

< Search in Rotated Sorted Array

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Search in Rotated Sorted Array

LeetCode 👤 Admin 📅 Jan 10, 2019

Solution

Approach 1: Binary search

The problem is to implement a search in $\mathcal{O}(\log N)$ time that gives an idea to use a binary search.

The algorithm is quite straightforward :

- Find a rotation index `rotation_index` , i.e. index of the smallest element in the array. Binary search works just perfect here.
- `rotation_index` splits array in two parts. Compare `nums[0]` and `target` to identify in which part one has to look for `target` .
- Perform a binary search in the chosen part of the array.

Target = 5

4 5 6 7 8 1 2 3

4 5 6 7 8 1 2 3

4 5 6 7 8 1 2 3

4 5 6 7 8 1 2 3

4 5 6 7 8 1 2 3

4 5 6 7 8 1 2 3

① Find rotation index = index of the smallest element

1. Pick the element in the middle as a pivot. $7 > 8 = \text{False}$, hence 8 is not the smallest element.

2. $7 > 4$, continue the search on the right side.

3. Pick the element in the middle as a pivot. $1 > 2 = \text{False}$, hence 2 is not the smallest element.

4. $8 > 1$, continue the search on the left side.

5. Pick the element in the middle as a pivot. $8 > 1 = \text{True}$, hence 1 is the smallest element and **rotation_index = 5**.

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Copy

```
Java Python
1 class Solution {
2     int [] nums;
3     int target;
4
5     public int find_rotate_index(int left, int right) {
6         if (nums[left] < nums[right])
7             return 0;
8
9         while (left <= right) {
10             int pivot = (left + right) / 2;
11             if (nums[pivot] > nums[pivot + 1])
12                 return pivot + 1;
13             else {
14                 if (nums[pivot] < nums[left])
15                     right = pivot - 1;
16                 else
17                     left = pivot + 1;
18             }
19         }
20         return 0;
21     }
22
23     public int search(int left, int right) {
24         /*
25          * Binary search
26          */
27         while (left <= right) {
```

Complexity Analysis

- Time complexity : $\mathcal{O}(\log N)$.
- Space complexity : $\mathcal{O}(1)$.

Approach 2: One-pass Binary Search

Instead of going through the input array in two passes, we could achieve the goal in one pass with an *revised* binary search.

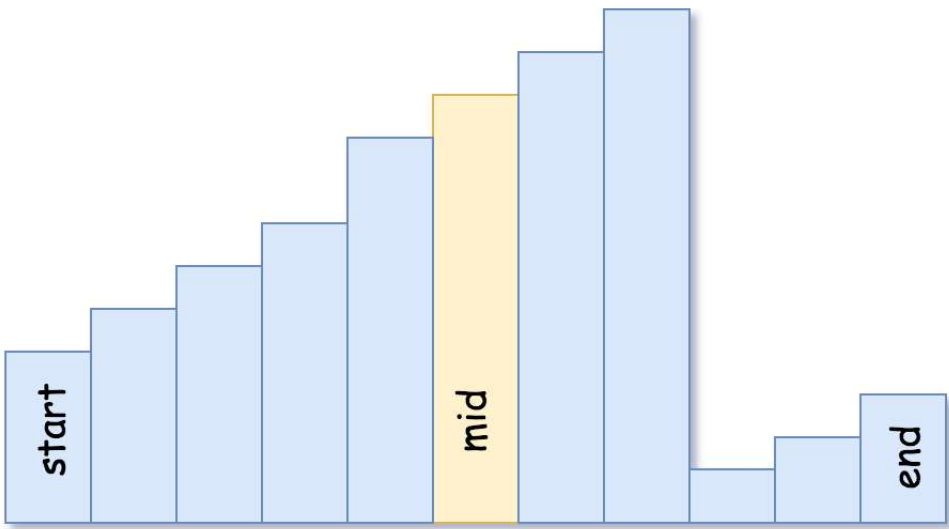
The idea is that we add some additional *condition checks* in the normal binary search in order to better *narrow down* the scope of the search.

Algorithm

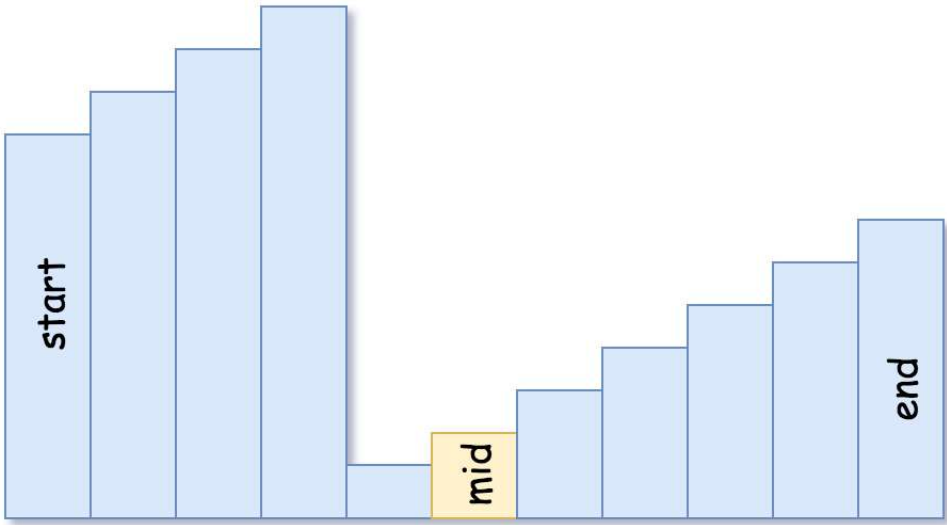
As in the normal binary search, we keep two pointers (*i.e.* `start` and `end`) to track the search scope. At each iteration, we reduce the search scope into half, by moving either the `start` or `end` pointer to the middle (*i.e.* `mid`) of the previous search scope.

Here are the detailed breakdowns of the algorithm:

- Initiate the pointer `start` to `0` , and the pointer `end` to `n - 1` .
- Perform standard binary search. While `start <= end` :
 - Take an index in the middle `mid` as a pivot.
 - If `nums[mid] == target` , the job is done, return `mid` .
 - Now there could be two situations:
 - Pivot element is larger than the first element in the array, *i.e.* the subarray from the first element to the pivot is non-rotated, as shown in the following graph.



- If the target is located in the non-rotated subarray:
go left: `end = mid - 1`.
 - Otherwise: go right: `start = mid + 1`.
- o Pivot element is smaller than the first element of the array, *i.e.* the rotation index is somewhere between `0` and `mid`. It implies that the sub-array from the pivot element to the last one is non-rotated, as shown in the following graph.



- If the target is located in the non-rotated subarray:
go right: `start = mid + 1`.
- Otherwise: go left: `end = mid - 1`.

- We're here because the target is not found. Return -1.

Implementation

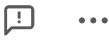
JavaPythonCopy

```
1 class Solution {
2   public int search(int[] nums, int target) {
3     int start = 0, end = nums.length - 1;
4     while (start <= end) {
5       int mid = start + (end - start) / 2;
6       if (nums[mid] == target) return mid;
7       else if (nums[mid] >= nums[start]) {
8         if (target >= nums[start] && target < nums[mid]) end = mid - 1;
9         else start = mid + 1;
10      }
11      else {
12        if (target <= nums[end] && target > nums[mid]) start = mid + 1;
13        else end = mid - 1;
14      }
15    }
16    return -1;
17  }
18 }
```

Complexity Analysis

- Time complexity: $\mathcal{O}(\log N)$.
- Space complexity: $\mathcal{O}(1)$.

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reyou

Apr 27, 2019

Formula: If a sorted array is shifted, if you take the middle, always one side will be sorted. Take the recursion according to that rule.

- 1- take the middle and compare with target, if matches return.
- 2- if middle is bigger than left side, it means left is sorted
 - 2a- if [left] < target < [middle] then do recursion with left, middle - 1 (right)
 - 2b- left side is sorted, but target not in here, search on right side middle + 1 (left), right
- 3- if middle is less than right side, it means right is sorted
 - 3a- if [middle] < target < [right] then do recursion with middle + 1 (left), right
 - 3b- right side is sorted, but target not in here, search on left side left, middle -1 (right)

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kakrafoon

Jan 02, 2021

Took me a while to sort out the cases. I always get stuck figuring out whether to use '<=' or '<' in these type of problems. In this case, while the binary search itself was fine (<=), while looking for the pivot, I used '<'. Need to practice until this becomes second nature I think. Either that or just fix (<=) as the convention and take it from there.

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sutirtho

Jun 25, 2020

It is an moderately easy concept but can be very tricky to code.

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lim142857

Jul 27, 2019

Python 3 modified binary search (readable)(fast)(short)

```
# Initilize two pointers
begin = 0
end = len(nums) - 1
while begin <= end:
    mid = (begin + end)//2
    if nums[mid] == target:
        return mid
    if nums[mid] > nums[end]: # left side of mid is sorted
```

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Jul 14, 2021

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Jan 25, 2019

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Sithis

Jan 11, 2019

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softwareshortcut

Apr 19, 2019

Creating helper methods is key otherwise you'd easily get lost dealing with indices. Great exercise, thanks for sharing. I see shorter solutions in the comments, but they might be tricky to come up with during an interview.

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layak

Jun 06, 2020

why do we have equal to sign in this condition? `nums[mid] >= nums[start]`

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rubalwanchoo

Apr 03, 2019

In one of the test case, the input array is `[1,3]`, target 0
How is `[1,3]` a rotated array?

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