of every i, j combination s.t. $i < j$ and store it in a list.
- Create a **segment tree** representing sums of many segments of the list.

Let's explore the segment tree since it has a good balance of pre-processing time and query time.

Basically, a segment tree's nodes would have these fields:

- `start_ind` and `end_ind` which represents the segment in the list this node represents.
- `val`, which represents the sum of this segment.
- `left`, for the left node.
- `right`, for the right node.

With these fields we can see how we might query for a segment.

- If the node's segment is equivalent to the query, return its `val`
- Otherwise, query for the left node if it encloses the starting index and add it to our result.
- Also query for the right node if it encloses the ending index and add it to our result.

```
result = 0

if left.start_ind <= start_ind <= left.end_ind:
    result += query(left, start_ind, min(end_ind, left.end_ind))

if right.start_ind <= end_ind <= right.end_ind:
    result += query(right, max(start_ind, right.start_ind), end_ind)
```

The full code:

```

class Node:
    def __init__(self, val, start_ind, end_ind, left=None, right=None):
        self.val = val
        self.start_ind = start_ind
        self.end_ind = end_ind
        self.left = left
        self.right = right

    @property
    def interval(self):
        return (self.start_ind, self.end_ind)

def make_segment_tree(lst):
    return _make_segment_tree(lst, 0, len(lst) - 1)

def _make_segment_tree(lst, start_ind, end_ind):
    if start_ind == end_ind:
        assert(len(lst) == 1)
        val = lst[0]
        return Node(val, start_ind, end_ind)

    mid = len(lst) // 2

    left = _make_segment_tree(lst[:mid], start_ind, start_ind + mid - 1)
    right = _make_segment_tree(lst[mid:], start_ind + mid, end_ind)

    root_val = left.val + right.val

    return Node(root_val, start_ind, end_ind, left, right)

def query(node, start_ind, end_ind):
    if node.start_ind == start_ind and node.end_ind == end_ind:
        return node.val

    result = 0

```

```
left = node.left
right = node.right

if left.start_ind <= start_ind <= left.end_ind:
    result += query(left, start_ind, min(end_ind, left.end_ind))

if right.start_ind <= end_ind <= right.end_ind:
    result += query(right, max(start_ind, right.start_ind), end_ind)

return result
```

This takes $O(N \log N)$ time and space during pre-processing while querying takes $O(\log N)$ time.

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