**MP3**

**Team: Cache Money Records (Young Moolah Baby)**

**Checkpoint 1 Report**

**Progress**

This is the first checkpoint of MP3, so it’s relatively straightforward. We used two wishbone modules to connect between our CPU datapath (Fetch and Memory stage respectively) and magic memory. For datapath, it’s similar to MP1 design except that we added pipeline registers after each stage for pipelining purpose. We created the components that needed for all instructions expect for STI/LDI. Another change is that instead of having a state machine to control the control signals in each state, we used control words for different opcodes. In this checkpoint, we didn’t consider data/control hazards situations.

Functionalities implemented and tested: ADD, AND, NOT, STR, LDR, BR.

We wrote test programs to test each instruction by itself, then we tested them together using the provided cp1 test code.

**Contributions**

We all coded the checkpoint together and debugged together

**Roadmap**

Ryan: Implement all the remaining instructions

Subhash: Cach design/data forwarding paper design

Chaohua: Cach interconnect

**Cache design**

ICache and DCache have wishbone master port and they both need to connect to L2 cache, sharing one slave port. So we use mux/demux to select data coming from or send to icache and dcache according to control signals.

If it’s an instruction miss (icache read/write), L2 communicates with icache, L1L2 cache wishbone connecting to icache wishbone. If it’s a data miss (dcache read/write), L1L2 cache wishbone connects to dcache. If address coming from icache and dcache is same with L1L2 cache’s address, then physical memory also communicates with icache and dcache which means if it’s a icache/dcache miss and also L2 cache miss, icache/dcache can get data from physical memory directly.

Below is the paper design for interconnect and cache. For icache, dcache and L2 cache, We used the same cache design from mp2.



