

Thm 6 : Let a & b be integers, not both zero. For a positive integer d , $d = \gcd(a, b)$ if and only if

(a) $d|a$ and $d|b$.

(b) whenever, $c|a$ and $c|b$, then $c|d$.

Proof. Try to prove!

The Euclidean Algorithm

The Euclidean Algorithm is described as follows:
Let a and b be two integers whose greatest common divisor is desired.

Since $\gcd(|a|, |b|) = \gcd(a, b)$, so there is no harm in assuming that $a \geq b > 0$.

The first step is to apply the Division Algorithm to a and b to get

$$a = q_1 b + r_1, \quad 0 \leq r_1 < b.$$

If $r_1 = 0$, then $b|a$ and $\gcd(a, b) = b$.

When $r_1 \neq 0$, divide b by r_1 to produce integers q_2 and r_2 satisfying

$$b = q_2 r_1 + r_2, \quad 0 \leq r_2 < r_1.$$

If $r_2 = 0$, then stop, otherwise, proceed as before to obtain

$$r_1 = q_3 r_2 + r_3, \quad 0 \leq r_3 < r_2.$$

This division process continues until some zero remainder appears, say, at the $(n+1)^{\text{th}}$ stage where r_{n-1} is divided by r_n . [a zero remainder occurs since the decreasing sequence $b > r_1 > r_2 > \dots \geq 0$ cannot contain more than 1 integers].

So, we get the following system of equations: (10)

$$\begin{aligned} a &= q_1 b + r_1, & 0 < r_1 < b \\ b &= q_2 r_1 + r_2, & 0 < r_2 < r_1 \\ r_1 &= q_3 r_2 + r_3, & 0 < r_3 < r_2 \\ &\vdots \\ r_{n-2} &= q_n r_{n-1} + r_n, & 0 < r_n < r_{n-1} \\ r_{n-1} &= q_{n+1} r_n + 0. \end{aligned} \quad \rightarrow (*)$$

In this manner, r_n , the last non-zero remainder that appears is equal to $\gcd(a, b)$.

Lemma: If $a = qb + r$, then $\gcd(a, b) = \gcd(b, r)$.

Proof:- If $d = \gcd(a, b)$, then the relations $d|a$ and $d|b$ together imply that $d|a - qb$, or, $d|r$.

Then, d is a common divisor of both b and r . On the other hand, if c is an arbitrary common divisor of b and r , then $c|qb + r$, hence $c|a$.

This gives c a common divisor of a and b , so $c \leq d$.

It follows ~~that~~ from the definition of $\gcd(b, r)$ that $d = \gcd(b, r)$.

Note: Using the above lemma, we simply work down the displayed system of equations, obtaining $\gcd(a, b) = \gcd(b, r_1) = \dots = \gcd(r_{n-1}, r_n) = \gcd(r_n, 0) = r_n$ as claimed.

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Ex: Check how the Euclidean Algorithm works in a concrete case by calculating, say $\gcd(12378, 3054)$. The appropriate applications of the Division Algorithm produce the equations

$$12378 = 4 \cdot 3054 + 162$$

$$3054 = 18 \cdot 162 + 138$$

$$162 = 1 \cdot 138 + 24$$

$$138 = 5 \cdot 24 + 18$$

$$24 = 1 \cdot 18 + \boxed{6}$$

$$18 = 3 \cdot 6 + 0$$

So, the last non-zero remainder appearing in these equations, namely, the integer 6, is the greatest common divisor of 12378 and 3054:

$$6 = \gcd(12378, 3054).$$

To represent 6 as a linear combination of integers 12378 and 3054.

We start with the ~~and last~~ following:

$$6 = 24 - 18$$

$$= 24 - (138 - 5 \cdot 24)$$

$$= 6 \cdot 24 - 138$$

$$= 6(162 - 138) - 138$$

$$= 6 \cdot 162 - 7 \cdot 138$$

$$= 6 \cdot 162 - 7(3054 - 18 \cdot 162)$$

$$= 132 \cdot 162 - 7 \cdot 3054$$

$$= 132(12378 - 4 \cdot 3054) - 7 \cdot 3054$$

$$= 132 \cdot 12378 + (-535) \cdot 3054$$

9 has, —

$$6 = \gcd(12378, 3054) = 12378x + 3054y,$$

where $x = 132$ and $y = -535$.

Note: This is not the only way to express the integer 6 as a linear combination of 12378 & 3054; among other possibilities, we could add and subtract $3054 \cdot 12378$ to get

$$\begin{aligned} 6 &= (132 + 3054)12378 + (-535 - 12378)3054 \\ &= 3286 \cdot 12378 + (-12913)(3054). \end{aligned}$$

Note: (a) In the above example, the smaller integer is (3054). 3054 has 4 digits. So the total number of divisions cannot be greater than 20; in actuality only six divisions were needed. This is by the French Mathematician, Gabriel Lame.

(b) One observation is that for each $n > 0$, it is possible to find integers a_n & b_n such that exactly n divisions are required to compute $\gcd(a_n, b_n)$ by the Euclidean algorithm. (Prove later!)

Thm 7: If $k > 0$, then $\gcd(ka, kb) = k \gcd(a, b)$.

Proof: If each of the equations appearing in the Euclidean algorithm for a & b (see pag (10) *) is multiplied by k , we get —

$$ak = q_1 bk + r_1 k, \quad 0 < r_1 k < bk$$

$$bk = q_2(r_1 k) + r_2 k, \quad 0 < r_2 k < r_1 k$$

⋮

$$r_{n-2}k = q_n (r_{n-1}k) + r_n k, \quad 0 < r_n k < r_{n-1}k$$

$$r_{n-1}k = q_{n+1} (r_n k) + 0.$$

It is clearly the Euclidean Algorithm applied to the integers ak and bk , so that their gcd is the last non-zero remainder $r_n k$; i.e.

$$\gcd(ka, kb) = r_n k = k \gcd(a, b).$$

Corollary: For any integer $k \neq 0$, $\gcd(ka, kb) = |k| \gcd(a, b)$.

Proof: Prove it!

Note: $\gcd(ak, bk)$ is the smallest positive integer of the form $(ak)x + (bk)y$, which, is equal to k times the smallest (+ve) integer of the form $ax + by$; the latter value is equal to $k \gcd(a, b)$.

We see that —

$$\gcd(12, 30) = 3 \gcd(4, 10) = 3 \cdot 2 \gcd(2, 5) = 6 \cdot 1 = 6.$$

→ An integer c is said to be a common multiple of two non-zero integers a and b whenever $a|c$ & $b|c$.
Zero is a common multiple of a and b .

For nontrivial existence of common multiples, note that the products ab and $-(ab)$ are both common multiples of a and b , one of these is (+ve).

By Well-Ordering Principle, the set of positive common multiples of a and b must contain a smallest integer call it the least common multiple of a and b .

Defⁿ: The least common multiple of two non-zero integers a and b , denoted by $\text{lcm}(a, b)$, is the (+ve) integer m satisfying the following:

(a) $a|m$ and $b|m$

(b) If $a|c$ and $b|c$, with $c > 0$, then $m \leq c$.

eg: The positive common multiples of integers -12 and 30 are $60, 120, 180, \dots$, hence, $\text{lcm}(-12, 30) = 60$.

Remark: Given non-zero integers a and b , $\text{lcm}(a, b)$ always exist and $\text{lcm}(a, b) \leq |ab|$.

Thm ⑧: For positive integers a and b

$$\text{gcd}(a, b) \text{ lcm}(a, b) = ab.$$

Proof: Put $d = \text{gcd}(a, b)$ and write $a = dx$, $b = dy$ for integers x and y .

If $m = \frac{ab}{d}$, then $m = as = tb$, the effect of which is to make m a common multiple (positive) of a & b .

Now, let c be any positive integer that is a common multiple of a and b ; say, $c = au = bv$.
we know, \exists integers x and y satisfying

$$d = ax + by.$$

$$\text{As a result, } \frac{c}{m} = \frac{cd}{ab} = \frac{c(ax+by)}{ab} = \frac{c}{b}(x) + \frac{c}{a}(y) = ux + vy.$$

$$\Rightarrow m|c \dots \text{Then by Defⁿ of lcm, } m = \text{lcm}(a, b).$$

$$\text{So } m \leq c \dots \Rightarrow \boxed{\text{lcm}(a, b) = \frac{ab}{\text{gcd}(a, b)}}.$$