

FONS

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*The mathematician does not study mathematics because it is useful; he studies it because he delights in it and he delights in it because it is beautiful.*

— Henri Poincaré (1854–1912)  
French mathematician and theoretical physicist

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This is (still!) an incomplete draft. Please send corrections and comments to [ryanjooruian18@gmail.com](mailto:ryanjooruian18@gmail.com), or pull-request at <https://github.com/Ryanjoo18/latex>.

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# Preface

*Fons*, derived from the Latin word for source or fountain, introduces the core concepts of university-level mathematics. Just as a fountain provides a continuous wellspring of water, *Fons* aims to be a continuous source of knowledge for you.

## About the author

At this moment of writing, I am a high school student working on my A Level studies in Singapore. I have about 11 years of participating in mathematics competitions, including three years of experience in mental arithmetic and the rest few years in mathematics olympiad.

## About this book

This book mainly serves as my notes when studying mathematics at the university level. Feel free to refer to it too.

## References

- Lecture notes by the University of Oxford.
- Lecture notes on MIT OpenCourseWare.
- Rudin, W. (1953). *Principles of Mathematical Analysis*. McGraw-Hill.

## Problem Solving

In his book “How to Solve It”, George Pólya outlined the following problem solving cycle:

### 1. Understand the problem

Ask yourself the following questions:

- Do you understand all the words used in stating the problem?
- Is it possible to satisfy the condition? Is the condition sufficient to determine the unknown? Or is it insufficient? Or redundant? Or contradictory?
- What are you asked to find or show? Can you restate the problem in your own words?
- Draw a figure. Introduce suitable notation.
- Is there enough information to enable you to find a solution?

### 2. Devise a plan

A partial list of strategies is included:

- Guess and check
- Look for a pattern
- Make an orderly list
- Draw a picture
- Eliminate possibilities
- Solve a simpler problem
- Use symmetry
- Use a model
- Consider special cases
- Work backwards
- Use direct reasoning
- Use a formula
- Solve an equation
- Be ingenious

### 3. Execute the plan

This step is usually easier than devising the plan. In general, all you need is care and patience, given that you have the necessary skills. Persist with the plan that you have chosen. If it continues not to work discard it and choose another. Don't be misled, this is how mathematics is done, even by professionals.

- Carrying out your plan of the solution, check each step. Can you see clearly that the step is correct? Can you prove that it is correct?

### 4. Check and expand

Pólya mentions that much can be gained by taking the time to reflect and look back at what you have done, what worked, and what didn't. Doing this will enable you to predict what strategy to use to solve future problems.

Look back reviewing and checking your results. Ask yourself the following questions:

- Can you check the result? Can you check the argument?
- Can you derive the solution differently? Can you see it at a glance?
- Can you use the result, or the method, for some other problem?

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# Acronyms

**H2FM** H2 Further Mathematics. 28

**H3M** H3 Mathematics. 123, 125, 126, 128



**Part I**

**Preliminaries**

# 1 Mathematical Reasoning and Logic

## §1.1 Logical statements and notation

It is useful to be familiar with the following terminology.

- A **definition** is a precise and unambiguous description of the meaning of a mathematical term. It characterises the meaning of a word by giving all the properties and only those properties that must be true.
- A **theorem** is a true mathematical statement that can be proven mathematically. In a mathematical paper, the term theorem is often reserved for the most important results.
- A **lemma** is a minor result whose sole purpose is to help in proving a theorem. It is a stepping stone on the path to proving a theorem. Very occasionally lemmas can take on a life of their own.
- A **corollary** is a result in which the (usually short) proof relies heavily on a given theorem. We often say that “this is a corollary of Theorem A”.
- A **proposition** is a proven and often interesting result, but generally less important than a theorem.
- A **conjecture** is a statement that is unproved, but is believed to be true.
- An **axiom** is a statement that is assumed to be true without proof. These are the basic building blocks from which all theorems are proven.
- An **identity** is a mathematical expression giving the equality of two (often variable) quantities.
- A **paradox** is a statement that can be shown, using a given set of axioms and definitions, to be both true and false. Paradoxes are often used to show the inconsistencies in a flawed theory.

### §1.1.1 Notation

A **proposition** is a sentence which has exactly one truth value, i.e. it is either true or false, but not both and not neither. A proposition is denoted by uppercase letters such as  $P$  and  $Q$ . If the proposition  $P$  depends on a variable  $x$ , it is sometimes helpful to denote it by  $P(x)$ .

We can do some algebra on propositions, which include

- (i) **equivalence**, denoted by  $P \equiv Q$ , which means  $P$  and  $Q$  are logically equivalent statements;
- (ii) **conjunction**, denoted by  $P \wedge Q$ , which means “ $P$  and  $Q$ ”;
- (iii) **disjunction**, denoted by  $P \vee Q$ , which means “ $P$  or  $Q$ ”;
- (iv) **negation**, denoted by  $\neg P$ , which means “not  $P$ ”.

Here are some useful properties when handling logical statements. You can easily prove all of them using truth tables.

- Double negation law:

$$P \equiv \neg(\neg P)$$

- Commutative property:

$$P \wedge Q \equiv Q \wedge P, \quad P \vee Q \equiv Q \vee P$$

- Associative property for conjunction:

$$(P \wedge Q) \wedge R \equiv P \wedge (Q \wedge R)$$

- Associative property for disjunction:

$$(P \vee Q) \vee R \equiv P \vee (Q \vee R)$$

- Distributive property for conjunction across disjunction:

$$P \wedge (Q \vee R) \equiv (P \wedge Q) \vee (P \wedge R)$$

- Distributive property for disjunction across conjunction:

$$P \vee (Q \wedge R) \equiv (P \vee Q) \wedge (P \vee R)$$

- **De Morgan's Laws:**

$$\neg(P \vee Q) \equiv (\neg P \wedge \neg Q)$$

$$\neg(P \wedge Q) \equiv (\neg P \vee \neg Q)$$

#### Exercise 1.1.1

Assume that  $x$  is a fixed real number. What is the negation of the statement  $1 < x < 2$ ?

**Solution** The negation of  $1 < x < 2$  is “it is not the case that  $1 < x < 2$ ”. However this is not useful.

Note that  $1 < x < 2$  means  $1 < x$  and  $x < 2$ . Let  $P : 1 < x$  and  $Q : x < 2$ . Then the statement  $1 < x < 2$  is  $P \wedge Q$ .

By De Morgan's Laws, we have  $\neg(P \wedge Q) \equiv \neg P \vee \neg Q$ .

The *Trichotomy Axiom of real numbers* states that given fixed real numbers  $a$  and  $b$ , exactly one of the statements  $a < b, a = b, b < a$  is true. Hence  $\neg P \equiv \neg(1 < x) \equiv (x \leq 1)$  and  $\neg Q \equiv \neg(x < 2) \equiv (x \geq 2)$ .

Thus

$$\neg(1 < x < 2) \equiv \neg(P \wedge Q) \equiv \neg P \vee \neg Q \equiv (1 \geq x) \vee (x \geq 2).$$

Therefore the negation of  $1 < x < 2$  is logically equivalent to the statement  $x \leq 1$  or  $x \geq 2$ .  $\square$

#### Exercise 1.1.2

Assume that  $n$  is a fixed positive integer. Find a useful denial of the statement

$$n = 2 \text{ or } n \text{ is odd.}$$

**Solution** Using De Morgan's Laws,

$$\begin{aligned} \neg[(n = 2) \vee (n \text{ is odd})] &\equiv \neg(n = 2) \wedge \neg(n \text{ is odd}) \\ &\equiv (n \neq 2) \wedge (n \text{ is even}) \end{aligned}$$

where we are using the fact that every integer is either even or odd, but not both.

Thus a useful denial of the given statement is:  $n$  is an even integer other than 2.  $\square$

### §1.1.2 If, only if, $\implies$

**Implication** is denoted by  $P \implies Q$ , which means “ $P$  implies  $Q$ ”, i.e. if  $P$  holds then  $Q$  also holds. It is equivalent to saying “If  $P$  then  $Q$ ”. The only case when  $P \implies Q$  is false is when the hypothesis  $P$  is true and the conclusion  $Q$  is false.

$P \implies Q$  is known as a **conditional statement**.  $P$  is known as the **hypothesis**,  $Q$  is known as the **conclusion**.

Statements of this form are probably the most common, although they may sometimes appear quite differently. The following all mean the same thing:

- (i) if  $P$  then  $Q$ ;
- (ii)  $P$  implies  $Q$ ;
- (iii)  $P$  only if  $Q$ ;
- (iv)  $P$  is a sufficient condition for  $Q$ ;
- (v)  $Q$  is a necessary condition for  $P$ .

The **converse** of  $P \implies Q$  is given by  $Q \implies P$ ; both are not logically equivalent.

The **inverse** of  $P \implies Q$  is given by  $\neg P \implies \neg Q$ , i.e. the hypothesis and conclusion of the statement are both negated.

The **contrapositive** of  $P \implies Q$  is given by  $\neg Q \implies \neg P$ ; both are logically equivalent.

**How to prove:** To prove  $P \implies Q$ , start by assuming that  $P$  holds and try to deduce through some logical steps that  $Q$  holds too. Alternatively, start by assuming that  $Q$  does not hold and show that  $P$  does not hold (that is, we prove the contrapositive).

### §1.1.3 If and only if, iff, $\iff$

**Bidirectional implication** is denoted by  $P \iff Q$ , which means both  $P \implies Q$  and  $Q \implies P$ . We can read this as “ $P$  if and only if  $Q$ ”. The letters “iff” are also commonly used to stand for ‘if and only if’.

$$P \iff Q \equiv (P \implies Q) \wedge (Q \implies P)$$

$P \iff Q$  is true exactly when  $P$  and  $Q$  have the same truth value.

$P \iff Q$  is known as a **biconditional statement**.

These statements are usually best thought of separately as ‘if’ and ‘only if’ statements.

**How to prove:** To prove  $P \iff Q$ , prove the statement in both directions, i.e. prove both  $P \implies Q$  and  $Q \implies P$ . Remember to make very clear, both to yourself and in your written proof, which direction you are doing.

### §1.1.4 Quantifiers

The **universal quantifier** is denoted by  $\forall$ , which means “for all” or “for every”. An universal statement has the form  $\forall x \in X, P(x)$ .

The **existential quantifier** is denoted by  $\exists$ , which means “there exists”. An existential statement has the form  $\exists x \in X, P(x)$ , where  $X$  is known as the **domain**.

These are versions of De Morgan’s laws for quantifiers:

$$\neg \forall x \in X, P(x) \equiv \exists x \in X, \neg P(x)$$

$$\neg \exists x \in X, P(x) \equiv \forall x \in X, \neg P(x)$$

#### Exercise 1.1.3

Find a useful denial of the statement

for all real numbers  $x$ , if  $x > 2$ , then  $x^2 > 4$

**Solution** In logical notation, this statement is  $(\forall x \in \mathbb{R})[x > 2 \implies x^2 > 4]$ .

$$\begin{aligned} \neg\{(\forall x \in \mathbb{R})[x > 2 \implies x^2 > 4]\} &\equiv (\exists x \in \mathbb{R})\neg[x > 2 \implies x^2 > 4] \\ &\equiv (\exists x \in \mathbb{R})\neg[(x > 2) \vee (x^2 > 4)] \\ &\equiv (\exists x \in \mathbb{R})[(x > 2) \wedge (x^2 \leq 4)] \end{aligned}$$

Therefore a useful denial of the statement is:

there exists a real number  $x$  such that  $x > 2$  and  $x^2 \leq 4$ .

□

#### Exercise 1.1.4

Negate surjectivity.

**Solution** If  $f : X \rightarrow Y$  is not surjective, then it means that there exists  $y \in Y$  not in the image of  $X$ , i.e. for all  $x$  in  $X$  we have  $f(x) \neq y$ .

$$\begin{aligned} \neg \forall y \in Y, \exists x \in X, f(x) = y &\iff \exists y \in Y, \neg(\exists x \in X, f(x) = y) \\ &\iff \exists y \in Y, \forall x \in X, \neg(f(x) = y) \\ &\iff \exists y \in Y, \forall x \in X, f(x) \neq y \end{aligned}$$

□

**How to prove:** To prove a statement of the form  $\forall x \in X$  s.t.  $P(x)$ , start the proof with ‘Let  $x \in X$ .’ or ‘Suppose  $x \in X$  is given.’ to address the quantifier with an arbitrary  $x$ ; provided no other assumptions about  $x$  are made during the course of proving  $P(x)$ , this will prove the statement for all  $x \in X$ .

**How to prove:** To prove a statement of the form  $\exists x \in X$  s.t.  $P(x)$ , there is not such a clear steer about how to continue: you may need to show the existence of an  $x$  with the right properties; you may need to demonstrate logically that such an  $x$  must exist because of some earlier assumption, or it may be that you can show constructively how to find one; or you may be able to prove by contradiction, supposing that there is no such  $x$  and consequently arriving at some inconsistency.

**Remark** Read from left to right, and as new elements or statements are introduced they are allowed to depend on previously introduced elements but cannot depend on things that are yet to be mentioned.

**Remark** To avoid confusion, it is a good idea to keep to the convention that the quantifiers come first, before any statement to which they relate.

## §1.2 Proofs

### §1.2.1 Direct Proof

A direct proof of  $P \implies Q$  is a series of valid arguments that start with the hypothesis  $P$  and end with the conclusion  $Q$ . It may be that we can start from  $P$  and work directly to  $Q$ , or it may be that we make use of  $P$  along the way.

### §1.2.2 Proof by Contrapositive

To prove  $P \implies Q$ , we can instead prove  $\neg Q \implies \neg P$ .

#### Exercise 1.2.1

For every integer  $a$ , prove that if  $3a^2 + 1$  is even, then  $a$  is odd.

**Proof** We prove this by contrapositive.

Suppose  $a$  is not odd. So  $a = 2k$  for some integer  $k$ . Then

$$3a^2 + 1 = 3(2k)^2 + 1 = 2(6k^2) + 1.$$

Since  $3a^2 + 1 = 2q + 1$  for some integer  $q$ , hence  $3a^2 + 1$  is odd. □

#### Exercise 1.2.2

For  $m \in \mathbb{Z}$ , prove that if  $3 \mid m^2$  then  $3 \mid m$ .

**Proof** We prove this by contrapositive.

Suppose  $3 \nmid m$ . We shall prove  $3 \nmid m^2$ .

**Case 1:**  $m = 3k + 1$

Then  $m^2 = (3k + 1)^2 = 3(3k^2 + 2k) + 1$  so  $m^2$  has remainder 1 when divided by 3, hence  $3 \nmid m^2$ .

**Case 2:**  $m = 3k + 2$

This case shall be left as an exercise. □

### §1.2.3 Disproof by Counterexample

Providing a counterexample is the best method for refuting, or disproving, a conjecture.

In seeking counterexamples, it is a good idea to keep the cases you consider simple, rather than searching randomly. It is often helpful to consider “extreme” cases; for example, something is zero, a set is empty, or a function is constant.

The counterexample must make the hypothesis a true statement, and the conclusion a false statement.

### §1.2.4 Proof by Cases

You can sometimes prove a statement by:

1. Dividing the situation into cases which exhaust all the possibilities; and
2. Showing that the statement follows in all cases.

**Remark** It is important to cover all the possibilities.

### §1.2.5 Proof by Contradiction

To prove  $P$  by contradiction, suppose that  $P$  is false, i.e.  $\neg P$ . Similarly, to prove  $P \implies Q$  by contradiction, suppose that  $Q$  is false, i.e.  $P \wedge \neg Q$ .

Then show through some logical reasoning that this leads to a contradiction or inconsistency. We may arrive at something that contradicts the hypothesis  $P$ , or something that contradicts the initial supposition that  $Q$  is not true, or we may arrive at something that we know to be universally false.

#### Exercise 1.2.3: Irrationality of $\sqrt{2}$

Prove that  $\sqrt{2}$  is irrational.

**Proof** We prove by contradiction. Suppose otherwise, that  $\sqrt{2}$  is rational. Using the definition of rational numbers, we can write it as  $\sqrt{2} = \frac{a}{b}$  for some  $a, b \in \mathbb{Z}, b \neq 0$ .

We also assume that  $\frac{a}{b}$  is simplified to lowest terms, since that can obviously be done with any fraction. Notice that in order for  $\frac{a}{b}$  to be in simplest terms, both  $a$  and  $b$  cannot be even; one or both must be odd, otherwise we could simplify the fraction further.

Squaring both sides gives us

$$a^2 = 2b^2.$$

Since RHS is even, LHS must also be even. Hence it follows that  $a$  is even. Let  $a = 2k$  where  $k \in \mathbb{Z}$ . Substituting  $a = 2k$  into the above equation and simplifying it gives us

$$b^2 = 2k^2.$$

This means that  $b^2$  is even, from which follows again that  $b$  is even.

This is a contradiction, as we started out assuming that  $\frac{a}{b}$  was simplified to lowest terms, and now it turns out that  $a$  and  $b$  both would be even. Hence proven.  $\square$

#### Exercise 1.2.4

For any integer  $n$ , prove that there is no integer  $a > 1$  such that  $a \mid n$  and  $a \mid (n+1)$ .

**Proof** Suppose there is an integer  $n$  and integer  $a > 1$  such that  $a \mid n$  and  $a \mid (n+1)$ .

Then  $n = ak$  and  $n+1 = ah$  for some integers  $k$  and  $h$ .

$$ak + 1 = ah \implies 1 = a(h - k) \implies a \mid 1 \implies a = \pm 1$$

This contradicts  $a > 1$ .

Hence we conclude that, for any  $n$ , there is no integer  $a > 1$  such that  $a \mid n$  and  $a \mid (n+1)$ .  $\square$

### §1.2.6 Proof of Uniqueness

$\exists!$  means “there exists a unique”.

To prove uniqueness, we can do one of the following:

- Assume  $\exists x, y \in S$  such that  $P(x) \wedge P(y)$  is true and show  $x = y$ .
- Argue by assuming that  $\exists x, y \in S$  are distinct such that  $P(x) \wedge P(y)$ , then derive a contradiction.

To prove uniqueness and existence, we also need to show that  $\exists x \in S$  s.t.  $P(x)$  is true.



### §1.2.7 Proof of Existence

To prove existential statements, we can adopt two approaches:

1. Constructive proof (direct proof)
2. Non-constructive proof (indirect proof)

#### Constructive Proof

To prove statements of the form  $\exists x \in X$  s.t.  $P(x)$ , find or construct *a specific example* for  $x$ . To prove statements of the form  $\forall y \in Y, \exists x \in X$  s.t.  $P(x, y)$ , construct example for  $x$  *in terms of*  $y$  (since  $x$  is dependent on  $y$ ).

In both cases, you have to justify that your example  $x$

1. belongs to the domain  $X$ , and
2. satisfies the condition  $P$ .

#### Exercise 1.2.5

Prove that we can find 100 consecutive positive integers which are all composite numbers.

**Proof** We can prove this existential statement via constructive proof.

Our goal is to find integers  $n, n+1, n+2, \dots, n+99$ , all of which are composite.

Take  $n = 101! + 2$ . Then  $n$  has a factor of 2 and hence is composite. Similarly,  $n+k = 101! + (k+2)$  has a factor  $k+2$  and hence is composite for  $k = 1, 2, \dots, 99$ .

Hence the existential statement is proven.  $\square$

#### Exercise 1.2.6

Prove that for all rational numbers  $p$  and  $q$  with  $p < q$ , there is a rational number  $x$  such that  $p < x < q$ .

**Proof** We prove this by construction. Our goal is to find such a rational  $x$  *in terms of*  $p$  and  $q$ .

We take the average. Let  $x = \frac{p+q}{2}$  which is a rational number.

Since  $p < q$ ,

$$x = \frac{p+q}{2} < \frac{q+q}{2} = q \implies x < q$$

Similarly,

$$x = \frac{p+q}{2} > \frac{p+p}{2} = p \implies p < x$$

Hence we have shown the existence of rational number  $x$  such that  $p < x < q$ .

**Remark** For this type of question, there are two parts to prove: firstly,  $x$  satisfies the given statement; secondly,  $x$  is within the domain (for this question we do not have to prove  $x$  is rational since  $\mathbb{Q}$  is closed under addition).  $\square$

#### Exercise 1.2.7

Prove that for all rational numbers  $p$  and  $q$  with  $p < q$ , there is an irrational number  $r$  such that  $p < r < q$ .

**Proof** We prove this by construction. Similarly, our goal is to find an irrational  $r$  in terms of  $p$  and  $q$ .

Note that we cannot simply take  $r = \frac{p+q}{2}$ ; a simple counterexample is the case  $p = -1, q = 1$  where  $r = 0$  is clearly not irrational.

Since  $p$  lies in between  $p$  and  $q$ , let  $r = p + c$  where  $0 < c < q - p$ . Since  $c < q - p$ , we have  $c = \frac{q-p}{k}$  for some  $k > 1$ ; to make  $c$  irrational, we take  $k$  to be irrational.

Take  $r = p + \frac{q-p}{\sqrt{2}}$ . We need to show  $r$  is irrational and  $p < r < q$ .

**Part 1:**  $p < r < q$

Since  $q < p$ ,  $r = p + (\text{positive number}) > p$ . On the other hand,  $\frac{q-p}{\sqrt{2}} < q - p$  so  $r < p + (q - p) = q$ .

**Part 2:**  $r$  is irrational

We prove by contradiction. Suppose  $r$  is rational. We have  $\sqrt{2} = \frac{q-p}{r-p}$ . Since  $p, q, r$  are all rational (and  $r - p \neq 0$ ), RHS is rational. This implies that LHS is rational, i.e.  $\sqrt{2}$  is rational, a contradiction.  $\square$

### Non-constructive Proof

Use when specific examples are not easy or not possible to find or construct. Make arguments why such objects have to exist. May need to use proof by contradiction. Use definition, axioms or results that involve existential statements.

#### Exercise 1.2.8

Prove that every integer greater than 1 is divisible by a prime.

**Proof** If  $n$  is prime, then we are done as  $n \mid n$ .

If  $n$  is not prime, then  $n$  is composite. So  $n$  has a divisor  $d_1$  such that  $1 < d_1 < n$ . If  $d_1$  is prime then we are done as  $d_1 \mid n$ . If  $d_1$  is not prime then  $d_1$  is composite, has divisor  $d_2$  such that  $1 < d_2 < n$ .

If  $d_2$  is prime, then we are done as  $d_2 \mid d_1$  and  $d_1 \mid n$  imply  $d_2 \mid n$ . If  $d_2$  is not prime then  $d_2$  is composite, has divisor  $d_3$  such that  $1 < d_3 < d_2$ .

Continuing in this manner after  $k$  times, we will get

$$1 < d_k < d_{k-1} < \cdots < d_2 < d_1 < n$$

where  $d_i \mid n$  for all  $i$ .

This process must stop after finite steps, as there can only be a finite number of  $d_i$ 's between 1 and  $n$ . On the other hand, the process will stop only if there is a  $d_i$  which is a prime.

Hence we conclude that there must be a divisor  $d_i$  of  $n$  that is prime.  $\square$

**Remark** This proof is also known as *proof by infinite descent*, a method which relies on the well-ordering principle of the positive integers.

#### Exercise 1.2.9

Prove that the equation  $x^2 + y^2 = 3z^2$  has no solutions  $(x, y, z)$  in integers where  $z \neq 0$ .

**Proof** Suppose we have a solution  $(x, y, z)$ . Without loss of generality, we may assume that  $z > 0$ . By the least integer principle, we may also assume that our solution has  $z$  minimal. Taking remainders modulo 3, we see that

$$x^2 + y^2 \equiv 0 \pmod{3}$$

Recalling that squares may only be congruent to 0 or 1 modulo 3, we conclude that

$$x^2 \equiv y^2 \equiv 0 \implies x \equiv y \equiv 0 \pmod{3}$$

Writing  $x = 3a$  and  $y = 3b$  we obtain

$$9a^2 + 9b^2 = 3z^2 \implies 3(a^2 + b^2) = z^2 \implies 3 \mid z^2 \implies 3 \mid z$$

Now let  $z = 3c$  and cancel 3's to obtain

$$a^2 + b^2 = 3c^2.$$

We have therefore constructed another solution  $(a, b, c) = (\frac{x}{3}, \frac{y}{3}, \frac{z}{3})$  to the original equation. However  $0 < c < z$  contradicts the minimality of  $z$ .  $\square$

### Exercise 1.2.10

An odd prime  $p$  may be written as a sum of two squares if and only  $p \equiv 1 \pmod{4}$ .

**Proof** We again use the method of descent, though this time *constructively*.

( $\implies$ ) If  $p = x^2 + y^2$ , then both  $x$  and  $y$  are non-zero modulo  $p$ . Taking Legendre symbols, we see that

$$1 = \left(\frac{x^2}{p}\right) = \left(\frac{-y^2}{p}\right) = \left(\frac{-1}{p}\right) \implies p \equiv 1 \pmod{4}$$

( $\impliedby$ ) Suppose that  $p$  is a prime congruent to 1 modulo 4. We must show that there exist integers  $x, y$  such that  $x^2 + y^2 = p$ . We do this by descent:

1. Modulo  $p$ , the congruence  $x^2 + 1 \equiv 0$  has a solution  $x$  since  $-1$  is a quadratic residue. By taking  $y = 1$ , we may therefore assume the existence of a solution to an equation  $x^2 + y^2 = mp$  for some integer  $1 \leq m < p$ . If  $m = 1$  we are done. Otherwise ...

2. Define

$$\begin{cases} u \equiv x \pmod{m} \\ v \equiv y \pmod{m} \end{cases} \quad \text{such that } |u|, |v| \leq \frac{m}{2}.$$

Since  $xu + yv$ ,  $xv - yu$  and  $u^2 + v^2$  are all divisible by  $m$ , we may divide the identity

$$(u^2 + v^2)(x^2 + y^2) = (xu + yv)^2 + (xv - yu)^2$$

by  $m^2$  to obtain an equation in integers:

$$kp = \left(\frac{xu + yv}{m}\right)^2 + \left(\frac{xv - yu}{m}\right)^2 \quad \text{where } k = \frac{u^2 + v^2}{m} \leq \frac{m}{2}$$

3. We have therefore constructed an integer solution to  $X^2 + Y^2 = kp$  with  $k < m$ . If  $k \geq 2$ , simply repeat the process from step 2: by descent, we must eventually reach  $k = 1$ .  $\square$

## §1.2.8 Pigeonhole Principle

**Theorem 1.2.1** (Pigeonhole Principle (naive)). If  $m$  objects are placed into  $n$  boxes and  $m > n$ , then at least one box must contain more than one object.

**Theorem 1.2.2** (Pigeonhole Principle (general)). If more than  $k \cdot n$  objects are placed into  $n$  boxes, then at least one box must contain more than  $k$  objects.

### §1.2.9 Proof by Mathematical Induction

Induction is an extremely powerful method of proof used throughout mathematics. It deals with infinite families of statements which come in the form of lists. The idea behind induction is in showing how each statement follows from the previous one on the list – all that remains is to kick off this logical chain reaction from some starting point.

**Theorem 1.2.3** (Principle of Mathematical Induction (PMI)). Let  $P(n)$  be a family of statements indexed by  $\mathbb{Z}^+$ . Suppose that

- (i) (**base case**)  $P(1)$  is true and
- (ii) (**inductive step**) for all  $k \in \mathbb{Z}^+$ ,  $P(k) \implies P(k+1)$ .

Then  $P(n)$  is true for all  $n \in \mathbb{Z}^+$ .

Using logic notation, this is written as

$$\{P(1) \wedge (\forall n \in \mathbb{Z}^+)[P(k) \implies P(k+1)]\} \implies (\forall n \in \mathbb{Z}^+)P(n)$$

Induction is often visualised like toppling dominoes. The inductive step (ii) corresponds to placing each domino sufficiently close that it will be hit when the previous one falls over, and base case (i) corresponds to knocking over the first one.

$$P(1) \implies P(2) \implies \dots \implies P(k) \implies P(k+1) \implies \dots$$

#### Exercise 1.2.11

Prove that for any  $n \in \mathbb{Z}^+$ ,

$$\sum_{k=1}^n k = \frac{n(n+1)}{2}$$

**Proof** Let  $P(n)$  be the statement  $\sum_{k=1}^n k = \frac{n(n+1)}{2}$ .

Clearly  $P(1)$  holds because for  $n = 1$ , the sum on the LHS is 1 and the expression on the RHS is also 1.

Now suppose  $P(n)$  holds. Then we have

$$\sum_{k=1}^n k = \frac{n(n+1)}{2}$$

Adding  $n+1$  to both sides,

$$\begin{aligned} \sum_{k=1}^{n+1} k &= \frac{n(n+1)}{2} + (n+1) \\ &= \frac{(n+1)(n+2)}{2} \\ &= \frac{(n+1)[(n+1)+1]}{2} \end{aligned}$$

thus  $P(n+1)$  is true.

By PMI,  $P(n)$  is true for all  $n \in \mathbb{Z}^+$ . □

**Remark** Do not write  $P(n) = \frac{n(n+1)}{2}$ , as  $P(n)$  is a statement, not an expression (which does not have truth values).

A corollary of induction is if the family of statements holds for  $n \geq N$ , rather than necessarily  $n \geq 0$ :

**Corollary 1.2.4.** Let  $N$  be an integer and let  $P(n)$  be a family of statements indexed by integers  $n \geq N$ . Suppose that

- (i) (**base case**)  $P(N)$  is true and

- (ii) (**inductive step**) for all  $k \geq N$ ,  $P(k) \implies P(k+1)$ .

Then  $P(n)$  is true for all  $n \geq N$ .

**Proof** This follows directly by applying the above theorem to the statement  $Q(n) = P(n+N)$  for  $n \in \mathbb{N}$ .  $\square$

### Strong Induction

Another variant on induction is when the inductive step relies on some earlier case(s) but not necessarily the immediately previous case. This is known as **strong induction**:

**Theorem 1.2.5** (Strong Form of Induction). Let  $P(n)$  be a family of statements indexed by the natural numbers. Suppose that

- (i) (**base case**)  $P(1)$  is true and  
 (ii) (**inductive step**) for all  $m \in \mathbb{Z}^+$ , if for integers  $k$  with  $1 \leq k \leq m$ ,  $P(k)$  is true then  $P(m+1)$  is true.

Then  $P(n)$  is true for all  $n \in \mathbb{N}$ .

Using logic notation, this is written as

$$\{P(1) \wedge (\forall m \in \mathbb{Z}^+)[P(1) \wedge P(2) \wedge \cdots \wedge P(m) \implies P(m+1)]\} \implies (\forall n \in \mathbb{Z}^+)P(n)$$

**Proof** We can this it to an instance of “normal” induction by defining a related family of statements  $Q(n)$ .

Let  $Q(n)$  be the statement “ $P(k)$  holds for  $k = 0, 1, \dots, n$ ”. Then the conditions for the strong form are equivalent to

- (i)  $Q(0)$  holds and  
 (ii) for any  $n$ , if  $Q(n)$  is true then  $Q(n+1)$  is also true.

It follows by induction that  $Q(n)$  holds for all  $n$ , and hence  $P(n)$  holds for all  $n$ .  $\square$

The following example illustrates how the strong form of induction can be useful:

#### Example 1.2.1: Fundamental Theorem of Arithmetic

Every natural number greater than 1 may be expressed as a product of one or more prime numbers.

**Proof** Let  $P(n)$  be the statement that  $n$  may be expressed as a product of prime numbers.

Clearly  $P(2)$  holds, since 2 is itself prime.

Let  $n \geq 2$  be a natural number and suppose that  $P(m)$  holds for all  $m < n$ .

- If  $n$  is prime then it is trivially the product of the single prime number  $n$ .
- If  $n$  is not prime, then there must exist some  $r, s > 1$  such that  $n = rs$ . By the inductive hypothesis, each of  $r$  and  $s$  can be written as a product of primes, and therefore  $n = rs$  is also a product of primes.

Thus, whether  $n$  is prime or not, we have have that  $P(n)$  holds. By strong induction,  $P(n)$  is true for all natural numbers. That is, every natural number greater than 1 may be expressed as a product of one or more primes.  $\square$

### Cauchy Induction

**Theorem 1.2.6** (Cauchy Induction). Let  $P(n)$  be a family of statements indexed by  $\mathbb{Z}_{\geq 2}^+$ . Suppose that

- (i) (**base case**)  $P(2)$  is true and
- (ii) (**inductive step**) for all  $k \in \mathbb{Z}^+$ ,  $P(k) \implies P(2k)$  and  $P(k) \implies (k-1)$ .

Then  $P(n)$  is true for all  $n \in \mathbb{Z}_{\geq 2}^+$ .

#### Exercise 1.2.12

Using Cauchy Induction, prove the AM–GM Inequality for  $n$  variables, which states that for positive reals  $a_1, a_2, \dots, a_n$ ,

$$\frac{a_1 + a_2 + \dots + a_n}{n} \geq \sqrt[n]{a_1 a_2 \dots a_n}.$$

**Proof** Let  $P(n)$  be  $\frac{a_1 + a_2 + \dots + a_n}{n} \geq \sqrt[n]{a_1 a_2 \dots a_n}$ .

Base case  $P(2)$  is true because

$$\frac{a_1 + a_2}{2} \geq \sqrt{a_1 a_2} \iff (a_1 + a_2)^2 \geq 4a_1 a_2 \iff (a_1 - a_2)^2 \geq 0$$

Next we show that  $P(n) \implies P(2n)$ , i.e. if AM–GM holds for  $n$  variables, it also holds for  $2n$  variables:

$$\begin{aligned} \frac{a_1 + a_2 + \dots + a_{2n}}{2n} &= \frac{\frac{a_1 + a_2 + \dots + a_n}{n} + \frac{a_{n+1} + a_{n+2} + \dots + a_{2n}}{n}}{2} \\ &\geq \frac{\frac{a_1 + a_2 + \dots + a_n}{n} + \frac{a_{n+1} + a_{n+2} + \dots + a_{2n}}{n}}{2} \geq \frac{\sqrt[n]{a_1 a_2 \dots a_n} + \sqrt[n]{a_{n+1} a_{n+2} \dots a_{2n}}}{2} \\ &\geq \frac{\sqrt[n]{a_1 a_2 \dots a_n} + \sqrt[n]{a_{n+1} a_{n+2} \dots a_{2n}}}{2} \geq \sqrt{\sqrt[n]{a_1 a_2 \dots a_n} \sqrt[n]{a_{n+1} a_{n+2} \dots a_{2n}}} \\ &= \sqrt{\sqrt[n]{a_1 a_2 \dots a_n} \sqrt[n]{a_{n+1} a_{n+2} \dots a_{2n}}} = \sqrt[n]{a_1 a_2 \dots a_{2n}} \end{aligned}$$

The first inequality follows from  $n$ -variable AM–GM, which is true by assumption, and the second inequality follows from 2-variable AM–GM, which is proven above.

Finally we show that  $P(n) \implies P(n-1)$ , i.e. if AM–GM holds for  $n$  variables, it also holds for  $n-1$  variables. By  $n$ -variable AM–GM,  $\frac{a_1 + a_2 + \dots + a_n}{n} \geq \sqrt[n]{a_1 a_2 \dots a_n}$ . Let  $a_n = \frac{a_1 + a_2 + \dots + a_{n-1}}{n-1}$ . Then we have

$$\frac{a_1 + a_2 + \dots + a_{n-1} + \frac{a_1 + a_2 + \dots + a_{n-1}}{n-1}}{n} = \frac{a_1 + a_2 + \dots + a_{n-1}}{n-1}$$

So,

$$\begin{aligned} \frac{a_1 + a_2 + \dots + a_{n-1}}{n-1} &\geq \sqrt[n]{a_1 a_2 \dots a_{n-1}} \cdot \frac{a_1 + a_2 + \dots + a_{n-1}}{n-1} \\ \Rightarrow \left( \frac{a_1 + a_2 + \dots + a_{n-1}}{n-1} \right)^n &\geq a_1 a_2 \dots a_{n-1} \cdot \frac{a_1 + a_2 + \dots + a_{n-1}}{n-1} \\ &\Rightarrow \left( \frac{a_1 + a_2 + \dots + a_{n-1}}{n-1} \right)^{n-1} \geq a_1 a_2 \dots a_{n-1} \\ &\Rightarrow \frac{a_1 + a_2 + \dots + a_{n-1}}{n-1} \geq \sqrt[n-1]{a_1 a_2 \dots a_{n-1}} \end{aligned}$$

By Cauchy Induction, this proves the AM–GM inequality for  $n$  variables.  $\square$

**Other Variations**

Apart from proving  $P(n)$  indexed by  $\mathbb{Z}^+$ , we can also use PMI to prove statements of the form

- $(\forall n \in \mathbb{Z})P(n)$

**Base case:**  $P(0)$

**Inductive step:**  $(\forall k \in \mathbb{Z}_{\geq 0})P(k) \implies P(k+1)$  and  $(\forall k \in \mathbb{Z}_{\leq 0})P(k) \implies P(k-1)$

$$\dots \Longleftarrow P(-n) \Longleftarrow \dots \Longleftarrow P(-1) \Longleftarrow P(0) \implies P(1) \implies \dots \implies P(n) \implies \dots$$

- $(\forall n \in \mathbb{Q})P(n)$

**Base case:**  $P(0)$

**Inductive step:**  $P(x) \implies P(-x)$  and  $P\left(\frac{a}{b}\right) \implies P\left(\frac{a+1}{b}\right)$  and  $P\left(\frac{a}{b}\right) \implies P\left(\frac{a}{b+1}\right)$

**§1.2.10 Symmetry Principle**

**§1.2.11 Combinatorial Arguments and Proofs**



## Exercises

Some of the exercise problems here are from the “Number and Proofs” topic of H3 Mathematics, so the reader is assumed to have some basic knowledge in Number Theory, in particular modular arithmetic.

**Problem 1.** Let  $a, b$  be integers, not both 0. Prove that  $\gcd(a + b, a - b) \leq \gcd(2a, 2b)$ .

**Proof** Direct proof.

Let  $e = \gcd(a + b, a - b)$ . Then  $e \mid (a + b)$  and  $e \mid (a - b)$ . So

$$e \mid (a + b) + (a - b) \implies e \mid 2a$$

and

$$e \mid (a + b) - (a - b) \implies e \mid 2b$$

This implies  $e$  is a common divisor of  $2a$  and  $2b$ . So  $e \leq \gcd(2a, 2b)$ .  $\square$

**Problem 2** (Division Algorithm). Let  $c$  and  $d$  be integers, not both 0. If  $q$  and  $r$  are integers such as  $c = dq + r$ , then  $\gcd(c, d) = \gcd(d, r)$ .

**Proof** Let  $m = \gcd(c, d)$  and  $n = \gcd(d, r)$ . To prove  $m = n$ , we will show  $m \leq n$  and  $n \leq m$ .

(i) Show  $n \leq m$

Since  $n = \gcd(d, r)$ ,  $n \mid d$  and  $n \mid r$ . There exists integers  $x$  and  $y$  such that  $d = nx$  and  $r = ny$ .

From  $c = dq + r$ , we have  $c = (nx)q + ny = n(xq + y)$  thus  $n \mid c$ .  $n$  is a common divisor of  $c$  and  $d$ , so  $n \leq \gcd(c, d)$ . Hence  $n \leq m$ .

(ii) Show  $m \leq n$

This is left as an exercise.  $\square$

**Problem 3** (Euclid’s Lemma). Let  $a, b, c$  be any integers. If  $a \mid bc$  and  $\gcd(a, b) = 1$ , then  $a \mid c$ .

**Proof** Since  $a \mid bc$ ,  $bc = ak$  for some  $k \in \mathbb{Z}$ .

Since  $\gcd(a, b) = 1$ ,

$$\begin{aligned} ax + by &= 1 \quad \text{for some } x, y \in \mathbb{Z} \\ cax + cby &= c \\ acx + ak y &= c \\ a(cx + ky) &= c \end{aligned}$$

thus  $a \mid c$ .  $\square$

**Problem 4.** Let  $a$  and  $b$  be integers, not both 0. Show that  $\gcd(a, b)$  is the smallest possible positive linear combination of  $a$  and  $b$ . (i.e. There is no positive integer  $c < \gcd(a, b)$  such that  $c = ax + by$  for some integers  $x$  and  $y$ .)

**Proof** Prove by contradiction.

Suppose there is a positive integer  $c < \gcd(a, b)$  such that  $c = ax + by$  for some integers  $x$  and  $y$ .

Let  $d = \gcd(a, b)$ . Then  $d \mid a$  and  $d \mid b$ , and hence  $d \mid ax + by$ . This means  $d \mid c$ .

Since  $c$  is positive, this implies  $\gcd(a, b) = d \leq c$ . This contradicts  $c < \gcd(a, b)$ .

Hence we conclude that there is no positive integer  $c < \gcd(a, b)$  such that  $c = ax + by$  for some integers  $x$  and  $y$ .  $\square$

**Problem 5.** Use the Unique Factorisation Theorem to prove that, if a positive integer  $n$  is not a perfect square, then  $\sqrt{n}$  is irrational.

[The Unique Factorisation Theorem states that every integer  $n > 1$  has a unique standard factored form, i.e. there is exactly one way to express  $n = p_1^{k_1} p_2^{k_2} \cdots p_t^{k_t}$  where  $p_1 < p_2 < \cdots < p_t$  are distinct primes and  $k_1, k_2, \dots, k_t$  are some positive integers.]

**Proof** Prove by contradiction.

Suppose  $n$  is not a perfect square and  $\sqrt{n}$  is rational.

Then  $\sqrt{n} = \frac{a}{b}$  for some integers  $a$  and  $b$ . Squaring both sides and clearing denominator gives

$$nb^2 = a^2. \quad (*)$$

Consider the standard factored forms of  $n$ ,  $a$  and  $b$ :

$$\begin{aligned} n &= p_1^{k_1} p_2^{k_2} \cdots p_t^{k_t} \\ a &= q_1^{e_1} q_2^{e_2} \cdots q_u^{e_u} \implies a^2 = q_1^{2e_1} q_2^{2e_2} \cdots q_u^{2e_u} \\ b &= r_1^{f_1} r_2^{f_2} \cdots r_v^{f_v} \implies b^2 = r_1^{2f_1} r_2^{2f_2} \cdots r_v^{2f_v} \end{aligned}$$

i.e. the powers of primes in the standard factored form of  $a^2$  and  $b^2$  are all even integers.

This means the powers  $k_i$  of primes  $p_i$  in the standard factored form of  $n$  are also even by Unique Factorisation Theorem (UFT):

Note that all  $p_i$  appear in the standard factored form of  $a^2$  with even power  $2c_i$ , because of (\*). By UFT,  $p_i$  must also appear in the standard factored form of  $nb^2$  with the same even power  $2c_i$ .

If  $p_i \nmid b$ , then  $k_i = 2c_i$  which is even. If  $p_i \mid b$ , then  $p_i$  will appear in  $b^2$  with even power  $2d_i$ . So  $k_i + 2d_i = 2c_i$ , and hence  $k_i = 2(c_i - d_i)$ , which is again even.

$$\text{Hence } n = p_1^{k_1} p_2^{k_2} \cdots p_t^{k_t} = \left( p_1^{\frac{k_1}{2}} p_2^{\frac{k_2}{2}} \cdots p_t^{\frac{k_t}{2}} \right)^2.$$

Since  $\frac{k_i}{2}$  are all integers,  $p_1^{\frac{k_1}{2}} p_2^{\frac{k_2}{2}} \cdots p_t^{\frac{k_t}{2}}$  is an integer and  $n$  is a perfect square. This contradicts the given hypothesis that  $n$  is not a perfect square.

So we conclude that when a positive integer  $n$  is not a perfect square, then  $\sqrt{n}$  is irrational.  $\square$

**Problem 6** (Sieve of Eratosthenes). If  $p > 1$  is an integer and  $n \mid p$  for each integer  $n$  for which  $2 \leq n \leq \sqrt{p}$ , then  $p$  is prime.

**Proof** Prove by contrapositive.

Suppose that  $p$  is not prime, so it factors as  $p = mn$  for  $1 < m, n < p$ .

Observe that it is not the case that both  $m > \sqrt{p}$  and  $n > \sqrt{p}$ , because if this were true the inequalities would multiply to give  $mn > \sqrt{p}\sqrt{p} = p$ , which contradicts  $p = mn$ .

Therefore  $m \leq \sqrt{p}$  or  $n \leq \sqrt{p}$ . Without loss of generality, say  $n \leq \sqrt{p}$ . Then the equation  $p = mn$  gives  $n \mid p$ , with  $1 < n \leq \sqrt{p}$ . Hence it is not true that  $n \nmid p$  for each integer  $n$  for which  $2 \leq n \leq \sqrt{p}$ .  $\square$

**Problem 7** (Euclid's proof). There are infinitely many primes.

**Proof** Prove by contradiction.

Suppose otherwise, that the list of primes is finite. Let  $p_1, \dots, p_r$  be our finite list of primes. We want to show this is not the full list of the primes.

Consider the number

$$N = p_1 \cdots p_r + 1.$$

Since  $N > 1$ , it has a prime factor  $p$ . The prime  $p$  cannot be any of  $p_1, \dots, p_r$  since  $N$  has remainder 1 when divided by each  $p_i$ . Therefore  $p$  is a prime not on our list, so the set of primes cannot be finite.  $\square$

**Problem 8.** If  $n$  is an integer, prove that 3 divides  $n^3 - n$ .

**Proof** Prove by cases. This is done by partitioning  $\mathbb{Z}$  according to remainders when divided by  $d$  (i.e. equivalence classes).

We prove the three cases:  $n = 3k$ ,  $n = 3k + 1$ , and  $n = 3k + 2$ .

**Case 1:**  $n = 3k$  for some integer  $k$

Then

$$n^3 - n = (3k)^3 - (3k) = 3(9k^3 - k).$$

Since  $9k^3 - k$  is an integer,  $3 \mid n^3 - n$ .

**Case 2:**  $n = 3k + 1$  for some integer  $k$

Then

$$n^3 - n = (3k + 1)^3 - (3k + 1) = 3(9k^3 + 9k^2 + 2k).$$

Since  $9k^3 + 9k^2 + 2k$  is an integer,  $3 \mid n^3 - n$ .

**Case 3:**  $n = 3k + 2$  for some integer  $k$

The proof is similar and shall be left as an exercise. □

**Problem 9.** Prove that for every pair of irrational numbers  $p$  and  $q$  such that  $p < q$ , there is an irrational  $x$  such that  $p < x < q$ .

**Proof** Consider the average of  $p$  and  $q$ :  $p < \frac{p+q}{2} < q$ .

If  $\frac{p+q}{2}$  is irrational, take  $x = \frac{p+q}{2}$  and we are done.

If  $\frac{p+q}{2}$  is rational, call it  $r$ , take the average of  $p$  and  $r$ :  $p < \frac{p+r}{2} < r < q$ . Since  $p$  is irrational and  $r$  is rational,  $\frac{p+r}{2}$  is irrational. In this case, we take  $x = \frac{3p+q}{4}$ .  $\square$

**Problem 10.** Given  $n$  real numbers  $a_1, a_2, \dots, a_n$ . Show that there exists an  $a_i$  ( $1 \leq i \leq n$ ) such that  $a_i$  is greater than or equal to the mean (average) value of the  $n$  numbers.

**Proof** Prove by contradiction.

Let  $\bar{a}$  denote the mean value of the  $n$  given numbers. Suppose  $a_i < \bar{a}$  for all  $a_i$ . Then

$$\bar{a} = \frac{a_1 + a_2 + \dots + a_n}{n} < \frac{\bar{a} + \bar{a} + \dots + \bar{a}}{n} = \frac{n\bar{a}}{n} = \bar{a}.$$

We derive  $\bar{a} < \bar{a}$ , which is a contradiction.

Hence there must be some  $a_i$  such that  $a_i \geq \bar{a}$ .  $\square$

**Problem 11.** Prove that the following statement is false: there is an irrational number  $a$  such that for all irrational number  $b$ ,  $ab$  is rational.

**Thought process:** prove the negation of the statement: for every irrational number  $a$ , there is an irrational number  $b$  such that  $ab$  is irrational.

**Proving technique:** constructive proof (note that we can consider multiple cases and construct more than one  $b$ )

**Proof** Given an irrational number  $a$ , let us consider  $\frac{\sqrt{2}}{a}$ .

**Case 1:**  $\frac{\sqrt{2}}{a}$  is irrational.

Take  $b = \frac{\sqrt{2}}{a}$ . Then  $ab = \sqrt{2}$  which is irrational.

**Case 2:**  $\frac{\sqrt{2}}{a}$  is rational.

Then the reciprocal  $\frac{a}{\sqrt{2}}$ . Since  $\sqrt{6}$  is irrational, the product  $\left(\frac{a}{\sqrt{2}}\right)\sqrt{6} = a\sqrt{3}$  is irrational. Take  $b = \sqrt{3}$ , which is irrational. Then  $ab = a\sqrt{3}$  which is irrational.  $\square$

**Problem 12.** Prove that there are infinitely many prime numbers that are congruent to 3 modulo 4.

**Proof** Prove by contradiction.

Suppose there are only finitely many primes that are congruent to 3 modulo 4. Let  $p_1, p_2, \dots, p_m$  be the list of all the primes that are congruent to 3 modulo 4.

We construct an integer  $M$  by  $M = (p_1 p_2 \dots p_m)^2 + 2$ .

We have the following observation:

- (i)  $M \equiv 3 \pmod{4}$ .
- (ii) Every  $p_i$  divides  $M - 2$ .
- (iii) None of the  $p_i$  divides  $M$ . [Otherwise, together with (ii), this will imply  $p_i$  divides 2, which is impossible.]

- (iv)  $M$  is not a prime number. [Otherwise, by (i),  $M$  is a prime number congruent to 3 modulo 4. But  $M \neq p_i$  for all  $1 \leq i \leq m$ . This contradicts the assumption that  $p_1, p_2, \dots, p_m$  are all the prime numbers congruent to 3 modulo 4.]

From the above discussion, we know that  $M$  is a composite number by (iv). So it has a prime factorization  $M = q_1 q_2 \cdots q_k$ .

Since  $M$  is odd, all these prime factors  $q_j$  must be odd, and hence  $q_j$  must be congruent to either 1 or 3 modulo 4.

By (iii),  $q_j$  cannot be any of the  $p_i$ . So all  $q_j$  must be congruent to 1 modulo 4. Then  $M$ , which is the product of  $q_j$ , must also be congruent to 1 modulo 4.

This contradicts (i) that  $M$  is congruent to 3 modulo 4.

Hence we conclude that there must be infinitely many primes that are congruent to 3 modulo 4.  $\square$

**Problem 13.** Prove that, for any positive integer  $n$ , there is a perfect square  $m^2$  ( $m$  is an integer) such that  $n \leq m^2 \leq 2n$ .

**Proof** Prove by contradiction.

Suppose otherwise, that  $n > m^2$  and  $(m+1)^2 > 2n$  so that there is no square between  $n$  and  $2n$ , then

$$(m+1)^2 > 2n > 2m^2.$$

Since we are dealing with integers and the inequalities are strict, we get

$$(m+1)^2 \geq 2m^2 + 2$$

which simplifies to

$$0 \geq m^2 - 2m + 1 = (m-1)^2$$

The only value for which this is possible is  $m = 1$ , but you can eliminate that easily enough.  $\square$

**Problem 14.** Prove that for every positive integer  $n \geq 4$ ,

$$n! > 2^n.$$

**Proof** Let  $P(n) : n! > 2^n$

**Base case:**  $P(4)$

LHS:  $4! = 4 \times 3 \times 2 \times 1 = 24$ , RHS:  $2^4 = 16 < 24$

So  $P(4)$  is true.

**Inductive step:**  $P(k) \implies P(k+1)$  for all  $k \in \mathbb{Z}_{\geq 4}^+$

$$\begin{aligned} k! &> 2^k \\ (k+1)k! &> 2^k(k+1) \\ &> 2^k 2 \quad \text{since from } k \geq 4, k+1 \geq 5 > 2 \\ &= 2^{k+1} \end{aligned}$$

hence proven  $P(k) \implies P(k+1)$  for integers  $k \geq 4$ .

By PMI, we have proven  $P(n)$  for all integers  $n \geq 4$ .  $\square$

**Problem 15** (H2FM TJC 2023). Prove by mathematical induction, for  $n \geq 2$ ,

$$\sqrt[n]{n} < 2 - \frac{1}{n}.$$

**Proof** Let  $P(n) : \sqrt[n]{n} < 2 - \frac{1}{n}$  for  $n \geq 2$ .

**Base case:**  $P(2)$

When  $n = 2$ ,  $\sqrt{2} = 1.41 \dots < 2 - \frac{1}{2} = 1.5$  which is true. Hence  $P(2)$  is true.

**Inductive step:**  $P(k) \implies P(k+1)$  for all  $k \in \mathbb{Z}_{\geq 2}^+$

Assume  $P(k)$  is true for  $k \geq 2, k \in \mathbb{Z}^+$ , i.e.

$$\sqrt[k]{k} < 2 - \frac{1}{k} \implies k < \left(2 - \frac{1}{k}\right)^k$$

We want to prove that  $P(k+1)$  is true, i.e.

$$k+1 < \left(2 - \frac{1}{k+1}\right)^{k+1}$$

Since  $k > 2$ , we have

$$\begin{aligned} \left(2 - \frac{1}{k+1}\right)^{k+1} &> \left(2 - \frac{1}{k}\right)^{k+1} \quad \because k > 2 \\ &= \left(2 - \frac{1}{k}\right)^k \left(2 - \frac{1}{k}\right) \\ &> k \left(2 - \frac{1}{k}\right) \quad [\text{by inductive hypothesis}] \\ &= 2k - 1 = k + k - 1 > k - 1 \because k > 2 \end{aligned}$$

Hence  $P(k+1)$  is true.

Since  $P(2)$  is true and  $P(k) \implies P(k+1)$ , by mathematical induction  $P(n)$  is true. □

**Problem 16.** Prove that for all integers  $n \geq 3$ ,

$$\left(1 + \frac{1}{n}\right)^n < n$$

**Proof Base case:**  $P(3)$

On the LHS,  $\left(1 + \frac{1}{3}\right)^3 = \frac{64}{27} = 2\frac{10}{27} < 3$ . Hence  $P(3)$  is true.

**Inductive step:**  $P(k) \implies P(k+1)$  for all  $k \in \mathbb{Z}_{\geq 3}^+$

Our inductive hypothesis is

$$\left(1 + \frac{1}{k}\right)^k < k$$

Multiplying both sides by  $\left(1 + \frac{1}{k}\right)$  (to get a  $k+1$  in the power),

$$\left(1 + \frac{1}{k}\right)^k \left(1 + \frac{1}{k}\right) = \left(1 + \frac{1}{k}\right)^{k+1} < k \left(1 + \frac{1}{k}\right) = k+1$$

Since  $k < k+1 \iff \frac{1}{k} > \frac{1}{k+1}$ ,

$$\left(1 + \frac{1}{k}\right)^{k+1} > \left(1 + \frac{1}{k+1}\right)^{k+1}$$

The rest of the proof follows easily. □

A sequence of integers  $F_i$ , where integer  $1 \leq i \leq n$ , is called the *Fibonacci sequence* if and only if it is defined recursively by  $F_1 = 1, F_2 = 1, F_n = F_{n-1} + F_{n-2}$  for  $n > 2$ .

**Problem 17.** Let  $F_i$  be the Fibonacci sequence. Prove that  $3 \nmid n$  if and only if  $F_n$  is odd.

**Proof Forward direction:**  $3 \nmid n \implies F_n$  is odd

**Backward direction:**  $F_n$  is odd  $\implies 3 \nmid n$  (We prove the contrapositive:  $3 \mid n \implies F_n$  is even)

Hence we only need to prove the following via PMI:

- $(\forall n \in \mathbb{Z}^+ \text{ and } 3 \nmid n), F_n \text{ is odd}$

**Base case:**  $P(1), P(2)$

**Inductive step:**  $P(k) \implies P(k+3)$  for all  $k \geq 1$

- $(\forall n \in \mathbb{Z}^+ \text{ and } 3 \mid n), F_n \text{ is even}$

**Base case:**  $P(3)$

**Inductive step:**  $P(k) \implies P(k+3)$  for all  $k \geq 3$

[Note that we have partitioned the domain into two.]

Hence to show  $\forall n \in \mathbb{Z}^+ P(n)$ ,

**Base case:**  $P(1), P(2), P(3)$

**Inductive step:**  $\forall k \in \mathbb{Z}^+ P(k) \implies P(k+3)$

□

**Problem 18.** Let  $a_i$  where integer  $1 \leq i \leq n$  be a sequence of integers defined recursively by initial conditions  $a_1 = 1$ ,  $a_2 = 1$ ,  $a_3 = 3$  and the recurrence relation  $a_n = a_{n-1} + a_{n-2} + a_{n-3}$  for  $n > 3$ .

For all  $n \in \mathbb{Z}^+$ , prove that

$$a_n \leq 2^{n-1}.$$

**Proof** Let  $P(n) : a_n \leq 2^{n-1}$ .

Given the recurrence relation, it could be possible to use  $P(k), P(k+1), P(k+2)$  to prove  $P(k+3)$  for all  $k \in \mathbb{Z}^+$ .

**Base case:**  $P(1), P(2), P(3)$

$P(1) : a_1 = 1 \leq 2^{1-1} = 1$  is true.

$P(2) : a_2 = 1 \leq 2^{2-1} = 2$  is true.

$P(3) : a_3 = 3 \leq 2^{3-1} = 4$  is true.

**Inductive step:**  $P(k) \wedge P(k+1) \wedge P(k+2) \implies P(k+3)$  for all  $k \in \mathbb{Z}^+$

By inductive hypothesis, for  $k \in \mathbb{Z}^+$  we have  $a_k \leq 2^k, a_{k+1} \leq 2^{k+1}, a_{k+2} \leq 2^{k+2}$ .

$$\begin{aligned} a_{k+3} &= a_k + a_{k+1} + a_{k+2} && \text{[start from recurrence relation]} \\ &\leq 2^k + 2^{k+1} + 2^{k+2} && \text{[use inductive hypothesis]} \\ &= 2^k(1 + 2 + 2^2) \\ &< 2^k(2^3) && \text{[approximation, since } 1 + 2 + 2^2 < 2^3\text{]} \\ &= 2^{k+3} \end{aligned}$$

which is precisely  $P(k+3) : a_{k+3} \leq 2^{k+3}$ . □



**Problem 19** (Bézout's lemma). Let  $a$  and  $b$  be integers, not both 0. Prove that  $\gcd(a, b) = ax_0 + by_0$  for some integers  $x_0$  and  $y_0$ .

**Solution** Given  $a$  and  $b$ , consider the set

$$S = \{z \in \mathbb{Z} \mid z > 0; \exists x, y \in \mathbb{Z}, z = ax + by\}.$$

$S$  satisfies the conditions of well-ordering principle, and hence has a smallest element  $c = ax_0 + by_0$ . We want to show that (i)  $c$  is a common divisor of  $a$  and  $b$ ; (ii)  $c = \gcd(a, b)$ .

(i)  $c$  is a common divisor of  $a$  and  $b$

Suppose  $c \nmid a$ . By quotient-remainder theorem,  $a = cq + r$  where  $0 < r < c$ .

Then

$$a = (ax_0 + by_0)q + r \implies r = a - (ax_0 + by_0)q \implies r = a(1 - x_0q) - b(y_0q)$$

So  $r$  is an element in  $S$ , and  $r < c$ . This contradicts the minimality of  $c$  in  $S$ . Hence  $c \mid a$ . Then  $a = (ax_0 + by_0)q + r$ .

Similarly,  $c \mid b$ .

(ii)  $c = \gcd(a, b)$

Suppose otherwise, that  $c$  is not the greatest common divisor of  $a$  and  $b$ .

Let there exists some  $d > c$  which satisfies  $d \mid a$  and  $d \mid b$ .

Then  $d \mid (ax + by)$  for any  $x$  and  $y$ . So  $d$  divides all elements in  $S$ . In particular,  $d \mid c$ , which means  $d \leq c$ , a contradiction.

Hence  $c = \gcd(a, b)$ .

This concludes the proof that  $\gcd(a, b) = ax_0 + by_0$  for some integers  $x_0$  and  $y_0$ . □

**Problem 20** (Wilson's Theorem). Let  $p$  be a prime number. Prove that  $(p-1)! + 1$  is divisible by  $p$ .

**Proof** We first prove the uniqueness of inverse modulo  $p$ : for each  $x \in Q = \{1, 2, \dots, p-1\}$  for some prime  $p$ , there is precisely one integer  $y$  such that  $xy \equiv 1 \pmod{p}$ . **Proof** Suppose otherwise, that there are two distinct inverses for  $x$  modulo  $p$ ; that is,  $xy_1 \equiv 1 \pmod{p}$  and  $xy_2 \equiv 1 \pmod{p}$ . Then  $x(y_1 - y_2) \equiv 0 \pmod{p}$ . Since  $x \nmid p$ , by Euclid's lemma,  $y_1 \equiv y_2 \pmod{p}$  so  $y_1 = y_2 + kp$  for some integer  $k$ . But we know that  $0 \leq y_1, y_2 < p$ , so  $kp = y_1 - y_2$ ,  $0 \leq kp < p$  thus  $k = 0$ . Hence  $y_1 = y_2$ .  $\square$

If  $y \neq x$ , we can pair up elements of  $Q$  such that their product is congruent to 1 modulo  $p$ .

If  $y = x$ , then  $x^2 \equiv 1 \pmod{p}$ . Thus

$$p \mid x^2 = 1 \implies p \mid (x+1)(x-1) \implies p \mid x+1 \text{ or } p \mid x-1 \implies x \equiv \pm 1 \pmod{p}$$

which is  $1^2 \equiv 1 \pmod{p}$  and  $(p-1)^2 \equiv 1 \pmod{p}$ . So aside 1 and  $p-1$ , all other elements can be paired up. Hence,

$$\begin{aligned} (p-1)! + 1 &\equiv 1(p-1) + 1 \pmod{p} \\ &\equiv p-1 + 1 \pmod{p} \\ &\equiv p \pmod{p} \end{aligned}$$

Hence  $(p-1)! + 1$  is divisible by  $p$ .  $\square$

**Problem 21.** For  $m, n \in \mathbb{N}$ , prove that

$$F_{n+m+1} = F_n F_m + F_{n+1} F_{m+1}.$$

**Proof** For  $n \in \mathbb{N}$ , take  $P(n) : F_{n+m+1} = F_n F_m + F_{n+1} F_{m+1}$  for all  $m \in \mathbb{N}$  in the cases  $k = n$  and  $k = n + 1$ .

So we are using induction to progress through  $n$  and dealing with  $m$  simultaneously at each stage.

To verify  $P(0)$ , we note that

$$F_{m+1} = F_0 F_m + F_1 F_{m+1}$$

and

$$F_{m+2} = F_1 F_m + F_2 F_{m+1}$$

for all  $m$ , as  $F_0 = 0$  and  $F_1 = F_2 = 1$ .

For the inductive step we assume  $P(n)$ , i.e. that for all  $m \in \mathbb{N}$ ,  $F_{n+m+1} = F_n F_m + F_{n+1} F_{m+1}$ ,  $F_{n+m+2} = F_{n+1} F_m + F_{n+2} F_{m+1}$ . To prove  $P(n+1)$  it remains to show that for all  $m \in \mathbb{N}$ ,

$$F_{n+m+3} = F_{n+2} F_m + F_{n+3} F_{m+1}.$$

From our  $P(n)$  assumptions and the definition of the Fibonacci numbers, LHS of (5) =  $F_{n+m+3} = F_{n+m+2} + F_{n+m+1} = F_{n+1} F_m + F_{n+2} F_{m+1} + F_n F_m + F_{n+1} F_{m+1} = (F_n + F_{n+1}) F_m + (F_{n+1} + F_{n+2}) F_{m+1} = F_{n+2} F_m + F_{n+3} F_{m+1} = \text{RHS of (5)}.$   $\square$

# 2 Set Theory

## §2.1 Basics

### §2.1.1 Notation

You should, by now, be familiar with the following definitions and notation:

- A **set**  $S$  can be loosely defined as a collection of objects.
- For a set  $S$ , we write  $x \in S$  to mean that  $x$  is an **element** of  $S$ , and  $x \notin S$  if otherwise.
- A set can be defined in terms of some property  $P(x)$  that the elements  $x \in S$  satisfy, denoted by the following **set builder notation**:

$$\{x \in S \mid P(x)\}$$

- Some basic sets (of numbers) you should be familiar with:
  - $\mathbb{N} = \{0, 1, 2, 3, \dots\}$  denotes the natural numbers (non-negative integers).
  - $\mathbb{Z} = \{\dots, -2, -1, 0, 1, 2, \dots\}$  denotes the integers.
  - $\mathbb{Q} = \{\frac{p}{q} \mid p, q \in \mathbb{Z}, q \neq 0\}$  denotes the rational numbers.
  - $\mathbb{R}$  denotes the real numbers, which can be expressed in terms of decimal expansion.
  - $\mathbb{C} = \{x + yi \mid x, y \in \mathbb{R}\}$  denotes the of complex numbers.
- The **empty set** is the set with no elements, denoted by  $\emptyset$ .
- $A$  is a **subset** of  $B$  if every element of  $A$  is in  $B$ , denoted by  $A \subseteq B$ .

$$A \subseteq B \iff \forall x, x \in A \implies x \in B$$

$\subseteq$  is transitive, i.e. if  $A \subseteq B$  and  $B \subseteq C$ , then  $A \subseteq C$ . **Proof** Let  $x \in A$ . Since  $A \subseteq B$  and  $x \in A$ ,  $x \in B$ . Since  $B \subseteq C$  and  $x \in B$ ,  $x \in C$ . Hence  $A \subseteq C$ .  $\square$

$A$  is a **proper subset** of  $B$  if  $A \subseteq B$  and  $A \neq B$ , denoted by  $A \subset B$ .

Using this definition, we have the relationship

$$\mathbb{N} \subset \mathbb{Z} \subset \mathbb{Q} \subset \mathbb{R}$$

- $A$  and  $B$  are **equal** if and only if they contain the same elements, denoted by  $A = B$ .  
To prove that  $A$  and  $B$  are equal, we simply need to prove that  $A \subseteq B$  and  $A \supseteq B$ .

**Proof** We have

$$\begin{aligned} A = B &\iff (\forall x)[x \in A \iff x \in B] \\ &\iff (\forall x)[(x \in A \implies x \in B) \wedge (x \in B \implies x \in A)] \\ &\iff \{(\forall x)[x \in A \implies x \in B]\} \wedge \{(\forall x)[x \in B \implies x \in A]\} \\ &\iff (A \subseteq B) \wedge (B \subseteq A) \end{aligned}$$

$\square$

- Some frequently occurring subsets of the real numbers are known as **intervals**, which can be visualised as sections of the real line:

- Open interval

$$(a, b) = \{x \in \mathbb{R} \mid a < x < b\}$$

- Closed interval

$$[a, b] = \{x \in \mathbb{R} \mid a \leq x \leq b\}$$

- Half open interval

$$(a, b] = \{x \in \mathbb{R} \mid a < x \leq b\}$$

- The **power set**  $\mathcal{P}(A)$  of  $A$  is the set of all subsets of  $A$  (including the set itself and the empty set).
- An **ordered pair** is denoted by  $(a, b)$ , where the order of the elements matters. Two pairs  $(a_1, b_1)$  and  $(a_2, b_2)$  are equal if and only if  $a_1 = a_2$  and  $b_1 = b_2$ .

Similarly, we have ordered triples  $(a, b, c)$ , quadruples  $(a, b, c, d)$  and so on. If there are  $n$  elements it is called an  $n$ -tuple.

- The **Cartesian product** of sets  $A$  and  $B$ , denoted by  $A \times B$ , is the set of all ordered pairs with the first element of the pair coming from  $A$  and the second from  $B$ :

$$A \times B = \{(a, b) \mid a \in A, b \in B\} \quad (2.1)$$

More generally, we define  $A_1 \times A_2 \times \cdots \times A_n$  to be the set of all ordered  $n$ -tuples  $(a_1, a_2, \dots, a_n)$ , where  $a_i \in A_i$  for  $1 \leq i \leq n$ . If all the  $A_i$  are the same, we write the product as  $A^n$ .

**Example 2.1.1.**  $\mathbb{R}^2$  is the Euclidean plane,  $\mathbb{R}^3$  is the Euclidean space, and  $\mathbb{R}^n$  is the  $n$ -dimensional Euclidean space.

$$\begin{aligned} \mathbb{R} \times \mathbb{R} &= \mathbb{R}^2 = \{(x, y) \mid x, y \in \mathbb{R}\} \\ \mathbb{R} \times \mathbb{R} \times \mathbb{R} &= \mathbb{R}^3 = \{(x, y, z) \mid x, y, z \in \mathbb{R}\} \\ \mathbb{R}^n &= \{(x_1, x_2, \dots, x_n) \mid x_1, x_2, \dots, x_n \in \mathbb{R}\} \end{aligned}$$

## §2.1.2 Algebra of Sets

Given  $A \subset S$  and  $B \subset S$ .

- The **union**  $A \cup B$  is the set consisting of elements that are in  $A$  or  $B$  (or both):

$$A \cup B = \{x \in S \mid x \in A \vee x \in B\}$$

- The **intersection**  $A \cap B$  is the set consisting of elements that are in both  $A$  and  $B$ :

$$A \cap B = \{x \in S \mid x \in A \wedge x \in B\}$$

$A$  and  $B$  are **disjoint** if both sets have no element in common:

$$A \cap B = \emptyset$$

More generally, we can take unions and intersections of arbitrary numbers of sets, even infinitely many. If we have a family of subsets  $\{A_i \mid i \in I\}$ , where  $I$  is an **indexing set**, we write

$$\bigcup_{i \in I} A_i = \{x \mid \exists i \in I (x \in A_i)\}$$

and

$$\bigcap_{i \in I} A_i = \{x \mid \forall i \in I (x \in A_i)\}$$

- The **complement** of  $A$ , denoted by  $A^c$  or  $A'$ , is the set containing elements that are not in  $A$ :

$$A^c = \{x \in S \mid x \notin A\}$$

- The **set difference**, or complement of  $B$  in  $A$ , denoted by  $A \setminus B$ , is the subset consisting of those elements that are in  $A$  and not in  $B$ :

$$A \setminus B = \{x \in A \mid x \notin B\}$$

Note that  $A \setminus B = A \cap B^c$ .

**Proposition 2.1.2** (Double Inclusion). Let  $A \subset S$  and  $B \subset S$ . Then

$$A = B \iff A \subseteq B \text{ and } B \subseteq A \quad (2.2)$$

**Proof** We prove both directions.

**Forward direction:**

If  $A = B$ , then every element in  $A$  is an element in  $B$ , so certainly  $A \subseteq B$ , and similarly  $B \subseteq A$ .

**Backward direction:**

Suppose  $A \subseteq B$ , and  $B \subseteq A$ . Then for every element  $x \in S$ , if  $x \in A$  then  $A \subseteq B$  implies that  $x \in B$ , and if  $x \notin A$  then  $B \subseteq A$  means  $x \notin B$ . So  $x \in A$  if and only if  $x \in B$ , and therefore  $A = B$ .  $\square$

**Proposition 2.1.3** (Distributive Laws). Let  $A \subset S$ ,  $B \subset S$  and  $C \subset S$ . Then

$$(A \cup B) \cap C = (A \cap C) \cup (B \cap C) \quad (2.3)$$

$$(A \cap B) \cap C = (A \cup C) \cap (B \cup C) \quad (2.4)$$

**Proof** For the first one, suppose  $x$  is in the LHS, that is  $x \in A \cup (B \cap C)$ . This means that  $x \in A$  or  $x \in B \cap C$  (or both). Thus either  $x \in A$  or  $x$  is in both  $B$  and  $C$  (or  $x$  is in all three sets). If  $x \in A$  then  $x \in A \cup B$  and  $x \in A \cup C$ , and therefore  $x$  is in the RHS. If  $x$  is in both  $B$  and  $C$  then similarly  $x$  is in both  $A \cup B$  and  $A \cup C$ . Thus every element of the LHS is in the RHS, which means we have shown  $A \cup (B \cap C) \subseteq (A \cup B) \cap (A \cup C)$ .

Conversely suppose that  $x \in (A \cup B) \cap (A \cup C)$ . Then  $x$  is in both  $A \cup B$  and  $A \cup C$ . Thus either  $x \in A$  or, if  $x \notin A$ , then  $x \in B$  and  $x \in C$ . Thus  $x \in A \cup (B \cap C)$ . Hence  $(A \cup B) \cap (A \cup C) \subseteq A \cup (B \cap C)$ .

By double inclusion,  $(A \cup B) \cap (A \cup C) = A \cup (B \cap C)$ .

The proof of the second one follows similarly and is left as an exercise.  $\square$

**Proposition 2.1.4** (De Morgan's Laws). Let  $A \subset S$  and  $B \subset S$ . Then

$$(A \cup B)^c = A^c \cap B^c \quad (2.5)$$

$$(A \cap B)^c = A^c \cup B^c \quad (2.6)$$

**Proof** For the first one, suppose  $x \in (A \cup B)^c$ . Then  $x$  is not in either  $A$  or  $B$ . Thus  $x \in A^c$  and  $x \in B^c$ , and therefore  $x \in A^c \cap B^c$ .

Conversely, suppose  $x \in A^c \cap B^c$ . Then  $x \notin A$  and  $x \notin B$ , so  $x$  is in neither  $A$  nor  $B$ , and therefore  $x \in (A \cup B)^c$ .

By double inclusion, the first result holds. The second result follows similarly and is left as an exercise.  $\square$

De Morgan's laws extend naturally to any number of sets, so if  $\{A_i \mid i \in I\}$  is a family of subsets of  $S$ , then

$$\left( \bigcap_{i \in I} A_i \right)^c = \bigcup_{i \in I} A_i^c \quad \text{and} \quad \left( \bigcup_{i \in I} A_i \right)^c = \bigcap_{i \in I} A_i^c$$

### Exercise 2.1.1

Prove the following:

1.  $(\bigcup_{i \in I} A_i) \cup B = \bigcup_{i \in I} (A_i \cup B)$
2.  $(\bigcap_{i \in I} A_i) \cup B = \bigcap_{i \in I} (A_i \cup B)$
3.  $(\bigcup_{i \in I} A_i) \cup (\bigcup_{j \in J} B_j) = \bigcup_{(i,j) \in I \times J} (A_i \cup B_j)$
4.  $(\bigcap_{i \in I} A_i) \cup (\bigcap_{j \in J} B_j) = \bigcap_{(i,j) \in I \times J} (A_i \cup B_j)$

**Exercise 2.1.2**

Let  $S \subset A \times B$ . Express the set  $A_S$  of all elements of  $A$  which appear as the first entry in at least one of the elements in  $S$ .  
( $A_S$  here may be called the projection of  $S$  onto  $A$ .)

### §2.1.3 Cardinality

Informally, the **cardinality** of  $S$ , denoted  $|S|$ , is a measure of its “size”. We will be able to give a nicer definition of cardinality later, once we have discussed bijections, but the following provides a recursive definition of the cardinality for a finite set:

**Definition 2.1.5** (Finiteness and the cardinality of a finite set). The empty set  $\emptyset$  is finite, with  $|\emptyset| = 0$ .  $S$  is **finite** with  $|S| = n + 1$ , if there exists  $s \in S$  such that  $|S \setminus \{s\}| = n$  for some  $n \in \mathbb{Z}^+$ . We call  $|S|$  the **cardinality** of  $S$ .

Any set that is not finite is said to be **infinite**.

It is not hard to see that this means that if  $S = \{x_1, x_2, \dots, x_n\}$ , and  $x_i \neq x_j$  whenever  $i \neq j$ , then  $|S| = n$ . Conversely, if  $|S| = n$  then  $S$  is a set with  $n$  elements.

**Proposition 2.1.6.** Let  $A$  and  $B$  be finite sets. Then  $|A \cup B| = |A| + |B| - |A \cap B|$ .

**Proof** The proof is left as an exercise. □

**Proposition 2.1.7** (Subsets of a finite set). If a set  $A$  is finite with  $|A| = n$ , then its power set has  $|\mathcal{P}(A)| = 2^n$ .

**Proof** We use induction. For the initial step, note that if  $|A| = 0$  then  $A = \emptyset$  has no elements, so there is a single subset  $\emptyset$ , and therefore  $|\mathcal{P}(A)| = 1 = 2^0$ .

Now suppose that  $n \geq 0$  and that  $|\mathcal{P}(S)| = 2^n$  for any set  $S$  with  $|S| = n$ . Let  $A$  be any set with  $|A| = n + 1$ . By definition, this means that there is an element  $a$  and a set  $A_0 = A \setminus \{a\}$  with  $|A_0| = n$ . Any subset of  $A$  must either contain the element  $a$  or not, so we can partition  $\mathcal{P}(A) = \mathcal{P}(A_0) \cup \{S \cup \{a\} \mid S \in \mathcal{P}(A_0)\}$ . These two sets are disjoint, and each of them has cardinality  $|\mathcal{P}(A_0)| = 2^n$  by the inductive hypothesis. Hence  $|\mathcal{P}(A)| = 2^n + 2^n = 2^{n+1}$ .

Thus, by induction, the result holds for all  $n$ . □

Another way to see this is through combinatorics: Consider the process of creating a subset. We can do this systematically by going through each of the  $|A|$  elements in  $A$  and making the yes/no decision whether to put it in the subset. Since there are  $|A|$  such choices, that yields  $2^{|A|}$  different combinations of elements and therefore  $2^{|A|}$  different subsets.



## §2.2 Relations

### §2.2.1 Definition

**Definition 2.2.1.**  $R$  is a **relation** between  $A$  and  $B$  if and only if  $R$  is a subset of the Cartesian product  $A \times B$ , i.e.  $R \subseteq A \times B$ .

$a \in A$  and  $b \in B$  are **related** if  $(a, b) \in R$ , denoted by  $aRb$ .

**Remark** A relation is a set of ordered pairs.

Visually speaking, a relation is uniquely determined by a simple bipartite graph over  $A$  and  $B$ . On the bipartite graph, this is usually represented by an edge between  $a$  and  $b$ .

**Definition 2.2.2.** A **binary relation** in  $A$  is a relation between  $A$  and itself, i.e.  $R \subseteq A \times A$ .

$A$  and  $B$  are the **domain** and **range** of  $R$  respectively, denoted by  $\text{dom } R$  and  $\text{ran } R$  respectively, if and only if  $A \times B$  is the smallest Cartesian product of which  $R$  is a subset.

**Example 2.2.3.** Given  $R = \{(1, a), (1, b), (2, b), (3, b)\}$ , then  $\text{dom } R = \{1, 2, 3\}$  and  $\text{ran } R = \{a, b\}$ .

In many cases we do not actually use  $R$  to write the relation because there is some other conventional notation:

**Example 2.2.4.**

- The “less than or equal to” relation  $\leq$  on the set of real numbers is  $\{(x, y) \in \mathbb{R}^2 \mid x \leq y\}$ . We write  $x \leq y$  if  $(x, y)$  is in this set.
- The “divides” relation  $\mid$  on  $\mathbb{N}$  is  $\{(m, n) \in \mathbb{N}^2 : m \text{ divides } n\}$ . We write  $m \mid n$  if  $(m, n)$  is in this set.
- For a set  $S$ , the “subset” relation  $\subseteq$  on  $\mathcal{P}(S)$  is  $\{(A, B) \in \mathcal{P}(S)^2 \mid A \subseteq B\}$ . We write  $A \subseteq B$  if  $(A, B)$  is in this set.

### §2.2.2 Properties of relations

Let  $A$  be a set,  $R$  a relation on  $A$ , and  $x, y, z \in A$ . We say that

- $R$  is **reflexive** if  $xRx$  for all  $x \in A$ ;
- $R$  is **symmetric** if  $xRy \implies yRx$ ;
- $R$  is **anti-symmetric** if  $xRy$  and  $yRx \implies x = y$ ;
- $R$  is **transitive** if  $xRy$  and  $yRz \implies xRz$ .

**Example 2.2.5** (Less than or equal to). The relation  $\leq$  on  $R$  is reflexive, anti-symmetric, and transitive, but not symmetric.

**Definition 2.2.6.** Any relation on  $A$  that is reflexive, anti-symmetric, and transitive is called a **partial order**, denoted by  $\leq$ . It is called a **total order** if for every  $x, y \in A$ , either  $xRy$  or  $yRx$  (or both).

**Example 2.2.7** (Less than). The relation  $<$  on  $R$  is not reflexive, symmetric, or anti-symmetric, but it is transitive.

**Example 2.2.8** (Not equal to). The relation  $\neq$  on  $R$  is not reflexive, anti-symmetric or transitive, but it is symmetric.

#### Exercise 2.2.1: Congruence modulo $n$

Let  $n \geq 2$  be an integer, and define  $R$  on  $\mathbb{Z}$  by saying  $aRb$  if and only if  $a - b$  is a multiple of  $n$ . Prove that  $R$  is reflexive, symmetric and transitive.

**Proof**

- Reflexivity: For any  $a \in \mathbb{Z}$  we have  $aRa$  as  $0$  is a multiple of  $n$ .
- Symmetry: If  $aRb$  then  $a - b = kn$  for some integer  $k$ . So  $b - a = -kn$ , and hence  $bRa$ .
- Transitivity: If  $aRb$  and  $bRc$  then  $a - b = kn$  and  $b - c = ln$  for integers  $k, l$ . So then  $a - c = (a - b) + (b - c) = (k + l)n$ , and hence  $aRc$ .

□

### §2.2.3 Equivalence relations, equivalence classes, and partitions

One important type of relation is an equivalence relation. An equivalence relation is a way of saying two objects are, in some particular sense, “the same”.

**Definition 2.2.9.** A binary relation  $R$  on  $A$  is an **equivalence relation** if it is reflexive, symmetric and transitive.

**Notation** We use the symbol  $\sim$  to denote the equivalence relation  $R$  in  $A \times A$ : whenever  $(a, b) \in R$  we denote  $a \sim b$ .

An equivalence relation provides a way of grouping together elements that can be viewed as being the same:

**Definition 2.2.10.** Given an equivalence relation  $\sim$  on a set  $A$ , and given  $x \in A$ , the **equivalence class** of  $x$ , denoted  $[x]$ , is the subset

$$[x] = \{y \in A \mid y \sim x\}$$

**Example 2.2.11** (Congruence modulo  $n$ ). For the equivalence relation of congruence modulo  $n$ , the equivalence class of 1 is the set  $1 = \{\dots, -n+1, 1, n+1, 2n+1, \dots\}$ ; that is, all the integers that are congruent to 1 modulo  $n$ .

Properties of equivalence classes:

- Every two equivalence classes are disjoint
- The union of equivalence classes form the entire set

You can translate these properties into the point of view from the elements: Every element belongs to one and only one equivalence class.

- No element belongs to two distinct classes
- All elements belong to an equivalence class

**Definition 2.2.12.** The **set of equivalence classes** (quotient sets) are the set of all equivalence classes, denoted by  $A/\sim$ .

Grouping the elements of a set into equivalence classes provides a partition of the set, which we define as follows:

**Definition 2.2.13.** A **partition** of a set  $A$  is a collection of subsets  $\{A_i \subseteq A \mid i \in I\}$ , where  $I$  is an indexing set, with the property that

- (i)  $A_i \neq \emptyset$  for all  $i \in I$  (that is, all the subsets are non-empty),
- (ii)  $\bigcup_{i \in I} A_i = A$  (that is, every member of  $A$  lies in one of the subsets),
- (iii)  $A_i \cap A_j = \emptyset$  for every  $i \neq j$  (that is, the subsets are disjoint).

The subsets are called the parts of the partition.

**Example 2.2.14** (Odd and even natural numbers).  $\{\{n \in \mathbb{N} \mid n \text{ is divisible by } 2\}, \{n \in \mathbb{N} \mid n+1 \text{ is divisible by } 2\}\}$  forms a partition of the natural numbers, into evens and odds.

## §2.3 Functions

### §2.3.1 Definition

**Definition 2.3.1.** Given two sets  $X$  and  $Y$ , a **function**  $f$  from  $X$  to  $Y$  is a mapping of every element of  $X$  to some element of  $Y$ , denoted by  $f : X \rightarrow Y$ .

$X$  and  $Y$  are known as the **domain** and **codomain** of  $f$  respectively.

**Remark** The definition requires that a unique element of the codomain is assigned for every element of the domain. For example, for a function  $f : \mathbb{R} \rightarrow \mathbb{R}$ , the assignment  $f(x) = \frac{1}{x}$  is not sufficient as it fails at  $x = 0$ . Similarly,  $f(x) = y$  where  $y^2 = x$  fails because  $f(x)$  is undefined for  $x < 0$ , and for  $x > 0$  it does not return a unique value; in such cases, we say the function is **ill-defined**. We are interested in the opposite; functions that are **well-defined**.

**Definition 2.3.2.** Given a function  $f : X \rightarrow Y$ , the **image** (or range) of  $f$  is

$$f(X) = \{f(x) \mid x \in X\} \subseteq Y$$

More generally, given  $A \subseteq X$ , the image of  $A$  under  $f$  is

$$f(A) = \{f(x) \mid x \in A\} \subseteq Y$$

Given  $B \subseteq Y$ , the **pre-image** of  $B$  under  $f$  is

$$f^{-1}(B) = \{x \mid f(x) \in B\} \subseteq X$$

**Remark** Beware the potentially confusing notation: for  $x \in X$ ,  $f(x)$  is a single element of  $Y$ , but for  $A \subseteq X$ ,  $f(A)$  is a set (a subset of  $Y$ ). Note also that  $f^{-1}(B)$  should be read as “the pre-image of  $B$ ” and not as “ $f$ -inverse of  $B$ ”; the pre-image is defined even if no inverse function exists (in which case  $f^{-1}$  on its own has no meaning; we discuss invertibility of a function below).

#### Exercise 2.3.1

Prove the following statements:

- (a)  $f(A \cup B) = f(A) \cup f(B)$
- (b)  $f(A_1 \cup \dots \cup A_n) = f(A_1) \cup \dots \cup f(A_n)$
- (c)  $f(\bigcup_{\lambda \in A} A_\lambda) = \bigcup_{\lambda \in A} f(A_\lambda)$
- (d)  $f(A \cap B) \subseteq f(A) \cap f(B)$
- (e)  $f^{-1}(f(A)) \supseteq A$
- (f)  $f(f^{-1}(A)) \subseteq A$
- (g)  $f^{-1}(A \cup B) = f^{-1}(A) \cup f^{-1}(B)$
- (h)  $f^{-1}(A \cap B) = f^{-1}(A) \cap f^{-1}(B)$
- (i)  $f^{-1}(A_1 \cup \dots \cup A_n) = f^{-1}(A_1) \cup \dots \cup f^{-1}(A_n)$
- (j)  $f^{-1}(\bigcup_{\lambda \in A} A_\lambda) = \bigcup_{\lambda \in A} f^{-1}(A_\lambda)$

If a function is defined on some larger domain than we care about, it may be helpful to restrict the domain:

**Definition 2.3.3** (Restriction). Given a function  $f : X \rightarrow Y$  and a subset  $A \subseteq X$ , the **restriction** of  $f$  to  $A$  is the map  $f|_A : A \rightarrow Y$  defined by  $f|_A(x) = f(x)$  for all  $x \in A$ .

The restriction is almost the same function as the original  $f$  – just the domain has changed.

Another rather trivial but nevertheless important function is the identity map:

**Definition 2.3.4** (Identity map). Given a set  $X$ , the **identity**  $\text{id}_X : X \rightarrow X$  is defined by  $\text{id}_X(x) = x$  for all  $x \in X$ .

**Notation** If the domain is unambiguous, the subscript may be removed.

### §2.3.2 Injectivity, Surjectivity, Bijectivity

**Definition 2.3.5.**  $f : X \rightarrow Y$  is **injective** if each element of  $Y$  has at most one element of  $X$  that maps to it.

$$\forall x_1, x_2 \in X, f(x_1) = f(x_2) \implies x_1 = x_2$$

**Proposition 2.3.6.** If  $f : X \rightarrow Y$  is injective and  $g : Y \rightarrow Z$  is injective, then  $g \circ f : X \rightarrow Z$  is injective.

**Proof** Let  $f : X \rightarrow Y$  and  $g : Y \rightarrow Z$  be arbitrary injective functions. We want prove that the function  $g \circ f : X \rightarrow Z$  is also injective.

To do so, we will prove  $\forall x, x' \in X$  that

$$(g \circ f)(x) = (g \circ f)(x') \implies x = x'$$

Suppose that  $(g \circ f)(x) = (g \circ f)(x')$ . Expanding out the definition of  $g \circ f$ , this means that  $g(f(x)) = g(f(x'))$ .

Since  $g$  is injective and  $g(f(x)) = g(f(x'))$ , we know  $f(x) = f(x')$ .

Similarly, since  $f$  is injective and  $f(x) = f(x')$ , we know that  $x = x'$ , as required.  $\square$

**Proposition 2.3.7.**  $f$  is injective if and only if for any set  $Z$  and any functions  $g_1, g_2 : Z \rightarrow X$  we have  $f \circ g_1 = f \circ g_2 \implies g_1 = g_2$ .

**Proof Forward direction:**

If  $f$  is injective, we ultimately wish to show that  $g_1 = g_2$ , so in order to do this we consider all possible inputs  $z \in Z$ , hoping to show that  $g_1(z) = g_2(z)$ .

But this is quite simple because we are given that  $f \circ g_1 = f \circ g_2$  and that  $f$  is injective, so

$$f \circ g_1(z) = f \circ g_2(z) \implies g_1(z) = g_2(z)$$

**Backward direction:**

We specifically pick  $Z = \{1\}$ , basically some random one-element set.

Then  $\forall x, y \in X$ , we define

$$\begin{aligned} g_1 : Z \rightarrow X, g_1(1) &= x \\ g_2 : Z \rightarrow X, g_2(1) &= y \end{aligned}$$

Then

$$f(x) = f(y) \implies f(g_1(1)) = f(g_2(1)) \implies g_1(1) = g_2(1) \implies x = y$$

$\square$

**Definition 2.3.8.**  $f : X \rightarrow Y$  is **surjective** if every element of  $Y$  is mapped to at least one element of  $X$ .

$$\forall y \in Y, \exists x \in X \text{ s.t. } f(x) = y$$

**Proposition 2.3.9.** If  $f : X \rightarrow Y$  is surjective and  $g : Y \rightarrow Z$  is surjective, then  $g \circ f : X \rightarrow Z$  is surjective.

**Proof** Let  $f : X \rightarrow Y$  and  $g : Y \rightarrow Z$  be arbitrary surjective functions. We want to prove that the function  $g \circ f : X \rightarrow Z$  is surjective.

To do so, we want to prove that for any  $z \in Z$ , there is some  $x \in X$  such that  $(g \circ f)(x) = z$ . Equivalently, we want to prove that for any  $z \in Z$ , there is some  $x \in X$  such that  $g(f(x)) = z$ .

Consider any  $z \in Z$ . Since  $g : Y \rightarrow Z$  is surjective, there is some  $y \in Y$  such that  $g(y) = z$ . Similarly, since  $f : X \rightarrow Y$  is surjective, there is some  $x \in X$  such that  $f(x) = y$ . This means that there is some  $x \in X$  such that  $g(f(x)) = g(y) = z$ , as required.  $\square$

**Proposition 2.3.10.**  $f$  is surjective if and only if for any set  $Z$  and any functions  $g_1, g_2 : Y \rightarrow Z$  we have  $g_1 \circ f = g_2 \circ f \implies g_1 = g_2$ .

**Proof Forward direction:**

First we suppose that  $f$  is surjective. Again, we wish to show that  $g_1 = g_2$ , so we need to consider every possible input  $y$  in  $Y$ . Then, since  $f$  is injective, we can always pick  $x \in X$  such that  $f(x) = y$ .

Then

$$g_1 \circ f = g_2 \circ f \implies g_1 \circ f(x) = g_2 \circ f(x) \implies g_1(y) = g_2(y)$$

On the other hand, if  $f$  is not surjective, then there exists  $y \in Y$  such that for all  $x \in X$  we have  $f(x) \neq y$ . We then aim to construct set  $Z$  and  $g_1, g_2 : Y \rightarrow Z$  such that

$$(i) \quad g_1(y) \neq g_2(y)$$

$$(ii) \quad \forall y' \neq y, g_1(y') = g_2(y')$$

Because if this is satisfied, then  $\forall x \in X$ , since  $f(x) \neq y$  we have from (ii) that  $g_1(f(x)) = g_2(f(x))$ ; thus  $g_1 \circ f = g_2 \circ f$ , and yet from (i) we have  $g_1 \neq g_2$ .

**Backward direction:**

We construct  $Z = Y \cup \{1, 2\}$  for some random  $1, 2 \notin Y$ .

Then we define

$$g_1 : Y \rightarrow Z, g_1(y) = 1, g_1(y') = y' \qquad g_2 : Y \rightarrow Z, g_2(y) = 2, g_2(y') = y'$$

Then when  $y$  is not in the image of  $f$ , these two functions will satisfy  $g_1 \circ f = g_2 \circ f$  but not  $g_1 = g_2$ .

So conversely, if for any set  $Z$  and any functions  $g_i : Y \rightarrow Z$  we have  $g_1 \circ f = g_2 \circ f \implies g_1 = g_2$ , such a value  $y$  that is in the codomain but not in the range of  $f$  cannot appear, and hence  $f$  must be surjective.  $\square$

**Definition 2.3.11.**  $f : X \rightarrow Y$  is **bijective** if it is both injective and surjective: each element of  $Y$  is mapped to a unique element of  $X$ .

### §2.3.3 Cardinality and Countable Sets

When do two sets have the same size? Cantor answered this question in the 1800s, stating that two sets have the same size when you can pair each element in one set with a unique element in the other.

**Definition 2.3.12** (Cardinality).  $X$  and  $Y$  have the same **cardinality** if there exists a bijection  $f : X \rightarrow Y$ , denoted by  $|X| = |Y|$ .

**Definition 2.3.13** (Cardinality of finite sets). The empty set  $\emptyset$  is finite and has cardinality  $|\emptyset| = 0$ . A non-empty set  $A$  is said to be **finite** and have cardinality  $|A| = n \in \mathbb{Z}^+$  if and only if there exists a bijection from  $A$  to the set  $\{1, 2, \dots, n\}$ .

**Remark** Note that for finite sets  $X$  and  $Y$ , a function  $f : X \rightarrow Y$  can only be **injective** if  $|Y| \geq |X|$ , since for any injective function the number of elements in the image  $f(X)$ , is equal to the number of elements in the domain, and  $f(X) \subseteq Y$ . In other words, the codomain of an injective function cannot be smaller than the domain.<sup>1</sup>

Similarly, a function  $f : X \rightarrow Y$  can only be **surjective** if  $|Y| \leq |X|$ . Hence if  $f$  is bijective, then  $|X| = |Y|$ ; that is, the domain and codomain of a bijection have equal cardinality. (These results hold true for infinite sets too, though less obviously).

**Theorem 2.3.14** (Cantor–Schröder–Bernstein). If  $|X| \leq |Y|$  and  $|Y| \leq |X|$  then  $|X| = |Y|$ .

**Definition 2.3.15** (Countably infinite). A set  $X$  is **countably infinite** if it has the same cardinality as the set  $\mathbb{Z}^+$ .

**Definition 2.3.16** (Countable). A set  $X$  is **countable** if it is either finite or infinitely countable.

**Example 2.3.17.** The set  $2\mathbb{Z}^+ = \{2n \mid n \in \mathbb{Z}^+\}$  is countably infinite, i.e.  $|2\mathbb{Z}^+| = |\mathbb{Z}^+|$ .

To prove this, define the function  $f : \mathbb{Z}^+ \rightarrow 2\mathbb{Z}^+$  as  $f(n) = 2n$ . Then,  $f$  is injective – if  $f(n) = f(m)$  then  $2n = 2m \implies n = m$ . Furthermore,  $f$  is surjective, as if  $m \in 2\mathbb{Z}^+$  then  $\exists n \in \mathbb{Z}^+$  such that  $m = 2n = f(n)$ .

**Example 2.3.18.**  $\mathbb{Z}$  is countable since we have a bijection  $f : \mathbb{Z}^+ \rightarrow \mathbb{Z}$  given by

$$f(k) = \begin{cases} \frac{k}{2} & \text{if } k \text{ is even} \\ \frac{1-k}{2} & \text{if } k \text{ is odd} \end{cases}$$

In other words,  $f(1) = 0, f(2) = 1, f(3) = -1, f(4) = 2, f(5) = -2, f(6) = -3, \dots$  where our function  $f$  stretches in both positive and negative directions.

**Theorem 2.3.19** (Cantor). For a set  $A$ ,  $|A| < |\mathcal{P}(A)|$ .

**Proof** Define the function  $f : A \rightarrow \mathcal{P}(A)$  by  $f(x) = \{x\}$ . Then,  $f$  is injective as  $\{x\} = \{y\} \implies x = y$ . Thus  $|A| \leq |\mathcal{P}(A)|$ . To finish the proof now all we need to show is that  $|A| \neq |\mathcal{P}(A)|$ . We will do so through contradiction. Suppose that  $|A| = |\mathcal{P}(A)|$ . Then, there exists a surjection  $g : A \rightarrow \mathcal{P}(A)$ . We define the set  $B$  as

$$B := \{x \in A \mid x \notin g(x)\} \in \mathcal{P}(A)$$

Since  $g$  is surjective, there exists a  $b \in A$  such that  $g(b) = B$ . There are two cases:

1.  $b \in B$ . Then  $b \notin g(b) = B \implies b \notin B$ .
2.  $b \notin B$ . Then  $b \in g(b) = B \implies b \in B$ .

In either case we obtain a contradiction. Thus,  $g$  is not surjective so  $|A| \neq |\mathcal{P}(A)|$ . □

**Corollary 2.3.20.** For all  $n \in \mathbb{N} \cup \{0\}$ ,  $n < 2^n$ .

**Proof** This can be easily proven through induction. □

<sup>1</sup>This is also referred to as the pigeonhole principle: if  $n$  letters are placed in  $m$  pigeonholes and  $n > m$ , then at least one hole must contain more than one letter; the non-injective function in that case is the assignment of pigeonholes to letters.



**Lemma 2.3.21.** If  $X$  is a countable set and  $B \subseteq A$  then  $B$  is countable.

**Lemma 2.3.22.** If  $\{A_1, A_2, \dots\}$  is a collection of countably many countable sets then the set  $\bigcup_{i=1}^{\infty} A_i$  is countable.

**Lemma 2.3.23.** If  $\{A_1, A_2, \dots, A_n\}$  is a collection of finitely many countable sets then the set  $A_1 \times \dots \times A_n$  is countable.

### §2.3.4 Composition of functions and invertibility

**Definition 2.3.24.** Given two functions  $f : X \rightarrow Y$  and  $g : Y \rightarrow Z$ , the **composition**  $g \circ f : X \rightarrow Z$  is defined by

$$(g \circ f)(x) = g(f(x)) \quad \forall x \in X.$$

The composition of functions is not commutative. However, composition is associative, as the following results shows:

**Proposition 2.3.25** (Associativity). Let  $f : X \rightarrow Y$ ,  $g : Y \rightarrow Z$ ,  $h : Z \rightarrow W$  be three functions. Then

$$f \circ (g \circ h) = (f \circ g) \circ h.$$

**Proof** Let  $x \in X$ . Then, by the definition of composition, we have

$$(f \circ (g \circ h))(x) = f((g \circ h)(x)) = f(g(h(x))) = (f \circ g)(h(x)) = ((f \circ g) \circ h)(x).$$

□

The following proposition addresses the extent to which composition of functions preserves injectivity and surjectivity:

**Proposition 2.3.26.** Let  $f : X \rightarrow Y$  and  $g : Y \rightarrow Z$  be functions.

- (i) If  $f$  and  $g$  are injective then so is  $g \circ f$ . Conversely, if  $g \circ f$  is injective, then  $f$  is injective, but  $g$  need not be.
- (ii) If  $f$  and  $g$  are surjective then so is  $g \circ f$ . Conversely, if  $g \circ f$  is surjective, then  $g$  is surjective, but  $f$  need not be.

**Proof** For the first part of (i), suppose  $(g \circ f)(x_1) = (g \circ f)(x_2)$  for some  $x_1, x_2 \in X$ . From the injectivity of  $g$  we know that  $g(f(x_1)) = g(f(x_2))$  implies  $f(x_1) = f(x_2)$ , and then from the injectivity of  $f$  we know that this implies  $x_1 = x_2$ . So  $g \circ f$  is injective.

For the second part of (i), suppose  $f(x_1) = f(x_2)$  for some  $x_1, x_2 \in X$ . Then applying  $g$  gives  $g(f(x_1)) = g(f(x_2))$ , and by the injectivity of  $g \circ f$  this means  $x_1 = x_2$ . So  $f$  is injective. To see that  $g$  need not be injective, a counterexample is  $X = Z = \{0\}$ ,  $Y = \mathbb{R}$ , with  $f(0) = 0$  and  $g(y) = 0$  for all  $y \in \mathbb{R}$ . □

Recalling that  $\text{id}_X$  is the identity map on  $X$ , we can define invertibility:

**Definition 2.3.27.** A function  $f : X \rightarrow Y$  is **invertible** if there exists a function  $g : Y \rightarrow X$  such that  $g \circ f = \text{id}_X$  and  $f \circ g = \text{id}_Y$ . The function  $g$  is the **inverse** of  $f$ , denoted by  $g = f^{-1}$ .

Note that directly from the definition, if  $f$  is invertible then  $f^{-1}$  is also invertible, and  $(f^{-1})^{-1} = f$ .

An immediate concern we might have is whether there could be multiple such functions  $g$ , in which case the inverse  $f^{-1}$  would not be well-defined. This is resolved by the following result:

**Proposition 2.3.28** (Uniqueness of inverse). If  $f : X \rightarrow Y$  is invertible then its inverse is unique.

**Proof** Let  $g_1$  and  $g_2$  be two functions for which  $g_i \circ f = \text{id}_X$  and  $f \circ g_i = \text{id}_Y$ . Using the fact that composition is associative, and the definition of the identity maps, we can write

$$g_1 = g_1 \circ \text{id}_Y = g_1 \circ (f \circ g_2) = (g_1 \circ f) \circ g_2 = \text{id}_X \circ g_2 = g_2$$

□

The following result shows how to invert the composition of invertible functions:

**Proposition 2.3.29.** Let  $f : X \rightarrow Y$  and  $g : Y \rightarrow Z$  be functions. If  $f$  and  $g$  are invertible, then  $g \circ f$  is invertible, and  $(g \circ f)^{-1} = f^{-1} \circ g^{-1}$ .

**Proof** Making repeated use of the fact that function composition is associative, and the definition of the inverses  $f^{-1}$  and  $g^{-1}$ , we note that

$$\begin{aligned}
 (f^{-1} \circ g^{-1}) \circ (g \circ f) &= ((f^{-1} \circ g^{-1}) \circ g) \circ f \\
 &= (f^{-1} \circ (g^{-1} \circ g)) \circ f \\
 &= (f^{-1} \circ \text{id}_Y) \circ f \\
 &= f^{-1} \circ f \\
 &= \text{id}_X
 \end{aligned}$$

and similarly,

$$\begin{aligned}
 (g \circ f) \circ (f^{-1} \circ g^{-1}) &= g \circ (f \circ (f^{-1} \circ g^{-1})) \\
 &= g \circ ((f \circ f^{-1}) \circ g^{-1}) \\
 &= g \circ (\text{id}_X \circ g^{-1}) \\
 &= g \circ g^{-1} \\
 &= \text{id}_Y
 \end{aligned}$$

which shows that  $f^{-1} \circ g^{-1}$  satisfies the properties required to be the inverse of  $g \circ f$ .  $\square$

The following result provides an important and useful criterion for invertibility:

**Theorem 2.3.30.** A function  $f : X \rightarrow Y$  is invertible if and only if it is bijective.

**Proof** We prove this in both directions.

**Forward direction:**

Suppose  $f$  is invertible, so it has an inverse  $f^{-1} : Y \rightarrow X$ . To show  $f$  is injective, suppose that for some  $x_1, x_2 \in X$  we have  $f(x_1) = f(x_2)$ . Then applying  $f^{-1}$  to both sides and noting that by definition  $f^{-1} \circ f = \text{id}_X$ , we see that  $x_1 = f^{-1}(f(x_1)) = f^{-1}(f(x_2)) = x_2$ . So  $f$  is injective. To show that  $f$  is surjective, let  $y \in Y$ , and note that  $f^{-1}(y) \in X$  has the property that  $f(f^{-1}(y)) = y$ . So  $f$  is surjective. Therefore  $f$  is bijective.

**Backward direction:**

Suppose that  $f$  is bijective, we aim to show that there is a well-defined  $g : Y \rightarrow X$  such that  $g \circ f = \text{id}_X$  and  $f \circ g = \text{id}_Y$ . Since  $f$  is surjective, we know that for any  $y \in Y$ , there is an  $x \in X$  such that  $f(x) = y$ . Furthermore, since  $f$  is injective, we know that this  $x$  is unique. So for each  $y \in Y$  there is a unique  $x \in X$  such that  $f(x) = y$ . This recipe provides a well-defined function  $g(y) = x$ , for which we have  $g(f(x)) = x$  for any  $x \in X$  and  $f(g(y)) = y$  for any  $y \in Y$ . So  $g$  satisfies the property required to be an inverse of  $f$  and therefore  $f$  is invertible.  $\square$

It is also possible to define left-inverse and right-inverse functions as functions that partially satisfy the definition of the inverse:

**Definition 2.3.31.** A function  $f : X \rightarrow Y$  is **left invertible** if there exists a function  $g : Y \rightarrow X$  such that  $g \circ f = \text{id}_X$ , and is **right invertible** if there exists a function  $h : Y \rightarrow X$  such that  $f \circ h = \text{id}_Y$ .

As may be somewhat apparent from the previous proof, being left- and right-invertible is equivalent to being injective and surjective, respectively. We leave this as an exercise to show.

### §2.3.5 Strictly Monotonic Functions / Increasing and Decreasing Functions

**Definition 2.3.32.** A function  $f$  is **strictly monotonic increasing** if for all  $x_1, x_2 \in D_f$ ,

$$x_2 > x_1 \iff f(x_2) > f(x_1).$$

**Definition 2.3.33.** A function  $f$  is **strictly monotonic decreasing** if for all  $x_1, x_2 \in D_f$ ,

$$x_2 > x_1 \iff f(x_2) < f(x_1).$$

**Definition 2.3.34.** A function  $f$  is **increasing** on an interval  $I$  if for all  $x_1, x_2 \in I$  with  $x_1 < x_2$ ,  $f(x_1) \leq f(x_2)$ .

**Definition 2.3.35.** A function  $f$  is **decreasing** on an interval  $I$  if for all  $x_1, x_2 \in I$  with  $x_1 < x_2$ ,  $f(x_1) \geq f(x_2)$ .

Locate roots of an equation: suppose  $f(x)$  is continuous in the interval  $[a, b]$ ,

- If  $f(a)$  and  $f(b)$  have *opposite* signs, i.e.  $f(a)f(b) < 0$ , then there is an odd number of real roots (counting repeated) in  $[a, b]$ .

Furthermore, if  $f$  is either strictly increasing or decreasing in  $[a, b]$ , then  $f(x) = 0$  has *exactly one* real root in  $[a, b]$ .

- If  $f(a)$  and  $f(b)$  have *same* signs, i.e.  $f(a)f(b) > 0$ , then there is an even number of roots (counting repeated) in  $[a, b]$ .

### §2.3.6 Convex and Concave Functions

**Definition 2.3.36.** A function  $f$  is **convex** if for all  $x_1, x_2 \in D_f$  and  $0 \leq t \leq 1$ , we have

$$f(tx_1 + (1-t)x_2) \leq tf(x_1) + (1-t)f(x_2).$$

Note that equality holds when  $x_1 = x_2$ .

**Definition 2.3.37.** A function  $f$  is **strictly convex** if for all  $x_1, x_2 \in D_f$  with  $x_1 \neq x_2$  and  $0 < t < 1$ , we have

$$f(tx_1 + (1-t)x_2) < tf(x_1) + (1-t)f(x_2).$$

**Definition 2.3.38.** A function  $f$  is **concave** if for all  $x_1, x_2 \in D_f$  and  $0 \leq t \leq 1$ , we have

$$f(tx_1 + (1-t)x_2) \geq tf(x_1) + (1-t)f(x_2).$$

Note that equality holds when  $x_1 = x_2$ .

**Definition 2.3.39.** A function  $f$  is **strictly concave** if for all  $x_1, x_2 \in D_f$  with  $x_1 \neq x_2$  and  $0 < t < 1$ , we have

$$f(tx_1 + (1-t)x_2) > tf(x_1) + (1-t)f(x_2).$$

### §2.3.7 Other Functions

#### Piecewise Functions

A function that has its domain divided into *separate partitions* and each partition of the domain given a different formula or rule is known as a **piecewise function**, i.e. a function defined “piece-wise”.

- Absolute Value Function

The **absolute function** of  $x$ , denoted by  $|x|$ , is the distance of  $x$  from 0 on the real number line. Distances are always non-negative, so for every  $x \in \mathbb{R}$  we have

$$|x| \geq 0.$$

The most basic absolute value function is  $f(x) = |x|$ :

$$|x| = \begin{cases} -x & x < 0 \\ x & x \geq 0 \end{cases}$$

- Floor Function

We define the **floor function**  $f(x) = \lfloor x \rfloor$  as the greatest integer smaller than or equal to  $x$ .

For  $x \in \mathbb{R}$  and  $n \in \mathbb{Z}$ ,

$$\lfloor x \rfloor = n \iff n \leq x < n + 1.$$

- Ceiling Function

The ceiling function  $f(x) = \lceil x \rceil$  is the direct opposite of the floor function; it maps all real numbers in the domain to the smallest integer not smaller than it.

$$\lceil x \rceil = \begin{cases} \lfloor x \rfloor + 1 & x \notin \mathbb{Z} \\ \lfloor x \rfloor & x \in \mathbb{Z} \end{cases}$$

### Exercise 2.3.2

Prove that

(a)  $\lfloor \sqrt{x} \rfloor = \lfloor \sqrt{\lfloor x \rfloor} \rfloor$

(b)  $\lceil \sqrt{x} \rceil = \lceil \sqrt{\lceil x \rceil} \rceil$

### Solution

(a)

$$\begin{aligned} \lfloor \sqrt{x} \rfloor &= n \\ \iff n &\leq \sqrt{x} < n + 1 \quad [\text{by definition of floor function}] \\ \iff n^2 &\leq x < (n + 1)^2 \quad [\text{square both sides}] \\ \iff n^2 &\leq \lfloor x \rfloor \leq x < (n + 1)^2 \\ \iff n &\leq \sqrt{\lfloor x \rfloor} < n + 1 \quad [\text{take square root throughout}] \\ \iff \lfloor \sqrt{\lfloor x \rfloor} \rfloor &= n \quad [\text{by definition of floor function}] \end{aligned}$$

(b)

$$\begin{aligned} \lceil \sqrt{x} \rceil &= n + 1 \\ \iff n &< \sqrt{x} \leq n + 1 \quad [\text{by definition of ceiling function}] \\ \iff n^2 &< x \leq (n + 1)^2 \quad [\text{square both sides}] \\ \iff n^2 &< x \leq \lceil x \rceil \leq (n + 1)^2 \\ \iff n &< \sqrt{\lceil x \rceil} \leq n + 1 \quad [\text{take square root throughout}] \\ \iff \lceil \sqrt{\lceil x \rceil} \rceil &= n + 1 \quad [\text{by definition of ceiling function}] \end{aligned}$$

□

### Symmetrical Functions

There are special functions with some form of geometric symmetry.

- Even Functions

$f$  is **even** if  $f(-x) = f(x)$  for every  $x \in D_f$ .

The graph of an even function is symmetric about the  $y$ -axis.

- Odd Functions

$f$  is **odd** if  $f(-x) = -f(x)$  for every  $x \in D_f$ .

The graph of an odd function is symmetric about the origin.

- Periodic Functions

$f$  is **periodic** if  $f(x+p) = f(x)$  for every  $x \in D_f$ , where  $p$  is a positive constant. The smallest such  $p$  is known as the period.

**Exercise 2.3.3**

For a triangle  $ABC$  with corresponding angles  $a, b$  and  $c$ , show that

$$\sin a + \sin b + \sin c \leq \frac{3\sqrt{3}}{2}$$

and determine when equality holds. (Hint:  $y = \sin x$  is concave)

**Solution** Since  $f(x) = \sin x$  is strictly concave on  $[0, \pi]$ ,

$$\begin{aligned} & \frac{1}{3}f(a) + \frac{1}{3}f(b) + \frac{1}{3}f(c) \\ &= \frac{1}{3}f(a) + \frac{2}{3}\left(\frac{1}{2}f(b) + \frac{1}{2}f(c)\right) \\ &\leq \frac{1}{3}f(a) + \frac{2}{3}\left(f\left(\frac{b+c}{2}\right)\right) \quad [\text{Concavity Inequality}] \\ &\leq f\left(\frac{a}{3} + \frac{2}{3}\left(\frac{b+c}{2}\right)\right) \quad [\text{Concavity Inequality}] \\ &= f\left(\frac{a+b+c}{3}\right) \end{aligned}$$

Hence

$$\sin a + \sin b + \sin c = f(a) + f(b) + f(c) \leq 3f\left(\frac{a+b+c}{3}\right) = 3\sin \frac{\pi}{3} = \frac{3\sqrt{3}}{2}.$$

Equality holds when  $a = b = c$ , i.e. when  $ABC$  is an equilateral triangle. □

## Exercises

**Problem 22.** Let  $A$  be the set of all complex polynomials in  $n$  variables. Given a subset  $T \subset A$ , define the *zeros* of  $T$  as the set

$$Z(T) = \{P = (a_1, \dots, a_n) \mid f(P) = 0 \text{ for all } f \in T\}$$

A subset  $Y \in \mathbb{C}^n$  is called an algebraic set if there exists a subset  $T \subset A$  such that  $Y = Z(T)$ .

Prove that the union of two algebraic sets is an algebraic set.

**Proof** We would like to consider  $T = \{f_1, f_2, \dots\}$  expressed as indexed sets  $T = \{f_i\}$ . Then  $Z(T)$  can also be expressed as  $\{P \mid \forall i, f_i(P) = 0\}$ .

Suppose that we have two algebraic sets  $X$  and  $Y$ . Let  $X = Z(S)$ ,  $Y = Z(T)$  where  $S, T$  are subsets of  $A$  (basically, they are certain sets of polynomials). Then

$$X = \{P \mid \forall f \in S, f(P) = 0\}$$

$$Y = \{P \mid \forall g \in T, g(P) = 0\}$$

We imagine that for  $P \in X \cap Y$ , we have  $f(P) = 0$  or  $g(P) = 0$ . Hence we consider the set of polynomials

$$U = \{f \cdot g \mid f \in S, g \in T\}$$

For any  $P \in X \cup Y$  and for any  $fg \in U$  where  $f \in S$  and  $f \in g$ , either  $f(P) = 0$  or  $g(P) = 0$ , hence  $fg(P) = 0$  and thus  $P \in Z(U)$ .

On the other hand if  $P \in Z(U)$ , suppose otherwise that  $P$  is not in  $X \cup Y$ , then  $P$  is neither in  $X$  nor in  $Y$ . This means that there exists  $f \in S, g \in T$  such that  $f(P) \neq 0$  and  $g(P) \neq 0$ , hence  $fg(P) \neq 0$ . This is a contradiction as  $P \in Z(U)$  implies  $fg(P) = 0$ . Hence we have  $X \cup Y = Z(U)$  and thus  $X \cup Y$  is an algebraic set.

Now the other direction is simpler and can actually be generalised: The intersection of arbitrarily many algebraic sets is algebraic.

The basic result is that if  $X = Z(S)$  and  $Y = Z(T)$  then  $X \cap Y = Z(S \cup T)$ . □

**Problem 23** (Modular Arithmetic). Define the ring of integers modulo  $n$ :

$$\mathbb{Z}/n\mathbb{Z} = \mathbb{Z}/\sim \text{ where } x \sim y \iff x - y \in n\mathbb{Z}.$$

The equivalence classes are called congruence classes modulo  $n$ .

- (a) Define the sum of two congruence classes modulo  $n$ ,  $[x], [y] \in \mathbb{Z}/n\mathbb{Z}$ , by

$$[x] + [y] = [x + y]$$

Show that the above definition is well-defined.

- (b) Define the product of two congruence classes modulo  $n$  and show that such a definition is well-defined.

### Solution

- (a) We often define such concepts by considering the **representatives** of the equivalence classes.

For example, here we define  $[x] + [y] = [x + y]$  for two elements  $[x]$  and  $[y]$  in  $\mathbb{Z}/n\mathbb{Z}$ . So what we know here are the classes  $[x]$  and  $[y]$ . But what exactly are  $x$  and  $y$ ? They are just some element in the equivalence classes that was arbitrarily picked out. We then perform the sum  $x + y$ , and consequently, we used this to point towards the class  $[x + y]$ .

However,  $x$  and  $y$  are arbitrarily picked. We want to show that, regardless of which representatives are chosen from the equivalence classes  $[x]$  and  $[y]$ , we will always obtain the same result.

In the definition itself, we have defined that, for the two representatives  $x$  and  $y$  we define  $[x] + [y] = [x + y]$ . So now, let's say that we take two other arbitrary representatives,  $x' \in [x]$  and  $y' \in [y]$ . Then by definition, we have

$$[x] + [y] = [x' + y']$$

Thus, our goal is to show that  $x' + y' = [x + y]$ . This expression means that the two sides of the equation are referring to the same equivalence class. Therefore, the expression above is completely equivalent to the condition.

$$x' + y' \sim x + y$$

We then check that this final expression is indeed true: Since  $x' \in [x]$  and  $y' \in [y]$ , we have

$$\begin{aligned} x' \sim x \text{ and } y' \sim y \\ \implies x' - x, y' - y \in n\mathbb{Z} \\ \implies (x' + y') - (x + y) = (x' - x) + (y' - y) \in n\mathbb{Z} \end{aligned}$$

- (b) The product of two congruence classes is defined by

$$[x][y] = [xy]$$

For any other representatives  $x', y'$  we have

$$\begin{aligned} x'y' - xy \\ = x'y' - xy' + xy' - xy \\ = (x' - x)y' + x(y' - y) \in n\mathbb{Z} \end{aligned}$$

Thus  $[x'y'] = [xy]$  and the product is well-defined.

□



**Problem 24.** Let  $A = \mathbb{R}$  and for any  $x, y \in A$ ,  $x \sim y$  if and only if  $x - y \in \mathbb{Z}$ . For any two equivalence classes  $[x], [y] \in A/\sim$ , define

$$[x] + [y] = [x + y] \text{ and } -[x] = [-x]$$

- (a) Show that the above definitions are well-defined.  
 (b) Find a one-to-one correspondence  $\phi : X \rightarrow Y$  between  $X = A/\sim$  and  $Y : |z| = 1$ , i.e. the unit circle in  $\mathbb{C}$ , such that for any  $[x_1], [x_2] \in X$  we have

$$\phi([x_1])\phi([x_2]) = \phi([x_1 + x_2])$$

- (c) Show that for any  $[x] \in X$ ,

$$\phi(-[x]) = \phi([x])^{-1}$$

### Solution

- (a)

$$(x' + y') - (x + y) = (x' - x) + (y' - y) \in \mathbb{Z}$$

Thus  $[x' + y'] = [x + y]$

$$(-x') - (-x) = -(x' - x) \in \mathbb{Z}$$

Thus  $[-x'] = [-x]$ .

- (b) Complex numbers in the polar form:  $z = re^{i\theta}$

Then the correspondence is given by  $\phi([x]) = e^{2\pi ix}$

$$[x] = [y] \iff x - y \in \mathbb{Z} \iff e^{2\pi i(x-y)} = 1 \iff e^{2\pi ix} = e^{2\pi iy}$$

Hence this is a bijection.

Before that, we also need to show that  $\phi$  is well-defined, which is almost the same as the above.

If we choose another representative  $x'$  then

$$\phi([x]) = e^{2\pi ix'} = e^{2\pi ix} \cdot e^{2\pi i(x'-x)} = e^{2\pi ix}$$

- (c) You can either refer to the specific correspondence  $\phi([x]) = e^{2\pi ix}$  or use its properties.

$$\phi(-[x])\phi([x]) = \phi([-x])\phi([x]) = \phi([-x + x]) = \phi([0]) = 1$$

□

**Problem 25** (Set of Rational Numbers). Let  $\mathbb{Z}$  be the set of integers, and let  $\mathbb{Z}^*$  be the set of nonzero integers. We define

$$\mathbb{Q} = \{(a, b) \mid a \in \mathbb{Z}, b \in \mathbb{Z}^*\} / \sim$$

where

$$(a, b) \sim (c, d) \iff ad = bc.$$

Let  $\frac{a}{b}$  denote the equivalence class for  $(a, b)$ . Such an equivalence class is called a rational number.

- (a) For any two rational numbers  $\frac{a}{b}$  and  $\frac{c}{d}$ , their sum is determined by

$$\frac{a}{b} + \frac{c}{d} = \frac{ad + bc}{bd}$$

Show that the above definition is well-defined.

- (b) Define the product of two rational numbers and show that such a definition is well-defined.

- (c) Prove that for every equivalence class  $\frac{a}{b} \in \mathbb{Q}$ , there exists a unique integer pair  $(p, q)$  satisfying the following properties:

$$q > 0, (p, q) = 1 \text{ and } (p, q) \in \frac{a}{b}.$$

- (d) Using the partial order of  $\mathbb{Z}$ , define the partial order of  $\mathbb{Q}$ .

### Solution

- (a) For this problem, we are dealing with a “hidden” equivalence class.

The expressions  $\frac{a}{b}$  and  $\frac{c}{d}$  themselves are derived from their representatives  $(a, b)$  and  $(c, d)$ .

So suppose that we choose other representatives  $(a', b')$  and  $(c', d')$ , then the sum would be

$$\frac{a'}{b'} + \frac{c'}{d'} = \frac{a'd' + b'c'}{b'd'}$$

We now have to show that  $\frac{ad+bc}{bd} = \frac{a'd'+b'c'}{b'd'}$ :

$$\begin{aligned} \frac{ad+bc}{bd} &\iff (ad+bc, bd) \sim (a'd'+b'c', b'd') \\ &\iff (ad+bc)b'd' = (a'd'+b'c')bd \end{aligned}$$

$$\begin{aligned} \frac{a}{b} &= \frac{a'}{b'} \\ (a, b) &\sim (a', b') \\ ab' &= a'b \end{aligned}$$

Hence

$$\begin{aligned} (ad+bc)b'd' &= ab'dd' + bb'cd' \\ &= a'bdd' + bb'c'd \\ &= (a'd' + b'c')bd \end{aligned}$$

- (b) The definition would be  $\frac{a}{b} \cdot \frac{c}{d} = \frac{ac}{bd}$ .

This is actually a lot simpler to check.

$$a'c'bd = (a'b)(c'd) = (ab')(cd') = acb'd'$$

Hence  $\frac{a'c'}{b'd'} = \frac{ac}{bd}$ .

- (c) We basically try to do this step by step as we would in simplifying fractions.

First pick  $b$  to be positive, otherwise we swap  $a$  and  $b$  with  $-a$  and  $-b$ .

Then simplify the common factors. For this one we let  $(a, b) = d$ , and  $a = dp, b = dq$ . Then  $(p, q)$  is the pair that we need

- (d) In order to define the partial order we need to account for whether the denominators are negative.

$\frac{a}{b} \leq \frac{c}{d}$ , and if  $b, d > 0$  then we can safely draw a connection to the expression  $ad \leq bc$

In order to show that this does in fact give a partial order we check that

(a) 1:  $ab \leq ab$  and hence  $\frac{a}{b} \leq \frac{a}{b}$

(b) 2: If  $\frac{a}{b} \leq \frac{c}{d}$  and  $\frac{c}{d} \leq \frac{e}{f}$ , then  $ad \leq bc$  and  $bc \leq df$ , hence  $ad \leq df$  and thus  $\frac{a}{b} \leq \frac{e}{f}$

(c) 3: This is trickier due to complications arising from inequalities and multiplication

If  $\frac{a}{b} \leq \frac{c}{d}$  and  $\frac{c}{d} \leq \frac{e}{f}$ , note that  $b, d, f > 0$  and so  $ad \leq bc$  and  $cf \leq de$ .

i)  $e < 0$ , then  $c < 0$  and  $a < 0$ , thus  $-ad \geq -bc$ ,  $-cf \geq -de$  and we have  $acdf \geq bcde$

$af \leq be$  ( $c < 0, d > 0$ )

Thus  $\frac{a}{b} \leq \frac{e}{f}$

ii)  $e \geq 0$  but  $a < 0$ , then  $af < 0 \leq be$  and thus  $\frac{a}{b} < \frac{e}{f}$

iii)  $a \geq 0$ , then  $c \geq 0$  and  $e \geq 0$ , and we have the ordinary case.

Hence proven.

□

**Problem 26** (Complex Numbers). Let  $\mathbb{R}[x]$  denote the set of real polynomials. Define

$$\mathbb{C} = \mathbb{R}[x]/(x^2 + 1)\mathbb{R}[x]$$

where

$$f(x) \sim g(x) \iff x^2 + 1 \text{ divides } f(x) - g(x).$$

The complex number  $a + bi$  is defined to be the equivalence class of  $a + bx$ .

- (a) Define the sum and product of two complex numbers and show that such definitions are well-defined.
- (b) Define the reciprocal of a complex number.

**Part II**

**Real Analysis**

# 3 The Real Number System

## §3.1 Rational numbers $\mathbb{Q}$

Rational numbers  $\mathbb{Q}$  can be introduced following a general procedure called the construction of field of fractions. Every rational number can be written as

$$\frac{m}{n}, \quad m, n \in \mathbb{Z}, n \neq 0.$$

Moreover, every nonzero rational number can be uniquely written as

$$\frac{p}{q}, \quad p \in \mathbb{Z}^+, q \in \mathbb{Z}, \gcd(p, q) = 1.$$

### §3.1.1 $\mathbb{Q}$ is a field

We explain the meaning of a **field** using  $\mathbb{Q}$  as an example.

**Proposition 3.1.1.**  $(\mathbb{Q}, +, \cdot)$  is a field, which means

- $(\mathbb{Q}, +)$  is an abelian group with identity 0:
  - $\mathbb{Q}$  is closed under addition  $+$  ( $+$  is a binary operation over  $\mathbb{Q}$ ).
  - Addition  $+$  is associative:  $(a + b) + c = a + (b + c)$ .
  - Addition  $+$  has identity element 0:  $a + 0 = 0 + a = a$ .
  - Any element has inverse element:  $a + (-a) = (-a) + a = 0$ .
  - Addition  $+$  is commutative:  $a + b = b + a$ .
- $(\mathbb{Q}, +, \cdot)$  is a commutative, unital ring:
  - $\mathbb{Q}$  is closed under multiplication  $\cdot$  ( $\cdot$  is a binary operation over  $\mathbb{Q}$ ).
  - Multiplication  $\cdot$  is associative:  $(a \cdot b) \cdot c = a \cdot (b \cdot c)$ .
  - Multiplication  $\cdot$  has unity element 1:  $a \cdot 1 = 1 \cdot a = a$ .
  - Multiplication  $\cdot$  is commutative:  $a \cdot b = b \cdot a$ .
  - Addition and multiplication satisfy distribution law:  $(a + b) \cdot c = a \cdot c + b \cdot c$ .
- Any nonzero element has a multiplicative inverse:  $a \cdot \frac{1}{a} = \frac{1}{a} \cdot a = 1$  for any  $a \in \mathbb{Q}, a \neq 0$ .

#### Exercise 3.1.1

Prove that  $(\mathbb{Z}, +, \cdot)$  is a commutative, unital ring, but it is not a field.

**Proof**

- Check  $(\mathbb{Z}, +, \cdot)$  is a commutative, unital ring.
- The number  $2 \in \mathbb{Z}$  (in fact, every nonzero number except  $\pm 1$ ) has no multiplication inverse in  $\mathbb{Z}$ .

□

### §3.1.2 $\mathbb{Q}$ is an ordered set

Let  $A$  be a set.

**Definition 3.1.2.** An **order** on  $A$  is a relation, denoted by  $<$ , with the following two properties:

- (i)  $\forall x, y \in A$ , one and only one of the following statements is true:  $x < y, x = y, y < x$ .
- (ii)  $\forall x, y, z \in A$ , if  $x < y$  and  $y < z$ , then  $x < z$ .

**Notation** The notation  $x \leq y$  indicates that  $x < y$  or  $x = y$ , without specifying which of these two is to hold. In other words,  $x \leq y$  is the negation of  $x > y$ .

**Definition 3.1.3.** An **ordered set** is a set  $S$  in which an order is defined.

**Example 3.1.4.**  $\mathbb{Q}$  is an ordered set if  $r < s$  is defined to mean that  $s - r$  is a positive rational number.

Let  $A \subset \mathbb{R}$ .

**Definition 3.1.5.**  $A$  is **bounded from above** if there exists an **upper bound**  $M \in \mathbb{R}$  such that  $x \leq M$  for all  $x \in A$ .  $A$  is **bounded from below** if there exists a **lower bound**  $m \in \mathbb{R}$  such that  $x \geq m$  for all  $x \in A$ .  $A$  is **bounded** in  $\mathbb{R}$  if it is bounded above and below.

**Definition 3.1.6.** The **supremum** (or *least upper bound*) of  $A$ , denoted by  $\sup A$ , is the smallest real number  $M$  such that  $x \leq M$  for all  $x \in A$ .

- (i)  $M$  is an upper bound for  $A$ .
- (ii) If  $N$  is an upper bound for  $A$ , then  $M \leq N$ .

**Remark** If  $M \in A$ , then  $M$  is the *maximum value* of  $A$ .

The following proposition is convenient in working with suprema.

**Proposition 3.1.7.** Let  $A$  be a nonempty subset of  $\mathbb{R}$  that is bounded above. Then  $M = \sup A$  if and only if

- (i)  $x \leq M$  for all  $x \in A$
- (ii) For any  $\varepsilon > 0$ , there exists  $a \in A$  such that  $M - \varepsilon < a$ .

**Proof** Suppose first that  $M = \sup A$ . Then clearly (i) holds (since this is identical to condition (1) in the definition of supremum). Now let  $\varepsilon > 0$ . Since  $M - \varepsilon < a$ , condition (ii) in the definition of supremum implies that  $M - \varepsilon$  is not an upper bound of  $A$ . Therefore, there must exist an element  $a$  in  $A$  such that  $M - \varepsilon < a$ , as desired. □

**Definition 3.1.8.** The **infimum** (or *greatest lower bound*) of  $A$ , denoted by  $\inf A$ , is defined as the largest real number  $m$  such that  $x \geq m$  for all  $x \in A$ .

- (i)  $m$  is a lower bound for  $A$ .
- (ii) If  $n$  is a lower bound for  $A$ , then  $m \geq n$ .

**Remark** If  $m \in A$ , then  $m$  is the *minimum value* of  $A$ .

**Proposition 3.1.9** (Uniqueness of supremum). If a set  $A \subset \mathbb{R}$  has a supremum, then it is unique.

**Proof** Assume that  $M$  and  $N$  are suprema of a set  $A$ .

Since  $N$  is a supremum, it is an upper bound for  $A$ . Since  $M$  is a supremum, then it is the least upper bound and thus  $M \leq N$ .

Similarly, since  $M$  is a supremum, it is an upper bound for  $A$ ; since  $N$  is a supremum, it is a least upper bound and thus  $N \leq M$ .

Since  $N \leq M$  and  $M \leq N$ , thus  $M = N$ . Therefore, a supremum for a set is unique if it exists.  $\square$

**Theorem 3.1.10** (Comparison Theorem). Let  $S, T \subset \mathbb{R}$  be non-empty sets such that  $s \leq t$  for every  $s \in S$  and  $t \in T$ . If  $T$  has a supremum, then so does  $S$ , and  $\sup S \leq \sup T$ .

**Proof** Let  $\tau = \sup T$ . Since  $\tau$  is a supremum for  $T$ , then  $t \leq \tau$  for all  $t \in T$ . Let  $s \in S$  and choose any  $t \in T$ . Then, since  $s \leq t$  and  $t \leq \tau$ , then  $s \leq \tau$ . Thus,  $\tau$  is an upper bound for  $S$ .

By the Completeness Axiom,  $S$  has a supremum, say  $\sigma = \sup S$ . We will show that  $\sigma \leq \tau$ . Notice that, by the above,  $\tau$  is an upper bound for  $S$ . Since  $\sigma$  is the least upper bound for  $S$ , then  $\sigma \leq \tau$ . Therefore,

$$\sup S \leq \sup T.$$

$\square$

Let's explore some useful properties of  $\sup$  and  $\inf$ .

**Proposition 3.1.11.** Let  $S, T$  be non-empty subsets of  $\mathbb{R}$ , with  $S \subseteq T$  and with  $T$  bounded above. Then  $S$  is bounded above, and  $\sup S \leq \sup T$ .

**Proof** Since  $T$  is bounded above, it has an upper bound, say  $b$ . Then  $t \leq b$  for all  $t \in T$ , so certainly  $t \leq b$  for all  $t \in S$ , so  $b$  is an upper bound for  $S$ .

Now  $S, T$  are non-empty and bounded above, so by completeness each has a supremum. Note that  $\sup T$  is an upper bound for  $T$  and hence also for  $S$ , so  $\sup T \geq \sup S$  (since  $\sup S$  is the least upper bound for  $S$ ).  $\square$

**Proposition 3.1.12.** Let  $T \subseteq \mathbb{R}$  be non-empty and bounded below. Let  $S = \{-t \mid t \in T\}$ . Then  $S$  is non-empty and bounded above. Furthermore,  $\inf T$  exists, and  $\inf T = -\sup S$ .

**Proof** Since  $T$  is non-empty, so is  $S$ . Let  $b$  be a lower bound for  $T$ , so  $t \geq b$  for all  $t \in T$ . Then  $-t \leq -b$  for all  $t \in T$ , so  $s \leq -b$  for all  $s \in S$ , so  $-b$  is an upper bound for  $S$ .

Now  $S$  is non-empty and bounded above, so by completeness it has a supremum. Then  $s \leq \sup S$  for all  $s \in S$ , so  $t \geq -\sup S$  for all  $t \in T$ , so  $-\sup S$  is a lower bound for  $T$ .

Also, we saw before that if  $b$  is a lower bound for  $T$  then  $-b$  is an upper bound for  $S$ . Then  $-b \geq \sup S$  (since  $\sup S$  is the least upper bound), so  $b \leq -\sup S$ . So  $-\sup S$  is the greatest lower bound.

So  $\inf T$  exists and  $\inf T = -\sup S$ .  $\square$

**Proposition 3.1.13** (Approximation Property). Let  $S \subseteq \mathbb{R}$  be non-empty and bounded above. For any  $\varepsilon > 0$ , there is  $s_\varepsilon \in S$  such that  $\sup S - \varepsilon < s_\varepsilon \leq \sup S$ .

**Proof** Take  $\varepsilon > 0$ .

Note that by definition of the supremum we have  $s \leq \sup S$  for all  $s \in S$ . Suppose, for a contradiction, that  $\sup S - \varepsilon \geq s$  for all  $s \in S$ .

Then  $\sup S - \varepsilon$  is an upper bound for  $S$ , but  $\sup S - \varepsilon < \sup S$ , which is a contradiction.

Hence there is  $s_\varepsilon \in S$  with  $\sup S - \varepsilon < s_\varepsilon$ .  $\square$

**Theorem 3.1.14.** Any set bounded from above/below must have a supremum/infimum.

**Proof** We prove the above theorem for supremum, using Dedekind cuts.

Let  $S$  be a real number set. We consider the rational number set<sup>1</sup>

$$A = \{x \in \mathbb{Q} \mid \exists y \in S, x < y\}$$

Now we go through the definitions to check that  $(A|B)$  is a Dedekind cut

<sup>1</sup>very important that you remember this for the definition of Dedekind cuts



1. Since  $S \neq \emptyset$ , pick  $y \in S$ , then  $[y] - 1$  is a real number smaller than some element in  $S$ , hence  $[y] - 1 \in A$  and thus  $A \neq \emptyset$ .

Also, since we are given that  $S$  is bounded, there exists a positive integer  $M$  as an upper bound for  $S$ , thus  $B \neq \emptyset$ . (Note that an upper bound is simply a number that is bigger than anything from the set, and is not the supremum.)

2. We define  $B$  to be the complement of  $A$  in  $\mathbb{Q}$  so condition 2 is trivial.
3. For any  $x, y \in A$ , if  $x < y$  and  $y \in A$ , then there exists  $z \in S$  such that  $y < z$ , hence  $x < z$  and thus  $x \in A$ .
4. Suppose otherwise that  $x \in A$  is the largest element in  $A$ , then there exists  $y \in S$  such that  $x < y$ . We then pick a rational number  $z$  between  $x$  and  $y$ . Since we still have  $z < y$ , we have  $z \in A$  but  $z > x$ , contradictory to  $x$  being the largest.

Now there is actually an issue with the proof for property 4 here: How exactly are we finding  $z$ ?

First  $x \in \mathbb{Q}$ . Then  $y \in \mathbb{R}$  so we rewrite it as  $y = (C|D)$  via definition.

$x < y$  translates to the fact that  $x \in C$ . Now, since  $y$  is real, by definition we know that  $C$  must not have a largest element. In particular,  $x$  is not largest and we can pick  $z \in C$  such that  $z > x$ . This is in fact the  $z$  that we need.

Now that all the properties of a real number are validated, we may finally conclude that  $\alpha = (A|B)$  is indeed a real number.

Now we need to show that  $\alpha$  is in fact the supremum of  $S$ .

Let  $x \in S$ . If  $x$  is not the maximum value of  $S$ , i.e.  $\exists y \in S$  such that  $x < y$ . Then  $x \in A$  and thus  $x < \alpha$ .

If  $x$  is the maximum value of  $S$ , then for any rational number  $y < x$  we have  $y \in A$ , and for any rational number  $y \geq x$  we have  $y \in B$ . Thus  $x = (A|B) = \alpha$ .

In conclusion,  $x \leq \alpha$  for all  $x \in S$ .

For any upper bound  $x$  of  $S$ , since  $\forall y \in S, x \geq y$  we have  $x \in B$  and thus  $x \geq \alpha$ .

Therefore,  $\alpha$  is the smallest upper bound of  $S$  and thus  $\sup S = \alpha$  exists.  $\square$

**Definition 3.1.15.** An ordered set  $S$  is said to have the **least-upper-bound property** if the following is true: If  $E \subset S$ ,  $E$  is not empty, and  $E$  is bounded above, then  $\sup E$  exists in  $S$ .

Example 3.2.1 shows that  $\mathbb{Q}$  does not have the least-upper-bound property.

We shall now show that there is a close relation between greatest lower bounds and least upper bounds, and that every ordered set with the least-upper-bound property also has the greatest-lower-bound property.

**Theorem 3.1.16.** Suppose  $S$  is an ordered set with the least-upper-bound property,  $B \subset S$ ,  $B$  is not empty, and  $B$  is bounded below. Let  $L$  be the set of all lower bounds of  $B$ . Then

$$\alpha = \sup L$$

exists in  $S$ , and  $\alpha = \inf B$ .

In particular,  $\inf B$  exists in  $S$ .

**Proof** Since  $B$  is bounded below,  $L$  is not empty. Since  $L$  consists of exactly those  $y \in S$  which satisfy the inequality  $y \leq x$  for every  $x \in B$ , we see that every  $x \in B$  is an upper bound of  $L$ . Thus  $L$  is bounded above. Our hypothesis about  $S$  thus implies that  $L$  has a supremum in  $S$ ; call it  $\alpha$ .

If  $\gamma < \alpha$  then  $\gamma$  is not an upper bound of  $L$ , hence  $\gamma \notin B$ . It follows that  $\alpha \leq x$  for every  $x \in B$ . Thus  $\alpha \in L$ .

If  $\alpha < \beta$  then  $\beta \notin L$ , since  $\alpha$  is an upper bound of  $L$ .

We have shown that  $\alpha \in L$  but  $\beta \notin L$  if  $\beta > \alpha$ . In other words,  $\alpha$  is a lower bound of  $B$ , but  $\beta$  is not if  $\beta > \alpha$ . This means that  $\alpha = \inf B$ .  $\square$

**Problem 27.** Consider the set  $\{\frac{1}{n} \mid n \in \mathbb{Z}^+\}$ .

- (a) Show that  $\max S = 1$ .
- (b) Show that if  $d$  is a lower bound for  $S$ , then  $d \leq 0$ .
- (c) Use (b) to show that  $0 = \inf S$ .

**Proof**

□

If we are dealing with rational numbers, the sup/inf of a set may not exist. For example, a set of numbers in  $\mathbb{Q}$ , defined by  $\{[\pi \cdot 10^n]/10^n\}$ . 3,3.1,3.14,3.141,3.1415,3.14159,... But this set does not have an infimum in  $\mathbb{Q}$ .

By ZFC, we have the Completeness Axiom, which states that any non-empty set  $A \subset \mathbb{R}$  that is bounded above has a supremum; in other words, if  $A$  is a non-empty set of real numbers that is bounded above, there exists a  $M \in \mathbb{R}$  such that  $M = \sup A$ .

**Problem 28.** Find, with proof, the supremum and/or infimum of  $\{\frac{1}{n}\}$ .

**Proof** For the supremum,

$$\sup \left\{ \frac{1}{n} \right\} = \max \left\{ \frac{1}{n} \right\} = 1.$$

For the infimum, for all positive  $a$  we can pick  $n = [\frac{1}{a}] + 1$ , then  $a > \frac{1}{n}$ . Hence

$$\inf \left\{ \frac{1}{n} \right\} = 0.$$

□

**Problem 29.** Find, with proof, the supremum and/or infimum of  $\{\sin n\}$ .

**Proof** The answer is easy to guess:  $\pm 1$

For the supremum, we need to show that 1 is the smallest we can pick, so for any  $a = 1 - \varepsilon < 1$  we want to find an integer  $n$  close enough to  $2k\pi + \frac{\pi}{2}$  so that  $\sin n > a$ .

Whenever we want to show the approximations between rational and irrational numbers we should think of the **pigeonhole principle**.

$$2k\pi + \frac{\pi}{2} = 6k + (2\pi - 6)k + \frac{\pi}{2}$$

Consider the set of fractional parts  $\{(2\pi - 6)k\}$ . Since this an infinite set, for any small number  $\delta$  there is always two elements  $\{(2\pi - 6)a\} < \{(2\pi - 6)b\}$  such that

$$|\{(2\pi - 6)b\} - \{(2\pi - 6)a\}| < \varepsilon$$

Then  $\{(2\pi - 6)(b - a)\} < \delta$

We then multiply by some number  $m$  (basically adding one by one) so that

$$0 \leq \{(2\pi - 6) \cdot m(b - a)\} - \left(2 - \frac{\pi}{2}\right) < \delta$$

Picking  $k = m(b - a)$  thus gives

$$\begin{aligned} 2k\pi + \frac{\pi}{2} &= 6k + (2\pi - 6)k + \frac{\pi}{2} \\ &= 6k + [(2\pi - 6)k] + 2 + (2\pi - 6)k - \left(2 - \frac{\pi}{2}\right) \end{aligned}$$

Thus  $n = 6k + [(2\pi - 6)k] + 2$  satisfies  $\left|2k\pi + \frac{\pi}{2} - n\right| < \delta$

Now we're not exactly done here because we still need to talk about how well  $\sin n$  approximates to 1.

We need one trigonometric fact:  $\sin x < x$  for  $x > 0$ . (This simply states that the area of a sector in the unit circle is larger than the triangle determined by its endpoints.)

$$\begin{aligned}\sin n &= \sin \left( n - \left( 2k\pi + \frac{\pi}{2} \right) + \left( 2k\pi + \frac{\pi}{2} \right) \right) \\ &= \cos \left( n - \left( 2k\pi + \frac{\pi}{2} \right) \right) \\ &= \cos \theta\end{aligned}$$

$$1 - \sin n = 2 \sin^2 \frac{\theta}{2} = 2 \sin^2 \left| \frac{\theta}{2} \right| \leq \frac{\theta^2}{2} < \delta$$

Hence we simply pick  $\delta = \varepsilon$  to ensure that  $1 - \sin n < \varepsilon$ , and we're done.  $\square$

**Theorem 3.1.17.** Archimedean Principle If  $a, b \in \mathbb{R}$  with  $a > 0$ , then there exists  $n \in \mathbb{N}$  such that  $na > b$ .

**Proof** We prove by contradiction: suppose otherwise, that the Archimedean Property is false. Then there exists  $a, b \in \mathbb{R}, a > 0$  such that  $na \leq b$  for all  $n \in \mathbb{N}$ .

For these particular  $a$  and  $b$ , we can say that  $b$  is an upper bound of  $S := \{na \mid n \in \mathbb{N}\}$ . From the completeness axiom,  $s_0 := \sup S$  exists. Let  $n \in \mathbb{N}$ , we have  $n+1 \in \mathbb{N}$ . So  $s_0 \geq (n+1)a = na + a$ .

Then we have  $s_0 - a \geq na$ . This is true for all  $n \in \mathbb{N}$ . So  $s_0 - a$  is an upper bound of  $S$ . However,  $s_0 - a < s_0$ , which contradicts that  $s_0$  is the least upper bound of  $S$ . This contradiction shows that the Archimedean Property is true.  $\square$

## §3.2 Real Numbers $\mathbb{R}$

### §3.2.1 Dedekind cuts

The rational number system is inadequate for many purposes, both as a field and as an ordered set (which we will discuss later). For instance, there is no rational  $p$  such that  $p^2 = 2$ . (We shall prove this presently.) This leads to the introduction of so-called “irrational numbers” which are often written as infinite decimal expansions and are considered to be “approximated” by the corresponding finite decimals. Thus the sequence

$$1, 1.4, 1.41, 1.414, 1.4142, \dots$$

“tends to  $\sqrt{2}$ ”. But unless the irrational number  $\sqrt{2}$  has been clearly defined, the question must arise: Just what is it that this sequence “tends to”?

This sort of question can be answered as soon as the so-called “real number system” is constructed.

**Example 3.2.1.** We now show that the equation

$$p^2 = 2$$

is not satisfied by any rational  $p$ . If there were such a  $p$ , we could write  $p = \frac{m}{n}$  where  $m$  and  $n$  are integers that are not both even. Let us assume this is done. Then we have

$$m^2 = 2n^2.$$

This shows that  $m^2$  is even. Hence  $m$  is even, and so  $m^2$  is divisible by 4. It follows that the RHS is divisible by 4, so  $n^2$  is even, which implies that  $n$  is even. Thus both  $m$  and  $n$  are even, a contradiction. Hence  $p^2 = 2$  is impossible for rational  $p$ .

We now examine this situation a little more closely. Let  $A$  be the set of all positive rationals  $p$  such that  $p^2 < 2$  and let  $B$  consist of all positive rationals  $p$  such that  $p^2 > 2$ . We shall show that  $A$  contains no largest number and  $B$  contains no smallest.

More explicitly, for every  $p \in A$  we can find a rational  $q \in A$  such that  $p < q$ , and for every  $p \in B$  we can find a rational  $q \in B$  such that  $q < p$ .

To do this, we associate with each rational  $p > 0$  the number

$$q = p - \frac{p^2 - 2}{p + 2} = \frac{2p + 2}{p + 2}. \quad (1)$$

Then

$$q^2 - 2 = \frac{2(p^2 - 2)}{(p + 2)^2}. \quad (2)$$

If  $p \in A$  then  $p^2 - 2 < 0$ , (1) shows that  $q > p$ , and (2) shows that  $q^2 < 2$ , thus  $q \in A$ ; if  $p \in B$  then  $p^2 - 2 > 0$ , (1) shows that  $0 < q < p$ , and (2) shows that  $q^2 > 2$ , thus  $q \in B$ .

We shall construct  $\mathbb{R}$  from  $\mathbb{Q}$ .

**Step 1** The members of  $\mathbb{R}$  will be certain subsets of  $\mathbb{Q}$ , called *cuts*. A cut is defined as any set  $\alpha \subset \mathbb{Q}$  with the following properties:

- (i)  $\alpha \neq \emptyset$ ,  $\alpha \neq \mathbb{Q}$ .
- (ii) If  $p \in \alpha$ ,  $q \in \mathbb{Q}$  and  $q < p$ , then  $q \in \alpha$ .
- (iii) If  $p \in \alpha$ , then  $p < r$  for some  $r \in \alpha$ .

Rational numbers are denoted by  $p, q, r, \dots$ , and cuts are denoted by  $\alpha, \beta, \gamma, \dots$

Note that (iii) simply says that  $\alpha$  has no largest member; (ii) implies two facts which will be used freely:

- If  $p \in \alpha$  and  $q \notin \alpha$  then  $p < q$ .
- If  $r \notin \alpha$  and  $r < s$  then  $s \notin \alpha$ .

**Step 2**

In fact in 1872, German mathematician Richard Dedekind pointed out that a real number  $x$  can be determined by its lower set  $A$  and upper set  $B$ :

$$A := \{a : \mathbb{Q} \mid a < x\}$$

$$B := \{b : \mathbb{Q} \mid x < b\}$$

He defined a “real number” as a pair of sets of rational numbers, the lower and upper sets shown above. Such a pair of sets of rational numbers are known as a Dedekind cut.

- $A$  is a **lower set**:  $\forall a, b \in \mathbb{R}$ , if  $a < b$  where  $b \in A$ , then  $a \in A$ .
- $B$  is an **upper set**:  $\forall a, b \in \mathbb{R}$ , if  $a < b$  where  $a \in B$ , then  $b \in B$ .

**Definition 3.2.2.** Given that  $B$  is the complement of  $A$  in the reals, a non-empty subset  $(A, B) \subset \mathbb{Q}$  is a **Dedekind cut** if:

- (i)  $A$  is non-empty:  $A \neq \emptyset$
- (ii)  $A$  and  $B$  are disjoint:  $A \cap B = \emptyset, A \cup B = \mathbb{Q}$
- (iii)  $A$  is closed downwards:  $\forall x, y \in \mathbb{Q}$  with  $x < y, y \in A$ , then  $x \in A$
- (iv)  $A$  does not contain a greatest element:  $\forall x \in A, \exists y \in A$  such that  $x < y$

Perhaps a not-so-intuitive fact here is that there are two possible things happening to  $B$ :

1.  $B$  contains a least element
2.  $B$  does not contain a least element

Case 1 and 2 are known as rational and irrational Dedekind cuts respectively.

**Definition 3.2.3** (Real numbers). The set of real numbers  $\mathbb{R}$  is defined to be the set of all Dedekind cuts.

**Remark** The way we think about this is that Dedekind cuts are real numbers, and real numbers are Dedekind cuts.

**Order relations**

Given real numbers  $\alpha$  and  $\beta$ , let  $\alpha = (A, B)$  and  $\beta = (C, D)$ . Then

$$\alpha < \beta \iff A \subset C$$

**Remark** Since  $B$  is the complement of  $A$ ,  $\alpha$  is completely determined by  $A$  itself.

This ordering on the real numbers satisfies the following properties:

- $x < y$  and  $y < z \implies x < z$
- Exactly one of  $x < y$ ,  $x = y$  or  $x > y$  holds
- $x < y \implies x + z < y + z$

**Proposition 3.2.4** (Ordering property). For any two real numbers  $\alpha$  and  $\beta$ , one of the following must hold:

$$\alpha < \beta \quad \alpha = \beta \quad \alpha > \beta$$

**Proof** We prove by contradiction.

Note that  $\alpha \leq \beta \iff A \subseteq C$  ( $A = C$  is possible).

Suppose otherwise, that all three of the above are false, then neither of the sets  $A$  and  $C$  can be a subset of the other.

We pick two rational numbers from each set: Pick  $p$  where  $p \in A$ ,  $p \notin C$ , pick  $q$  where  $q \in C$ ,  $q \notin A$

- Obviously we cannot have  $p = q$ .
- If  $p < q$ , then since  $q \in C$ , according to property 3, we have  $p \in C$ , a contradiction.
- Similarly for  $p > q$ , we would find that  $q \in A$ , a contradiction.

Hence our assumption is false.

$\therefore$  One of the three cases  $\alpha < \beta$ ,  $\alpha = \beta$ ,  $\alpha > \beta$  must hold. □

**Addition**

**Property 3.2.5** (Addition). Let  $\alpha = (A, B)$ ,  $\beta = (C, D)$ , then  $\alpha + \beta = (X, Y)$  where

$$X = \{a + c \mid a \in A, c \in C\}$$

**Proof** To show that  $(X, Y)$  is a Dedekind cut, we simply need to check the conditions for Dedekind cuts.

- Property 1 is trivial.

- Property 2 is by definition.

- Property 3:

Let  $x, y \in X$  satisfy  $x < y$ ,  $y \in X$ .

Let  $y = a + c$ ,  $a \in A$ ,  $c \in C$ .

Let  $\varepsilon = y - x$ .

Let  $a' = a - \frac{\varepsilon}{2}$ ,  $c' = c - \frac{\varepsilon}{2}$ .

Then

$$a' + c' = a + c - \varepsilon = x$$

$a' < a$ ,  $a \in A \implies a' \in A$ . Similarly,  $c' \in C$ .

$\therefore x = a' + c' \in X$ .

- Property 4:

$\forall a + c \in X, a \in A, c \in C, \exists a' \in A, c' \in C$  such that  $a < a', c < c'$ .

$\therefore a' + c' \in X$  satisfies  $a + c < a' + c'$ .

□

**Property 3.2.6** (Commutativity). Addition is **commutative**:

$$\alpha + \beta = \beta + \alpha$$

**Proof** The proof is trivial.

□

**Property 3.2.7** (Associativity). Addition is **associative**:

$$\alpha + (\beta + \gamma) = (\alpha + \beta) + \gamma$$

**Proof** Let  $\alpha = (A, A')$ ,  $\beta = (B, B')$ ,  $\gamma = (C, C')$

$$\beta + \gamma = (B + C, (B + C)')$$

In this notation we only need to show that  $A + (B + C) = (A + B) + C$ .

$$\begin{aligned} x \in A + (B + C) &\iff \exists a \in A, p \in B + C \text{ s.t. } x = a + p \\ &\iff \exists a \in A, b \in B, c \in C \text{ s.t. } x = a + b + c \\ &\iff x \in (A + B) + C \end{aligned}$$

Hence proven.

□

**Exercise 3.2.1**

Prove that

$$\alpha + 0 = \alpha = 0 + \alpha$$

**Proof** Let  $0 = (O, O')$  where  $O = \{x \mid x < 0\}$ ,  $O' = \{x \mid x \geq 0\}$ .

Let  $\alpha = (A, B)$ , then  $\alpha + 0 = (C, D)$  where

$$\begin{aligned} C &= \{a + \varepsilon \mid a \in A, \varepsilon < 0\} \\ &= \{a - \varepsilon \mid a \in A, \varepsilon > 0\} \end{aligned}$$

$$a - \varepsilon < a, a \in A \implies a - \varepsilon \in A \implies C \subseteq A$$

According to Property 4,  $\forall a \in A, \exists a' \in A$  such that  $a < a'$ .

Let  $\varepsilon = a' - a > 0$ , then

$$a = a' - \varepsilon, a' \in A, \varepsilon > 0 \implies a \in C$$

So  $A = C$ .

$\therefore \alpha + 0 = \alpha$  □

### Exercise 3.2.2

Express  $-\alpha$  in terms of  $\alpha$ ; show

$$\alpha + (-\alpha) = 0 = (-\alpha) + \alpha$$

**Proof** We split this into two cases.

**Case 1:**  $\alpha$  is a rational number, then  $\alpha = (A, B)$  where  $A = \{x \mid x < \alpha\}$ ,  $B = \{x \mid x \geq \alpha\}$ .

Let  $-\alpha = (A', B')$ , where  $A' = \{x \mid x < -\alpha\}$ ,  $B' = \{x \mid x \geq -\alpha\}$ . We see that  $\alpha + (-\alpha) \leq 0$  is obvious.

On the other hand, since  $0 = (O, O')$ , for any  $\varepsilon < 0$  we have

$$\varepsilon = \left(\alpha + \frac{\varepsilon}{2}\right) + \left(-\alpha + \frac{\varepsilon}{2}\right) \in A + A'$$

Hence  $\alpha + (-\alpha) = 0$ .

**Case 2:**  $\alpha$  is irrational, let  $\alpha = (A, B)$  where  $B$  does not have a lowest value. Then  $-B = \{-x \mid x \in B\}$  does not have a highest value.

We wish to define  $-\alpha = (-B, -A)$ , but first we need to show that this is well-defined by checking through all the conditions.

- Property 1: This is trivial.
- Property 2: Prove that  $-A$  and  $B$  are disjoint.  
Note that  $\forall x \in \mathbb{R}$ , if  $x = -y$ , then exactly one out of  $y \in A$  and  $y \in B$  is true  $\implies$  exactly one out of  $x \in -B$  and  $x \in -A$  is true.
- Property 3: Prove  $-B$  is closed downwards.  
Suppose otherwise, that  $x < y, y \in -B$  but  $x \notin -B$ . Then  $-y \in B$ ,  $-x \notin B$ . Since  $A$  is the complement of  $B$ ,  $-y \notin A$ ,  $-x \in A$ . But  $-y < -x$ , which is a contradiction.
- Property 4 is already guaranteed by the irrationality of  $\alpha$ .

All of these properties imply that the real numbers form a commutative group by addition. □

### Negation

Given any set  $X \subset \mathbb{R}$ , let  $-X$  denote the set of the negatives of those rational numbers. That is  $x \in X$  if and only if  $-x \in -X$ .

If  $(A, B)$  is a Dedekind cut, then  $-(A, B)$  is defined to be  $(-B, -A)$ .

This is pretty clearly a Dedekind cut. - proof



### Signs

A Dedekind cut  $(A, B)$  is **positive** if  $0 \in A$  and **negative** if  $0 \in B$ . If  $(A, B)$  is neither positive nor negative, then  $(A, B)$  is the cut representing 0.

If  $(A, B)$  is positive, then  $-(A, B)$  is negative. Likewise, if  $(A, B)$  is negative, then  $-(A, B)$  is positive. The cut  $(A, B)$  is non-negative if it is either positive or 0.

### Multiplication

Let  $\alpha = (A, B)$  and  $\beta = (C, D)$  where  $\alpha, \beta$  are both non-negative.

We define  $\alpha \times \beta$  to be the pair  $(X, Y)$  where

$X$  is the set of all products  $ac$  where  $a \in A, c \in C$  and at least one of the two numbers is non-negative.  $Y$  is the set of all products  $bd$  where  $b \in B, d \in D$ .

Intermediate Value Theorem

Bolzano-Weiersstrass Theorem

Connectedness of  $\mathbb{R}$

### §3.2.2 $\mathbb{R}$ is archimedean

**Theorem 3.2.8** (Archimedean Principle). For any  $x \in \mathbb{R}^+$  and  $y \in \mathbb{R}$ , there exists some  $n \in \mathbb{Z}^+$  such that

$$n \cdot x > y.$$

In particular, if we take  $n = 1$  from this theorem, we immediately get the following statement.

**Proposition 3.2.9.** For any  $y \in \mathbb{R}$ , there exists some positive integer  $n$  such that  $n > y$ .

We now give a proof of Proposition 3.2.9 directly without using Theorem 3.2.8, and then we prove 3.2.8 from Proposition 3.2.9. This shows that these two statements are in fact equivalent, though Proposition 3.2.9 looks much simpler.

### §3.2.3 $\mathbb{Q}$ is dense in $\mathbb{R}$

### §3.2.4 $\mathbb{R}$ is closed under taking roots

### §3.2.5 The extended real number system

## §3.3 The Euclidean Plane $\mathbb{R}^2$

We consider the Cartesian product of  $\mathbb{R}$  with  $\mathbb{R}$ ; that is,

$$\mathbb{R}^2 := \mathbb{R} \times \mathbb{R} := \{(x_1, x_2) \mid x_1, x_2 \in \mathbb{R}\}.$$

Over  $\mathbb{R}^2$ , we can define operations

- Addition  $+$ :  $(x_1, x_2) + (y_1, y_2) = (x_1 + y_1, x_2 + y_2)$ ;
- Scalar multiplication  $\mathbb{R} \times \mathbb{R}^2 \rightarrow \mathbb{R}^2$ :  $c \cdot (x_1, x_2) = (c \cdot x_1, c \cdot x_2)$ .

These two operations make  $\mathbb{R}^2$  a 2-dimensional vector space (linear space) over the real field  $\mathbb{R}$ . We also say  $\mathbb{R}^2$  is a  $\mathbb{R}$ -linear space of real dimension 2. For example,  $\{(1, 0), (0, 1)\}$  form a basis of  $\mathbb{R}^2$ .

Moreover, over the linear space  $\mathbb{R}^2$ , one can define an inner product as

$$\langle (x_1, x_2), (y_1, y_2) \rangle = x_1 y_1 + x_2 y_2.$$

The inner product induces a norm

$$|(x_1, x_2)| = \sqrt{\langle (x_1, x_2), (x_1, x_2) \rangle} = \sqrt{x_1^2 + x_2^2}.$$

From now on, we use  $\vec{x}$  to denote  $(x_1, x_2)$ .

**Proposition 3.3.1.**

- $|\vec{x}| \geq 0$ , where equality holds if and only if  $\vec{x} = \vec{0}$ .
- $|c \cdot \vec{x}| = |c| |\vec{x}|$
- $|\vec{x} + \vec{y}| \leq |\vec{x}| + |\vec{y}|$
- $|\langle \vec{x}, \vec{y} \rangle| \leq |\vec{x}| |\vec{y}|$

All constructions here can be easily generalised to any  $\mathbb{R}^n$  with  $n \in \mathbb{Z}^+$ .

### §3.4 The Complex Numbers $\mathbb{C}$

Over  $\mathbb{R}^2$ , we can define a multiplication  $\cdot$  as

$$(a, b) \cdot (c, d) = (ac - bd, ad + bc).$$

If we identify  $\mathbb{R}^2$  with

$$\mathbb{C} := \{x + yi \mid x, y \in \mathbb{R}\}$$

via  $(x, y) \mapsto x + yi$ , then all structures defined above are induced to  $\mathbb{C}$ . In particular, the multiplication is induced to  $\mathbb{C}$  via requiring  $i^2 = -1$ . A nontrivial fact is that  $(\mathbb{C}, +, \cdot)$  is a field. A element in  $\mathbb{C}$  is called a complex number. Usually, people prefer to use  $z = x + yi$ ,  $x, y \in \mathbb{R}$ , to denote a complex number. Here  $x$  is called the real part of  $z$  and  $y$  is called the imaginary part of  $z$ . We use  $|z|$  to denote its norm.

## §3.5 Completeness

### §3.5.1 Completeness axiom

**Theorem 3.5.1** (Completeness axiom for the real numbers). Let  $A$  be a non-empty subset of  $\mathbb{R}$  that is bounded above. Then  $A$  has a supremum.

Any set in the reals bounded from above/below must have a supremum/infimum.

**Proof** We prove this using Dedekind cuts.

Let  $S$  be a real number set. We consider the rational number set  $A = \{x \in \mathbb{Q} \mid \exists y \in S\}$ . Set  $B$  is defined to be the complement of  $A$  in  $\mathbb{Q}$ .

We go through the definitions to check that  $(A|B)$  is a Dedekind cut.

1. Since  $S \neq \emptyset$ , pick  $y \in S$ , then  $[y] - 1$  is a real number smaller than some element in  $S$ , hence  $[y] - 1 \in A$  and thus  $A \neq \emptyset$ .

Since we're given that  $S$  is bounded,  $\exists M > 0$  as the upper bound for  $S$ , thus  $B \neq \emptyset$ .

(Note that an upper bound is simply a number that is bigger than anything from the set, and is not the supremum)

2. We defined  $B$  to be the complement of  $A$  in  $\mathbb{Q}$ , so this condition is trivial.
3. For any  $x, y \in A$ , if  $x < y$  and  $y \in A$ , then  $\exists z \in S$  such that  $y < z \implies x < z \implies x \in A$ .
4. Suppose otherwise that  $x \in A$  is the largest element in  $A$ , then  $\exists y \in S$  such that  $x < y$ . We then pick a rational number  $z$  between  $x$  and  $y$ . Since we still have  $z < y$ , we have  $z \in A$  but  $z > x$ , contradictory to  $x$  being the largest.

Now there's actually an issue with the proof for property 4 here How exactly are we finding  $z$ ?

First  $x \in \mathbb{Q}$ . Then  $y \in \mathbb{R}$  so we rewrite it as  $y = (C|D)$  via definition.

$x < y$  translates to the fact that  $x \in C$ .

Since  $y$  is real, by definition we know that  $C$  must not have a largest element.

In particular,  $x$  is not largest and we can pick  $z \in C$  such that  $z > x$ . This is in fact the  $z$  that we need

Now that all the properties of a real number are validated, we may finally conclude that  $\alpha = (A|B)$  is indeed a real number.

Now we need to show that  $\alpha = \sup S$ .

Let  $x \in S$ . If  $x$  is not the maximum value of  $S$ , i.e.  $\exists y \in S, x < y$ , then  $x \in A$  and thus  $x < \alpha$ .

If  $x$  is the maximum value of  $S$ , then for any rational number  $y < x$  we have  $y \in A$ , and for any rational number  $y \geq x$  we have  $y \in B$ . Thus  $x = (A|B) = \alpha$ .

In conclusion,  $x \leq \alpha$  for all  $x \in S$ .

For any upper bound  $x$  of  $S$ , since  $\forall y \in S, x \geq y$  we have  $x \in B$  and thus  $x \geq \alpha$ .

$\therefore \alpha$  is the smallest upper bound of  $S$  and thus  $\sup S = \alpha$  exists. □

**Theorem 3.5.2** (Archimedean property of  $\mathbb{N}$ ).  $\mathbb{N}$  is not bounded above.

**Proof** Suppose, for a contradiction, that  $\mathbb{N}$  is bounded above. Then  $\mathbb{N}$  is non-empty and bounded above, so by completeness (of  $\mathbb{R}$ )  $\mathbb{N}$  has a supremum.

By the Approximation property with  $\varepsilon = \frac{1}{2}$ , there is a natural number  $n \in \mathbb{N}$  such that  $\sup \mathbb{N} - \frac{1}{2} < n \leq \sup \mathbb{N}$ .

Now  $n + 1 \in \mathbb{N}$  and  $n + 1 > \sup \mathbb{N}$ . This is a contradiction. □

# 4 Basic Topology

## §4.1 Metric Space

**Definition 4.1.1.** A set  $X$ , whose elements we shall call *points*, is a **metric space** if for any two points  $p, q \in X$  there is associated a real value function (called distance function or *metric*)  $d : X \times X \rightarrow \mathbb{R}$  which satisfies the following properties:

- (i) (**positive definiteness**)  $d(p, q) \geq 0$ , where equality holds if and only if  $x = y$ ;
- (ii) (**symmetry**)  $d(p, q) = d(q, p)$ ;
- (iii) (**triangle inequality**)  $d(p, q) \leq d(p, r) + d(r, q)$  for any  $r \in X$ .

**Example 4.1.2.** The most important examples of metric spaces are the euclidean spaces  $\mathbb{R}^n$ , especially  $\mathbb{R}^1$  (the real line) and  $\mathbb{R}^2$  (the complex plane); the distance in  $\mathbb{R}^n$  is defined by

$$d(\mathbf{x}, \mathbf{y}) = |\mathbf{x} - \mathbf{y}|.$$

A metric space  $(X, d)$  naturally induces a metric on any of its subsets.

**Definition 4.1.3.** For any  $x \in X$ ,  $r > 0$ , the subset  $B_r(x) := \{y \in X \mid d(y, x) < r\}$  is called the **open ball** centred at  $x$  with radius  $r$ .

Similarly, the subset  $\bar{B}_r(x) := \{y \in X \mid d(y, x) \leq r\}$  is called the **closed ball** centred at  $x$  with radius  $r$ .

An open ball centred at  $x$  is also called a **neighbourhood** of  $x$ .

**Example 4.1.4.** An open (closed) ball in  $\mathbb{R}$  is equivalent to a finite open (closed) interval, i.e.  $(a, b)$  ( $[a, b]$ ),  $a, b \in \mathbb{R}$ .

**Definition 4.1.5.** Let  $X$  be a metric space. All points and sets mentioned below are understood to be elements and subsets of  $X$ .

- (1)  $p$  is a **limit point** of  $E$  if every neighborhood of  $p$  contains  $q \neq p$  such that  $q \in E$ :

$$\forall r > 0, \exists q \in E, q \neq p \text{ s.t. } q \in B_r(p).$$

- (2)  $p$  is an **isolated point** of  $E$  if it not a limit point of  $E$ .

- (3)  $E$  is **closed** if every limit point of  $E$  is a point of  $E$ .

- (4)  $p$  is an **interior point** of  $E$  if there is a neighborhood  $N$  of  $p$  such that  $N \subset E$ :

$$\exists r > 0, B_r(p) \subset E.$$

- (5)  $E$  is **open** if every point of  $E$  is an interior point of  $E$ .

- (6)  $E$  is **perfect** if  $E$  is closed and if every point of  $E$  is a limit point of  $E$ .

(7)  $E$  is **bounded** if

$$\exists M \in \mathbb{R}, q \in X \text{ s.t. } \forall p \in E, d(p, q) < M.$$

(8)  $E$  is **dense** in  $X$  if every point of  $X$  is a limit point of  $E$ , or a point of  $E$  (or both).

**Proposition 4.1.6.** Any open ball is open.

**Proof** Assume  $B_r(x)$  is an open ball in a metric space  $(X, d)$ . Then for any point  $y \in B_r(x)$ , there is

$$d(y, x) < r.$$

Now we define  $r' := r - d(y, x)$ , which is positive.

Consider the ball  $B_{r'}(y)$ . We shall show it lives in  $B_r(x)$ . For this, take any point  $z \in B_{r'}(y)$ . Using the triangle inequality of a metric, we have

$$\begin{aligned} d(z, x) &\leq d(z, y) + d(y, x) \\ &< r' + d(y, x) \\ &= r. \end{aligned}$$

Hence  $z \in B_r(x)$ , and  $B_{r'}(y) \subset B_r(x)$ . □

**Proposition 4.1.7.** Assume  $(X, d)$  is a metric space.

- (1) Both  $\emptyset$  and  $X$  are open.
- (2) If  $S_1, S_2$  are open, then  $S_1 \cap S_2$  is open.
- (3) For any set  $\Lambda$  such that for any  $\alpha \in \Lambda$ ,  $S_\alpha$  is an open subset of  $X$ , the union  $\bigcup_{\alpha \in \Lambda} S_\alpha$  is open.

**Proof**

- (1) Obvious by definition.
- (2) Take a point  $x \in S_1 \cap S_2$ , we need to find an open ball with radius  $r > 0$  such that  $x \in B_r(x) \subset S_1 \cap S_2$ . To find such  $r > 0$ , notice that since both  $S_1$  and  $S_2$  are open, there are open balls

$$\begin{aligned} x \in B_{r_1}(x) &\subset S_1 \\ x \in B_{r_2}(x) &\subset S_2 \end{aligned}$$

Take  $r := \min\{r_1, r_2\}$ . Then  $B_r(x) \subset B_{r_1}(x) \subset S_1$  and  $B_r(x) \subset B_{r_2}(x) \subset S_2$ , and hence  $B_r(x) \subset S_1 \cap S_2$ .

- (3) Take a point  $x \in \bigcup_{\alpha \in \Lambda} S_\alpha$ , then we can assume  $x$  lives in some  $S_{\alpha_0}$ ,  $\alpha_0 \in \Lambda$ . Since  $S_{\alpha_0}$  is open, take an open ball

$$B_r(x) \subset S_{\alpha_0}.$$

It follows

$$B_r(x) \subset S_{\alpha_0} \subset \bigcup_{\alpha \in \Lambda} S_\alpha.$$

Hence  $\bigcup_{\alpha \in \Lambda} S_\alpha$  is open. □

**Example 4.1.8.** We know  $I_n := (-\frac{1}{n}, \frac{1}{n}) \subset \mathbb{R}$  is open for any  $n \in \mathbb{Z}^+$ . However,  $\bigcap_{n \in \mathbb{Z}^+} I_n = \{0\}$  is not open.

**Proposition 4.1.9.** Assume  $(X, d)$  is a metric space.

- (1) Both  $\emptyset$  and  $X$  are closed.
- (2) If  $S_1$  and  $S_2$  are closed, then  $S_1 \cup S_2$  is closed.
- (3) For any set  $\Lambda$  so that any  $\alpha \in \Lambda$ ,  $S_\alpha$  is a closed subset of  $X$ , the intersection  $\bigcap_{\alpha \in \Lambda} S_\alpha$  is closed.

**Proof**

(1) It follows immediately from  $\emptyset = X^c$  and  $X = \emptyset^c$ .

(2) It follows from above that

$$(S_1 \cup S_2)^c = S_1^c \cap S_2^c$$

is open, and hence  $S_1 \cup S_2$  is closed.

(3) It follows from above that

$$\left( \bigcap_{\alpha \in \Lambda} S_\alpha \right)^c = \bigcup_{\alpha \in \Lambda} S_\alpha^c$$

is open, and hence  $\bigcap_{\alpha \in \Lambda} S_\alpha$  is closed.

□

**Proposition 4.1.10.** If  $p$  is a limit point of  $E$ , then every neighbourhood of  $p$  contains infinitely many points of  $E$ .

**Proof** Prove by contradiction. Suppose there is a neighborhood  $B_r(p)$  which contains only a finite number of points of  $E$ :  $q_1, \dots, q_n$ , which are distinct from  $p$ . Define

$$r = \min_{1 \leq m \leq n} d(p, q_m).$$

The minimum of a finite set of positive numbers is clearly positive, so that  $r > 0$ .

The neighborhood  $B_r(p)$  contains no point  $q \in E$ ,  $q \neq p$  so that  $p$  is not a limit point of  $E$ , a contradiction.

□

**Corollary 4.1.11.** A finite point set has no limit points.

**Proposition 4.1.12.**  $E$  is open if and only if its complement  $E^c$  is closed.

**Proof**

( $\implies$ ) Suppose  $E$  is open. Let  $x$  be a limit point of  $E^c$ . Then every neighbourhood of  $x$  contains a point of  $E^c$ , so that  $x$  is not an interior point of  $E$ . Since  $E$  is open, this means that  $x \in E^c$ . It follows that  $E^c$  is closed.

( $\impliedby$ ) Suppose  $E^c$  is closed. Choose  $x \in E$ . Then  $x \notin E^c$ , and  $x$  is not a limit point of  $E^c$ . Hence there exists  $B_r(x)$  such that  $E^c \cap B_r(x)$  is empty, that is,  $B_r(x) \subset E$ . Thus  $x$  is an interior point of  $E$ , and  $E$  is open. □

**Definition 4.1.13.**  $E \subset X$ ,  $E'$  denotes the set of all limits points of  $E$  in  $X$ . Then the **closure** of  $E$  is the set  $\bar{E} = E \cup E'$ .

**Proposition 4.1.14.** If  $X$  is a metric space and  $E \subset X$ , then

- (1)  $\bar{E}$  is closed;
- (2)  $E = \bar{E}$  if and only if  $E$  is closed;
- (3)  $\bar{E} \subset F$  for every closed set  $F \subset X$  such that  $E \subset F$ .

By (1) and (3),  $\bar{E}$  is the *smallest* closed subset of  $X$  that contains  $E$ .

**Proof**

- (1)
- (2)
- (3)

□

## §4.2 Compact Sets

**Definition 4.2.1.** Assume  $(X, d)$  is a metric space. A collection of open sets  $\{U_\alpha \mid \alpha \in \Lambda\}$  is called an **open cover** of a subset  $S$  of  $X$ , if

$$S \subset \bigcup_{\alpha \in \Lambda} U_\alpha.$$

For  $\Lambda' \subset \Lambda$ , if the subcollection  $\{U_\alpha \mid \alpha \in \Lambda'\}$  is also an open cover of  $S$ ; that is,

$$S \subset \bigcup_{\alpha \in \Lambda'} U_\alpha,$$

then  $\{U_\alpha \mid \alpha \in \Lambda'\}$  is called a **subcover**. If moreover,  $\Lambda'$  is finite, then it is called a finite subcover.

## §4.3 Perfect Sets

## §4.4 Connected Sets

**Definition 4.4.1.** Two subsets  $A$  and  $B$  of a metric space  $X$  are said to be **separated** if both  $A \cap \bar{B}$  and  $\bar{A} \cap B$  are empty, i.e. no point of  $A$  lies in the closure of  $B$  and no point of  $B$  lies in the closure of  $A$ .

A set  $E \subset X$  is said to be **connected** if  $E$  is not a union of two non-empty separated sets.

**Remark** Separated sets are of course disjoint, but disjoint sets need not be separated. For example, the interval  $[0, 1]$  and the segment  $(1, 2)$  are not separated, since 1 is a limit point of  $(1, 2)$ . However, the segments  $(0, 1)$  and  $(1, 2)$  are separated.

The connected subsets of the line have a particularly simple structure:

**Proposition 4.4.2.** A subset  $E \subset \mathbb{R}^1$  is connected if and only if it has the following property: if  $x, y \in E$  and  $x < z < y$ , then  $z \in E$ .

**Proof** ( $\Leftarrow$ ) If there exists  $x, y \in E$  and some  $z \in (x, y)$  such that  $z \notin E$ , then  $E = A_z \cup B_z$  where

$$A_z = E \cap (-\infty, z), \quad B_z = E \cap (z, \infty).$$

Since

( $\Rightarrow$ )

□

# 5 Numerical Sequences and Series

## §5.1 Sequences in $\mathbb{R}$

### §5.1.1 Convergent Sequences

**Definition 5.1.1.** A sequence, which we denote by  $\{x_n\}$ , in  $\mathbb{R}$  is a map from  $\mathbb{Z}^+ \rightarrow \mathbb{R}$ , which maps  $n \in \mathbb{Z}^+$  to  $x_n \in \mathbb{R}$ . The range of the map is called the range of the sequence.

A **subsequence** of  $\{x_n\}$  is defined via an injective map  $s$  from  $\mathbb{Z}^+$  to a subset of  $\mathbb{Z}^+$  satisfying

$$s(k_1) < s(k_2) \quad \forall k_1, k_2 \in \mathbb{Z}^+, k_1 < k_2,$$

and denoted as  $\{x_{n_k}\}$  with  $x_{n_k} = x_{s(k)}$ .

A sequence  $\{x_n\}$  in  $\mathbb{R}$  is called **convergent**, if there exists some  $L \in \mathbb{R}$  such that for any  $\varepsilon > 0$ , there exists some  $N \in \mathbb{Z}^+$  so that for all  $n > N$ ,  $|x_n - L| < \varepsilon$ . We denote it as  $\lim_{n \rightarrow \infty} x_n = L$ , and call  $L$  the **limit** of the sequence  $\{x_n\}$ .

A sequence  $\{x_n\}$  in  $\mathbb{R}$  is called **divergent**, if it has no limit in  $\mathbb{R}$ .

**Remark** Take note of the use of logical statements:

- $\varepsilon$  is independent, so it is literally for all  $\varepsilon > 0$ .
- $N$  is dependent on  $\varepsilon$ ; if  $\varepsilon$  is very small we would expect the sequence  $\{a_n\}$  to get close enough to  $L$  further down the line.
- The order of the quantifiers matters.

**Proposition 5.1.2.**

- The limit of a convergent sequence in  $\mathbb{R}$  is unique.
- The sequence  $\{x_n\}$  converges to  $L \in \mathbb{R}$  if and only if every open disk centred at  $L$  contains all but finitely many of terms in the sequence.
- The sequence  $\{x_n\}$  converges to  $L \in \mathbb{R}$  if and only if every subsequence of it converges to  $L \in \mathbb{R}$ .
- If a sequence  $\{x_n\}$  in  $\mathbb{R}$  is convergent, then it must be bounded.
- The set of all subsequential limits of a sequence  $\{x_n\}$  in  $\mathbb{R}$  is closed.



**Exercise 5.1.1**

What do we really mean by saying that  $\frac{1}{n} \rightarrow 0$  as  $n \rightarrow \infty$ ?

We mean that the sequence of numbers  $\frac{1}{n}$  converges to 0, proven as follows:

**Proof**  $\forall \varepsilon > 0$ , pick  $N = \frac{1}{\varepsilon} + 1$ . Then  $\forall n > N$ ,

$$\frac{1}{n} < \frac{1}{N} < \frac{1}{\frac{1}{\varepsilon}} = \varepsilon.$$

□

We shall cover some characteristics of limits.

**Proposition 5.1.3.** Given a sequence of points  $\{x_k\}$  and a point  $x \in \mathbb{R}^n$ ,  $x_k$  converges to  $x$  if and only if all neighbourhoods of  $x$  “eventually” contain all  $x_k$ .

By “eventually” we mean something similar to the definition above: there exists some large  $N$  such that the property is satisfied for all  $n > N$ .

**Proof****Forward direction:**

If  $\{x_k\}$  converges to  $x$ , we wish to prove: given any neighbourhood  $U$  of  $x$ ,  $U$  eventually contains all  $x_k$ . Since  $U$  is a neighbourhood of  $x$ , we pick a ball of radius  $\varepsilon$  centered at  $x$ ,  $B(x, \varepsilon)$ , so that  $B(x, \varepsilon)$  is contained in  $U$ .

Then since  $B(x, \varepsilon)$  is precisely the set of points whose distance to  $x$  is no larger than  $\varepsilon$ , we then apply the fact that  $\{x_k\}$  converges to  $x$ .

So for this particular  $\varepsilon$ , we take a natural number  $N$  so that  $|x_k - x| < \varepsilon$ , or  $x_k \in B(x, \varepsilon)$ , for all  $k > N$ .

Then simultaneously  $x_k$  are in  $U$  since  $B(x, \varepsilon)$  is a subset of  $U$ , thus we’ve shown that  $U$  will contain all  $x_k$  after a certain point  $N$ .

**Backward direction:**

Suppose that all neighbourhoods of  $x$  will eventually contain all  $x_k$ , then in particular for any  $\varepsilon > 0$ , since  $B(x, \varepsilon)$  is a neighbourhood of  $x$ , it will also eventually contain all  $x_k$ .

This then easily translates to the fact that  $\{x_k\}$  converges to  $x$ . □

**Proposition 5.1.4** (Uniqueness of the limit). Suppose that  $\{x_k\}$  converges to both  $x$  and  $x'$ , then  $x = x'$ .

**Proof**  $\forall \varepsilon > 0$ , we know that the terms in  $\{x_k\}$  must be less than  $\varepsilon$  away from its limit after a certain point.

However, this certain point may not be the same for both limits; for the two limits  $x$  and  $x'$ , we must first assume two separate numbers  $N$  and  $N'$  so that  $|x_k - x| < \varepsilon$  when  $k > N$ , and  $|x_k - x'| < \varepsilon$  when  $k > N'$ .

Now if you look at the book here, it says that we have a stronger requirement:  $|x_k - x| < \frac{\varepsilon}{2}$  when  $k > N$ ,  $|x_k - x'| < \frac{\varepsilon}{2}$  when  $k > N'$ . This is simply because we want to prove certain statements strictly by definition

There is an important detail to take note, regarding  $\max\{N, N'\}$ .

We’re taking the larger one of these, so it means that, after this certain point, we in fact have  $|x_k - x| < \frac{\varepsilon}{2}$  and  $|x_k - x'| < \frac{\varepsilon}{2}$  at the same time.

Therefore by triangle inequality,

$$|x - x'| \leq |x_k - x| + |x_k - x'| < \varepsilon$$

The choice of  $k$  actually vanished in the final statement; you can think of this as if picking this particular choice of  $k$  helps us to establish some kind of property for the original objects

Finally, since we've in fact proven that  $|x - x'| < \varepsilon$  holds for any given positive  $\varepsilon > 0$ , we must have  $|x - x'| = 0$  and therefore  $x = x'$ .

Strictly speaking, for the first part we need to explain why  $a < \varepsilon$  for any positive  $\varepsilon$  implies that  $a \leq 0$ . This is very easy to prove (by contradiction) so let's not be too redundant. The second part simply relies on the fact that  $|x - y|$  is the Euclidean metric and so by positive definiteness  $|x - y| = 0$  if and only if  $x = y$ .  $\square$

**Proposition 5.1.5** (Boundedness of converging sequences). If  $\{x_k\}$  converges, then  $\{x_k\}$  is bounded.

We simply take the limit  $x$  and note that the sequence is eventually contained in some ball centered at  $x$ , say  $B(x, 1)$ .

There are several outlying points prior to this, but since there are only a finite number of these, it doesn't change the fact that the sequence (viewed as a set) is bounded nevertheless.

This argument is precisely expressed by the construction of  $r$  given in the book: let  $|x_k - x| < 1$  whenever  $k > N$ , then  $\{x_k\}$  is in  $B(x, r)$  where  $r = \max\{1, |x_1 - x|, \dots, |x_N - x|\}$

1. We talk about the relationship between the limit of a sequence and the limit points of a set.

Generally, limit points are a weaker construction.

Suppose that  $\{x_k\}$  converges to  $x$ . If we view  $\{x_k\}$  as a set, then  $x$  will be a limit point of this set.

The converse, however, is not true.

Exercise 1: Construct a sequence in  $\mathbb{R}$  that is bounded and contains a single limit point but is divergent (not convergent).

The thing about convergence of a series is that, unlike for limit points where we only require that there are other points that get arbitrarily close, but moreover we have to ensure that this pattern ensues for each and every term in the sequence.

Me: Suppose that  $\{x_k\}$  converges to  $x$ . If we view  $\{x_k\}$  as a set, then  $x$  will be a limit point of this set." - - - - - Sorry I forgot something crucial about this: There is the strange possibility that the sequence  $\{x_k\}$  is constant: (or at least eventually constant): Then in fact  $x$  by definition is not a limit point of  $x_k$  because you can find a ball around  $x$  that only contains the element  $x$  itself, since that point is merely what the entire sequence  $\{x_k\}$  amounts to: Anyways, we simply can't say that a sequence  $\{x_k\}$  converges to  $x$  if we're only provided with the fact that  $x$  is a limit point of  $\{x_k\}$ .

However, we can say the following: (d) If  $x$  is a limit point of  $E$ , then there exists a sequence  $\{x_n\}$  in  $E \setminus x$  such that  $\{x_n\}$  converges to  $x$ .

In fact this is correct in both ways so let's rewrite this as follows: (d)  $x$  is a limit point of  $E$ , if and only if there exists a sequence  $\{x_n\}$  in  $E \setminus x$  such that  $\{x_n\}$  converges to  $x$ .

( $E \setminus x$  is important here, otherwise we simply pick the constant sequence  $x_k = x$ )

$\rightarrow$ : If  $x$  is a limit point, then for all  $\varepsilon > 0$ ,  $B_0(x, \varepsilon)$  contains points in  $E$ . We then construct such a sequence  $\{x_k\}$  in  $E \setminus x$ : pick any  $x_k \in E$  so that  $x_k$  is contained in  $B_0(x, 1/k)$ .

Then it is easy to show that  $\{x_k\}$  is a sequence in  $E \setminus x$  which converges to  $x$ .

$\leftarrow$ : Suppose that there exists a sequence  $\{x_n\}$  in  $E \setminus x$  such that  $\{x_n\}$  converges to  $x$ . We wish to show that  $B_0(x, \varepsilon)$  contains points in  $E$  for all  $\varepsilon > 0$ .

Since  $\{x_n\}$  converges to  $x$ , for all  $\varepsilon > 0$  the sequence is eventually contained in  $B(x, \varepsilon)$ . However because we have the precondition that  $\{x_n\}$  has to be in  $E \setminus x$ , the sequence is in fact eventually contained in  $B_0(x, \varepsilon)$ .

**Proposition 5.1.6.**  $\{x_k\}$  converges to  $x$  if and only if every subsequence of  $\{x_k\}$  converges to  $x$ .

**Proof** We only need to prove this in the forwards direction. Every subsequence of  $\{x_k\}$  can be written in the form  $\{x_{k_i}\}$  where  $k_1 < k_2 < \dots$  is a strictly increasing sequence of natural numbers.

Intuitively, if every neighbourhood of  $x$  eventually contains all  $x_k$ , then since  $\{x_{k_i}\}$  is just a subset of  $\{x_k\}$  they should all be contained in the neighbourhood eventually as well. For every  $\varepsilon > 0$ , pick  $N$  such that for  $k > N$ ,  $|x_k - x| < \varepsilon$ . Pick  $M$  such that  $k_M > N$ , then for all  $i > M$  we have  $|x_{k_i} - x| < \varepsilon$ .  $\square$

**Proposition 5.1.7.** Subsequential limits of a sequence are precisely the limit points of the sequence (viewed as a set)

**Proof** This is just part (d) of the previous section.

Again, to make this work, we need to assume that nothing funny is going on at subsequential limits. If the limits appear due to eventually constant subsequences, then they need not be limit points of the original sequence when viewed as a set.

3.6, 3.7 are precisely the statements we've prepared for last week □

**Proposition 5.1.8.** If  $\{x_n\}$  is a sequence in a compact set (bounded closed set), then there exists a convergent subsequence of  $\{x_n\}$

**Proof** This is Weierstrass-Bolzano together with part (b)

Ah yes, regarding compact sets I need to emphasize this again, but the definition that we are currently using for compact sets is not the actual definition.

I've sent a video before the lesson which talks about the real definition for compact sets. Essentially, compact sets satisfy the property akin to the statement in Heine-Borel: Given a topological space  $(X, \tau)$ , a compact set  $K$  in  $X$  is a set satisfying that, given any open covering  $\{U_i\}$  of  $X$ , there exists a finite open cover  $\{U_1, \dots, U_n\}$  of  $X$ .

This is difficult to process at this stage. Since we're currently only working with Euclidean spaces it would be more beneficial if you consider the Heine-Borel Theorem as a property first. It would be a lot easier to accept the definition after you're more accustomed to applying the theorem. □

**Proposition 5.1.9.** (Rudin 3.7) Subsequential limits form a closed subset.

**Proof** Actually we've done this two weeks before, it is simply saying that  $A''$  is a subset of  $A'$ .

( $A''$  is not always  $A'$ ; consider the set in  $\mathbb{R}^2$  given by  $(1/n, 1/m) | n, m \in \mathbb{N}$ . Then  $(1, 0), (0, 1)$  are in  $A'$  but not in  $A''$ .) □

## §5.1.2 Cauchy Sequences

This is a very very helpful way to determine whether a sequence is convergent or divergent, as it does not require the limit to be known. In the future you will see many instances where the convergence of all sorts of limits are compared with similar counterparts; generally we describe such properties as **Cauchy criteria**.

Cauchy sequences have to deal with the differences between terms within the sequence itself.

**Definition 5.1.10.** A sequence  $\{x_k\}$  in  $\mathbb{R}^n$  is **Cauchy**, if the distances between any two terms is sufficiently small after a certain point.

Strictly speaking,  $\forall \varepsilon > 0$ , there exists integer  $N$  such that

$$\forall k, l > N, |x_k - x_l| < \varepsilon.$$

It is easy to prove that a converging sequence is Cauchy using the triangle inequality. The idea is that, if all the points are becoming arbitrarily close to a given point  $p$ , then they are also becoming close to each other. The converse is not always true, however.

**Proposition 5.1.11.** A sequence  $\{x_k\}$  in  $\mathbb{R}^n$  is convergent if and only if it is Cauchy.

**Proof Forward direction:**

Suppose that  $\{x_k\}$  converges to  $x$ , then there exists  $N$  such that for  $k > N$ ,  $|x_k - x| < \frac{\varepsilon}{2}$ . Then for  $k, l > N$ ,

$$|x - k - x_l| \leq |x_k - x| + |x_l - x| < \varepsilon$$

**Backward direction:**

First, we show that  $\{x_k\}$  must be bounded. Pick  $N$  such that for all  $k, l > N$  we have  $|x_k - x_l| < 1$ . Centered at  $x_k$ , we show that  $\{x_k\}$  is bounded; to do this we pick

$$r = \max\{1, |x_k - x_1|, \dots, |x_k - x_N|\}$$

Then the sequence  $x_k$  is in  $B(x_k, r)$  and thus is bounded.

Since  $\{x_k\}$  is bounded, by the collorary of Bolzano-Weierstrass we know that  $\{x_k\}$  contains a subsequence  $\{x_{k_i}\}$  that converges to a limit  $x$ .

Then for all  $\varepsilon > 0$ , pick  $N_1$  such that for all  $k, l > N$ ,  $|x_k - x_l| < \frac{\varepsilon}{2}$ . Simultaneously, since  $\{x_{k_i}\}$  converges to  $x$ , pick  $M$  such that for  $i > M$ ,  $|x_{k_i} - x| < \frac{\varepsilon}{2}$ .

Now, since  $k_1 < k_2 < \dots$  is a sequence of strictly increasing natural numbers, we can pick  $i > M$  such that  $k_i > N$ . Then for all  $k > N$ , by setting  $l = k_i$  we obtain

$$|x_k - x_{k_i}| < \frac{\varepsilon}{2}, \quad |x_{k_i} - x| < \frac{\varepsilon}{2}$$

and hence

$$|x_k - x| \leq |x_k - x_{k_i}| + |x_{k_i} - x| < \varepsilon$$

□

### §5.1.3 Upper and Lower Limits

### §5.1.4 Limits of Multiple Sequences

We shall cover some of the more basic aspects of limits in this section.

#### Inequalities

First let's consider two converging sequences  $\{a_n\}$  and  $\{b_n\}$

If  $a_n \leq b_n$ , then  $\lim a_n \leq \lim b_n$ .

**Remark** One important thing to take note for limits is that, even if you have  $a_n < b_n$ , you cannot say that  $\lim a_n < \lim b_n$ ; for example,  $\frac{1}{n} > -\frac{1}{n}$  but their limits are both 0.

**Proof** Let's say that  $A = \lim a_n$  and  $B = \lim b_n$ . Suppose otherwise that  $A > B$ , then we try to cause some chaos with  $\varepsilon = A - B > 0$ .

Since  $\frac{\varepsilon}{2} > 0$ , then there exists  $N_1$  such that for  $n > N_1$  we have  $|a_n - A| < \frac{\varepsilon}{2}$ ; and there exists  $N_2$  such that for  $n > N_2$  we have  $|b_n - B| < \frac{\varepsilon}{2}$ .

Let  $N = \max\{N_1, N_2\}$ , then for any  $n > N$ , the two inequalities above will hold simultaneously. But then we would have

$$a_n > A - \frac{\varepsilon}{2}, b_n < B + \frac{\varepsilon}{2}$$

and thus

$$a_n - b_n > A - B - \varepsilon = 0,$$

so  $a_n > b_n$ , a contradiction

□

A corollary is that limits essentially preserve signs, if you include 0 in your consideration

A converging sequence of nonnegative numbers will always be nonnegative, and same goes to nonpositive numbers: Now as we can see in the proof above, there is actually a place where the restrictions of limits overpower the statement itself: What I mean by that is, suppose that you want to form a proof by contradiction: What you need here is just one term  $a_n > b_n$ . But you actually have  $a_n > b_n$  eventually for all terms in the sequence: In fact, a better exercise would have been to show that limsups and liminfs also preserves inequalities

I'll just use limsups for example. If  $a_n \leq b_n$ , let  $A = \limsup a_n$ ,  $B = \limsup b_n$ . Suppose otherwise that  $A > B$ . Let  $\varepsilon = A - B > 0$ ; since  $\frac{\varepsilon}{2} > 0$ , then for all  $N_1$ , there exists  $n > N_1$  such that  $a_n > A - \frac{\varepsilon}{2}$ ; and there exists  $N_2$  such that for all  $n > N_2$ ,  $b_n < B + \frac{\varepsilon}{2}$ .

Now we arrange our thoughts logically First, we pick  $N_2 = N$  such that for all  $n > N$ ,  $b_n < B + \frac{\varepsilon}{2}$ . Then we may fix  $N_1 = N$ .

Due to the first condition, we see that it is possible to pick  $n_0 > N$  such that  $a_{n_0} > A - \frac{\varepsilon}{2}$ . Now due to the second condition, since  $n_0 > N$ , this exact same  $n_0$  would satisfy  $b_{n_0} < B + \frac{\varepsilon}{2}$ .

Therefore,  $n_0$  satisfies  $a_{n_0} - b_{n_0} > A - B - \varepsilon = 0$  and we are done.

### Sandwich Theorem

**Theorem 5.1.12** (Sandwich Theorem). Let  $a_n \leq c_n \leq b_n$  where  $\{a_n\}, \{b_n\}$  are converging sequences such that  $\lim a_n = \lim b_n = L$ , then  $\{c_n\}$  is also a converging sequence and  $\lim c_n = L$ .

Now, one very very very important thing about this theorem

The purpose of this theorem is to investigate some difficult sequence  $\{c_n\}$  with two simpler sequences  $\{a_n\}$  and  $\{b_n\}$  which bounds it from below and from above respectively. If you look closely at the statement, you may realize that we're only working under the condition that  $\{a_n\}$  and  $\{b_n\}$  are converging sequences.

In other words, at this point we don't know whether  $\{c_n\}$  is convergent.

In fact, this is supposed to be the main implication

Of course,  $\lim c_n = L$  is proven at the exact same time, so both implications constitute the two parts of the conclusion

What I want to say is that you cannot simply take  $\lim$  over  $a_n \leq c_n \leq b_n$  and say that  $\lim$  preserves inequalities, because in order to apply this inequality-preserving property, you need to ensure that all sequences are converging before you can apply it; clearly, this does not work here since we have not shown that  $c_n$  is convergent, therefore this idea does not work.

There are two ways to circumvent this. One is to use  $\varepsilon - N$ ; basically, just do it.

But if you're really lazy, then the second method is to use the idea above except you first take  $\limsup$  and  $\liminf$ .

The advantage of these two is that you don't need the original sequences to be convergent in order to apply them, and that they preserve inequalities even if the original sequences show no signs of convergence.

So basically,

$$\limsup a_n \leq \limsup c_n \leq \limsup b_n,$$

and

$$\liminf a_n \leq \liminf c_n \leq \liminf b_n.$$

Then since  $\{a_n\}$  and  $\{b_n\}$  actually converge to  $L$ , all the  $\liminf$ s and  $\limsup$ s of  $a_n$  and  $b_n$  are  $L$ , so we obtain  $\limsup c_n = L$  and  $\liminf c_n = L$ .

In particular,  $\limsup c_n = \liminf c_n$ , thus  $c_n$  is convergent and it follows that  $\lim c_n = L$ .

### Arithmetic properties

**Proposition 5.1.13.** For converging  $\{a_n\}$  and real constant  $k$ ,

$$\lim ka_n = k \lim a_n.$$

**Proof** The proof is left as an exercise. You will need to  $k$  into cases where it is positive, negative or 0. □

**Remark** In multivariable calculus there's a similar property that is more interesting:

If  $T$  is a linear map on  $\mathbb{R}^n$ , and  $\{x_n\}$  is a converging sequence of points, then  $\{Tx_n\}$  is also converging; moreover if  $x_n \rightarrow x_0$  then  $Tx_n \rightarrow Tx_0$ .

**Proposition 5.1.14.** If  $\{a_n\}$  and  $\{b_n\}$  are converging sequences of real numbers, then

$$\lim(a_n + b_n) = \lim a_n + \lim b_n.$$

**Proof** Let  $A = \lim a_n$  and  $B = \lim b_n$ , then for all  $\varepsilon > 0$ , there exists  $N_1$  such that for all  $n > N_1$ ,  $|a_n - A| < \frac{\varepsilon}{2}$ ; there exists  $N_2$  such that for all  $n > N_2$ ,  $|b_n - B| < \frac{\varepsilon}{2}$ .

Let  $N = \max\{N_1, N_2\}$ , then for all  $n > N$ , by the triangle inequality we have

$$|(a_n + b_n) - (A + B)| \leq |a_n - A| + |b_n - B| < \varepsilon.$$

□

**Remark** This proof is simple enough to generalise to any normed vector spaces.

The following corollary can be easily derived from the above.

**Corollary 5.1.15.** If  $\{a_n\}$  and  $\{b_n\}$  are converging sequences of real numbers, then

$$\lim(a_n - b_n) = \lim a_n - \lim b_n.$$

**Proposition 5.1.16.** If  $\{a_n\}$  and  $\{b_n\}$  are converging, then

$$\lim(a_n b_n) = \lim a_n \cdot \lim b_n.$$

**Proof** Let  $A = \lim a_n$  and  $B = \lim b_n$ .

Consider the limit  $\lim(a_n b_n - AB)$ , as it would be sufficient to prove that this is equal to 0.

Now we will use a common technique to deal with such products:

$$\lim(a_n b_n - AB) = \lim(a_n b_n - Ab_n + Ab_n - AB)$$

The idea is to show that this is equal to

$$\lim(a_n b_n - Ab_n) + \lim(Ab_n - AB)$$

(Note that we cannot write this yet because we have not shown that these two sequences are convergent)

So let's examine these two sequences. The second one is easier since we have proved proposition 5.1.14:

$$\lim b_n = B \implies \lim(b_n - B) = 0$$

Thus  $\lim(Ab_n - AB) = A \lim(b_n - B) = 0$ .

As for the first one, we want to show that  $\lim(a_n - A)b_n = 0$ . Since we know that  $b_n$  is itself a converging sequence, thus in particular  $b_n$  is bounded, so suppose that  $M > 0$  is a bound of  $b_n$ , i.e. for all natural number  $n$ ,  $|b_n| \leq M$ .

Since  $\lim a_n = A$ , for all  $\varepsilon > 0$ , there exists  $N$  such that for all  $n > N$ ,  $|a_n - A| < \frac{\varepsilon}{M}$ .

Combining the two above, we then conclude that for all  $\varepsilon > 0$ , there exists  $N$  such that for all  $n > N$ ,

$$|a_n b_n - Ab_n| = |(a_n - A)b_n| < \frac{\varepsilon}{M} \cdot M = \varepsilon.$$

Therefore, this implies that  $\lim(a_n b_n - Ab_n) = 0$ .

Since we have shown that the two parts are equal to 0, we can conclude that  $\lim(a_n b_n - AB) = 0$ . □

**Proposition 5.1.17.** If  $\{a_n\}$  and  $\{b_n\}$  are converging,  $b_n$  is never 0 and  $\lim b_n \neq 0$ , then

$$\lim \frac{a_n}{b_n} = \frac{\lim a_n}{\lim b_n}.$$

**Proof** Since we already have third proposition, it is sufficient for us to show that  $\lim \frac{1}{b_n} = \frac{1}{\lim b_n}$ .

Let  $b = \lim b_n$ , then we consider the limit

$$\lim \left( \frac{1}{b_n} - \frac{1}{b} \right) = \lim \left( \frac{b - b_n}{b_n b} \right).$$

Again, the important term here is  $b - b_n$ , but there is an extra term of  $\frac{1}{b_n b}$ , so we'll need to control this.

Since we need this to be bounded, we actually cannot have  $b_n$  to be close to 0. The good thing here is that  $b \neq 0$ , so we can restrict  $b_n$  to be close enough to  $b$  so that it stays away from 0.

So we can first pick  $N_1$  such that for all  $n > N_1$ ,

$$|b_n - b| < \frac{|b|}{2}.$$

Then

$$\begin{aligned} |b_n b - b^2| &< \frac{b^2}{2} \\ \frac{b^2}{2} &< b_n b < \frac{3b^2}{2} \end{aligned}$$

This shows that if  $n > N_1$ ,  $b_n b$  would always be positive, and  $\frac{1}{b_n b} < \frac{2}{b^2}$ .

Let  $M = \frac{2}{b^2}$ , then we may refer back to the original statement

$$\left| \frac{b - b_n}{b_n b} \right| < M |b - b_n|$$

We pick  $N_2$  such that for all  $n > N_2$ ,  $|b_n - b| < \frac{\varepsilon}{M}$ .

Let  $N = \max\{N_1, N_2\}$ , then for all  $n > N$ ,

$$\left| \frac{b - b_n}{b_n b} \right| < M \cdot \frac{\varepsilon}{M} = \varepsilon.$$

□

Now let's talk a little bit about the arithmetic properties of limsups and liminfs: There are quite a number of differences for this; essentially the arithmetical properties aren't as well-behaved as the more specific case of limits: (i)  $\limsup k a_n = k \limsup a_n$  holds if  $k > 0$ . However, if  $k < 0$ , then  $\limsup k a_n = k \liminf a_n$ .

(ii)  $\limsup(a_n + b_n)$  is in general not equal to  $\limsup a_n + \limsup b_n$ . However, we do have the following:

$$\limsup(a_n + b_n) \leq \limsup a_n + \limsup b_n$$

Moreover,  $\limsup(a_n + b_n)$  may be bounded from below as follows:

$$\limsup(a_n + b_n) \geq \limsup a_n + \liminf b_n$$

Your homework for today is to write down the analogous properties for liminf, and to prove (i) and (ii).

Now you should try to prove (i) for liminf as well; as for (ii), try to explain why properties (i), (ii) for limsup and property (i) for liminf would imply property (ii) for liminf.

**Problem 30.** Let  $\{x_n\}$  be a sequence of real numbers and let  $\alpha \geq 2$  be a constant. Define the sequence  $\{y_n\}$  as follows:

$$y_n = x_n + \alpha x_{n+1}, n = 1, 2, \dots$$

Show that if  $\{y_n\}$  is convergent, then  $\{x_n\}$  is also convergent.

## §5.2 Series in $\mathbb{R}$ ( $\mathbb{C}$ )

### §5.2.1 Definition and basic properties

### §5.2.2 Comparison test

### §5.2.3 Root and ratio tests

### §5.2.4 Addition and multiplication of series

### §5.2.5 Rearrangement

# 6 Continuity

## §6.1 Limit of Functions

Assume  $(X, d_X)$  is metric space and  $E \subset X$  is a subset of  $X$ . Then the metric  $d_X$  induces a metric on  $E$ . We now consider another metric space  $(Y, d_Y)$ . A map  $f : E \rightarrow Y$  is also called a function over  $E$  with values in  $Y$ . In particular, if  $Y = \mathbb{R}$ , then  $f$  is called a real-valued function; and if  $Y = \mathbb{C}$ ,  $f$  is called a complex-valued function.

**Definition 6.1.1.** Consider a limit point  $p \in E$  and a point  $q \in Y$ . We say the **limit** of the function  $f(x)$  at  $p$  is  $q$ , denoted as

$$\lim_{x \rightarrow p} f(x) = q$$

if for any  $\varepsilon > 0$ , there exists some  $\delta > 0$  such that for any  $x \in E$  with  $0 < d_X(x, p) < \delta$ , there is

$$d_Y(f(x), q) < \varepsilon.$$

We can recast this definition in terms of limits of sequences:

$$\lim_{n \rightarrow \infty} f(p_n) = q$$

for every sequence  $(p_n) \in E$  so that  $p_n \neq p$  and  $\lim_{n \rightarrow \infty} p_n = p$ .

By the same proofs as for sequences, limits are unique, and in  $\mathbb{R}$  they add/multiply/divide as expected.

**Definition 6.1.2.**  $f$  is **continuous** at  $p$  if

$$\lim_{x \rightarrow p} f(x) = f(p).$$

In the case where  $p$  is not a limit point of the domain  $E$ , we say  $f$  is continuous at  $p$ . If  $f$  is continuous at all points of  $E$ , then we say  $f$  is continuous on  $E$ .

The sequential definition of continuity follows almost directly from the sequential definition of limits:  $f$  is continuous at  $p$  if for every sequence  $x_n$  converging to  $p$ , the sequence  $f(x_n)$  converges to  $f(p)$ .



**§6.2 Continuous Functions**

**§6.3 Continuity and Compactness**

**§6.4 Continuity and Connectedness**

**§6.5 Discontinuities**

**§6.6 Monotonic Functions**

**§6.7 Infinite Limits and Limits at Infinity**

# 7 Differentiation

We focus on real valued functions defined on open or closed intervals.

## §7.1 The Derivative of a Real Function

**Definition 7.1.1.** A function  $f : [a, b] \rightarrow \mathbb{R}$  is called **differentiable** at  $x_0 \in [a, b]$ , if the limit of the function

$$\phi(t) := \frac{f(t) - f(x_0)}{t - x_0}, \quad a < t < b, t \neq x_0$$

exists as  $t \rightarrow x_0$ . For this case, we write

$$f'(x_0) = \lim_{t \rightarrow x_0} \phi(t) = \lim_{t \rightarrow x_0} \frac{f(t) - f(x_0)}{t - x_0}. \quad (7.1)$$

The function  $f$  is differentiable over  $[a, b]$  if it is differentiable for each  $x \in [a, b]$ . It induces the function

$$\frac{df}{dx} = f' : [a, b] \rightarrow \mathbb{R},$$

which is called the **derivative** of  $f$ .

**Theorem 7.1.2.** If  $f : [a, b] \rightarrow \mathbb{R}$  is differentiable at  $x_0 \in [a, b]$ , then it must be continuous at  $x_0$ .

**Proof** As  $t \rightarrow x$ ,

$$f(t) - f(x) = \frac{f(t) - f(x)}{t - x} \cdot (t - x) \rightarrow f'(x) \cdot 0 = 0.$$

□

**Remark** The converse of this theorem is not true. It is easy to construct continuous functions which fail to be differentiable at isolated points.

**Notation** We use  $C_1[a, b]$  to denote the set of differentiable functions over  $[a, b]$  whose derivative is continuous. More generally, we use  $C_k[a, b]$  to denote the set of functions whose  $k$ -th ordered derivative is continuous. In particular,  $C_0[a, b]$  is the set of continuous functions over  $[a, b]$ .

Later on when we talk about properties of differentiation such as the intermediate value theorems, we usually have the following requirement on the function:

$f$  is a continuous function on  $[a, b]$  which is differentiable in  $(a, b)$ .

**Theorem 7.1.3** (Differentiation rules). Suppose  $f, g : [a, b] \rightarrow \mathbb{R}$  are differentiable at  $x_0 \in [a, b]$ . Then  $f \pm g$ ,  $fg$  and  $\frac{f}{g}$  (when  $g(x_0) \neq 0$ ) are differentiable at  $x_0$ . Moreover,

1.  $(f \pm g)'(x_0) = f'(x_0) \pm g'(x_0)$ ;
2.  $(fg)'(x_0) = f'(x_0)g(x_0) + f(x_0)g'(x_0)$ ;

$$3. \left( \frac{f}{g} \right)'(x_0) = \frac{f'(x_0)g(x_0) - f(x_0)g'(x_0)}{g(x_0)^2}$$

**Proof** We take (2) as an example.

We calculate

$$\begin{aligned} \frac{f(x)g(x) - f(x_0)g(x_0)}{x - x_0} &= \frac{(f(x) - f(x_0))g(x) + f(x_0)(g(x) - g(x_0))}{x - x_0} \\ &= \frac{f(x) - f(x_0)}{x - x_0} \cdot g(x) + f(x_0) \cdot \frac{g(x) - g(x_0)}{x - x_0} \\ &\rightarrow f'(x_0)g(x_0) + f(x_0)g'(x_0) \text{ as } x \rightarrow x_0 \end{aligned}$$

where we use  $f$  and  $g$  are differentiable at  $x_0$  and Theorem 7.1.2.  $\square$

**Theorem 7.1.4** (Chain rule). Let  $f : [a, b] \rightarrow \mathbb{R}$  be a real-valued function that is differentiable at  $x_0 \in [a, b]$ . Let  $g$  be a real-valued function defined on an interval that contains  $f([a, b])$ , and  $g$  is differentiable at  $f(x_0)$ . Then the composition

$$h(x) := g \circ f(x) := g(f(x)) : [a, b] \rightarrow \mathbb{R}$$

is differentiable at  $x_0$  and the derivative at  $x_0$  can be calculated as

$$h'(x_0) = g'(f(x_0)) f'(x_0).$$

**Proof** We know that

$$f'(x) = \lim_{t \rightarrow x} \frac{f(t) - f(x)}{t - x},$$

so under the assumption that  $t$  stays within the domain of  $f$ ,  $\frac{f(t) - f(x)}{t - x}$  should be a good approximation to  $f'(x)$ .

To actually quantify this, let  $u(t) = \frac{f(t) - f(x)}{t - x} - f'(x)$ .

Then the differentiability of  $f$  tells us that  $\lim_{t \rightarrow x} u(t) = 0$ .

Similarly, let  $v(s) = \frac{g(s) - g(y)}{s - y} - g'(y)$ , then  $\lim_{s \rightarrow y} v(s) = 0$ , as long as  $s$  stays in the domain of  $g$ .

What's nice here is that we can let  $s = f(t)$ , then by our assumption  $s$  always stays in the domain of  $g$ , so nothing fishy will happen.

Ah I forgot a small detail here. Additionally we also need to define  $u(x) = 0$  and  $v(y) = 0$ .

Now let  $h(t) = g(f(t))$ , then  $h$  is defined on  $[a, b]$ , and we deduce that

$$h(t) - h(x) = (t - x)[f'(x) + u(t)][g'(y) + v(s)]$$

We then check that

$$\lim_{t \rightarrow x} \frac{h(t) - h(x)}{t - x} = \lim_{t \rightarrow x} [f'(x) + u(t)][g'(y) + v(s)] = f'(x)g'(f(x))$$

and we are done.  $\square$

**Example 7.1.5.** One of the best (worst?) family of pathological examples in calculus are functions of the form

$$f(x) = x^p \sin \frac{1}{x}.$$

- For  $p = 1$ , the function is continuous and differentiable everywhere other than  $x = 0$ .
- For  $p = 2$ , the function is differentiable everywhere, but the derivative is discontinuous.

Other more advanced pathological results (just for fun):

- The graph for  $y = \sin \frac{1}{x}$  on  $(0, 1]$ , together with the interval  $[-1, 1]$  on the  $y$ -axis, is a connected closed set that is not path-connected.

- For  $0 < p < 1$ , we obtain functions that are continuous and bounded, but the graphs are of infinite length (ps. I think that this is also true for  $p = 1$ ).

Regarding continuous but not differentiable functions, a more pathological example is the Weierstrass function, which is continuous everywhere over  $\mathbb{R}$  but differentiable nowhere.

## §7.2 Mean Value Theorems

**Definition 7.2.1.** Let  $f$  be a real valued function defined over a metric space  $X$ . We say  $f$  has a **local maximum** at  $x_0 \in X$  if  $\exists \delta > 0$  s.t.  $\forall x \in B_\delta(x_0)$ ,

$$f(x_0) \geq f(x).$$

Similarly, we say  $f$  has **local minimum** at  $x_0 \in X$  if  $\exists \delta > 0$  s.t.  $\forall x \in B_\delta(x_0)$ ,

$$f(x_0) \leq f(x).$$

**Definition 7.2.2.** For a function  $f : (a, b) \rightarrow \mathbb{R}$ , a point  $x_0 \in [a, b]$  is called a **critical point** if  $f$  is not differentiable at  $x_0$  or  $f'(x_0) = 0$ .

**Theorem 7.2.3.** Assume  $f$  is defined over  $[a, b]$ . If  $f$  has a local maximum or local minimum at some  $x_0 \in (a, b)$ , then  $x_0$  is a critical point of  $f$ .

**Proof** If  $f$  is not differentiable at  $x_0$ , we are done. Assume now  $f$  is differentiable at  $x_0$  and  $x_0$  is a local maximum.

Then  $\exists \delta > 0$  s.t.  $\forall x \in B_\delta(x_0)$ ,

$$f(x_0) \geq f(x).$$

It follows

$$\frac{f(x) - f(x_0)}{x - x_0} \begin{cases} \geq 0 & x_0 - \delta < x < x_0 + \delta \\ \leq 0 & x_0 < x < x_0 + \delta \end{cases}$$

Further since  $f'(x_0)$  exists, there is

$$f'(x_0-) \geq 0, \quad f'(x_0+) \leq 0,$$

but  $f'(x_0-) = f'(x_0+) = f'(x_0)$ . Hence  $f'(x_0) = 0$ . □

**Theorem 7.2.4** (Fermat's Theorem (Interior Extremum Theorem)). If the differential exists, then by comparing the left and right limits it is easy to see that the differential for a local maximum/maximum can only be 0.

To summarize in four words: Local extrema are stationary

There are three mean value theorems, from specific to general:

1. Rolle's Theorem
2. (Lagrange's) Mean Value Theorem
3. Generalised (Cauchy's) Mean Value Theorem

**Theorem 7.2.5** (Rolle's Theorem). If  $f$  is continuous on  $[a, b]$ , differentiable in  $(a, b)$  and  $f(a) = f(b)$ , then there exists  $c \in (a, b)$  such that

$$f'(c) = 0.$$

**Proof** Let  $h(x)$  be a function defined on  $[a, b]$  where  $h(a) = h(b)$ .

The idea is to show that  $h$  has a local maximum/minimum, then by Fermat's Theorem this will then be the stationary point that we're trying to find.

First note that  $h$  is continuous on  $[a, b]$ , so  $h$  must have a maximum  $M$  and a minimum  $m$ .

If  $M$  and  $m$  were both equal to  $h(a) = h(b)$ , then  $h$  is just a constant function and so  $h'(x) = 0$  everywhere. Otherwise,  $h$  has a maximum/minimum that is not  $h(a) = h(b)$ , so this extremal point lies in  $(a, b)$ .

In particular, this extremal point is also a local extremum. Since  $h$  is differentiable on  $(a, b)$ , by Fermat's theorem this extremum point is stationary, thus Rolle's Theorem is proven.  $\square$

**Theorem 7.2.6** (Mean Value Theorem). If  $f$  is continuous on  $[a, b]$  and differentiable in  $(a, b)$ , then there exists  $c \in (a, b)$  such that

$$f'(c) = \frac{f(b) - f(a)}{b - a}.$$

Exercise 2: Show that the Mean Value Theorem results directly from Rolle's Theorem (the other direction is trivial) : This isn't a very significant exercise because we're going to prove something more general

**Theorem 7.2.7** (Generalised Mean Value Theorem). If  $f$  and  $g$  are continuous on  $[a, b]$  and differentiable in  $(a, b)$ , then there exists  $c \in (a, b)$  such that

$$\frac{f'(c)}{g'(c)} = \frac{f(b) - f(a)}{g(b) - g(a)}.$$

Now we return to the proof of the generalized MVT

We set the function  $h(t) = [f(b) - f(a)]g(t) - [g(b) - g(a)]f(t)$ , then  $h$  is continuous on  $[a, b]$  and differentiable on  $(a, b)$

Moreover,  $h(a) = f(b)g(a) - f(a)g(b) = h(b)$ , thus by Rolle's Theorem, there exists  $c \in (a, b)$  such that  $h'(c) = 0$ , i.e.  $[g(b) - g(a)]f'(c) = [f(b) - f(a)]g'(c)$

Corollary: If  $f$  and  $g$  are continuous on  $[a, b]$  and differentiable in  $(a, b)$ , and  $g'(x) \neq 0$  for all  $x \in (a, b)$ , then there exists

$$c \in (a, b) \text{ s.t. } f'(c)/g'(c) = [f(b) - f(a)]/[g(b) - g(a)]$$

This form of the generalized MVT will be used to prove the most beloved rule of high school students exercises for the Mean Value Theorem

#### Exercise 7.2.1

Let  $f$  and  $g$  be continuous on  $[a, b]$  and differentiable on  $(a, b)$ . If  $f'(x) = g'(x)$ , then  $f(x) = g(x) + C$ .

#### Exercise 7.2.2

Given that  $f(x) = x^\alpha$  where  $0 < \alpha < 1$ . Prove that  $f$  is uniformly continuous on  $[0, +\infty)$ .

#### Exercise 7.2.3: Olympiad level

Let  $f$  be a function continuous on  $[0, 1]$  and differentiable on  $(0, 1)$  where  $f(0) = f(1) = 0$ . Prove that there exists  $c \in (0, 1)$  such that

$$f(x) + f'(x) = 0.$$

## §7.3 Darboux's Theorem

Darboux's Theorem implies some sort of a 'intermediate value' property of derivatives that is similar to continuous functions

This is Theorem 5.12 in the book

Now first and foremost, the requirement for this statement is that  $f$  must be differentiable on  $[a, b]$ , not just in  $(a, b)$  Otherwise  $f'(a)$  and  $f'(b)$  may not make sense : One common theme in many of these problems is to construct auxiliary functions Suppose that  $f'(a) < \lambda < f'(b)$ , then we construct the auxiliary function

$g(x) = f(x) - \lambda x$  : Then we only need to find a point  $x \in (a, b)$  such that  $g'(x)=0$  : This means that we only need to find a local maximum/minimum, which by Fermat's Theorem has to be a stationary point as well : Now we look at the values of  $g$  near  $a$  and  $b$  : Exercise 1: Using the fact that  $g'(a) < 0$  and  $g'(b) > 0$ , show that  $a$  and  $b$  are local maxima of  $g$

Here we regard  $g$  as simply a function on  $[a, b]$ , so we only need to show that  $a, b$  are maximum and corresponding semi-open neighbourhoods  $[a, a + \varepsilon)$  and  $(b - \varepsilon, b]$  : Let  $m=g'(a)<0$  be the slope of the tangent at  $a$  : Then  $\lim_{h \rightarrow 0^+} [g(a+h)-g(a)]/h=m<0$  : This means that there should exist  $\delta > 0$  such that for  $0 < h < \delta$ ,  $[g(a+h)-g(a)]/h < m/2 < 0$  : Now we can rewrite the above as  $g(a+h) < g(a) + mh/2$  : Since  $m < 0$  and  $h > 0$ , we obtain  $g(a+h) < g(a)$  for  $0 < h < \delta$  : Thus this proves that  $x=a$  is a local maximum of  $g$  A similar proof applies for  $x=b$  : Now since  $g$  is differentiable on  $[a, b]$ , in particular it has to be continuous on  $[a, b]$  : Since  $[a, b]$  is compact,  $g([a, b])$  is compact in  $\mathbb{R}$  and thus  $g$  has both maximum and minimum values in  $[a, b]$  : Here we'll just focus on the minimum value : As we've shown,  $x=a$  is a 'strict' local maxima, in the sense that for any point  $x \in (a, a + \varepsilon)$ , we actually have the strict inequality  $g(x) < g(a)$  : This means that  $x=a$  cannot be a local minimum : Similarly,  $x=b$  cannot be a local minimum, and therefore  $g$  achieves its minimum strictly inside  $(a, b)$  : Only then we can say that this local minimum is stationary (This will not work otherwise; note that  $a$  and  $b$  are both local maxima but are not stationary points of  $g$ ) : An interesting implication of Darboux's Theorem is that if  $f$  is differentiable on  $[a, b]$ , then  $f'$  cannot have simple discontinuities (removable or jump discontinuities), simply because these discontinuities do not allow this 'intermediate value' property : However, we should recall certain pathological examples like  $f(x)=x^2 \sin 1/x$  ( $f(0)=0$ ) Here  $f'(0)=\lim_{h \rightarrow 0} [x^2 \sin 1/x-0]/x=0$ , but  $f'(x)=2x \sin 1/x - \cos 1/x$ , so  $f'$  is discontinuous at  $x=0$

## §7.4 L'Hopital's Rule

First, a counterexamples : Let's say that we apply this rule to  $\lim_{x \rightarrow \infty} \frac{\sin x}{x}$  : Then we have

$$\lim_{x \rightarrow \infty} \frac{\sin x}{x} = \lim_{x \rightarrow \infty} \frac{\cos x}{1}$$

The limit on the RHS doesn't exist because  $\cos x$  oscillates between  $-1$  and  $1$  : However, the limit on the LHS does in fact exist and is equal to  $0$  : So this tells us that there are certain cases where we can apply L'Hopital, and other cases where we can't That being said, the case that we can apply the rule is actually the more useful case, so this situation does not jeopardize the effectiveness of L'Hopital

The entire statement is consequently rather long, so we'll split it into a few sections

1.  $f$  and  $g$  are differentiable in  $(a, b)$  and  $g'(x) \neq 0$  in  $(a, b)$  (or at least in a small neighbourhood of  $a$ )
2.  $f(x)/g(x)$  is an indeterminate of the form  $0/0$  or  $\frac{\infty}{\infty}$  (Now for the second one here we only really need  $g(x) \rightarrow \infty$ , but if  $f(x)$  does not approach infinity then the limit would simply be zero, so L'Hopital's Rule would not be required here)
3.  $\lim_{x \rightarrow a} f'(x)/g'(x) = A$  This is the most important one : From this we obtain  $\lim_{x \rightarrow a} f(x)/g(x) = A$  : So for example, let's say that we want to calculate the following limit:  $\lim_{x \rightarrow 0} (\sin x - x)/x^3$  : Repeated application of L'Hopital gives  $\lim_{x \rightarrow 0} (\sin x - x)/x^3 = \lim_{x \rightarrow 0} (\cos x - 1)/3x^2 = \lim_{x \rightarrow 0} -\sin x/6x = -1/6$  : Now what we're really doing here is that, first we know that  $\lim_{x \rightarrow 0} \sin x/x = 1$ , so  $\lim_{x \rightarrow 0} -\sin x/6x = -1/6$  : Then by L'Hopital,  $\lim_{x \rightarrow 0} (\cos x - 1)/3x^2 = -1/6$  : Finally, again by L'Hopital,  $\lim_{x \rightarrow 0} (\sin x - x)/x^3 = -1/6$  : So, one very important thing to take note is that if you're calculating some complicated limit and you end up with the conclusion that it doesn't exist, you must make sure that you have not used L'Hopital during the process, because the rule never applies in such situations : As a side note, from the above calculation we see that as  $x \rightarrow 0$ ,  $\sin x \approx x - \frac{x^3}{6}$  This will later lead to the discussion of the Taylor series of  $\sin x$

Now the entire proof is quite tedious because there's actually eight main cases to think of 1.  $\frac{0}{0}$  or  $\frac{\infty}{\infty}$  2.  $a$  is normal or  $a = -\infty$  3.  $A$  is normal or  $A = \pm\infty$

We'll only prove the most basic one here:  $0/0$ ,  $a$  and  $A$  are normal This is the case which will be required for Taylor series

First we define  $f(a)=g(a)=0$ , so that  $f$  and  $g$  are continuous at  $x = a$

Now let  $x \in (a, b)$ , then  $f$  and  $g$  are continuous on  $[a, x]$  and differentiable in  $(a, x)$  : Thus by Cauchy's

Mean Value Theorem, there exists  $\xi \in (a, x)$  such that

$$\frac{f'(\xi)}{g'(\xi)} = \frac{f(x) - f(a)}{g(x) - g(a)} = \frac{f(x)}{g(x)}$$

For each  $x$ , we pick  $\xi$  which satisfies the above, so that  $\xi$  may be seen as a function of  $x$  satisfying  $a < \xi(x) < x$

Then by squeezing we have  $\lim_{x \rightarrow a^+} \xi(x) = a$ .

Since  $\frac{f'}{g'}$  is continuous near  $a$ , the theorem regarding the limit of composite functions give

$$\lim_{x \rightarrow a^+} \frac{f(x)}{g(x)} = \lim_{x \rightarrow a^+} \frac{f'(\xi)}{g'(\xi)} = \lim_{x \rightarrow a^+} \left( \frac{f'}{g'} \right) (\xi(x)) = A$$

Now the same reasoning can be used for  $b$  where we will use  $\lim_{x \rightarrow b^-}$  to replace all the  $\lim_{x \rightarrow a^+}$ , and  $\xi$  will be a function which maps to  $(x, b)$ .

## §7.5 Taylor Expansion

The main expression is as follows:

$$f(x) = f(a) + \frac{f'(a)}{1!}(x-a) + \frac{f''(a)}{2!}(x-a)^2 + \frac{f'''(a)}{3!}(x-a)^3 + \dots \quad (7.2)$$

So for example we have the following (we've used the ones for  $e^x$  and  $\ln x$  for generating functions):

$$\begin{aligned} e^x &= 1 + x + \frac{x^2}{2!} + \frac{x^3}{3!} + \dots \\ \sin x &= x - \frac{x^3}{3!} + \frac{x^5}{5!} - \frac{x^7}{7!} + \dots \\ \cos x &= 1 - \frac{x^2}{2!} + \frac{x^4}{4!} - \frac{x^6}{6!} + \dots \\ \ln(1+x) &= x - \frac{x^2}{2} + \frac{x^3}{3} - \frac{x^4}{4} + \dots \end{aligned}$$

There's a lot of things to say about these equations, for example the one for  $\ln(1+x)$  only works for  $|x| < 1$

Also, if you want the RHS of the expression to be an infinite power series,  $f(x)$  has to be smooth (infinitely differentiable)

Even then, the power series may never converge to  $f(x)$  at any interval, no matter how small. The most common example given here is  $f(x) = e^{-\frac{1}{x^2}}$  ( $f(0)=0$ ); the Taylor series for  $f(x)$  is just 0

Now sometimes we don't actually have that nice of a property for  $f$ , we're often given that fact that  $f$  is only finitely differentiable

Then we will have something along the lines of

$$f(x) \approx f(a) + \frac{f'(a)}{1!}(x-a) + \frac{f''(a)}{2!}(x-a)^2 + \dots + \frac{f^{(n)}(a)}{n!}(x-a)^n$$

where  $f^{(n)}$  denotes the  $n$ -th differential.

There are two main forms of the statement regarding the error between the original function and the Taylor series estimate

The simpler form is what's known as the Peano form: Given that  $f$  is  $n$  times differentiable at  $a$ , then

$$f(x) = f(a) + \frac{f'(a)}{1!}(x-a) + \frac{f''(a)}{2!}(x-a)^2 + \dots + \frac{f^{(n)}(a)}{n!}(x-a)^n + o((x-a)^n)$$

To show this, we only need to show that we have the following limit:

$$\lim_{x \rightarrow a} \frac{f(x) - f(a) + \frac{f'(a)}{1!}(x-a) + \frac{f''(a)}{2!}(x-a)^2 + \dots + \frac{f^{(n)}(a)}{n!}(x-a)^n}{(x-a)^n} = 0$$

The basic idea is to use the L'Hopital Rule  $n$  times. The numerator becomes  $f^{(n)}(x) - f^{(n)}(a)$  which approaches 0, whereas the denominator is just  $n!$ , so the limit exists and is equal to 0.

However, we need to verify all the necessary conditions for L'Hopital : Here the main problem is that we don't know if we have the  $0/0$  indeterminate at each step, so we'll need to check this for the  $k$ -th step where  $k=1, \dots, n$

Fortunately, the  $k$ -th derivative of the numerator is  $f^{(k)}(x) - f^{(k)}(a) - (x-a)F_k(x)$  where  $F_k$  is just a bunch of random stuff, so the numerator approaches 0 as  $x \rightarrow a$ . The  $k$ -th derivative of the denominator is  $n(n-1)\dots(n-k+1)(x-a)^{n-k}$  so it also approaches 0, and we're done

The other form is actually a family of similar statements which gives more precise values for the error. The Peano form has a fundamental obstacle when used in approximation, we don't have any control on the size of the final term other than its asymptotic behaviour : We'll be talking about the one given in the book, known as the Lagrange form : Given that  $f$  is  $n$  times differentiable on  $(a, b)$  such that  $f^{(n-1)}$  is continuous on  $[a, b]$ , then

$$f(x) = f(a) + \frac{f'(a)}{1!}(x-a) + \frac{f''(a)}{2!}(x-a)^2 + \dots + \frac{f^{(n-1)}(a)}{(n-1)!}(x-a)^{n-1} + \frac{f^{(n)}(\xi)}{n!}(x-a)^n$$

Just like in L'Hopital, we intuitively think of  $(a, b)$  as just a very small interval at the right hand side of  $x=a$  : Here we are giving up on the second final term of Peano by combining it with the infinitesimal (small  $o$ ) term to give an accurate description of the error

For the proof of this one we'll be using Cauchy's MVT

Fix any  $x \in (a, b)$ , then we construct the functions

$$F(t) = f(x) - \left( f(t) + \frac{f'(t)}{1!}(x-t) + \frac{f''(t)}{2!}(x-t)^2 + \dots + \frac{f^{(n-1)}(t)}{(n-1)!}(x-t)^{n-1} \right)$$

$$G(t) = (x-t)^n$$

We calculate  $F'(t)$  as follows:

$$-[f'(t) + \frac{f''(t)}{1!} - f'(t) + \frac{f'''(t)}{2!} - \frac{f''(t)}{1!} + \dots + \frac{f^{(n)}(t)}{(n-1)!}(x-t)^{n-1} - \frac{f^{(n-1)}(t)}{(n-2)!}(x-t)^{n-2}] = -\frac{f^{(n)}(t)}{(n-1)!}(x-t)^{n-1}$$

$G'(t) = -n(x-t)^{n-1}$ , so we have

$$\frac{F'(t)}{G'(t)} = \frac{f^{(n)}(t)}{n!}$$

The main reason for why we come up with the strange-looking  $F$  and  $G$  is that we specifically swap out  $a$  for  $t$  so that  $F(x) = G(x) = 0$ , in hopes of getting rid of  $x$ :

We apply Cauchy's MVT to  $F$  and  $G$  on  $[a, x]$ , so that we obtain  $\xi \in (a, x)$  satisfying

$$\frac{F'(\xi)}{G'(\xi)} = \frac{F(x) - F(a)}{G(x) - G(a)} = \frac{F(a)}{G(a)}.$$

Thus the Lagrange form of the remainder is given by

$$F(a) = \frac{f^{(n)}(\xi)}{n!}G(a).$$

Theorem 5.19 is important, so do go through that proof as an exercise



# 8

## The Riemann–Stieltjes Integral

### §8.1 Definition of Riemann–Stieltjes Integral

Assume  $[a, b]$  is a closed interval in  $\mathbb{R}$ . By a **partition**  $P$ , we mean a finite set of points

$$P : a = x_0 \leq x_1 \leq \cdots \leq x_{n-1} \leq x_n = b.$$

Assume  $f$  is a bounded real-valued function over  $[a, b]$  and  $\alpha$  is an increasing function over  $[a, b]$ . Denote by

$$M_i = \sup_{[x_{i-1}, x_i]} f(x), \quad m_i = \inf_{[x_{i-1}, x_i]} f(x)$$

and by

$$\Delta\alpha_i = \alpha(x_i) - \alpha(x_{i-1}).$$

Define the **upper sum** of  $f$  with respect to the partition  $P$  and  $\alpha$  as

$$U(f, \alpha; P) = \sum_{i=1}^n M_i \Delta\alpha_i$$

and the **lower sum** of  $f$  with respect to the partition  $P$  and  $\alpha$  as

$$L(f, \alpha; P) = \sum_{i=1}^n m_i \Delta\alpha_i.$$

Define the upper Riemann–Stieltjes integral as

$$\int_a^b f(x) d\alpha(x) := \inf_P U(f, \alpha; P)$$

and the lower Riemann–Stieltjes integral as

$$\int_a^b f(x) d\alpha(x) := \sup_P L(f, \alpha; P).$$

It is easy to see from definition that

$$\int_a^b f(x) d\alpha(x) \leq \int_a^b f(x) d\alpha(x).$$

**Definition 8.1.1.** A function  $f$  is **Riemann–Stieltjes integrable** with respect to  $\alpha$  over  $[a, b]$ , if

$$\int_a^b f(x) d\alpha(x) = \int_a^b f(x) d\alpha(x).$$

**Notation** We use  $\int_a^b f(x) d\alpha(x)$  to denote the common value, and call it the Riemann–Stieltjes of  $f$  with respect to  $\alpha$  over  $[a, b]$ .

**Notation** We use the notation  $R_\alpha[a, b]$  to denote the set of Riemann–Stieltjes integrable functions with respect to  $\alpha$  over  $[a, b]$ .

In particular, when  $\alpha(x) = x$ , we call the corresponding Riemann–Stieltjes integration the **Riemann integration**, and use  $R[a, b]$  to denote the set of Riemann integrable functions.

Now we move on to integrability conditions for  $f$ . The first one looks a lot like the  $\varepsilon$ – $N$  or  $\varepsilon$ – $\delta$  definition of limits:

**Theorem 8.1.2.**  $f \in R_\alpha[a, b]$  if and only if for each  $\varepsilon > 0$ , there exists some partition  $P$  such that

$$U(f, \alpha; P) - L(f, \alpha; P) < \varepsilon.$$

**Proof** If  $f \in R_\alpha[a, b]$ , then the two Darboux integrals would be equal and so

$$\inf_P U(f, P, \alpha) = \sup_P L(f, P, \alpha) = I$$

where  $I$  is the RS-integral of  $f$  wrt  $\alpha$ .

Then for any  $\varepsilon > 0$ , we can find  $P_1$  and  $P_2$  such that

$$I \leq U(f, P_1, \alpha) < I + \frac{\varepsilon}{2}$$

and

$$I - \frac{\varepsilon}{2} < L(f, P_2, \alpha) \leq I.$$

Let  $P'$  be the common partition of  $P_1$  and  $P_2$ , then

$$L(f, P_2, \alpha) \leq L(f, P', \alpha) \leq U(f, P', \alpha) \leq U(f, P_1, \alpha)$$

And so from the above inequalities we have

$$I - \frac{\varepsilon}{2} < L(f, P', \alpha) \leq U(f, P', \alpha) < I + \frac{\varepsilon}{2}$$

Therefore  $U(f, P', \alpha) - L(f, P', \alpha) < \varepsilon$  (Sorry for the inconsistent notation) : On the other hand, for any given partition  $P$  we can always bound the Darboux integrals between the lower and upper Darboux sums over  $P$ .

So whenever we find a partition  $P$  such that

$$0 \leq U(f, P, \alpha) - L(f, P, \alpha) < \varepsilon$$

We immediately obtain that

$$0 \leq U(f, \alpha) - L(f, \alpha) < \varepsilon$$

(I'm not sure if this notation down here is universal so please stick to the original Darboux integral notations)

Since  $\varepsilon > 0$  is arbitrary, this gives us that the two Darboux integrals must be equal.  $\square$

#### Example 8.1.1: Dirichlet function

The Dirichlet function is given by

$$f(x) = \begin{cases} 1 & x \in \mathbb{Q} \\ 0 & x \notin \mathbb{Q} \end{cases}$$

We try to calculate the two on the interval  $[0, 1]$ .

The Dirichlet function is pathological because for each subinterval  $[x_{i-1}, x_i]$ , the supremum is always 1 and the infimum is always 0.

So no matter what partition we use,  $U(f, P)$  is always 1 whereas  $L(f, P)$  is always 0. This means that  $U(f) = 1$  and  $L(f) = 0$ , so there are two different values for “the integral of  $f$ ”.

This is like the case where we try to find the limit of the Dirichlet function where  $x$  is approaching any given real number  $r$ , there exists two sequences approaching  $r$  whose image approaches two different values.

Now, a very important and fun case about the more general RS-integral, which we'll discuss next week (do try the exercise yourself first)

### Exercise 8.1.1

The Heaviside step function  $H$  is a real-valued function defined by the following:

$$H(x) = \begin{cases} 0 & x < 0 \\ 1 & x \geq 0 \end{cases}$$

For the purpose of this question we assume the convention  $\infty \cdot 0 = 0$ .

- (a) Let  $f$  be a real-valued function over  $\mathbb{R}$ . Show that  $f \in \mathbb{R}_H[a, b]$  iff  $f$  is continuous at 0, and find the RS-integral  $\int_{-\infty}^{\infty} f dH$ .
- (b) Suppose that the definition for  $H$  is changed for  $x = 0$ , say  $H(0) = \frac{1}{2}$ . Show that the above result still holds.
- (c) Examine the RS-integral of  $f$  over  $\mathbb{R} \setminus \{0\}$  wrt  $H$ , where  $f$  is a real-valued function over  $\mathbb{R} \setminus \{0\}$  such that  $\lim_{x \rightarrow 0} f(x) = \infty$  or  $-\infty$ .

(You may read up on more information regarding the Heaviside function, and the (in)famous Dirac delta function)

To to recap last time, we defined when a function is Riemann integrable over a closed interval  $[a, b]$ , or more generally, when it is Riemann–Stieltjes integrable over  $[a, b]$  with respect to  $\alpha$

Now we'll need a key idea regarding partitions which will enable us to develop much deeper properties of integrals, namely refinements

This is quite simple: given that

$$P : a \leq x_0 \leq x_1 \leq \dots \leq x_n = b$$

and

$$P' : a \leq y_0 \leq y_1 \leq \dots \leq y_m = b,$$

we say that  $P'$  is a refinement of  $P$  if all points in  $P$  are in  $P'$ , i.e.  $\{x_i\}$  is a subset of  $\{y_j\}$ .

Furthermore, given two partitions  $P_1$  and  $P_2$ , we can construct a refinement  $P'$  which collects all partition points of  $P_1$  and  $P_2$ , giving what's called the **common partition** of  $P_1$  and  $P_2$ .

Intuitively, a refinement will give a better estimation than the original partition, so the upper and lower sums of a refinement should be more restrictive

So there's Theorem 6.4 which states that, indeed, one has

$$L(P, f, \alpha) \leq L(P', f, \alpha) \leq U(P', f, \alpha) \leq U(P, f, \alpha)$$

This can be proven without induction: Suppose that

$$P : a \leq x_0 \leq x_1 \leq \dots \leq x_n = b$$

and

$$P' : a \leq y_0 \leq y_1 \leq \dots \leq y_m = b.$$

Then there exists a strictly increasing sequence of indices  $j_0 = 0, j_1, \dots, j_n = m$  such that  $y_{j_k} = x_k$ .

Now consider each closed interval  $[x_{i-1}, x_i]$

Focusing on the upper sum, we have

$$\sup_{[x_{i-1}, x_i]} f \geq \sup_{[y_{k-1}, y_k]} f$$

for  $k = j_{i-1} + 1, \dots, j_i$ . This is because  $[y_{k-1}, y_k]$  is contained in  $[x_{i-1}, x_i]$

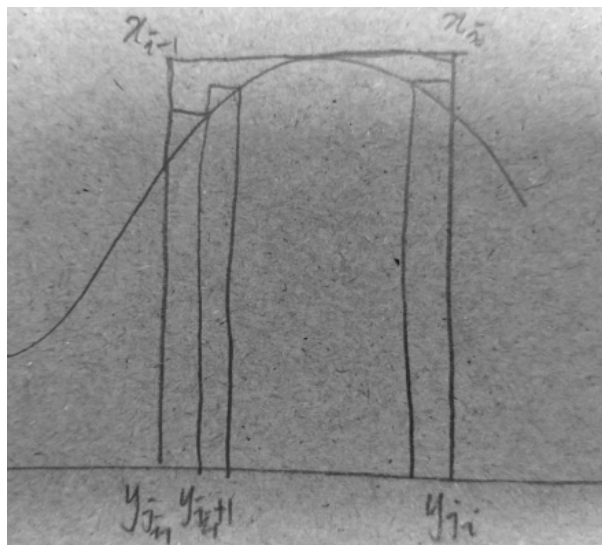


Figure 8.1: Partitions

Continuing from

$$\sup_{[x_{i-1}, x_i]} f \geq \sup_{[y_{k-1}, y_k]} f,$$

We then multiply by  $\alpha(y_k) - \alpha(y_{k-1})$  on both sides and then take the sum from  $k = j_{i-1} + 1$  to  $k = j_i$  : The RHS corresponds to the (weighted) sum of the thin rectangles that you see in the above picture : The LHS is actually a telescoping sum, and the sum would be

$$\left( \sup_{[x_{i-1}, x_i]} f \right) \cdot [\alpha(y_{j_i}) - \alpha(y_{j_{i-1}})] = \left( \sup_{[x_{i-1}, x_i]} f \right) \cdot [\alpha(x_i) - \alpha(x_{i-1})]$$

Finally, we take the sum from  $i = 1$  to  $i = n$  of the above inequality  $\text{LHS} \geq \text{RHS}$  (sorry I don't know of a better way to put it) We then obtain  $U(P, f, \alpha) \geq U(P', f, \alpha)$

(On the LHS we're collecting all the rectangles for the upper sum wrt  $P$ , but on the RHS we're collecting up collections of upper rectangles to obtain the entire collective of upper rectangles for the upper sum wrt  $P'$ ) : Lower sum is similar : Now, a lemma used to prove 6.5 Given any two partitions  $P_1$  and  $P_2$ , we have

$$L(P_1, f, \alpha) \leq U(P_2, f, \alpha)$$

So a lower sum will always be no larger than any other upper sum : So this includes the cases where we have the most refined of  $P_1$ 's and  $P_2$ 's, with no information regarding the partition points whatsoever To be honest, the result seems to be both intuitive and unclear at the same time

The key here is to use common refinements as a link for both sums The idea is stated in the proof of 6.5 and I don't think I need to elaborate further

What's nice here is that now we have two completely independent partitions  $P_1$  and  $P_2$ , so by fixing one partition, say  $P_2$ , and taking the 'limit' over the other (here we take the supremum over all possible  $P_1$ ) we then obtain an inequality between a Darboux integral and a Darboux sum (here it's the lower integral and an upper sum)

Since the Darboux integral is just a number, we can then safely take the 'limit' over the other partition to obtain the inequality in 6.5

As we've already defined,  $f$  is Riemann–Stieltjes integrable wrt  $\alpha$  if these two Darboux integrals coincide

Now we've been talking a lot about upper and lower sums because they're arguably the simplest way to define integrals, in the sense that there's not a whole lot of things that we could go wrong here By considering only upper and lower bound, we're essentially picking the most conservative route possible

It would be nice if we could just pick like one random point within each interval and consequently calculate the Riemann(-Stieltjes) sums

This method, of course, fails to be well defined for pathological functions like the Dirichlet function. On the other hand, by using upper and lower sums, we could give a persuasive explanation as to why the Dirichlet function is not Riemann integrable.

However, instead of throwing this idea away, there's actually a way for us to make this into a strict definition.

When we were talking about the sequential definition for limits of functions, we noted that there are certain scenarios where the limit cannot exist because there may be two distinct sequences that give different limits. Based on this observation, we then gave a reasonable condition as follows: " $\lim_{x \rightarrow a} f(x)$  exists and is equal to  $L$  iff for all sequences  $x_n$  converging but not containing  $a$ ,  $f(x_n)$  converges to  $L$ ".

Well here, it's actually the same kind of scenario. Given any partition  $P$ , we consider the Riemann sum  $\sum f(\xi_i)\Delta x_i$  where  $\xi_i$  is any point where  $x_{i-1} \leq \xi_i \leq x_i$ .

For the Dirichlet function over  $[0,1]$ , given any partition  $P$  (here we may assume that the partition points are distinct), we will always be able to specifically pick  $\xi_i, \eta_i \in [x_{i-1}, x_i]$  such that  $\xi_i$  is rational but  $\eta_i$  is irrational.

Then  $\sum f(\xi_i)\Delta x_i = 1$  but  $\sum f(\eta_i)\Delta x_i = 0$ .

Now be very mindful that this alone cannot be evidence that  $f$  is non-integrable. The key is that this somehow occurred for all partitions  $P$ , no matter how refined they are; for every single partition  $P$ , there exists two sets of 'representing points'  $\xi_i, \eta_i$  such that the two Riemann sums are constantly far apart (1 and 0 in this case).

Let  $\varepsilon_0 = 1$ , then this ultimately translates to the following: The Dirichlet function cannot be Riemann integrable because there exists some  $\varepsilon_0 > 0$ , such that for any given partition  $P$ , there exists two sets of representing points  $\xi_i, \eta_i$  such that their corresponding Riemann sums satisfy that

$$|\sum f(\xi_i)\Delta x_i - \sum f(\eta_i)\Delta x_i| \geq \varepsilon_0.$$

Now if we always pick the representatives such that  $\xi_i > \eta_i$  then we can neglect the absolute value.

So now, let's take the converse. A function  $f$  is said to be RS-integrable if for every  $\varepsilon > 0$ , there exists a partition  $P$ , such that for any two sets of representing points  $\xi_i, \eta_i$ , their corresponding Riemann sums satisfy that

$$\sum [f(\xi_i) - f(\eta_i)]\Delta x_i < \varepsilon$$

(The last one should be  $\Delta \alpha_i$  for RS-integrals, not  $\Delta x_i$ )

Unfortunately this is still not quite the correct definition according to Apostol, but we're pretty close. The problem with this definition is that it is too weak if we're considering general  $\alpha$  of bounded variation; if we were only talking about monotonically increasing  $\alpha$  then this will actually be an equivalent definition.

The official definition for the RS-integral wrt  $\alpha$  of bounded variation is as follows:

**Definition 8.1.3.** For every  $\varepsilon > 0$ , there exists a partition  $P$ , such that [For any refinement  $P'$  of  $P$ , and] For any two sets of representing points  $\xi_i, \eta_i$  [of  $P'$ ], their corresponding Riemann sums satisfy that

$$\sum [f(\xi_i) - f(\eta_i)]\Delta x_i < \varepsilon.$$

Now this definition is what mathematicians would refer to as a 'Cauchy' definition, since it defines a notion by comparing a pair of arbitrary values that are similar to one another, and if they agree in some sense then we say that that something satisfies some property.

The integral is then obtained as follows: If  $f$  were to satisfy the above Cauchy definition, then we may pick an arbitrary sequence of refinements

$$P_1 \subset P_2 \subset P_3 \subset \dots;$$

and for each partition we pick a set of representatives to obtain a sequence RS-sum  $I_1, I_2, I_3, \dots$ : This sequence will be a Cauchy sequence of real numbers, and so will converge to a specific value  $I$  which we consider to be RS-integral of  $f$ : Now the reason why Apostol needed to strengthen the definition is that, otherwise this value  $I$  may not be unique: So if you look at the statement you see in 6.7(b)(c), then they correspond to the Cauchy definition and the 'value-based' definition respectively. For monotonically increasing  $\alpha$ , it is much easier to discuss them using upper and lower sums. So your exercise today will be to read the statements and proofs in Theorem 6.7.

## §8.2 Properties of the Integral

### Theorem 8.2.1.

1. Linearity of  $f$ :

(a) If  $f_1, f_2 \in R_\alpha[a, b]$ , then  $f_1 + f_2 \in R_\alpha[a, b]$  and

$$\int_a^b (f_1 + f_2) d\alpha = \int_a^b f_1 d\alpha + \int_a^b f_2 d\alpha.$$

(b) If  $f \in R_\alpha[a, b]$  and  $c \in \mathbb{R}$ , then  $cf \in R_\alpha[a, b]$  and

$$\int_a^b (cf) d\alpha = c \int_a^b f d\alpha.$$

2. Linearity of  $\alpha$ :

(a) If  $f \in R_{\alpha_1}[a, b] \cap R_{\alpha_2}[a, b]$ , then  $f \in R_{(\alpha_1 + \alpha_2)}[a, b]$  and

$$\int_a^b f d(\alpha_1 + \alpha_2) = \int_a^b f d\alpha_1 + \int_a^b f d\alpha_2.$$

(b) If  $f \in R_\alpha[a, b]$  and  $c \geq 0$ , then  $f \in R_{c\alpha}[a, b]$  and

$$\int_a^b f d(c\alpha) = c \int_a^b f d\alpha.$$

3. If  $f_1, f_2 \in R_\alpha[a, b]$  and  $f_1 \leq f_2$ , then

$$\int_a^b f_1 d\alpha \leq \int_a^b f_2 d\alpha.$$

4. If  $f \in R_\alpha[a, b]$  and  $c \in [a, b]$ , then  $f \in R_\alpha[a, c] \cap R_\alpha[c, b]$  and

$$\int_a^b f d\alpha = \int_a^c f d\alpha + \int_c^b f d\alpha.$$

5. If  $f \in R_\alpha[a, b]$  and  $|f| \leq M$ , then

$$\left| \int_a^b f d\alpha \right| \leq M (\alpha(b) - \alpha(a)).$$

6. If  $f, g \in R_\alpha[a, b]$ , then  $fg \in R_\alpha[a, b]$ .

7. If  $f \in R_\alpha[a, b]$ , then  $|f| \in R_\alpha[a, b]$  and

$$\left| \int_a^b f d\alpha \right| \leq \int_a^b |f| d\alpha.$$

### Proof

- 1.
- 2.
- 3.
- 4.
- 5.
- 6.

7.

□

theorem 6.13

**Theorem 8.2.2** (Triangle inequality).  $f \in R_\alpha[a, b]$ , then  $|f| \in R_\alpha[a, b]$ ,

$$\left| \int_a^b f \, d\alpha \right| \leq \int_a^b |f| \, d\alpha.$$

theorem 6.13

**Theorem 8.2.3.**  $f \in R_\alpha[a, b]$ ,  $\phi$  is uniformly continuous on  $\mathbb{R}$ , then

$$\phi \circ f \in R_\alpha[a, b].$$

refer to book, split into two cases

6.14 6.15 Heaviside step function

6.16 corollary for infinite sum, need  $\sum c_n$  to converge (23) comparison test

6.17 integration by substitution

**Theorem 8.2.4.**  $\alpha$  increasing,  $\alpha' \in R[a, b]$ ,  $f$  bounded on  $[a, b]$ , then

$$f \in R_\alpha[a, b] \iff f\alpha' \in R[a, b].$$

6.19 change of variables

## §8.3 Fundamental Theorem of Calculus

6.20 6.21

**Theorem 8.3.1.**

6.22 integration by parts

# 9 Sequence and Series of Functions

## §9.1 Uniform Convergence

**Definition 9.1.1.** Suppose  $\{f_n\}$ ,  $n = 1, 2, 3, \dots$  is a sequence of functions defined on a set  $E$ , and suppose that the sequence of numbers  $\{f_n(x)\}$  converges for every  $x \in E$ . We can then define a function  $f$  by

$$f(x) = \lim_{n \rightarrow \infty} f_n(x).$$

We say that  $\{f_n\}$  **converges pointwise** to  $f$  on  $E$ , denoted by  $f_n \rightarrow f$ .

Similarly, if  $\sum f_n(x)$  converges for every  $x \in E$ , and if we define

$$f(x) = \sum_{n=1}^{\infty} f_n(x)$$

the function  $f$  is called the **sum of the series**  $\sum f_n$ .

pointwise convergence

**Definition 9.1.2.** Assume  $\{f_n\}$  is a sequence of functions defined over a set  $X$  and  $f$  is also a function defined over  $X$ . We say  $\{f_n\}$  **uniformly converges** to  $f$  over  $X$ , if for any  $\varepsilon > 0$ , there exists  $N > 0$  (which is independent of  $x$ ) so that for any  $x \in X$ ,

$$|f_n(x) - f(x)| < \varepsilon.$$

**Notation** We denote this uniform convergence over  $X$  by  $f_n \rightrightarrows f$ .

## §9.2 Uniform Convergence and Continuity

## §9.3 Uniform Convergence and Integration

**Theorem 9.3.1.** Assume  $\{f_n\}$  is a sequence of functions defined over  $[a, b]$  and each  $f_n \in R_\alpha[a, b]$ . If  $f_n \rightarrow f$ , then  $f \in R_\alpha[a, b]$ , and

$$\lim_{n \rightarrow \infty} \int_a^b f_n d\alpha = \int_a^b f d\alpha.$$

**Proof** Define □

**Corollary 9.3.2.** Assume  $a_n \in R_\alpha[a, b]$  and

$$f(x) := \sum_{n=0}^{\infty} a_n(x)$$

converges uniformly. Then it follows

$$\int_a^b f d\alpha = \sum_{n=0}^{\infty} \int_a^b a_n d\alpha.$$



**Proof** Consider the sequence of partial sums

$$f_n(x) := \sum_{k=0}^n a_k(x), \quad n = 0, 1, \dots$$

It follows  $f_n \in R_\alpha[a, b]$  and  $f_n \Rightarrow f$ . Apply above theorem to  $\{f_n\}$  and the conclusion follows.  $\square$

## §9.4 Uniform Convergence and Differentiation

**Theorem 9.4.1.** Assume  $\{f_n\}$  is a sequence of functions defined over  $[a, b]$  and differentiable. If  $\{f'_n\}$  uniformly converges on  $[a, b]$  and  $\{f_n\}$  converges at some point  $x_0 \in [a, b]$ , then  $\{f_n\}$  uniformly converges on  $[a, b]$  to some function  $f$ . Moreover,  $f$  is differentiable and

$$f'(x) = \lim_{n \rightarrow \infty} f'_n(x)$$

for any  $x \in [a, b]$ .

**Proof**

$\square$

# Part III

## Appendices

# A H3 Mathematics

## §A.1 A Level past year papers

2023

1. (a) Prove that, for any real numbers  $a_1, a_2, \dots, a_n$ ,

$$a_1 + a_2 + \dots + a_n \leq \sqrt{n} \sqrt{a_1^2 + a_2^2 + \dots + a_n^2}.$$

- (b) Prove that, for any positive real numbers  $x, y$  and  $z$ ,

$$\sqrt{\frac{x+y}{x+y+z}} + \sqrt{\frac{y+z}{x+y+z}} + \sqrt{\frac{z+x}{x+y+z}} \leq \sqrt{6}.$$

- (c) Hence solve the equation

$$2\sqrt{\frac{x+3}{x+6}} + \sqrt{\frac{6}{x+6}} = \sqrt{6}.$$

**Solution:**

- (a) Square both sides, apply Cauchy–Schwarz.

- (b) Let

$$a_1 = \sqrt{\frac{x+y}{x+y+z}}, \quad a_2 = \sqrt{\frac{y+z}{x+y+z}}, \quad a_3 = \sqrt{\frac{z+x}{x+y+z}},$$

then apply (a) to the above three real numbers.

- (c) Let  $y = 3, z = 3$ , then apply (b).

Equality in the Cauchy–Schwarz inequality holds if and only if

$$\left( \sqrt{\frac{x+3}{x+6}}, \sqrt{\frac{6}{x+6}}, \sqrt{\frac{x+3}{x+6}} \right) = \lambda (1, 1, 1)$$

for some  $\lambda > 0$ . This happens exactly when  $x = 3$ .

2.

3.

4. Let  $n$  stones be placed in fixed positions on a line. Each stone is painted using one of four colours (red, white, yellow or blue) in such a way that no two adjacent stones are the same colour. Let  $r_n$  be the number of ways of painting the stones such that the first and last stones are both red. Let  $s_n$  be the number of ways of painting the stones so that the first stone is red but the last stone is not red.

- (a) Explain why  $r_1 = 1, r_2 = 0, s_1 = 0$  and  $s_2 = 3$ .

- (b) Find a formula for  $r_n + s_n$  and explain why  $r_{n+1} = s_n$ .  
(c) Using mathematical induction, or otherwise, prove that for all  $n \geq 4$ ,

$$r_n = \frac{3^{n-1} + 3(-1)^{n-1}}{4}.$$

- (d) Now let  $n$  stones, where  $n > 1$ , be placed on a circle with numbered positions. Find the number of ways of painting these stones, using at most four distinct colours, in such a way that no two adjacent stones are the same colour.

**Solution:**

- (a) If  $n = 1$ , then the first stone is also the last stone, and there is thus only 1 way to paint the stone red. So

$$r_1 = 1.$$

Since the first stone is red and the last stone, which is the first stone, is red, there is no way to paint the stone such that the last (first) stone is not red. So

$$s_1 = 0.$$

If  $n = 2$ , then the first and last (second) stones must be painted red but in doing so they will become adjacent stones painted red, which violates the condition that no two adjacent stones are the same colour. Hence

$$r_2 = 0.$$

- (b) We use the following notation:

$R := \{\text{painting arrangements such that first and last stones are red}\};$

$S := \{\text{painting arrangements such that first stone is red and last stone is not red}\};$

$T := \{\text{painting arrangements such that first stone is red}\}.$

By definition of  $r_n$  and  $s_n$ ,

$$|R| = r_n, \quad |S| = s_n.$$

Since  $R$  and  $S$  are mutually exclusive such that  $R \cup S = T$ , we have

$$|T| = |R \cup S| = |R| + |S| = r_n + s_n.$$

On the other hand,

$$T = 1 \times \underbrace{3 \times 3 \times \cdots \times 3}_{n-1} = 3^{n-1}.$$

Thus

$$r_n + s_n = 3^{n-1}.$$

Observe that if  $n + 1$  stones are painted in such a way that the first and last stones are both red, then the  $n$ -th stone must necessarily be non-red since no two adjacent stones can share the same colour. Therefore, the number of ways to paint  $n + 1$  stones such that the first and last stones are both red is equal to the number of ways to paint  $n$  stones such that the first stone is red and the last stone is not red. Thus,  $r_{n+1} = s_n$ .

- (c)  
(d)

## 2022

- 1.
- 2.
- 3.
4. (a) Let  $a$  and  $b$  be positive numbers such that  $a + b = 1$ . Using a sketch graph of  $y = \ln x$ , for  $x > 0$ , show that

$$u^a v^b \leq au + bv$$

for positive  $u$  and  $v$ .

- (b) Let  $a_1, a_2, a_3, \dots$  be a sequence of positive numbers. Define

$$G_n = \sqrt[n]{a_1 a_2 \cdots a_n} \text{ and } A_n = \frac{a_1 + a_2 + \cdots + a_n}{n}.$$

- (i) Use the result of part (a) to prove that the sequence  $(n(G_n - A_n))$  is non-increasing.
- (ii) Let the first 3 terms of the sequence  $(a_n)$  be 1, 2 and 4. Define a suitable  $a_n$ , for  $n \geq 4$ , so that  $(n(G_n - A_n))$  is constant for  $n \geq 3$ .

**Solution:**

- (a) Sketch the graph  $y = \ln x$  for  $x > 0$ , defined on the closed bounded interval  $[u, v]$  for two positive real numbers  $u < v$ .

Since  $a + b = 1$ , the real number  $au + bv$  lies in the interval  $[u, v]$ ; that is,

$$u \leq au + bv \leq v.$$

Note that the equation of the straight line joining  $(u, \ln u)$  and  $(v, \ln v)$  is given by

$$y - \ln u = \frac{\ln v - \ln u}{v - u} \cdot (x - u).$$

When  $x = au + bv$ , the  $y$ -value on this straight line reads off

$$y = a \ln u + b \ln v.$$

Since  $y = \ln x$  is concave, the  $y$ -value read off the straight line is at most the  $y$ -value read off the curve  $y = \ln x$ , and thus it follows that

$$a \ln u + b \ln v \leq \ln(au + bv)$$

or equivalently,

$$u^a v^b \leq au + bv.$$

(b)

- 5.
- 6.
7. (a) The diagram below shows a  $3 \times 3$  array of circles, five of which are shaded. Of the 20 edges linking pairs of adjacent (including diagonally adjacent) circles, 11 link a shaded and an unshaded circle.
  - (i) Describe or draw a  $3 \times 3$  array of circles for which more than 11 edges link a shaded and an unshaded circle and state the number of such edges.

The second diagram shows how the edges for a  $4 \times 4$  array of circles can be grouped into square blocks (consisting of 6 edges) along a diagonal and arrowhead shapes (consisting of 4 edges) elsewhere. For clarity the circles are not shown.

  - (ii) For the edges of an  $n \times n$  array of circles grouped as in the second diagram, state the number of square blocks and arrowhead shapes that would be required.

- (iii) Explain why at most 3 of the edges in an arrowhead shape can link a shaded and an unshaded circle.
- (b) In the  $3 \times 3$  grid below, some of the squares are shaded. The number in each unshaded square shows the number of shaded squares with which the unshaded square shares a vertex. The sum of all the numbers, 12, is the score of this arrangement of shaded and unshaded squares.
  - (i) Explain, why, for any such arrangement, the score is unaltered by shading each unshaded square and vice versa.
  - (ii) Find the maximum possible score for an  $n \times n$  grid and prove that it can be attained.

**2021**

- 1.
2. Let  $a, b, c$  and  $r$  be positive real numbers.

(a) Prove that

$$a^r(a-b)(a-c) + b^r(b-c)(b-a) + c^r(c-a)(c-b) \geq 0.$$

(b) Hence, or otherwise, prove that

$$\begin{aligned} \text{(i)} \quad & a^3 + b^3 + c^3 + 3abc \geq a^2(b+c) + b^2(c+a) + c^2(a+b), \\ \text{(ii)} \quad & \frac{1}{a^5} + \frac{1}{b^5} + \frac{1}{c^5} + \frac{a+b+c}{a^2b^2c^2} \geq \frac{b^2+c^2}{a^3b^2c^2} + \frac{c^2+a^2}{a^2b^3c^2} + \frac{a^2+b^2}{a^2b^2c^3}. \end{aligned}$$

**Solution:**

(a) WLOG assume  $a \geq b \geq c$ . Then the given expression can be rewritten as

$$\underbrace{a^r(a-b)(a-c)}_{(1)} + \underbrace{(b-c)[b^r(b-a) - c^r(c-a)]}_{(2)}.$$

For (1), since  $a \geq b$  and  $a \geq c$ , we have  $a-b \geq 0$  and  $a-c \geq 0$ . Since  $a > 0$ , we then have

$$a^r(a-b)(a-c) \geq 0.$$

Now for (2), since  $b \geq c$  we have

$$b^r(b-a) - c^r(c-a) \geq c^r(b-a-c+a) = c^r(b-c) \geq 0.$$

Thus we have proven the given statement.

(b) Choose  $r = 1$ , by (a), we have

$$a(a-b)(a-c) + b(b-c)(b-a) + c(c-a)(c-b) \geq 0.$$

Expanding LHS gives us the desired statement.

(c) The term  $\frac{a+b+c}{a^2b^2c^2}$  can be seen as  $\frac{1}{ab^2c^2} + \frac{1}{bc^2a^2} + \frac{1}{ca^2b^2}$ . This term can be compared to the term  $3abc$  in the first part. From this observation, we consider Schur's inequality exhibited by

$$\frac{1}{a} \left( \frac{1}{a^2} - \frac{1}{b^2} \right) \left( \frac{1}{a^2} - \frac{1}{c^2} \right) + \frac{1}{b} \left( \frac{1}{b^2} - \frac{1}{c^2} \right) \left( \frac{1}{b^2} - \frac{1}{a^2} \right) + \frac{1}{c} \left( \frac{1}{c^2} - \frac{1}{a^2} \right) \left( \frac{1}{c^2} - \frac{1}{b^2} \right) \geq 0.$$

Expanding gives us

$$\begin{aligned} & \frac{1}{a^5} - \frac{1}{a^3c^2} - \frac{1}{a^3b^2} + \frac{1}{ab^2c^2} \\ & + \frac{1}{b^5} - \frac{1}{b^3a^2} - \frac{1}{b^3c^2} + \frac{1}{bc^2a^2} \\ & + \frac{1}{c^5} - \frac{1}{c^3b^2} - \frac{1}{c^3a^2} + \frac{1}{ca^2b^2} \geq 0. \end{aligned}$$

Thus

$$\begin{aligned} & \left( \frac{1}{a^5} + \frac{1}{b^5} + \frac{1}{c^5} \right) + \left( \frac{1}{ab^2c^2} + \frac{1}{bc^2a^2} + \frac{1}{ca^2b^2} \right) \\ & \geq \left( \frac{1}{a^3c^2} + \frac{1}{a^3b^2} \right) + \left( \frac{1}{b^3a^2} + \frac{1}{b^3c^2} \right) + \left( \frac{1}{c^3b^2} + \frac{1}{c^3a^2} \right). \end{aligned}$$

which gives us the desired inequality.

3. Let  $u$  and  $v$  be quadratic functions of  $x$  and let

$$y = \frac{u}{v}.$$

- (a) Use mathematical induction to prove that

$$v \frac{d^{n+2}y}{dx^{n+2}} + (n+2) \frac{dv}{dx} \frac{d^{n+1}y}{dx^{n+1}} + \binom{n+2}{2} \frac{d^2v}{dx^2} \frac{d^n y}{dx^n} = 0,$$

for  $n \geq 1$ .

- (b) Now assume that  $v = (\alpha - x)^2$  for some real number  $\alpha$  and, for all positive integers  $n$ , define

$$z_n = \frac{(\alpha - x)^{n+2}}{n!} \frac{d^n y}{dx^n}.$$

Use the result of part (a) to prove that  $z_1, z_2, z_3, \dots$  is an arithmetic progression. By writing  $y$  as partial fractions, or otherwise, show that the common difference is  $u(\alpha)$ .



**2020**

1. (i) For any positive integer  $n$  and positive numbers  $x$  and  $y$ , prove that

$$((n-1)x + y)^n \geq n^n x^{n-1} y.$$

- (ii) Hence, for any positive numbers  $a$ ,  $b$  and  $c$  such that  $abc = 1$ , prove that

$$(1+a)^2(1+b)^3(1+c)^4 > 256.$$

**Solution:**

- (i) Apply AM–GM on  $\underbrace{x + \cdots + x}_{n-1} + y$ .

- (ii) Watching out for the various powers of 2, 3 and 4, we rewrite

$$\begin{aligned} & (1+a)^2(1+b)^3(1+c)^4 \\ &= \left((2-1)1+a\right)^2 \cdot \left((3-1)\frac{1}{2}+b\right)^3 \cdot \left((4-1)\frac{1}{3}+c\right)^4 \\ &\geq \left(2^2 \cdot 1^{2-1} \cdot a\right) \cdot \left(3^3 \cdot \left(\frac{1}{2}\right)^{3-1} \cdot b\right) \cdot \left(4^4 \cdot \left(\frac{1}{3}\right)^{4-1} \cdot c\right) \\ &= 256abc = 256 \end{aligned}$$

where equality (for AM–GM) holds if and only if  $a = \frac{1}{2}$ ,  $b = \frac{1}{3}$  and  $c = \frac{1}{4}$ , which is impossible as it contradicts the given condition of  $abc = 1$ . Thus equality never holds, and so the inequality is a strict one.

2.

3. For any non-negative integer  $n$ , the function  $P_n$  is defined by

$$P_n(t) = \sum_{i=0}^n \frac{t^i}{i!}.$$

- (i) Use mathematical induction to prove that

$$\int_0^t x^n e^{-x} dx = n! (1 - e^{-t} P_n(t)).$$

- (ii) State the value of

$$\int_0^\infty x^n e^{-x} dx,$$

and briefly justify your answer.

- (iii) For  $n > t > 0$ , prove that

$$\left(1 + \frac{t}{n}\right)^n \leq P_n(t) < \left(1 - \frac{t}{n}\right)^{-n}.$$

**Solution:**

- (i) Formalise by stating the statement we want to prove:

$$P(n) : \int_0^t x^n e^{-x} dx = n! (1 - e^{-t} P_n(t)), \quad n = 0, 1, \dots$$

When  $n = 0$ , we must prove that

$$\int_0^t x^0 e^{-x} dx = 0! (1 - e^{-t} P_0(t)).$$

The working is direct:

$$\begin{aligned}
 \int_0^t x^0 e^{-x} dx &= \int_0^t e^{-x} dx \\
 &= [-e^{-x}]_0^t \\
 &= -e^{-t} + 1 \\
 &= \underbrace{0!}_{=1} (1 - e^{-1} \underbrace{P_0(t)}_{=1})
 \end{aligned}$$

Assume that  $P(n)$  holds, we want to prove that

$$P(n+1) : \int_0^t x^{n+1} e^{-x} dx = (n+1)! (1 - e^{-t} P_{n+1}(t))$$

holds.

Integrating by parts, let

$$u = x^{n+1}, \quad \frac{dv}{dx} = e^{-x},$$

we have

$$\begin{aligned}
 &\int_0^t x^{n+1} e^{-x} dx \\
 &= [-x^{n+1} e^{-x}]_0^t - \int_0^t (n+1)x^n (-e^{-x}) dx \\
 &= (-t^{n+1} e^{-t}) + (n+1) \int_0^t x^n e^{-x} dx \\
 &= (n+1)! \left( 1 - e^{-t} \left( P_n(t) + \frac{t^{n+1}}{(n+1)!} \right) \right) \\
 &= (n+1)! (1 - e^{-t} P_{n+1}(t))
 \end{aligned}$$

(ii)

$$\int_0^\infty x^n e^{-x} dx = \lim_{t \rightarrow \infty} n! (1 - e^{-t} P_n(t)) = n!$$

since  $\lim_{t \rightarrow \infty} \frac{P_n(t)}{e^t} = 0$ .

(iii) We start by proving the LHS:

$$\left(1 + \frac{t}{n}\right)^n \leq P_n(t).$$

Using binomial expansion,

$$\left(1 + \frac{t}{n}\right)^n = \sum_{k=0}^n \binom{n}{k} \frac{t^k}{n^k}.$$

We want to show that

$$\sum_{k=0}^n \binom{n}{k} \frac{t^k}{n^k} \leq \sum_{k=0}^n \frac{t^k}{k!}.$$

This can be achieved if we are able to prove that

$$\frac{\binom{n}{k}}{n^k} \leq \frac{1}{k!}$$

for  $k = 0, 1, 2, \dots, n$ . And this is best approached by working backwards: since for all  $j = 0, 1, 2, \dots, n$ , it holds that

$$\frac{n - (j-1)}{n} \leq 1,$$

we must have that

$$\frac{n \times (n-1) \times \dots \times (n - (k-1))}{n \times n \times \dots \times n} \leq 1.$$

Thus,

$$\frac{n!}{(n-k)!k!} \frac{1}{n^k} \leq \frac{1}{k!}$$

and we have proved that

$$\frac{\binom{n}{k}}{n^k} \leq \frac{1}{k!}$$

for  $k = 0, 1, 2, \dots, n$ , as planned. This then implies that

$$\left(1 + \frac{t}{n}\right)^n = \sum_{k=0}^n \binom{n}{k} \frac{t^k}{n^k} \leq \sum_{k=0}^n \frac{t^k}{k!}.$$

We now prove the RHS:

$$P_n(t) < \left(1 - \frac{t}{n}\right)^{-n}.$$

Note that  $n > t > 0$ , which implies that  $-n$  is a negative integer and  $|-t/n| < 1$ , flagging out the warrant for us to apply Newton's binomial expansion:

$$\begin{aligned} \left(1 - \frac{t}{n}\right)^{-n} &= 1 + (-n) \left(-\frac{t}{n}\right) + \frac{(-n)(-n-1)}{2!} \left(-\frac{t}{n}\right)^2 + \frac{(-n)(-n-1)(-n-2)}{3!} \left(-\frac{t}{n}\right)^3 + \dots \\ &\quad + \frac{(-n)(-n-1)(-n-2)\dots(-n-(k-1))}{k!} \left(-\frac{t}{n}\right)^k + \dots \\ &= 1 + t + \frac{n(n+1)}{2!} \left(\frac{t}{n}\right)^2 + \frac{n(n+1)(n+2)}{3!} \left(\frac{t}{n}\right)^3 + \dots \\ &\quad + \frac{n(n+1)(n+2)\dots(n+(k-1))}{k!} \left(\frac{t}{n}\right)^k + \dots \\ &= 1 + t + \frac{1\left(1 + \frac{1}{n}\right)}{2!} t^2 + \frac{1\left(1 + \frac{1}{n}\right)\left(1 + \frac{2}{n}\right)}{3!} t^3 + \dots + \frac{1\left(1 + \frac{1}{n}\right)\left(1 + \frac{2}{n}\right)\dots\left(1 + \frac{k-1}{n}\right)}{k!} t^k + \dots \\ &> 1 + t + \frac{t^2}{2!} + \frac{t^3}{3!} + \dots + \frac{t^n}{n!} \\ &= P_n(t) \end{aligned}$$

**2019**

- 1.
- 2.
3. A sequence is defined by

$$x_1 = 1 \text{ and } x_{i+1} = \left(\frac{i+a}{i+1}\right)x_i, \quad i \geq 1.$$

- (i) Assume that  $a \geq 0$ .
  - (a) Prove that  $x_i \geq \frac{1}{i}$ , for all positive integers  $i$ .
  - (b) Prove that

$$\sum_{i=n+1}^{2n} x_i \geq \frac{1}{2},$$

for all positive integers  $n$ .

- (c) Hence prove that  $\sum_{i=1}^{\infty} x_i$  is unbounded.
- (ii) Assume that  $a < 0$ .
  - (a) Prove that

$$a \sum_{i=m}^n x_i = (n+1)x_{n+1} - mx_m$$

for all positive integers  $m$  and  $n$  such that  $n > m$ .

- (b) For any sufficiently large integers  $m$  and  $n$ , prove that  $x_m x_n \geq 0$ .

**Solution:**

- (i) (a) We proceed to prove the statement

$$P(n) : x_n \geq \frac{1}{n}, \quad n = 1, 2, 3, \dots$$

$P(1)$  is true since  $x_1 = 1 \geq \frac{1}{1}$ .

Assume that  $P(k)$  holds for some  $k \in \mathbb{Z}^+$ ; that is,

$$x_k \geq \frac{1}{k}.$$

We want to prove  $P(k+1)$  holds; that is,

$$x_{k+1} \geq \frac{1}{k+1}.$$

Since  $x_{k+1} = \left(\frac{k+a}{k+1}\right)x_k$ , by the inductive hypothesis  $P(k)$  we deduce that

$$x_{k+1} = \left(\frac{k+a}{k+1}\right)x_k \geq \left(\frac{k+a}{k+1}\right) \cdot \frac{1}{k} = \underbrace{\left(1 + \frac{a}{k}\right)}_{\geq 1} \cdot \frac{1}{k+1} \geq \frac{1}{k+1}.$$

Since  $P(1)$  holds and  $P(k) \implies P(k+1)$  for all  $k \in \mathbb{Z}^+$ , by mathematical induction,  $P(n)$  holds for all  $n \in \mathbb{Z}^+$ .

- (b)

$$\begin{aligned} \sum_{i=n+1}^{2n} x_i &= x_{n+1} + x_{n+2} + \dots + x_{2n} \\ &\geq \frac{1}{n+1} + \frac{1}{n+2} + \dots + \frac{1}{2n} \\ &\geq \frac{1}{2n} + \frac{1}{2n} + \dots + \frac{1}{2n} \\ &= \frac{1}{2n} \times n = \frac{1}{2} \end{aligned}$$

(c) Informally, we see that

$$\begin{aligned}\sum_{i=1}^{\infty} &= x_1 + x_2 + (x_3 + x_4) + (x_5 + x_6 + x_7 + x_8) + \cdots \\ &\geq 1 + \frac{1}{2} + \frac{1}{2} + \frac{1}{2} + \cdots\end{aligned}$$

which provides compelling evidence that  $\sum_{i=1}^{\infty} x_i$  is unbounded.

We not write this argument in a rigorous manner. To show that  $\sum_{i=1}^{\infty} x_i$  is unbounded, one must show that for any  $M > 0$ , there exists  $N \in \mathbb{Z}^+$  such that

$$\sum_{i=1}^N x_i > M.$$

Indeed, given any  $M > 0$ , there exists  $k \in \mathbb{Z}^+$  so large that

$$k > 2(M - 1) \iff 1 + k \cdot \frac{1}{2} > M.$$

Thus it follows that

$$\begin{aligned}\sum_{i=1}^{2^k} x_i &= x_1 + x_2 + (x_3 + x_4) + (x_5 + x_6 + x_7 + x_8) + \cdots + (x_{2^{k-1}+1} + \cdots + x_{2^k}) \\ &> 1 + \underbrace{\frac{1}{2} + \frac{1}{2} + \cdots + \frac{1}{2}}_{k \text{ terms}} \\ &= 1 + k \cdot \frac{1}{2} > M.\end{aligned}$$

(ii) (a) By definition,

$$x_{i+1} = \left( \frac{i+a}{i+1} \right) x_i, \quad i \geq 1.$$

So

$$(i+1)x_{i+1} = (i_1)x_i \iff (i_1)x_{i+1} - ix_i = ax_i.$$

Thus, if  $n > m$  we have

$$\begin{aligned}\sum_{i=m}^n ax_i &= \sum_{i=m}^n [(i+1)x_{i+1} - ix_i] \\ &= (m+1)x_{m+1} - mx_m \\ &\quad + (m+2)x_{m+2} - (m+1)x_{m+1} \\ &\quad + \vdots \\ &\quad + (n+1)x_{n+1} - nx_n \\ &= (n+1)x_{n+1} - mx_m\end{aligned}$$

by the method of difference.

(b) For We want to prove that there exists a large enough  $N \in \mathbb{N}$  such that if  $m, n \geq N$  then  $x_m x_n \geq 0$ . Notice that in the situation when neither of  $x_m$  or  $x_n$  is zero these two numbers will have the same sign, i.e. either they are both negative or both positive.

By the recursive definition of  $x_i$ 's, if  $n > m$  then

$$x_n = \frac{n-1+a}{n} \cdot \frac{n-2+a}{n-1} \cdots \frac{m+a}{m+1} x_m.$$

Since  $a < 0$  is fixed, there exists a sufficiently large positive integer  $N$  such that  $N - a > 0$ . Consequently, if  $n > m \geq N$ , we have

$$x_n = \underbrace{\frac{n-1+a}{n}}_{>0} \cdot \underbrace{\frac{n-2+a}{n-1}}_{>0} \cdots \underbrace{\frac{m+a}{m+1}}_{>0} x_m.$$

Hence  $x_n$  and  $x_m$  are of the same sign.

**Remark** The question whether one can make use of (a) to solve (b) remains open.

4. An  $n$ -digit number uses no digits other than 1, 2 and 3. It does not have any 2s adjacent to each other, and it does not have any 3s adjacent to each other. Let there be  $T_n$  such numbers, with  $X_n$  of these having first digit 1 and  $Y_n$  having first digit 2.

(a) Prove that, for any  $n \geq 2$ ,

- (i)  $Y_n = X_{n-1} + Y_{n-1}$ ,
- (ii)  $X_n = X_{n-1} + 2Y_{n-1}$ ,
- (iii)  $X_{n+1} = 2X_n + X_{n-1}$ .

(b) Use mathematical induction to prove that, for  $n \geq 1$ ,

$$X_n \equiv n^2 - n + 1 \pmod{4}.$$

(c) Find and simplify an expression for  $T_n \pmod{4}$ .

5. (i) Use the substitution  $t = \frac{du}{dx}$  to find the general solution of the equation

$$\frac{d^2u}{dx^2} = \frac{du}{dx}.$$

(ii) Show that the differential equation can be transformed into the equation

$$f(x) \frac{d^2u}{dx^2} - (f'(x) + f(x)g(x)) \frac{du}{dx} = 0$$

by the substitution

$$u = e^{-\int f(x)g(x)dx}.$$

(iii) A solution curve of the differential equation

$$\frac{dy}{dx} = e^{-2x}y^2 + 3y$$

passes through the point  $(0, -\frac{1}{4})$ . Find the equation of the curve.

## 2018

1. A triangle has sides of lengths  $a$ ,  $b$  and  $c$  units. In each of the following cases, prove that there is a triangle having sides of the given lengths.

- (i)  $\frac{a}{1+a}$ ,  $\frac{b}{1+b}$  and  $\frac{c}{1+c}$  units.
- (ii)  $\sqrt{a}$ ,  $\sqrt{b}$  and  $\sqrt{c}$  units.
- (iii)  $\sqrt{a(b+c-a)}$ ,  $\sqrt{b(c+a-b)}$  and  $\sqrt{c(a+b-c)}$  units.

2.

3.

4. A clothes shop sells a particular make of T-shirt in four different colours. The shopkeeper has a large number of T-shirts of each colour.

- (i) A customer wishes to buy seven T-shirts.
  - (a) In how many ways can he do this?
  - (b) In how many ways can he do this if he buys at least one of each colour.
- (ii) The shopkeeper places seven T-shirts in a line.
  - (a) In how many ways can she do this?
  - (b) In how many ways can she do this if no two T-shirts of the same colour are to be next to each other?
  - (c) Use the principle of inclusion and exclusion to find the number of ways in which she can do this if she has to use at least one T-shirt of each colour but with no other restrictions.

5. A  $p \times q$  chessboard can be tessellated with  $a \times b$  tiles.

A unit square  $(x, y)$  is shaded if and only if  $x \equiv y \pmod{a}$ .

- (i) Explain why the following are necessary conditions for such a tessellation
  - (a)  $ab$  is a factor of  $pq$ .
  - (b)  $p$  and  $q$  can be written in the form  $ma + nb$  where  $m$  and  $n$  are non-negative integers.
  - (c) The  $p \times q$  chessboard has  $\frac{pq}{a}$  shaded squares.
- (ii) Let  $t$  be the smaller of  $r$  and  $s$  such that

$$\begin{aligned} p &\equiv r \pmod{a} & 0 \leq r < a \\ q &\equiv s \pmod{a} & 0 \leq s < a \end{aligned}$$

- (a) Explain why the number of shaded squares in the  $p \times q$  chessboard is  $\frac{pq - rs}{a} + t$ .
- (b) Hence prove that for a tessellation, either  $a \mid p$  or  $a \mid q$ .

**Solution:**

- (i) (a) A  $p \times q$  chessboard has  $pq$  squares, a  $a \times b$  tile has  $ab$  squares.  
Suppose  $k$  tiles are used to tessellate the board. Then  $pq = kab$ . Hence  $ab \mid pq$ .
- (b)  $p$  and  $q$  are the height and base of the  $p \times q$  chessboard respectively,  $a$  and  $b$  are the height and base of each  $a \times b$  tile respectively. Each tile can be placed horizontally or vertically in the tessellation.  
If we tessellate the board at the bottom from left to right with  $m$  vertical and  $n$  horizontal tiles, there will be  $ma + nb$  squares at the bottom row of the board. Each row of the board is made up of  $q$  squares. So we get  $q = ma + nb$ .  
Similarly, if we tessellate the board on the left from bottom to top, we will get  $p = sa + tb$  (with  $s$  horizontal and  $t$  vertical tiles).
- (ii) (a)

6. (Dirichlet's approximation theorem) Let  $x$  be any positive real numbers and  $n$  be any positive integer. Prove that there are integers  $a$  and  $b$  with  $1 \leq b \leq n$ , such that

$$\left| x - \frac{a}{b} \right| < \frac{1}{bn}.$$

**Solution:**

For any real number  $y$ , we write  $y = [y] + \{y\}$ , where  $[y]$  denotes the integer part of  $y$  and  $\{y\}$  denotes the fractional part of  $y$ ,  $0 \leq \{y\} < 1$ .

We divide the interval  $[0, 1)$  into  $n$  smaller intervals of measure  $\frac{1}{n}$ . Consider  $\{x\}, \{2x\}, \dots, \{nx\}$ . Let  $I_i$  denote the interval  $[\frac{i-1}{n}, \frac{i}{n}]$ , where  $1 \leq i \leq n$ .

We now consider two cases:

**Case 1:** Some  $\{kx\}$  falls in  $I_1$

Then  $kx - [kx] = \{kx\} < \frac{1}{n}$ .

Dividing both sides by  $k$ ,

$$\left| x - \frac{[kx]}{k} \right| < \frac{1}{kn}.$$

By taking  $a = [kx]$  and  $b = k$ , we have the inequality.

**Case 2:** None of  $\{kx\}$  falls in  $I_1$

This means all  $\{kx\}$  fall into  $I_2, I_3, \dots, I_n$ . By Pigeonhole Principle, at least two  $\{kx\}$  fall in the same  $I_i$ .

Let  $\frac{i-1}{n} \leq \{px\} < \frac{i}{n}$  and  $\frac{i-1}{n} \leq \{qx\} < \frac{i}{n}$ . Then

$$\begin{aligned} |\{px\} - \{qx\}| &< \frac{1}{n} \\ |(px - [px]) - (qx - [qx])| &< \frac{1}{n} \\ |(px - qx) - ([px] - [qx])| &< \frac{1}{n} \\ |(p - q)x - ([px] - [qx])| &< \frac{1}{n} \end{aligned}$$

Dividing both sides by  $p - q$ ,

$$\left| x - \frac{([px] - [qx])}{p - q} \right| < \frac{1}{(p - q)n}.$$

WLOG assume  $p > q$ . Then  $1 \leq p - q < n$ . By taking  $a = [px] - [qx]$  and  $b = p - q$ , we have the inequality.

7. The differential equation

$$y \frac{dy}{dx} = x \left( \frac{dy}{dx} \right)^2 + 1, \quad \text{for } x > 0 \tag{1}$$

has a solution curve  $S$  such that  $\frac{d^2y}{dx^2}$  is non-zero for all points of  $S$ .

- (i) By substituting  $t = \frac{dy}{dx}$  into equation (1) and differentiating with respect to  $x$ , show that  $S$  has equation  $y^2 = 4x$ .
  - (ii) Show that a straight line is tangent to the curve  $S$  if and only if it is itself a solution of the equation.
8. For any positive real number  $x$ ,  $n(x)$  is defined as the nearest integer to  $x$ , with halves rounded up.
- For example,  $n(3.5) = 4$ , and  $n(\pi) = 3$ .

- (a) Show that  $\sum_{r=1}^3 n\left(\frac{11}{7}r\right) = 10$ .

The diagram shows the line  $y = \frac{7}{11}x + \frac{1}{2}$  and the integer  $(x, y)$  such that  $1 \leq x \leq 5$ ,  $1 \leq y \leq 3$ .



- (b) Find  $\sum_{r=1}^5 n\left(\frac{7}{11}r\right)$  and explain the connection between your answer and the points underneath the line  $y = \frac{7}{11}x + \frac{1}{2}$ .
- (c) The line  $y = \frac{7}{11}x + \frac{1}{2}$  is rotated through  $180^\circ$  about  $(3, 2)$ . Find the equation of the new line in the form  $x = my + c$  and hence comment on the connection between

$$\sum_{r=1}^3 n\left(\frac{11}{7}r\right) = \sum_{r=1}^5 n\left(\frac{7}{11}r\right).$$

- (d) Let  $p$  and  $q$  be odd integers greater than 1 and consider the integer points  $(x, y)$  such that  $1 \leq x \leq \frac{p-1}{2}$ ,  $1 \leq y \leq \frac{q-1}{2}$ . Let  $N$  be the number of points which lie in between the lines  $y = \frac{q}{p}x + \frac{1}{2}$  and  $x = \frac{p}{q}y + \frac{1}{2}$ .

Explain why  $N + \left(\frac{p-1}{2}\right)\left(\frac{q-1}{2}\right) \equiv 0 \pmod{2}$ .

**2017**

- 1.
2. (i) Let  $y$  be a differentiable function of  $x$ . For any positive integer  $n$ , prove that

$$\frac{d^n}{dx^n}(xy) = x \frac{d^n y}{dx^n} + n \frac{d^{n-1} y}{dx^{n-1}}.$$

- (ii) For any non-negative integer  $n$ , define

$$y_n = e^{x^2} \frac{d^n}{dx^n} (e^{-x^2}).$$

- (a) Find  $y_0$ ,  $y_1$  and  $y_2$ .
- (b) Prove that  $y_{n+2} + 2xy_{n+1} + 2(n+1)y_n = 0$ , for  $n \geq 0$ .
- (c) Hence prove that  $\frac{d}{dx}(y_{n+1}) = -2(n+1)y_n$ , for  $n \geq 0$ .
3. (a) Consider integer solutions of the equation

$$1591x + 3913y = 9331.$$

Show that there is no solution with  $x$  prime.

- (b) Let  $a$ ,  $b$ ,  $r$  and  $s$  be integers such that

$$ra + sb = 1.$$

- (i) Prove that, if  $a$  and  $b$  are both factors of an integer  $n$ , then  $ab$  is a factor of  $n$ .
- (ii) Given that any integers  $u$  and  $v$ , prove by construction that there is an integer  $x$  such that both

$$x \equiv u \pmod{a} \quad \text{and} \quad x \equiv v \pmod{b}.$$

**Solution:**

- (a) First we find  $\gcd(1591, 3913)$  using the Euclidean Algorithm.

$$3913 = 2 \times 1591 + 731$$

$$1591 = 2 \times 731 + 129$$

$$731 = 5 \times 129 + 86$$

$$129 = 1 \times 86 + 43$$

$$86 = 2 \times 43 + 0$$

Thus  $\gcd(1591, 3913) = 43$ . By Bezout's Lemma, there are integer solutions for  $1591x + 3913y = 43$ . Since  $43 \mid 9331$ , multiplying both sides by some constant, there are also integer solutions for  $1591x + 3913y = 9331$ .

To prove by contradiction, we assume that  $x$  is prime, and there exists some integer  $y$  such that  $1591x + 3913y = 9331$ . Dividing both sides by 43,

$$37x + 91y = 217. \tag{*}$$

Observe that  $7 \mid 91y$  and  $7 \mid 217$ , so  $7 \mid 37x$ .

Since  $\gcd(7, 37) = 1$  so  $7 \mid x$ . By our assumption,  $x$  is a prime so  $x = 7$ .

Substituting  $x = 7$  into (\*), we get  $y = -\frac{6}{13}$ , which contradicts  $y$  being an integer.

Hence we conclude that  $x$  cannot be a prime.

- (b) (i) If  $a$  and  $b$  are both factors of  $n$ , then we have  $n = pa$  and  $n = qb$  for some integers  $p$  and  $q$ . Given  $ra + sb = 1$ , we have

$$rna + snb = n$$

$$r(qb)a + s(pa)b = n$$

$$(rq + sp)ab = n$$

and hence  $ab$  is a factor of  $n$ .

- (ii) Prove by construction.

Given that  $ra + sb = 1$ . Multiplying both sides by  $v - u$  gives

$$\begin{aligned} ra(v - u) + sb(v - u) &= v - u \\ ra(v - u) + u &= sb(u - v) + v \end{aligned}$$

We define  $x = ra(v - u) + u = sb(u - v) + v$ . Then  $x \equiv u \pmod{a}$  and  $x \equiv v \pmod{b}$ .

**Remark** The above proof shows the *existence* of solution by a construction.

4. Let  $I_n = \int_0^{\frac{\pi}{4}} \tan^n x \, dx$ .

(i) For  $n > 1$ , prove that  $I_n + I_{n+2} = \frac{1}{n-1}$ .

(ii) Justify the statement that  $\tan x \leq \frac{4}{\pi}x$  on  $\left[0, \frac{\pi}{4}\right]$ .

(iii) Hence, prove that  $I_n$  tends to zero as  $n$  tends to infinity.

(iv) Find the sum of the infinite series

$$1 - \frac{1}{3} + \frac{1}{5} - \frac{1}{7} + \cdots$$

**Solution:**

(i)

$$\begin{aligned} I_n + I_{n-2} &= \int_0^{\frac{\pi}{4}} (\tan^n x + \tan^{n-2} x) \, dx \\ &= \int_0^{\frac{\pi}{4}} \tan^{n-2} x (\tan^2 x + 1) \, dx \\ &= \int_0^{\frac{\pi}{4}} \tan^{n-2} x \cdot \sec^2 x \, dx \\ &= \left[ \frac{\tan^{n-1} x}{n-1} \right]_0^{\frac{\pi}{4}} = \frac{1}{n-1}. \end{aligned}$$

(ii) Sketch the graphs of  $y = \tan x$  and  $y = \frac{4}{\pi}x$  over the interval  $\left[0, \frac{\pi}{4}\right]$ .

Since  $y = \tan x$  is convex over  $\left[0, \frac{\pi}{4}\right]$ , it follows that

$$\tan x \leq \frac{4}{\pi}x$$

for all  $x \in \left[0, \frac{\pi}{4}\right]$ .

(iii)

(iv)

5. (i) Explain why the number of ways to distribute  $r$  distinct objects, where  $r \geq 2$ , into 2 distinct boxes such that neither is empty is  $2^r - 2$ .

(ii) Let  $S(r, n)$  denote the number of ways to distribute  $r$  objects into  $n$  identical boxes such that no box is empty.

(a) Explain why, for  $r \geq 3$ ,

$$S(r, 3) = 2^{r-2} - 1 + 3S(r-1, 3).$$

(b) Prove that, for  $r \geq 3$ ,

$$S(r, 3) = \begin{cases} 0 \pmod{6} & \text{if } r \text{ is even,} \\ 1 \pmod{6} & \text{if } r \text{ is odd.} \end{cases}$$

## Specimen

- 1.
- 2.
- 3.
- 4.
- 5.
- 6.
7. The figures below show, respectively, a square board of 4 unit squares with one unit square covered, and a triomino consisting of 3 unit squares.  
Irrespective of which unit square is covered, a triomino can cover the remaining 3 unit squares of the square board as shown.  
Consider a square board made up of  $4^n$  squares, where  $n \geq 1$ , with one of the unit squares covered. An example of such a square, with  $n = 3$ , is shown below.
  - (i) Explain how, irrespective of unit square is covered, a triomino can be placed on the board in such a way that each quarter of the board now has one unit square covered.
  - (ii) Use mathematical induction to prove that, irrespective of which unit square is initially covered, the remaining squares can be covered by triominoes. State the number of triominoes required.

**Problem 31** (H3M 2021). Let  $Q = \{1, 2, \dots, p-1\}$  for some prime  $p$ , and let there be  $N$  integers in  $Q$  whose cubes are congruent to 1 modulo  $p$ .

- (a) Use the pigeonhole principle to prove that for each integer  $x \in Q$  there is precisely one integer  $y \in Q$  such that  $xy \equiv 1 \pmod{p}$ .
- (b) Explain why the number of choices of integers  $x, y, z \in Q$  such that  $xyz \equiv 1 \pmod{p}$  is  $(p-1)^2$ .
- (c) Use the principle of inclusion and exclusion to prove that the number of choices of three different integers  $x, y, z \in Q$  such that  $xyz \equiv 1 \pmod{p}$  is  $(p-1)(p-4) + 2N$ .
- (d) Hence prove that  $N \equiv (p-1)^2 \pmod{3}$ .
- (e) Given that  $p \equiv 1 \pmod{3}$ , prove that there is an integer  $x \in Q$  such that  $x^2 + x + 1 \equiv 0 \pmod{p}$ .

## Proof

- (a) Note that, by Quotient Remainder Theorem, every integer not divisible by  $p$  is congruent to an integer in  $Q$  modulo  $p$ , and no two integers in  $Q$  are congruent to each other modulo  $p$ .

We have two parts to prove: existence and uniqueness of inverse modulo  $p$

**Existence:** prove by contradiction

Suppose there is an  $x \in Q$  such that for all  $y \in Q$ ,  $xy \not\equiv 1 \pmod{p}$ .

There are  $p-1$  possible  $y \in Q$ , but there are less than  $p-1$  possible  $xy \in Q$  (since  $xy \equiv 1 \pmod{p}$  is excluded).

By Pigeonhole Principle, there are two different  $y_1, y_2 \in Q$  such that  $xy_1 \equiv xy_2 \pmod{p}$  ( $\not\equiv 1$ ). Then

$$p \mid xy_1 - xy_2 \implies p \mid x(y_1 - y_2) \implies p \mid y_1 - y_2 \implies y_1 \equiv y_2 \pmod{p} \implies y_1 = y_2$$

which is a contradiction. Hence every  $x \in Q$  has a  $y \in Q$  such that  $xy \equiv 1 \pmod{p}$ .

**Uniqueness:** prove by contradiction

Suppose there are two different  $y_1, y_2 \in Q$  such that  $xy_1 \equiv xy_2 \pmod{p}$  ( $\equiv 1$ ).

The rest is similar to the above, and thus left as an exercise to the reader.

(b) Use combinatorics.

There are  $p - 1$  ways each to choose  $x$  and  $y$ .

By (a), there is only 1 way to choose  $z \in Q$ , the modular inverse of  $xy \bmod p$ , such that  $(xy)z \equiv 1 \pmod{p}$ .

Hence there is a total number of  $(p - 1)^2$  choices of  $x, y, z$  such that  $xyz \equiv 1 \pmod{p}$ .

(c) Let  $U$  contain all  $(x, y, z)$  such that  $xyz \equiv 1 \pmod{p}$ ,  $A$  is a subset of  $U$  such that  $x \equiv y \pmod{p}$ ,  $B$  is a subset of  $U$  such that  $x \equiv z \pmod{p}$ ,  $C$  is a subset of  $U$  such that  $y \equiv z \pmod{p}$ .

Note that  $A \cap B = A \cap C = B \cap C = A \cap B \cap C$  are all subsets of  $U$  such that  $x \equiv y \equiv z \pmod{p}$ , i.e. this subset of  $U$  contains all  $(x, x, x)$  such that  $x^3 \equiv 1 \pmod{p}$ .

We have  $|U| = (p - 1)^2$  from (b),  $|A| = |B| = |C| = p - 1$ , and  $|A \cap B \cap C| = N$ .

By principle of inclusion and exclusion,

$$|A \cup B \cup C| = 3(p - 1) - 2N.$$

To find the complement of  $A \cup B \cup C$ ,

$$|U - (A \cup B \cup C)| = (p - 1)^2 - (3(p - 1) - 2N) = (p - 1)(p - 4) + 2N.$$

(d) From (c), the number of choices of three different  $x, y, z \in Q$  such that  $xyz \equiv 1 \pmod{p}$  is  $(p - 1)(p - 4) + 2N$ .

Since the number of combinations of such  $x, y, z$  are symmetrical, this number is divisible by 3. That is,

$$\begin{aligned} (p - 1)(p - 4) + 2N &\equiv 0 \pmod{3} \\ (p - 1)(p - 1) - N &\equiv 0 \pmod{3} \\ (p - 1)^2 &\equiv N \pmod{3} \end{aligned}$$

(e) From (d),  $N \equiv (p - 1)^2 \equiv 0 \pmod{3} \implies 3 \mid N \implies N \geq 3$ .

There are at least 3 different  $x$  such that  $x^3 \equiv 1 \pmod{p}$ . Choose such an  $x \in Q$  such that  $x \neq 1$ .

$$x^3 - 1 \equiv 0 \pmod{p}$$

Factorising this gives

$$(x - 1)(x^2 + x + 1) \equiv 0 \pmod{p}$$

Hence

$$p \mid (x - 1)(x^2 + x + 1)$$

Since  $p \nmid x - 1$  as  $x \in \{1, 2, \dots, p - 1\}$ ,

$$p \mid x^2 + x + 1$$

thus  $x^2 + x + 1 \equiv 0 \pmod{p}$

□

**Problem 32** (H3M Specimen Q3, Fermat's Little Theorem).

- (i) Let  $p$  be an odd prime and let  $a$  be an integer not divisible by  $p$ .
- (a) Let  $T$  be the set of remainders for  $a, 2a, \dots, (p-1)a$ , when divided by  $p$ . Show that  $T = \{1, 2, \dots, p-1\}$ .
- (b) Hence prove that  $a^{p-1} \equiv 1 \pmod{p}$ .
- (ii) Let  $x$  and  $y$  be two integers such that  $x^5 + y^5$  is divisible by 5. Prove that  $x^5 + y^5$  is divisible by 25.

**Proof**

- (i) (a) Let  $S = \{1, 2, 3, \dots, p-1\}$ , the set of all non-zero positive remainders obtained when integers are divided by  $p$ .

**Known fact:**  $p \nmid k$  for all  $k \in S = \{1, 2, 3, \dots, p-1\}$ .

Given that  $T$  is the set of remainders when  $a, 2a, 3a, \dots, (p-1)a$  are divided by  $p$ .

Clearly,  $T \subseteq S \cup \{0\}$ .

**Claim 1:**  $0 \notin T$ . **Proof** Prove by contradiction.

Suppose  $0 \in T$ . Then  $p \mid ka$  for some  $k \in S = \{1, 2, 3, \dots, p-1\}$ . Since  $p$  is prime and  $p \nmid a$ , we apply Euclid's Lemma to conclude that  $p \mid k$ , which contracts Fact 1.]

Thus  $T \subseteq S$ . □

**Claim 2:**  $T = S$  itself. **Proof** Prove by contradiction.

Suppose, on the contrary, that  $T \neq S$ .

Then,  $T \subset S$  (i.e.  $T$  is a proper subset of  $S$ ).

Since the sets are finite sets,  $n(T) < n(S) = p-1$ . By the Pigeonhole Principle, there are (at least) two distinct  $ia$  and  $ja$  (from the list of  $p-1$  terms:  $a, 2a, 3a, \dots, (p-1)a$  – the “pigeons”), where  $1 \leq i \neq j \leq p-1$  that share the same remainder when divided  $p$ . The “holes” are the elements in  $T$ ; here we get less holes:  $n(T) < p-1$  based on our (wrong) assumption.

$$\begin{aligned} ia &\equiv ja \pmod{p} \\ ia - ja &\equiv 0 \pmod{p} \\ (i - j)a &\equiv 0 \pmod{p} \end{aligned}$$

We can cancel  $a$  on both sides due to Euclid's lemma. Hence  $i \equiv j \pmod{p}$ .

Since both  $i$  and  $j$  belong to  $S$ , having them share the same remainder when divided by  $p$  means that they are actually the same. Thus  $i = j$ . This contradicts our initial choice of distinct  $ia$  and  $ja$ .

Hence  $T = S = \{1, 2, 3, \dots, p-1\}$ . □

- (b) Let

$$\begin{aligned} a \cdot 1 &\equiv r_1 \pmod{p} \\ a \cdot 2 &\equiv r_2 \pmod{p} \\ a \cdot 3 &\equiv r_3 \pmod{p} \\ &\vdots \\ a \cdot (p-1) &\equiv r_{p-1} \pmod{p} \end{aligned}$$

where  $r_1, r_2, r_3, \dots, r_{p-1}$  are distinct elements of  $T = S = \{1, 2, 3, \dots, p-1\}$ .

So multiplying the LHS and RHS respectively of these congruence equations,

$$a^{p-1}(p-1)! \equiv r_1 r_2 r_3 \dots r_{p-1} \pmod{p}$$

Since  $r_1, r_2, r_3, \dots, r_{p-1}$  is just a rearrangement of  $1, 2, 3, \dots, p-1$ ,

$$a^{p-1}(p-1)! \equiv 1 \cdot 2 \cdot 3 \dots (p-1) \pmod{p}$$

or

$$a^{p-1}(p-1)! \equiv (p-1)! \pmod{p}$$

But  $p \nmid (p-1)!$  so by Euclid's lemma,

$$a^{p-1} \equiv 1 \pmod{p}$$

as desired.

(ii) Prove by cases

Given 5 divides  $x^5 + y^5$ .

**Case 1:** either  $x$  or  $y$  is divisible by 5

WLOG, assume  $5 \mid x$ . Then  $x = 5k$  for some integer  $k$ .

Then  $x^5 = (5k)^5 = 5^2(5^3k^5) = 25t$  so  $25 \mid x^5$ .

Since we can write  $y^5 = (x^5 + y^5) - x^5$ ,  $5 \mid y^5$  so  $5 \mid y$ . We can then similarly show that  $25 \mid y^5$ .

Hence  $25 \mid x^5 + y^5$ .

**Case 2:** both  $x$  and  $y$  are not divisible by 5

Since 5 is a prime, by Fermat's Little Theorem,  $x^5 \equiv x \pmod{5}$  and  $y^5 \equiv y \pmod{5}$ , so  $x^5 + y^5 \equiv x + y \pmod{5}$ .

Since  $5 \mid x^5 + y^5$ , we have also  $5 \mid x + y$ , i.e.  $x + y = 5k$  for some integer  $k$ . We rewrite  $y = 5k - x$ .

Then by binomial expansion,

$$y^5 = (5k - x)^5 = \sum_{i=0}^5 \binom{5}{i} (5k)^{5-i} (-x)^i$$

which gives  $y^5 \equiv (-x)^5 \pmod{25}$  as all the other terms are divisible by 25.

Hence  $x^5 + y^5 \equiv 0 \pmod{25}$ .

□

**Problem 33** (H3M Specimen). For any positive integer  $n$ , if one square is removed from a  $2^n \times 2^n$  checkerboard, the remaining squares can be completely covered by triominoes (an L-shaped domino consisting of three squares).

**Proof** Prove by induction.

**Base case:**  $P(1)$  is clearly true.

**Inductive step:**  $P(k) \implies P(k+1)$  is true for all  $k$ , i.e. if a  $2^k \times 2^k$  checkerboard with a square removed can be completely covered by triominoes, then a  $2^{k+1} \times 2^{k+1}$  checkerboard with a square removed can be completely covered by triominoes.

- (i) Divide the  $2^{k+1} \times 2^{k+1}$  checkerboard into four  $2^k \times 2^k$  sub-boards.
- (ii) One of the sub-boards include the removed square.
- (iii) WLOG, assume the top left sub-board has the removed square.
- (iv) By induction hypothesis, this sub-board can be covered by triominoes.
- (v) For the top right sub-board, we cover it with triominoes with a remaining square at the bottom left corner.
- (vi) For the bottom right sub-board, we cover it with triominoes with a remaining square at the top left corner.
- (vii) For the bottom left sub-board, we cover it with triominoes with a remaining square at the top right corner.
- (viii) The remaining three squares from (v) to (vii) are connected and can be covered by one triomino.

□

**Remark** Although it is easy to visualise this by drawing it out, always produce a written proof.

**Problem 34** (H3M 2017 Q8). The Fibonacci sequence is defined recursively by  $F_{n+1} = F_n + F_{n-1}$  and  $F_1 = 1, F_2 = 1$ .

- (i) Find the periods of Fibonacci sequences modulo 3 and 4.

- (ii) For any positive integer  $m$ , show that we can find two pairs  $(F_j, F_{j+1})$  and  $(F_k, F_{k+1})$  which are the same modulo  $m$  with  $1 \leq j < k \leq m^2 + 1$ .
- (iii) For  $m$ ,  $j$  and  $k$  as in (ii), explain why the Fibonacci sequence modulo  $m$  is periodic with period dividing  $k - j$ .
- (iv) For any positive integer  $m$ , prove that there is a Fibonacci number which is a multiple of  $m$ .

**Solution:**

- (i) Modulo 3:  $1, 1, 2, 0, 2, 2, 1, 0, 1, 1, \dots$  has period 8.  
Modulo 4:  $1, 1, 2, 3, 1, 0, 1, 1, \dots$  has period 6.
- (ii) Modulo  $m$ , there are  $m$  possible values  $0, 1, 2, \dots, m - 1$ . So there are exactly  $m^2$  possible distinct pairs  $(a, b)$ .  
If we consider  $m^2 + 1$  pairs of  $(F_i, F_{i+1})$  modulo  $m$  where  $1 \leq i \leq m^2 + 1$ , we can find two pairs  $(F_j, F_{j+1})$  and  $(F_k, F_{k+1})$  which are the same modulo  $m$ , by Pigeonhole Principle.
- (iii) This is the same as showing  $F_{j+n} \equiv F_{k+n} \pmod{m}$  for all non negative integer  $n$ .  
We prove using mathematical induction.  
Basis step:  $P(0)$  and  $P(1)$

$$F_j \equiv F_k \pmod{m} \quad F_{j+1} \equiv F_{k+1} \pmod{m}$$

Inductive step:  $P(q - 1) \wedge P(q) \implies P(q + 1)$  for all  $q \geq 1$

Given  $F_{j+q-1} \equiv F_{k+q-1} \pmod{m}$  and  $F_{j+q} \equiv F_{k+q} \pmod{m}$ . Then  $F_{j+q-1} + F_{j+q} \equiv F_{k+q-1} + F_{k+q} \pmod{m}$  so  $F_{j+q+1} \equiv F_{k+q+1} \pmod{m}$ .

By mathematical induction, the sequence repeats itself after  $k - j$  terms. This implies the period of the sequence divides  $k - j$ .

- (iv) For any positive  $m$ , by part (iii), the Fibonacci sequence modulo  $m$  is periodic. That is,  $(F_1, F_2)$  is congruent to  $(F_i, F_{i+1})$  modulo  $m$  for some  $i > 2$ :

$$F_i \equiv F_1 \equiv 1 \pmod{m} \quad F_{i+1} \equiv F_2 \equiv 1 \pmod{m}$$

Then  $F_{i-1} = F_{i+1} - F_i \equiv 1 - 1 \equiv 0 \pmod{m}$ , which means  $m \mid F_{i-1}$ .

We have proven that there is a Fibonacci number which is a multiple of  $m$ .



**Problem 35** (H3M Specimen N03). Functions  $f$  and  $g$  are defined for  $x \in \mathbb{R}$  by

$$f(x) = ax + b, \quad g(x) = cx + d$$

where  $a, b, c, d$  are constants with  $a \neq 0$ . Given that  $gf = f^{-1}g$ , show that

- either  $g$  is a constant function, i.e.  $g(x)$  is constant for all  $x \in \mathbb{R}$ ,
- or  $f^2$  is the identity function, i.e.  $ff(x) = x$  for all  $x \in \mathbb{R}$ ,
- or  $g^2$  is the identity function.

[9]

**Proof** Given that  $gf = f^{-1}g$ ,

$$\begin{aligned} cf(x) + d &= f^{-1}(cx + d) \\ c(ax + b) + d &= \frac{(cx + d) - b}{a} \\ a^2cx + abc + ad &= cx + d - b \end{aligned}$$

Comparing coefficients,

$$\begin{cases} a^2 = c \\ c(a - 1)(a + 1) = 0 \\ abc + ad = d - b \end{cases}$$

and we have three cases to work with. □

## §A.2 Selected problems from school papers

### §A.2.1 Number Theory

### §A.2.2 Analysis

### §A.2.3 Counting

1. (2019 DHS–EJC Prelim Q6) You have an unlimited supply of  $1 \times 1$ ,  $1 \times 2$  and  $2 \times 2$  tiles. Tiles of the same size are indistinguishable.

- (i) Let  $T_n$  is the number of ways of tiling a  $1 \times n$  path.

State the value of  $T_1$  and  $T_2$ . Write down an appropriate recurrence relation between  $T_{n+2}$ ,  $T_{n+1}$  and  $T_n$ . [1]

Consider the tilings of a  $2 \times n$  path. (The  $1 \times 2$  tiles can be rotated in the tilings.)

Let  $P_n$  be the number of tilings of

Let  $Q_n$  be the number of tilings of

- (ii) Show that  $P_{n+1} = P_n + Q_n$  for  $n \geq 1$ . Explain your reasoning clearly. [2]

- (iii) Show that  $Q_{n+1} = 2P_{n+1} + 2Q_{n+1}$  for  $n \geq 2$ . Explain your reasoning clearly. [4]

- (iv) Use (ii) and (iii) to show that  $P_{n+2} + 2P_{n-1} = 3P_{n+1} + 2P_n$  for  $n \geq 2$ . [2]

It is given that the solution to the above recurrence relation is

$$P_n = -\frac{(-1)^n}{7} + \frac{1+2\sqrt{2}}{14}(2+\sqrt{2})^n + \frac{1-2\sqrt{2}}{14}(2-\sqrt{2})^n.$$

- (v) Find the number of distinct ways of tiling a  $2 \times n$  path. [2]

#### Solution

- (i)

- (ii)

- (iii)

- (iv)

- (v) Number of tilings of  $2 \times n$  path is  $Q_n$ . Thus

$$\begin{aligned} Q_n &= P_{n+1} - P_n \\ &= \end{aligned}$$

□

2.