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Lecture 6-W4
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2021年2月4日 星期四 10:07

Duality Theory

max
$$Z = 4 \cdot \chi_1 + \chi_2 + 3 \cdot \chi_3$$

subject to: $\chi_1 + 4 \cdot \chi_2 \leq 1$
 $3 \cdot \chi_1 - \chi_2 + \chi_3 \leq 3$
 $\chi_1, \chi_2, \chi_3 \geqslant 0$

take constraints 0 x 2:

+ ---- (5)
$$\times 3$$
:
 $3(3x_1 - x_2 + x_3) \le 3.3$

$$(x_1, x_2, x_3)$$
 is admissible, then
$$(1x_1 + 5x_2 + 3x_3 \le 1)$$

$$(4x_1 + x_2 + 3x_3 = Z$$

$$(4x_1 + x_2 + 3x_3 = Z)$$

$$(2x_1 + x_2 + 3x_3 = Z)$$

$$(3x_1 + x_2 + 3x_3 = Z)$$

$$(4x_1 + x_2 + 3x_3 = Z)$$

What's special about factor 2 and 3?

To understand, we play the same game with two parameters 1, and 1/2.

$$y_1 (x_1 + y_2) \le y_1 \cdot 1$$

+ $y_2 (3x_1 - x_2 + x_3) \le y_2 \cdot 3$

$$\Rightarrow$$
 $(Y_1 + 3Y_2) X_1 + (+ Y_1 - Y_2) \cdot \chi_2 + \chi_2 \cdot \chi_3 \leq Y_1 + 3Y_2$

otherwise the inequalities flip!

$$\Rightarrow$$
 Y₁ and Y₂ sortisfy
$$\begin{cases} y_1 + 3y_2 > 4 \\ 4y_1 - y_2 \geq 1 \\ y_2 \geq 3 \end{cases}$$

$$\begin{cases} y_1, y_2 \geq 0 \\ y_1 + 3y_2 \geq 2 \end{cases}$$
this is satisfied
$$\begin{cases} y_1 + 3y_2 \geq 2 \end{cases}$$

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y_1+3y_2 \geqslant 2
y_1+3y_2 \geqslant \max 2
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Each choice of Y1 and Y2 is going to generate a different estimate:

 $\max z \leq y_1 + 3y_2$

what is the best possible estimate? max Z = min S

Dual Optimization problem

Ans :

minimum $y_1 + 3y_2 = 6$ subject to $y_1 + 3y_2 \ge 4$ $4y_1 - y_2 \ge 1$ $y_2 \ge 3$ $y_{1,1} y_2 \ge 0$ Linear optima fation parablem

max Z & min &

Rmk:

$$\frac{2(x_{min}) \leq 7(x) \leq 8(x_{max})}{4}$$

$$-\frac{2(x_{min}) \geq -8(x) \geq -8(x_{max})}{8}$$

meaning that:

· Xmin is the max ot - S

· Xmax is the min of - 8

• $\min S = \max - S$; $\min - S = \max S$

ex 3:
$$\max \ 7 = -3x_3 + x_4$$

S.T $\Rightarrow \begin{cases} x_1 = 6 + 3x_3 - 3x_4 \\ x_2 = 4 + 8x_3 - 4x_4 \\ x_1 - x_4 = 20 \end{cases}$

ex: max = 5 4 C1-X1 - ... + Cn-Xn + Co

$$\begin{cases} a_{11} \times_{1} + \cdots + a_{1n} \cdot \chi_{n} \leq b_{1} & y_{1} \left(a_{11} \times_{1} + \cdots + a_{1n} \cdot \chi_{n} \right) \leq y_{1} b_{1} \\ a_{21} \times_{1} + \cdots + a_{2n} \cdot \chi_{n} \leq b_{2} \\ \vdots & \vdots & \vdots \\ a_{m1} \cdot \chi_{1} + \cdots + a_{mn} \cdot \chi_{n} \leq b_{m} \\ \chi_{1} \cdots \chi_{n} \geq 0 \end{cases} + y_{m} \left(a_{m1} \cdot \chi_{1} + \cdots \cdot a_{mn} \cdot \chi_{n} \right) \leq y_{n} \cdot b_{m}$$

RHS = 4. b1 + 12 b2 + ... /m. bm

$$\left(\sum_{i=1}^{m} \gamma_{i} \alpha_{j_{i}}\right) \cdot \gamma_{i} + \dots + \left(\sum_{j=1}^{m} \gamma_{j} \cdot \alpha_{j_{n}}\right) \gamma_{n} \leq \sum_{j=1}^{m} \gamma_{j} \cdot b_{j}$$

$$(\sum_{j=1}^{m} \gamma_{j} \alpha_{j1}) \cdot \chi_{1} + \dots + (\sum_{j=1}^{m} \gamma_{j} \cdot \alpha_{jn}) \chi_{n} \leq \sum_{j=1}^{m} \gamma_{j} \cdot b_{j}$$

$$\frac{1}{2} \cdot \frac{1}{2} \cdot \frac{1}{2}$$

and However we choose the admissible vectors (X1 ... Xn) and (Y1 ... Xm), we have

$$Z = \sum_{i=1}^{n} C_i \cdot x_i \le \sum_{j=1}^{m} b_j \cdot y_j = 0$$

In particular X = X maxY = Y min we get

max Z = Z(Xmax) & B(Ymin) = min 8) weak duality theorem.

This leads to $\begin{cases}
\frac{m}{2} \text{ Yiaji } \geq c, & \text{min } S = \frac{m}{2} \text{ bi. Yj. + co} \\
\frac{m}{3} \text{ Yi.ajn } \geq c, & \text{Duality Problem.} \\
\text{Yr... } \text{Ym} \geq 0
\end{cases}$

max z min 8

Duality theorem. There is no gap: max = min 8