## Operators [*H2010* 3]

 Expressions can built up as expected in many programming languages

```
3 x x+y (x<=y) a + c * d - (e * (a / b))
```

- ➤ Some operators are left-associative like + \* /: a + b + c parses as (a + b) + c
- ➤ Some operators are right-associative like : . ^ && ||: a:b:c:[] parses as a:(b:(c:[]))
- Other operators are non-associative like == /= < <= >= >: a <= b <= c is illegal, but (a <= b) && (b <= c) is ok.</p>
- ► The minus sign is tricky: e f parses as "e subtract f", (- f) parses as "minus f", but e (- f) parses as "function e applied to argument minus f"

## Prelude extracts (I)

► Infix declarations

```
infixr 9 .
infixr 8 ^, ^^, ..
infixl 7 *, /, 'quot', 'rem', 'div', 'mod'
infixl 6 +, -
infixr 5 :, ++
infix 4 ==, /=, <, <=, >=, >
infixr 3 &&
infixr 2 ||
infixl 1 >>, >>=
infixr 1 =<<
infixr 0 $, $!, 'seq'</pre>
```

Higher precedence numbers bind tighter. Function application binds tightest of all

## The Haskell Prelude [H2010 9]

- ► The "Standard Prelude" is a library of functions loaded automatically (by default) into any Haskell program.
- ► Contains most commonly used datatypes and functions
- ► [H2010 9] is a *specification* of the Prelude the actual code is compiler dependent

# Prelude extracts (II)

Numeric Functions

```
subtract :: (Num a) => a -> a -> a
even, odd :: (Integral a) => a -> Bool
gcd :: (Integral a) => a -> a
lcm :: (Integral a) => a -> a -> a
(^) :: (Num a, Integral b) => a -> b -> a
(^^) :: (Fractional a, Integral b) => a -> b -> a
```

The Num, Integral and Fractional annotations have to do with *type-classes* — see later.

## Prelude extracts (III)

► Boolean Type & Functions

```
data Bool = False | True
(&&), (||) :: Bool -> Bool -> Bool
not :: Bool -> Bool
otherwise :: Bool
```

## Prelude extracts (V)

► Function Functions

```
id :: a -> a
const :: a -> b -> a
(.) :: (b -> c) -> (a -> b) -> a -> c
flip :: (a -> b -> c) -> b -> a -> c
seq :: a -> b -> b
($), ($!) :: (a -> b) -> a -> b
```

We will re-visit these later — note that type-polymorphism here means that the possible implementations of some of these are extremely constrained!

## Prelude extracts (IV)

▶ List Functions

```
map :: (a -> b) -> [a] -> [b]
(++) :: [a] -> [a] -> [a]
filter :: (a -> Bool) -> [a] -> [a]
concat :: [[a]] -> [a]
head :: [a] -> a
tail :: [a] -> [a]
null :: [a] -> Bool
length :: [a] -> Int
(!!) :: [a] -> Int -> a
repeat :: a -> [a]
take :: Int -> [a] -> [a]
elem :: Eq a => a -> [a] -> Bool
```

# Lists [H2010 3.7,3.10]

- ► Fundamentally lists are built from "nil" ([]) and "cons" (:)
- We use square brackets to provide syntactical sugar in a variety of ways
  - Enumeration:
    [a,b,c,d] for a:b:c:d:[]
    Ranges:
    [4..9] for [4,5,6,7,8,9]
    also [4,7..20] for [4,7,10,13,16,19]
  - ► Comprehension:

    [ x\*x | x <- [1..10], even x] for [4,16,36,64,100]

    Comprehensions are more complex than this (see later, or [H2010 3.11])
- ► Strings are a special notation of lists of characters "Hello" for ['H','e','l','l','o']

#### Function: head

head xs returns the first element of xs, if non-empty

Type Signature

```
head :: [a] -> a
```

Non-Empty List

```
head (x:_) = x
```

**Empty List** 

```
head [] = error "Prelude.head: empty list"
```

We have to fail in the last case because there is no way to generate a value  $\mathbf{v}$  of type  $\mathbf{a}$ , where  $\mathbf{a}$  can be any possible type, if there is no such value input to the function.

Empty list [], of type a, contains no value of type a!

### Why Not define a default value for head []?

► Why don't we define a default value for each type (default ::a) so that we can define (possible using Haskell classes):

```
head [] = default -- for any given type a
```

rather than having head  $[] = \bot$ ?

- ▶ Why not have default for Int equal to 0?
- ► A key design principle behind Haskell libraries and programs is to have programs (functions!) that obey nice obvious laws:

```
xs = head xs : tail xs
sum (xs ++ ys) = sum xs + sum ys
product (xs ++ ys) = product xs * product ys
```

► Consider the product law if default = 0 and xs = [], and assume that both sum and product use default for the empty list case. Lefthand side is then product ys while thre righthand side is zero.

#### Undefinedness in Haskell

- ► Sometimes a Haskell function is *partial*: it doesn't return a value for some input, because it can't without violating type restrictions.
- ► Haskell provides two ways to explicitly define such a undefined "value":

```
undefined :: a
error :: String -> a
```

Evaluating either of these results in a run-time error

- ▶ There are two ways in which "undefined" can occur implicitly:
  - ▶ If we use pattern-matching that is incomplete, so that some input values fail to match.
  - if a recursive function fails to terminate
- ▶ When talking about the meaning of Haskell, it is traditional to use the symbol  $\bot$ , a.k.a. "bottom", to denote undefinedness.

## Function: tail

tail xs, for non-empty xs returns it with first element removed Type Signature

```
tail :: [a] -> [a]
```

Non-Empty List

```
tail (_:xs) = xs
```

**Empty List** 

```
tail [] = error "Prelude.tail: empty list"
```

Here again, we have tail  $[] = \bot$ .

```
tail [] /= [] — Why Not?
```

- ▶ Why don't we define tail [] = []? The typing allows it.
- ► Re-consider the following law, given xs = tail []

```
xs = head xs : tail xs
```

```
tail [] = head ( tail[]) : (tail : tail [])
= [] = head [] : tail []
= [] = ⊥ : []
```

We have managed to show that the empty list is the same as a singleton list containing an undefined element.

▶ "Obvious" fixes can have unexpected consequences.

### Function: init

init xs, for non-empty xs returns it with last element removed Type Signature

```
init :: [a] -> [a]
```

Singleton List

```
init[x] = []
```

Non-Empty List

```
init (x:xs) = x : init xs
```

**Empty List** 

```
init [] = error "Prelude.init: empty list"
```

