

CHAPTER 18 LINKED LISTS, STACKS, QUEUES, AND PRIORITY QUEUES

OBJECTIVES

- To design and implement linked lists (§18.2).
- To explore variations of linked lists (§18.2).
- ~~– To define and create iterators for traversing elements in a container (§18.3).~~
- ~~– To generate iterators using generators (§18.4).~~
- To design and implement stacks (§18.5).
- To design and implement queues (§18.6).
- To design and implement priority queues (§18.7).

WHAT IS A DATA STRUCTURE?

- A data structure is a collection of data organized in some fashion. A data structure not only stores data, but also supports the operations for manipulating data in the structure.

OBJECT-ORIENTED DATA STRUCTURE

In object-oriented thinking, a data structure is an object that stores other objects, referred to as data or elements. So some people refer a data structure as a container object or a collection object. To define a data structure is essentially to declare a class. The class for a data structure should use data fields to store data and provide methods to support operations such as insertion and deletion. To create a data structure is therefore to create an instance from the class. You can then apply the methods on the instance to manipulate the data structure such as inserting an element to the data structure or deleting an element from the data structure.

FOUR CLASSIC DATA STRUCTURES

Four classic dynamic data structures to be introduced in this chapter are lists, stacks, queues, and priority queues. A list is a collection of data stored sequentially. It supports insertion and deletion anywhere in the list. A stack can be perceived as a special type of the list where insertions and deletions take place only at the one end, referred to as the top of a stack. A queue represents a waiting list, where insertions take place at the back (also referred to as the tail of) of a queue and deletions take place from the front (also referred to as the head of) of a queue. In a priority queue, elements are assigned with priorities. When accessing elements, the element with the highest priority is removed first.

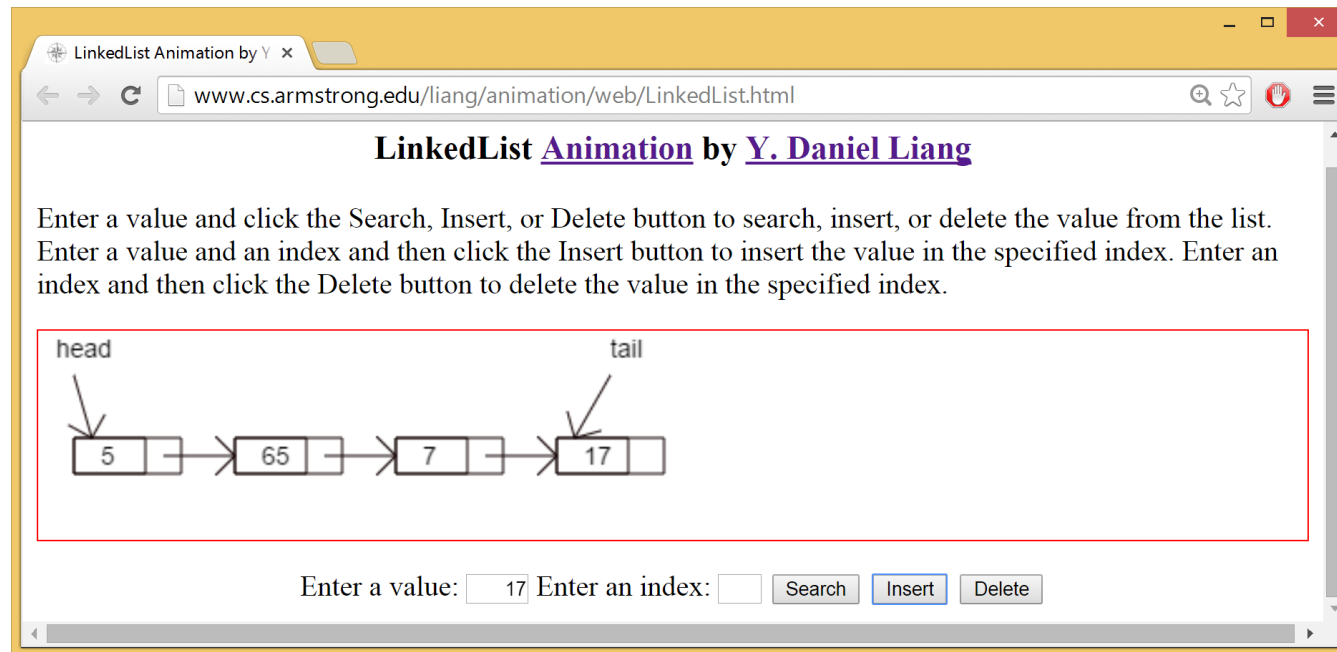
LISTS

A list is a popular data structure to store data in sequential order. For example, a list of students, a list of available rooms, a list of cities, and a list of books, etc. can be stored using lists. The common operations on a list are usually the following:

- Retrieve an element from this list.
- Insert a new element to this list.
- Delete an element from this list.
- Find how many elements are in this list.
- Find if an element is in this list.
- Find if this list is empty.

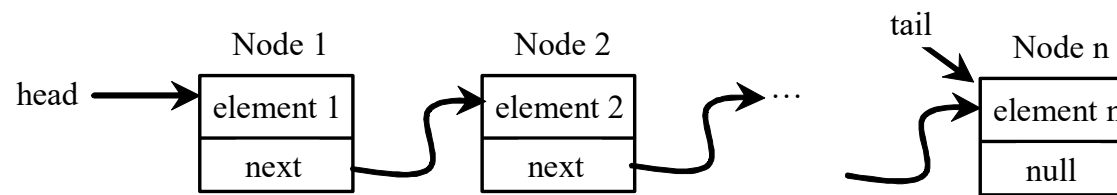
LINKED LIST ANIMATION

<https://liveexample.pearsoncmg.com/dsanimation/LinkedListBook.html>



NODES IN LINKED LISTS

- A linked list consists of nodes. Each node contains an element, and each node is linked to its next neighbor. Thus a node can be defined as a class, as follows:



class Node:

def __init__(self, element):

self.elmenet = element

self.next = None

ADDING THREE NODES

The variable head refers to the first node in the list, and the variable tail refers to the last node in the list. If the list is empty, both are None. For example, you can create three nodes to store three strings in a list, as follows:

Step 1: Declare head and tail:

```
head = None  
tail = None
```

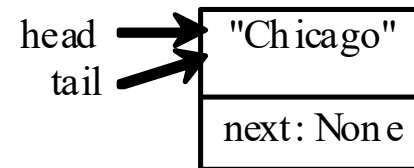
The list is empty now

ADDING THREE NODES, CONT.

Step 2: Create the first node and insert it to the list:

```
head = Node("Chicago")  
last = head
```

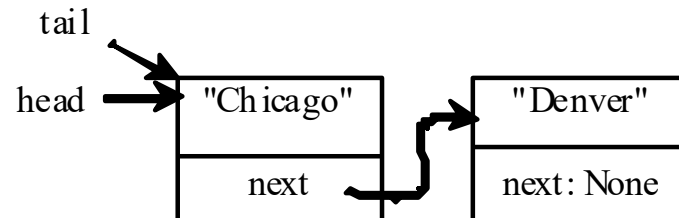
After the first node is inserted



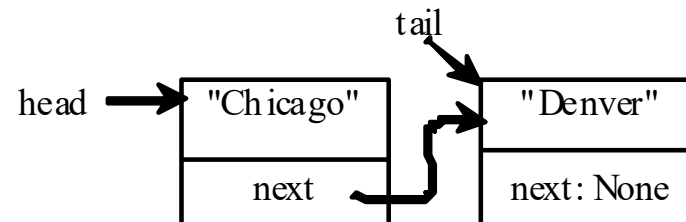
ADDING THREE NODES, CONT.

Step 3: Create the second node and insert it to the list:

```
tail.next = Node("Denver")
```



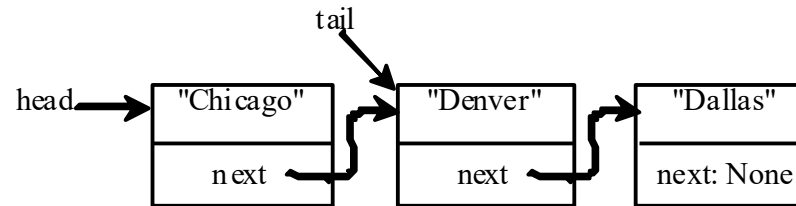
```
tail = tail.next
```



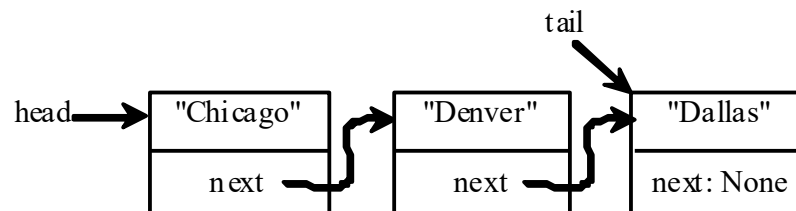
ADDING THREE NODES, CONT.

Step 4: Create the third node and insert it to the list:

```
tail.next = Node("Dallas")
```



```
tail = tail.next
```

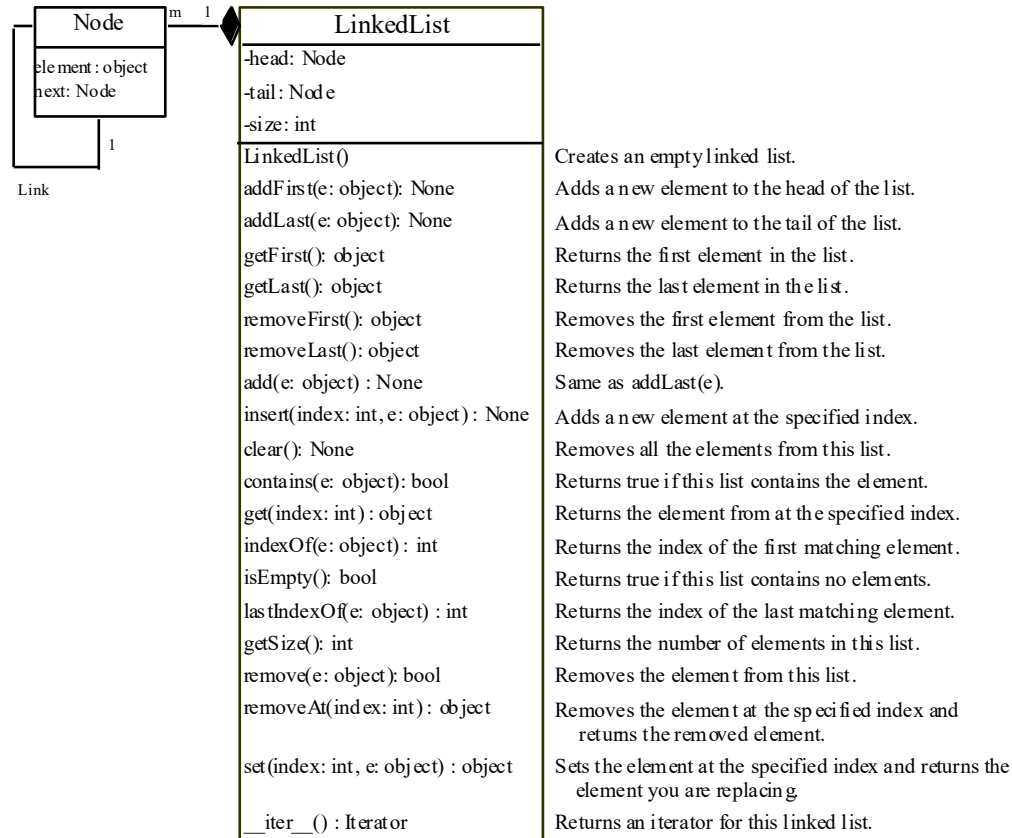


TRAVERSING ALL ELEMENTS IN THE LIST

Each node contains the element and a data field named *next* that points to the next element. If the node is the last in the list, its pointer data field next contains the value None. You can use this property to detect the last node. For example, you may write the following loop to traverse all the nodes in the list.

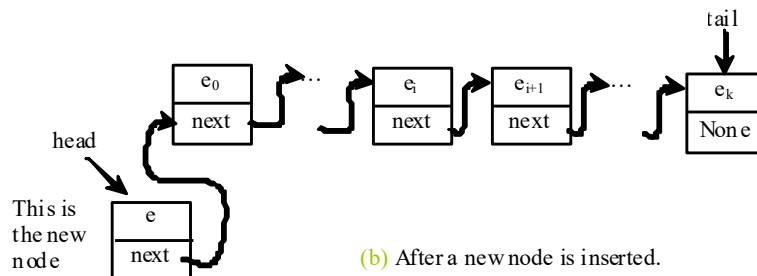
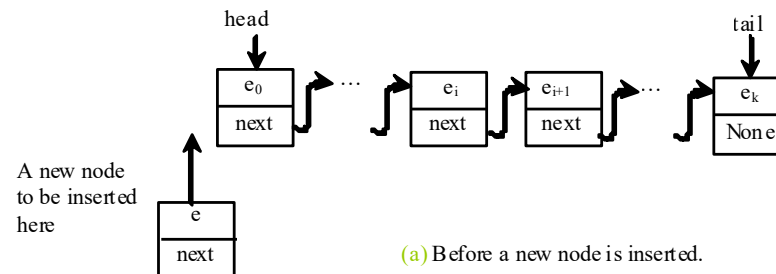
```
current = head
while current != None:
    print(current.element)
    current = current.next
```

LINKEDLIST



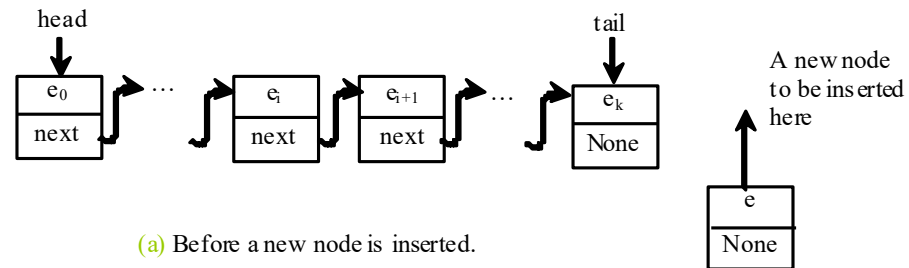
IMPLEMENTING ADDFIRST(E)

```
def addFirst(self, e):  
    newNode = Node(e) # Create a new node  
    newNode.next = self.__head # link the new node with the head  
    self.__head = newNode # head points to the new node  
    self.__size += 1 # Increase list size  
    if self.__tail == None: # the new node is the only node in list  
        self.__tail = self.__head
```

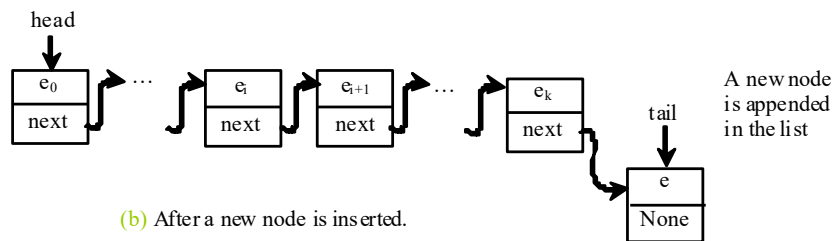


IMPLEMENTING ADDLAST(E)

```
def addLast(self, e):  
    newNode = Node(e) # Create a new node for e  
  
    if self.__tail == None:  
        self.__head = self.__tail = newNode # The only node in list  
    else:  
        self.__tail.next = newNode # Link the new with the last node  
        self.__tail = self.__tail.next # tail now points to the last node  
  
    self.__size += 1 # Increase size
```



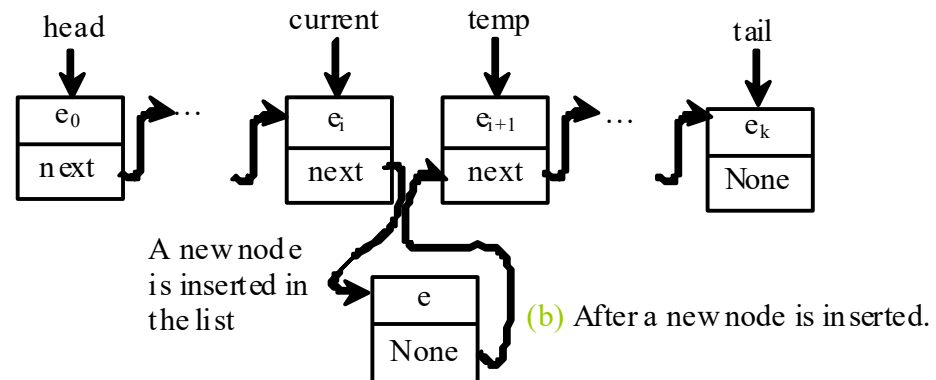
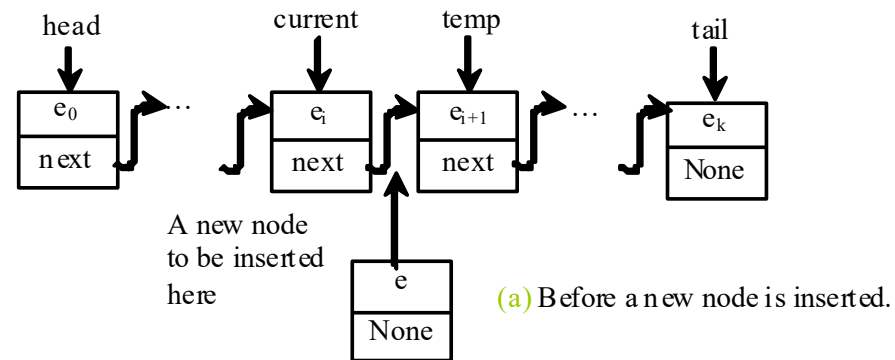
(a) Before a new node is inserted.



(b) After a new node is inserted.

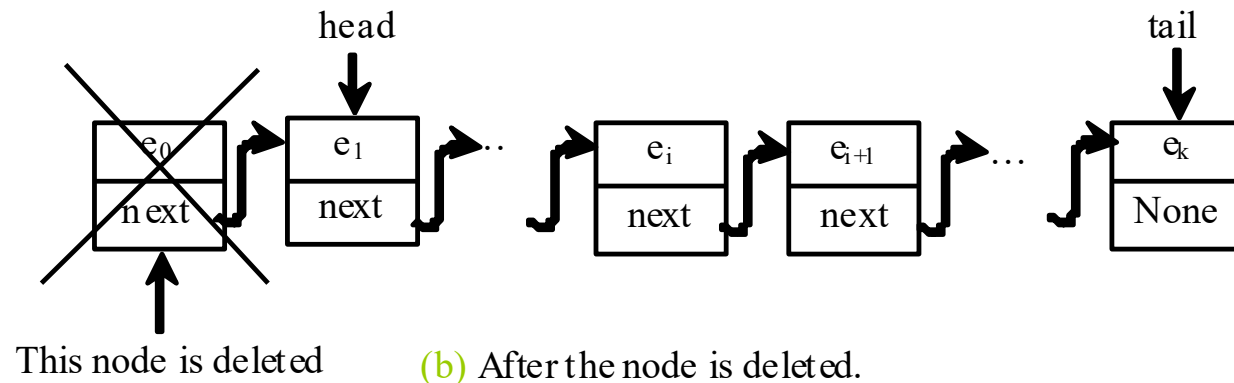
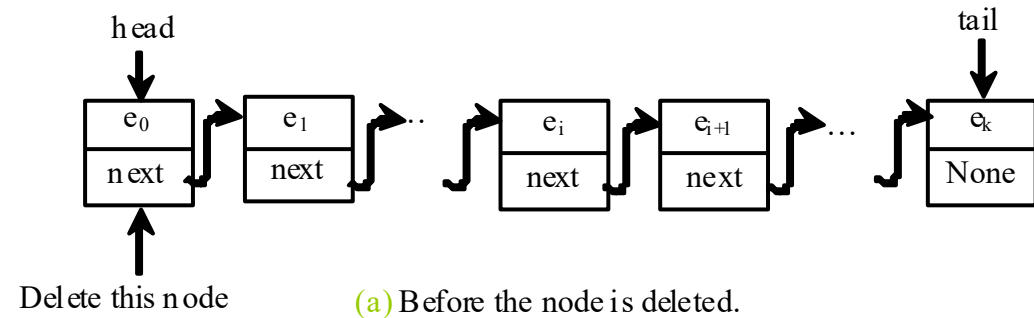
IMPLEMENTING ADD(INDEX, E)

```
def insert(self, index, e):  
    if index == 0:  
        self.addFirst(e) # Insert first  
    elif index >= size:  
        self.addLast(e) # Insert last  
    else: # Insert in the middle  
        current = head  
        for i in range(1, index):  
            current = current.next  
        temp = current.next  
        current.next = Node(e)  
        (current.next).next = temp  
        self.__size += 1
```



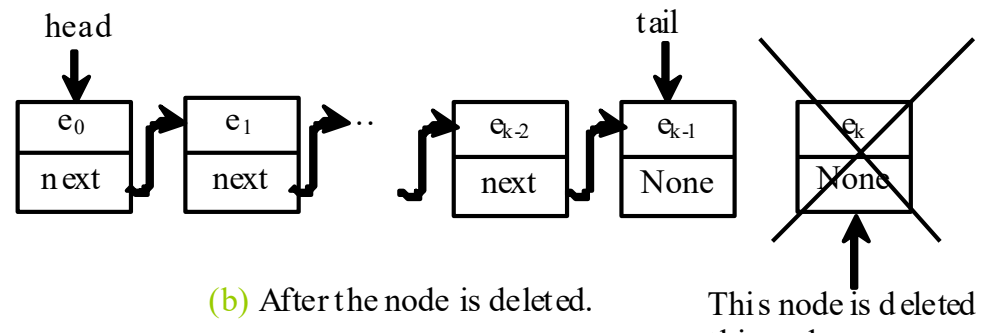
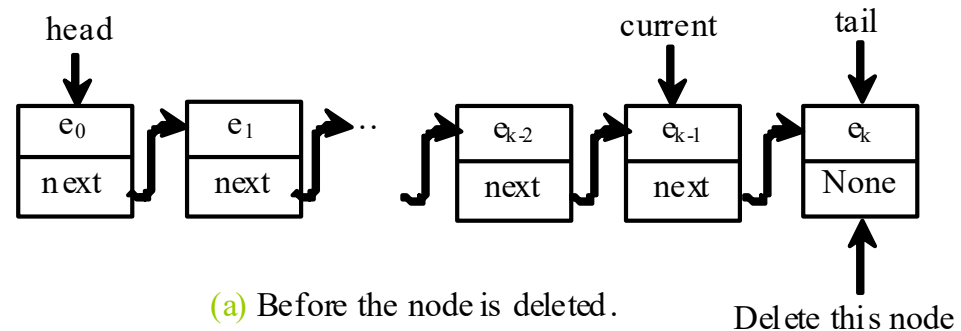
IMPLEMENTING REMOVEFIRST()

```
def removeFirst(self):  
    if self.__size == 0:  
        return None  
    else:  
        temp = self.__head  
        self.__head =  
        self.__head.next  
        self.__size -= 1  
    if self.__head == None:  
        self.__tail = None  
    return temp.element
```



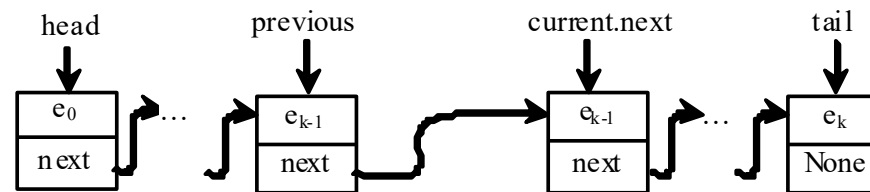
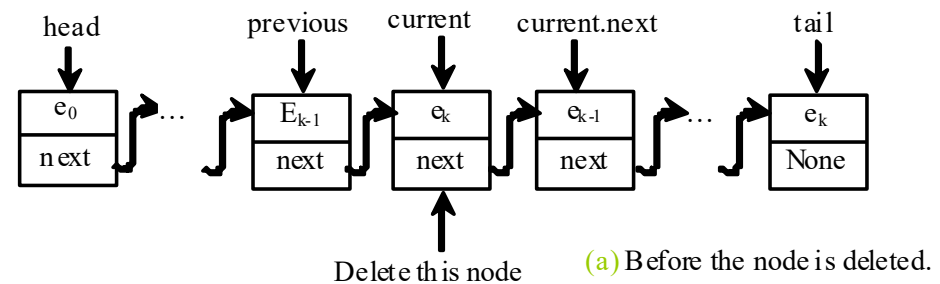
IMPLEMENTING REMOVE LAST()

```
def removeLast(self):  
    if self.__size == 0:  
        return None  
    elif self.__size == 1:  
        temp = self.__head  
        self.__head = self.__tail = None  
        self.__size = 0  
        return temp.element  
    else:  
        current = self.__head  
  
        for i in range(self.__size - 2):  
            current = current.next  
  
        temp = self.__tail  
        self.__tail = current  
        self.__tail.next = None  
        self.__size -= 1  
        return temp.element
```



IMPLEMENTING REMOVEAT(INDEX)

```
def removeAt(self, index):  
    if index < 0 or index >= self.__size:  
        return None # Out of range  
    elif index == 0:  
        return self.removeFirst() # Remove first  
    elif index == self.__size - 1:  
        return self.removeLast() # Remove last  
    else:  
        previous = self.__head  
  
        for i in range(1, index):  
            previous = previous.next  
  
        current = previous.next  
        previous.next = current.next  
        self.__size -= 1  
        return current.element
```

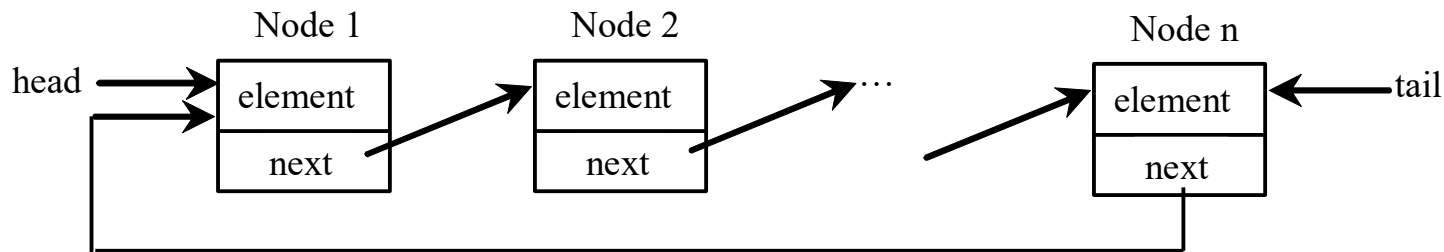


TIME COMPLEXITY FOR LIST AND LINKEDLIST

Methods for list/Complexity		Methods for LinkedList/Complexity	
<code>append(e: E)</code>	$O(1)$	<code>add(e: E)</code>	$O(1)$
<code>insert(index: int, e: E)</code>	$O(n)$	<code>insert(index: int, e: E)</code>	$O(n)$
N/A		<code>clear()</code>	$O(1)$
<code>e in myList</code>	$O(n)$	<code>contains(e: E)</code>	$O(n)$
<code>list[index]</code>	$O(1)$	<code>get(index: int)</code>	$O(n)$
<code>index(e: E)</code>	$O(n)$	<code>indexOf(e: E)</code>	$O(n)$
<code>len(x) == 0?</code>	$O(1)$	<code>isEmpty()</code>	$O(1)$
N/A		<code>lastIndexOf(e: E)</code>	$O(n)$
<code>remove(e: E)</code>	$O(n)$	<code>remove(e: E)</code>	$O(n)$
<code>len(x)</code>	$O(1)$	<code>getSize()</code>	$O(1)$
<code>del x[index]</code>	$O(n)$	<code>removeAt(index: int)</code>	$O(n)$
<code>x[index] = e</code>	$O(n)$	<code>set(index: int, e: E)</code>	$O(n)$
<code>insert(0, e)</code>	$O(n)$	<code>addFirst(e: E)</code>	$O(1)$
<code>del x[0]</code>	$O(n)$	<code>removeFirst()</code>	$O(1)$

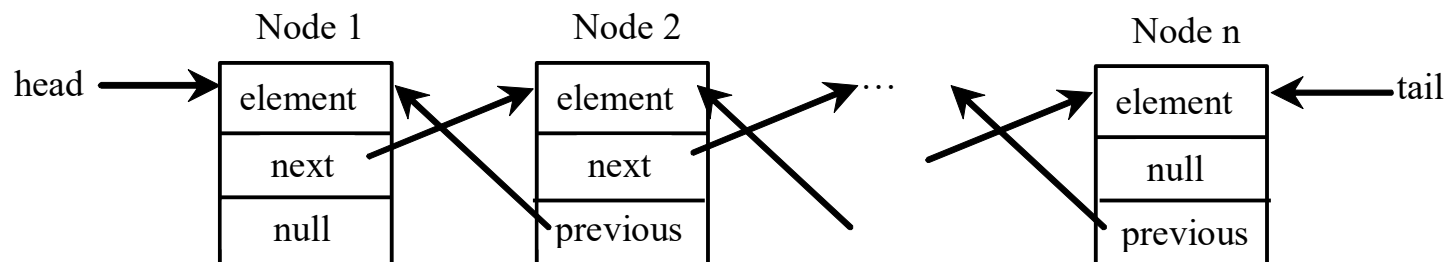
CIRCULAR LINKED LISTS

A circular, singly linked list is like a singly linked list, except that the pointer of the last node points back to the first node.



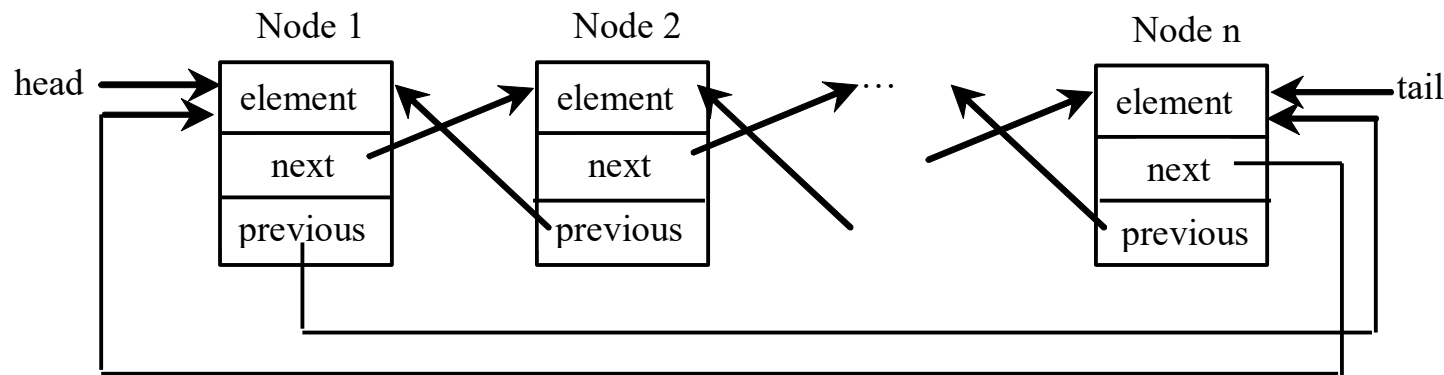
DOUBLY LINKED LISTS

A doubly linked list contains the nodes with two pointers. One points to the next node and the other points to the previous node. These two pointers are conveniently called a forward pointer and a backward pointer. So, a doubly linked list can be traversed forward and backward.



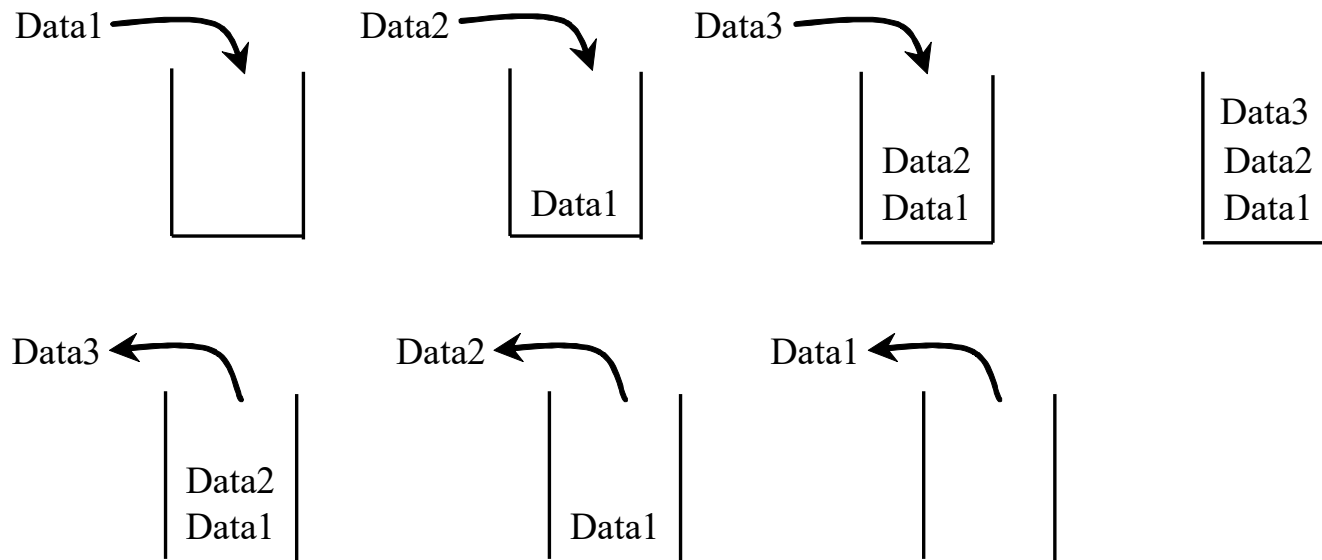
CIRCULAR DOUBLY LINKED LISTS

A circular, doubly linked list is doubly linked list, except that the forward pointer of the last node points to the first node and the backward pointer of the first pointer points to the last node.



STACKS

A stack can be viewed as a special type of list, where the elements are accessed, inserted, and deleted only from the end, called the top, of the stack.



STACK ANIMATION

<https://liveexample.pearsoncmg.com/dsanimation/StackeBook.html>



Stack Animation by Y. Daniel Liang

Enter a value and click the Push button to push the value into the stack. Click the Pop button to remove the top element from the stack.

Top →

5
3
3
4

Enter a value:

STACK

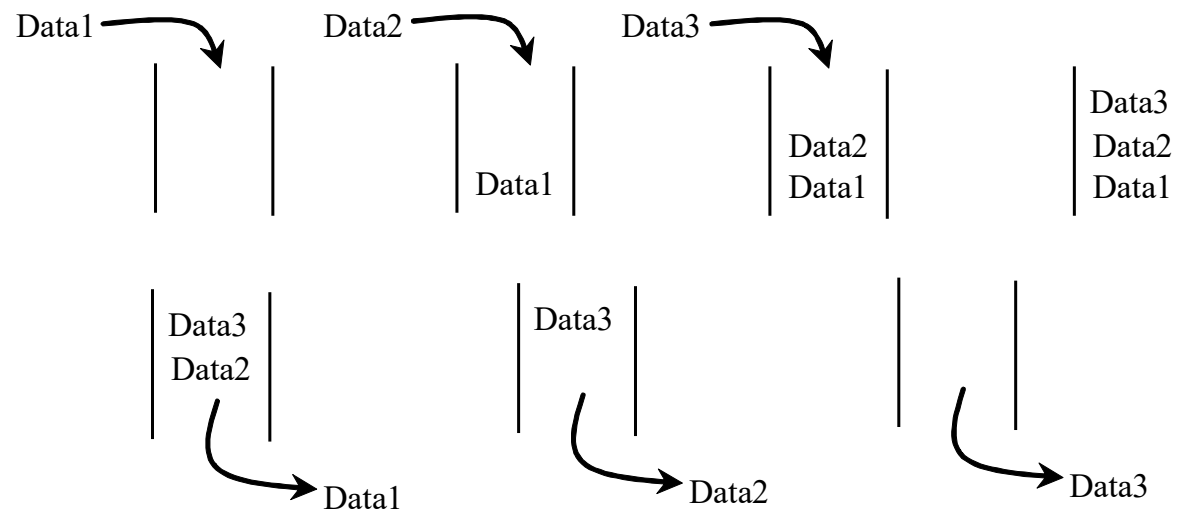
Stack	
-elements: list	A list to store the elements in the stack.
Stack()	Constructs an empty stack.
isEmpty(): bool	Returns true if the stack is empty.
peek(): object	Returns the element at the top of the stack without removing it from the stack.
push(value: object): None	Stores an element into the top of the stack.
pop(): object	Removes the element at the top of the stack and returns it.
getSize(): int	Returns the number of elements in the stack.

Stack

TestStack

QUEUES

A queue represents a waiting list. A queue can be viewed as a special type of list, where the elements are inserted into the end (tail) of the queue, and are accessed and deleted from the beginning (head) of the queue.



QUEUE ANIMATION

<https://liveexample.pearsoncmg.com/dsanimation/QueueBook.html>



Stack Animation by Y. Daniel Liang

www.cs.armstrong.edu/liang/animation/web/Queue.html

Queue Animation by Y. Daniel Liang

Enter a value and click the Enqueue button to append the value into the tail of the queue. Click the Dequeue button to remove the element from the head of the queue.

head tail

5 45 2 4 21

Enter a value: 21 Enqueue Dequeue

QUEUE

Queue	
-elements LinkedList	Stores queue's elements in a list.
Queue()	Creates an empty queue.
enqueue(e: object): None	Adds an element to this queue.
dequeue(): object	Removes an element from this queue.
getSIZE(): int	Returns the number of elements from this queue.
isEmpty(): bool	Returns true if the queue is empty.
__str__(): str	Returns a string representation of the queue.

Queue

TestQueue

PRIORITY QUEUE

A regular queue is a first-in and first-out data structure. Elements are appended to the end of the queue and are removed from the beginning of the queue. In a priority queue, elements are assigned with priorities. When accessing elements, the element with the highest priority is removed first. A priority queue has a largest-in, first-out behavior. For example, the emergency room in a hospital assigns patients with priority numbers; the patient with the highest priority is treated first.

PriorityQueue	
-heap: Heap	Elements are stored in a heap.
enqueue(element: object): None	Adds an element to this queue.
dequeue(): object	Removes an element from this queue.
getSize(): int	Returns the number of elements from this queue.