Blog

## Daily Coding Problem #218

## **Problem**

This problem was asked by Yahoo.

Write an algorithm that computes the reversal of a directed graph. For example, if a graph consists of  $A \rightarrow B \rightarrow C$ , it should become  $A \leftarrow B \leftarrow C$ .

## **Solution**

Our first instinct might be to traverse the graph, finding all the edges. For each edge, we can swap its vertices and add the updated edge to a new graph.

For example, if we find an edge directed from A to B, with a weight of 5, we would add a corresponding edge to our transposed graph with the same weight, but from B to A. Traversing the graph can be done with depth-first search.

```
from collections import defaultdict

def helper(graph, node, visited, transpose):
   if node not in visited:
```

```
visited.add(node)

for n, weight in graph[node]:
    transpose[n] += [(node, weight)]
    helper(graph, n, visited, transpose)

return transpose

def reverse(graph):
    start = graph.keys()[0]
    transpose = helper(graph, start, set(), defaultdict(list))
    return transpose
```

This will take O(|V| + |E|) time, and building up the reversed graph will use O(|V| + |E|) space.

However, there is a simpler approach. If we already have our graph represented as an adjacency list, all we need to do is generate a new adjacency list with swapped keys and values.

In other words, suppose we our graph is represented as {0: [2, 3], 1: [3, 4], 2: [1]}. We iterate through each key k, and every value v. If v is not a key in our transposed graph, we add it, along with the value k. Otherwise, we append v to the existing list of k's edges. In this example the first few edges would be {2: 0, 3: 0, 3: 1, ...}.

```
def reverse(graph):
    transpose = defaultdict(list)

for node in graph:
    for neighbor, weight in graph[node]:
        transpose[neighbor].append((node, weight))

return transpose
```

The time and space complexity is unchanged from above.

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