

Theory of Computation (60-354), Fall 2010
Solution to Midterm 2, 23 Nov., 2010
Total Marks : 40
Time: 80 minutes

Qn.1 Show that the following context-free grammar is ambiguous.

$$\begin{aligned}S &\rightarrow ABA|aBaa \\ A &\rightarrow a|aA \\ B &\rightarrow b\end{aligned}$$

by finding a string that has two leftmost or rightmost derivations.

[8 marks]

Ans: The string $abaa$ has the following two leftmost derivations.

1. $S \Rightarrow aBaa \Rightarrow abaa$
2. $S \Rightarrow ABA \Rightarrow aBA \Rightarrow abA \Rightarrow abaA \Rightarrow abaa$

Qn.2 Find a context-free grammar that generates the language accepted by the DPDA $P = (\{q_0, q_1\}, \{0, 1, 2\}, \{0, 1, 2, Z_0\}, \delta, q_0, Z_0)$ whose transition function is given below:

$$\begin{aligned}\delta(q_0, 0, Z_0) &= \{(q_0, 0Z_0)\} \\ \delta(q_0, 0, 0) &= \{(q_0, 00)\} \\ \delta(q_0, 1, Z_0) &= \{(q_1, Z_0)\} \\ \delta(q_0, 1, 0) &= \{(q_1, 0)\} \\ \delta(q_1, 2, 0) &= \{(q_1, \epsilon)\} \\ \delta(q_1, \epsilon, Z_0) &= \{(q_1, \epsilon)\}\end{aligned}$$

It is enough to provide templates for the productions corresponding to the first two transitions.

[8 marks]

Ans:

$$[q_0 Z_0 ?] \rightarrow 0[q_0 0!][!Z_0 ?]$$

$$[q_0 0 ?] \rightarrow 0[q_0 0!][!0 ?]$$

$$[q_0 Z_0 ?] \rightarrow 1[q_1 Z_0 ?]$$

$$[q_0 0 ?] \rightarrow 1[q_1 0 ?]$$

$$[q_1 0 q_1] \rightarrow 2$$

$$[q_1 Z_0 q_1] \rightarrow \epsilon$$

Qn.3 Design a context-free grammar that generates the language $L = \{a^n b^m c^{n+m} | n \geq 0, m \geq 0\}$. (**Hint:** We can rewrite the string $a^n b^m c^{n+m}$ as $a^n b^m c^m c^n$, which shows that the outer (a, c) -pairs and the inner (b, c) -pairs have to be generated by different variables). You must provide explanations for the productions of your grammar.

[8 marks]

Ans:

$$S \rightarrow aSc | B$$

$$B \rightarrow bBc | \epsilon$$

Qn.4 Use the pumping lemma (for context-free languages) to show that the language $L = \{a^i b^{2i} c^{3i} \mid i \geq 0\}$ over $\Sigma = \{a, b, c\}$ is not context-free. (**Hint:** You may choose the string $a^n b^{2n} c^{3n}$ in L to which to apply the PL). Clearly mention the i that you choose to show that $z' = uv^i wx^i y$ is not in L in each of the different cases that arises.

[8 marks]

Ans: Choose $z = a^n b^{2n} c^{3n}$. Let $uvwxy$ be an adversarial decomposition of z .

Case 1: vwx consists of a 's (b 's or c 's) alone. Setting $i = 0$ we get fewer a 's (b 's or c 's), disturbing the $1 : 2 : 3$ ratio of the a 's, b 's and c 's.

Case 2: vwx consists of a 's and b 's (b 's and c 's). Setting $i = 0$ may maintain the ratio of a 's to the b 's but not their ratios with respect to the c 's.

Qn.5 Use the CYK algorithm to determine whether the string $w = ababa$ is in the language generated by the following context-free grammar in CNF.

$$\begin{aligned}S &\rightarrow AB \\A &\rightarrow BB|a \\B &\rightarrow AB|b\end{aligned}$$

[8 marks]

Ans:

$$S_{11} = \{A\}, S_{22} = \{B\}, S_{33} = \{A\}, S_{44} = \{B\}, S_{55} = \{A\}$$

$$S_{12} = \{S, B\}, S_{23} = \{\}, S_{34} = \{S, B\}, S_{45} = \{\}$$

$$S_{13} = \{\}, S_{24} = \{A\}, S_{35} = \{\}$$

$$S_{14} = \{A\}, S_{25} = \{\}$$