

Graphs

Shortest Path

Dijkstra's Shortest Path Algorithm

- Dijkstra's algorithm finds the shortest path from one vertex v_0 to each other vertex in a digraph.
- When it has finished, the length of the shortest distance from v_0 to v is stored in the vertex v , and the shortest path from v_0 to v is recorded in the back pointers of v and the other vertices along that path.
- The algorithm uses a priority queue, initializing it with all the vertices and then dequeuing one vertex on each iteration.

Dijkstra's Shortest Path Algorithm

- (Precondition: $G = (V, w)$ is a weighted graph with initial vertex v_0 .)
 - (Postcondition: Each vertex v in V stores the shortest distance from v_0 to v and a back reference to the preceding vertex along that shortest path.)
1. Initialize the distance field to 0 for v_0 and to ∞ for each of the other vertices.
 2. Enqueue all the vertices into a priority queue Q with highest priority being the lowest distance field value.
 3. Repeat steps 4 to 10 until Q is empty.

Dijkstra's Shortest Path Algorithm

4. (Invariant: The distance and back reference fields of every vertex that is not in Q are correct.)
5. Dequeue the highest priority vertex into x.
6. Do steps 7 to 10 for each vertex y that is adjacent to x and in the priority queue.
7. Let s be the sum of the x's distance field plus the weight of the edge from x to y.
8. If s is less than y's distance field, do steps 9 10; otherwise go back to Step 3.
9. Assign s to y's distance field.
10. Assign x to y's back reference field.

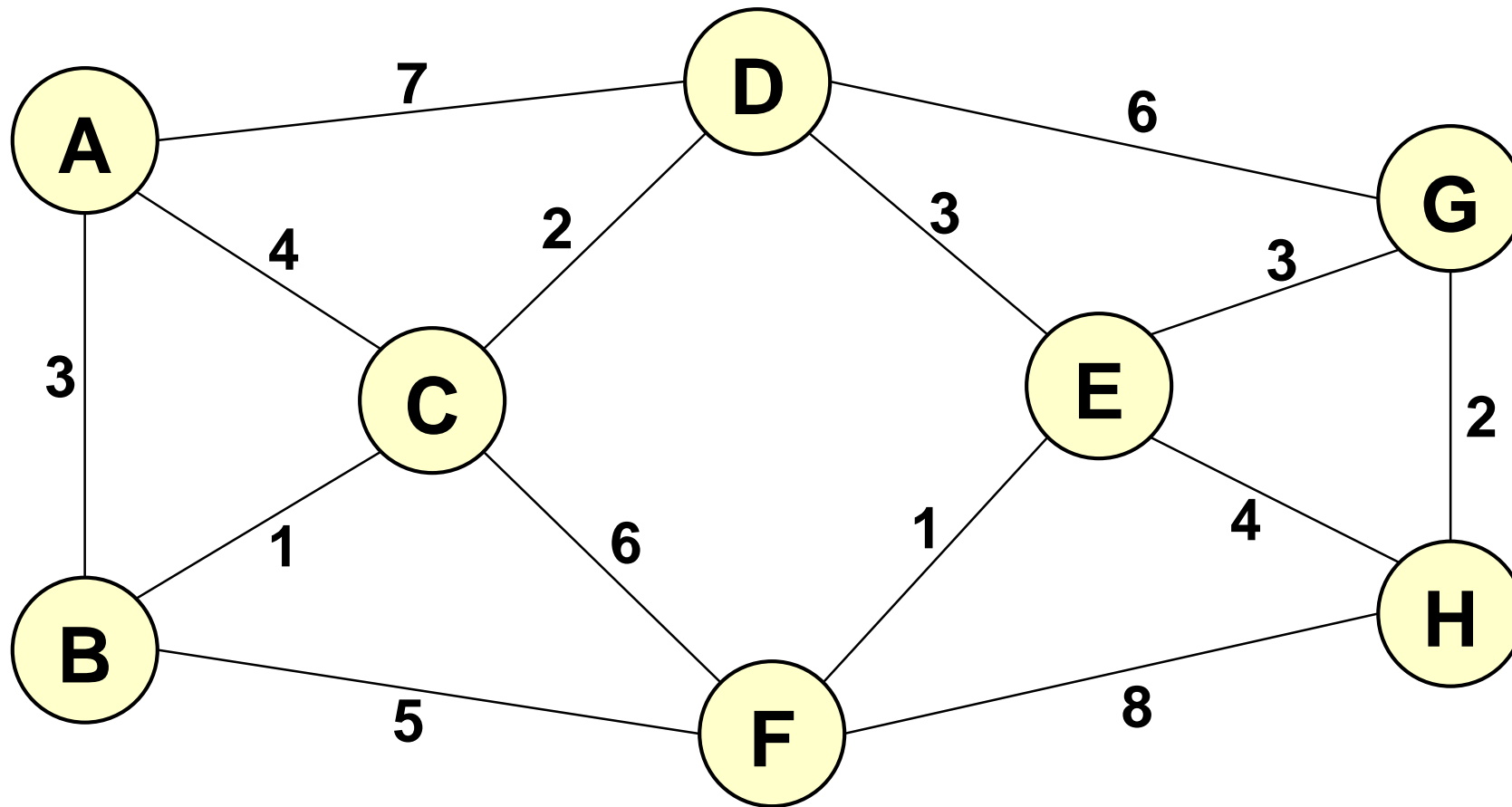
Dijkstra's Shortest Path Algorithm

Pseudocode
without using
priority queue

```
1 function Dijkstra(Graph, source):
2     create vertex set Q
3     for each vertex v in Graph:
4         dist[v] ← INFINITY
5         prev[v] ← UNDEFINED
6         add v to Q
7     dist[source] ← 0
8     while Q is not empty:
9         u ← vertex in Q with min dist[u]
10        remove u from Q
11        for each neighbor v of u:
12            alt ← dist[u] + length(u, v)
13            if alt < dist[v]:
14                dist[v] ← alt
15                prev[v] ← u
16    return dist[], prev[]
```

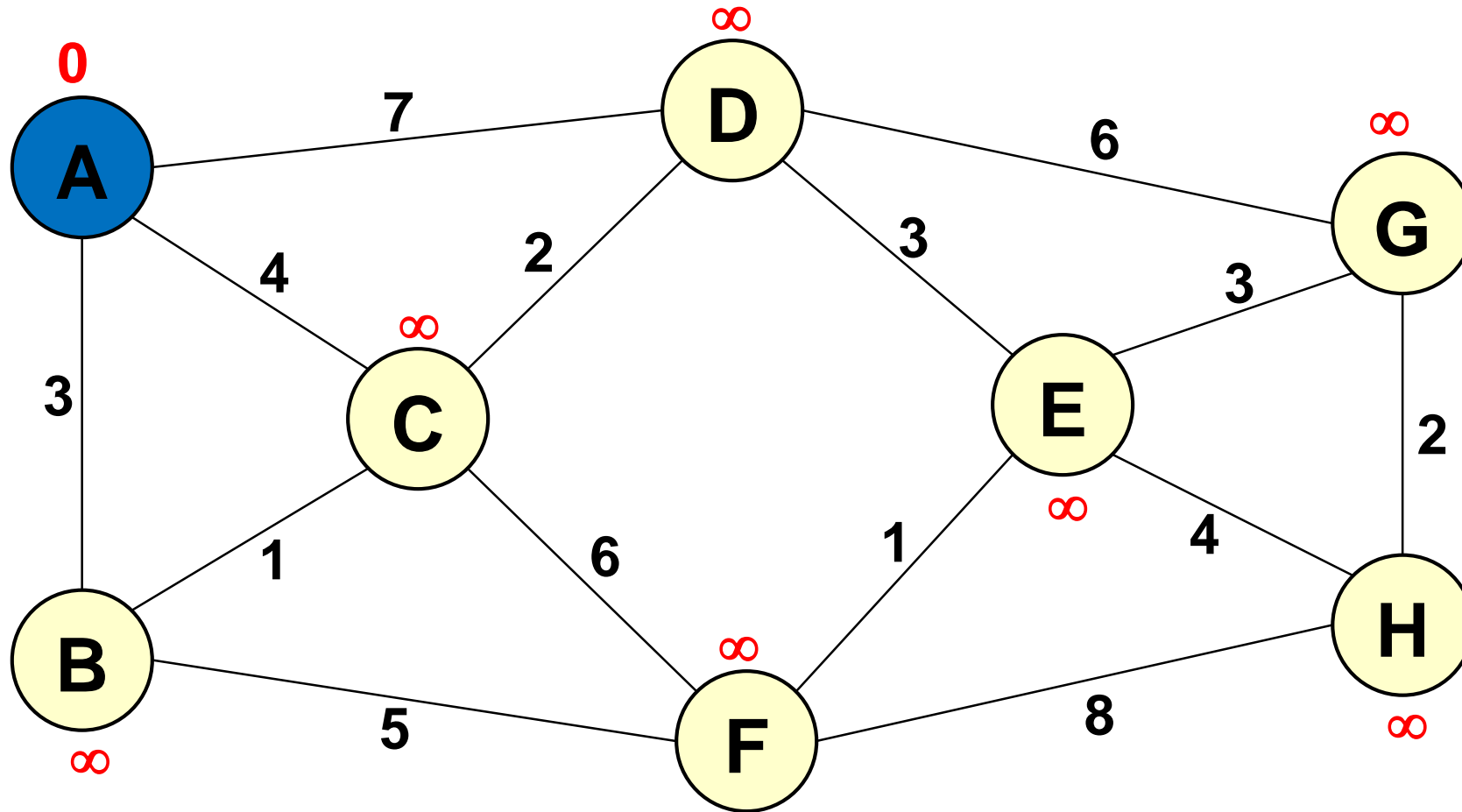

Dijkstra's Shortest Path Algorithm

Example 1:



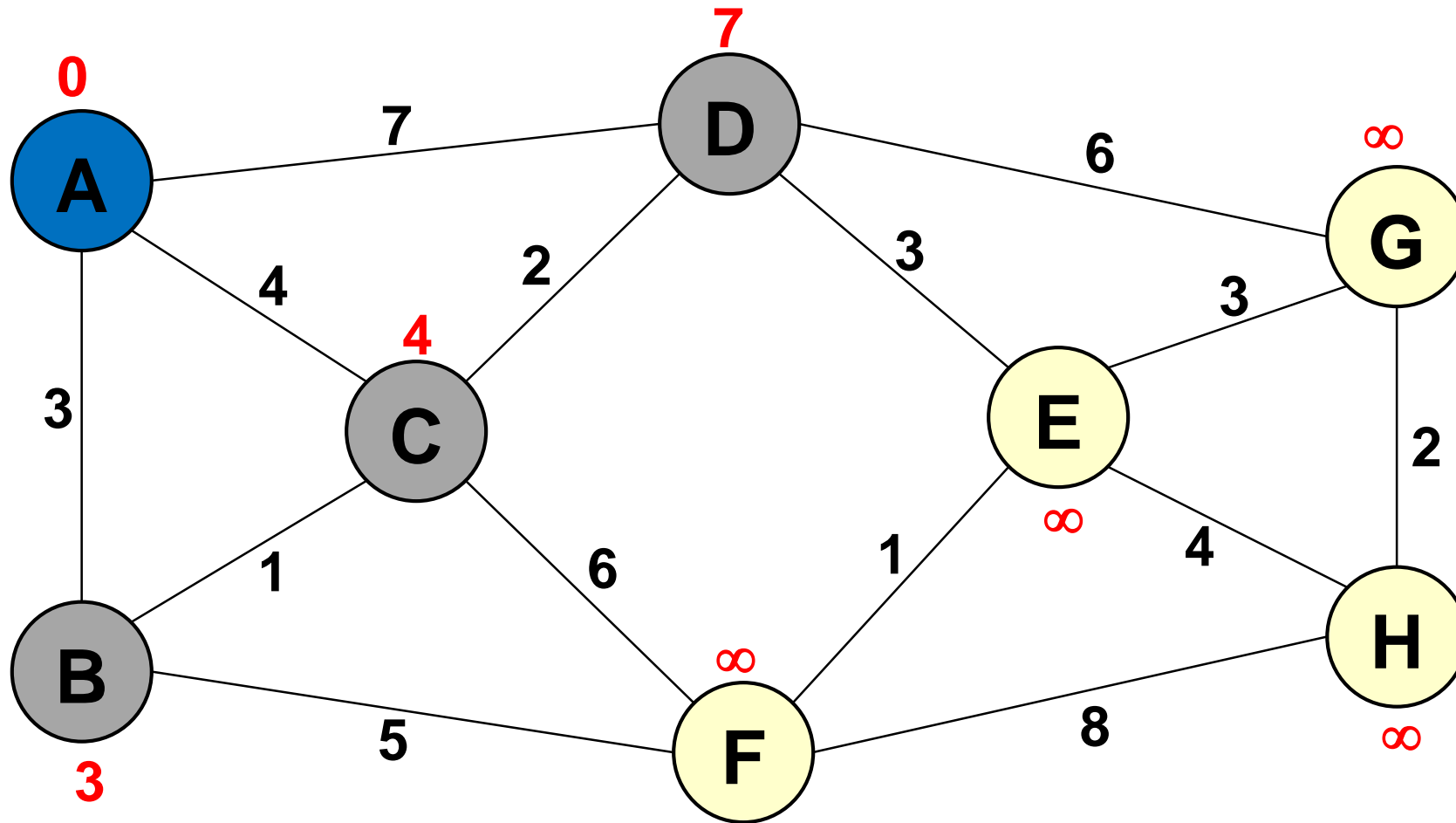
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Example 1:



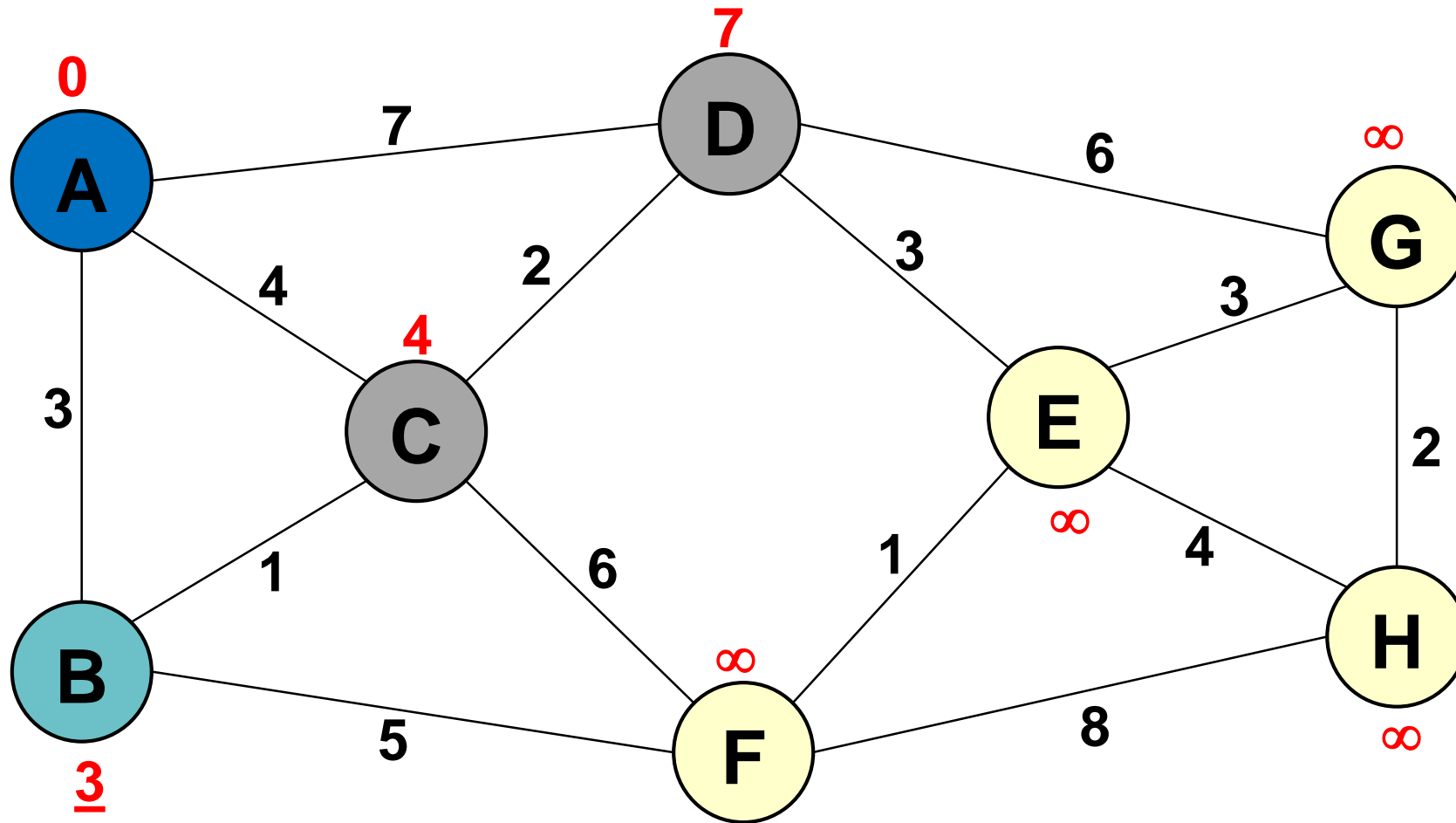
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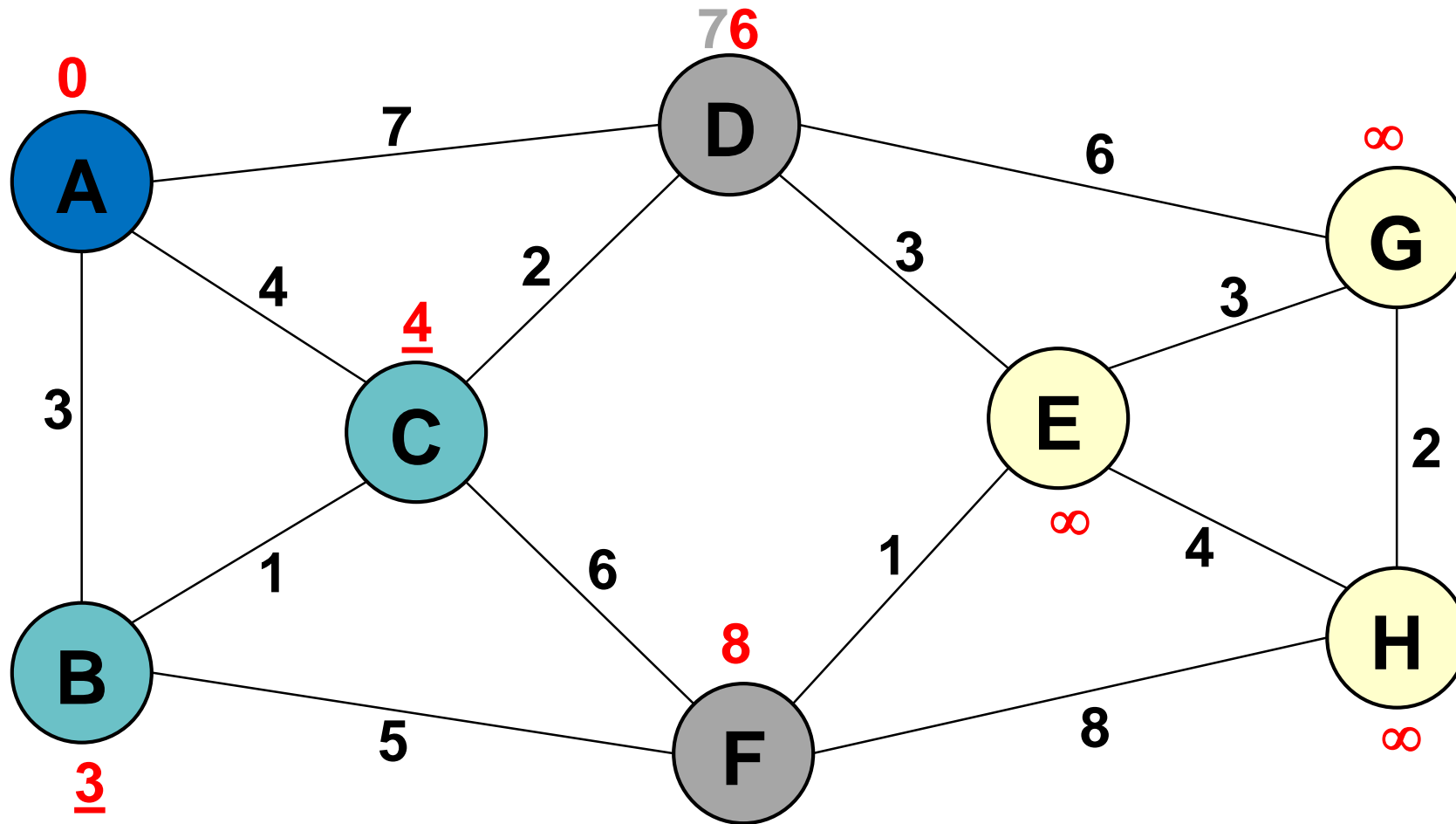
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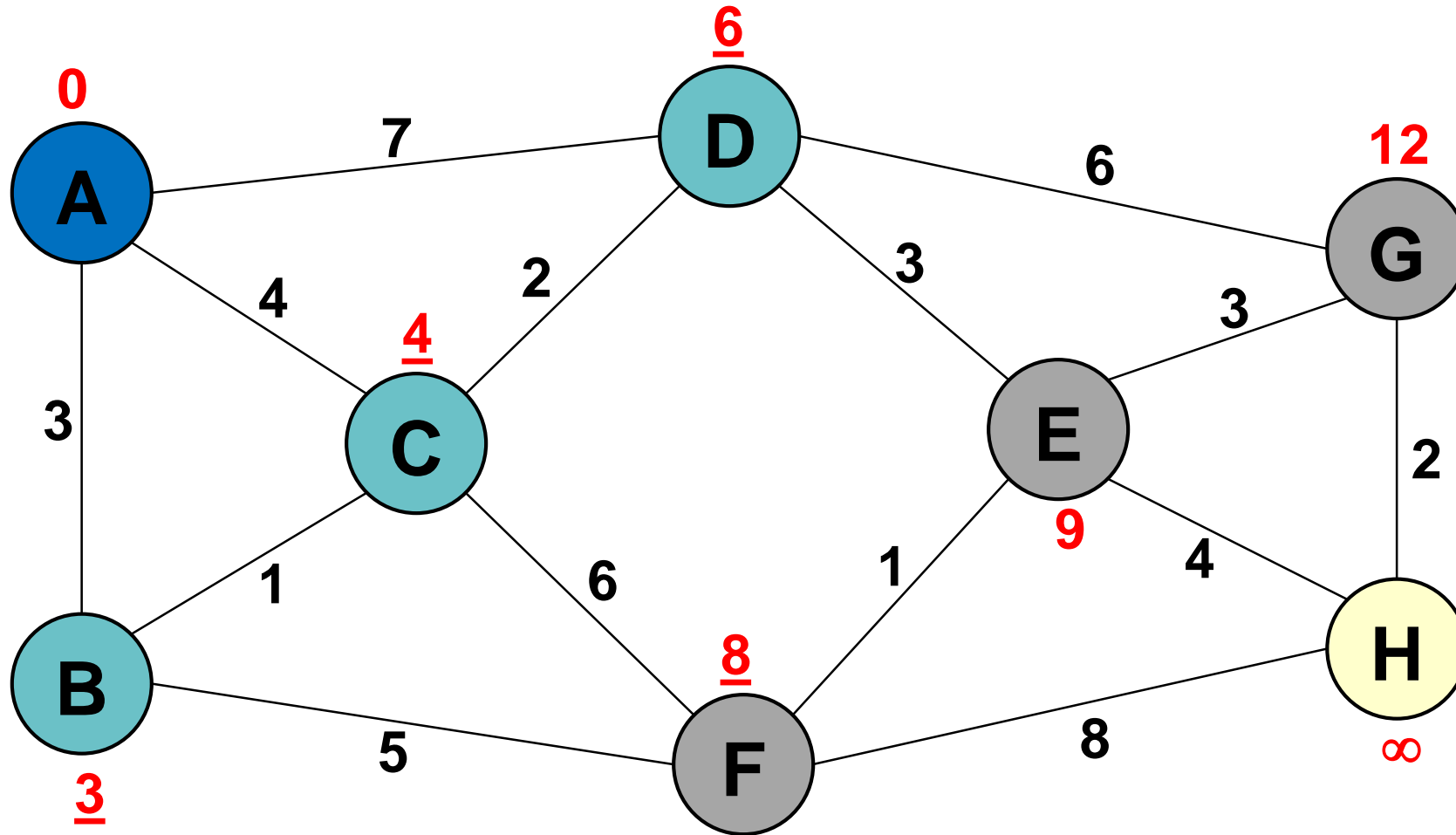
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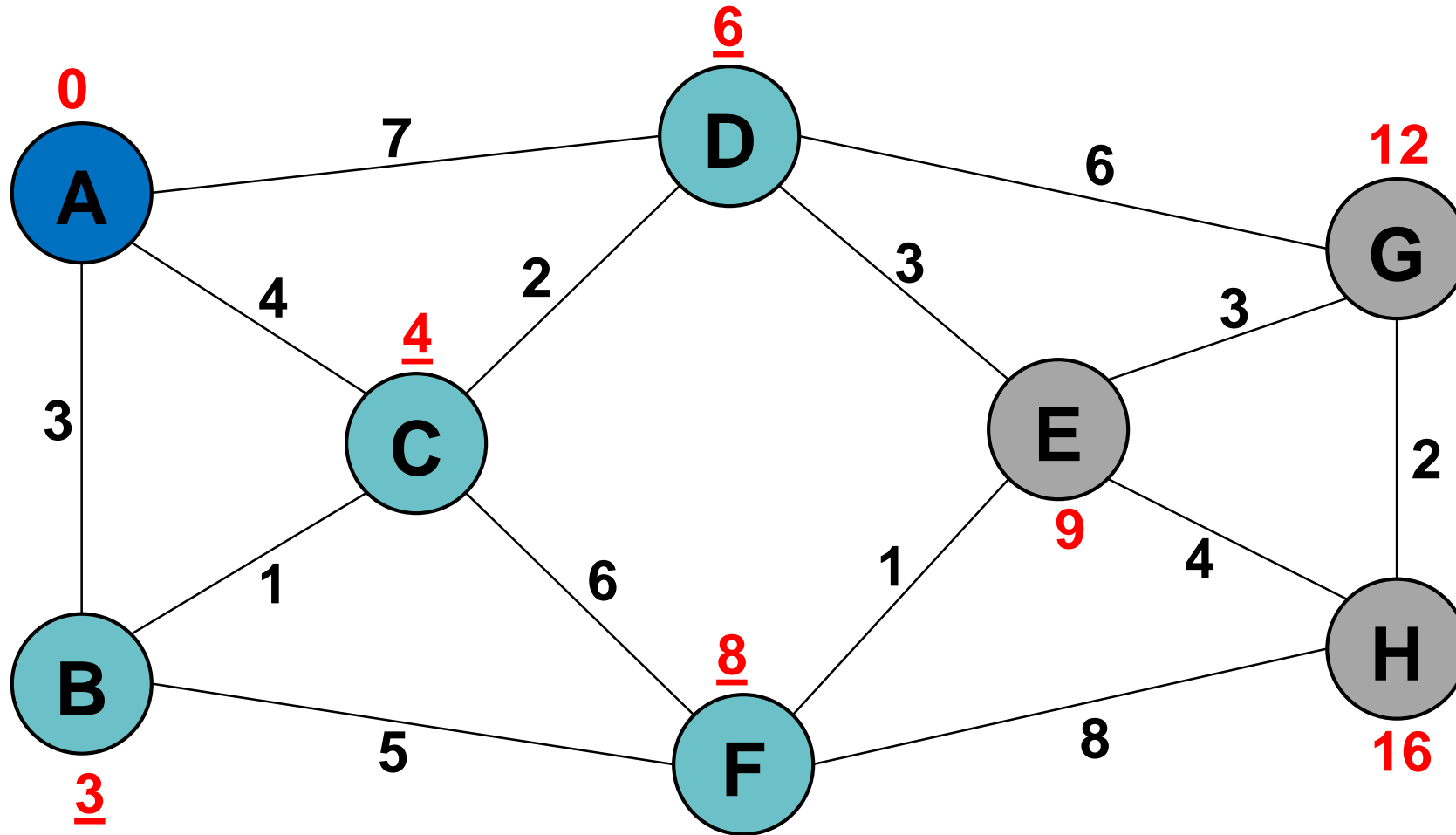
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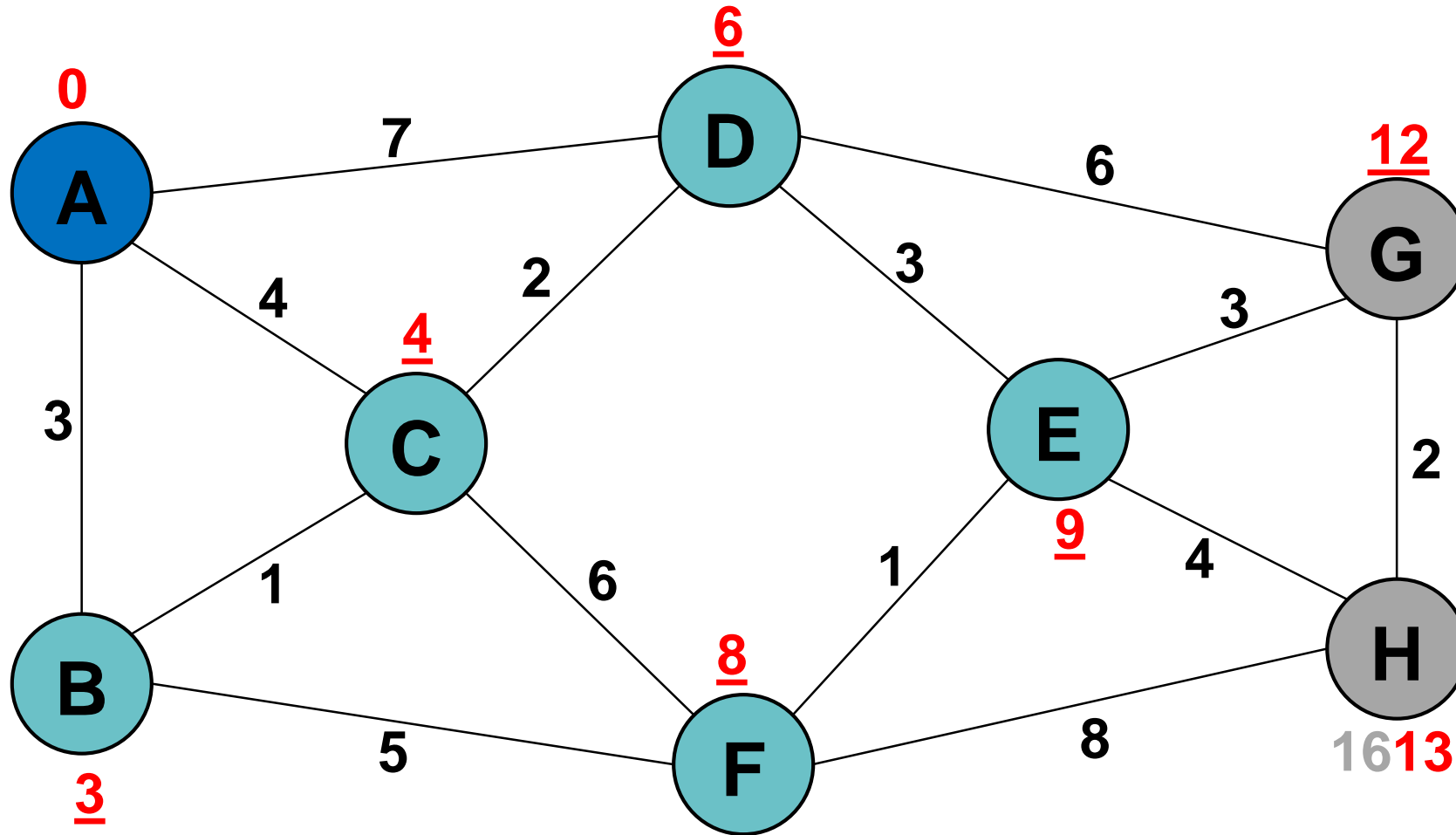
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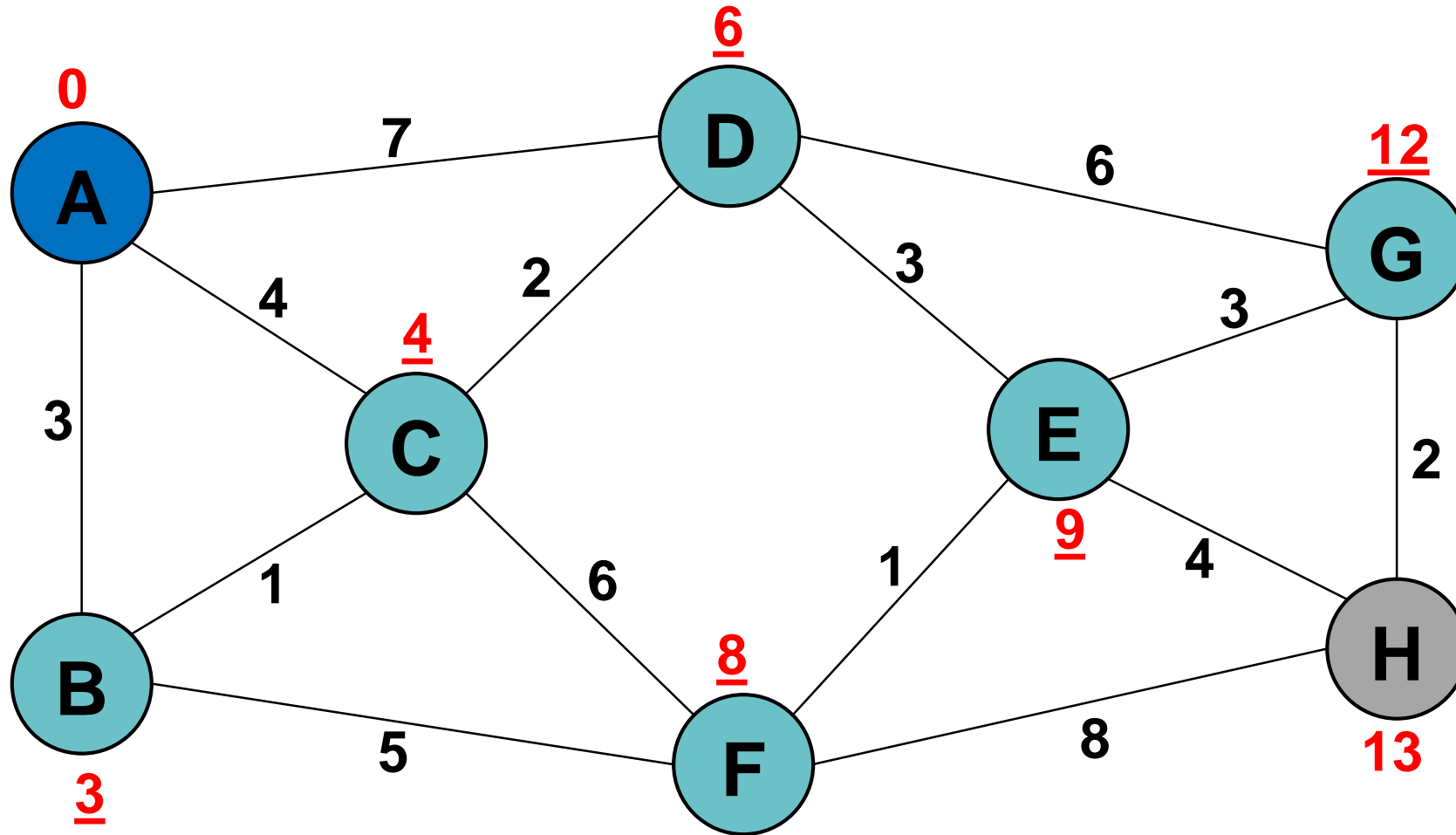
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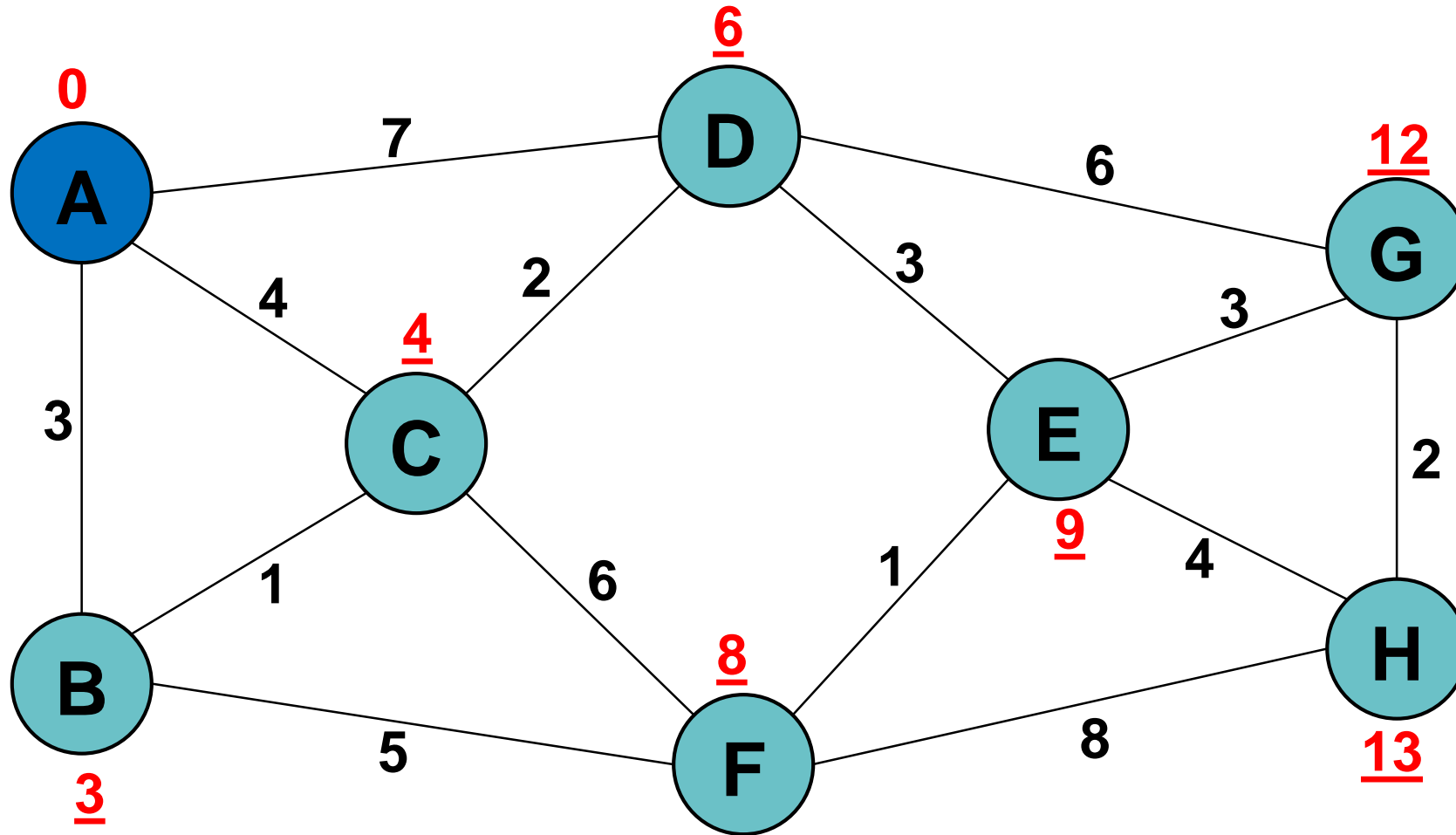
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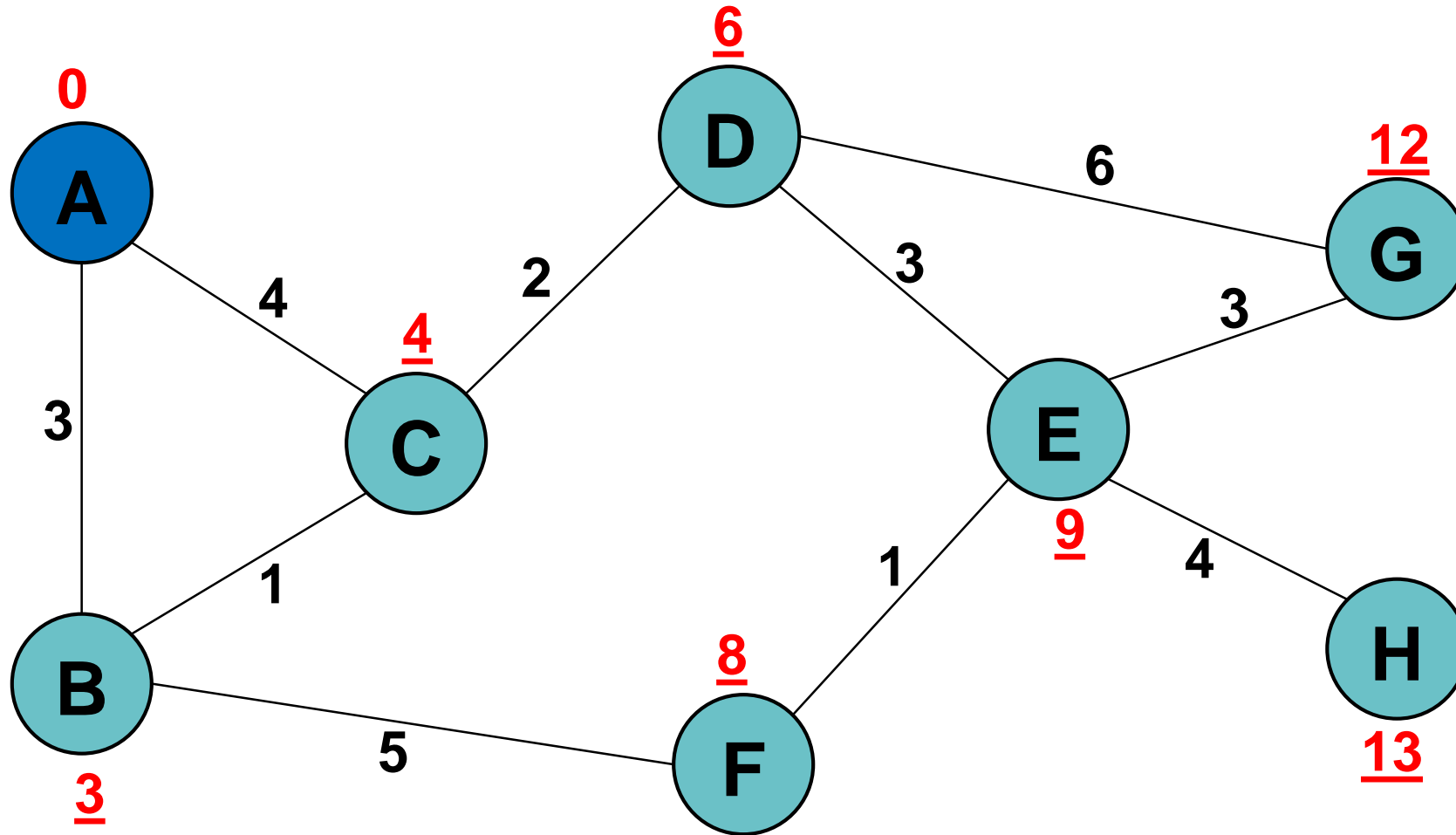
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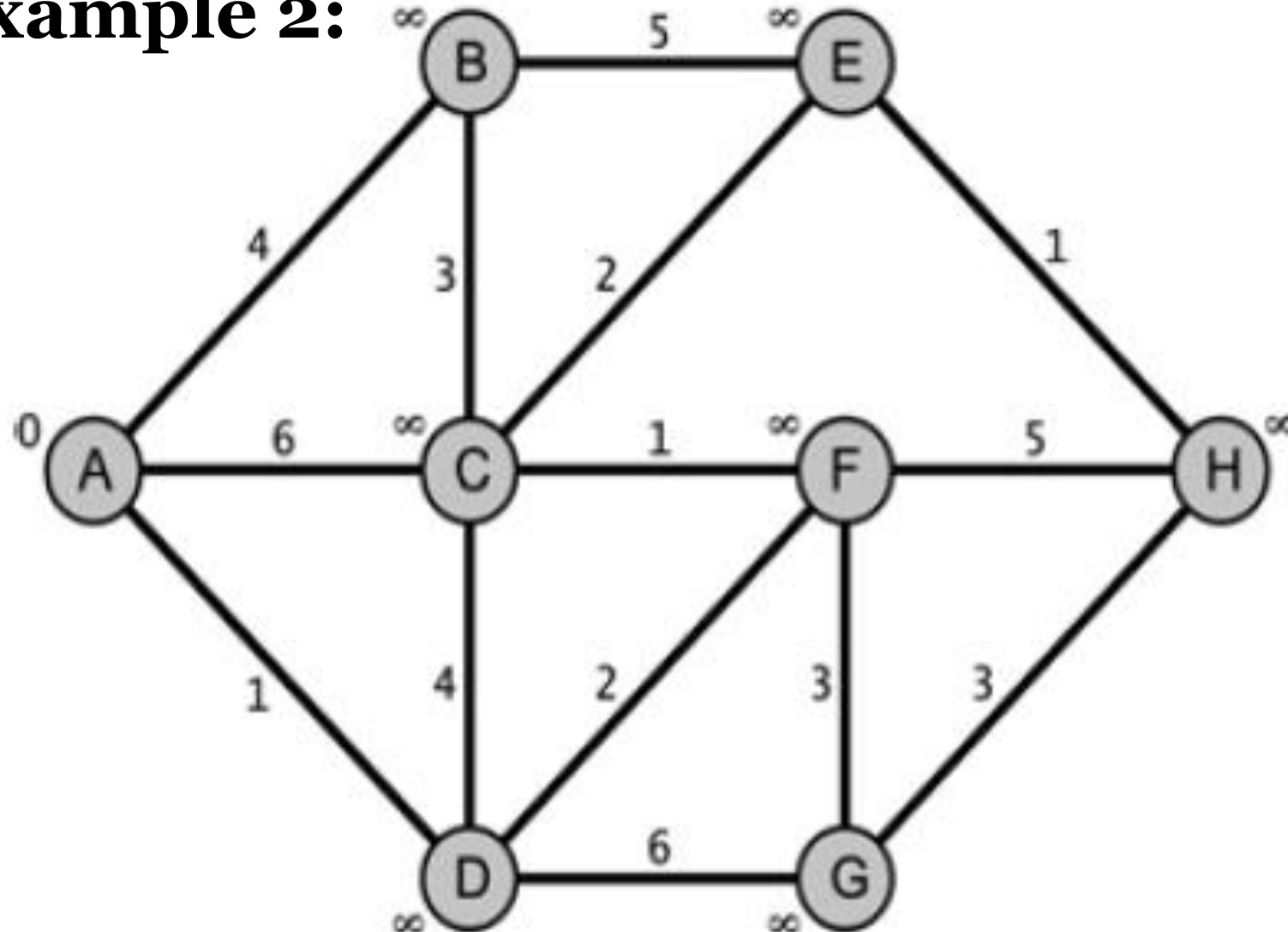
Dijkstra's Shortest Path Algorithm

Example 1: Only shortest paths are shown



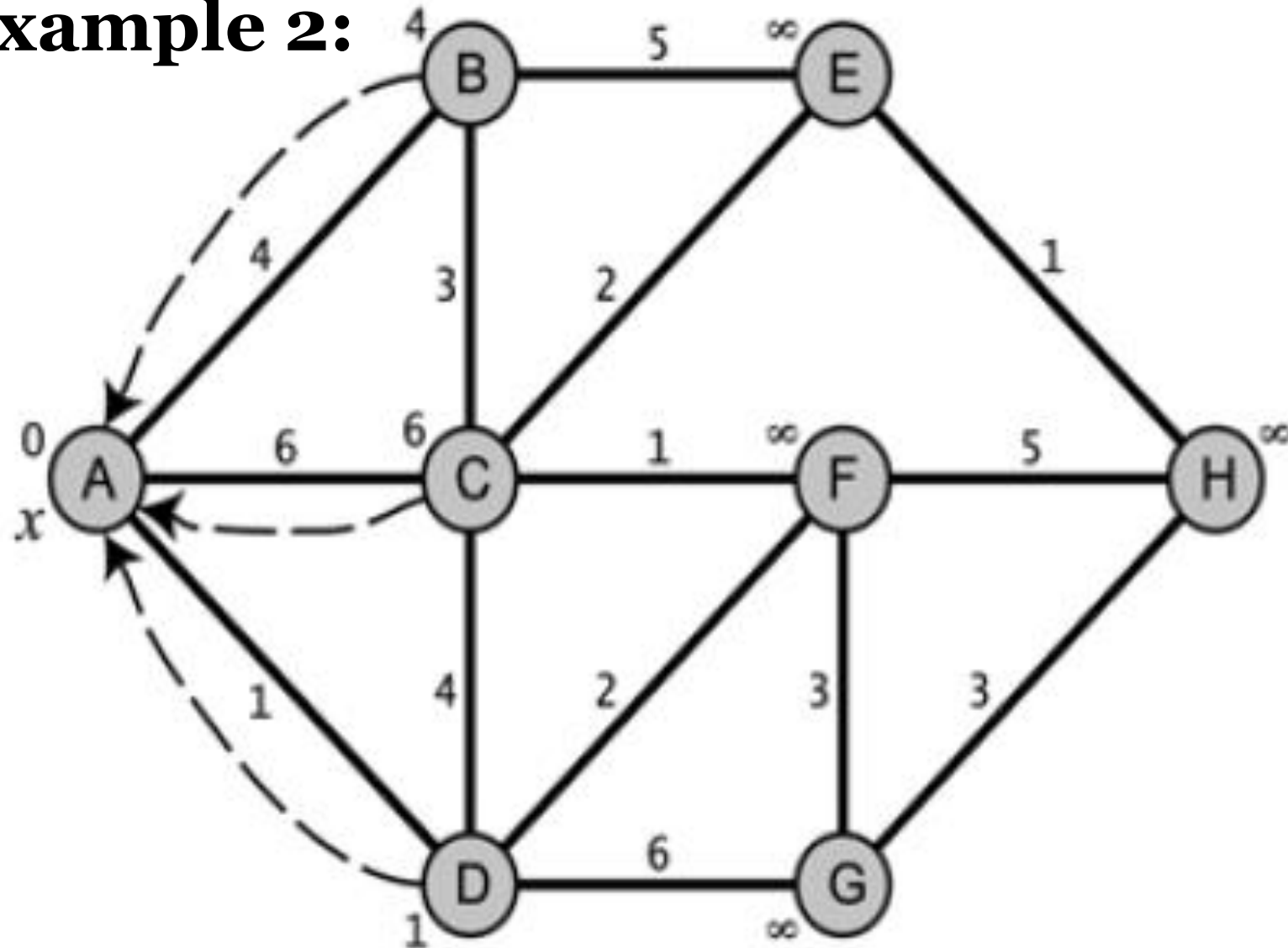
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Example 2:



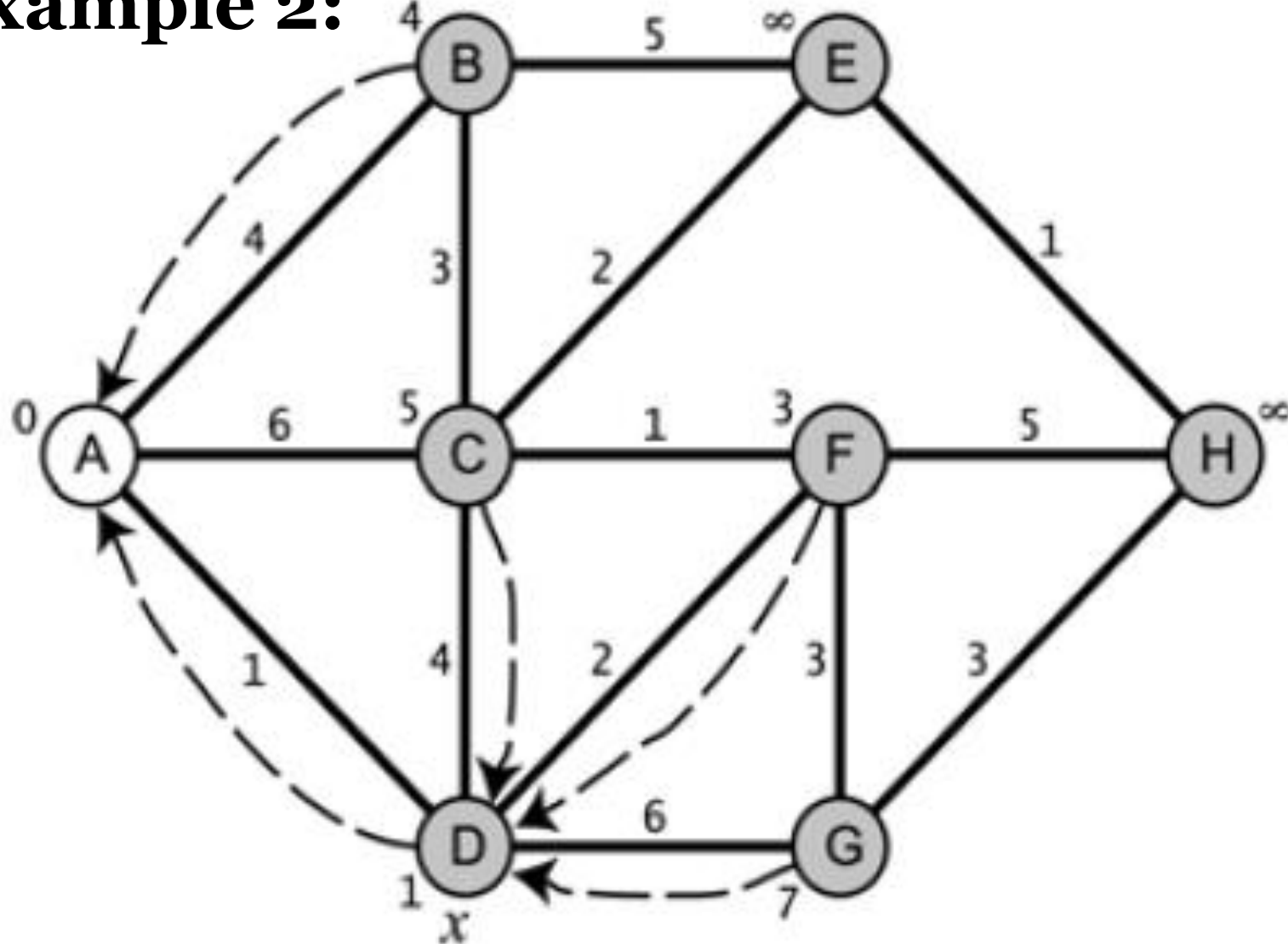
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Example 2:



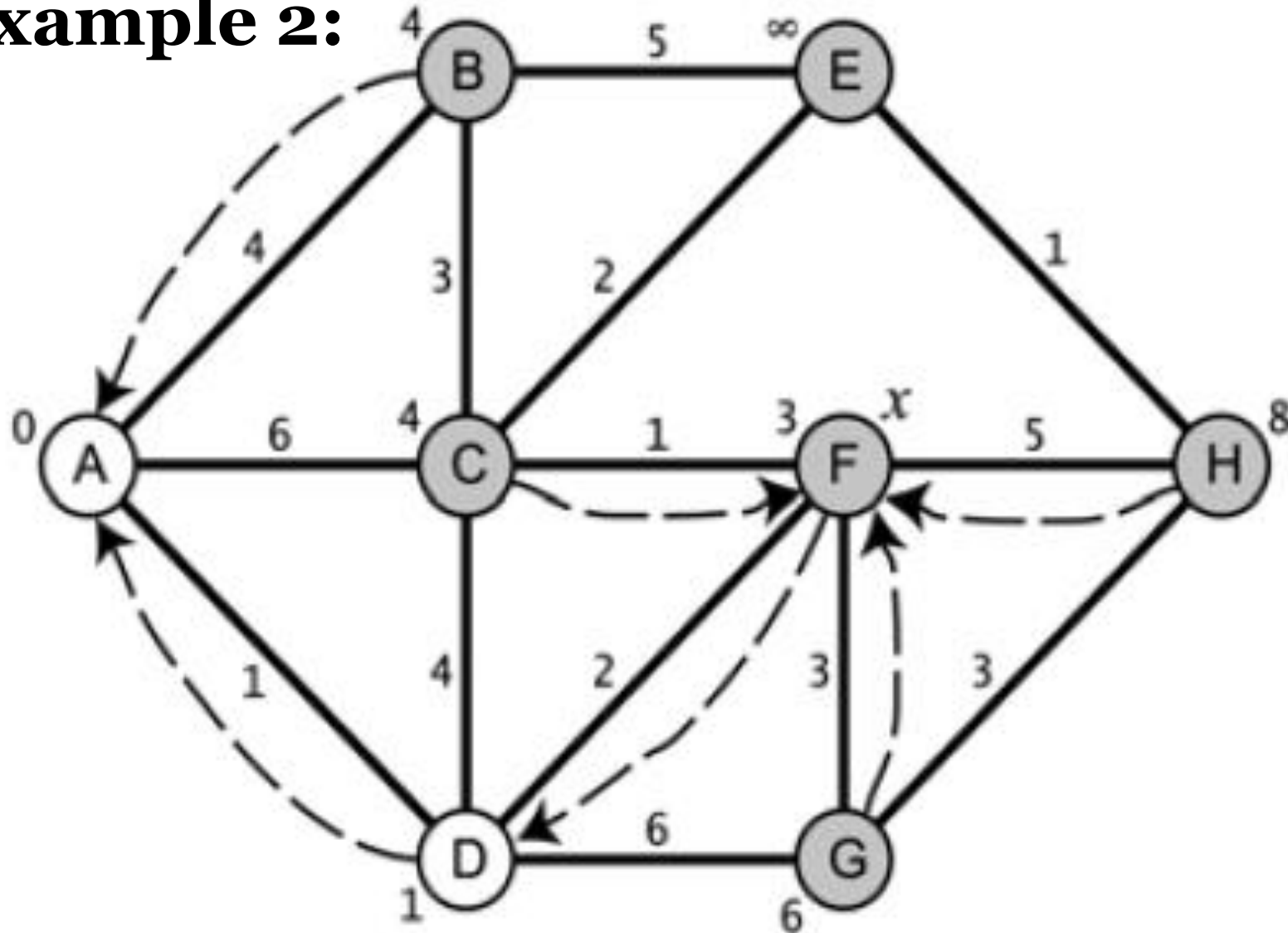
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Example 2:



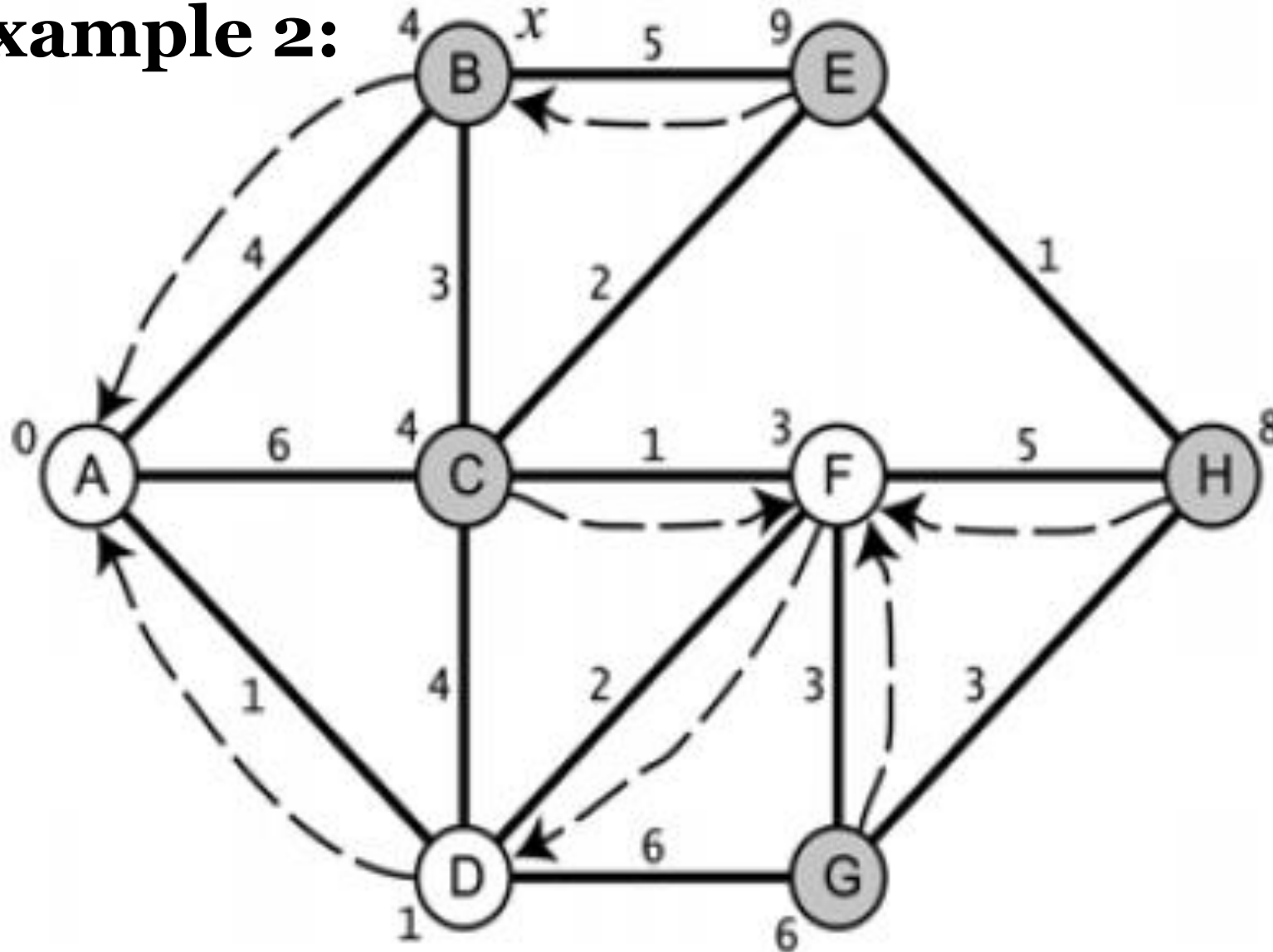
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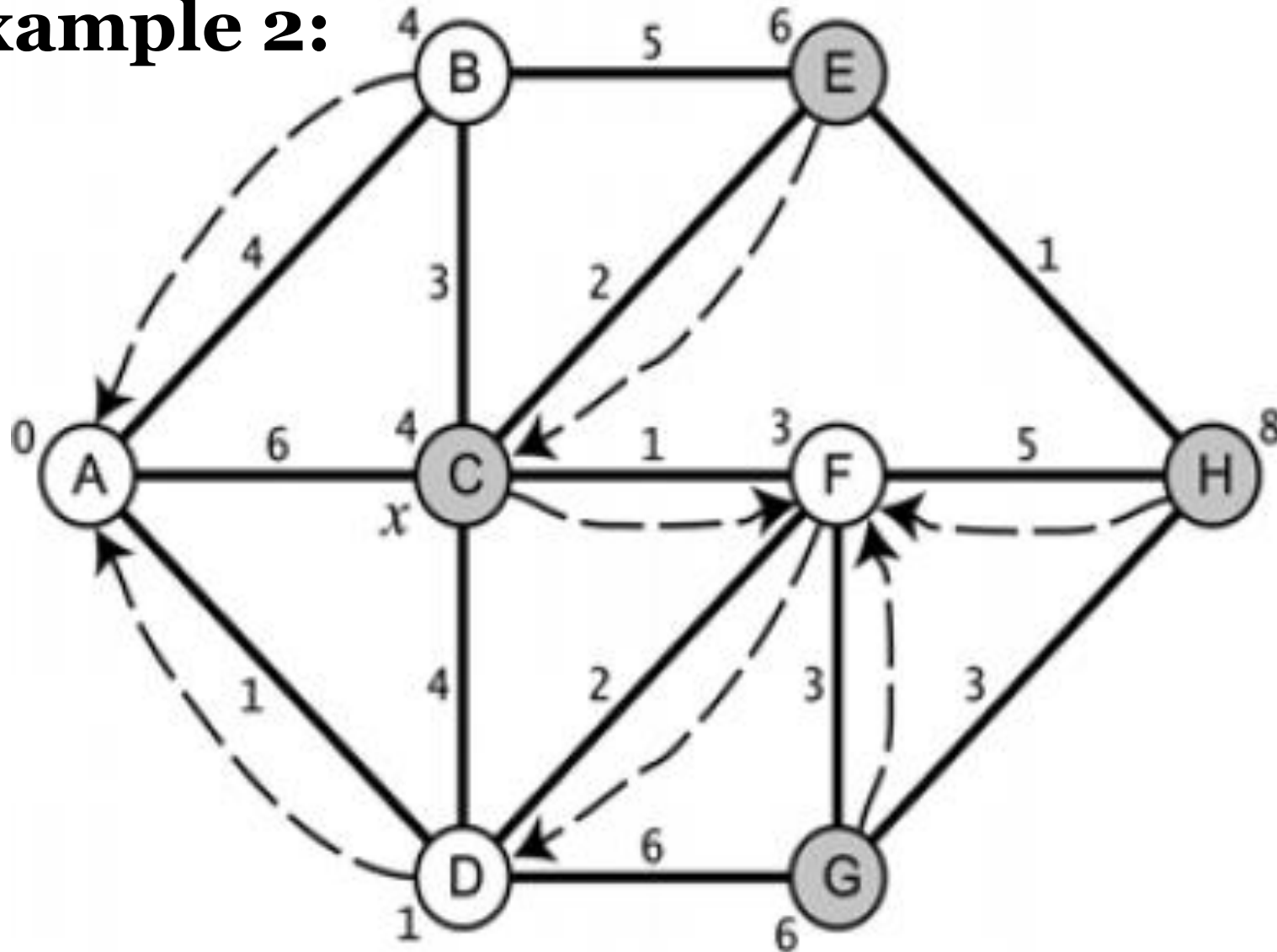
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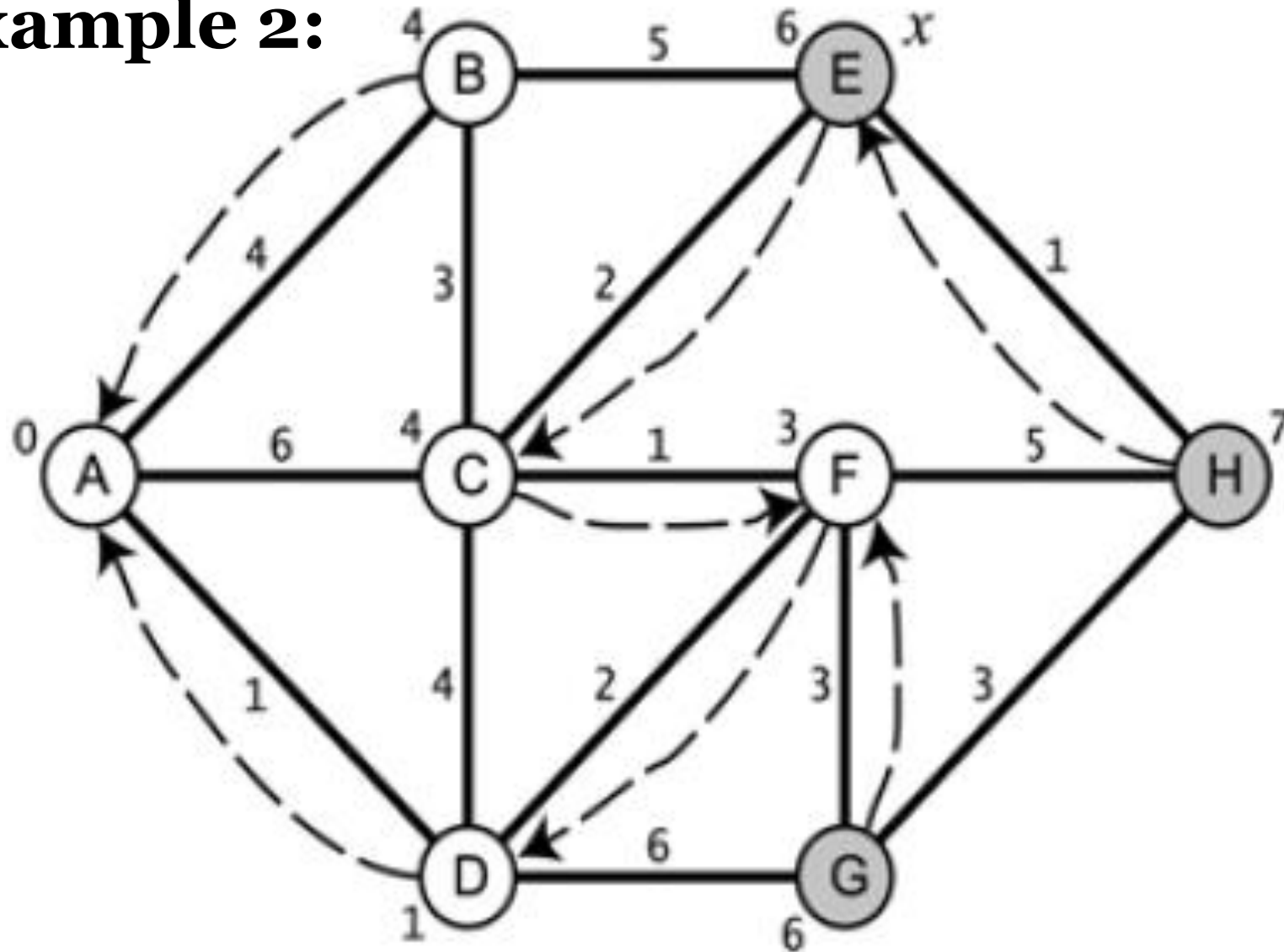
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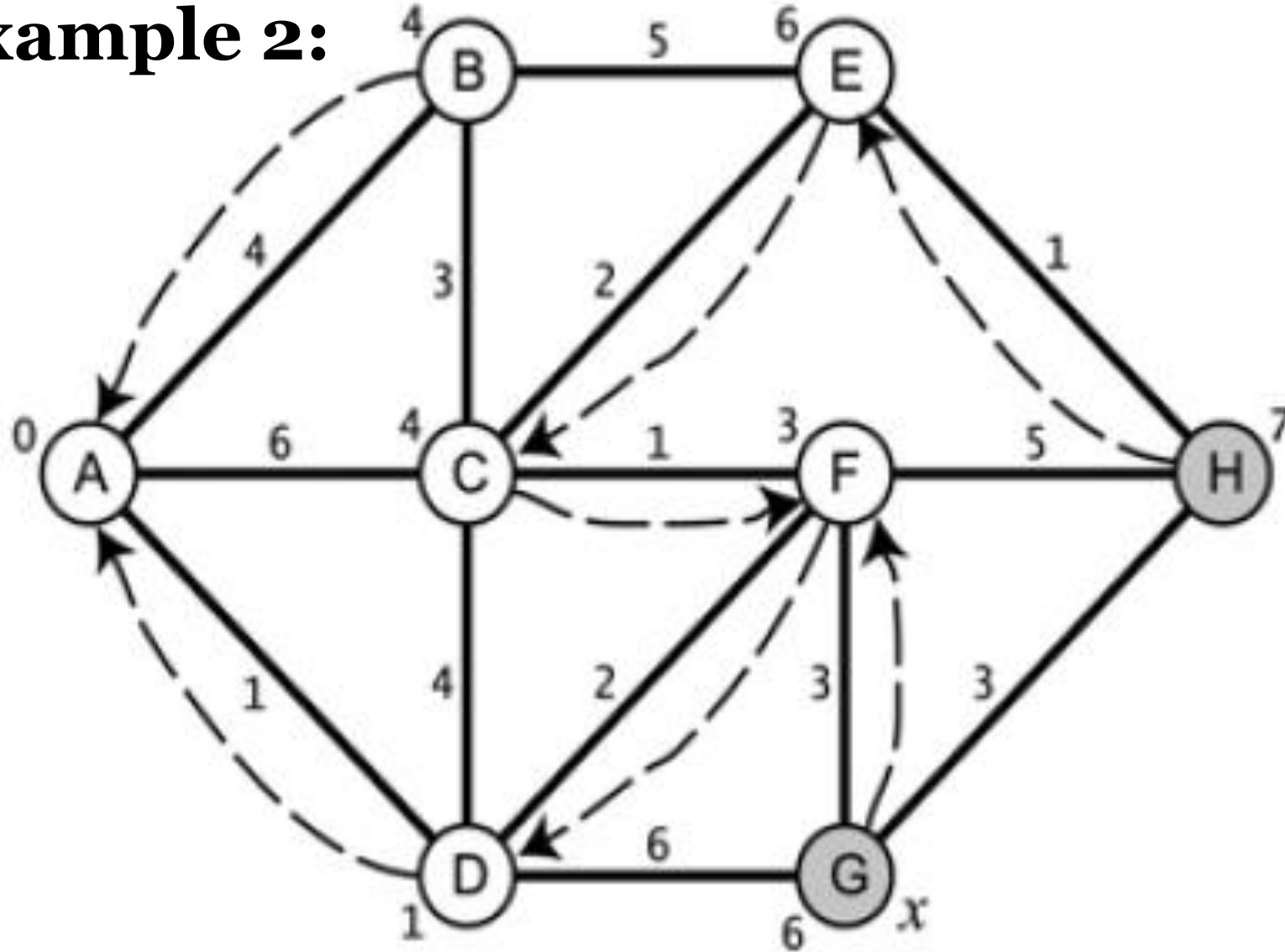
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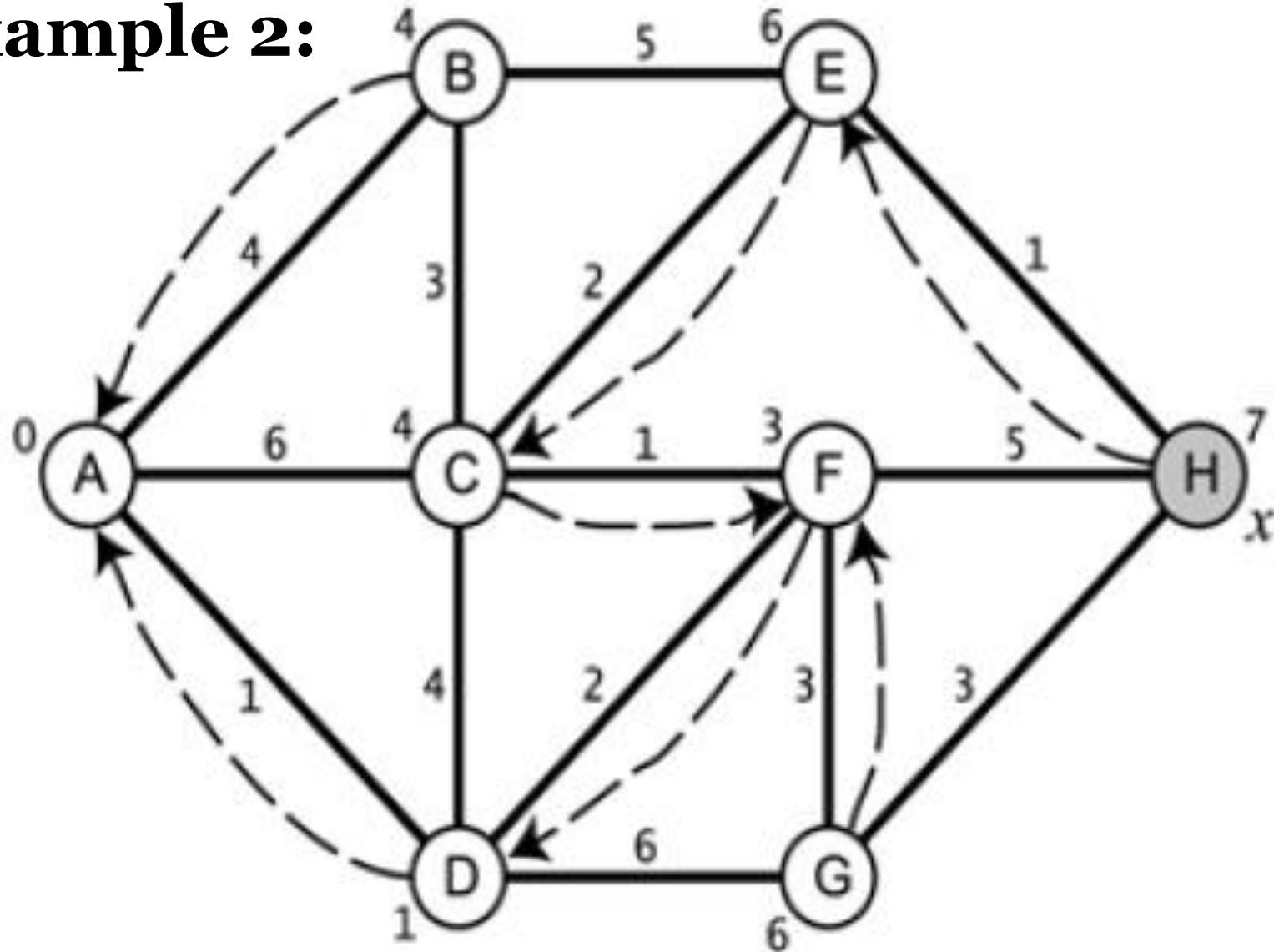
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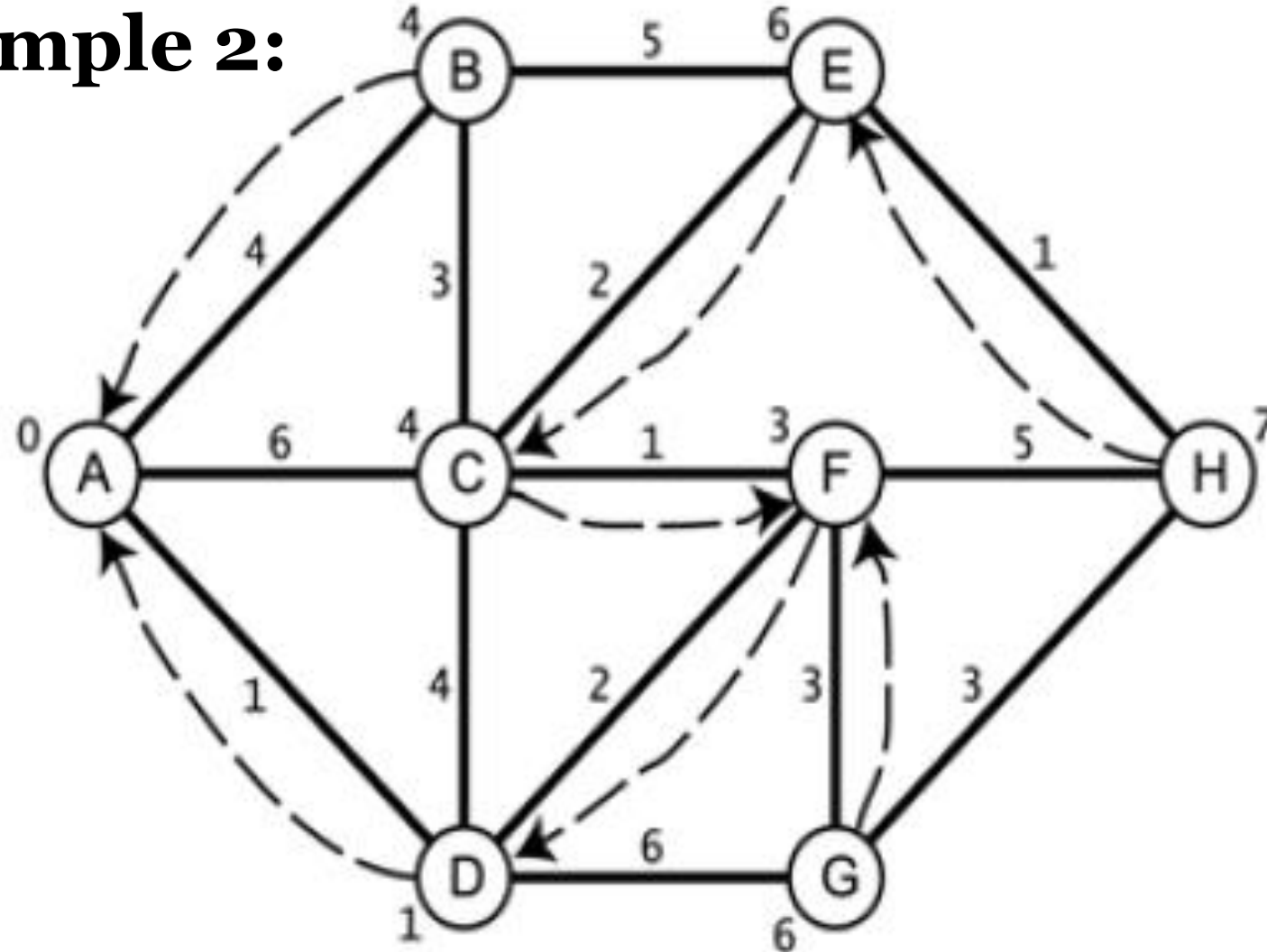
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Example 2:



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Example 2:



for example, that the shortest path from A to E is ADFCE with length 6.

Dijkstra's SPA- Time Complexity

- If the input graph is implemented using **adjacency matrix**, it is **$O(V^2)$** .
- If the input graph is represented using **adjacency list**, it can be reduced to **$O(E \log V)$** with the help of **binary heap (priority queue)**.

Thank you!