

For return on 24 January 2014 (late submission: 7 February 2014)

Electronic submission: pdf files only

1. (3%) Construct the truth-table for the Boolean function given by the Boolean formula

$$\neg(A \wedge (B \rightarrow C)) \wedge \neg B.$$

Use the truth-table to realise this function by a formula with the connectives  $\neg$ ,  $\vee$ ,  $\wedge$  only. Simplify the formula in such a way that the corresponding Boolean circuit contains a minimal number of gates. Show the Boolean circuit.

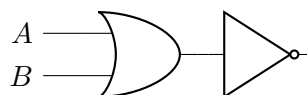
### Truth Table

| $A$ | $B$ | $C$ | $\neg$ | $(A \wedge (B \rightarrow C))$ | $\wedge$ | $\neg B$ |
|-----|-----|-----|--------|--------------------------------|----------|----------|
| 0   | 0   | 0   | 1      | 0                              | 1        | 1        |
| 0   | 0   | 1   | 1      | 0                              | 1        | 1        |
| 0   | 1   | 0   | 1      | 0                              | 0        | 0        |
| 0   | 1   | 1   | 1      | 0                              | 0        | 0        |
| 1   | 0   | 0   | 0      | 1                              | 0        | 1        |
| 1   | 0   | 1   | 0      | 1                              | 0        | 1        |
| 1   | 1   | 0   | 1      | 0                              | 0        | 0        |
| 1   | 1   | 1   | 0      | 1                              | 0        | 0        |

### Simplification

1.  $(\neg A \wedge \neg B \wedge \neg C) \vee (\neg A \wedge \neg B \wedge C)$
2.  $(\neg(A \vee B) \wedge \neg C) \vee (\neg(A \vee B) \wedge C)$
3.  $\neg(A \vee B) \vee (\neg C \wedge C)$
4.  $\neg(A \vee B) \vee 0$
5.  $\neg(A \vee B)$

### Boolean Circuit



2. (3%) Are the following Boolean formulas equivalent? Explain your answer.

(a)  $A \rightarrow (B \wedge C)$  and  $(A \rightarrow B) \wedge (A \rightarrow C)$

| $A$ | $B$ | $C$ | $A$ | $\rightarrow$ | $(B \wedge C)$ |
|-----|-----|-----|-----|---------------|----------------|
| 0   | 0   | 0   |     | 1             | 0              |
| 0   | 0   | 1   |     | 1             | 0              |
| 0   | 1   | 0   |     | 1             | 0              |
| 0   | 1   | 1   |     | 1             | 1              |
| 1   | 0   | 0   |     | 0             | 0              |
| 1   | 0   | 1   |     | 0             | 0              |
| 1   | 1   | 0   |     | 0             | 0              |
| 1   | 1   | 1   |     | 1             | 1              |

| $A$ | $B$ | $C$ | $(A \rightarrow B)$ | $\wedge$ | $(A \rightarrow C)$ |
|-----|-----|-----|---------------------|----------|---------------------|
| 0   | 0   | 0   | 1                   | 1        | 1                   |
| 0   | 0   | 1   | 1                   | 1        | 1                   |
| 0   | 1   | 0   | 1                   | 1        | 1                   |
| 0   | 1   | 1   | 1                   | 1        | 1                   |
| 1   | 0   | 0   | 0                   | 0        | 0                   |
| 1   | 0   | 1   | 0                   | 0        | 1                   |
| 1   | 1   | 0   | 1                   | 0        | 0                   |
| 1   | 1   | 1   | 1                   | 1        | 1                   |

Yes they are equivalent

(b)  $(A \wedge B) \rightarrow C$  and  $(\neg C \rightarrow \neg A) \wedge (\neg C \rightarrow \neg B)$

| $A$ | $B$ | $C$ | $(A \wedge B)$ | $\rightarrow$ | $C$ |
|-----|-----|-----|----------------|---------------|-----|
| 0   | 0   | 0   | 0              | 1             |     |
| 0   | 0   | 1   | 0              | 1             |     |
| 0   | 1   | 0   | 0              | 1             |     |
| 0   | 1   | 1   | 0              | 1             |     |
| 1   | 0   | 0   | 0              | 1             |     |
| 1   | 0   | 1   | 0              | 1             |     |
| 1   | 1   | 0   | 1              | 0             |     |
| 1   | 1   | 1   | 1              | 1             |     |

| $A$ | $B$ | $C$ | $(\neg C \rightarrow \neg A)$ | $\wedge$ | $(\neg C \rightarrow \neg B)$ |
|-----|-----|-----|-------------------------------|----------|-------------------------------|
| 0   | 0   | 0   | 1                             | 1        | 1                             |
| 0   | 0   | 1   | 0                             | 1        | 0                             |
| 0   | 1   | 0   | 1                             | 0        | 1                             |
| 0   | 1   | 1   | 0                             | 1        | 0                             |
| 1   | 0   | 0   | 1                             | 0        | 1                             |
| 1   | 0   | 1   | 0                             | 1        | 0                             |
| 1   | 1   | 0   | 1                             | 0        | 0                             |
| 1   | 1   | 1   | 0                             | 1        | 0                             |

No they are not equivalent

(c)  $(A \vee B) \rightarrow C$  and  $(A \rightarrow B) \vee (A \rightarrow C)$

| $A$ | $B$ | $C$ | $(A \vee B)$ | $\rightarrow$ | $C$ |
|-----|-----|-----|--------------|---------------|-----|
| 0   | 0   | 0   | 0            | 1             |     |
| 0   | 0   | 1   | 0            | 1             |     |
| 0   | 1   | 0   | 1            | 0             |     |
| 0   | 1   | 1   | 1            | 1             |     |
| 1   | 0   | 0   | 1            | 0             |     |
| 1   | 0   | 1   | 1            | 1             |     |
| 1   | 1   | 0   | 1            | 0             |     |
| 1   | 1   | 1   | 1            | 1             |     |

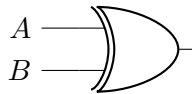
  

| $A$ | $B$ | $C$ | $(A \rightarrow B)$ | $\vee$ | $(A \rightarrow C)$ |
|-----|-----|-----|---------------------|--------|---------------------|
| 0   | 0   | 0   | 1                   | 1      | 1                   |
| 0   | 0   | 1   | 1                   | 1      | 1                   |
| 0   | 1   | 0   | 1                   | 1      | 1                   |
| 0   | 1   | 1   | 1                   | 1      | 1                   |
| 1   | 0   | 0   | 0                   | 0      | 0                   |
| 1   | 0   | 1   | 0                   | 1      | 1                   |
| 1   | 1   | 0   | 1                   | 1      | 0                   |
| 1   | 1   | 1   | 1                   | 1      | 1                   |

No they are not equivalent

3. **(3%)** A parity function is a Boolean function whose value is 1 if the input has an odd number of ones. Design a Boolean circuit for the 2-bit parity function. Show your working. (Hint: you may find XOR gates useful.)

| $A$ | $B$ |   |
|-----|-----|---|
| 0   | 0   | 0 |
| 0   | 1   | 1 |
| 1   | 0   | 1 |
| 1   | 1   | 0 |



4. **(3%)** Suppose  $\alpha = a_{31}a_{30}\dots a_1a_0$  is a 32-bit binary word. Consider the 32-bit binary word  $\beta = b_{31}b_{30}\dots b_1b_0$  computed by the following algorithm: scan  $\alpha$  from right to left and copy its bits to  $\beta$  until the first 1 is found (which is also copied to  $\beta$ ); after that, copy the Boolean negations of the bits in  $\alpha$ . For example,  $\alpha = 10100\dots 00$  is transformed to  $\beta = 01100\dots 00$ . Explain what this algorithm computes if  $\alpha$  and  $\beta$  are interpreted as binary numbers.

Given a two's complement binary number, the algorithm computes the conversion into a binary number.

5. (6%) Given the machine 32-bit word

1100 0001 0011 0000 0000 0000 0000

find the decimal number represented by this word assuming that it is

(a) a two's complement integer;

1. Read the first character from the left of the word to get the sign of the integer it represents.  
In this case it is a one so the number is negative.
2. Use the two's complement algorithm to convert it into a binary number. Starting from the right of the word, move to the first one :  
\* 1100 0001 0011 00...00
3. Move left, from the first on flipping the bit after the first one until you reach the far left:  
\* 1100 0001 0001 00...00  
\* 1100 0001 0101 00...00  
\* 1100 0001 1101 00...00  
\* 1100 0000 1101 00...00  
\* 1100 0010 1101 00...00  
\* 1100 0110 1101 00...00  
\* 1100 1110 1101 00...00  
\* 1101 1110 1101 00...00  
\* 1111 1110 1101 00...00  
\* 1011 1110 1101 00...00  
\* 0011 1110 1101 00...00
4. Then add all the base two 1's together to get the integer:  
\*  $2^{29} + 2^{28} + 2^{27} + 2^{26} + 2^{25} + 2^{23} + 2^{22} + 2^{20} = 1053818880$
5. And multiply it by the sign, which is -1:  
\*  $-1(1053818880) = -1053818880$

(b) an unsigned integer;

As an unsigned integer, add the base two ones together:  
 $2^{31} + 2^{30} + 2^{24} + 2^{21} + 2^{20} = 3241148416$

(c) a single precision IEEE 754 floating-point number.

$$(-1)^S \times (1 + F) \times 2^E$$

S = Sign

F = fraction ( $0 < F < 1$ )

E = Exponent - Bias

Bias = 127 for single precision

If the first character of the word is a one then the sign is negative:

S = 1

The next eight bits make up the exponent:

$$E = 100\ 0001\ 0 = 2^7 + 2^1 = 128 + 2 = 130 - \text{Bias} = 130 - 127 = 3$$

F = 011 0000 0000 0000 0000 0000

Going from left to right, with left as the zero point and going into the negative:

$$2^{-2} + 2^{-3} = 0.25 + 0.125 = 0.375$$

$$(-1)^1 \times (1 + 0.375) \times 2^3$$

$$-1 \times 1.375 \times 8$$

The answer is -11

6. (6%) Find computer representations of the following numbers:

(a) -1022 as a two's complement 32-bit binary number;

First turn it into a positive binary number:

S = -1

N = 1022

$$- 1022/2 = 511 \rightarrow 0$$

$$- 511/2 = 255.5 \rightarrow 1$$

$$- 255/2 = 127.5 \rightarrow 1$$

$$- 127/2 = 63.5 \rightarrow 1$$

$$- 63/2 = 31.5 \rightarrow 1$$

$$- 31/2 = 15.5 \rightarrow 1$$

$$- 15/2 = 7.5 \rightarrow 1$$

$$- 7/2 = 3.5 \rightarrow 1$$

$$- 3/2 = 1.5 \rightarrow 1$$

$$- 1/2 = 0.5 \rightarrow 1$$

Now we have 0...0011 1111 1110

Since it needs to be a two's complement number we need to use the algorithm from before to flip the bits.

Which becomes: 1...1100 0000 0010 Since it is a negative number we leave the first character as a one:

Answer = 1111 1111 1111 1111 1111 1100 0000 0010

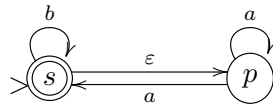
(b)  $-32.75$  as an IEEE 754 32-bit floating-point number;

□

(c) 77 as a two's complement 32-bit binary number.

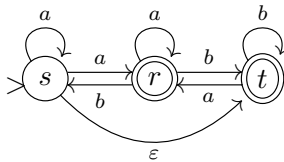
□

7. (6%) Consider the following finite automaton:



- Give all the computations of the automaton on the input strings  $aab$ ,  $aba$ , and  $\varepsilon$ , and determine if the strings are accepted.
- Describe the language accepted by the automaton in English.
- Describe the language accepted by the automaton by means of a regular expression.
- Describe the language accepted by the automaton by means of a context-free grammar.

8. (10%) Transform, using the subset construction, the following nondeterministic finite automaton into an equivalent deterministic finite automaton. Show your working.



What is the language of this automaton.

- (7%) Design a (deterministic or nondeterministic) finite automaton  $A$  such that  $L(A)$  consists of all strings over the alphabet  $\{0, 1\}$  that begin with 01 and do not end with 11. Find a regular expression representing the language  $L(A)$ .
- (6%) Design a (deterministic or nondeterministic) finite automaton  $A$  such that  $L(A)$  consists of all strings over the alphabet  $\{0, 1\}$  whose third symbol from the right end is 1 (for example, 100101 is in  $L(A)$ , but 100011 is not). Find a regular expression representing  $L(A)$ .
- (8%) Convert the regular language  $L[x((y \cup x)^*x)^*y]$  to a finite automaton accepting it.
- (4%) Consider the following context free grammar:

$$S \rightarrow SS, \quad S \rightarrow L0L0L, \quad L \rightarrow \varepsilon, \quad L \rightarrow 1L, \quad L \rightarrow 0L.$$

- Give a derivation for the string 101101.
- Describe in English the language of this grammar.

13. (6%) Construct context free grammars for the following languages

- (a)  $\{w \in \{0,1\}^* \mid w \text{ starts and ends with different symbols}\},$   
 (b)  $\{w \in \{0,1\}^* \mid \text{the length of } w \text{ is even}\}$

14. (15%) Construct a context free grammar and a pushdown automaton for the language of words over the alphabet  $\{0,1\}$  that start and end with the same symbol and have the same number of 0s as 1s.

15. (5%) Consider the following transition table of a Turing machine:

|     |                  |     |               |
|-----|------------------|-----|---------------|
| $s$ | 0                | $s$ | $\rightarrow$ |
| $s$ | 1                | $s$ | $\rightarrow$ |
| $s$ | $\sqcup$         | $p$ | $\leftarrow$  |
| $s$ | $\triangleright$ | $s$ | $\rightarrow$ |
| $p$ | 0                | $h$ | 1             |
| $p$ | 1                | $p$ | $\rightarrow$ |
| $p$ | $\sqcup$         | $h$ | 0             |
| $p$ | $\triangleright$ | $s$ | $\rightarrow$ |

(i) Give the computations of the machine starting with the configurations

- $(s, \triangleright \underline{0})$ ,
- $(s, \triangleright \underline{1}11)$ ,
- $(s, \triangleright \underline{1}00)$ .

(ii) Describe in English what this Turing machine does.

16. (9%) Consider the following  $\mathbb{N} \rightarrow \mathbb{N}$  function  $f$ :

$$f(n) = \begin{cases} 4n & \text{if } n \text{ is odd,} \\ n/2 & \text{if } n \text{ is even.} \end{cases}$$

(Don't forget that all numbers are represented in binary.)

- (i) Explain what it means to say that a Turing machine *computes* this function  $f$ .
- (ii) Give an implementation level description in English of a Turing machine that computes this  $f$ .
- (iii) Give the complete transition table of this Turing machine.
- (iv) Give the computations of your Turing machine on inputs 0, 11 and 100.