Product Rule, Bijection Rule

DISCRETE STRUCTURES II

DARRYL HILL

BASED ON THE TEXTBOOK:

DISCRETE STRUCTURES FOR COMPUTER SCIENCE: COUNTING, RECURSION, AND PROBABILITY

BY MICHIEL SMID

How Effective is My Password?

What makes a "good" password?

- Hard to guess
- Algorithm to guess a password:
 - Try every combination
- We want lots of possible combinations

How do we count all possible combinations?

Simple example

Strict password rules:

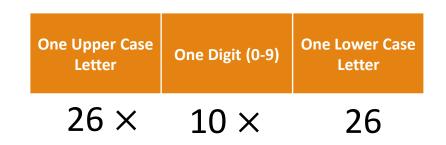
- 3 characters Upper case, digit, lower case
- How many combinations is that?

Examples: A7d, G4s, P0x, etc.

One Upper Case Letter		One Lower Case Letter
26 ×	10 ×	26

We have to do 3 tasks:

- 1. Choose an upper case letter (26 ways to do this)
- 2. Choose a digit (10 ways to do this)
- 3. Choose a lower case letter (26 ways to do this)



= 6760

Important! How we do Task 1 does not affect the number of ways to do Task 2!

Likewise, once we have finished Task 1 and Task 2, these do not affect the number of ways we can do Task 3

These Tasks are Independent

Product Rule in general:

Procedure: We have a sequence of m tasks to be done in *order*.

For i = 1, 2, 3, ..., m let N_i be the number of ways to do task i, and N_i does not depend on how the previous tasks were done.

Then the number of ways to do the Procedure is:

$$N_1 \times N_2 \times N_3 \times \cdots \times N_m$$

This rule is used *extensively* throughout the course.

 $n \ge 1$ we want bitstrings of length n. How many are there?

A bitstring is a string of two symbols, 0's and 1's.

Example: if n = 5 then possible bitstrings are

01001

00000

11110

11111

etc.

Does this fit the product rule?

Product Rule:

Procedure: We have a sequence of m tasks to be done in *order*.

For i = 1, 2, 3, ..., m let N_i be the number of ways to do task i, and N_i does not depend on how the previous tasks were done.

Then the number of ways to do the Procedure is:

$$N_1 \times N_2 \times N_3 \times \cdots \times N_m$$

How many bitstrings of length n, $n \ge 1$ are there?

What is the procedure?

Procedure: Write a bitstring one bit at a time from left to right.

Tasks: For i = 1, 2, ..., n task i = write 0 or 1.

 $N_i = 2$ since there are two ways to do task i

Notice if we have already done tasks 1 ... i - 1, then task i still has two ways to do it

1110001 - still two choices, 0 or 1

Since we have a sequence of n tasks and there is no dependency, we can apply the product rule

How many bitstrings of length n, $n \ge 1$ are there?

What is the procedure?

Procedure: Write a bitstring one bit at a time from left to right.

Tasks: For i = 1, 2, ..., n task i = write 0 or 1.

 $N_i = 2$ since there are two ways to do task i

Total number of ways to write a bitstring of length 5 is

$$= N_1 \cdot N_2 \cdot N_3 \cdot N_4 \cdot N_5$$

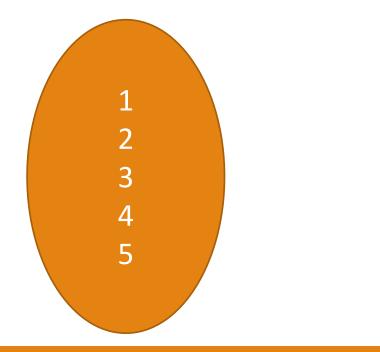
$$= 2 \cdot 2 \cdot 2 \cdot 2 \cdot 2$$

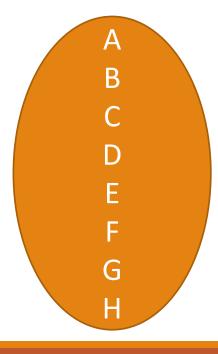
$$= 32$$

Sets A and B, |A| = m, |B| = n. How many functions $f: A \rightarrow B$?

(Recall the definition of a function)

Can we define all functions in a way that uses the product rule?

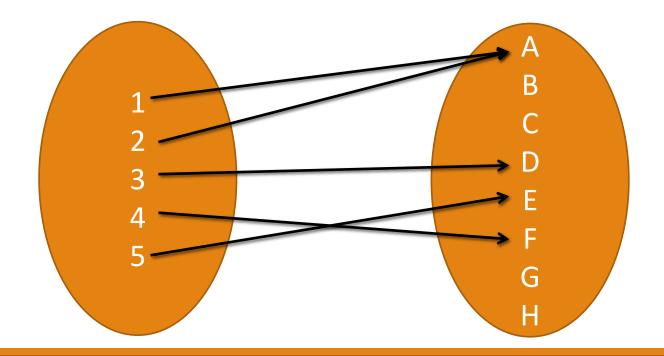




Sets A and B, |A| = m, |B| = n. How many functions $f: A \rightarrow B$?

Each element in A is mapped to exactly one element in B.

Example function:

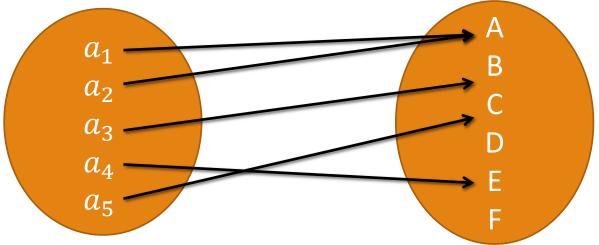


Sets A and B, |A| = m, |B| = n. How many functions $f: A \to B$? (Let's generalize the elements.)

What procedure did we use to create this? What are the tasks?

Procedure: $\forall x \in A$, specify f(x) (assign an element from B).

Tasks: For each element $a_i \in A$, $i = 1 \dots m$, task i = choose an element $b \in B$ such that $f(a_i) = b$. How many ways can we do this?



Sets A and B, |A| = m, |B| = n. How many functions $f: A \rightarrow B$?

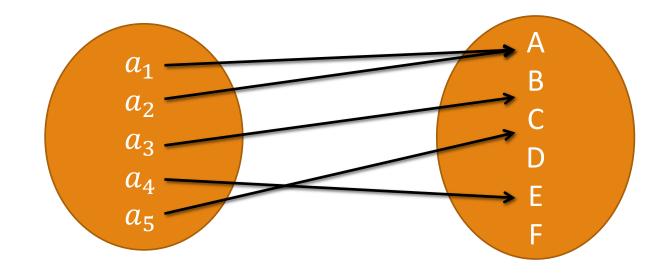
Procedure: $\forall x \in A$, specify f(x) (assign an element from B).

Tasks: For each element $a_i \in A$, $i = 1 \dots m$, task i = choose an element $b \in B$ such that $f(a_i) = b$.

$$N_1 = n, N_2 = n, N_3 = n, ..., N_m = n$$

$$N_1 \cdot N_2 \cdot N_3 \cdot \dots \cdot N_m = n \cdot n \cdot \dots \cdot n = n^m$$

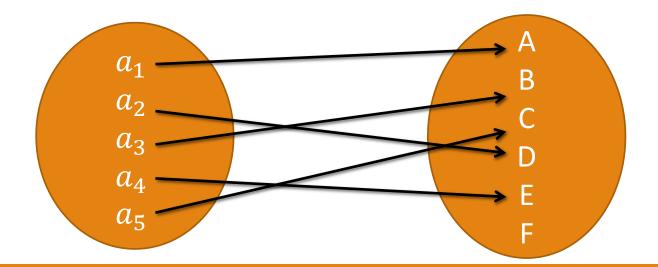
In general: $|B|^{|A|}$



Sets A and B, |A| = m, |B| = n. How many **injective** functions $f: A \to B$?

Procedure: $\forall x \in A$, specify f(x) (assign an element from B).

Tasks: For each element $a_i \in A$, $i = 1 \dots m$, task i = choose an element $b \in B$ such that $f(a_i) = b$ and b has not been selected yet. How many ways can we do this?

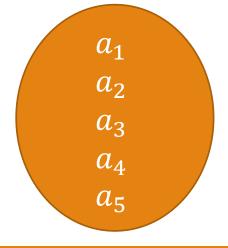


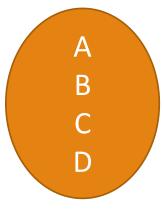
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Tasks: For each element $a_i \in A$, $i = 1 \dots m$, task i = choose an element $b \in B$ such that $f(a_i) = b$ and b has not been selected yet. How many ways can we do this?

What if |B| < |A|? How many 1-to-1 functions are there?

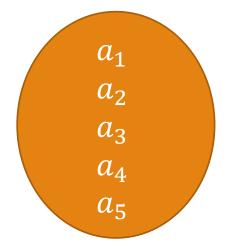


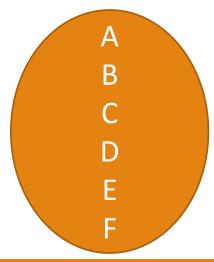


Sets A and B, |A| = m, |B| = n. How many **injective** functions $f: A \to B$? $|B| \ge |A|$

Procedure: $\forall x \in A$, specify f(x) (assign an element from B).

Tasks: For each element $a_i \in A$, $i = 1 \dots m$, task i = choose an element $b \in B$ such that $f(a_i) = b$ and b has not been selected yet. How many ways can we do this?





Sets A and B, |A| = m, |B| = n. How many **injective** functions $f: A \to B$? $|B| \ge |A|$

Procedure: $\forall x \in A$, specify f(x) (assign an element from B).

Tasks: For each element $a_i \in A$, $i = 1 \dots m$, task i = choose an element $b \in B$ such that $f(a_i) = b$ and b has not been selected yet.

$$N_1 = n$$

 $N_2 = n - 1$
 $N_3 = n - 2$
...
 $N_{m-1} = n - (m - 2) = n - m + 2$
 $N_m = n - m + 1$

Sets A and B, |A| = m, |B| = n. How many **injective** functions $f: A \to B$? $|B| \ge |A|$

$$N_1 \cdot N_2 \cdot N_3 \cdot ... \cdot N_m = n \cdot (n-1) \cdot ... \cdot (n-m+2) \cdot (n-m+1)$$

How can we express this? Recall factorials:

$$0! = 1$$

if $k \ge 1$, then $k! = 1 \cdot 2 \cdot \dots \cdot k$

$$5! = 1 \cdot 2 \cdot 3 \cdot 4 \cdot 5$$
$$3! = 1 \cdot 2 \cdot 3$$

Can we manipulate factorials into the expression $(n-m+1)\cdot (n-m+2)\cdot ...\cdot (n-1)\cdot n$?

Sets A and B, |A| = m, |B| = n. How many **injective** functions $f: A \to B$? $|B| \ge |A|$

How do we get from:

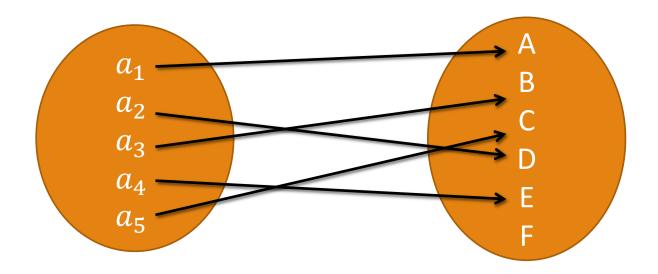
$$n! = 1 \cdot 2 \cdot ... \cdot (n-m) \cdot (n-m+1) \cdot (n-m+2) \cdot ... \cdot (n-1) \cdot n$$
 to
$$(n-m+1) \cdot (n-m+2) \cdot ... \cdot (n-1) \cdot n$$

Divide by the first part:

$$\frac{1 \cdot 2 \cdot \dots \cdot (n-m) \cdot (n-m+1) \cdot (n-m+2) \cdot \dots \cdot (n-1) \cdot n}{1 \cdot 2 \cdot \dots \cdot (n-m)} = \frac{n!}{(n-m)!}$$

Sets A and B, |A| = m, |B| = n. How many **injective** functions $f: A \to B$? $|B| \ge |A|$

There are $\frac{n!}{(n-m)!}$ injective functions $f: A \to B$, where |A| = m, |B| = n.



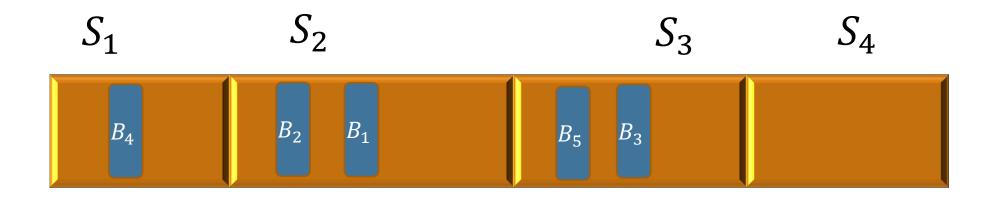
Review of Product Rule

Express as a procedure

Where a procedure is a **sequence of tasks**

The choice we make for Task i does not affect the number of ways to do Task i+1

Then we can apply the product rule



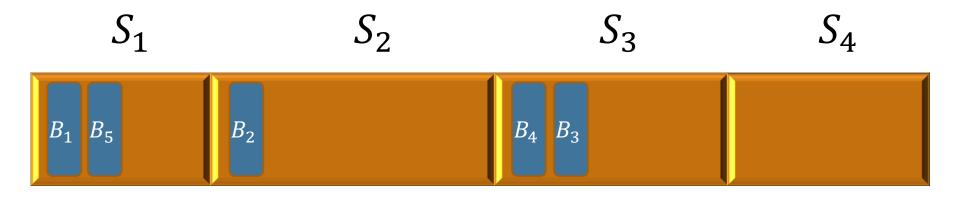
$$m \text{ books } B_1, B_2, B_3, ..., B_m$$

 $n \text{ shelves } S_1, S_2, S_3, ..., S_n$

Procedure: Place books on shelves – two things to keep track of

- 1. For each book, which shelf
- 2. For each shelf: what is the left-to-right ordering of books on this shelf

Observe that this order is different from

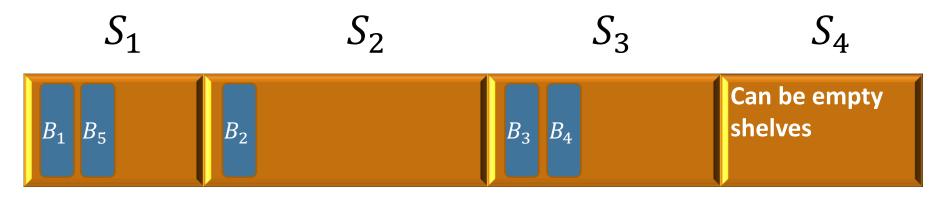


$$m$$
 books $B_1, B_2, B_3, ..., B_m$
 n shelves $S_1, S_2, S_3, ..., S_n$

Procedure: Place books on shelves – two things to keep track of

- 1. For each book, which shelf
- 2. For each shelf: what is the left-to-right ordering of books on this shelf

Observe that this order is different from this order (B_3 and B_4 are switched)

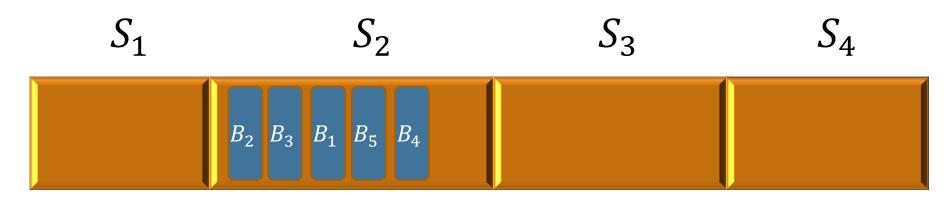


$$m$$
 books $B_1, B_2, B_3, \dots, B_m$
 n shelves $S_1, S_2, S_3, \dots, S_n$

Procedure: Place books on shelves – two things to keep track of

- 1. For each book, which shelf
- 2. For each shelf: what is the left-to-right ordering of books on this shelf

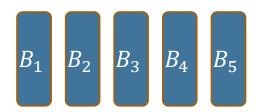
Can have all books on one shelf (order still matters)



$$m = 5$$
 books $B_1, B_2, B_3, ..., B_m$
 $n = 4$ shelves $S_1, S_2, S_3, ..., S_n$

Can use the Product rule

Procedure: Start with empty shelves, place books one by one



$$N_1 = 4$$

$$\mathcal{S}_1$$

$$S_2$$

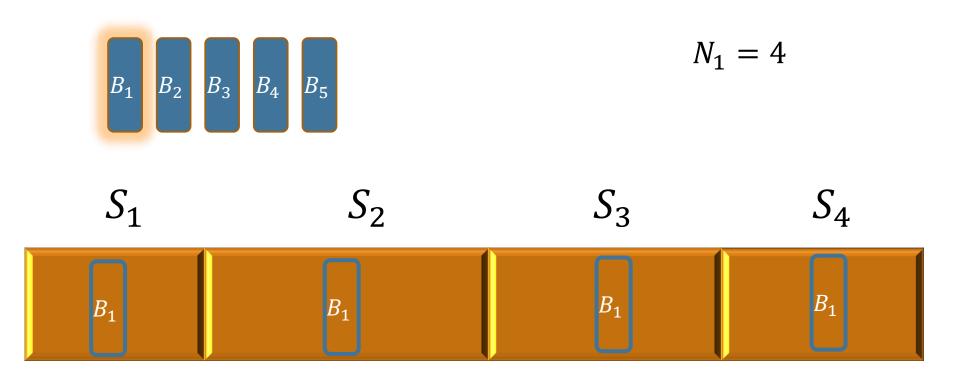
$$S_3$$

$$S_4$$

$$m = 5$$
 books $B_1, B_2, B_3, ..., B_m$
 $n = 4$ shelves $S_1, S_2, S_3, ..., S_n$

Can use the Product rule

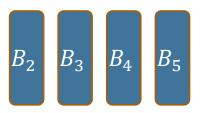
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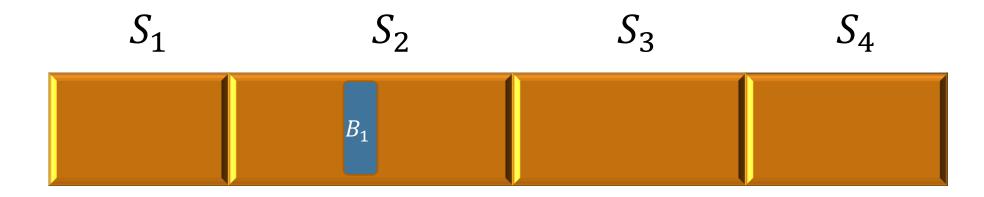
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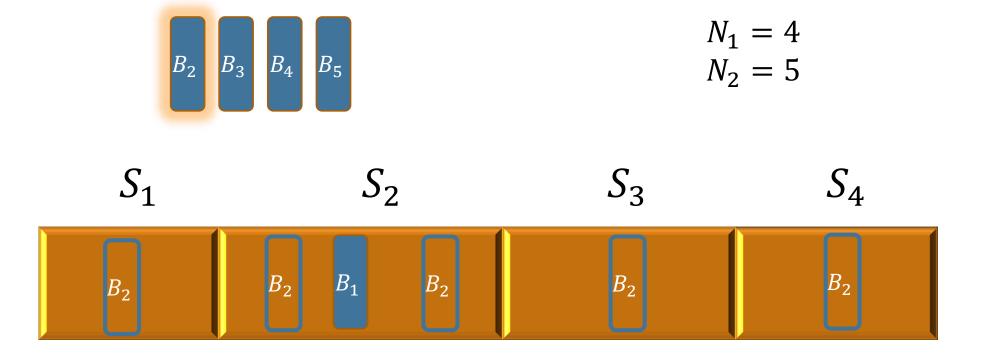
$$N_1 = 4$$
$$N_2 = 5$$



$$m = 5$$
 books $B_1, B_2, B_3, ..., B_m$
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Can use the Product rule

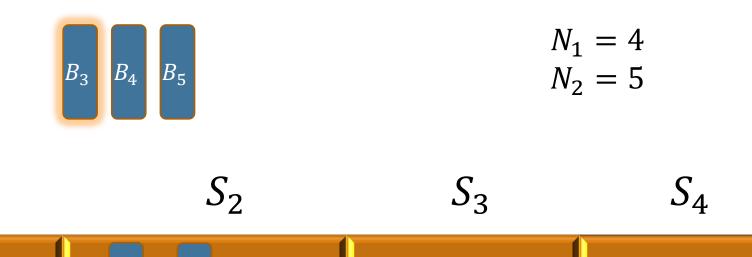
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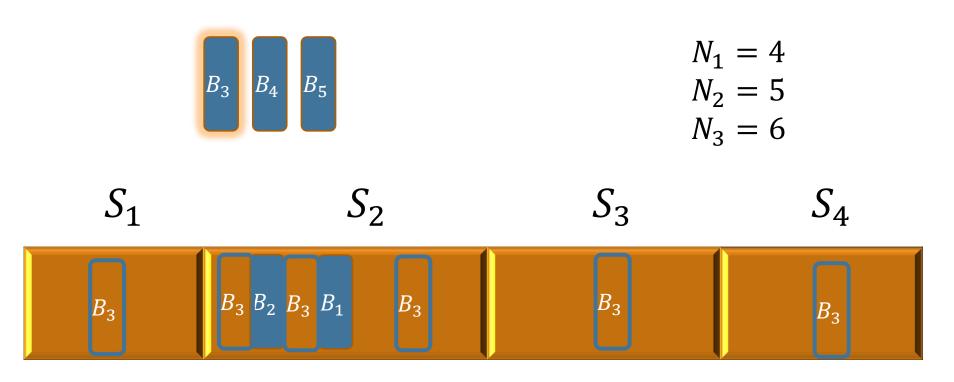
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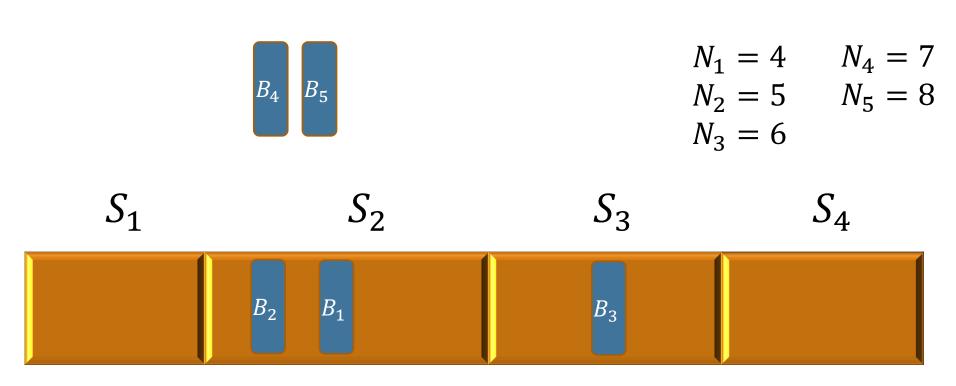
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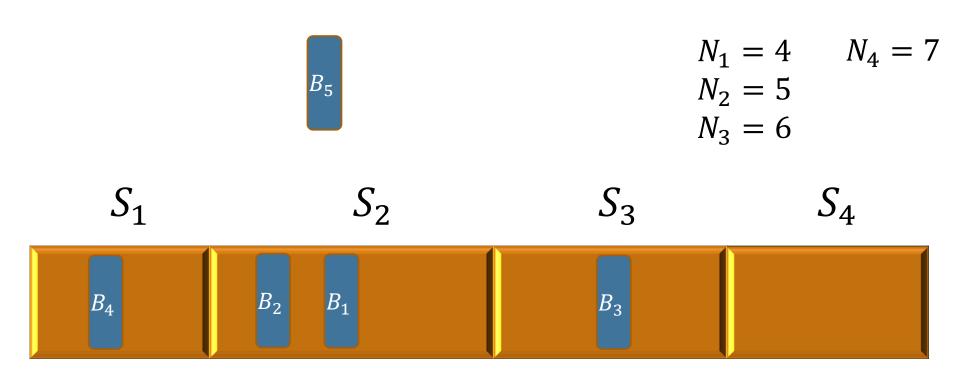
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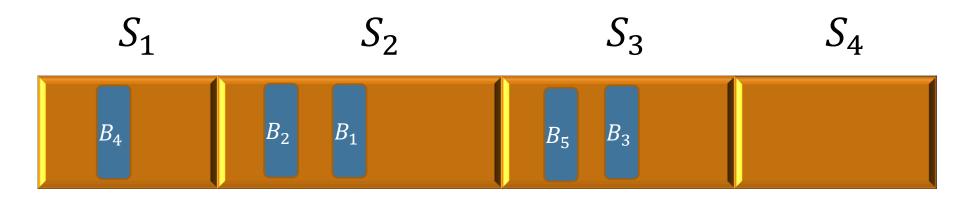
Procedure: Start with empty shelves, place books one by one



$$m = 5$$
 books $B_1, B_2, B_3, ..., B_m$
 $n = 4$ shelves $S_1, S_2, S_3, ..., S_n$

Using the product rule, the number of ways we can place 5 books on 4 shelves: $N_1 \cdot N_2 \cdot N_3 \cdot N_4 \cdot N_5$

$$N_1 = 4$$
 $N_4 = 7$
 $N_2 = 5$ $N_5 = 8$
 $N_3 = 6$



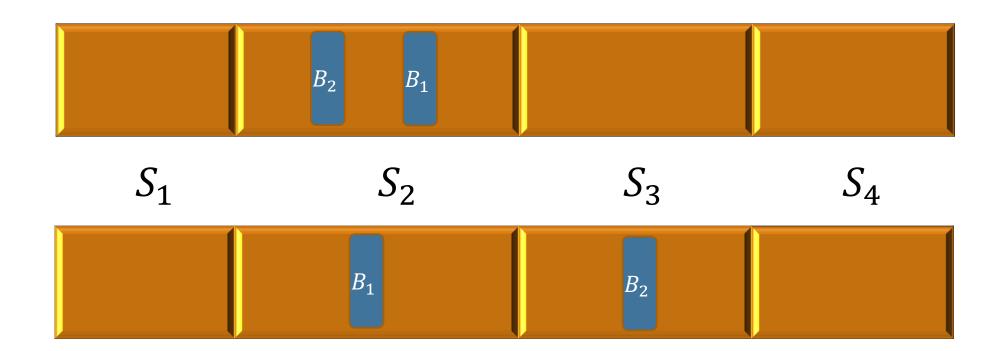
$$m = 5$$
 books $B_1, B_2, B_3, ..., B_m$
 $n = 4$ shelves $S_1, S_2, S_3, ..., S_n$

 B_3

What if we placed them differently? Where can B_3 go?

$$N_1 = 4$$

 $N_2 = 5$

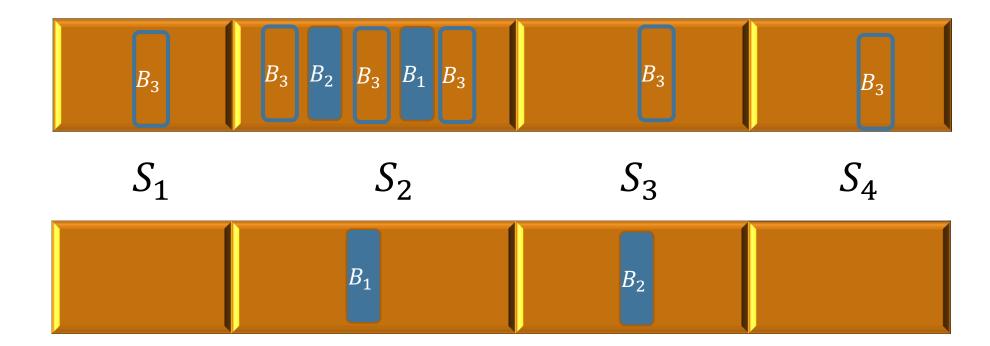


 $m = 5 \text{ books } B_1, B_2, B_3, ..., B_m$ $n = 4 \text{ shelves } S_1, S_2, S_3, ..., S_n$

 B_3

What if we placed them differently? Where can B_3 go?

 $N_1 = 4$ $N_2 = 5$

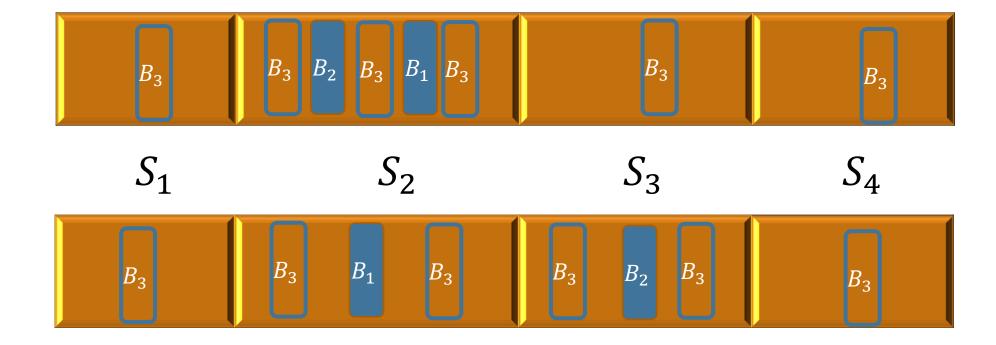


$$m = 5$$
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 B_3

What if we placed them differently? Where can B_3 go?

 $N_1 = 4$ $N_2 = 5$



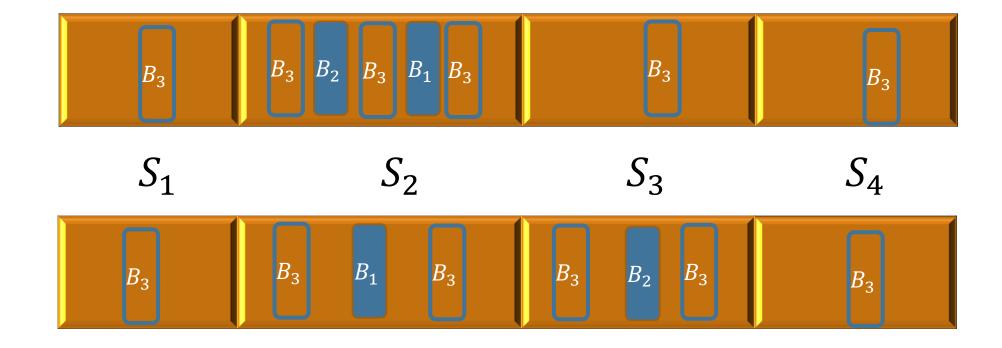
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 B_3

What if we placed them differently? Where can B_3 go?

$$N_1 = 4$$
$$N_2 = 5$$

$$N_3 = 6$$

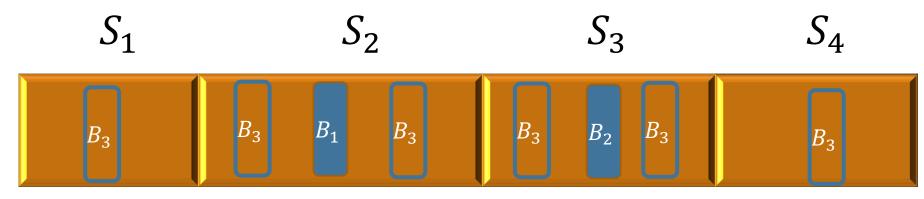


$$m = 5$$
 books $B_1, B_2, B_3, ..., B_m$
 $n = 4$ shelves $S_1, S_2, S_3, ..., S_n$

What if we placed them differently? Where can B_3 go?

Every time we place a book, we split an existing "slot" into 2 "slots", regardless of where we place the book.

Each book becomes a new "divider". So where we place B_i during N_i does not affect the value of N_{i+1}



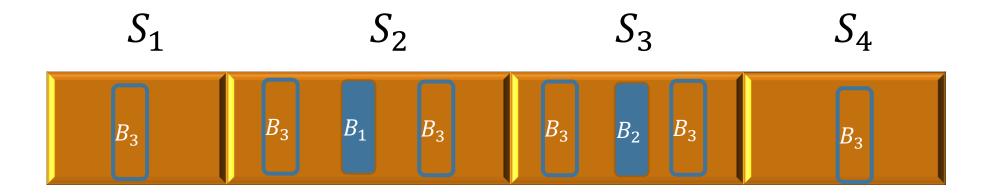
m books $B_1, B_2, B_3, \dots, B_m$ n shelves $S_1, S_2, S_3, \dots, S_n$

Task i: Place B_i

- 1. Place B_i on a shelf on the far left: n choices
- 2. Place B_i immediately to the right of a book $B_1, ..., B_{i-1}$ already on the shelf: i-1 choices.

Thus
$$N_i = n + i - 1$$

Therefore the total number of ways to place books on shelves is $N_1 \cdot N_2 \cdot ... \cdot N_m$



$$m$$
 books $B_1, B_2, B_3, \dots, B_m$
 n shelves $S_1, S_2, S_3, \dots, S_n$

$$N_1 = n$$

$$N_2 = n + 1$$

$$N_3 = n + 2$$

...

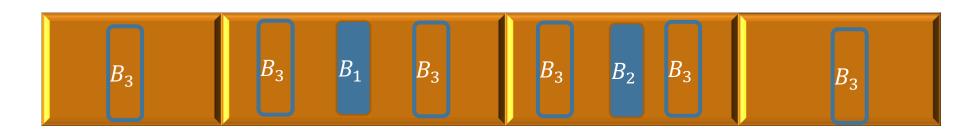
$$N_{m-1} = n + m - 2$$

$$N_m = n + m - 1$$

Therefore the total number of ways to place books on shelves is

$$N_1 \cdot N_2 \cdot \dots \cdot N_m$$

= $n \cdot (n+1) \cdot (n+2) \cdot \dots \cdot (n+m-2) \cdot (n+m-1)$



$$m$$
 books $B_1, B_2, B_3, \dots, B_m$
 n shelves $S_1, S_2, S_3, \dots, S_n$

Again we want to express $n \cdot (n+1) \cdot (n+2) \cdot ... \cdot (n+m-2) \cdot (n+m-1)$ as a factorial

$$(n+m-1)! = 1 \cdot 2 \cdot ... \cdot (n-1) \cdot n \cdot (n+1) \cdot ... \cdot (n+m-2) \cdot (n+m-1)$$
to
$$n \cdot (n+1) \cdot ... \cdot (n+m-2) \cdot (n+m-1)$$

Divide by the first part:

$$\frac{1 \cdot 2 \cdot \dots \cdot (n-1) \cdot n \cdot (n+1) \cdot \dots \cdot (n+m-2) \cdot (n+m-1)}{1 \cdot 2 \cdot \dots \cdot (n-1)} = \frac{(n+m-1)!}{(n-1)!}$$

$$m$$
 books $B_1, B_2, B_3, \dots, B_m$
 n shelves $S_1, S_2, S_3, \dots, S_n$

Therefore there are $\frac{(n+m-1)!}{(n-1)!}$ ways to place m books on n shelves

$$\frac{(n+m-1)!}{(n-1)!}$$

$$=\frac{(4+5-1)!}{(4-1)!}$$

$$S_1$$
 S_2 $=\frac{(8)!}{(3)!} = 6720$ S_3 S_4

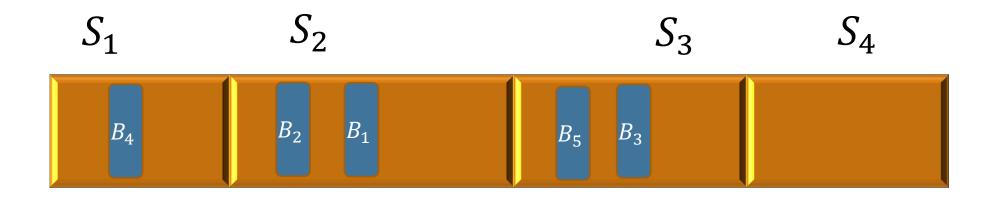
Review of Product Rule

Express as a procedure

Where a procedure is a **sequence of tasks**

The choice we make for Task i does not affect the number of ways to do Task i+1

Then we can apply the product rule



Set $S = \{a, b, c, d, e\}$ How many subsets? $\{a\}, \{a, b, c\}, \{b, c\}, \{a, b, c, d, e\}, \emptyset$, etc.

 $\emptyset \subseteq S \text{ if } \forall x \in \emptyset, x \in S$ $\neg \forall x \in \emptyset, x \in S$ $\exists x \in \emptyset, x \notin S$ False

Contradiction

Prove by contradiction, thus assume the negation

Set
$$S = \{a_1, a_2, a_3, ..., a_n\}$$

How many subsets?

Can we use the Product rule?

Procedure: specify a subset

for
$$i = 1, 2, 3, ..., n$$
:

Task i: specify whether or not a_i is in the subset

Note a_i is either there or it is not – two choices

 $N_i = 2$, and does not depend on first i - 1 tasks number of subsets = 2^n (even if n = 0)

Set
$$S = \{a, b, c, d, e\}$$

Very similar to bitstrings.

If a_i is in the subset, we write a 1 if a_i is not in the subset we write 0

Subsets of $S \leftrightarrow \text{bitstrings of length } |S|$

$$\{a\} \leftrightarrow 10000$$

$$\{a, b, d, e\} \leftrightarrow 11011$$

$$\emptyset \leftrightarrow 00000$$

$$\{a, b, c, d, e\} \leftrightarrow 11111$$

$$Set S = \{a, b, c, d, e\}$$

Can also go from a bitstring to a subset

If we read a 1 then a_i is in the subset If we read a 0 then a_i is not in the subset

bitstrings of length $\leftrightarrow |S|$

$$10000 \leftrightarrow \{a\}$$

$$11011 \leftrightarrow \{a, b, d, e\}$$

$$00000 \leftrightarrow \emptyset$$

Bijection Rule

Set
$$S = \{a, b, c, d, e\}$$

Can also go from a bitstring to a subset

That means these are basically the same problem

(What it also means is we can apply the same solution)