String Similarity Join

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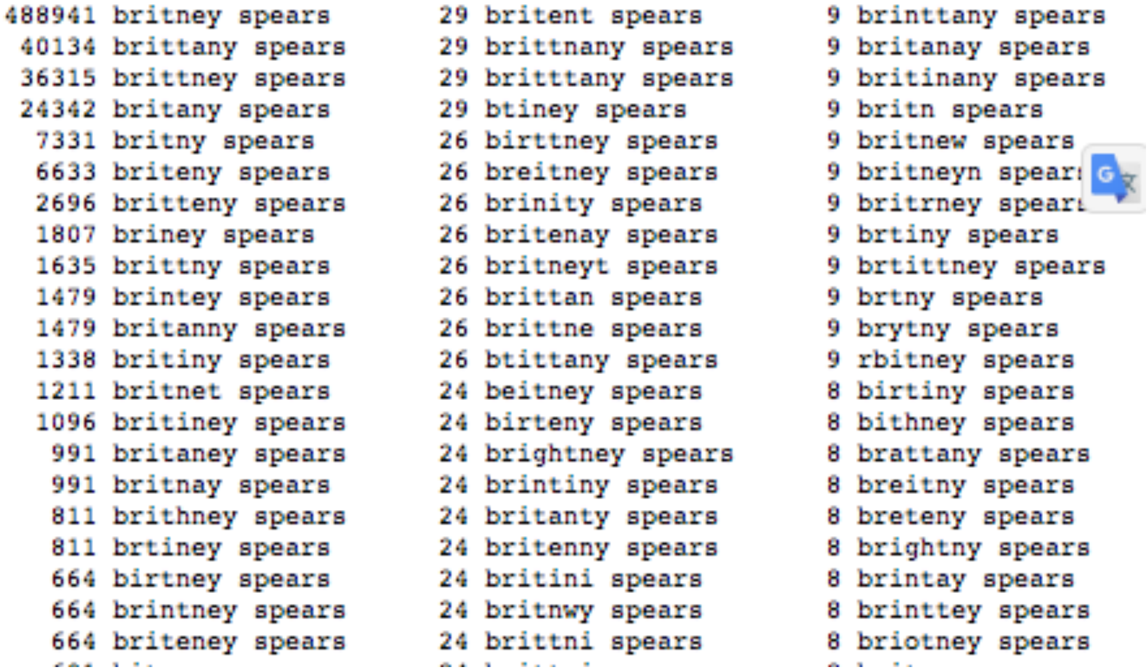
Real World Data is Dirty

- Misspellings of the query "*britney spears*" on Google



```
488941 britney spears        29 britent spears          9 brinttany spears
 40134 brittany spears       29 brittnany spears        9 britanay spears
 36315 brittney spears       29 britttany spears        9 britinany spears
 24342 britany spears        29 btiney spears           9 britn spears
  7331 britny spears         26 birttney spears         9 britnew spears
  6633 briteny spears        26 breitney spears         9 britneyn spear
  2696 britteny spears       26 brinity spears          9 britrney spear
  1807 briney spears         26 britenay spears         9 brtiny spears
  1635 brittny spears        26 britneyt spears         9 brtittney spears
  1479 brintey spears        26 brittan spears          9 brtny spears
  1479 britanny spears       26 brittne spears          9 brytny spears
  1338 britiny spears        26 btittany spears         9 rbitney spears
  1211 britnet spears        24 beitney spears          8 birtiny spears
  1096 britiney spears       24 birteny spears          8 bithney spears
   991 britaney spears       24 brightney spears        8 brattany spears
   991 britnay spears        24 brintiny spears         8 breitny spears
   811 brithney spears       24 britanty spears         8 breteny spears
   811 brtiney spears        24 britenny spears         8 brightny spears
   664 birtney spears        24 britini spears          8 brintay spears
   664 brintney spears       24 britnwy spears          8 brinttey spears
   664 briteney spears       24 brittni spears          8 briotney spears
```

http://marc.merlins.org/linux/talks/google/britney.html

Set Similarity Functions

- Overlap Size

$$overlap(x, y) = |x \cap y|$$

- Jaccard Similarity

$$J(x, y) = \frac{|x \cap y|}{|x \cup y|}$$

x **= {A,B,C,D,E}**

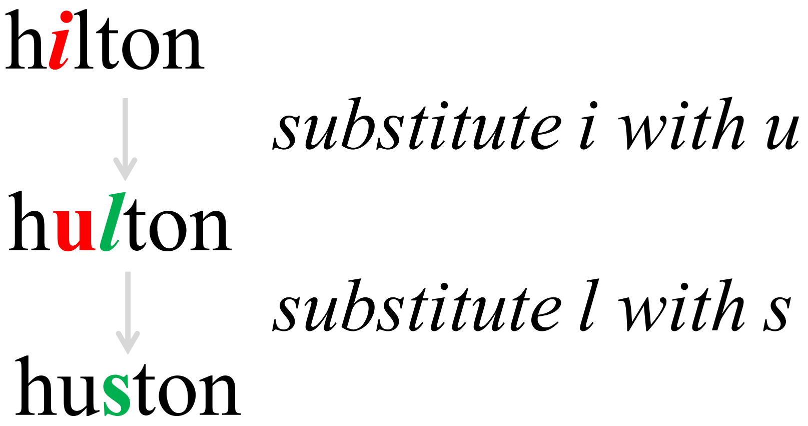
y **= {B,C,D,E,F}**

overlap(x, y) = 4

J(x, y) = 4/6 = 0.67

String Similarity Function

- Edit Distance ED(*r*, *s*): the minimum number of edit operations (insertion/deletion/substitution) needed to transform *r* to *s*.
- For example: ED(*hilton*, huston) = 2



h*i*lton

substitute i with u

h**u***l*ton

substitute l with s

hu**s**ton

- Edit Similarity: $\mathrm{EDS}(r,s) = 1 - \frac{ED(r,s)}{\max(|r|,|s|)}$

Calculating Edit Distance

Calculating ED(r,s)

Let r_n and s_m be the last characters in r and s.

a) Match r_n and s_m

ED(r,s)= 0 + ED(r[1,n-1],s[1,m-1])

b) Substitute r_n with s_m

ED(r,s)= 1 + ED(r[1,n-1],s[1,m-1])

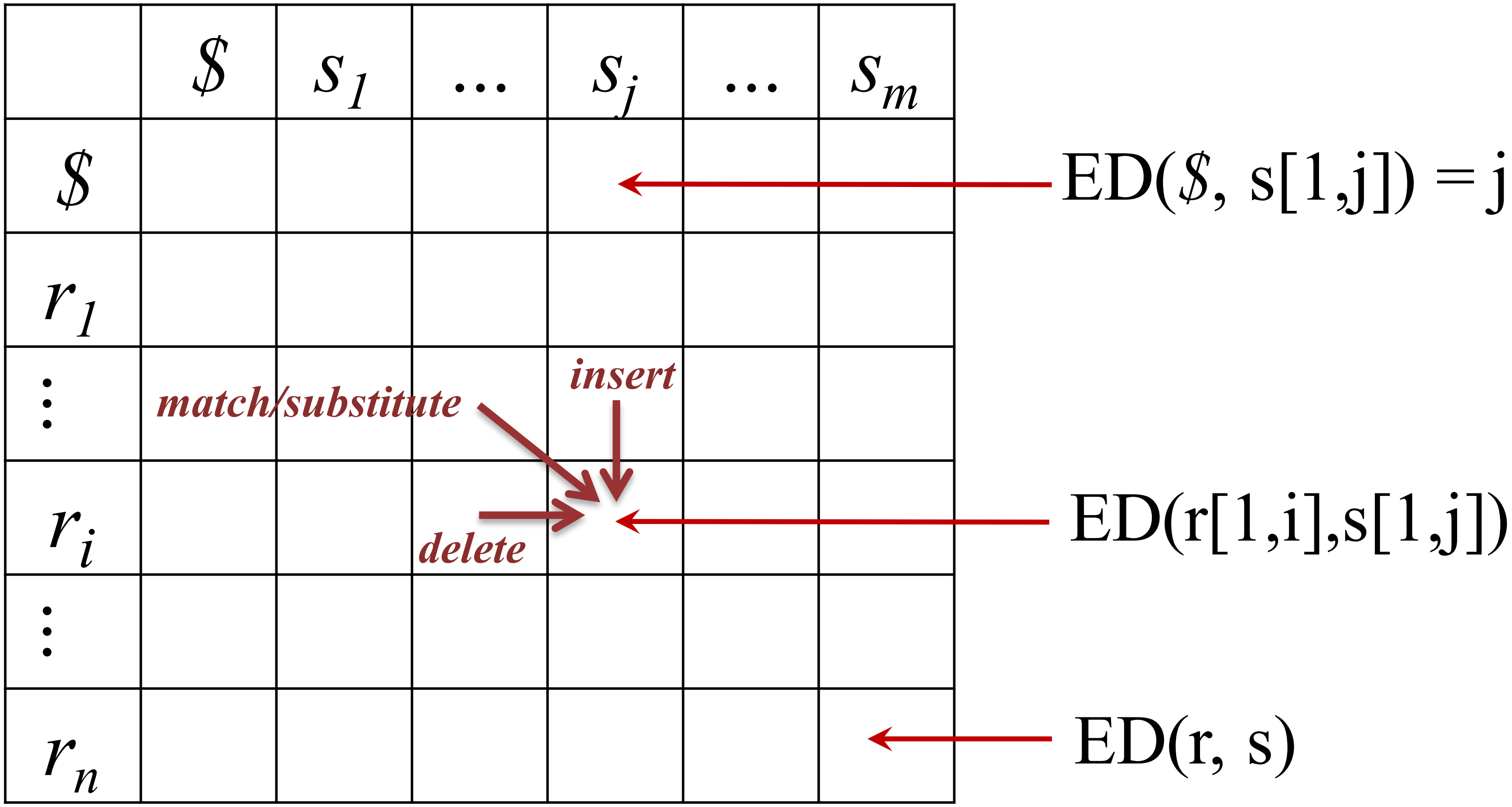
c) Delete r_n

ED(r,s)= 1 + ED(r[1,n-1],s[1,m])

d) Insert s_m

ED(r,s)= 1 + ED(r[1,n],s[1,m-1])

Calculating Edit Distance



ED($, s[1,j]) = j

ED(r[1,i],s[1,j])

ED(r, s)

Calculating Edit Distance

	$	b	r	i	t	n	e	y
$	0	1	2	3	4	5	6	7
b	1							
t	2							
i	3							
n	4							
e	5							
y	6							

Calculating Edit Distance

	$	b	r	i	t	n	e	y
$	0	$_{0}1_{2}$	2	3	4	5	6	7
b	1_{2}	0						
t	2							
i	3							
n	4							
e	5							
y	6							

Calculating Edit Distance

	$	**b**	**r**	**i**	**t**	**n**	**e**	**y**
$	0	1	2	3	4	5	6	7
b	1	0	1					
t	2							
i	3							
n	4							
e	5							
y	6							

Calculating Edit Distance

	$	b	r	i	t	n	e	y
$	0	1	2	3	4	5	6	7
b	1	0	1	2				
t	2							
i	3							
n	4							
e	5							
y	6							

Calculating Edit Distance

	$	b	r	i	t	n	e	y
$	0	1	2	3	4	5	6	7
b	1	0	1	2	3	4	5	6
t	2							
i	3							
n	4							
e	5							
y	6							

Calculating Edit Distance

	$	**b**	**r**	**i**	**t**	**n**	**e**	**y**
$	0	1	2	3	4	5	6	7
b	1	0	1	2	3	4	5	6
t	2	1	1	2	3	4	5	6
i	3							
n	4							
e	5							
y	6							

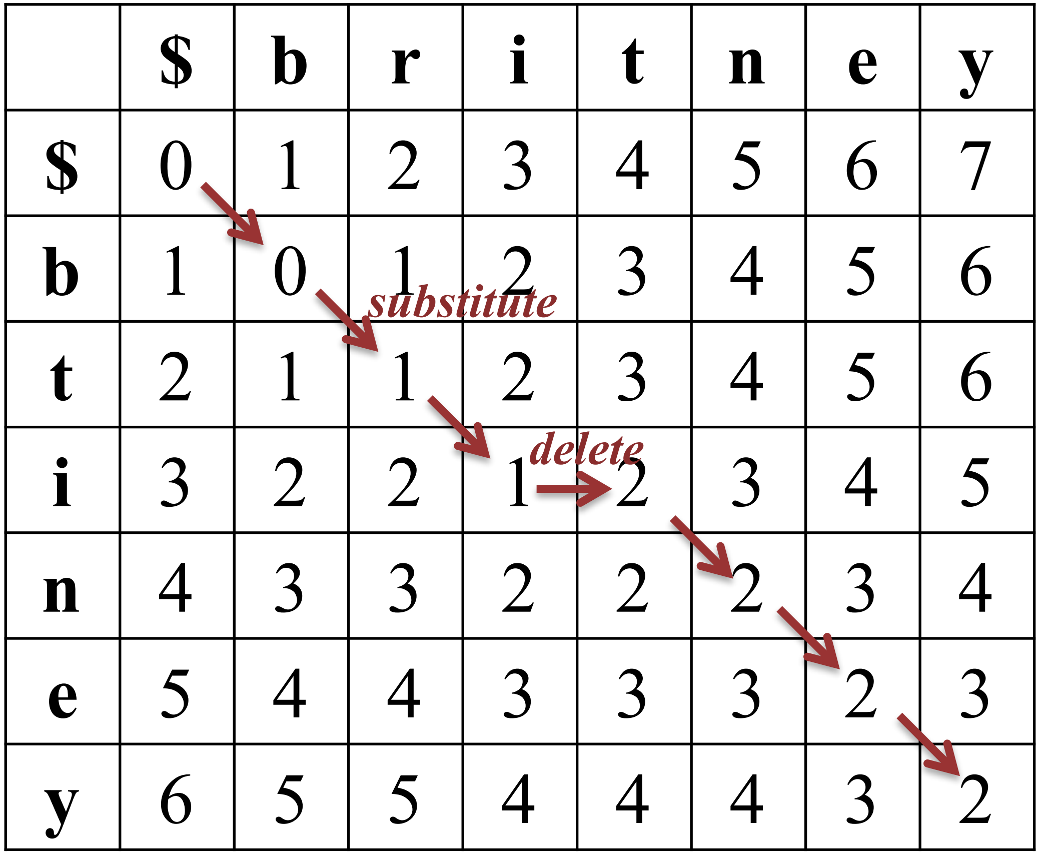
Calculating Edit Distance

	$	b	r	i	t	n	e	y
$	0	1	2	3	4	5	6	7
b	1	0	1	2	3	4	5	6
t	2	1	1	2	3	4	5	6
i	3	2	2	1	2	3	4	5
n	4	3	3	2	2	2	3	4
e	5	4	4	3	3	3	2	3
y	6	5	5	4	4	4	3	2

Calculating Edit Distance

	$	**b**	**r**	**i**	**t**	**n**	**e**	**y**
$	0	1	2	3	4	5	6	7
b	1	0	1	2	3	4	5	6
t	2	1	1	2	3	4	5	6
i	3	2	2	1	2	3	4	5
n	4	3	3	2	2	2	3	4
e	5	4	4	3	3	3	2	3
y	6	5	5	4	4	4	3	**2**

Calculating Edit Distance


substitute
delete

	$	b	r	i	t	n	e	y
$	0	1	2	3	4	5	6	7
b	1	0	1	2	3	4	5	6
t	2	1	1	2	3	4	5	6
i	3	2	2	1	2	3	4	5
n	4	3	3	2	2	2	3	4
e	5	4	4	3	3	3	2	3
y	6	5	5	4	4	4	3	2

From s=‘*btiney*’ to r=‘*britney*’

Match s[1]=‘b’ → *b*

Substitute s[2]=‘r’ with ‘t’→*bt*

Match s[3] = ‘i’ → *bti*

Delete s[4]=‘t’→ *bti*

Match s[5]=‘n’ → *btin*

Match s[6]=‘e’ → *btine*

Match s[7]=‘y’ → *btiney* = r

What’s the difference?

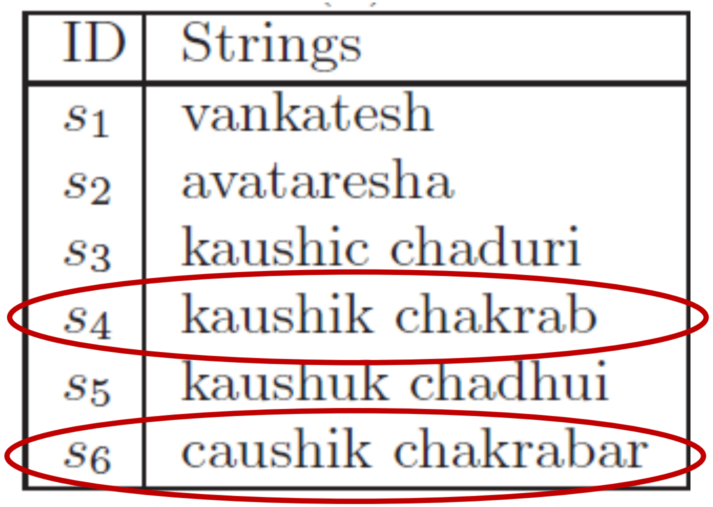
Functions (normalization)	**Edit Distance Edit Similarity**	**Overlap Size Jaccard Similarity**
Input	Sequences	Sets
Example Representations	DNA, String, Time Series	Image, Document, Vector, Friend List

String Similarity Join

- Input:
  - A collection of strings S
  - A threshold τ
- Output:
  - All string pairs $(s, r) \in S \times S\ such\ that\ ED(s, r) \leq \tau$

String Similarity Join

- Give threshold $\tau = 3$

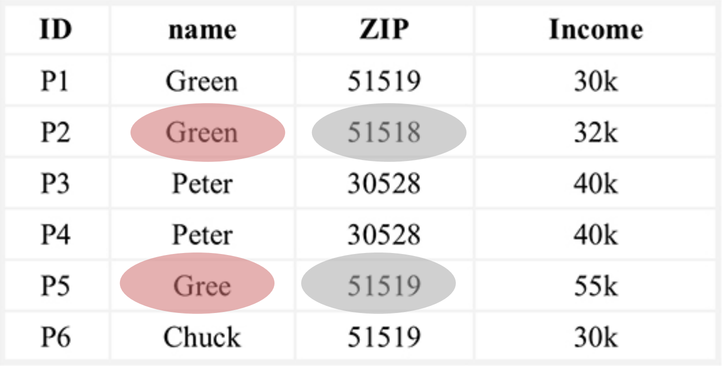


ID	Strings
s_1	vankatesh
s_2	avataresha
s_3	kaushic chaduri
s_4	kaushik chakrab
s_5	kaushuk chadhui
s_6	caushik chakrabar

ED(s_1 ,s_2)=5 ED(s_1 ,s_3)=13 ED(s_1 ,s_4)=12 ED(s_1 ,s_5)=12
ED(s_1 ,s_6)=14 ED(s_2 ,s_3)=12 ED(s_2 ,s_4)=12 ED(s_2 ,s_5)=12
ED(s_2 ,s_6)=14 ED(s_3 ,s_4)=5 ED(s_3 ,s_5)=4 ED(s_3 ,s_6)=8
ED(s_4 ,s_5)=4 **ED(s_4 ,s_6)=3** ED(s_5 ,s_6)=8

Data Cleaning & Integration

Relation with Duplicates



ID	name	ZIP	Income
P1	Green	51519	30k
P2	Green	51518	32k
P3	Peter	30528	40k
P4	Peter	30528	40k
P5	Gree	51519	55k
P6	Chuck	51519	30k

Challenges

$O(n^2)$ *pairs of strings.*

1 million strings result in 1 trillion pairs !

O($|r||s|$) *time to calculate ED(r, s)*

Filter-and-Refine Framework

- Basic idea
  - **Filter** a large number of dissimilar string pairs
  - **Verify** the remaining potentially similar pairs
- Good Filter Condition
  - Efficient to check
  - Effective for pruning dissimilar pairs

Length Filter

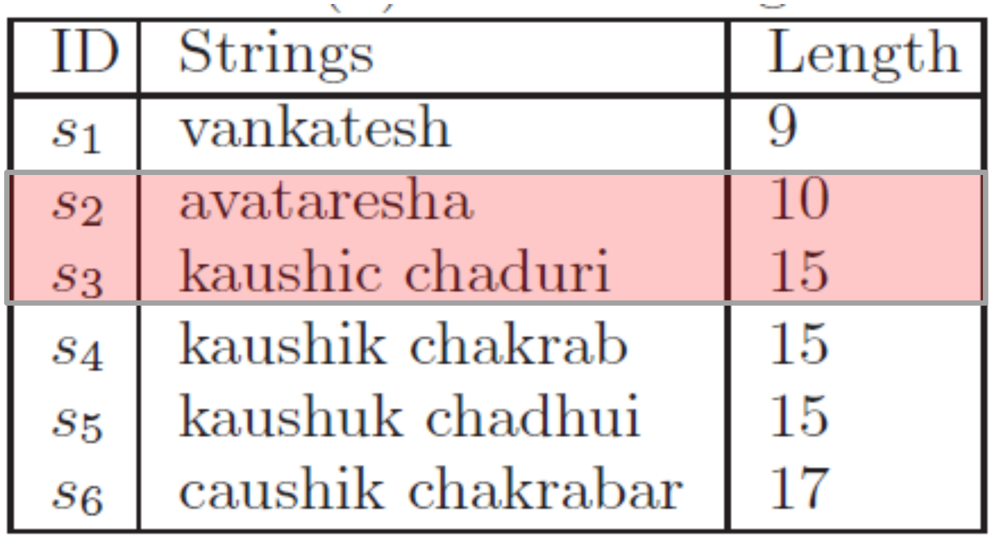
- ED(*r*, *s*): The minimum number of edit operations (insertion/deletion/substitution) needed to transform *r* to *s*.

- *Property:* $\mathrm{ED}(\boldsymbol{r}, \boldsymbol{s}) \geq \big||\boldsymbol{r}|-|\boldsymbol{s}|\big|$
  - it needs at least $||r|-|s||$ deletions to just make *r* and *s* have the same length

Applying the Length Filter

- Give threshold τ = 3
- Pruning Condition:

$$\left||s_i| - |s_j|\right| > 3$$

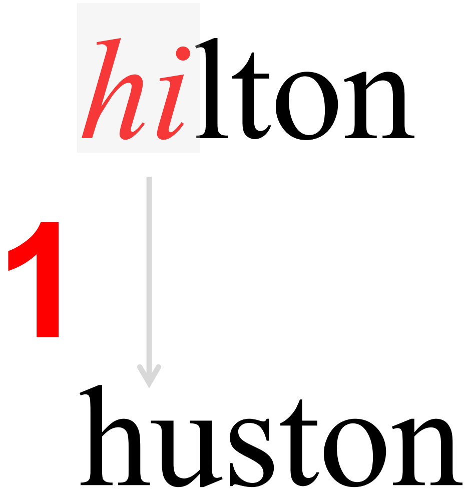


ID	Strings	Length
s_1	vankatesh	9
s_2	avataresha	10
s_3	kaushic chaduri	15
s_4	kaushik chakrab	15
s_5	kaushuk chadhui	15
s_6	caushik chakrabar	17

ED(s_1 ,s_2)=5 ~~ED(s_1 ,s_3)=13~~ ~~ED(s_1 ,s_4)=12~~ ~~ED(s_1 ,s_5)=12~~
~~ED(s_1 ,s_6)=14~~ ~~ED(s_2 ,s_3)=12~~ ~~ED(s_2 ,s_4)=12~~ ~~ED(s_2 ,s_5)=12~~
~~ED(s_2 ,s_6)=14~~ ED(s_3 ,s_4)=5 ED(s_3 ,s_5)=4 ED(s_3 ,s_6)=8
ED(s_4 ,s_5)=4 **ED(s_4 ,s_6)=3** ED(s_5 ,s_6)=8

Partitioning Filter

- Give threshold $\tau = 1$



*hi*lton

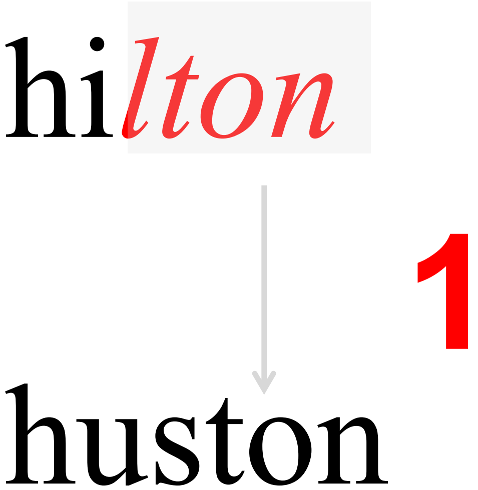
1

huston

hi dose not appear in huston and needs at least 1 edit

Partitioning Filter

- Give threshold $\tau = 1$



hi*lton*

1

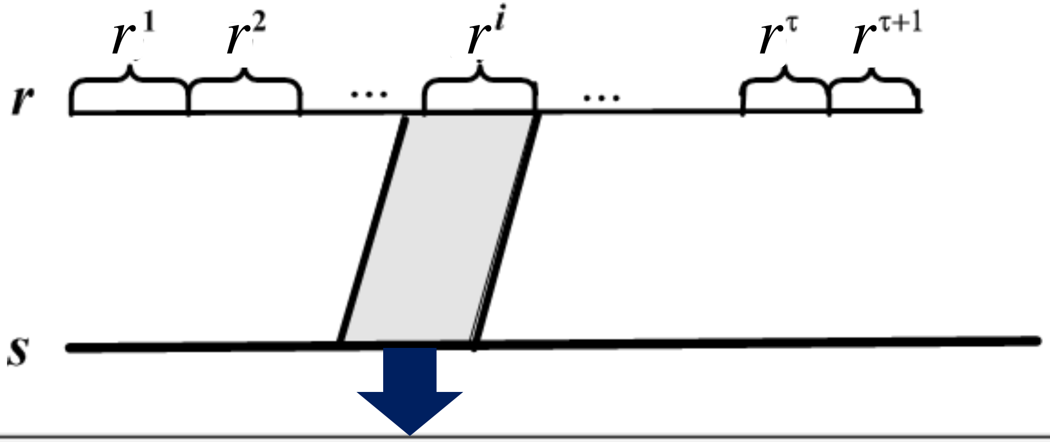
huston

minimum # of edit operations is 2. Prune!

Partitioning Filter

- Threshold ***τ***
- String r
- String s

split r to τ +1 disjoint segments



Is there any substring of *s* matching a segment of *r* ?

Yes

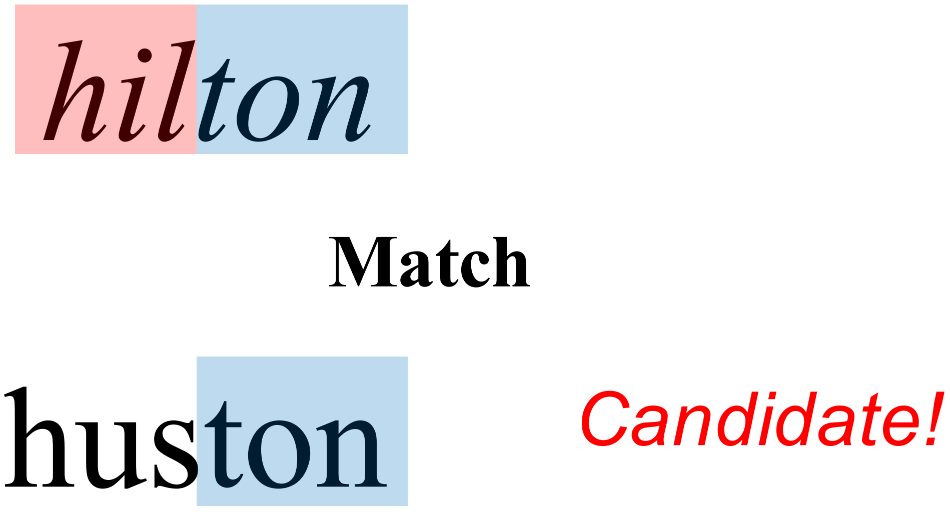
<*r*, *s*> are a candidate pair

No

<r, s> are dissimilar, prune

How to Partition?

- Give threshold $\tau = 1$



hilton

Match

huston

Candidate!

Partition Scheme

- Even Partition Scheme
  - Given $\tau = 3$, “*avataresha*” → {“*av*”, “*at*”, “*are*”, “*sha*”}
- Other Schemes
  - Select good partition strategies.
  - Adaptive partition scheme [Deng et al. 2012a].

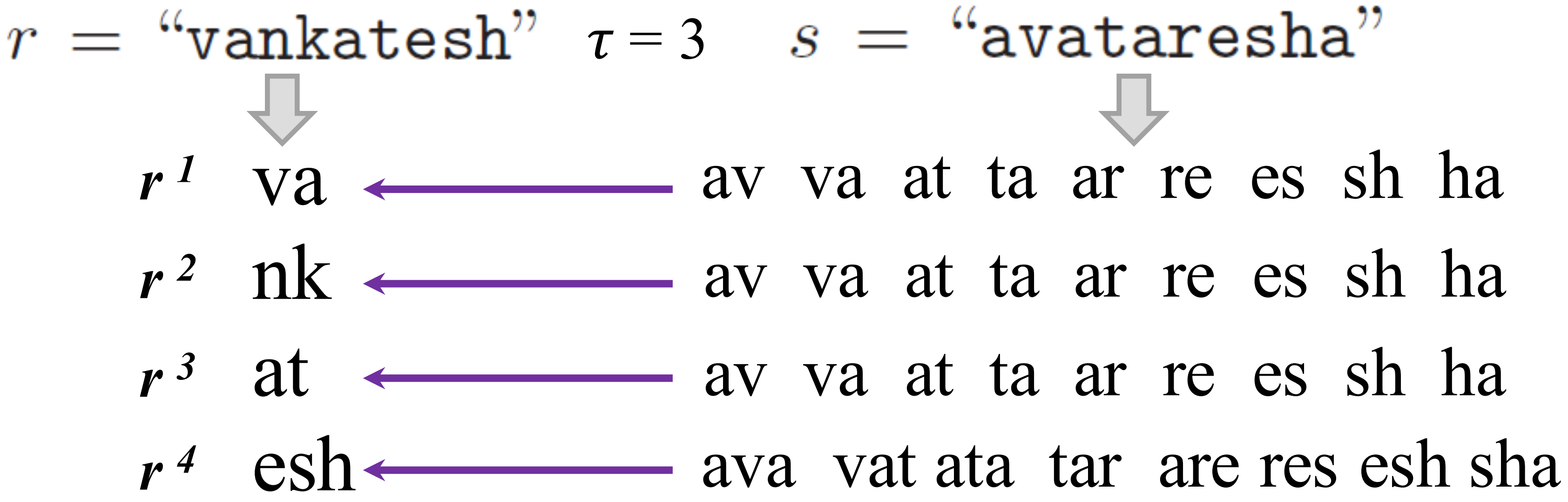
Challenge

there are $|s|^2$ *substrings in s*

how to reduce the number of substrings to compare with for each segment?

Length-based Method

- For each segment, only compare to the substrings with the same length as the segment.



$r =$ "vankatesh" $\tau = 3$ $s =$ "avataresha"

r^1 va ← av va at ta ar re es sh ha

r^2 nk ← av va at ta ar re es sh ha

r^3 at ← av va at ta ar re es sh ha

r^4 esh ← ava vat ata tar are res esh sha

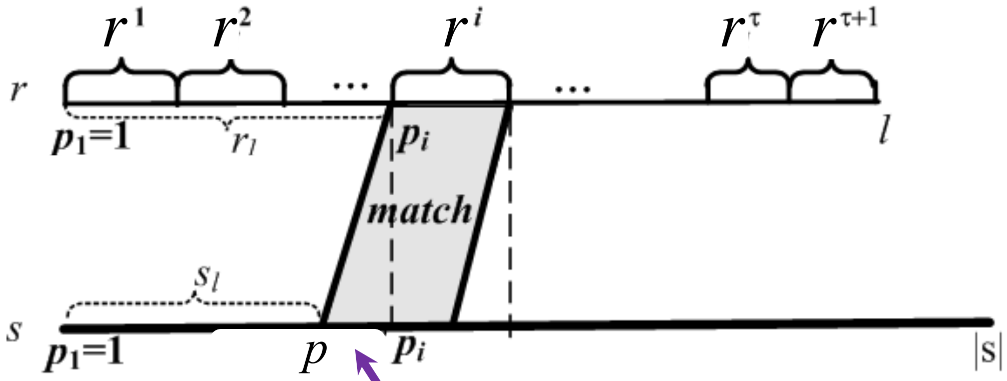
Length-based Method

- For any strings *r*, *s*, and *τ*, the number of comparisons:

$$(\tau+1)(|s|+1)-|r|$$

- For $r = $ “`vankatesh`” and $s = $ “`avataresha`” , the number is 35

Shift-based Method



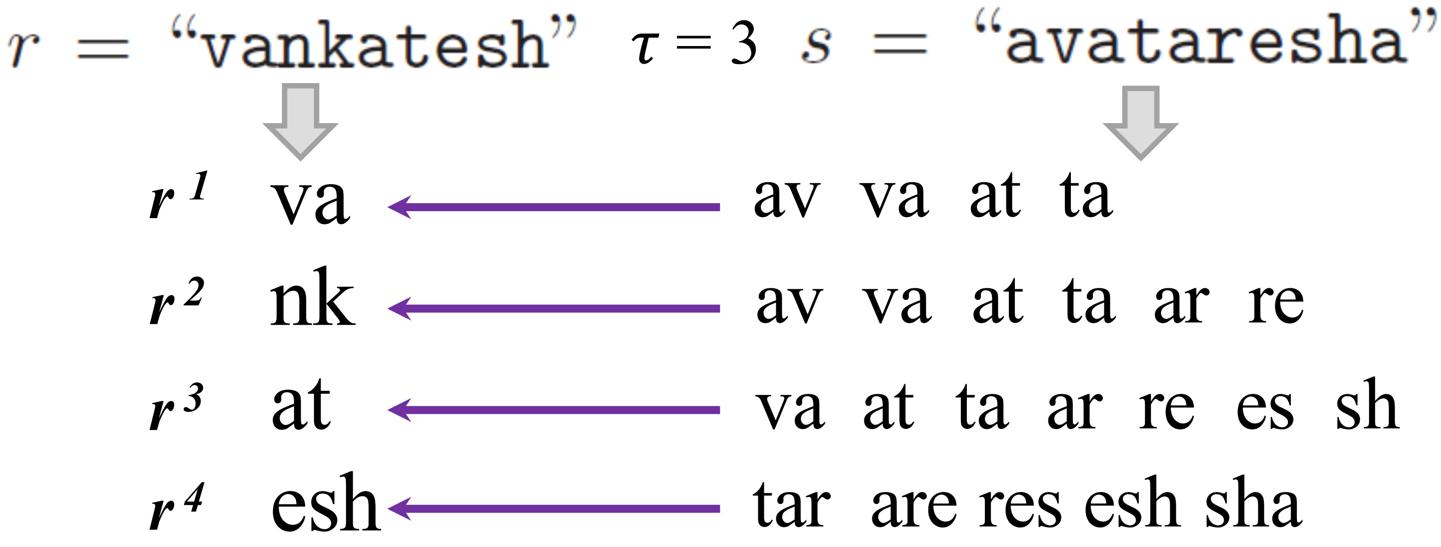
Just transform $\mathbf{r}_l$ *to* $\mathbf{s}_l$ *needs at least* $\left||s_l|-|r_l|\right|$ *edit operations*

Pruning Condition: $\left||s_l|-|r_l|\right| > \tau \longrightarrow |p - p_i| > \tau$

- For each segment r^i with the start position p_i , only compares to the substrings with start positions p in $[p_i - \tau, p_i + \tau]$

Shift-based Method

- For each segment r^i with the start position p_i, only compares to the substrings with start positions p in $[p_i - \tau, p_i+\tau]$



r = "vankatesh" $\tau = 3$ s = "avataresha"

r^1 va ← av va at ta

r^2 nk ← av va at ta ar re

r^3 at ← va at ta ar re es sh

r^4 esh ← tar are res esh sha

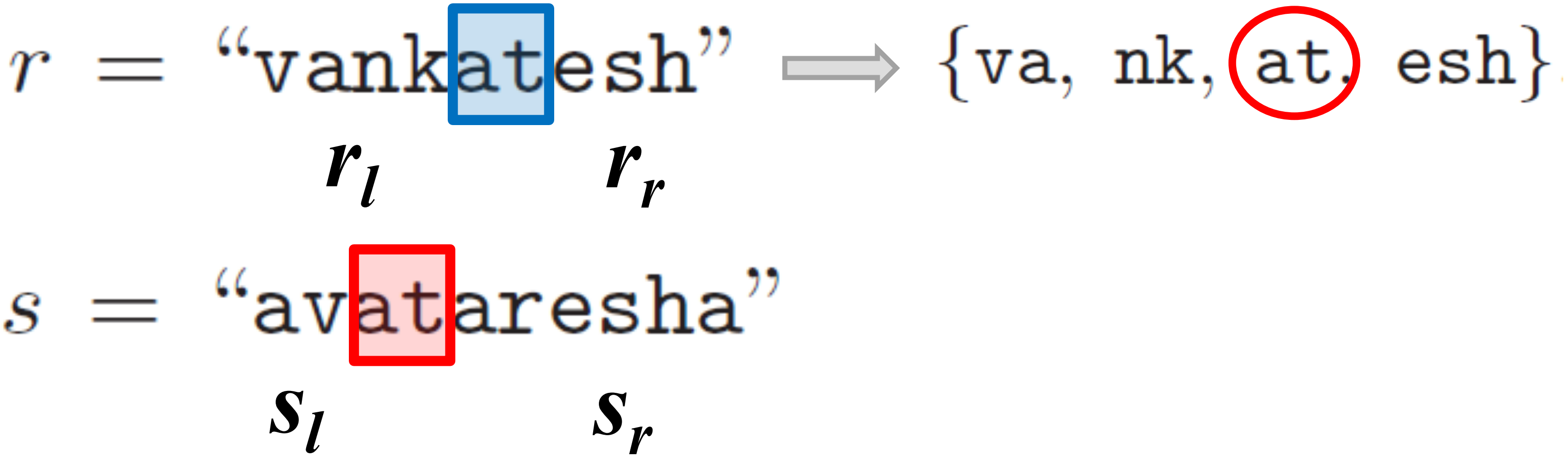
Shift-based Method

- For any strings r, s, and τ, the number of comparisons:

$$(\tau + 1)(2\tau + 1)$$

- For $r =$ “$\texttt{vankatesh}$” and $s =$ “$\texttt{avataresha}$” , the number is 22.

Position-aware Method



$r =$ "vankatesh" $\Longrightarrow$ {va, nk, at, esh}

r_l r_r

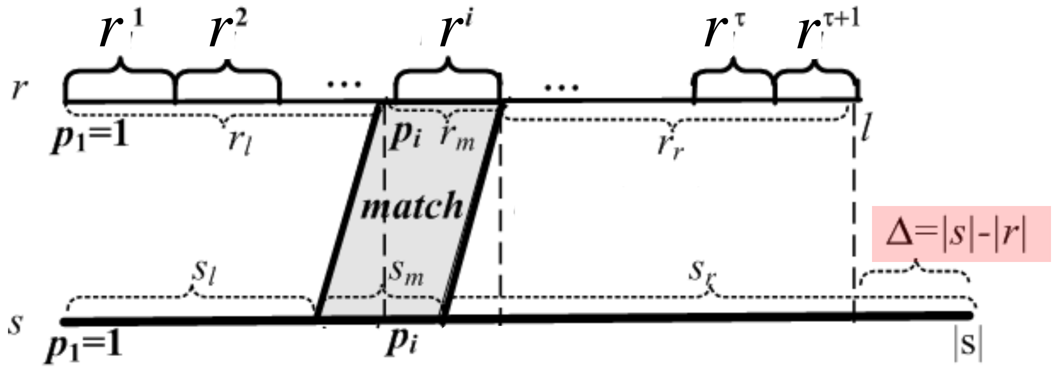
$s =$ "avataresha"

s_l s_r

Transform $\boldsymbol{r_l}$ *to* $\boldsymbol{s_l}$*, match "at", and transform* $\boldsymbol{r_r}$ *to* $\boldsymbol{s_r}$

$$\big||s_l|-|r_l|\big| + \big||s_r|-|r_r|\big| = 2 + 3 > \tau = 3$$

Position-aware Method



Transform $\boldsymbol{r}_l$ *to* $\boldsymbol{s}_l$, *match* $\boldsymbol{r}_m$ *and* $\boldsymbol{s}_m$ *then transform* $\boldsymbol{r}_r$ *to* $\boldsymbol{s}_r$

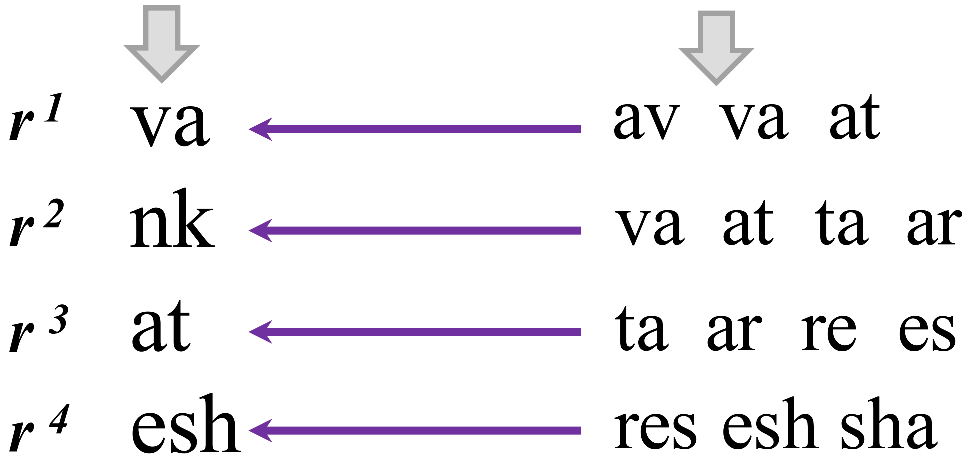
Pruning Condition: $||s_l|-|r_l|| + ||s_r|-|r_r|| > \tau$

- For each segment r^i with the start position p_i, only compare to the substrings with start position in $[P_i - \left\lfloor\frac{\tau-\Delta}{2}\right\rfloor, P_i + \left\lfloor\frac{\tau+\Delta}{2}\right\rfloor]$ where $\Delta = |s| - |r|$

Position-aware Method

- For each segment r^i with the start position $\boldsymbol{p_i}$, only compare to the substrings with start position in $[\boldsymbol{P}_i - \lfloor\frac{\tau-\Delta}{2}\rfloor, \boldsymbol{P}_i + \lfloor\frac{\tau+\Delta}{2}\rfloor]$ where $\Delta = |s| - |r|$

$\tau = 3$ r = “vankatesh” s = “avataresha”



r^1 va ← av va at

r^2 nk ← va at ta ar

r^3 at ← ta ar re es

r^4 esh ← res esh sha

Position-aware Method

- For any strings r, s, and τ, the number of comparisons:

$$(\tau + 1)^2$$

- For $r =$ “vankatesh” and $s =$ “avataresha” , the number is 14.

Multi-match-aware Method

-- Left-side Perspective

$r =$ "$\boxed{\text{va}}$nkatesh" $\Longrightarrow$ $\{\boxed{\text{va,}}\ \boxed{\text{nk}},\ \boxed{\text{at}},\ \boxed{\text{esh}}\}$

$\tau = 3$

r_l=""

3 unvisited segments

$s =$ "a$\boxed{\text{va}}$taresha"

s_l="*a*"

$1 \quad 1 \quad 1$

$+$

$||s_l|-|r_l||=1$

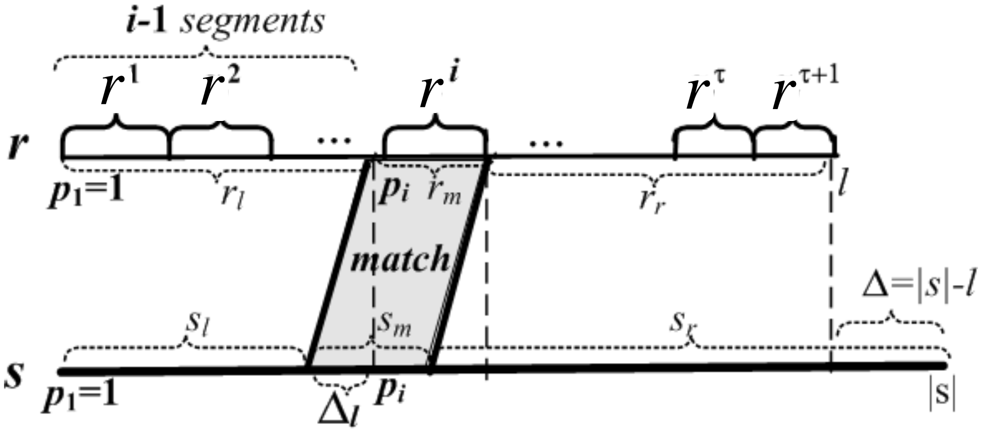
$=4 > \tau$

If none of the 3 unvisited segments having matching substring in s, then r and s cannot be similar

Thus we can safely skip the current matching segment and look for the next

Multi-match-aware Method

-- Left-side Perspective



$\Delta_l \le i-1$ *as there are i-1 segments in r_l*

Pruning Condition: $\left||s_l|-|r_l|\right|$ + (# of unvisited segments) > $\boldsymbol{\tau}$

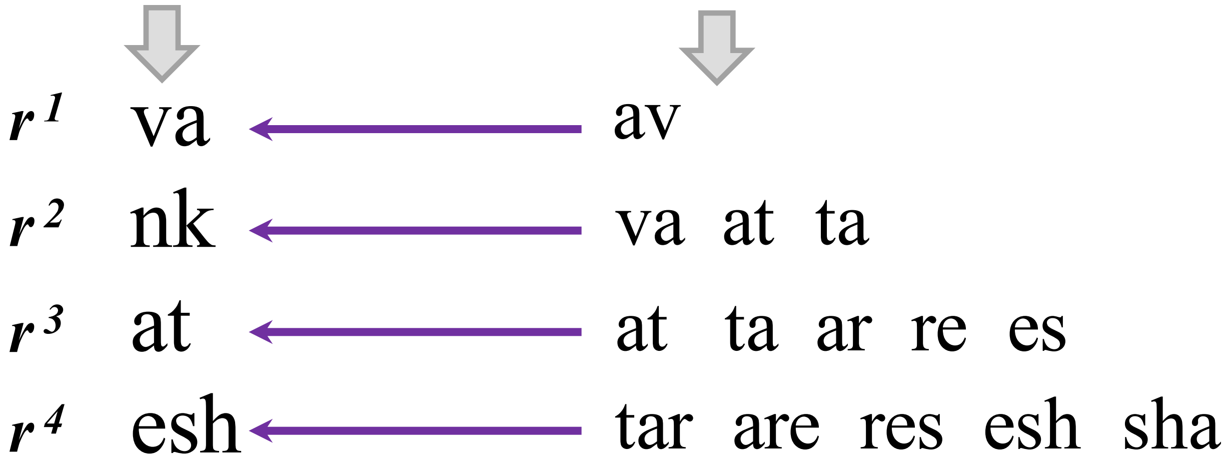
- For each segment r^i with the start position $\boldsymbol{p_i}$, only compare to the substrings with start position in $[\boldsymbol{P_i}-(\boldsymbol{i}-1), \boldsymbol{P_i}+(\boldsymbol{i}-1)]$

Multi-match-aware Method

-- *Left-side Perspective*

- For each segment r^i with the start position p_i , only compare to the substrings with start position in $[P_i - (i-1), P_i + (i-1)]$

$\tau = 3$ $r =$ "vankatesh" $s =$ "avataresha"



r^1 va ← av

r^2 nk ← va at ta

r^3 at ← at ta ar re es

r^4 esh ← tar are res esh sha

Multi-match-aware Method

-- Left-side Perspective

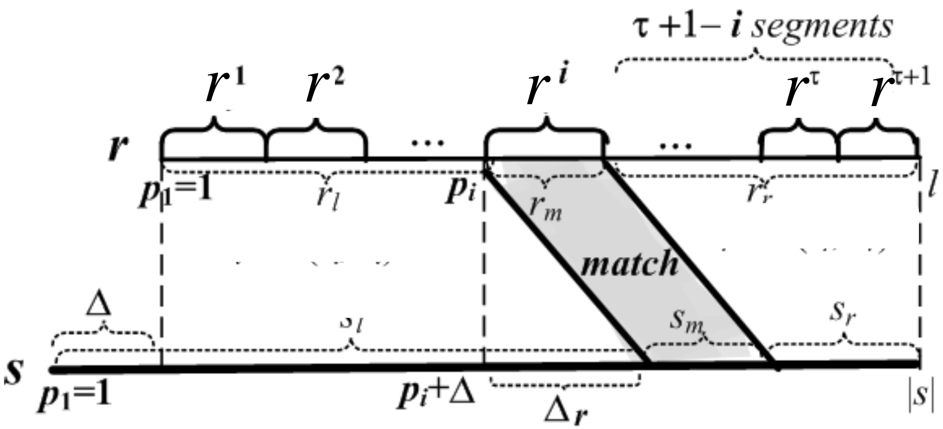
- For any strings r, s, and τ, the number of comparisons:

$$\tau^2 + 2\tau$$

- For $r = \text{“}\texttt{vankatesh}\text{”}$ and $s = \text{“}\texttt{avataresha}\text{”}$, the number is 14.

Multi-match-aware Method

-- *Right-side Perspective*



$\Delta_r \le \tau+1-i$ *as there are* $\tau+1-i$ *segments in* r_r

Pruning Condition: (# unvisited segments) + $\big|\,|s_r|-|r_r|\,\big| > \tau$

- For each segment r^i with the start position p_i, only compare to the substrings with start position in $[P_i+\Delta-(\tau+1-i), P_i+\Delta+(\tau+1-i)]$

Multi-match-aware Method

- Interestingly, we can apply the multi-match-aware method from left- and right-side perspectives simultaneously.

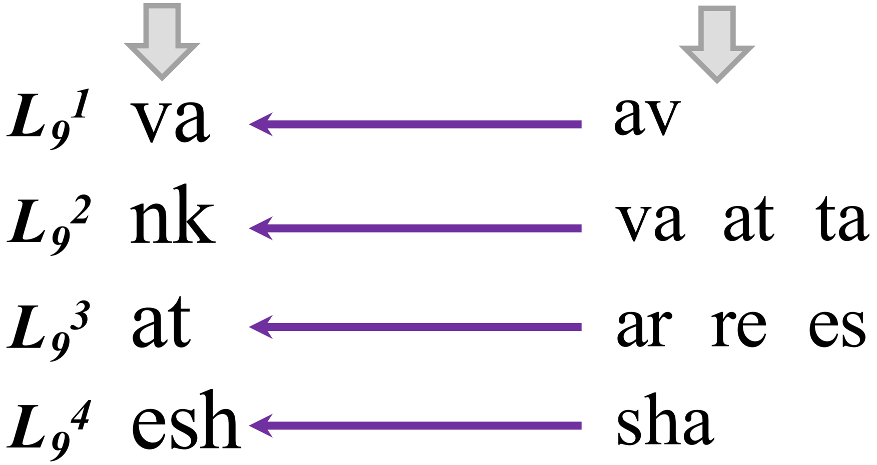
- For each segment $\boldsymbol{r}^{i}$ with start position $\boldsymbol{p}_{i}$, only compare to the substrings with start position in

$$[\boldsymbol{max}(\boldsymbol{P}_i - (\boldsymbol{i}-\mathbf{1}), \boldsymbol{P}_i + \boldsymbol{\Delta} - (\boldsymbol{\tau}+\mathbf{1}-\boldsymbol{i})),\quad \boldsymbol{min}(\boldsymbol{P}_i + (\boldsymbol{i}-\mathbf{1}), \boldsymbol{P}_i + \boldsymbol{\Delta} + (\boldsymbol{\tau}+\mathbf{1}-\boldsymbol{i}))]$$

Multi-match-aware Method

- For each segment $\boldsymbol{r^i}$ with start position $\boldsymbol{p_i}$, only compare to the substrings with start position in $$[\boldsymbol{max}(\boldsymbol{P_i} - (\boldsymbol{i}-1), \boldsymbol{P_i} + \Delta - (\tau + 1 - \boldsymbol{i})), \boldsymbol{min}(\boldsymbol{P_i} + (\boldsymbol{i}-1), \boldsymbol{P_i} + \Delta + (\tau + 1 - \boldsymbol{i}))]$$

$\tau = 3$ $r =$ “vankatesh” $s =$ “avataresha”



$\boldsymbol{L_9^1}$ va ← av

$\boldsymbol{L_9^2}$ nk ← va at ta

$\boldsymbol{L_9^3}$ at ← ar re es

$\boldsymbol{L_9^4}$ esh ← sha

Multi-match-aware Method

- For any strings r, s, and τ, the number of comparisons:

$$\left\lfloor \frac{\tau^2 - \Delta^2}{2} \right\rfloor + \tau + 1$$

- For $r = $ “vankatesh” and $s = $ “avataresha” , the number is 8.

Theoretical Results

- The number of comparisons by the multi-match-aware method is **minimum** while guarantees completeness

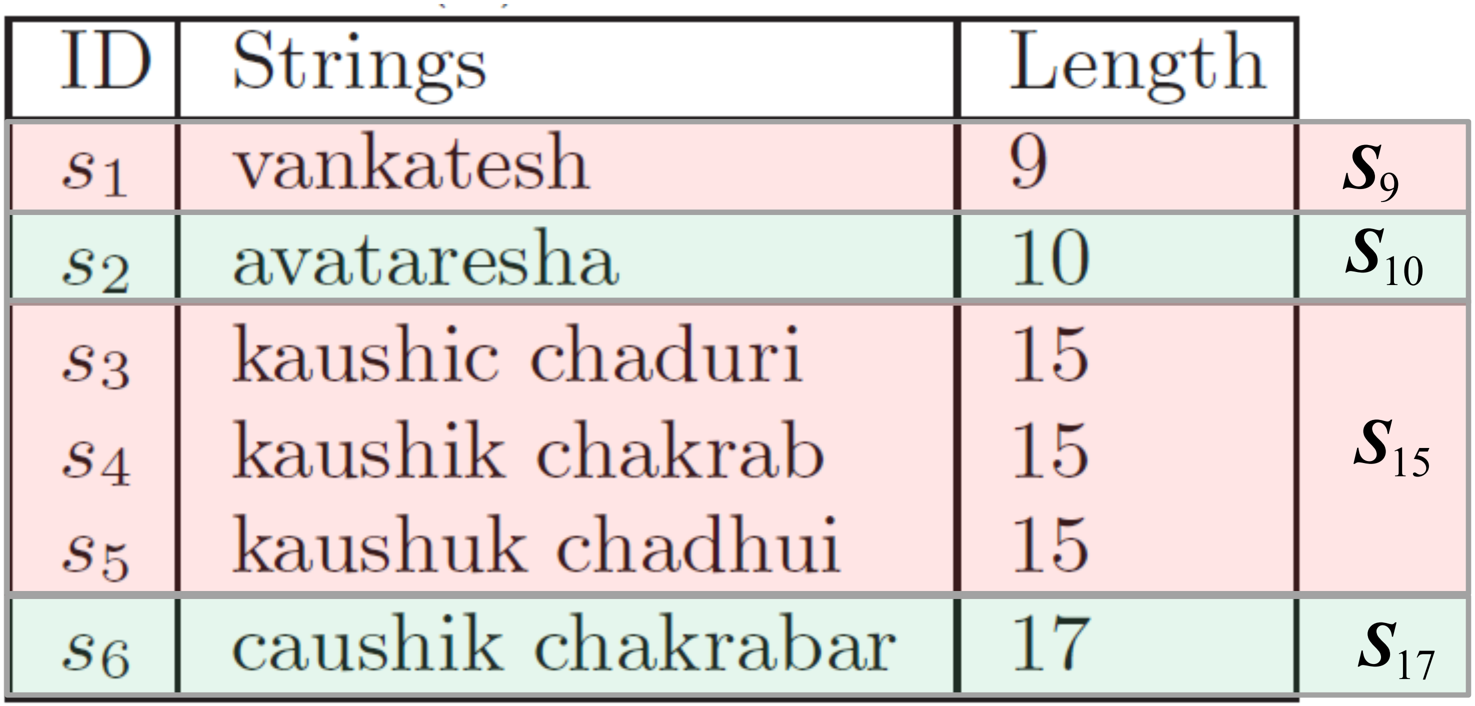
- For any s, r and τ,

$$W_{multi\text{-}match}(s, r, \tau) \subseteq W_{position}(s, r, \tau) \subseteq W_{shift}(s, r, \tau) \subseteq W_{length}(s, r, \tau)$$

Applying Partitioning Filter

Applying Partitioning Filter

- 1. Group all the strings by length: **S**

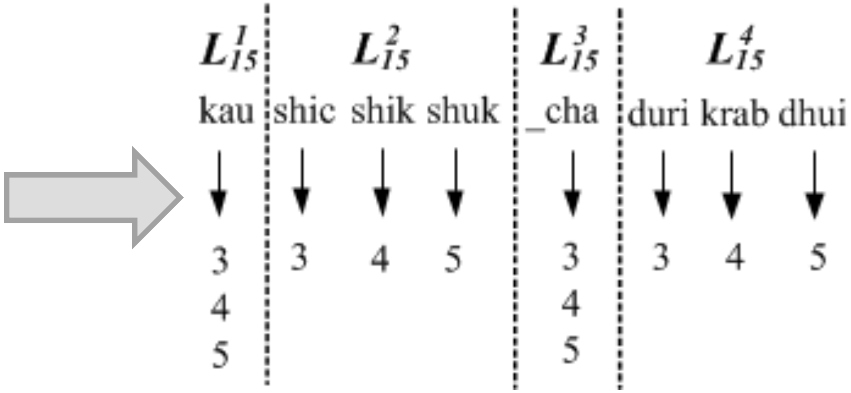


ID	Strings	Length	
s_1	vankatesh	9	$\boldsymbol{S}_9$
s_2	avataresha	10	$\boldsymbol{S}_{10}$
s_3	kaushic chaduri	15	$\boldsymbol{S}_{15}$
s_4	kaushik chakrab	15	
s_5	kaushuk chadhui	15	
s_6	caushik chakrabar	17	$\boldsymbol{S}_{17}$

Applying Partitioning Filter

- 2. For each group S_l, partition its strings into $\tau + 1$ segments and build $\tau + 1$ inverted indexes L_l^i

S_{15}	1	2	3	4
s_3=	kau	shic	_cha	duri
s_4=	kau	shik	_cha	krab
s_5=	kau	shuk	_cha	dhui



L_{15}^1: kau → 3, 4, 5

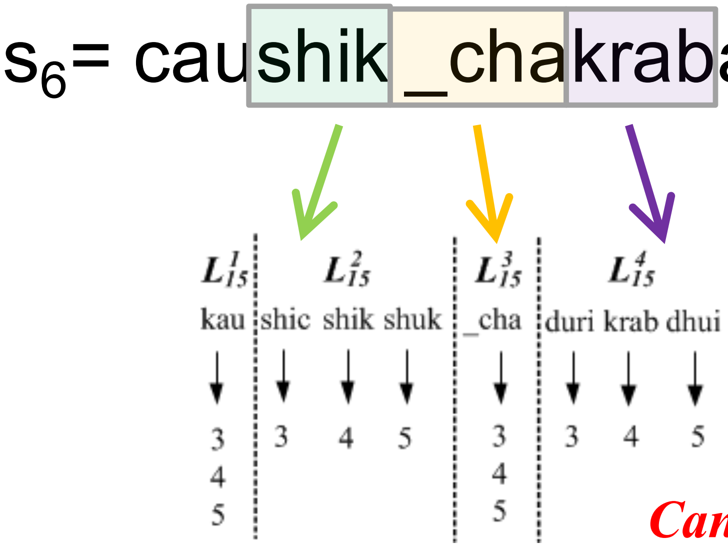
L_{15}^2: shic → 3; shik → 4; shuk → 5

L_{15}^3: _cha → 3, 4, 5

L_{15}^4: duri → 3; krab → 4; dhui → 5

Applying Partitioning Filter

- 3. For each string ***s*** and index $\boldsymbol{L_l^i}$, select substrings from ***s*** based on the partitioning filter to get candidates:



s_6= caushik_chakrabar

$$[\boldsymbol{max}(\boldsymbol{P_i} - (\boldsymbol{i}-\mathbf{1}), \boldsymbol{P_i} + \boldsymbol{\Delta} - (\boldsymbol{\tau} + \mathbf{1} - \boldsymbol{i})),$$
$$\boldsymbol{min}(\boldsymbol{P_i} + (\boldsymbol{i}-\mathbf{1}), \boldsymbol{P_i} + \boldsymbol{\Delta} + (\boldsymbol{\tau} + \mathbf{1} - \boldsymbol{i}))]$$

the segments in $\boldsymbol{L_l^i}$ have the same start positions p_i and same string lengths l

***Candidates*: <s_3, s_6> <s_4, s_6> <s_5, s_6>**

Applying Partitioning Filter

- 4. Verify the candidates

***Candidates*:** $<s_3, s_6>$ $<s_4, s_6>$ $<s_5, s_6>$

$\mathrm{ED}(s_3, s_6) > 3$ ✗

$\mathrm{ED}(s_4, s_6) = 3$ ✓

$\mathrm{ED}(s_5, s_6) > 3$ ✗

Improving Verification

Calculating Edit Distance

	$	**b**	**r**	**i**	**t**	**n**	**e**	**y**
$	0	1	2	3	4	5	6	7
b	1	0	1	2	3	4	5	6
t	2	1	1	2	3	4	5	6
i	3	2	2	1	2	3	4	5
n	4	3	3	2	2	2	3	4
e	5	4	4	3	3	3	2	3
y	6	5	5	4	4	4	3	**2**

Verification

	$	b	r	i	t	n	e	y
$	0	1	2	3	4	5	6	7
b	1	0	1	2	3	4	5	6
t	2	1	1	2	3	4	5	6
i	3	2	2	1	2	3	4	5
n	4	3	3	2	2	2	3	4
e	5	4	4	3	3	3	2	3
y	6	5	5	4	4	4	3	2

$ED(r[1,i],s[1,j]) \geq |i - j|$

Verification

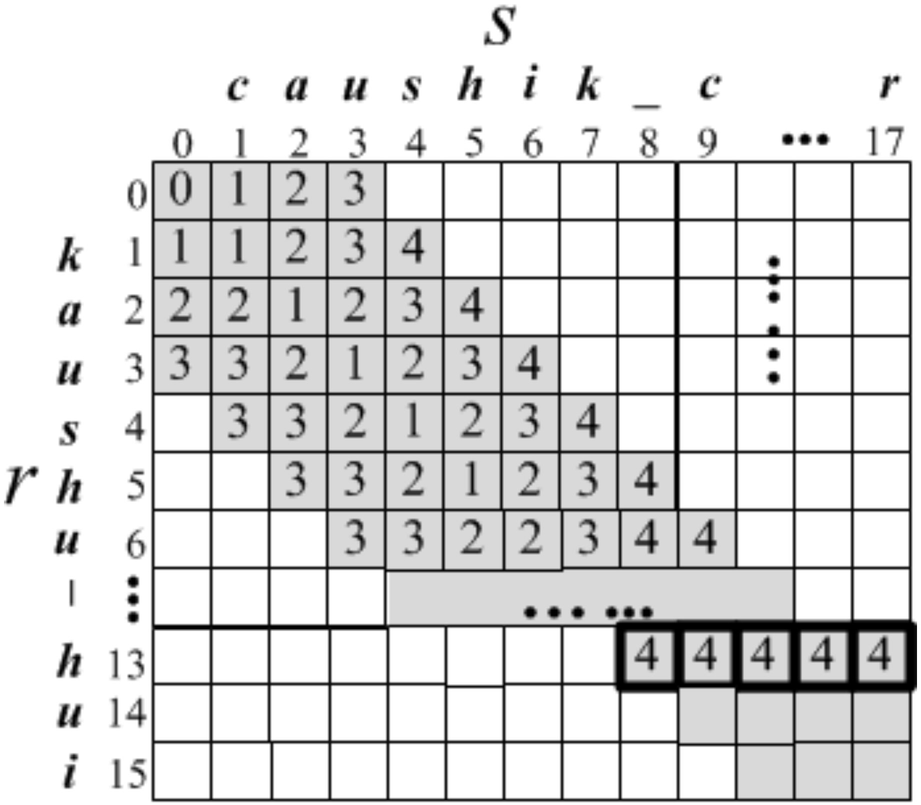
	$	b	r	i	t	n	e	y
$	0	1	2					
b	1	0	1	2				
t	2	1	1	2	3			
i		2	2	1	2	3		
n			3	2	2	2	3	
e				3	3	3	2	3
y					4	4	3	2

Verification

only need to calculate a band of width $2\tau+1$

	$	**b**	**r**	**i**	**t**	**n**	**e**	**y**
$	0	1	2					
b	1	0	1	2				
t	2	1	1	2	3			
i		2	2	1	2	3		
n			3	2	2	2	3	
e				3	3	3	2	3
y					4	4	3	2

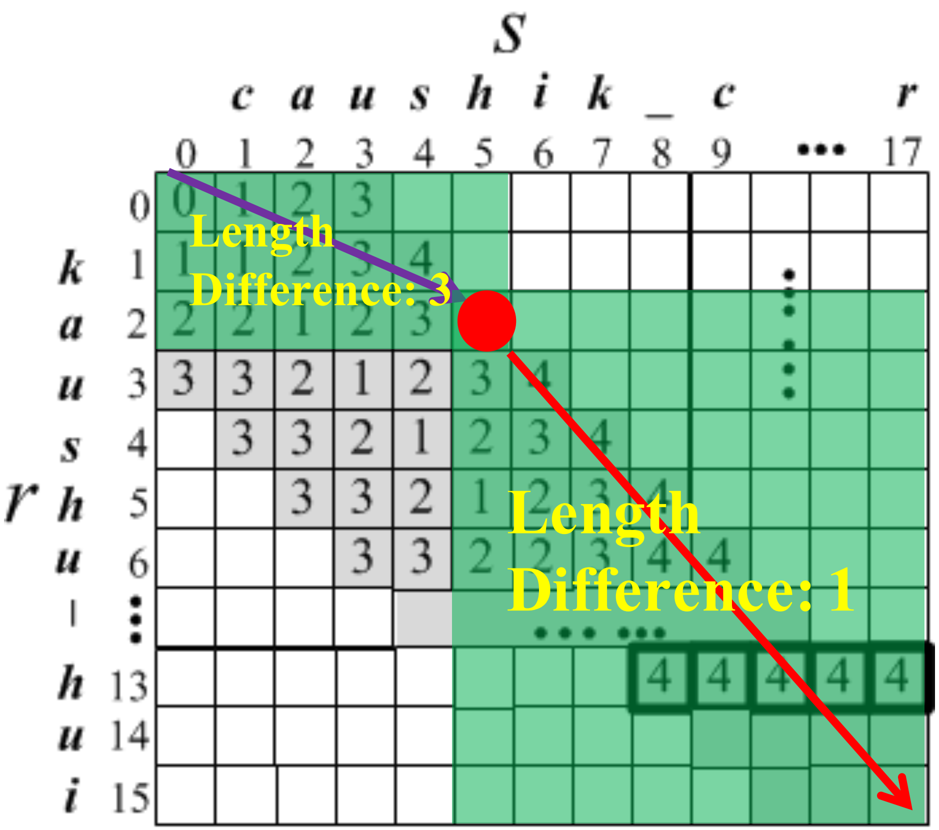
Early Termination



		c	a	u	s	h	i	k	_	c			r
		1	2	3	4	5	6	7	8	9		•••	17
	0	0	1	2	3								
k	1	1	1	2	3	4							
a	2	2	2	1	2	3	4				⋮		
u	3	3	3	2	1	2	3	4			⋮		
s	4		3	3	2	1	2	3	4				
h	5			3	3	2	1	2	3	4			
u	6				3	3	2	2	3	4	4		
⋮	⋮					•••	•••						
h	13								4	4	4	4	4
u	14												
i	15												

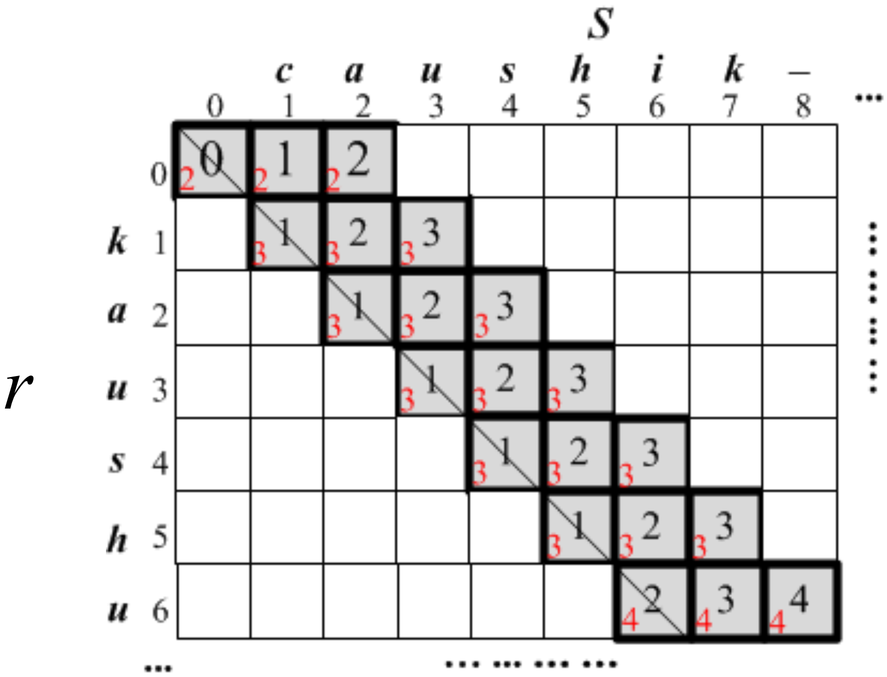
all derived cells must have values larger than τ.

Length-aware Verification



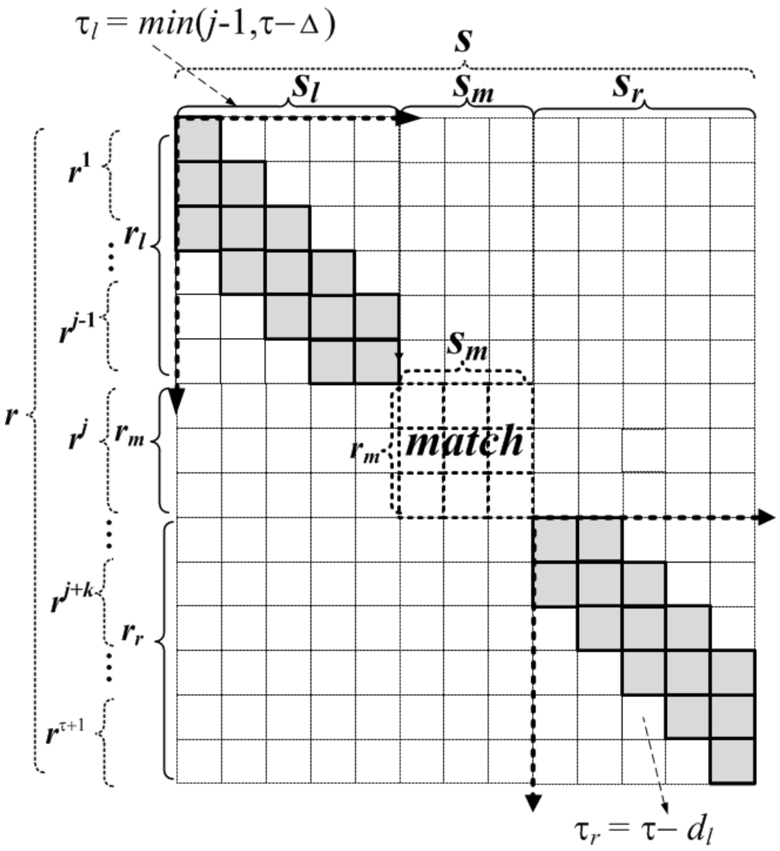
total length difference is 4 > τ. Thus no need to calculate M[2][5].

Length-aware Verification



only need to calculate a band of width τ

Extension-base Verification



Extension-base Verification

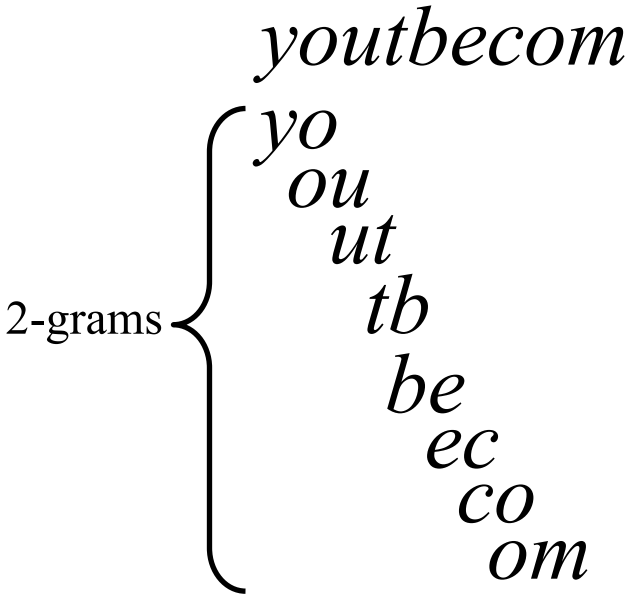
- We can verify a candidate pair using tighter thresholds:
  - For the left parts we can set $\tau_l = i - 1$.
  - For the right parts we can set $\tau_r = \tau + 1 - i$.

band widths are $\tau_l < \tau$ *and* $\tau_r < \tau$

Alternative: Prefix Filter

q-gram

- *q*-gram is the substring of length *q*



youtbecom

2-grams

yo
ou
ut
tb
be
ec
co
om

q-gram

- 1 edit operation destroies at most *q* grams.

*yout**d**ecom*

yo

ou

ut

*t**d***

***d**e*

ec

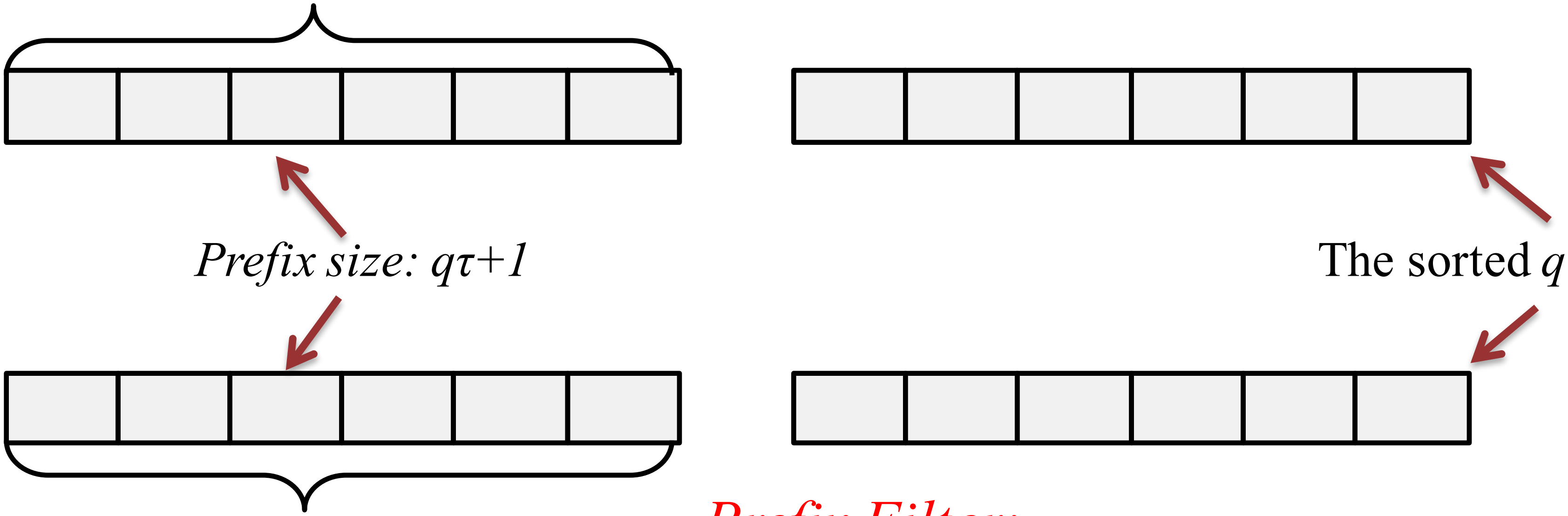
co

om

- τ edit operations destroy *at most* $q\tau$ q-grams.
- if r and s have more than $q\tau$ mismatch q-grams, ED(r, s)>τ

Prefix Filter

Sort all q-grams by global ordering, such as *alphabetic order*



Prefix(r)

Prefix size: $q\tau+1$

The sorted q-grams

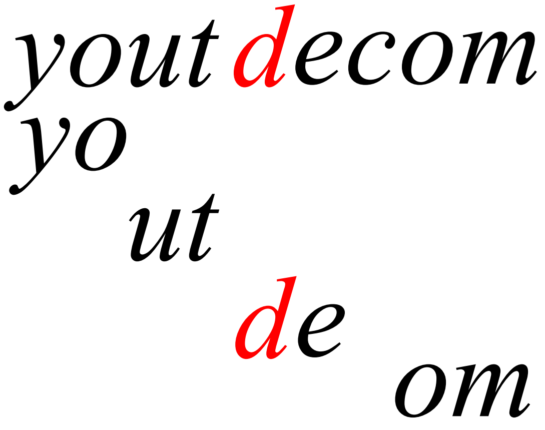
Prefix(s)

Prefix Filter:

If $prefix(r) \cap prefix(s) = \phi$, $ED(r,s) > \tau$

Disjoint q-gram

- One edit operation destroies at most *1* disjoint gram.


youtdecom

yo

ut

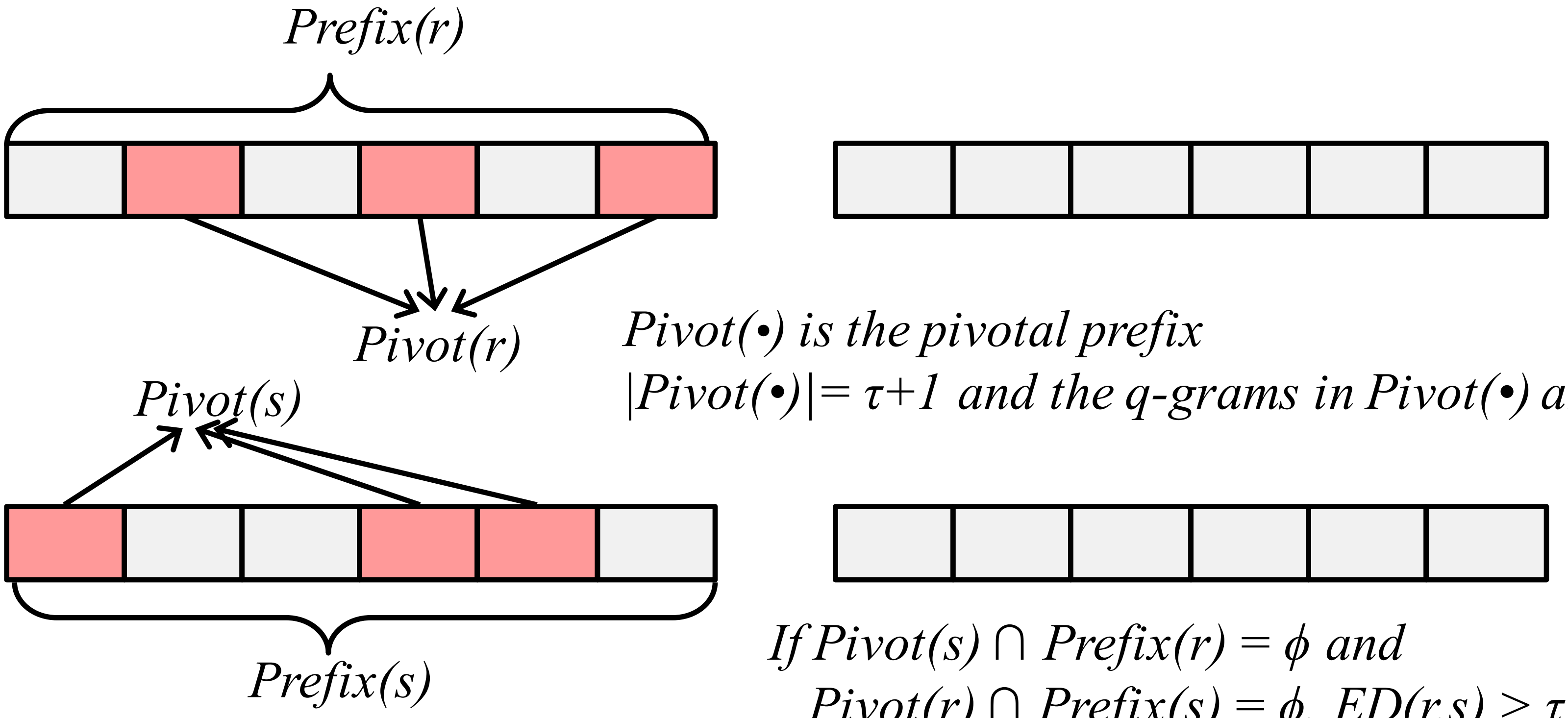
de

om

- τ edit operations destroy *at most* τ disjoint q-grams.
- if r and s have more than τ mismatch disjoint q-grams, ED(r, s)> τ

Pivotal Prefix Filter

Sort all q-grams by global ordering, such as *alphabetical order*



Pivot(•) is the pivotal prefix

$|Pivot(\bullet)| = \tau+1$ *and the q-grams in Pivot(•) are disjoint*

If $Pivot(s) \cap Prefix(r) = \phi$ *and*

$Pivot(r) \cap Prefix(s) = \phi,\ ED(r,s) > \tau$

Takeaways

(1) The partitioning filter for Edit Distance

(2) The multi-match-aware method

(3) The extension-based verification

(4) The prefix filters for Edit Distance

References

- Pass-Join: A Partition based Method for Similarity Joins. G. Li, D. Deng, J. Wang, J. Feng. VLDB 2012.
- A Pivotal Prefix Based Filtering Algorithm for String Similarity Search. D. Deng, G. Li, J. Feng. SIGMOD 2014.
- Ed-Join: An Efficient Algorithm for Similarity Joins with Edit Distance Constraints. C. Xiao, W. Wang, X. Lin. VLDB 2008