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# Daily Coding Problem #318

## Problem

This problem was asked by Apple.

You are going on a road trip, and would like to create a suitable music playlist. The trip will require  $N$  songs, though you only have  $M$  songs downloaded, where  $M < N$ . A valid playlist should select each song at least once, and guarantee a buffer of  $B$  songs between repeats.

Given  $N$ ,  $M$ , and  $B$ , determine the number of valid playlists.

## Solution

First let us consider a simpler problem: creating playlists of  $N$  songs from  $M$  downloaded options, but without any buffer requirement.

We can construct a solution using dynamic programming. In particular we would like to construct a matrix, ways, such that  $\text{ways}[i][j]$  represents the number of ways of making a playlist of length  $i$  with  $j$  unique songs.

Our base case is when  $i$  and  $j$  are both zero. Here, we can consider there to be only one trivial option: a playlist with no songs. Otherwise, we can find the value of  $\text{ways}[i][j]$  by dividing into two cases: when the last song is new, and when the last song is a repeat.

If the last song is new, it could be any of the  $m - (j - 1)$  unused songs. Each of these

songs can be combined with any of the playlists with one fewer song and one fewer unique song, or  $\text{ways}[i - 1][j - 1]$ . On the other hand, if the last song is a repeat, it must be one of the  $j$  songs already chosen. Each of these songs can be combined with any option provided by  $\text{ways}[i - 1][j]$ .

Once we built up our matrix in this way, we can take the value of  $\text{ways}[n][m]$  to be our solution. We can implement this as follows.

```
def valid_playlists(n, m, b):
    ways = [[0 for _ in range(m + 1)] for _ in range(n + 1)]
    ways[0][0] = 1

    for i in range(1, n + 1):
        for j in range(1, m + 1):
            ways[i][j] = ways[i - 1][j - 1] * (m - (j - 1)) + ways[i - 1][j] * j

    return ways[n][m]
```

Now let us try adding in the buffer  $B$  to our dynamic programming formula. If the last song is new, no change needs to be made, since it cannot possibly be a repeat.

If the song is old, on the other hand, there will be  $B$  options that cannot be used as the next song, specifically the last  $B$  songs in our playlist. Therefore the number of new playlist formed can be represented as  $\text{ways}[i - 1][j] * (j - b)$ . If the buffer is bigger than the number of distinct songs played so far, no repeat songs are possible, so we can use  $\max(j - b, 0)$  to handle this case.

```
def valid_playlists(n, m, b):
    ways = [[0 for _ in range(m + 1)] for _ in range(n + 1)]
    ways[0][0] = 1

    for i in range(1, n + 1):
        for j in range(1, m + 1):
            ways[i][j] = ways[i - 1][j - 1] * (m - (j - 1)) + ways[i - 1][j] * max(j - b,
0)

    return ways[n][m]
```

The time and space complexity of this algorithm will be  $O(M * N)$ , since we must loop through each cell of our  $M \times N$  matrix and perform a few calculations.

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