## Database Concepts- Assignment 8 Due: December 5 by 11:59pm

1. (a) Consider the following parameters:

block size = 4096 bytes block-address size = 8 bytes

block access time = 20 ms (micro seconds)

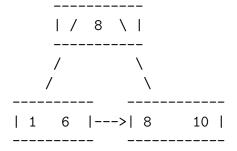
record size = 100 bytes record key size = 9 bytes

Assume that there is a B<sup>+</sup>-tree, adhering to these parameters, that indexes 10 million records on their primary key values.

- i. Specify (in ms) the minimum time to determine whether a record with key k is in the B<sup>+</sup>-tree.
- ii. Specify (in ms) the maximum time to determine whether a record with key k is in the B<sup>+</sup>-tree.
- iii. How many records would there need to be to increase the minimum time to determine whether a record with key k is in the B<sup>+</sup>-tree by approximately 50 ms?

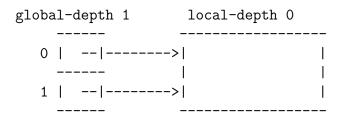
Show all the intermediate computations leading to your answer.

(b) Consider the following B<sup>+</sup>-tree of order 2 that holds records with keys 1, 6, 8, and 10. (Observe that (a) an internal node of a B<sup>+</sup>-tree of order 2 can have either 1 or 2 keys values, and 2 or 3 sub-trees, and (b) a leaf node can have either 1 or 2 key values.)



i. Show the contents of your  $B^+$ -tree after inserting records with keys 5, 9, 12, and 3, in that order.

- ii. Starting from your answer in question 1(b)i, show the contents of your  $B^+$ -tree after deleting records with keys 1, 12, 3, and 9, in that order.
- 2. (a) Consider an extensible hashing data structure wherein (1) the initial global depth is set at 1 and (2) all directory pointers point to the same **empty** block which has local depth 0. So the hashing structure looks like this:



Assume that a page can hold at most two records.

- i. Show the state of the hash data structure after each of the following insert sequences:
  - A. records with keys 1 and 5.
  - B. records with keys 2 and 3.
  - C. records with keys 4 and 7.
- ii. Starting from the answer you obtained for Question 2(a)i, show the state of the hash data structure after each of the following delete sequences:
  - A. records with keys 1 and 3.
  - B. records with keys 2 and 4.
  - C. records with keys 5 and 7.
- iii. Give an example where the insertion of a record in an extendible hash data structure can result in the recursive doubling of the directory of the hash data structure.

- 3. Consider a database wherein we maintain 2 data files:
  - (a) a file STUDENT which keeps student records of the form (SId, Sname, major; SId is a key),
  - (b) an ENROLLMENT file which keeps records of the form (SId, Cno).

Specify a file organization for each of these files which would enable efficient processing of the following query

"Find the course numbers of courses taken by students who major in a field with a name that begins with the letter 'A'." (For example African-Studies, Anthropology, Archeology etc.)

Argue why your file organization provides efficient processing of this query.

4. Let R(A, B) and S(B, C) be two relations and consider their natural join  $R \bowtie S$ .

Assume that R has 200000 records and that S has 10000 records.

Furthermore, assume that 20 records of R can fit in a block and that 10 records of S can fit in a block.

Assume that you have a main-memory buffer with 51 blocks. (One of these block is reserved for output purposes.) In the following questions, you should consider making maximum use of the buffer.

- (a) How many block IO's are necessary to perform  $R \bowtie S$  using the nested-loops join algorithm? Show your analysis.
- (b) How many block IO's are necessary to perform  $R \bowtie S$  using the merge-join algorithm? Show your analysis.
- (c) How many block IO's are necessary to perform  $R \bowtie S$  using the hash-join algorithm? Show your analysis.

- 5. Give an example of two conflict-equivalent, but different, non-serializable schedules.
- 6. Give an example of a serializable schedule which is conflict-equivalent with two different serial schedules.
- 7. Consider the following transactions:

```
T1: read(A);
    read(B);
    if A = 0 then B := B+1;
    write(B).

T2: read(B);
    read(A);
    if B = 0 then A := A+1;
    write(A).
```

Let the consistency requirement be  $A = 0 \lor B = 0$ , and let A = B = 0 be the initial values.

- (a) Show that each serial schedule involving transaction T1 and T2 preserves the consistency requirement of the database.
- (b) Construct a schedule on T1 and T2 that produces a non-serializable schedule.
- (c) Is there a non-serial schedule on T1 and T2 that produces a serializable schedule. If so, give an example.
- (d) i. Add lock and unlock instructions to T1 and T2, so that they observe the two-phase locking protocol, but in such a way that interleaving between operations in T1 and T2 is still possible.
  - ii. Can the execution of these transactions result in a deadlock? If so, give an example.