



→ End user

IDE

Chrome

media player.

shell

Hello.c

⇒ Appl<sup>n</sup>  
sw

OS ⇒ Core OS + utility sw + Appl<sup>n</sup> sw  
↓  
Kernel

(control panel, task mgr) (Notepad, paint, ms office, calendar...)

...

CPU

RAM

HDD

Key-board  
(primary) I/O

monitor  
(primary) O/P

etc.

⇒ Hardware  
Resources

① Interface.

② Control prog.

③ Resource Allocation / manager.

End user.

④ Bootable storage device (Has contain OS  
setup)  
CD / DVD / PD / HDD

## \* Functions of Core os / kernel

① Process management

② CPU Scheduling

③ Memory mg<sup>m</sup>

④ file & I/o mg<sup>m</sup>

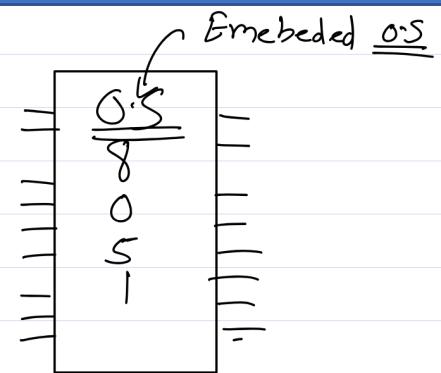
⑤ Hardware Abstraction.

⑥ Networking

⑦ Security and protection

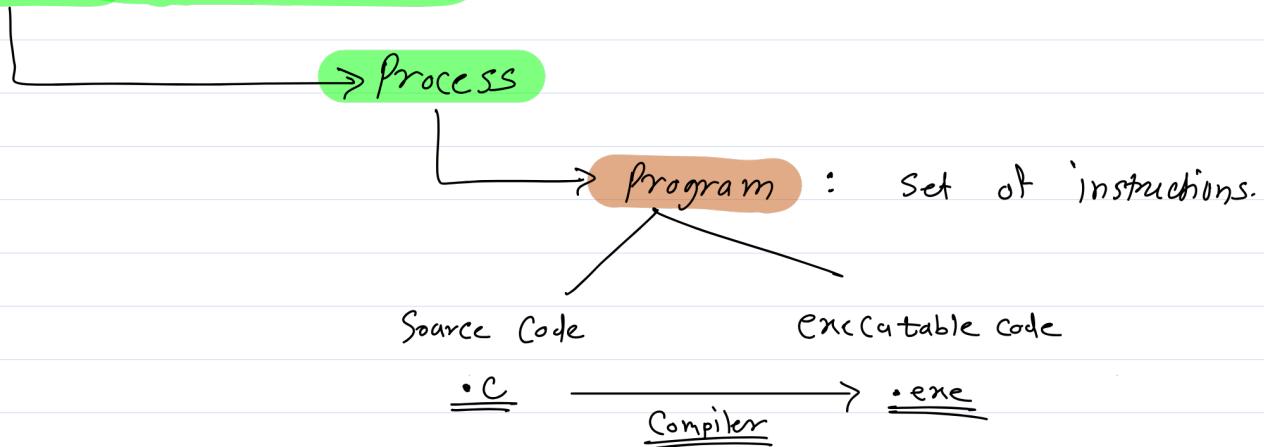
⑧ User interface.

Core Function of OS

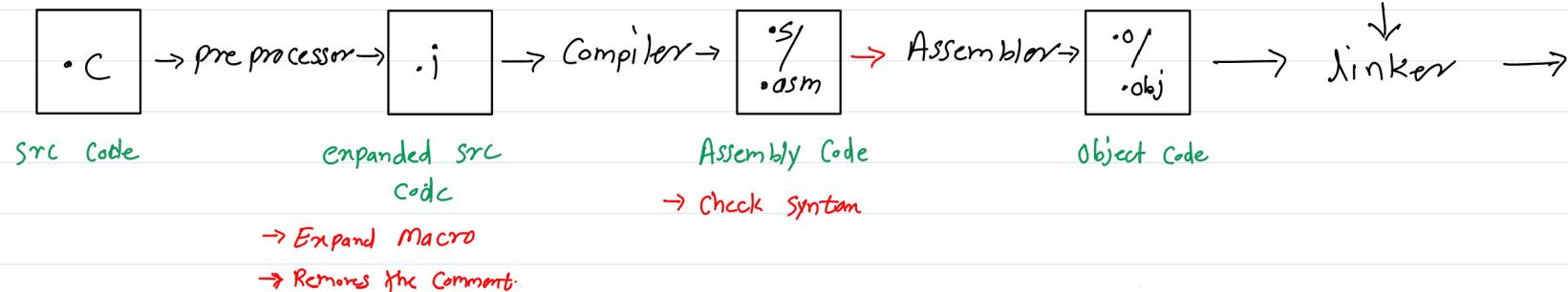


Extra utility fun of OS.

## \* **Process Management**



## \* Compilation steps:



a.exe => Windows  
a.out => Linux/ MAC



(section binary)

• out

exe header or  
primary header

text / code

data

BSS

RO data

Symbol  
table

→ Entry point function address (main)

Information about the Remaining section.

→ Magic Number → uniq Number of file format  
→ Identity of file format.

Char \*P = "Sunbeam"  $\Rightarrow$  RO

\*P = 'x' ; // Error  
putSC P )

machine code

initialised global & static variable

(block started by symbol)

- uninitialised global & static variable

" Constant string literals, 0,1,2, 'A','B' ...

→ The information Related of Variable & Function.

- Name, address, section, size, flags
- These info use for debugging purpose.

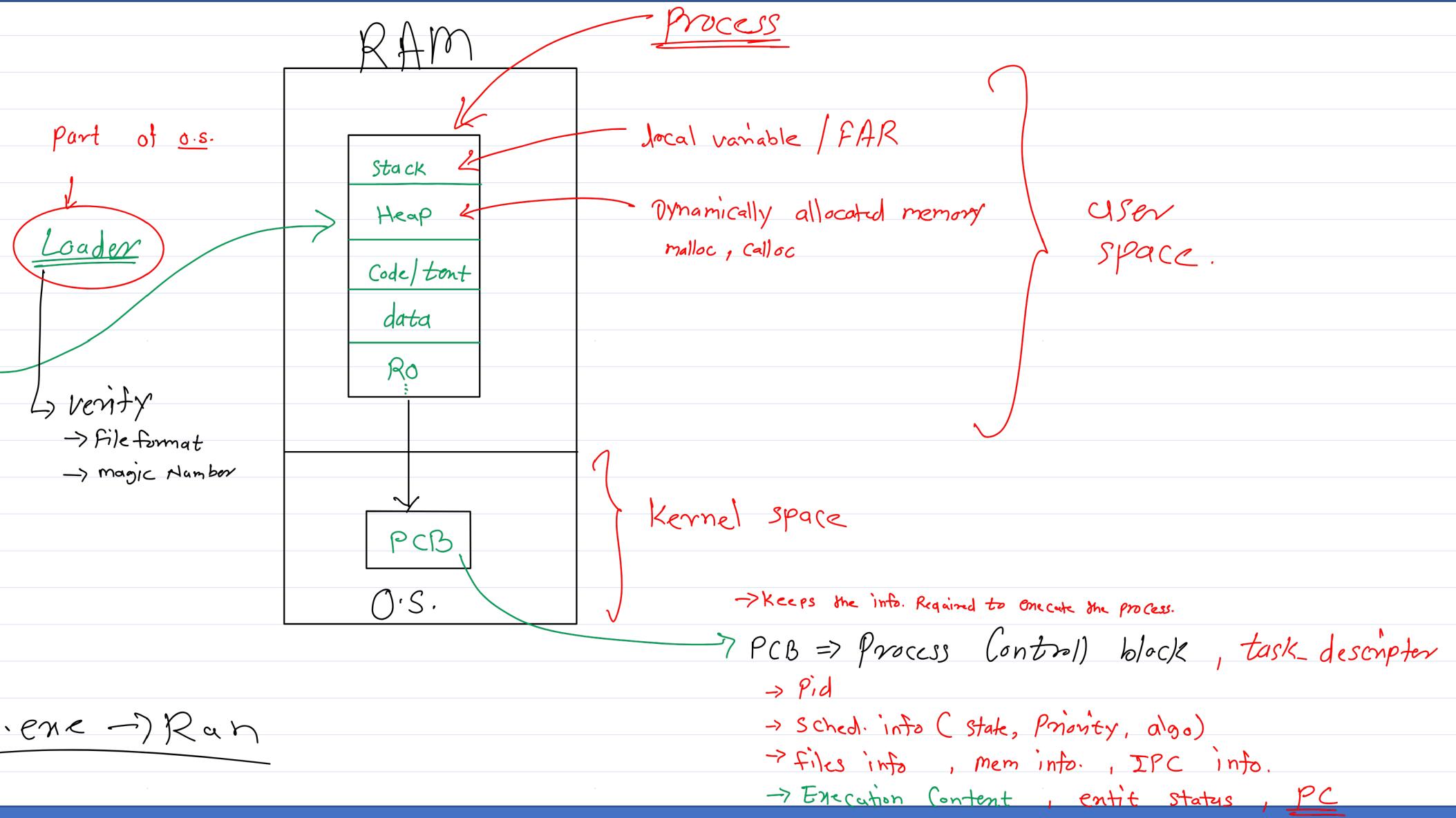
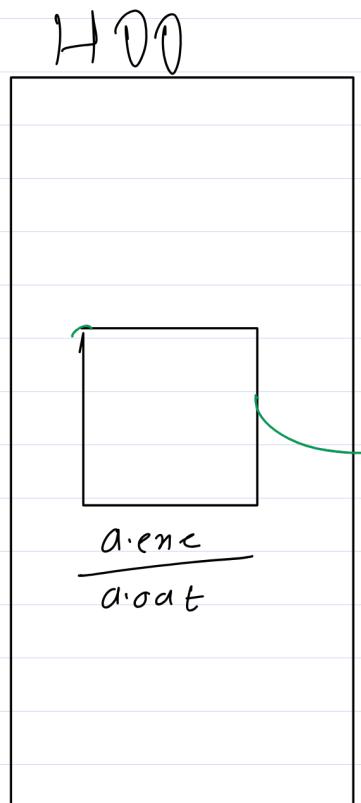
Binary file has magic Number

(It present on starting 2 or 4 bytes of  
of code file)

a • exe  $\rightarrow$  MZ  
a • out  $\rightarrow$  ?ELF  
• bmp  $\rightarrow$   
• pdf  $\rightarrow$   
• text  $\rightarrow$  X

Portable Executable (PE)

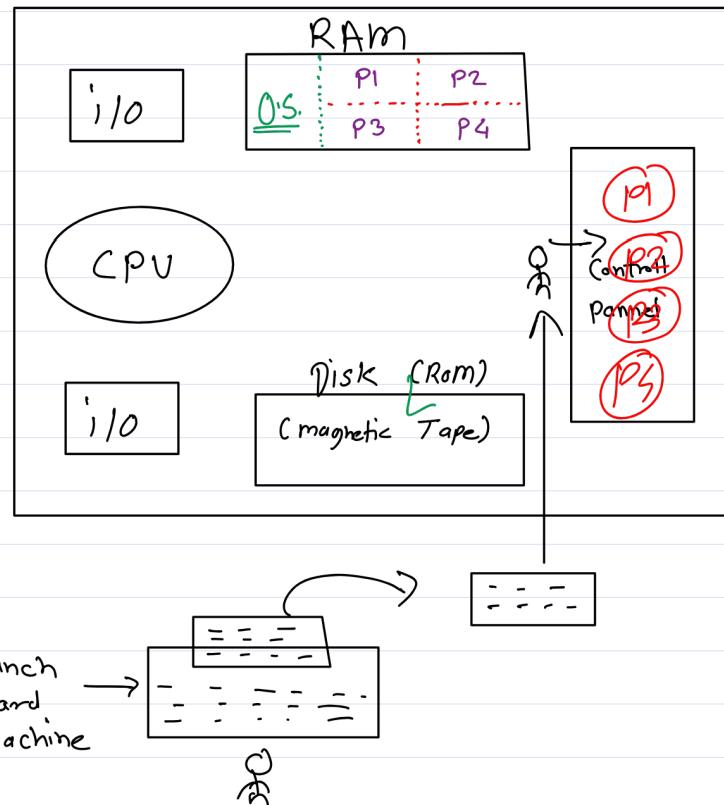
Executable linking format (ELF)



# \* OS Evolution.

## Mainframe Computer

(old computer)



① Resident monitor.

② Batch System

③ Multi-Programming.

- loading multiple programs in main memory.

- Here on RAM load mixed program

↳ CPU bound + I/O bound

- better utilization of CPU

- degree of M.P.

⇒ num of programs that can be kept in RAM.

$sf(n_1)$   
 $sf(n_2)$ 
} I/O instr

$add = n_1 + n_2$   
 $sub = n_1 - n_2$   
 $mul = n_1 * n_2$ 
} CPU instr

$pf(add)$   
 $pf(sub)$   
 $pf(mul)$ 
} I/O instr

$$\text{Program} = \frac{\text{CPU instr}}{\text{CPU burst Time}} + \frac{\text{I/O instr}}{\text{I/O burst time}}$$

$$\left( \frac{\text{CPU burst Time}}{\text{CPU burst Time}} \right) + \left( \frac{\text{I/O burst time}}{\text{I/O burst time}} \right)$$



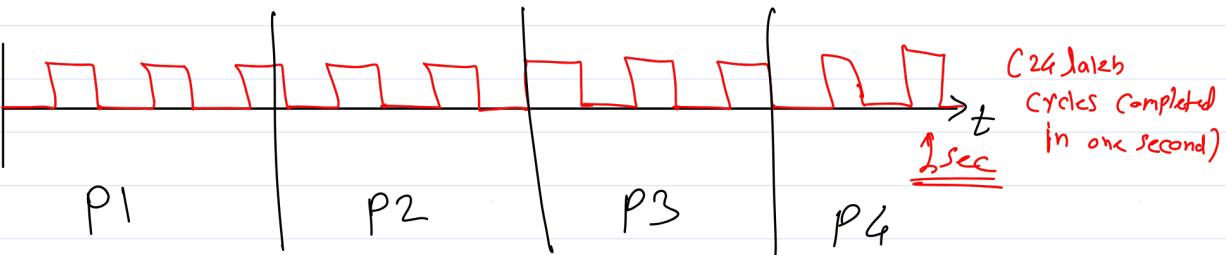
## ④ Time sharing / multitasking O.S.

(2.4GHz)

\* freq : Num of Cycles completed in one sec.

⇒ Sharing CPU time among multiple tasks present in main memory & Ready for execution.

⇒ Response time < 1 sec



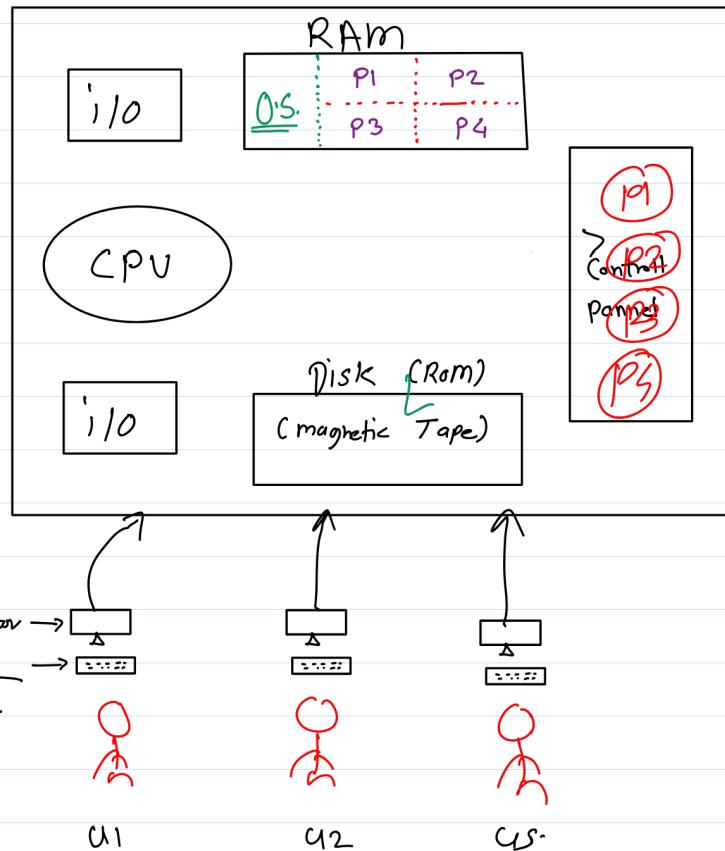
multi - tasking

① Process based  
M.T

② Thread based  
M.T

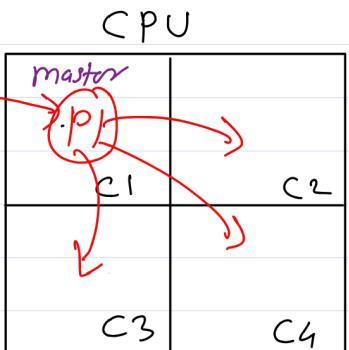
## \* Multi-user System.

puter)

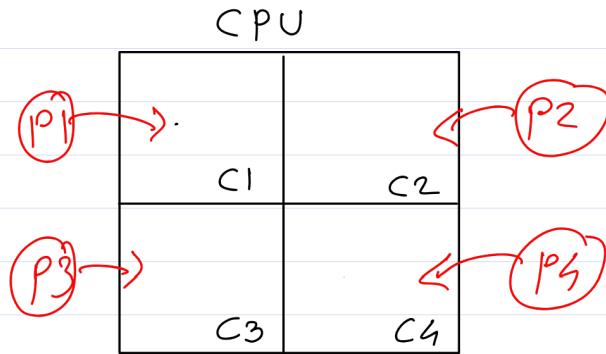


## \* Multi-core / multi-processor.

OS



Asymmetric M.P.

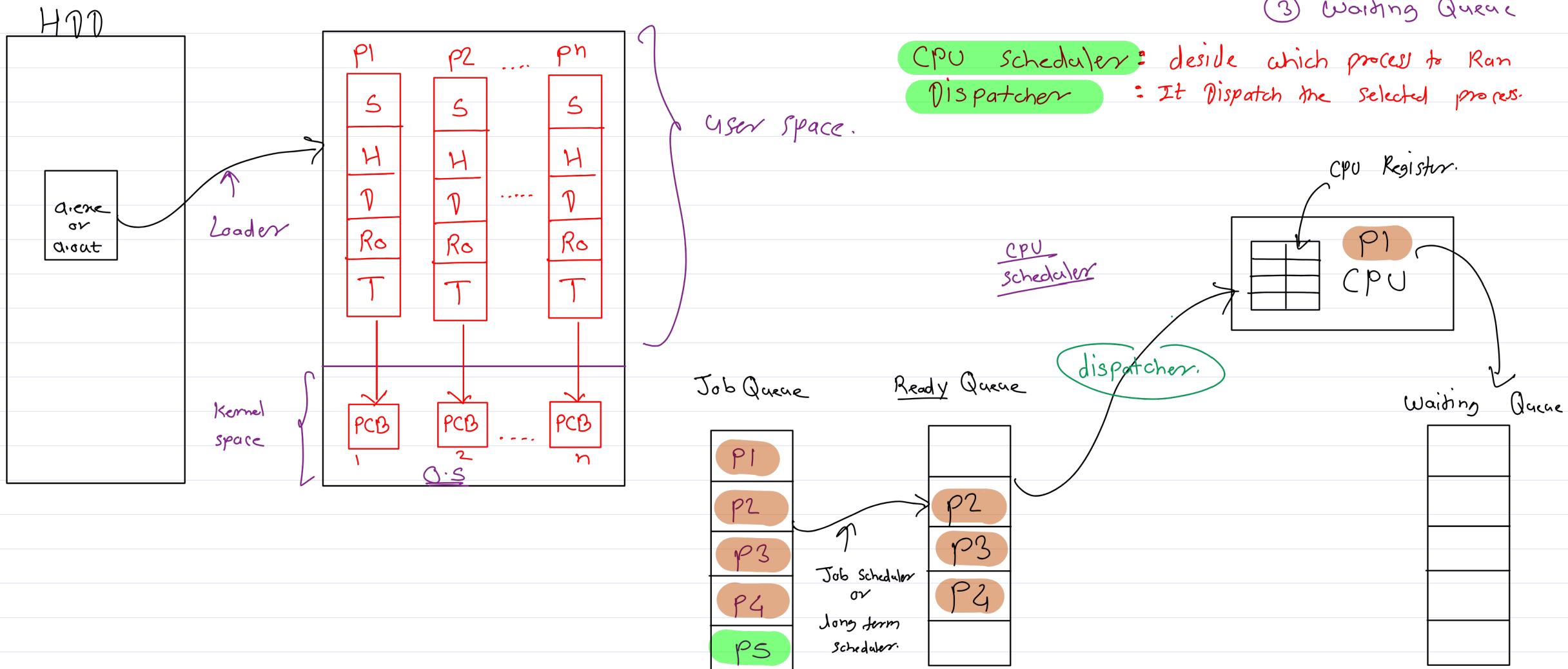


Symmetric M.P.

## Kernel data structure.

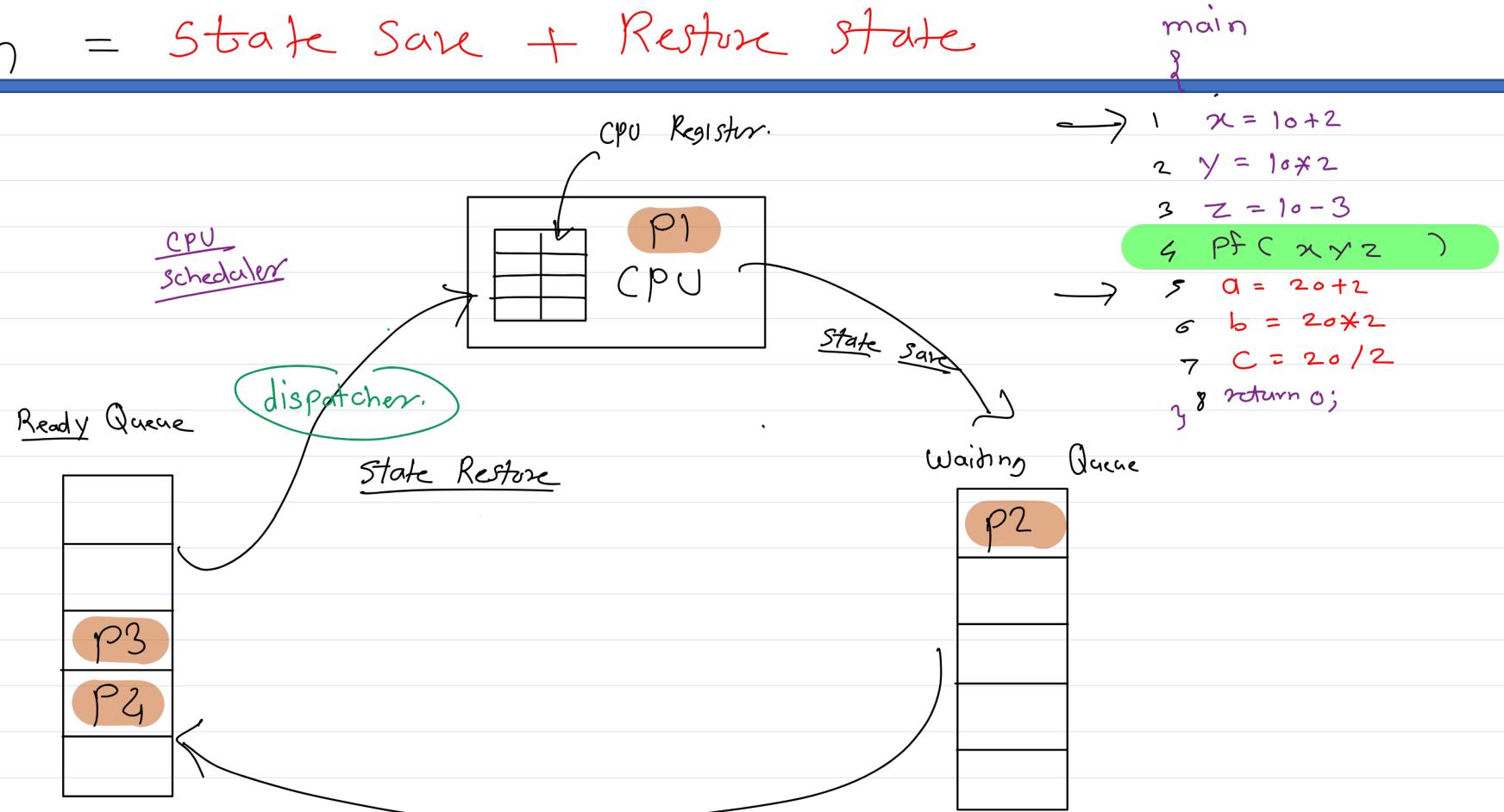
- ① Job Queue
- ② Ready Queue
- ③ Waiting Queue

## \* Process life Cycle



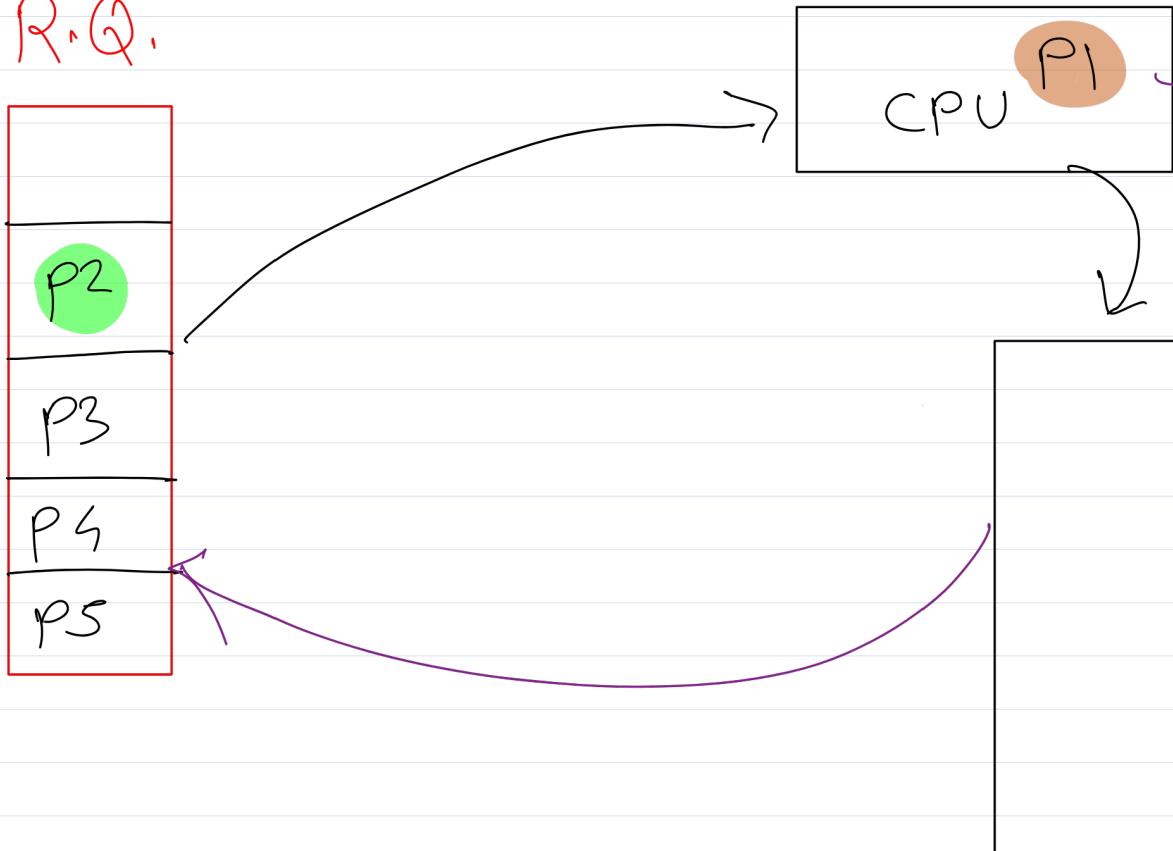
\* Content Switch = State Save + Restore State

- PCB
- Pid
  - Scheduling info
  - Mem info
  - PC
  - Execution Content.



# A) CPU Schedulers

R.Q.



W.Q.

GT  
PXO

- ① Running → Terminated
- ② Running → waiting state.
- ③ Running → Ready state
- ④ Waiting → Ready.

\* FCFS.

Convey effect  $\rightarrow$  If bigger process arrive first, Avg W.T. increase.

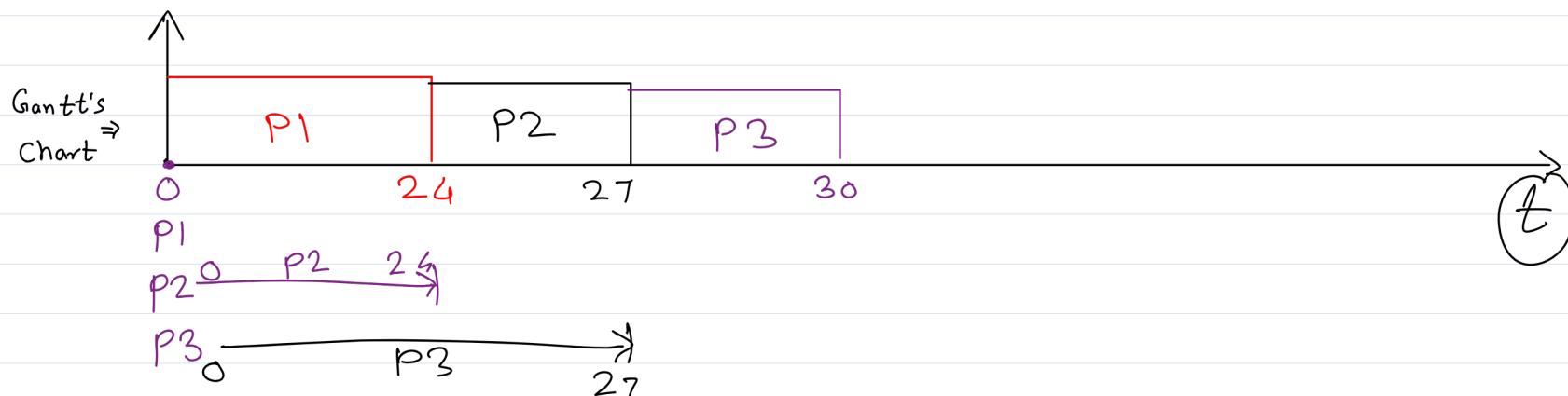
## ►FCFS Scheduling

Process	Arrival Time	CPU Burst	Wait Time	Turn Around Time
P1 X	0	24	0	24
(P2) X	0	3	-24	27
P3 X	0	3	-27	30

Avg W.T. =  $\frac{\text{Sum of all process W.T.}}{\text{Num of process}}$

$$= \frac{0 + 24 + 27}{3}$$

$$= \frac{51}{3} \Rightarrow 17$$



Avg. T.A.T =  $\frac{\text{Sum of all process T.A.T}}{\text{Num of process}}$

$$\Rightarrow \frac{24 + 27 + 30}{3}$$

$$\Rightarrow \frac{81}{3}$$

$$\Rightarrow 27$$

\* FCFS.

## ►FCFS Scheduling

Process	Arrival Time	CPU Burst	Wait Time	Turn Around Time
P1	0	3	0	3
P2	0	3	3	6
P3	0	24	6	30

No - Convoy effect.

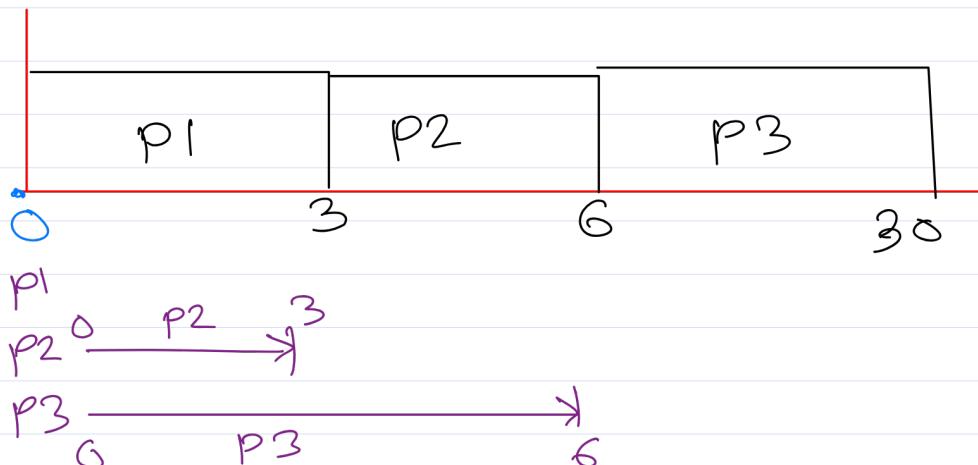
Completion time  
↓

$$\text{Avg. W.T.} \Rightarrow \frac{0+3+6}{3}$$

$$= \frac{9}{3} \\ = 3$$

$$\text{Avg. T.A.T} = \frac{3+6+30}{3}$$

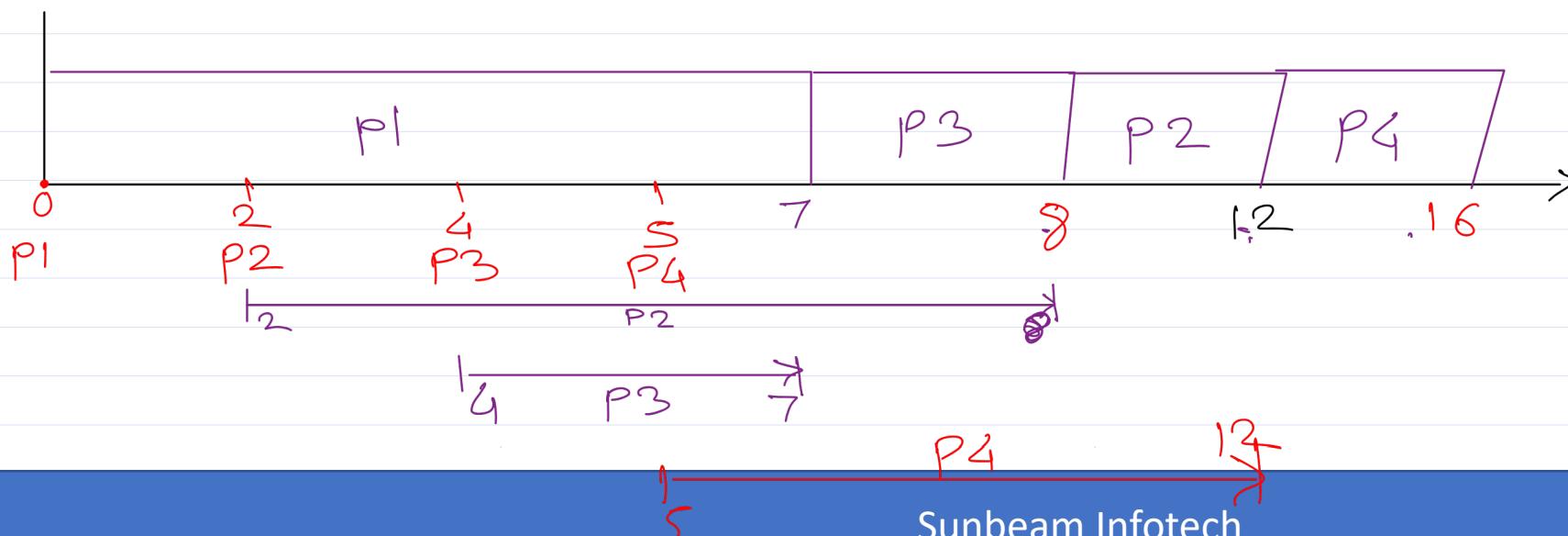
$$= \frac{39}{3} \\ \Rightarrow 13$$



SJF - gives min Avg wait time.

## ► SJF/SNTF Scheduling (Non-Premptive)

Process	Arrival Time	CPU Burst	Wait Time	Turn Around Time
P1	0	7 ✗	0	7
P2	2	4 ✗	6	10
P3	4	1 ✗	3	4
P4	5	4	7	11



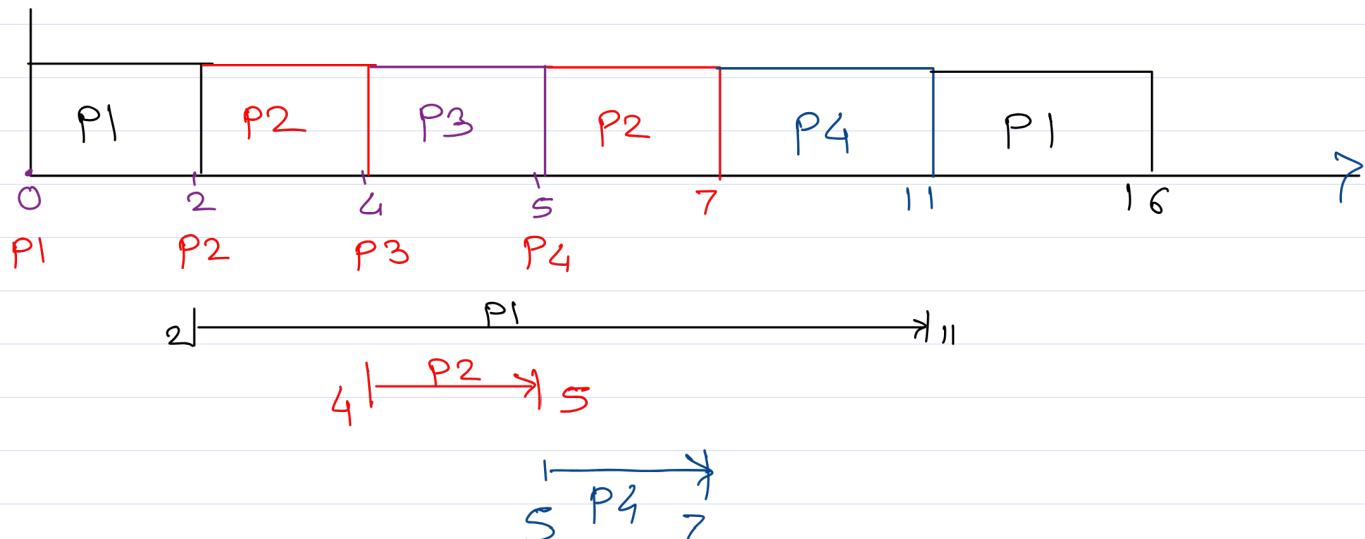
SJF  
↓

## ➤ SRTF Scheduling (Premptive)

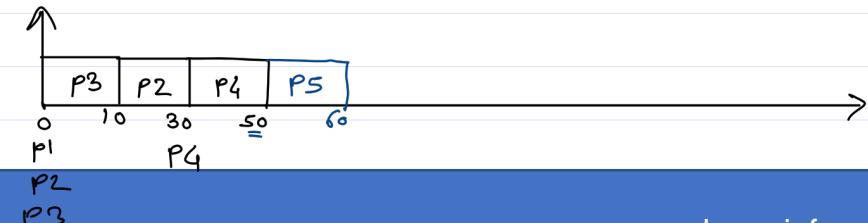
Process	Arrival Time	CPU Burst	Remaining Time	Wait Time	Turn Around Time
P1	0	7	5	9	16
P2	2	4	2 ✗	1	5
P3	4	1	✗	0	12
P4	5	4	✗	2	6

Avg. W.T. =>

Avg. T.A.T. =



A.T	Process	C.B.T	* Starvation
			blocked
0	P1	100	(Starvation)
0	P2	20 ✗	
0	P3	10 ✗	
30	P4	20 ✗	
50	P5	10	

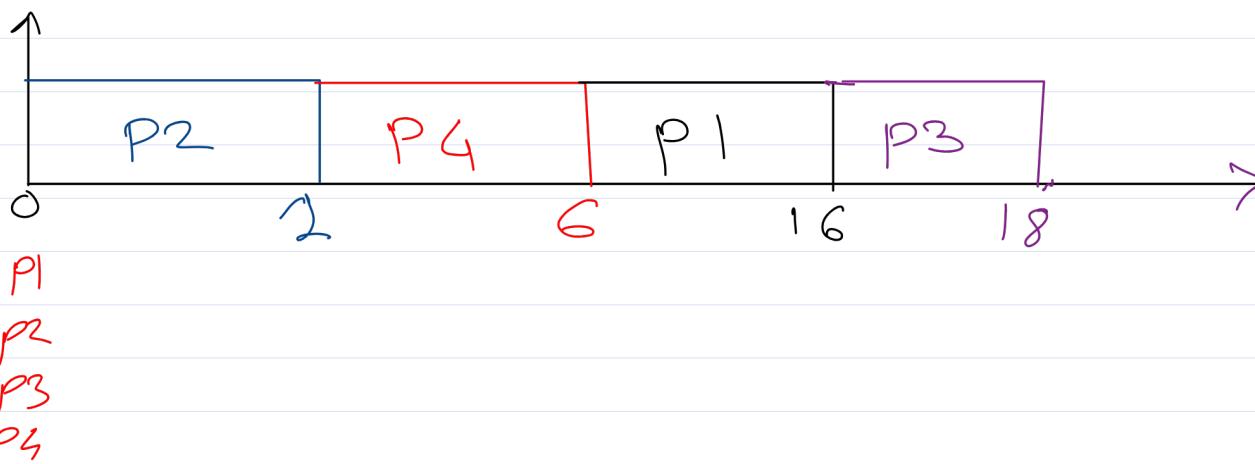


# ➤ Priority Scheduling

min → high priority

max → low priority

Process	Arrival Time	CPU Burst	Priority	Wait Time
P1	0	10 X	3	6
P2	0	1 X	1 - high	0
P3	0	2 ↴	4 - low	16
P4	0	5 X	2	1



T.A.T

16  
2  
18  
6

\* Starvation  $\xrightarrow{\text{soln}} \text{Ageing}$

blocked (starvation)

A.T.	Process	C.B.T	Priority
0	P1	10	20 $\rightarrow$ 19 $\rightarrow$ 18.
0	P2	10	5 X
0	P3	10	1 X
20	P4	10	7 X
30	P5	10	10



## ➤ Round Robin Scheduling

Process	CPU Burst	Remaining Time	Wait Time	Response Time
P1	53			
P2	17			
P3	68			
P4	24			



# SJF → Non-Premptive (SJTF)

Process	A.T.	C.B.T		W.T + C.B.T	W.T.	T.A.T
P1	1	7	X	0	7	
P2	2	5	3 X	9	14	
P3	3	1	1 X	5	6	
P4	4	2	2 X	5	7	
P5	5	8	4	11	19	

$$\text{Avg. T.A.T.} \Rightarrow \frac{7+14+6+7+19}{5}$$

$$= 53/5$$

$$= \underline{\underline{10.6}}$$

