



**Trinity College Dublin**

Coláiste na Tríonóide, Baile Átha Cliath

The University of Dublin

## 02 – Arrays

**CS1022 – Introduction to Computing II**

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Nothing very new here ... we have already been using arrays!

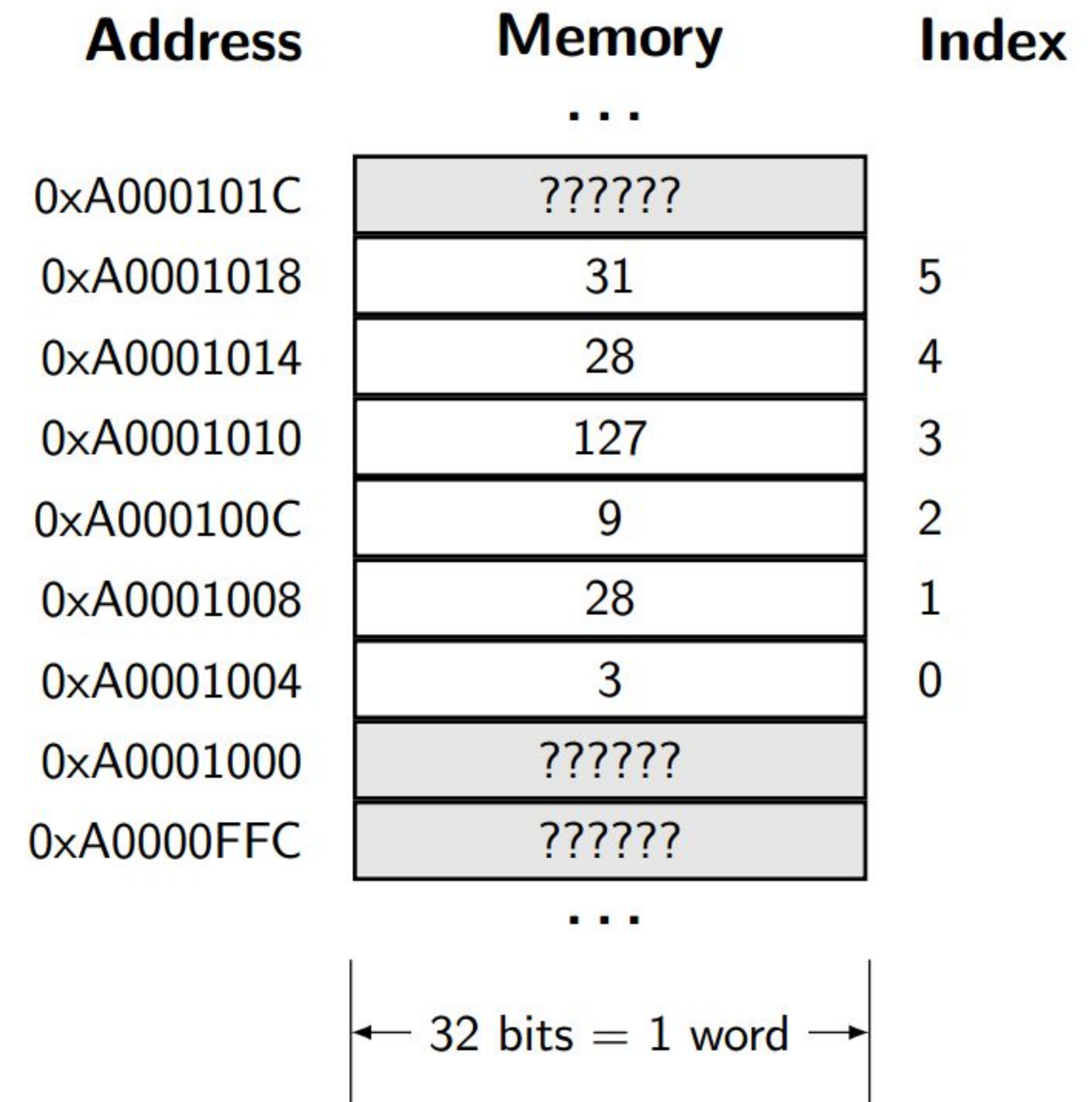
Array – an ordered collection of elements stored sequentially in memory

e.g. integers, ASCII characters, lottery numbers

Homogeneous elements?

(at least with respect to size)

Dimension: number of elements in array



Efficient access at a specific index is an important feature of arrays

Referred to as “random access”

Example: retrieve the 4th element (index=3) of an array of words

Step 1: translate array index into byte offset from start address of array in memory

$$\text{<byte offset>} = \text{<index>} \times \text{<elem size>}$$

Step 2: add byte offset to array base address to access element

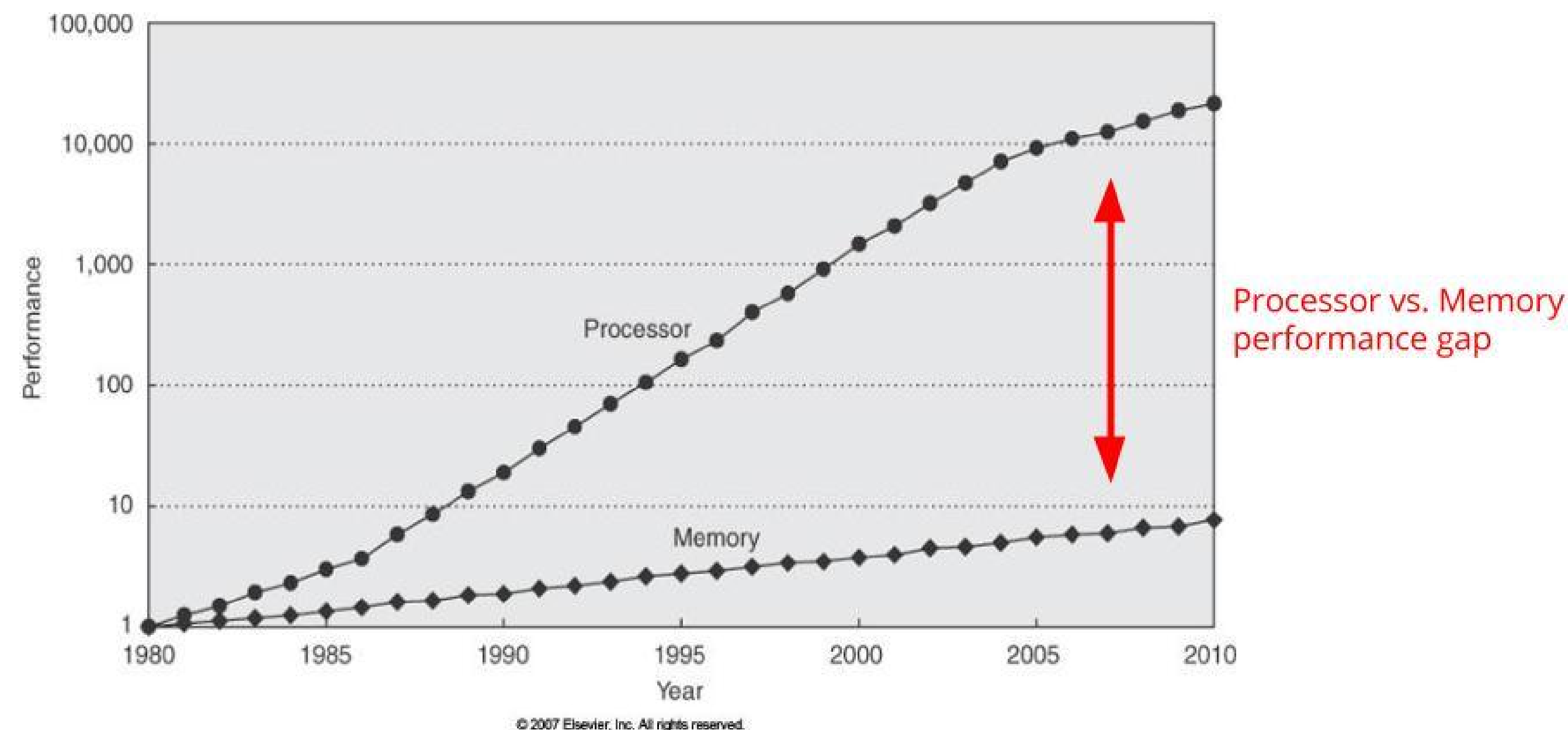
$$\text{<address>} = \text{<array start address>} + \text{<byte offset>}$$

Efficient random access using Scaled Register Offset addressing mode:

```
LDR r4, =array           ; pArr = start address of array
LDR r5, =3               ; index = 3 (4th element)
LDR r6, [r4, r5, LSL #2] ; elem = Memory.Word[pArr + (index * 4)]
```

Generally CPU speed increases much faster than memory access speed

Relatively speaking, memory is getting slower\*



Memory speed lags behind CPU speed

Déjà Vu!

```
LDR R1, =myArray ; start address of myArray
LDR R0, =0 ; sum = 0
LDR R4, =0 ; count = 0

whSum CMP R4, #10 ; while (count < 10)
      BHS eWhSum ; {
      LDR R6, [R1, R4, LSL #2] ; num = myArray[count]
      ADD R0, R0, R6 ; sum = sum + num
      ADD R4, R4, #1 ; count = count + 1
      B whSum ; }
eWhSum ;
```

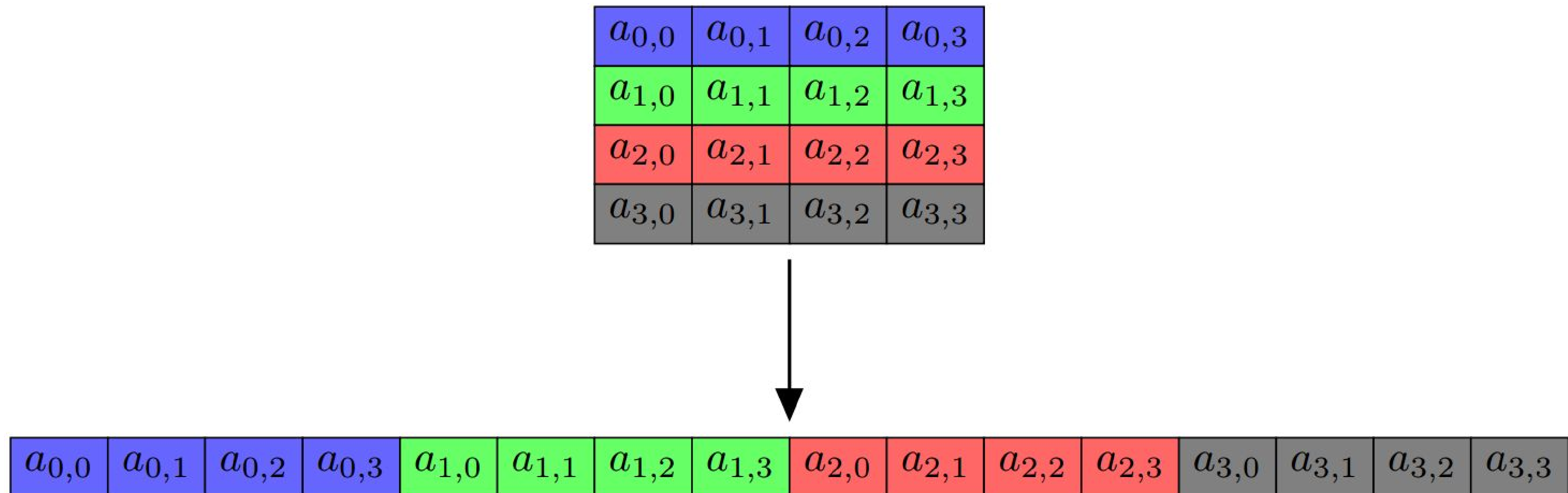
The pseudo-code comments have changed but the program is identical (See **Addressing Modes**)



Arrays can have more than one dimension

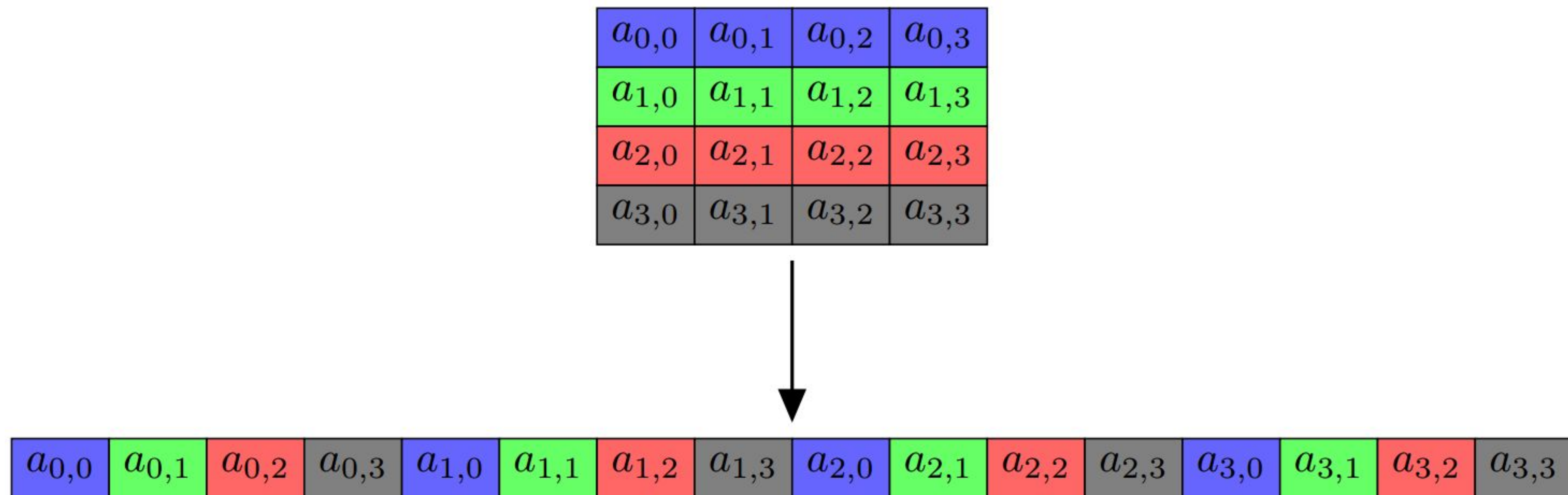
e.g. a two-dimensional array – analogous to a table containing elements arranged in rows and columns

Stored in memory by mapping the 2D array into 1D memory, e.g.



Row-major order: 2D array is stored in memory by storing each row contiguously in memory

Column-major order: 2D array is stored in memory by storing each column contiguously in memory (in image)



## 2D array declared in memory

```
AREA TestData, DATA, READWRITE
```

```
col_size EQU 6      ; just for convenience, not required
row_size EQU 8      ; just for convenience, not required
```

```
array DCD 6, 3, 8, 2, 5, 2, 9, 1 ; row 0
      DCD 3, 7, 2, 8, 5, 7, 2, 7 ; row 1
      DCD 2, 4, 7, 4, 2, 6, 7, 4 ; row 2
      DCD 1, 9, 3, 2, 9, 5, 6, 8 ; row 3
      DCD 7, 5, 3, 7, 5, 8, 2, 1 ; row 4
      DCD 6, 4, 8, 9, 0, 3, 2, 5 ; row 5
```

## ... or equivalently ...

```
AREA TestData, DATA, READWRITE
```

```
col_size EQU 6      ; just for convenience, not required
row_size EQU 8      ; just for convenience, not required
```

```
array DCD 6, 3, 8, 2, 5, 2, 9, 1, 3, 7, 2, 8, 5, 7, 2, 7, 2, 4, 7, 4, 2
      DCD 6, 7, 4, 1, 9, 3, 2, 9, 5, 6, 8, 7, 5, 3, 7, 5, 8, 2, 1, 6, 4
      DCD 8, 9, 0, 3, 2, 5
```



Example: retrieve the element at the 4th row and 3rd column of a 2D array of words with 6 rows and 8 columns – `array[3][2]`

Step 1: translate 2D array index into 1D array index

$$\text{<index>} = (\text{<row>} \times \text{<row size>}) + \text{<col>}$$

Step 2: translate 1D array index into byte offset from start address of array in memory

$$\text{<byte offset>} = \text{<index>} \times \text{<elem size>}$$

Step 3: add byte offset to array base address to access element

$$\text{<address>} = \text{<array base address>} + \text{<byte offset>}$$

Example: retrieve the element at the 4th row and 3rd column of a 2D array of words with 6 rows and 8 columns – array[3][2]

```
LDR r4, =array ; pArr = address of array start
```

```
LDR r5, =col_size; load col_size
```

```
LDR r6, =row_size; load row_size
```

```
; looking for array[3][2] (4th row, 3rd column)
```

```
LDR r1, =3 ; row = 3
```

```
LDR r2, =2 ; col = 4
```

```
; <byte offset> = ((row * <row_size>) + col) * <elem size>
```

```
MUL r7, r1, r6 ; index = row * row_size
```

```
ADD r7, r7, r2 ; index = index + col
```

```
LDR r0, [r4, r7, LSL #2] ; elem = Memory.Word[ pArr + (index*4) ]
```

e.g. a 3D array of size  $sz \times sy \times sx$

In general, the index of element  $a[z][y][x]$  is:

$$\text{index} = ((z \times sy \times sx) + (y \times sx) + x)$$

e.g. a 4D array of size  $sz \times sy \times sx \times sw$

In general, the index of element  $a[z][y][x][w]$  is:

$$\text{index} = ((z \times sy \times sx \times sw) + (y \times sx \times sw) + (x \times sw) + w)$$

Warning: an array of bytes with odd dimensions may cause the data following the array to be odd aligned, requiring padding of one byte (similarly for half-words)

Occasionally you will see the following warning when working with memory (particularly later on in the course)

“warning: A1581W: Added n bytes of padding at address 0xn timer”

The following demonstrates it

	Area	padding,	DATA,	READWRITE
chars	DCB	“abcde”		
nums	DCD	-1, 122, 13		

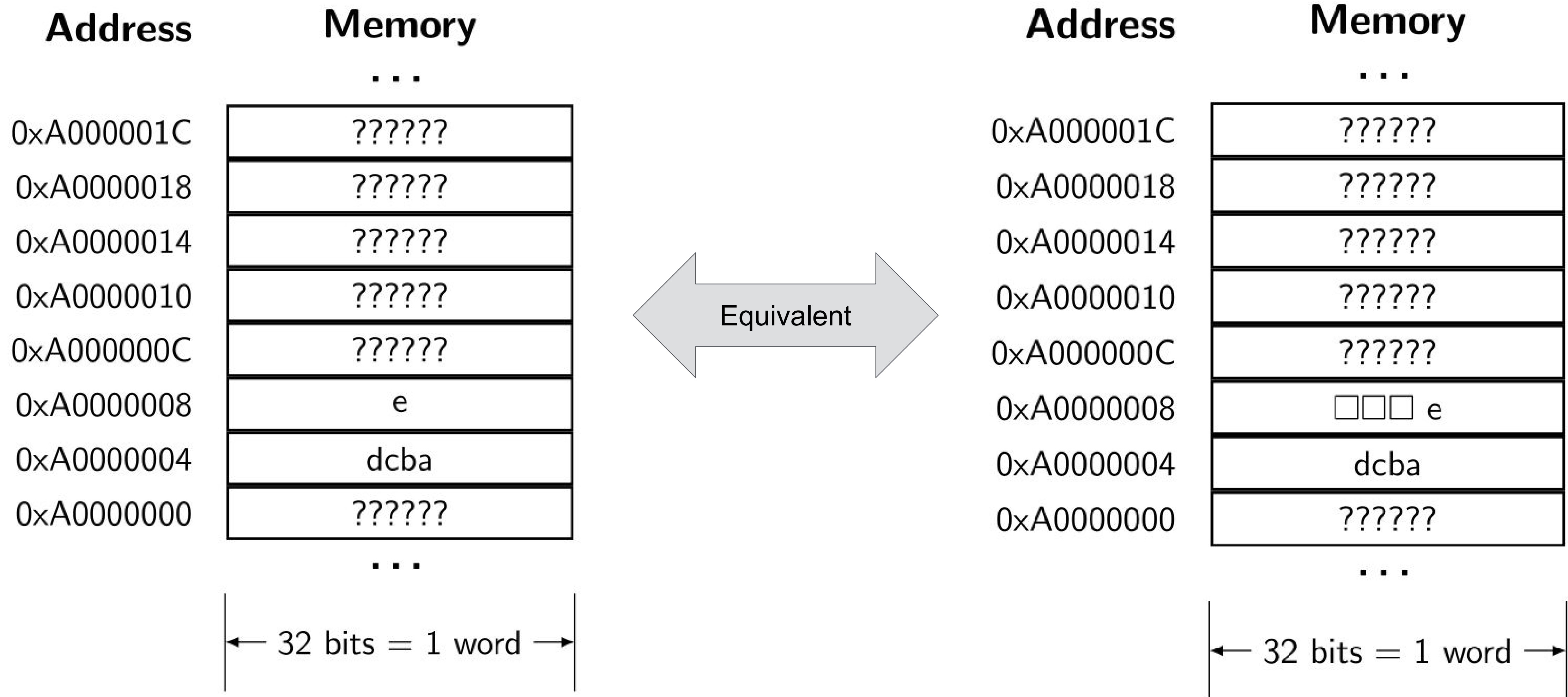
Loading just “chars” into memory for now

```

Area      padding,   DATA,   READWRITE
chars    DCB    "abcde"
;comment for now;nums    DCD    -1, 122, 13
    
```

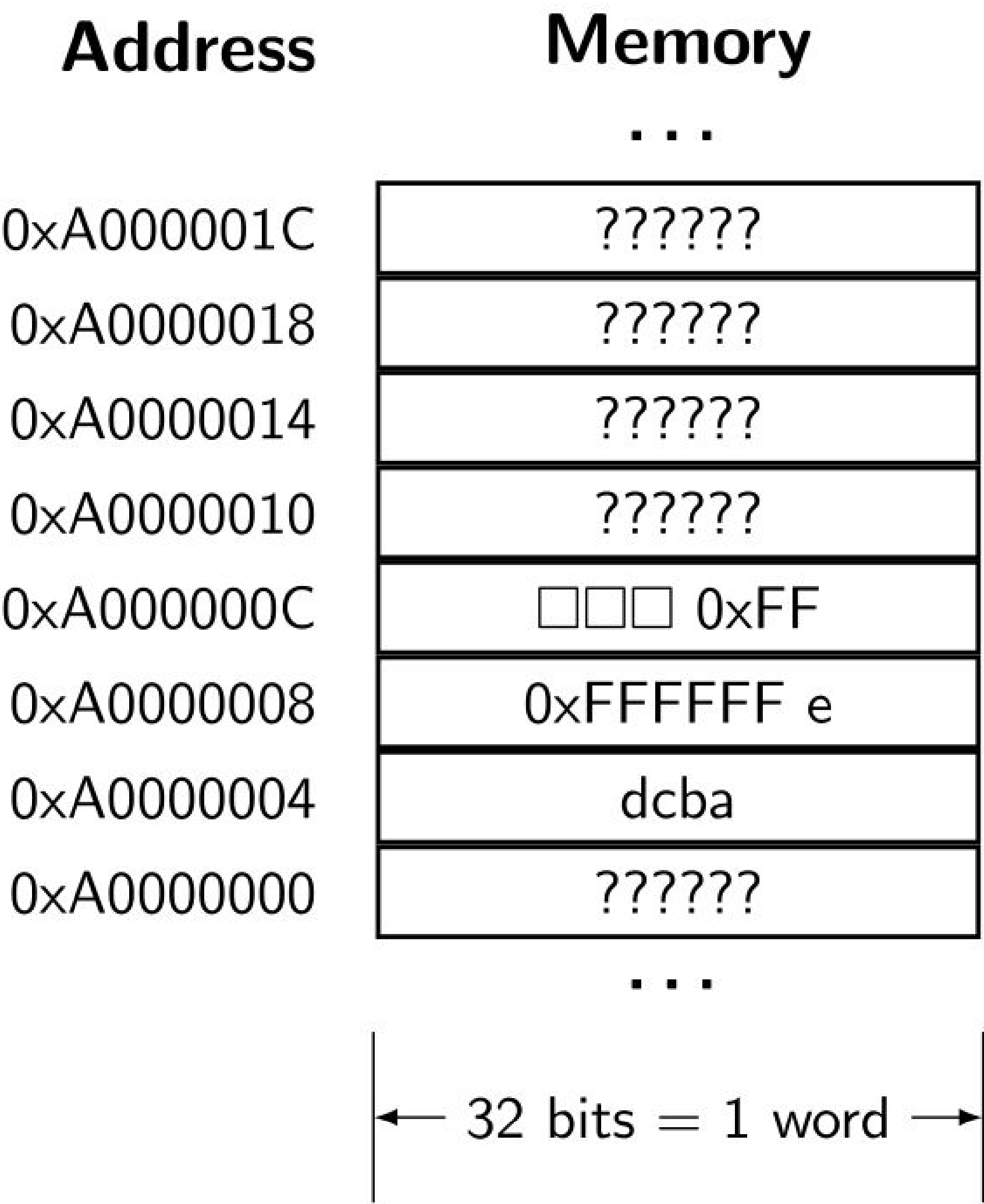
Address	Memory
	...
0xA000001C	??????
0xA0000018	??????
0xA0000014	??????
0xA0000010	??????
0xA000000C	??????
0xA0000008	e
0xA0000004	dcba
0xA0000000	??????
	...
← 32 bits = 1 word →	





Now store the numbers (not what actually happens)

	Area	padding,	DATA,	READWRITE
chars	DCB	"abcde"		
nums	DCD	-1;comment	for now;	122, 13



Now store the numbers (not what actually happens)

	Area	padding,	DATA,	READWRITE
chars	DCB	"abcde"		
nums	DCD	-1, 122, 13		

